

A3 written

Either

The Functor and Applicative for the Either data-type is as follows:

```
instance Functor (Either a) where
  fmap f (Right x) = Right (f x)
  fmap f (Left x)  = Left x
```

```
instance Applicative (Either e) where
  pure      = Right
  Left e <*> _ = Left e
  Right f <*> r = fmap f r
```

Question 1. [4 MARKS]

Show Either satisfies the *second* functor law:

$$\text{fmap } (g \circ h) = \text{fmap } g \circ \text{fmap } h$$

Application of composition

$$(f \circ g) x = f \circ g \$ x = f \$ g x = f \$ g \$ x$$

for any $x :: \text{Either}$

$$P(x) \iff (g.h) \langle \$ \rangle x = (g \langle \$ \rangle).(h \langle \$ \rangle) x$$

Suppose $x = \text{Right } a$

Suppose $x = \text{Left } a$

$$(g.h) \langle \$ \rangle (\text{Right } a)$$

$$(g.h) \langle \$ \rangle (\text{Left } a)$$

$$= \text{Right } (g.h \$ a) \quad \text{apply } (g.h) \langle \$ \rangle$$

$$= \text{Left } a \quad \text{apply } (g.h) \langle \$ \rangle$$

$$= \text{Right } (g \$ h \$ a) \quad \text{composition}$$

$$= (g \langle \$ \rangle) \text{Left } a \quad \text{unapply } g \langle \$ \rangle$$

$$= (g \langle \$ \rangle) \text{Right } (h \$ a) \quad \text{unapply } g \langle \$ \rangle$$

$$= (g \langle \$ \rangle) \$ (h \langle \$ \rangle) \text{Left } a \quad \text{unapply } h \langle \$ \rangle$$

$$= (g \langle \$ \rangle) \$ (h \langle \$ \rangle) \text{Right } a \quad \text{unapply } h \langle \$ \rangle$$

$$= ((g \langle \$ \rangle).(h \langle \$ \rangle)) \text{Left } a \quad \text{unapply composition}$$

$$= ((g \langle \$ \rangle) (h \langle \$ \rangle)) (\text{Right } a) \quad \text{unapply composition}$$

$$\text{Therefore, } P(\text{Left } a)$$

$$\text{Therefore, } P(\text{Right } a)$$

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    fmap f (Right x) = Right (f x)
    fmap f (Left x)  = Left x

instance Applicative (Either e) where
    pure      = Right
    Left e <*> _ = Left e
    Right f <*> r = fmap f r
```

Question 2. [6 MARKS]

Show Either satisfies the *third* applicative law:

$x \text{ <*> pure } y = \text{pure } (\backslash g \rightarrow g \ y) \text{ <*> } x$

$x \text{ <*> pure } y = \text{pure } (\$y) \text{ <*> } x$

Assuming $x = \text{Right } a$

$\text{pure } (\$y) \text{ <*> Right } a$

$= (\$y) \text{ <*> Right } a \quad \text{apply <*>}$

$= \text{Right } ((\$y) a) \quad \text{apply <*>}$

$= \text{Right } (a \ \$y) \quad \text{apply } (\$y)$

$= a \text{ <*> Right } y \quad \text{unapply <*>}$

$= \text{Right } a \text{ <*> Right } y \quad \text{unapply <*>}$

$= x \text{ <*> Right } y \quad \text{by assumption}$

$= x \text{ <*> pure } y \quad \text{unapply pure}$

ZipWith

Recall the alternate definition for a list applicative given in tutorial 8.

```

1 instance Functor [] where
2   fmap _ [] = []
3   fmap g (x:xs) = g x : (fmap g xs)
4
5 instance Applicative [] where
6   pure f = repeat f
7   [] <*> _ = []
8   _ <*> [] = []
9   (f:fs) <*> (x:xs) = (f x) : (fs <*> xs)

```

When writing your proofs use the line numbers given above when justifying your steps.

Question 3. [7 MARKS]

Show your Applicative satisfies the *second* applicative law:

$\text{pure } (g \ x) = \text{pure } g \ \<*> \ \text{pure } x$

Base case

$\text{pure } (g \ []) \quad P(x) : \text{pure } (g \ x) = \text{pure } g \ \<*> \ \text{pure } x$

$= [] \quad \text{line 2 apply } g$

$= (\text{pure } g) \ \<*> \ [] \quad \text{unapply } (\text{pure } g) \ \<*>$

$= (\text{pure } g) \ \<*> \ (\text{pure } []) \quad \text{unapply pure}$

Therefore $P([])$

Induction

suppose $P(xs)$ and prove $P(x:xs)$

H1 $P(xs) : \text{pure } (g \ xs) = \text{pure } g \ \<*> \ \text{pure } xs$

$\text{pure } g \ \<*> \ \text{pure } (x:xs)$

$\text{repeat } (g \ xs) = \text{repeat } g \ \<*> \ \text{repeat } xs$

$= \text{repeat } g \ \<*> \ \text{repeat } (x:xs) \quad \text{apply pure}$

Definition of repeat : $\text{repeat } x = x : (\text{repeat } x)$

$= \text{repeat } g \ \<*> \ (x : \text{repeat } xs) \quad \text{definition of repeat}$

$= \text{repeat } (g \ \$ \ x) : \text{repeat } g \ \<*> \ \text{repeat } xs \quad \text{apply } \<*>$

$= \text{repeat } (g \ \$ \ x) : \text{repeat } (g \ \$ \ xs) \quad \text{apply } \<*> + P(xs) \text{ H1}$

$= \text{repeat } (g \ (x:xs)) \quad \text{apply } g \ \$$

Therefore $P(x:xs)$

ZipWith

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```

When writing your proofs *use the line numbers given above* when justifying your steps.

Question 4. [8 MARKS]

Show your Applicative satisfies the *forth* applicative law:

$$x \text{ <*> } (y \text{ <*> } z) = (\text{pure } (.) \text{ <*> } x \text{ <*> } y) \text{ <*> } z$$

Base case

$$x = [] \quad y = bs \quad z = cs$$

$$[] \text{ <*> } (bs \text{ <*> } cs)$$

$$= [] \quad \text{line 7 apply <*>}$$

$$= [] \text{ <*> } (bs) \quad \text{line 7 unapply <*> bs}$$

$$= (\text{pure } (.) \text{ <*> } []) \text{ <*> } bs \quad \text{line 8 unapply pure } (.) \text{ <*>}$$

$$= (\text{pure } (.) \text{ <*> } []) \text{ <*> } bs \text{ <*> } cs \quad \text{line 7 unapply <*> cs}$$

Loss of Generality, Same progress when $y = []$ or $z = []$

$x \text{ <*> } (y \text{ <*> } z) = (\text{pure } (.) \text{ <*> } x \text{ <*> } y) \text{ <*> } z$ is true for Base case when one of the input is empty.

Induction	$x = (a : a5) \quad y = (b : b5) \quad z = (c : c5)$
	$x \<*\> (y \<*\> z) = (pure \<.\> \<*\> x \<*\> y) \<*\> z$
	Presume $a5 \<*\> (b5 \<*\> z) = (pure \<.\> \<*\> a5 \<*\> b5) \<*\> c5$
	$(pure \<.\> \<*\> (a : a5) \<*\> (b : b5) \<*\> (c : c5))$
	$((\<.\> a : repeat \<.\> \<*\> a5) \<*\> (b : b5) \<*\> (c : c5)) \text{ apply } \<*\>$
	$((\<.\> a b) : (repeat \<.\> \<*\> a5) \<*\> b5)) \<*\> (c : c5) \text{ apply } \<*\>$
	$(a.b) c : (((repeat \<.\> \<*\> a5) \<*\> b5) \<*\> c5) \text{ apply } \<*\>$
	$(a.b) c : a5 \<*\> b5 \<*\> c5 \text{ by H1}$
	$(a \$ b \$ c) : a5 \<*\> b5 \<*\> c5 \text{ composition}$
	$(a : a5) \<*\> ((b \$ c) : (b5 \<*\> c5)) \text{ line 9 unapply } \<*\>$
	$(a : a5) \<*\> ((b : b5) \<*\> (c : c5)) \text{ line 9 unapply } \<*\>$
	Therefore , $(a : a5) \<*\> ((b : b5) \<*\> (c : c5))$
	$= (pure \<.\> \<*\> (a : a5)) \<*\> (b : b5) \<*\> (c : c5)$
	$= x \<*\> (y \<*\> z)$