

Center for Information Services and High Performance Computing (ZIH)

Interim Diploma Presentation

Automatic Profile-Based Characterization
of Performance Traces
for Highly Parallel Applications

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Contents

- Introduction
 - Performance Analysis, Vampir, Challenges
 - Differences in Traces
- Classifying Process Traces
 - Current Approaches
 - Methodology
 - Evaluation
- A Novel Profile Display
 - State of the Art
 - Outlook
- Conclusion, Future Work, References





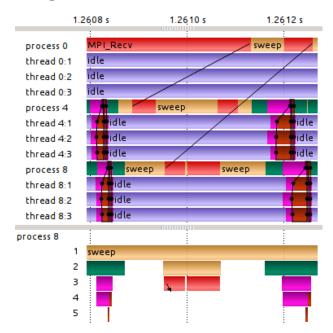
Introduction > Performance Analysis

- Profiling
 - Information accumulated per function and process
 - Sampled
 - Small overhead

cumulative	e self		
seconds	seconds	calls	name
0.31	0.05	25195	<pre>QList::isEmpty()</pre>
0.40	0.04	30239	QList::Node::t()
0.44	0.04	12294	QList::end()
0.55	0.03	3696	QList::end()
0.70	0.03	4939	handleEnter
0.73	0.03	36939	handleLeave
0.88	0.02	99207	<pre>void std::swap()</pre>

Tracing

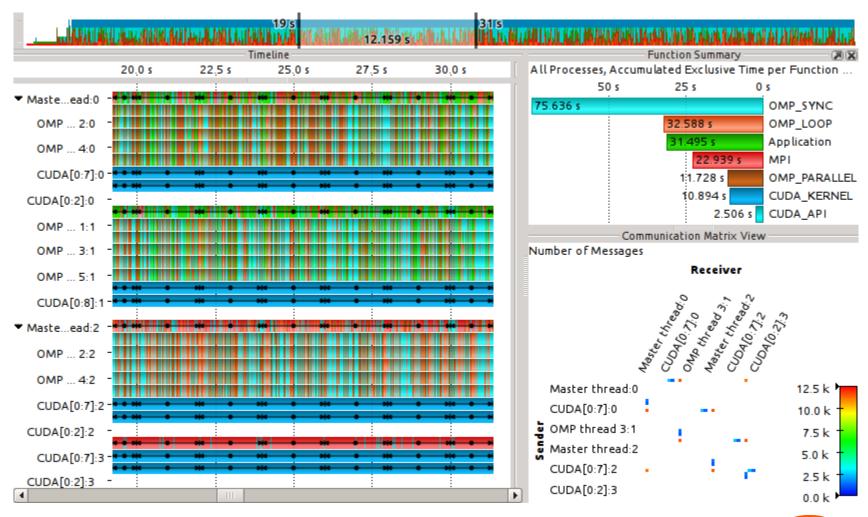
- Complete information about a program run
- Instrumented
- Large overhead







Introduction > Vampir







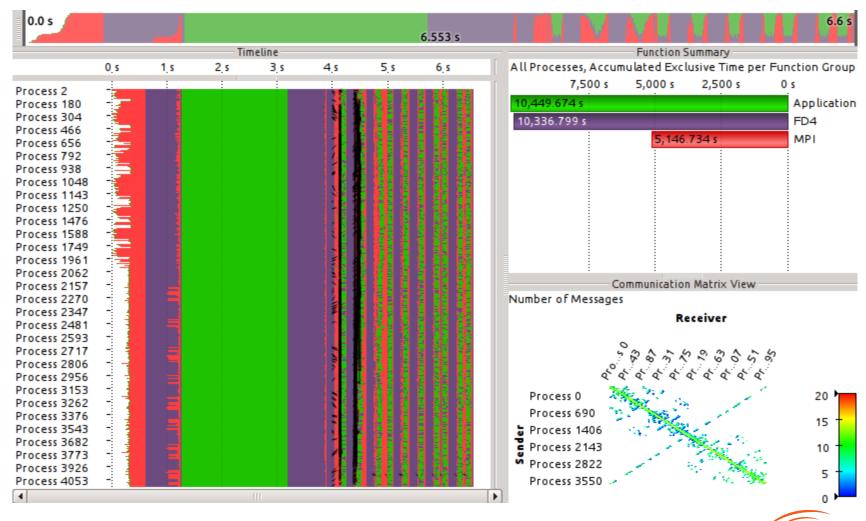
Introduction > Challenges

- Datastructures & algorithms:
 - Limited main memory size
 - Achieving scalibility wrt:
 - Number of processes in the trace
 - Number of processes used for analysis
 - Tracefile Size
- Visualisation:
 - Limited number of Pixels on a screen
 - Achieving scalibilty wrt the number of processes in the trace





Introduction > Challenges

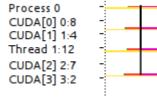






Introduction > Differences in Traces

- How does detecting differences between process traces aid performance analysis?
 - Present new information
 - Visualise timing differences between similar processes
 - Visualise the impact of optimisation
 - Compare runs of the same program on different platforms
 - Improve scalibity of existing views
 - Preserve screen real estate by e.g. merging similar processes





 Detect timing differences between structurally similar processes





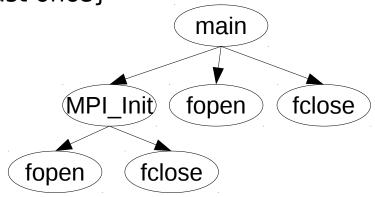
Classifying Process Traces > Current Approaches

- Take any similarity measure. Calculating it for all pairs of processes already requires an all-to-all comparison, which is not feasible for 50k+ processes.
- Clustering n processes using common clustering algorithms, e.g. k-means and DBSCAN, takes at least n² steps
- Some algorithms require a predefined number of clusters or make assumptions about the data layout, which leads to unnatural clusters.
- To handle large amounts of processes properly, we need something better





- A simple, structural similarity metric
 - Each process p is assigned a set of functions{A | A is called on p at least once}
 - Similarity(P1, P2) := |P1nP2| / |P1uP2|
 - Alternatively: each process p is assigned a set of function pairs
 {(A, B) : A calls B on process p at least once}
 - Similarity is calculated analogously



Next step: Calculate and store similarity between processes

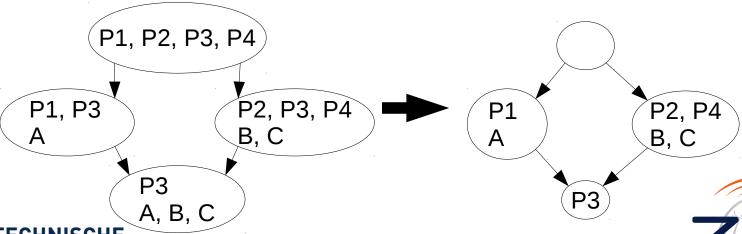




- Use a concept lattice
 - Total order, e.g. (\mathbb{N} , ≤)
 - Partial order, e.g. (ℕ, |)
 - Bounded lattice, e.g. ({A, B}, ⊆)
 - Concept lattice $(\mathcal{P}, \mathcal{F}, \mathcal{I} \subseteq \mathcal{P} \times \mathcal{F})$

where e.g. $P := \{P1, P2, P3\}, \mathcal{F} := \{A, B, C\},\$

 $I := \{ P1 : A P2 : B, C P3 : A, B, C P4 : B, C \}$



2

{A, B}

Ø

Zentrum für Informationsdienste

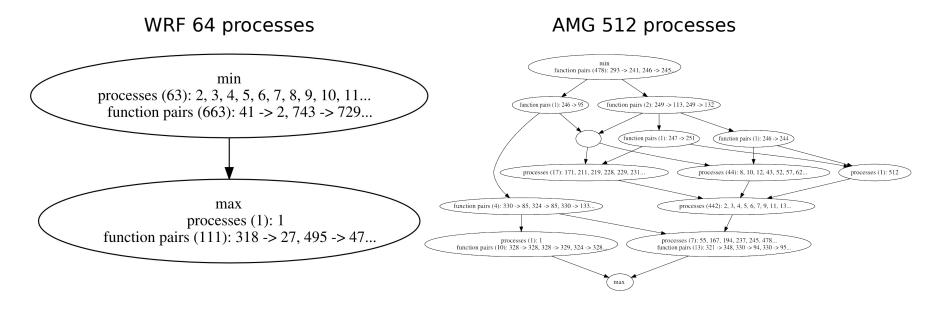
und Hochleistungsrechnen

{B}

{A}

1

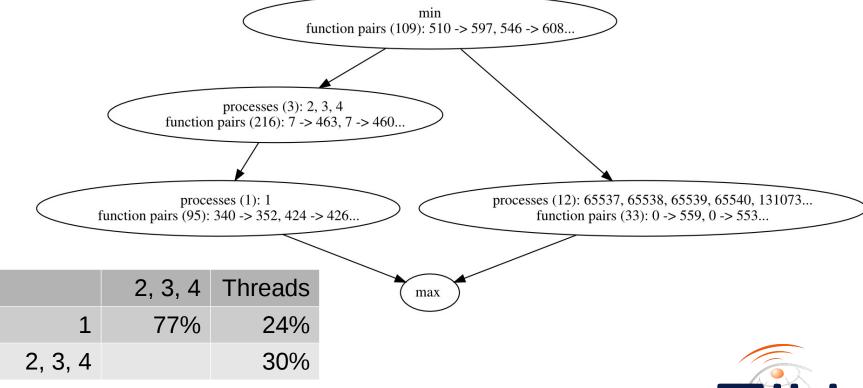
- The result is a directed acyclic graph of "function pair inheritance"
- Expected complexity for building and storing the graph is linear, worst case exponential



Next step: calculate similarity between "equivalence classes"



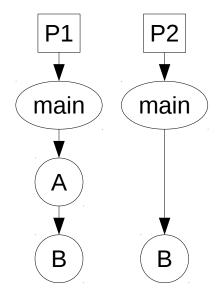
- Similarity(P1, P2) := |P1nP2| / |P1uP2|
 - Simple addition and division of set sizes. No set operations involved.
 - Similarity matrix can then be calculated in O(g²n) steps

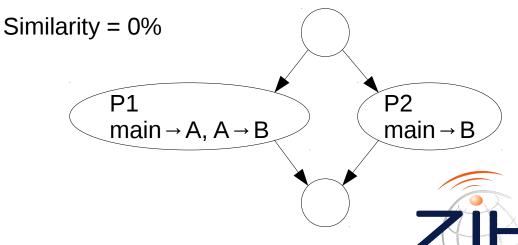






- Slightly modifying this technique lets us find out whether a process calls the same and more functions than another. This way we can:
 - Detect inlining
 - Detect wether processes are dissimilar because one performs additional calls, e.g. logs output
 - Detect different levels of coarsness of trace recording



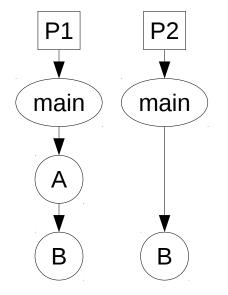


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und Hochleistungsrechnen

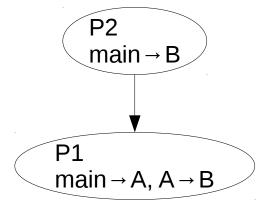


1st step: Calculate transitive closure of both P and Q



P1 =
$$\{(main, A), (A, B)\} \cup \{(main, B)\}$$

P2 = $\{(main, B)\}$



- 2nd step: Build the concept lattice
- 3rd step: SubsumptionMetric(P1, P2) := |P1nP2| / |P2|
- Example result is 100%

	P1	P2
P1		1
P2	1/3	





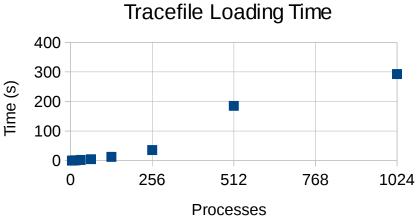
Trace			Result			
Application	Proc.	Size (MiB)	Construction Time (%)	Nodes	Groups	
LINPACK	832	145	≤ 0.1	3	3	
Gromacs	36	3,700	≤ 0.1	24	11	
COSMO-SPECS	100	976	≤ 0.1	1	1	
FD4	4096	166	≤ 0.1	2	2	
WRF	64	283	≤ 0.1	2	2	
HOMME	1024	179	≤ 0.1	3	3	
LULESH	432	354	0.27	182	35	
PIConGPU	39	286	≤ 0.1	60	17	
ВТ	16	150	≤ 0.1	5	3	



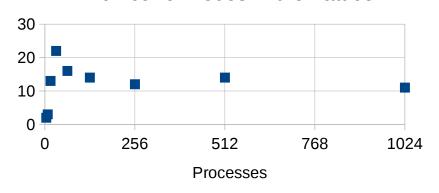


AMG2006

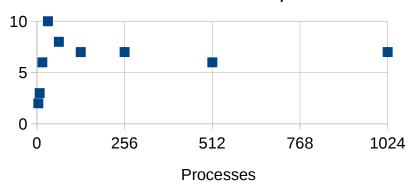




Number of Nodes in the Lattice



Number of Groups

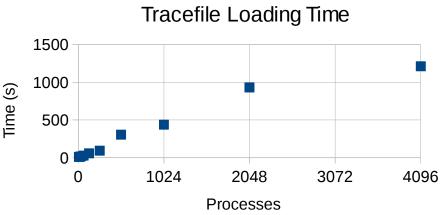




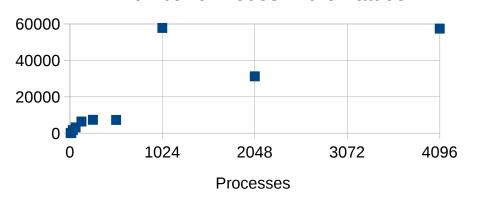


ParaDiS

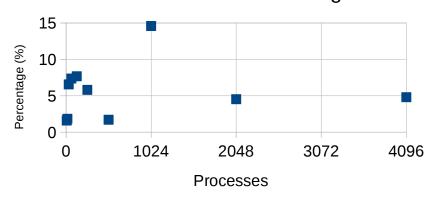




Number of Nodes in the Lattice



Construction Time div. Loading Time





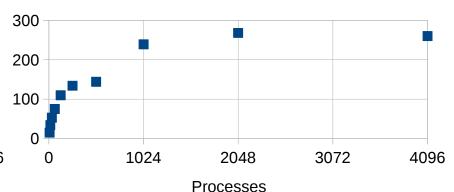


ParaDiS

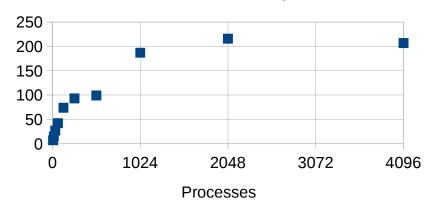
Number of Nodes in the Lattice

60000 20000 0 1024 2048 3072 4096 Processes

Number of non-empty Nodes in the Lattice



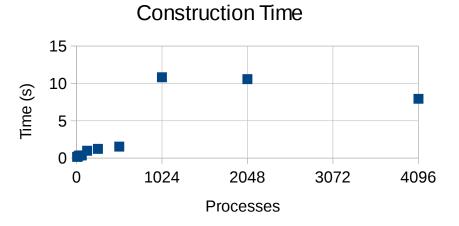
Number of Groups

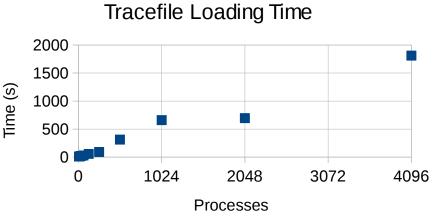




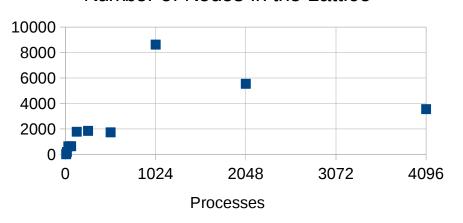


ParaDiS, using called functions instead of function pairs

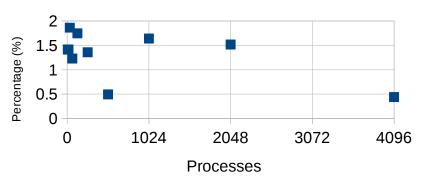




Number of Nodes in the Lattice



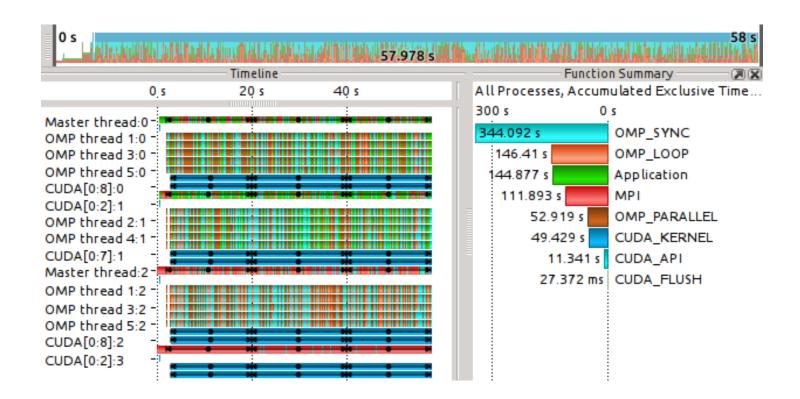
Construction Time div. Loading Time







- Gromacs
 - Decreasing tracing detail from Masterthread 0 to 3







Gromacs

 Subsumption metric matrix, function appearance (not call relations), for all master threads

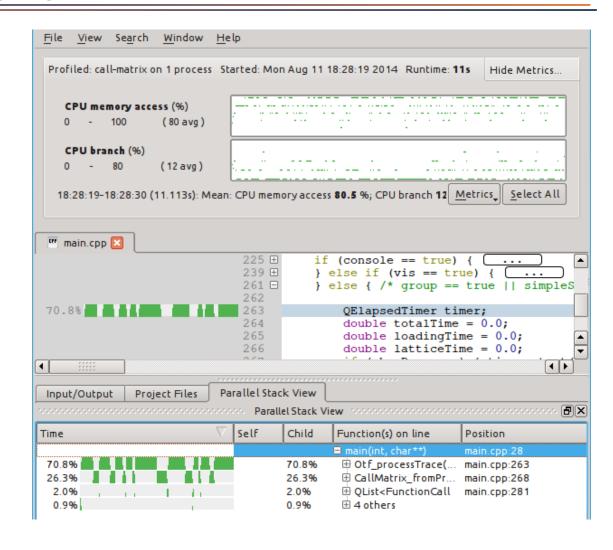
	0	1	2	3
0		100%	99%	98%
1	76%		99%	98%
2	24%	31%		100%
3	8%	10%	33%	





A Novel Profile Display > State of the Art

- Allinea Map
 - Call path profiling
 - Accumulates over all processes
 - Sliced information (vampir-esque)
 - Various metrics
 - Very intuitive and simple GUI

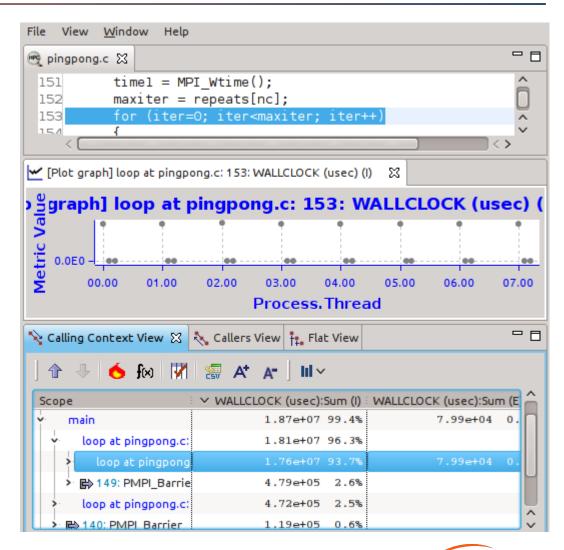






A Novel Profile Display > State of the Art

- HPCToolkit
 - Call path profiling
 - Accumulates over all processes
 - Plots for one function value x processes
 - Custom metrics
 - Usability is ok

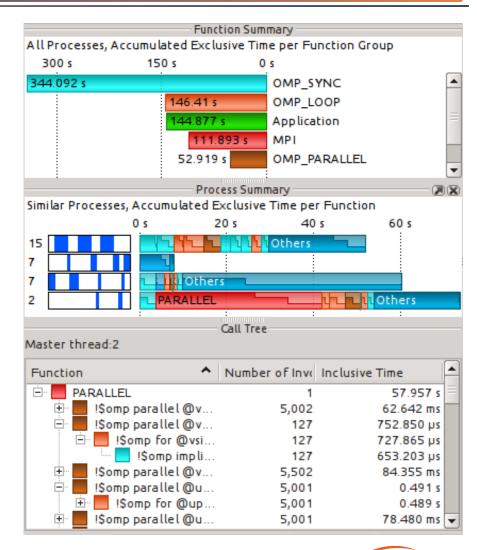






A Novel Profile Display > State of the Art

- Vampir
 - Powerful profile displays
 - Supports clustering similar profiles
 - Usability is good, but needs practice/knowledge







A Novel Profile Display > Outlook

- What can we improve?
 - Replace the Function Summary, Call Tree and Process Summary displays in Vampir with one unified display.
 - Usability
 - Comparative analysis





A Novel Profile Display > Outlook

- Ideas: General Usability
 - Fewer bars, maybe none
 - Improved search functionality
 - Proper undo stack
 - Context menu that is actually context-sensitive

- ...





A Novel Profile Display > Outlook

- Ideas:
 - Show useful and simple information at first sight
 - Enable easy regrouping of functions
 - Highlighting points of interest
 - Special sorting options
 - ?
 - Display timing variations between different calls or processes using e.g. box plots
 - Comparative Analysis:
 - Filter processes using the new clustering
 - ?





Conclusion

- Introduced a structural similarity metric
- Developed methods to cluster and compare processes in < O(n²) steps</p>
 - works with any set-based similarity metric
- First steps towards a better profile viewer





Future Work

- Further scalability testing of the clustering
- Develop a novel, scalable profile viewer
- Use the developed methods in an actual visualisation





References

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- [2] Matthias Weber, Kathryn Mohror, Martin Schulz, Bronis R. de Supinski, Holger Brunst, and Wolfgang E. Nagel. Alignment-Based Metrics for Trace Comparison. In *Euro-Par 2013 Parallel Processing*, pages 29–40. Springer, 2013.
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Thank You!



