

Tree Path Labeling of Path Hypergraphs

Generalization of the Consecutive-ones Property

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① Introduction

An Illustration

Terminology

Motivation

② Characterization of a feasible TPL

ICPPL

Filtering algorithm

③ Computing a feasible TPL on k -subdivided trees

Algorithm

④ Conclusion

Application

An Illustration

- To introduce the combinatorial problem of TPL.

Study Group Accommodation problem

- A set of n **students** arrive for a summer course, say $\{a, b, c, d, e, f, g, h, i, j, k\}$, $n = 11$.

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- There are n single occupancy **apartments** in the university campus for their accommodation.
- All these apartments are placed such that streets connecting them do not form loops - **streets form a tree**

Study Group Accommodation problem

The problem

How should the students be allocated apartments such that:

- students of each study group are neighbours?

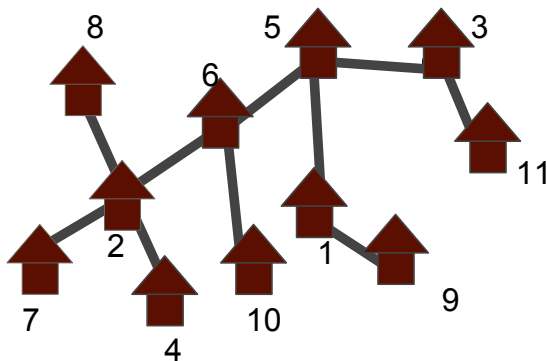
Study Group Accommodation problem

The problem

How should the students be allocated apartments such that:

- students of each study group are neighbours?
- i.e. a study group forms a **path in the tree**.

Study Group Accommodation problem



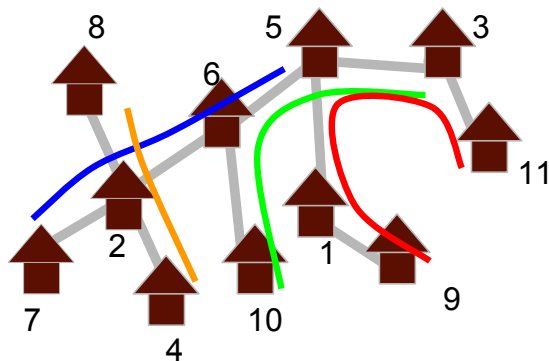
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Study Group Accommodation problem



$$R = \{g, h, i, j, k\}$$

$$\rightarrow \{9, 1, 5, 3, 11\}$$

$$B = \{a, b, e, g\}$$

$$\rightarrow \{7, 2, 6, 5\}$$

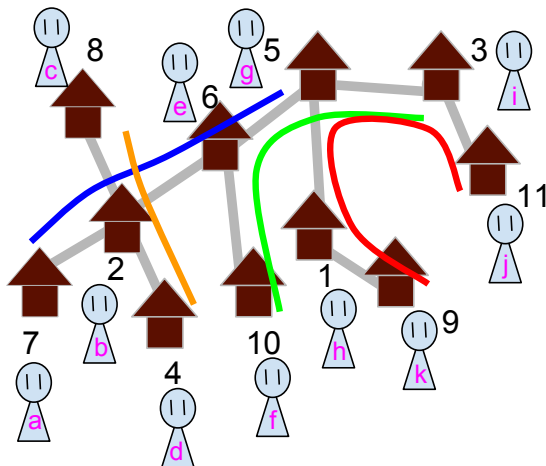
$$O = \{c, b, d\}$$

$$\rightarrow \{4, 2, 8\}$$

$$G = \{e, f, g, i\}$$

$$\rightarrow \{10, 6, 5, 3\}$$

Study Group Accommodation problem



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$\rightarrow \{9, 1, 5, 3, 11\}$

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$\rightarrow \{7, 2, 6, 5\}$

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Tree Path Labeling of Path Hypergraphs

- The set of study groups \rightarrow **Set system / Hypergraph**

Tree Path Labeling of Path Hypergraphs

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- The streets with apartments → **Target tree**

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- The path mapping to study groups → **Tree Path Labeling (TPL)**
- The apartment allocation → **Path Hypergraph Isomorphism**

Tree Path Labeling of Path Hypergraphs

There *exists* an apartment allocation that “fits” the path mapping

Tree Path Labeling of Path Hypergraphs

There *exists* a **hypergraph isomorphism** that “fits” the **TPL**

Tree Path Labeling of Path Hypergraphs

There *exists* a **hypergraph isomorphism** that “fits” the **TPL**
 \Rightarrow the TPL is **FEASIBLE**

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There *exists* an apartment allocation that gives some study
group path mapping

Tree Path Labeling of Path Hypergraphs

There *exists* a **hypergraph isomorphism** that “fits” the **TPL**
 \Rightarrow the TPL is **FEASIBLE**

There *exists* a **hypergraph isomorphism** that gives **at least one feasible TPL**

Tree Path Labeling of Path Hypergraphs

There *exists* a **hypergraph isomorphism** that “fits” the **TPL**

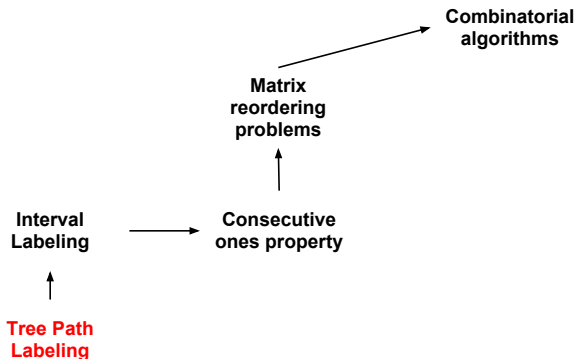
⇒ the TPL is **FEASIBLE**

There *exists* a **hypergraph isomorphism** that gives **at least one feasible TPL**

⇒ the hypergraph is a **PATH HYPERGRAPH**

Consecutive Ones \rightarrow Path Labeling

The motivation



2

Computing a feasible TPL

Given hypergraph \mathcal{F} with certain properties and a k -subdivided star T , can we find a feasible TPL ℓ to T ?

1

Characterization of feasible TPL

The characterization

ICPPL + a filtering algorithm

The characterization

ICPPL + a filtering algorithm

2

Computing a feasible TPL

Given hypergraph \mathcal{F} with certain properties and a k -subdivided star T , can we find a feasible TPL ℓ to T ?

Special case

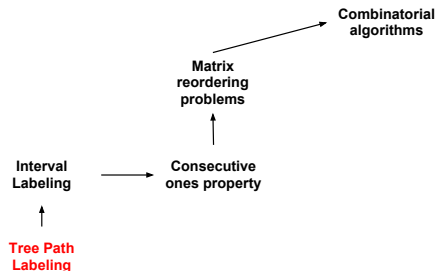
Interval assignment problem / COP

- ① T is a path \implies paths in T are intervals
- ② Only pairwise intersection cardinality needs to be preserved \implies ICPIA [NS09]
- ③ Higher level intersection cardinalities preserved by **Helly Property** – [Gol04]
- ④ $filter_1, filter_2$ do not need the **exit** conditions.

This problem is equivalent to Consecutive Ones Property of binary matrices [NS09]

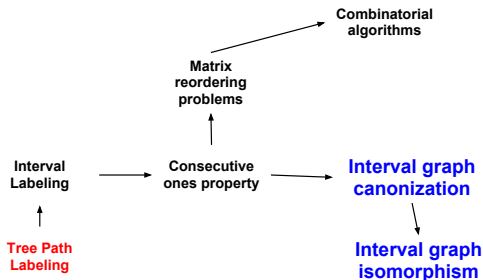
Path Labeling \rightarrow Graph Isomorphism

Application



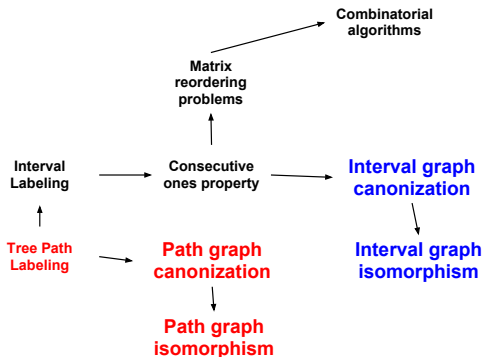
Path Labeling \rightarrow Graph Isomorphism

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Path Labeling → Graph Isomorphism

Application



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