



Added handouts for Week02

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more-typeclasses.md 5.75 KB

More on type classes and instances

In this section, we study in detail two important type classes:

1. The type class `Ord` for ordered types.
2. The type class `Num` for numeric types.

We also study an example of an *instance declaration with constraints*: `instance Ord a => Ord [a]`. This instance tells us that whenever `a` is an instance of `Ord`, then so is `[a]`.

The type `Ordering` and the typeclass `Ord`

There is also a [video](#) on this section.

The type class `Ord` implements the idea that elements of a type can be **compared not only for equality, but also for less/greater than**. A comparison on a type `a` is a map `compare :: a -> a -> Ordering`, where the type `Ordering` is defined as follows:

```
data Ordering = LT | EQ | GT
              deriving (Eq, Ord, Enum, Read, Show, Bounded)
```

Here, `LT` stands for *less than*, `EQ` stands for *equal*, and `GT` stands for *greater than*.

The following defines the `Ord` class for types `a` which are in the `Eq` class.

```
class (Eq a) => Ord a where
  compare :: a -> a -> Ordering
  (<), (<=), (>=), (>) :: a -> a -> Bool
  max, min :: a -> a -> a

  -- Minimal complete definition:
  --   (<=) or compare
  -- Using compare can be more efficient for complex types.
  compare x y
    | x == y    = EQ
    | x <= y    = LT
    | otherwise = GT

  x <= y        = compare x y /= GT
  x < y         = compare x y == LT
  x >= y        = compare x y /= LT
  x > y         = compare x y == GT

-- note that (min x y, max x y) = (x,y) or (y,x)
max x y
  | x <= y    = y
  | otherwise = x
min x y
  | x <= y    = x
  | otherwise = y
```

There is a seeming circularity between the definition of `Ordering` and `Ord`, in that each one refers to the other. We may see this as a mutually recursive definition.

For example, the `compare` function on lists over `a` (where `a` is already equipped with a function `compare`) is implemented as follows:

```
instance (Ord a) => Ord [a] where
  compare [] []      = EQ
  compare [] (_,_)   = LT
  compare (_,_) []   = GT
  compare (x:xs) (y:ys) = case compare x y of
```

```
EQ    -> compare xs ys
other -> other
```

Exercises

1. From reading the code, explain how two lists are compared.
2. Run some examples of comparisons between lists to confirm or refute your explanation.

This [video](#) gives an explanation of the implementation of `compare :: [a] -> [a] -> Ordering`.

The type class `Num`

Consider the following examples

```
Prelude> 5 + 3
8
Prelude> 3.14159 + 2.71828
5.85987
```

Here, the operator `+` acts on any type that is an instance of the `Num` typeclass: for any such type `a`, the function `(+)` is of type `a -> a -> a`, used as an infix operator.

We can use the `:info` command to find out what operations the `Num` typeclass provides:

```
Prelude> :info Num
class Num a where
  (+) :: a -> a -> a
  (-) :: a -> a -> a
  (*) :: a -> a -> a
  negate :: a -> a
  abs :: a -> a
  signum :: a -> a
  fromInteger :: Integer -> a
  {-# MINIMAL (+), (*), abs, signum, fromInteger, (negate | (-)) #-}
  -- Defined in ‘GHC.Num’
instance Num Word -- Defined in ‘GHC.Num’
instance Num Integer -- Defined in ‘GHC.Num’
instance Num Int -- Defined in ‘GHC.Num’
instance Num Float -- Defined in ‘GHC.Float’
instance Num Double -- Defined in ‘GHC.Float’
```

This shows that any type `a` that is an instance of `Num` comes with operations `(+)`, `(-)`, `(*)`, ..., `fromInteger`. (In particular, `1` without a type annotation is not an integer, but instead an element of any type `a` that is an instance of `Num`; specifically, the image of the integer `1 :: Integer` under the function `fromInteger :: Integer -> a`.)

We also see that there are five **instances** of the type class `Num` defined, for the types `Word`, `Integer`, `Int`, `Float`, and `Double`. When we want to consider `1` as an element of a particular numeric type, we can do that by **annotating** it with that type, e.g.,

```
Prelude> :type 1 :: Word
1 :: Word :: Word
```

Here, we **check** that `1 :: Word` instead of asking `ghci` to **infer** the type of `1` for us. `ghci` then just needs to check that `Word` is an instance of `Num`. Similarly, for `(+)` we can check

```
Prelude> :type (+) :: Integer -> Integer -> Integer
(+) :: Integer -> Integer -> Integer
      :: Integer -> Integer -> Integer
```

This check fails for types for which no instance of `Num` has been declared, e.g., for `Char`:

```
Prelude> :type (+) :: Char -> Char -> Char

<interactive>:1:1: error:
  • No instance for (Num Char) arising from a use of `+`
  • In the expression: (+) :: Char -> Char -> Char
```

Quiz time

Test your understanding by taking this [quiz](#). Don't worry, it is not marked, and you can take it as many times as you want.

See also

3. A blog post comparing Java and Haskell: <https://mmhaskell.com/blog/2019/1/28/why-haskell-iv-typeclasses-vs-inheritance>. Do you agree with it? Why (not)?

Summary

1. We have studied in detail one function that uses both pattern matching on lists and a `case` expression (a more general form of pattern matching).
2. We have seen how instances can be derived automatically from other instances, using the example `instance Ord a => Ord [a]`.
3. We have taken a look at the type class `Num` for number types.
4. We have seen how to use a type annotation to force a Haskell expression to have a particular type, e.g., `1 :: Word`.