4736

| | (a) (i |) 8 7 5 4 3 3 3 3 2 2 | MI | | For sorting the list into decreasing order |
|---|-----------|--|-----------------|-----|---|
| | | First bag 8 2 | MI | | For trying to apply first-fit to their list |
| | | Second bag 7 3 Third bag 5 4 | | | |
| | | Fourth bag 3 3 3 Fifth bag 2 | | | |
| | | rata bag 2 | Al | | For a completely correct solution |
| | (ii) | A packing that uses fewer bags could be | | | |
| | | First bag 8 2 | ВТ | | For any valid packing into four bags (may be as |
| | | Second bag 7 3 Third bag 5 3 2 | | | an incorrect answer to using algorithm, need no |
| | | Fourth bag 4 3 3 | | | be packed in this order) |
| | (b) | | | | For scaling 4 seconds by 5° or for an equivalent |
| | | $\left(\frac{500}{100}\right) \times 4$ or 125000000 + 0,000004 | MI | | valid and complete method. Condone minor |
| | | (100) - 500 | Al | | errors with the number of zeros. For 500 or 500 seconds or 500 s. |
| | | | ''' | 6 | Accept 8 minutes 20 seconds or 8.3 minutes |
| | (i) | cā | ВТ | | For any simple graph with 4 vertices and 5 ares |
| | | | | | Vertices need not be labelled |
| | | | | | Need not be planar |
| | (ii) | The Calling Court of the Section 1997 | | | |
| | (11) | The sum of the orders of the vertices is twice the number of arcs, and hence is even. | MI | | Or start from a null graph and successively add ares. |
| | | | 1 | | Each time an arc is added the number of odd |
| | | Hence the sum of the odd orders must be even and so there must be an even number of odd | l _{A1} | | vertices is either unchanged or it increases or decreases by 2. |
| | (iii) | vertices. | | | So the number of odd nodes is always even |
| | (111) | 5 arcs \Rightarrow sum of orders of vertices = 10 Simple graph connecting vertices so each vertex | MI | | |
| | | has order 1, 2 or 3 | | | |
| | | 1 + 3 + 3 + 3 + 10 or $2 + 2 + 3 + 3 + 10$ | ļ | | |
| | | But $1 \rightarrow 3 + 3 + 3$ is not possible since if three | AT | | Explaining why $1 + 3 + 3 + 3$ is not possible. |
| | | vertices have order 3 they are all connected to the fourth vertex so it also has order 3. | | ļ | |
| | | With $2 + 2 + 3 + 3$ the two vertices of order 2 | AL | | P. D. C. L. |
| | | cannot be adjacent, since otherwise two ares | / · · | | Explaining why there is only one graph with nodes of orders 2, 2, 3, 3. |
| | | connect the other two vertices so not simple. | | ŝ | |
| | | Hence only one possible graph. | | (3 | |
| - | (i) | | ļ | 6 | |
| | • | <u>, </u> | Мι | | For a correct tree (labels not required) |
| | | () | | ļ | |
| | | $E \longrightarrow F$ | l I | Í | |
| | | $G \bullet H$ | | | |
| | | Kruskal: DF, CD, BD and EF, FH, AC, EG | AI | } | For a valid order (using Prim or Kruskal) |
| | | 40 | BT | | For length 40 |
| | (ii) | ACDFEGHBA | MI | - 1 | At getting at least as far as $A \cap DFE$ |

| BI BI BI BI BI | (if shown on a diagram, needs direction shown) For both, vertices in any order For 5 For ABCD, vertices in any order Answer should be on insert sheet For using Dijkstra's algorithm |
|----------------------------|---|
| MI | For using Dijkstra's algorithm updating at E and E (even if incomplete) |
| | updating at E and F (even if incomplete) |
| A I | For all rarmonant labels around |
| | To an permanent rapers correct |
| B 1 | For valid order of assigning permanent labels |
| 31 | For copying their permanent labels, or correct values |
| MI MI | Correct answer only For identifying A and F or value 19 or their 19 For 120 + their 19 |
| AI M1 | For 139 (cao) For identifying F and G or value 10 as only extra For 130 (cao) |
| | 81 31 41 41 |

| 7 (i) | Minimise $70x + 80y + 50z$ | Bi | For 'minimise' a (non-zero) multiple of 7x+8y+5z |
|-------|---|----------|--|
| | No more than twice as many packs of type Y as packs of type X | 81 | For identifying this constraint from the list, or equivalent |
| | Other constraints $y \ge 200, 0 \ge z \le 50$ $y \ge z$ $x + z \le 220$ $x + y \ge 300$ | BI BI | Ignore extra 'constraints' unless contradictions For boundary constraints on x and z For this, or an equivalent correct answer For this, or an equivalent correct answer |
| | 3 1 1 300 | B1 | For this, or an equivalent correct answer Use of strict inequalities—penalise first time only |
| (ii |) | | cise of so tel medianties penanse first time only |
|) (a | Subject to $y \leq 2x$ $x \leq 200$ $y \geq 50$ | MI A1 | For replacing z by 50 For their $y \ge 50$ |
| | v + p <u>≥</u> 300 500 {_v to the left } √ | МІ | For at least two appropriate lines drawn on a graph with plausibly scaled axes. |
| | 300 (easible region) | MI | For boundary lines drawn correctly (follow through their equations provided there are at least two horizontal or vertical lines and at least (wo lines that 'slope') |
| | 100 200 300 400 500 | Al | Feasible region correctly identified (correct answer only, not follow through) |
| (b) | (200, 400), (200, 100), (250, 50) | MI Al | For reading off or calculating at least one of their vertices For getting these three vertices correct with no extras |
| | (200, 100) gives $70x + 80y = 22000 \text{ (£245)}$ (250, 50) gives $70x + 80y = 21500 \text{ (£240)}$ | МІ | For calculating their cost at one of their vertices or using an appropriate line of constant cost |
| · | Cost is minimised when $x = 250$, $y = 50$ Cost = £240 | AI BI | For identifying vertex (250, 50) For £240 or 24000 p (with units) |
| (iii | eg x = 300, y = 0 , z = 0 only costs £210 | MI | For finding a feasible point with z > 50 Or a written explanation For finding such a feasible point with a lower cost |
| - | | 18 | than that in (ii)(b) and showing that cost is lower. |

Mark Scheme