

# Tree Path Labeling of Set Systems

## A Generalization of Consecutive Ones Property

Anju Srinivasan

as part of **M. S.** by Research  
advised by **Dr. N. S. Narayanaswamy**  
CSE, IITM, Chennai - 36

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## 1 Introduction

An Illustration

Motivation

Definitions

## 2 Characterization of a feasible TPL

ICPPL

Filtering algorithm

## 3 Computing a feasible TPL on $k$ -subdivided trees

Algorithm

## 4 Conclusion

Application

# An Illustration

## Caveat

- A very simplistic example.
- Aims only to introduce the combinatorial problem of TPL.

# A Study Group Housing problem

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- There are  $n$  single occupancy apartments in the university campus for their accommodation.

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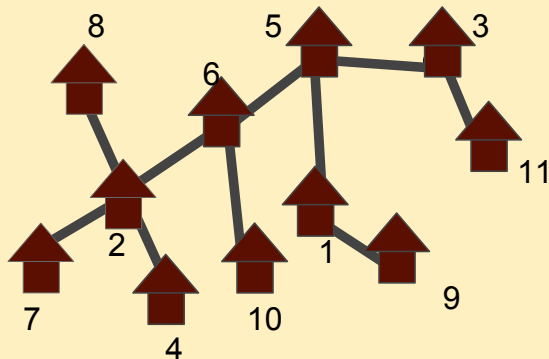
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- There are  $n$  single occupancy apartments in the university campus for their accommodation.
- All these apartments are placed such that streets connecting them do not form loops

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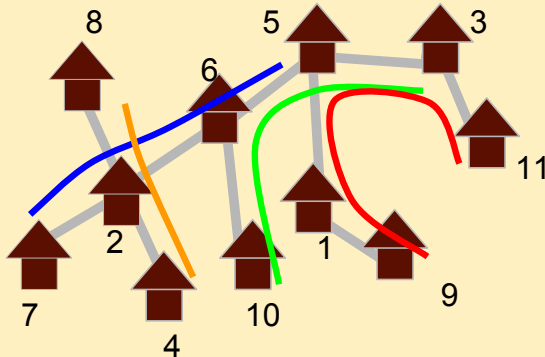
## The problem

How should the students be allocated apartments such that each study group has the least distance to travel for a discussion?

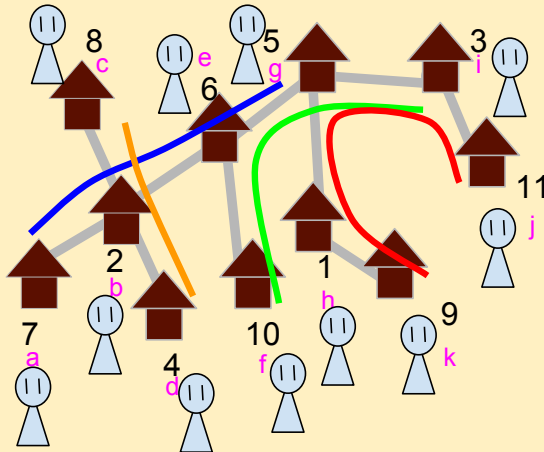
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# Tree Path Labeling of Set Systems

The combinatorial problem terminology

## Terminology

- The set of study groups i.e. sets of students  $\rightarrow$  SET SYSTEM / HYPERGRAPH
- The streets with apartments  $\rightarrow$  TARGET TREE
- The route mapping to study groups  $\rightarrow$  TREE PATH LABELING (TPL)
- The apartment allocation  $\rightarrow$  PATH HYPERGRAPH ISOMORPHISM

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The combinatorial problem

## Terminology [contd.]

There *exists* an apartment allocation that “fits” the route mapping

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There *exists* a hypergraph isomorphism that “fits” the TPL

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There *exists* a hypergraph isomorphism that “fits” the TPL  
→ the TPL is FEASIBLE



# Tree Path Labeling of Set Systems

The combinatorial problem

## Terminology [contd.]

There *exists* a hypergraph isomorphism that “fits” the TPL  
→ the TPL is FEASIBLE

There *exists* an apartment allocation that gives the optimal route mapping

# Tree Path Labeling of Set Systems

The combinatorial problem

## Terminology [contd.]

There *exists* a hypergraph isomorphism that “fits” the TPL  
→ the TPL is FEASIBLE

There *exists* a hypergraph isomorphism that gives  
paths/adjacent vertices in tree

# Tree Path Labeling of Set Systems

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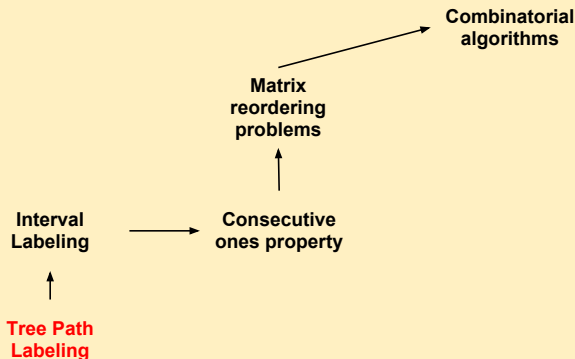
## Terminology [contd.]

There *exists* a hypergraph isomorphism that “fits” the TPL  
→ the TPL is FEASIBLE

There *exists* a hypergraph isomorphism that gives  
paths/adjacent vertices in tree  
→ the hypergraph is a PATH HYPERGRAPH

# Consecutive Ones $\rightarrow$ Path Labeling

The motivation



# Tree path labeling of path hypergraphs

The two problems

1

Characterization of a *feasible TPL* and finding the certificate for feasibility i.e. *hypergraph isomorphism*

2

Computation of a *feasible TPL* if any

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The two problems

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Characterization of a *feasible TPL* and finding the certificate for feasibility i.e. *hypergraph isomorphism*

2

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## 1

## Characterization of feasible TPL

Given

- i. a set system or hypergraph  $\mathcal{F}$ ,
- ii. a feasible TPL  $\ell : \mathcal{F} \rightarrow \mathcal{P}$  where  $\mathcal{P}$  is a path system from tree  $T$  and  $\text{supp}(\mathcal{P}) = V(T)$ ,

what is the hypergraph isomorphism

$$\underline{\phi : \text{supp}(\mathcal{F}) \rightarrow \text{supp}(\mathcal{P})}$$

such that the induced labeling  $\ell_\phi = \ell$ ?

# 2

## Computing a feasible TPL

Given hypergraph  $\mathcal{F}$  with certain properties and a  $k$ -subdivided star  $T$ , can we find a feasible TPL  $\ell$  to  $T$ ?



# 1

## Characterization of feasible TPL

# The characterization

ICPPL + a filtering algorithm

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## 2

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# Special case

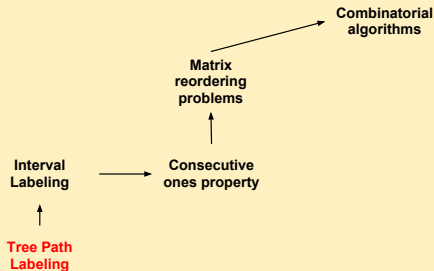
Interval assignment problem / COP

- 1  $T$  is a path  $\implies$  paths in  $T$  are intervals
- 2 Only pairwise intersection cardinality needs to be preserved  $\implies$  ICPIA [NS09]
- 3 Higher level intersection cardinalities preserved by **Helly Property** – [Gol04]
- 4  $filter\_1, filter\_2$  do not need the the **exit** conditions.

This problem is equivalent to Consecutive Ones Property of binary matrices [NS09]

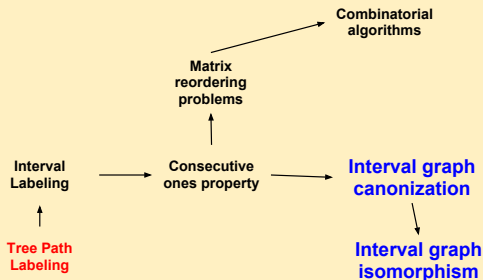
# Path Labeling $\rightarrow$ Graph Isomorphism

## Application



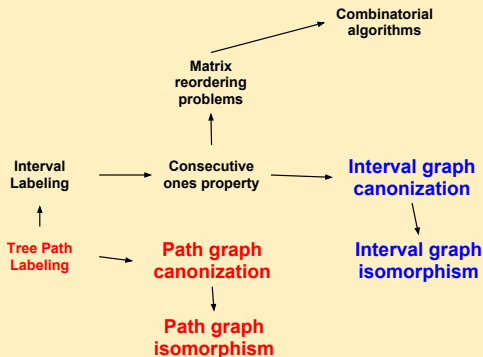
# Path Labeling $\rightarrow$ Graph Isomorphism

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# Path Labeling → Graph Isomorphism

## Application





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