Architecture Synthesis for Linear Time-Invariant Filters

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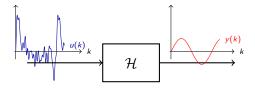
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Introduction: LTI Filters

Signal: temporal variable x(k) with $\{x(k)\}_{k\geq 0}\in\mathbb{R}$



LTI (Linear Time Invariant) filters are particular filters that are:

- Linear: outputs are linear combinations of inputs (allows to use linear algebra definitions)
- Time-Invariant: all coefficients are constant



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Purpose of this work

What is the purpose of this work?

- Lopez's PhD thesis states how to compute LTI filters in software:
 - scheduling issues
 - fixed size
- The goal here is to do such a work in hardware, where we have more flexibility:
 - full parallelism
 - arbitrary size

In this context, constraints become degrees of freedom.



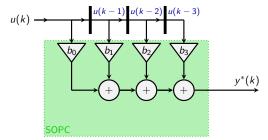
FIR, IIR

Signal processing and filters

FIR definition:

$$y(k) = \sum_{i=0}^{n} b_i u(k-i)$$
 (1)

Details of the implementation



Abstract architecture for the direct form realization of an FIR filter

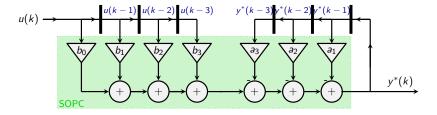


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Signal processing and filters

IIR definition:

$$y(k) = \sum_{i=0}^{n} b_i u(k-i) - \sum_{i=0}^{n} a_i y(k-i)$$
 (1)



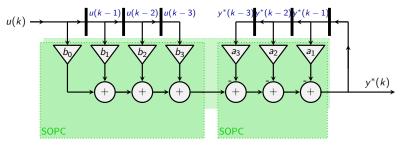
Abstract architecture for the direct form realization of an IIR filter



FIR, IIR

IIR definition:

$$y(k) = \sum_{i=0}^{n} b_i u(k-i) - \sum_{i=0}^{n} a_i y(k-i)$$
 (1)



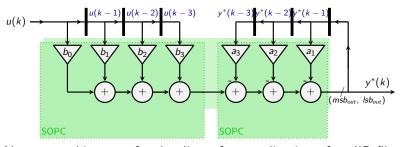
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IIR definition:

$$y(k) = \sum_{i=0}^{n} b_i u(k-i) - \sum_{i=0}^{n} a_i y(k-i)$$
 (1)



Abstract architecture for the direct form realization of an IIR filter



State-Space representation The "ABCD" form

Let's define x(k) a state vector (hardware register)

$$\begin{cases} x(k+1) = \mathbf{A}x(k) + \mathbf{B}\mathbf{u}(k) \\ y(k+1) = \mathbf{C}x(k) + \mathbf{D}\mathbf{u}(k) \end{cases}$$
(2)

Details of the implementation

With:

$$m{A} \in \mathbb{R}^{n_{x} imes n_{x}}$$
 , $m{B} \in \mathbb{R}^{n_{x} imes n_{u}}$, $m{C} \in \mathbb{R}^{n_{y} imes n_{x}}$, $m{D} \in \mathbb{R}^{n_{y} imes n_{u}}$

Equivalent matrix formulation:

$$\begin{pmatrix} \mathbf{x}(k+1) \\ \mathbf{y}(k+1) \end{pmatrix} = \begin{pmatrix} \mathbf{A} & \mathbf{B} \\ \mathbf{C} & \mathbf{D} \end{pmatrix} \begin{pmatrix} \mathbf{x}(k) \\ \mathbf{u}(k) \end{pmatrix}$$
 (3)

The SIF: a unified realization representation

Problems with ABCD:

- Isb and msb computations have to be rebuilt for each new filter in this form.
- this doesn't gives an explicit order in the operations

The SIF generalizes the state-space.

Addition: t(k) describes the operations order:

$$\begin{pmatrix} J & \mathbf{0} & \mathbf{0} \\ -K & I_{n_x} & \mathbf{0} \\ -L & \mathbf{0} & I_{n_y} \end{pmatrix} \begin{pmatrix} \mathbf{t}(k+1) \\ \mathbf{x}(k+1) \\ \mathbf{y}(k) \end{pmatrix} = \begin{pmatrix} \mathbf{0} & \mathbf{M} & \mathbf{N} \\ \mathbf{0} & P & Q \\ \mathbf{0} & R & S \end{pmatrix} \begin{pmatrix} \mathbf{t}(k) \\ \mathbf{x}(k) \\ \mathbf{u}(k) \end{pmatrix} \quad (4)$$

The SIF as an algorithm

end

for int
$$i = 0$$
; $i \le n_x$; $i++$ do
$$x_i(k+1) \leftarrow \sum_{j=1}^{n_t} K_{ij} t_j(k+1) + \sum_{j=1}^{n_x} P_{ij} x_j(k) + \sum_{j=1}^{n_u} Q_{ij} u_j(k)$$

end

end

Algorithm 1: Computation of SIF outputs from inputs



A more common notation

An easiest way to communicate SIFs: The Z matrix:

$$Z = \begin{pmatrix} -J & M & N \\ K & P & Q \\ L & R & S \end{pmatrix}$$
 (5)

Workflow

Signal processing and filters

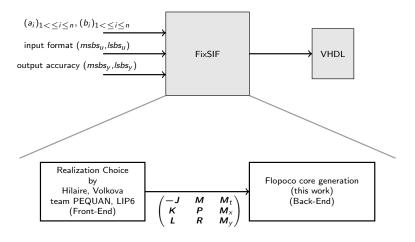


Figure: Workflow overview of tools usage



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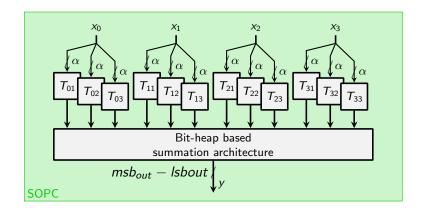
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FloPoCo:

- C++ framework
- Target: FPGAs
- Work: generating arithmetical cores in VHDL computing just right
- Reference: http://flopoco.gforge.inria.fr/

Basic block: SOPC



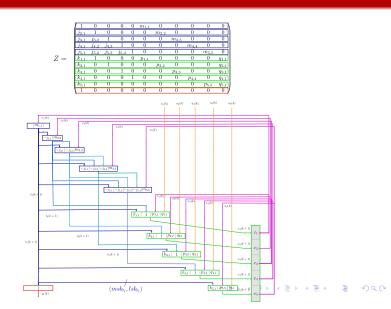
SOPC architecture: based on the KCM architecture, split in 3 chunks



Architecture generation algorithm

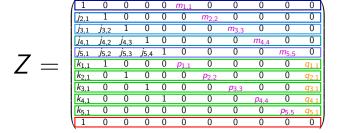
```
computeMSBSLSBS([msbs, lsbs][][]) //get the matrix of msbs lsbs.
for i=1; i=Z.size(); i++ do
   row[] = Z[i][] //pick first row of Z
   for j=1; j=1; j=Z.size() j++ do
      assign(SOPC[i], row[i], "T","X","U",[msbs,lsbs][i][j])
   end
   Second pass for wiring.
end
        Algorithm 2: Architecture Generation Algorithm
```

Example

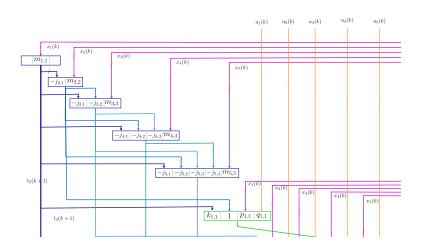


Example

Each line will be a SOPC



Example



Example

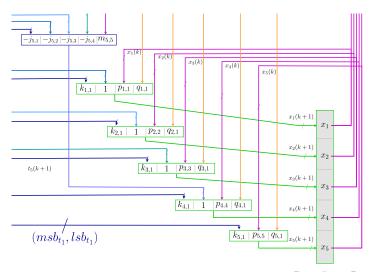
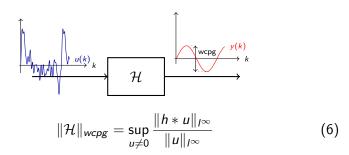


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Computing the MSB



This computation is done by colleagues in LIP6.

Computing the LSB: Error definition

Let's define:

$$\mathbf{v'} = \begin{pmatrix} \mathbf{t}(k+1) \\ \mathbf{x}(k+1) \\ \mathbf{y}(k) \end{pmatrix} \tag{7}$$

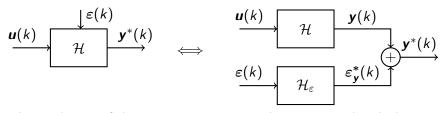
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Errors introduced by SOPCs:

$$\boldsymbol{\varepsilon}_{v'}(k) = \begin{pmatrix} \varepsilon_{t}(k) \\ \varepsilon_{t}(k) \\ \vdots \\ \varepsilon_{t_{n_{t}}}(k) \\ \varepsilon_{x}(k) \\ \varepsilon_{y}(k) \end{pmatrix} = \begin{pmatrix} \varepsilon_{t_{1}}(k) \\ \varepsilon_{x_{1}}(k) \\ \varepsilon_{x_{2}}(k) \\ \vdots \\ \varepsilon_{x_{n_{x}}}(k) \\ \varepsilon_{y_{1}}(k) \\ \varepsilon_{y_{2}}(k) \\ \vdots \\ \varepsilon_{y_{n_{v}}}(k) \end{pmatrix}$$

(8)

Computing the LSB: Point of view about error

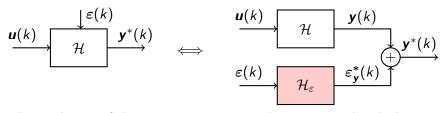


A signal view of the error propagation with respect to the ideal filter

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Signal processing and filters

Computing the LSB: Point of view about error



A signal view of the error propagation with respect to the ideal filter

Computing the LSB: Impact of errors

$$\begin{array}{l} \textbf{for int } i = 0 \text{ ; } i \leq n_t \text{; } i + + \textbf{ do} \\ & \quad t_i(k+1) \leftarrow \\ & \quad - \sum\limits_{j < i} \textbf{\textit{J}}_{ij} \textbf{\textit{t}}_j(k+1) + \sum\limits_{j=1}^{n_x} \textbf{\textit{M}}_{ij} \textbf{\textit{x}}_j(k) + \sum\limits_{j=1}^{n_u} \textbf{\textit{N}}_{ij} \textbf{\textit{u}}_j(k) + \boldsymbol{\varepsilon}_{t_i}(k) \end{array}$$

end

for int i = 0; $i \le n_x$; i++ do

$$m{x}_i(k+1) \leftarrow \sum\limits_{j=1}^{n_t} m{K}_{ij} m{t}_j(k+1) + \sum\limits_{j=1}^{n_x} m{P}_{ij} m{x}_j(k) + \sum\limits_{j=1}^{n_u} m{Q}_{ij} m{u}_j(k) + m{arepsilon}_{m{x}_i}(m{k})$$

end

for int i = 0; $i \le ny$; i++ do

$$m{y}_i(k) \leftarrow \sum\limits_{j=1}^{n_t} m{L}_{ij} m{t}_j(k+1) + \sum\limits_{j=1}^{n_{\mathsf{x}}} m{R}_{ij} m{x}_j(k) + \sum\limits_{j=1}^{n_u} m{S}_{ij} m{u}_j(k) + m{arepsilon}_{y_i}(k)$$

end

Algorithm 3: Computation of SIF outputs from inputs

Computing the LSB: current solution

Main constraint:

$$\varepsilon_{\nu} < \mathbf{2}^{-lsb_{\nu}}$$
 (9)

Transformed into the following constraint:

$$|\langle\langle\mathcal{H}_{\varepsilon}\rangle\rangle|\cdot\mathbf{2}^{lsb_{v'}+1}<\mathbf{2}^{-lsb_{y_i}}$$
 (10)

Lopez gives a simple solution that matches the fixed size constraint.

In the present case, a better solution is achievable at the end of this work.

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Future Work

- Removal of power of two: adapt KCM and SOPC FloPoCo cores
- Sub-filter detection: open question for either Front-End and Back-End
- Precision calculations improvement
- File format re-specification

Conclusion

To conclude

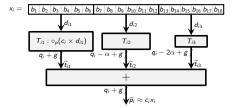
- Adapting LTI Filters generation from software to hardware
- Reuse of Lopez's calculations in our context
- Implementation for FPGAs in a parametric view

Any question?

$$\begin{pmatrix} \mathbf{J} & \mathbf{0} & \mathbf{0} \\ -\mathbf{K} & \mathbf{I}_{n_{x}} & \mathbf{0} \\ -\mathbf{L} & \mathbf{0} & \mathbf{I}_{n_{y}} \end{pmatrix} \begin{pmatrix} \mathbf{t}(k+1) \\ \mathbf{x}(k+1) \\ \mathbf{y}(k) \end{pmatrix} = \begin{pmatrix} \mathbf{0} & \mathbf{M} & \mathbf{N} \\ \mathbf{0} & \mathbf{P} & \mathbf{Q} \\ \mathbf{0} & \mathbf{R} & \mathbf{S} \end{pmatrix} \begin{pmatrix} \mathbf{t}(k) \\ \mathbf{x}(k) \\ \mathbf{u}(k) \end{pmatrix}$$
(11)

With:

$$\mathbf{J} \in \mathbb{R}^{n_t \times n_t}, \mathbf{M} \in \mathbb{R}^{n_t \times n_x}, \mathbf{N} \in \mathbb{R}^{n_t \times n_u},
\mathbf{K} \in \mathbb{R}^{n_x \times n_t}, \mathbf{P} \in \mathbb{R}^{n_x \times n_x}, \mathbf{Q} \in \mathbb{R}^{n_x \times n_u},
\mathbf{L} \in \mathbb{R}^{n_y \times n_t}, \mathbf{R} \in \mathbb{R}^{n_y \times n_x}, \mathbf{S} \in \mathbb{R}^{n_y \times n_u},$$
(12)



The FixRealKCM method when x_i is split in 3 chunks

Mathematical definition of a filter \mathcal{H}

Definition of a filter: $\mathbf{y} = \mathcal{H}(\mathbf{u})$ With $dim(\mathbf{y}) = n_{v}$ and, $dim(\mathbf{u}) = n_{ii}$ Linearity:

$$\mathcal{H}(\alpha \cdot \mathbf{u}_1 + \beta \cdot \mathbf{u}_2) = \alpha \cdot \mathcal{H}(\mathbf{u}_1) + \beta \cdot \mathcal{H}(\mathbf{u}_2)$$

Time invariance:

$$\{\mathcal{H}(\mathbf{u})(k-k_0)\}_{k\geq 0} = \mathcal{H}(\{\mathbf{u}(k-k_0)\}_{k\geq 0})$$

$$u = \sum_{i>0} u(1)\delta_i$$

Where δ_I is a Dirac impulsion centered in I:

$$\delta_I(k) = \begin{cases} 1 & \text{when } k = I \\ 0 & \text{else} \end{cases}$$
 (14)

Time invariance gives: $\mathcal{H}(\delta_I)(k) = h(k-I)$ Computation of the outputs:

$$y_i(k) = \sum_{j=1}^{n_u} \sum_{l=0}^{k} u_j(l) h_{i,j}(k-l), \ \forall 1 \leq i \leq n_y$$

From:

$$Z = \begin{pmatrix} -J & M & N \\ K & P & Q \\ L & R & S \end{pmatrix}$$
 (15)

The ABCD form is deducible from the SIF:

$$A_{Z} = KJ^{-1}M + P, \ B_{Z} = KJ^{-1}N + Q,$$
 $C_{Z} = LJ^{-1}M + R, \ D_{Z} = LJ^{-1}N + S,$
(16)

with:

$$\mathbf{A}_{\mathbf{Z}} \in \mathbb{R}^{n_{\mathbf{x}} \times n_{\mathbf{x}}}, \mathbf{B}_{\mathbf{Z}} \in \mathbb{R}^{n_{\mathbf{x}} \times n_{\mathbf{u}}},$$
 $\mathbf{C}_{\mathbf{Z}} \in \mathbb{R}^{n_{\mathbf{y}} \times n_{\mathbf{x}}}, \mathbf{D}_{\mathbf{Z}} \in \mathbb{R}^{n_{\mathbf{y}} \times n_{\mathbf{u}}},$

$$(17)$$

$$Z_{\varepsilon} = \begin{pmatrix} -J & M & M_t \\ K & P & M_x \\ L & R & M_y \end{pmatrix}$$
 (18)

with:

$$\mathbf{M}_{t} = (\mathbf{I}_{n_{t}} \ \mathbf{0}_{n_{t} \times n_{x}} \ \mathbf{0}_{n_{t} \times n_{y}}), \tag{19}$$

$$\mathbf{M}_{\mathsf{X}} = (\mathbf{0}_{n_{\mathsf{X}} \times n_{\mathsf{t}}} \ \mathbf{I}_{n_{\mathsf{X}}} \ \mathbf{0}_{n_{\mathsf{X}} \times n_{\mathsf{y}}}), \tag{20}$$

$$\mathbf{M}_{\mathbf{y}} = (\mathbf{0}_{n_{\mathbf{y}} \times n_{\mathbf{t}}} \ \mathbf{0}_{n_{\mathbf{y}} \times n_{\mathbf{x}}} \ \mathbf{I}_{n_{\mathbf{y}}}), \tag{21}$$

Interface specification:

X I c

x_1_1 x_1_2 ... x_1_c

 $x_2_1 x_2_2 \dots x_2_c$

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x_l_1 x_l_2 ... x_l_c

Where:

- X = name of the matrix
- $\bullet \ \ X \in \{ \ J, \ K, \ L, \ M, \ N, \ P, \ Q, \ R, \ S, \ T \}$
- x_i_j = the coefficient
- $i \in [1, I]$
- $j \in [1, c]$
- I = the number of lines
- c = the numer of columns.



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