Architecture Synthesis for Linear Time-Invariant Filters

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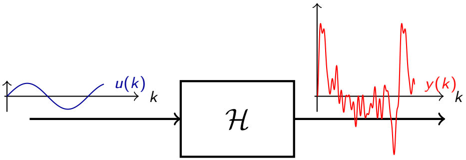
Introduction

Fundamentals: signal processing and filters

LTI Filters

LTI (Linear Time Invariant) filters are particular filters that are:

- Linear: outputs are linear combinations of inputs (allows to use linear algebra definitions)
- Time-Invariant: all coefficients are constant



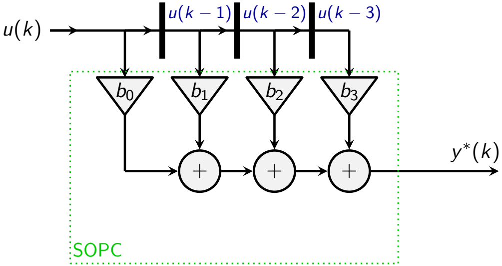
What is the purpose of this work?

- Lopez's PhD thesis states how to compute LTI filters in software:
  - ordonancing issues
  - fixed size
- The goal here is to do such a work in hardware, where we have more flexibility:
  - full parallelism
  - arbitrary size

In this context, constraints become degrees of freedom.

FIR definition:

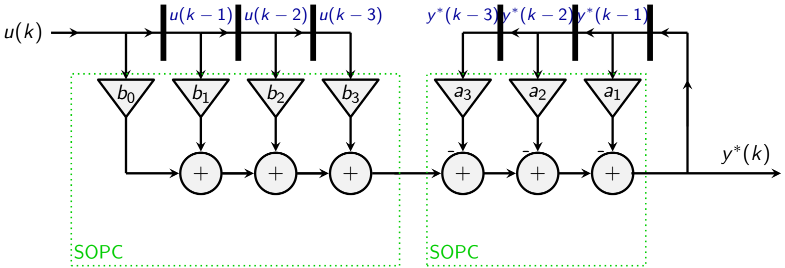
$$y(k) = \sum_{i=0}^{n} b_i u(k-i) \tag{1}$$



Abstract architecture for the direct form realization of an LTI filter

IIR definition:

$$y(k) = \sum_{i=0}^{n} b_i u(k-i) - \sum_{i=0}^{n} a_i y(k-i) \quad (1)$$



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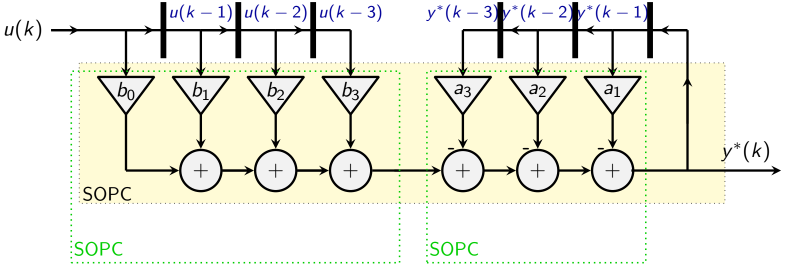

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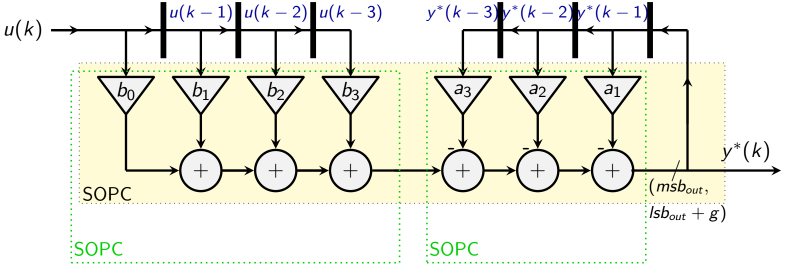


Abstract architecture for the direct form realization of an LTI filter

Here we have loop-backs, so registers:

$$\text{Loop-backs} \iff \text{Registers}$$

We can generalise by introducing state registers.

State-Space representation

The "ABCD" form

Let's define $\boldsymbol{x}(k)$ a state vector (hardware register)

$$\begin{cases} \boldsymbol{x}(k+1) = \boldsymbol{Ax}(k) + \boldsymbol{Bu}(k) \\ \boldsymbol{y}(k+1) = \boldsymbol{Cx}(k) + \boldsymbol{Du}(k) \end{cases} \tag{2}$$

With:

$$\boldsymbol{A} \in \mathbb{R}^{n_x \times n_x} \ , \ \boldsymbol{B} \in \mathbb{R}^{n_x \times n_u} \ , \\ \boldsymbol{C} \in \mathbb{R}^{n_y \times n_x} \ , \ \boldsymbol{D} \in \mathbb{R}^{n_y \times n_u}$$

Equivalent matrix formulation:

$$\begin{pmatrix} \boldsymbol{x}(k+1) \\ \boldsymbol{y}(k+1) \end{pmatrix} = \begin{pmatrix} \boldsymbol{A} & \boldsymbol{B} \\ \boldsymbol{C} & \boldsymbol{D} \end{pmatrix} \begin{pmatrix} \boldsymbol{x}(k) \\ \boldsymbol{u}(k) \end{pmatrix} \tag{3}$$

Definition of the SIF

Problems with ABCD:

- *lsb* and *msb* computations have to be rebuilt for each new filter in this form.
- this doesn't gives an explicit order in the operations

The SIF generalizes the state-space.
Addition: $\boldsymbol{t}(k)$ describes the operations order:

$$\begin{pmatrix} \boldsymbol{J} & \boldsymbol{0} & \boldsymbol{0} \\ -\boldsymbol{K} & \boldsymbol{I}_{n_x} & \boldsymbol{0} \\ -\boldsymbol{L} & \boldsymbol{0} & \boldsymbol{I}_{n_y} \end{pmatrix} \begin{pmatrix} \boldsymbol{t}(k+1) \\ \boldsymbol{x}(k+1) \\ \boldsymbol{y}(k) \end{pmatrix} = \begin{pmatrix} \boldsymbol{0} & \boldsymbol{M} & \boldsymbol{N} \\ \boldsymbol{0} & \boldsymbol{P} & \boldsymbol{Q} \\ \boldsymbol{0} & \boldsymbol{R} & \boldsymbol{S} \end{pmatrix} \begin{pmatrix} \boldsymbol{t}(k) \\ \boldsymbol{x}(k) \\ \boldsymbol{u}(k) \end{pmatrix} \quad (4)$$

The SIF as an algorithm

for *int* $i = 0$; $i \leq n_t$; i++ **do**

$$\boldsymbol{t}_i(k+1) \leftarrow -\sum_{j<i} \boldsymbol{J}_{ij}\boldsymbol{t}_j(k+1) + \sum_{j=1}^{n_x} \boldsymbol{M}_{ij}\boldsymbol{x}_j(k) + \sum_{j=1}^{n_u} \boldsymbol{N}_{ij}\boldsymbol{u}_j(k)$$

end

for *int* $i = 0$; $i \leq n_x$; i++ **do**

$$\boldsymbol{x}_i(k+1) \leftarrow \sum_{j=1}^{n_t} \boldsymbol{K}_{ij}\boldsymbol{t}_j(k+1) + \sum_{j=1}^{n_x} \boldsymbol{P}_{ij}\boldsymbol{x}_j(k) + \sum_{j=1}^{n_u} \boldsymbol{Q}_{ij}\boldsymbol{u}_j(k)$$

end

for *int* $i = 0$; $i \leq ny$; i++ **do**

$$\boldsymbol{y}_i(k) \leftarrow \sum_{j=1}^{n_t} \boldsymbol{L}_{ij}\boldsymbol{t}_j(k+1) + \sum_{j=1}^{n_x} \boldsymbol{R}_{ij}\boldsymbol{x}_j(k) + \sum_{j=1}^{n_u} \boldsymbol{S}_{ij}\boldsymbol{u}_j(k)$$

end

Algorithm 1: Computation of SIF outputs from inputs whatever

Error definition

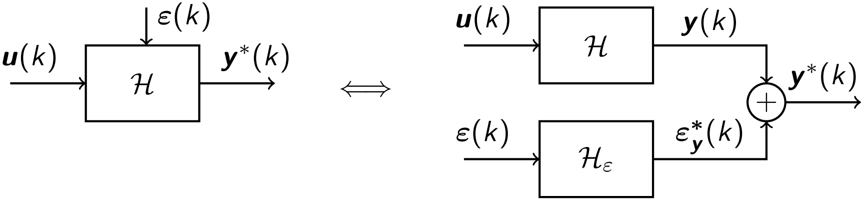
Let's define: $v' = n_t + n_x + n_y$

ξ will be the desired error. Errors introduced by SOPCs:

$$\varepsilon'_v(k) = \begin{pmatrix} \varepsilon_t(k) \\ \varepsilon_x(k) \\ \varepsilon_y(k) \end{pmatrix} = \begin{pmatrix} \varepsilon_{t_1}(k) \\ \varepsilon_{t_2}(k) \\ \vdots \\ \varepsilon_{t_{n_t}}(k) \\ \varepsilon_{x_1}(k) \\ \varepsilon_{x_2}(k) \\ \vdots \\ \varepsilon_{x_{n_x}}(k) \\ \varepsilon_{y_1}(k) \\ \varepsilon_{y_2}(k) \\ \vdots \\ \varepsilon_{y_{n_y}}(k) \end{pmatrix} \tag{5}$$

Total error: $\varepsilon^*_{v'}(k)$ defined as $\varepsilon'_v(k)$

Point of view about error



A signal view of the error propagation with respect to the ideal filter

$$|\langle\langle \mathcal{H}_\varepsilon \rangle\rangle| \cdot 2^{lsb_{v'}+1} < \xi_i \tag{6}$$

Definition of the Worst-Case-Peak-Gain (WCPG)

$$\|\mathcal{H}\|_{wcpg} = \sup_{u \neq 0} \frac{\|h * u\|_{l^\infty}}{\|u\|_{l^\infty}} \tag{7}$$

M'sieu on a rien fait c'est pas nous.

Here we use FloPoCo, which is a C++ framework which first purpose is to generate floating point cores in VHDL. It is described

in Florent de Dinechin and Bogdan Pasca. Designing custom arithmetic data paths with FloPoCo.
IEEE Design & Test of Computers, 28(4):18–27, July 2011
http://flopoco.gforge.inria.fr/

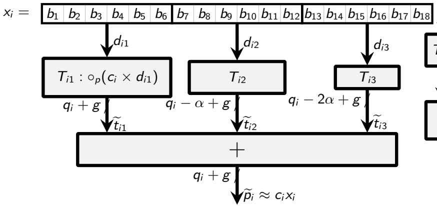


Figure: The FixRealKCM method when x_i is split in 3 chunks

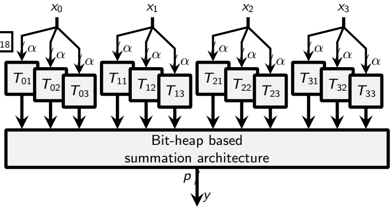


Figure: KCM-based SPOC architecture for $n_c = 4$, each input being split into 3 chunks

Architecture generation algorithm

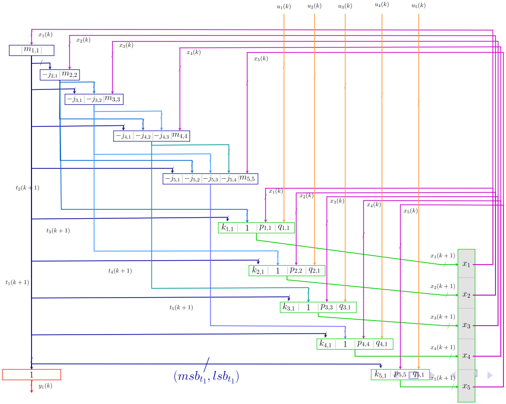
computePrecisions([*msbs*, *lsbs*][][]) //get the matrix of msbs lsbs, functions of the wcpg.

```
for i=1; i=Z.size(); i++ do
    row[ ]= Z[i][ ] //pick first row of Z
    for j=1; j=1; j=Z.size() j++ do
        assign(SOPC[i], row[j], "T","X","U",[msbs,lsbs][i][j])
    end
    Second pass for wiring.
end
```

Algorithm 2: Architecture Generation Algorithm

Example

$$Z = \begin{pmatrix}
1 & 0 & 0 & 0 & 0 & m_{1,1} & 0 & 0 & 0 & 0 & 0 \\
j_{2,1} & 1 & 0 & 0 & 0 & 0 & m_{2,2} & 0 & 0 & 0 & 0 \\
j_{3,1} & j_{3,2} & 1 & 0 & 0 & 0 & 0 & m_{3,3} & 0 & 0 & 0 \\
j_{4,1} & j_{4,2} & j_{4,3} & 1 & 0 & 0 & 0 & 0 & m_{4,4} & 0 & 0 \\
j_{5,1} & j_{5,2} & j_{5,3} & j_{5,4} & 1 & 0 & 0 & 0 & 0 & m_{5,5} & 0 \\
k_{1,1} & 1 & 0 & 0 & 0 & p_{1,1} & 0 & 0 & 0 & 0 & q_{1,1} \\
k_{2,1} & 0 & 1 & 0 & 0 & 0 & p_{2,2} & 0 & 0 & 0 & q_{2,1} \\
k_{3,1} & 0 & 0 & 1 & 0 & 0 & 0 & p_{3,3} & 0 & 0 & q_{3,1} \\
k_{4,1} & 0 & 0 & 0 & 1 & 0 & 0 & 0 & p_{4,4} & 0 & q_{4,1} \\
k_{5,1} & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & p_{5,5} & q_{5,1} \\
1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0
\end{pmatrix}$$



Example

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j_{3,1} & j_{3,2} & 1 & 0 & 0 & 0 & 0 & m_{3,3} & 0 & 0 & 0 \\
j_{4,1} & j_{4,2} & j_{4,3} & 1 & 0 & 0 & 0 & 0 & m_{4,4} & 0 & 0 \\
j_{5,1} & j_{5,2} & j_{5,3} & j_{5,4} & 1 & 0 & 0 & 0 & 0 & m_{5,5} & 0 \\
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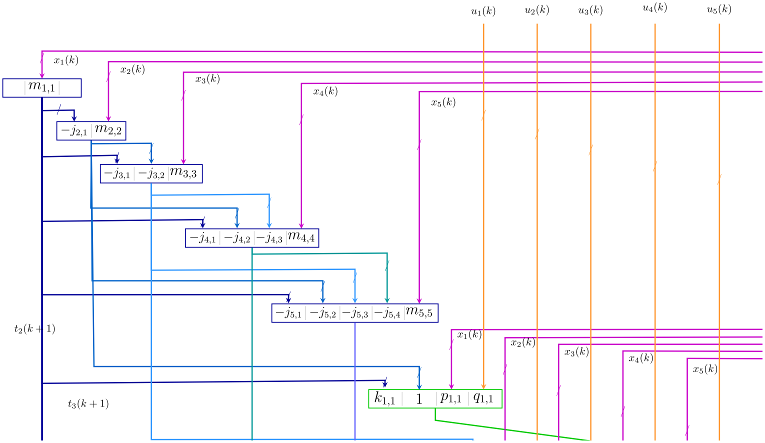
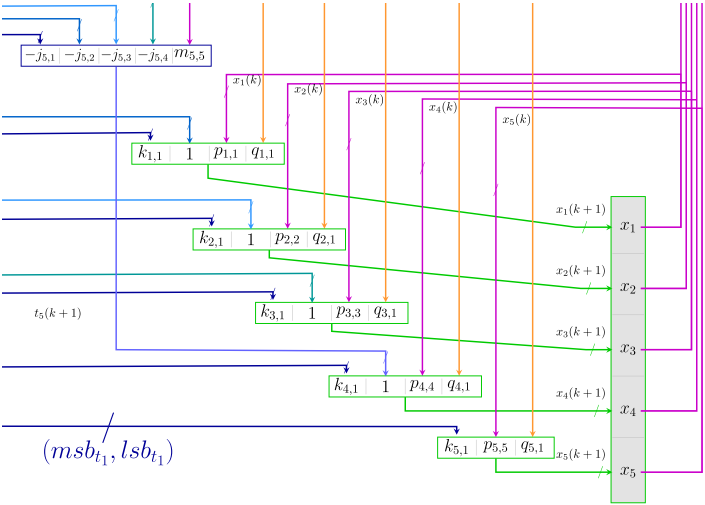
k_{2,1} & 0 & 1 & 0 & 0 & 0 & p_{2,2} & 0 & 0 & 0 & q_{2,1} \\
k_{3,1} & 0 & 0 & 1 & 0 & 0 & 0 & p_{3,3} & 0 & 0 & q_{3,1} \\
k_{4,1} & 0 & 0 & 0 & 1 & 0 & 0 & 0 & p_{4,4} & 0 & q_{4,1} \\
k_{5,1} & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & p_{5,5} & q_{5,1} \\
1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0
\end{pmatrix}$$

Example


$u_1(k)$
$u_2(k)$
$u_3(k)$
$u_4(k)$
$u_5(k)$
$x_1(k)$
$x_2(k)$
$x_3(k)$
$x_4(k)$
$x_5(k)$
$m_{1,1}$
$-j_{2,1}$ $m_{2,2}$
$-j_{3,1}$ $-j_{3,2}$ $m_{3,3}$
$-j_{4,1}$ $-j_{4,2}$ $-j_{4,3}$ $m_{4,4}$
$-j_{5,1}$ $-j_{5,2}$ $-j_{5,3}$ $-j_{5,4}$ $m_{5,5}$
$t_2(k+1)$
$t_3(k+1)$
$k_{1,1}$ 1 $p_{1,1}$ $q_{1,1}$

Example


$-j_{5,1}$ $-j_{5,2}$ $-j_{5,3}$ $-j_{5,4}$ $m_{5,5}$
$x_1(k)$
$x_2(k)$
$x_3(k)$
$x_4(k)$
$x_5(k)$
$k_{1,1}$ 1 $p_{1,1}$ $q_{1,1}$
$k_{2,1}$ 1 $p_{2,2}$ $q_{2,1}$
$k_{3,1}$ 1 $p_{3,3}$ $q_{3,1}$
$k_{4,1}$ 1 $p_{4,4}$ $q_{4,1}$
$k_{5,1}$ $p_{5,5}$ $q_{5,1}$
$t_5(k+1)$
(msb_{t_1}, lsb_{t_1})
$x_1(k+1)$
$x_2(k+1)$
$x_3(k+1)$
$x_4(k+1)$
$x_5(k+1)$
x_1
x_2
x_3
x_4
x_5

Future Work

- Removal of power of two: adapt KCM and SOPC FloPoCo cores
- Sub-filter detection: open question for either Front-End and Back-End
- Precision calculations improvement
- File format re-specification

Conclusion

To conclude

References

Small bibliography

Florent de Dinechin and Bogdan Pasca.
Designing custom arithmetic data paths with FloPoCo.
IEEE Design & Test of Computers, 28(4):18–27, July 2011.

Any question?