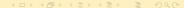
A Generalization of Consecutive Ones Property

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as part of M. S. by Research advised by Dr. N. S. Narayanaswamy CSE, IITM, Chennai - 36

XX Dec 2011



- Introduction
 - An Illustration Motivation Definitions
- Characterization of a feasible TPL ICPPL Filtering algorithm
- Computing a feasible TPL on k-subdivided trees Algorithm
- 4 Conclusion
 Application



An Illustration

[say something less apologetic!]



[update to the example in synopsis doc]

• A set of n students arrive for a summer course, say $\{a, b, c, d, e, f, g, h, i, j, k\}$, n = 11.



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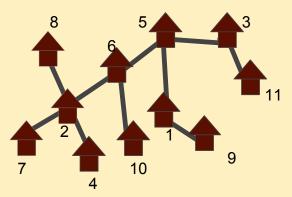
- $G = \{e, f, g, i\}$
- There are *n* single occupancy apartments in the university campus for their accommodation.
- All these apartments are placed such that streets



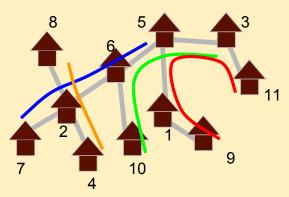
The problem

How should the students be allocated apartments such that each study group has the least distance to travel for a discussion?

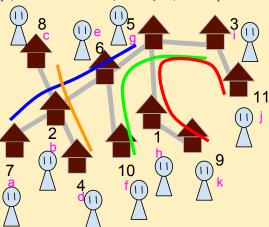


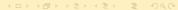












The combinatorial problem terminology

- \bullet The set of study groups i.e. sets of students $\to S{\rm ET}$ SYSTEM / HYPERGRAPH
- ullet The streets with apartments o TARGET TREE
- The route mapping to study groups \rightarrow TREE PATH LABELING (TPL)
- The apartment allocation → PATH HYPERGRAPH
 ISOMORPHISM



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The combinatorial problem

Terminology [contd.]

There *exists* an apartment allocation that "fits" the route mapping



The combinatorial problem

Terminology [contd.]

There exists a hypergraph isomorphism that "fits" the TPL



The combinatorial problem

Terminology [contd.]

There *exists* a hypergraph isomorphism that "fits" the TPL → the TPL is FEASIBLE



The combinatorial problem

Terminology [contd.]

There *exists* a hypergraph isomorphism that "fits" the TPL → the TPL is FEASIBLE

There *exists* an apartment allocation that gives the optimal route mapping



The combinatorial problem

Terminology [contd.]

There *exists* a hypergraph isomorphism that "fits" the TPL → the TPL is FEASIBLE

There *exists* a hypergraph isomorphism that gives paths/adjacent vertices in tree



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Terminology [contd.]

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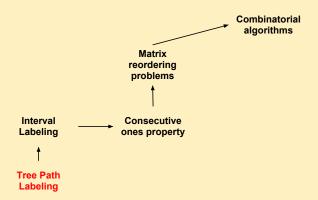
There *exists* a hypergraph isomorphism that gives paths/adjacent vertices in tree

 \rightarrow the hypergraph is a PATH HYPERGRAPH



Consecutive Ones → **Path Labeling**

The motivation





Tree path labeling of path hypergraphs

The two problems

1

Characterization of a *feasible TPL* and finding the certificate for feasibility i.e. *hypergraph isomorphism*

2

Computation of a feasible TPL if any



Tree path labeling of path hypergraphs

The two problems

1

Characterization of a *feasible TPL* and finding the certificate for feasibility i.e. *hypergraph isomorphism*

2

Computation of a *feasible TPL* if any



[do we need this frame at all?]

1

Characterization of feasible TPL

Given

- i. a set system or hypergraph ${\mathcal F}$,
- ii. a feasible TPL $\ell: \mathcal{F} \to \mathcal{P}$ where \mathcal{P} is a path system from tree T and $supp(\mathcal{P}) = V(T)$,

what is the hypergraph isomorphism

$$\phi: \mathsf{supp}\,(\mathcal{F}) \to \mathsf{supp}\,(\mathcal{P})$$

such that the induced labeling $l_{\phi} = l$?

_ _ _ _ _ _ _

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[do we need this frame at all?]

2

Computing a feasible TPL

Given hypergraph \mathcal{F} with certain properties and a k-subdivided star T, can we find a feasible TPL ℓ to T?



1

Characterization of feasible TPL



The characterization

ICPPL + a filtering algorithm

^{a:} [TBD Write the theorem]



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Special case

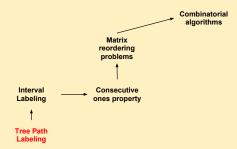
Interval assignment problem / COP

- **1** T is a path \Longrightarrow paths in T are intervals a^{a} [quick illustration]
- ② Only pairwise intersection cardinality needs to be preserved ⇒ ICPIA [NS09]
- Higher level intersection cardinalities preserved by Helly Property – [Gol04]
- filter_1, filter_2 do not need the the exit conditions. a: [is this cryptic?]

This problem is equivalent to Consecutive Ones Property of binary matrices [NS09]

Path Labeling → Graph Isomorphism

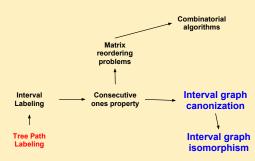
Application





Path Labeling → Graph Isomorphism

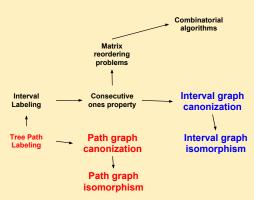
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Path Labeling → Graph Isomorphism

Application





Thank You

Q & A



[improve - add some jazz. this is a notional slide only for offline reference.]

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