

Tree Path Labeling of Set Systems

A Generalization of Consecutive Ones Property

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1 Introduction

An Illustration

Motivation

Definitions

2 Characterization of a feasible TPL

ICPPL

Filtering algorithm

3 Computing a feasible TPL on k -subdivided trees

Algorithm

4 Conclusion

Application

An Illustration

Caveat

- A very simplistic example.
- Aims only to introduce the combinatorial problem of TPL.

A Study Group Housing problem

- A set of n students arrive for a summer course, say $\{a, b, c, d, e, f, g, h, i, j, k\}$, $n = 11$.

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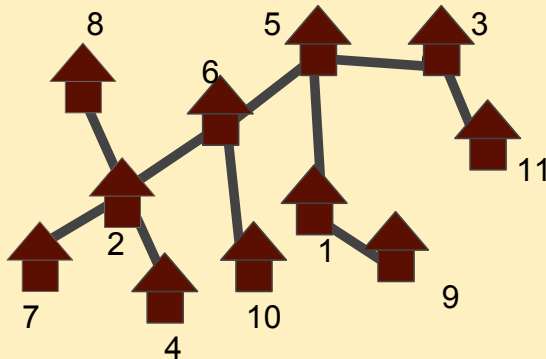
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- There are n single occupancy apartments in the university campus for their accommodation.
- All these apartments are placed such that streets connecting them do not form loops

A Study Group Housing problem

The problem

How should the students be allocated apartments such that each study group has the least distance to travel for a discussion?

A Study Group Housing problem



Tree Path Labeling of Set Systems

The combinatorial problem terminology

Terminology

- The set of study groups i.e. sets of students → SET SYSTEM / HYPERGRAPH
- The streets with apartments → TARGET TREE
- The route mapping to study groups → TREE PATH LABELING (TPL)
- The apartment allocation → PATH HYPERGRAPH ISOMORPHISM

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The combinatorial problem

Terminology [contd.]

There *exists* an apartment allocation that “fits” the route mapping

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Terminology [contd.]

There *exists* a hypergraph isomorphism that “fits” the TPL

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Terminology [contd.]

There *exists* a hypergraph isomorphism that “fits” the TPL
→ the TPL is FEASIBLE

Tree Path Labeling of Set Systems

The combinatorial problem

Terminology [contd.]

There *exists* a hypergraph isomorphism that “fits” the TPL
→ the TPL is **FEASIBLE**

There *exists* an apartment allocation that gives the optimal route mapping

Tree Path Labeling of Set Systems

The combinatorial problem

Terminology [contd.]

There *exists* a hypergraph isomorphism that “fits” the TPL
→ the TPL is FEASIBLE

There *exists* a hypergraph isomorphism that gives
paths/adjacent vertices in tree

Tree Path Labeling of Set Systems

The combinatorial problem

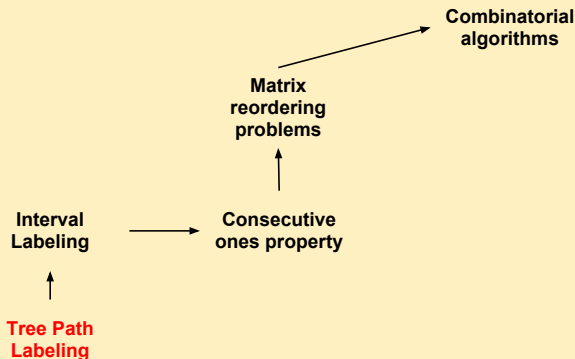
Terminology [contd.]

There *exists* a hypergraph isomorphism that “fits” the TPL
→ the TPL is FEASIBLE

There *exists* a hypergraph isomorphism that gives
paths/adjacent vertices in tree
→ the hypergraph is a PATH HYPERGRAPH

Consecutive Ones \rightarrow Path Labeling

The motivation



Tree path labeling of path hypergraphs

The two problems

1

Characterization of a *feasible TPL* and finding the certificate for feasibility i.e. *hypergraph isomorphism*

2

Computation of a *feasible TPL* if any

Tree path labeling of path hypergraphs

The two problems

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Characterization of a *feasible TPL* and finding the certificate for feasibility i.e. *hypergraph isomorphism*

2

Computation of a *feasible TPL* if any

1

Characterization of feasible TPL

Given

- i. a set system or hypergraph \mathcal{F} ,
- ii. a feasible TPL $\ell : \mathcal{F} \rightarrow \mathcal{P}$ where \mathcal{P} is a path system from tree T and $\text{supp}(\mathcal{P}) = V(T)$,

what is the hypergraph isomorphism

$$\underline{\phi : \text{supp}(\mathcal{F}) \rightarrow \text{supp}(\mathcal{P})}$$

such that the induced labeling $\ell_\phi = \ell$?

2

Computing a feasible TPL

Given hypergraph \mathcal{F} with certain properties and a k -subdivided star T , can we find a feasible TPL ℓ to T ?

1

Characterization of feasible TPL

The characterization

ICPPL + a filtering algorithm

^a: [TBD Write the theorem]

The characterization

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2

Computing a feasible TPL

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Special case

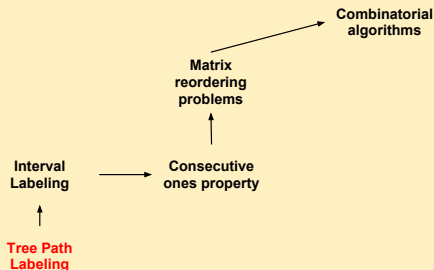
Interval assignment problem / COP

- 1 T is a path \implies paths in T are intervals ^{a:} [quick illustration]
- 2 Only pairwise intersection cardinality needs to be preserved \implies ICPIA [NS09]
- 3 Higher level intersection cardinalities preserved by **Helly Property** – [Gol04]
- 4 $filter_1, filter_2$ do not need the the **exit** conditions. ^{a:} [is this cryptic?]

This problem is equivalent to Consecutive Ones Property of binary matrices [NS09]

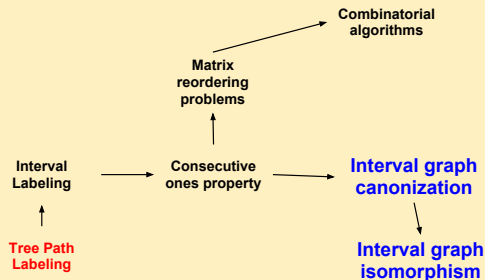
Path Labeling → Graph Isomorphism

Application



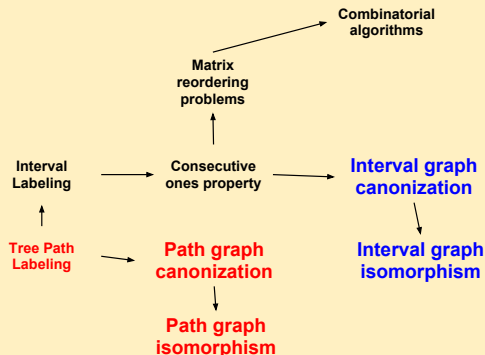
Path Labeling → Graph Isomorphism

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Path Labeling → Graph Isomorphism

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