

FTT 1047 - Assignment 1

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Tutorial: Tuesday 4-6pm

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Task 1.1

Given Truth Table

Table 1 shows the truth table that was provided in the Assignment1 question.

	X1	X2	X3	X4	Z1	Z2
1	0	0	0	0	1	1
2	0	0	0	1	0	1
3	0	0	1	0	1	0
4	0	0	1	1	1	0
5	0	1	0	0	0	1
6	0	1	0	1	1	0
7	0	1	1	0	1	0
8	0	1	1	1	1	0
9	1	0	0	0	0	1
10	1	0	0	1	1	1
11	1	0	1	0	1	0
12	1	0	1	1	1	0
13	1	1	0	0	0	0
14	1	1	0	1	0	1
15	1	1	1	0	1	1
16	1	1	1	1	0	1

Table 1

Based on the truth table provided, it is split into 2 separate tables, with output Z1(Table 2) and Z2(Table 3) separately for easy understanding and explanation.

Output for Z1

	X1	X2	X3	X4	Z1
1	0	0	0	0	1
2	0	0	0	1	0
3	0	0	1	0	1
4	0	0	1	1	1
5	0	1	0	0	0
6	0	1	0	1	1
7	0	1	1	0	1
8	0	1	1	1	1
9	1	0	0	0	0
10	1	0	0	1	1
11	1	0	1	0	1
12	1	0	1	1	1
13	1	1	0	0	0
14	1	1	0	1	0
15	1	1	1	0	1
16	1	1	1	1	0

Table 2



1	$Z1 = \overline{X1} \cdot \overline{X2} \cdot \overline{X3} \cdot \overline{X4}$
2	
3	$Z1 = \overline{X1} \cdot \overline{X2} \cdot X3 \cdot \overline{X4}$
4	$Z1 = \overline{X1} \cdot \overline{X2} \cdot X3 \cdot X4$
5	
6	$Z1 = \overline{X1} \cdot X2 \cdot \overline{X3} \cdot X4$
7	$Z1 = \overline{X1} \cdot X2 \cdot X3 \cdot \overline{X4}$
8	$Z1 = \overline{X1} \cdot X2 \cdot X3 \cdot X4$
9	
10	$Z1 = X1 \cdot \overline{X2} \cdot \overline{X3} \cdot X4$
11	$Z1 = X1 \cdot \overline{X2} \cdot X3 \cdot \overline{X4}$
12	$Z1 = X1 \cdot \overline{X2} \cdot X3 \cdot X4$
13	
14	
15	$Z1 = X1 \cdot X2 \cdot X3 \cdot \overline{X4}$
16	

Output for Z2

	X1	X2	X3	X4	Z2
1	0	0	0	0	1
2	0	0	0	1	1
3	0	0	1	0	0
4	0	0	1	1	0
5	0	1	0	0	1
6	0	1	0	1	0
7	0	1	1	0	0
8	0	1	1	1	0
9	1	0	0	0	1
10	1	0	0	1	1
11	1	0	1	0	0
12	1	0	1	1	0
13	1	1	0	0	0
14	1	1	0	1	1
15	1	1	1	0	1
16	1	1	1	1	1

Table 3



1	$Z2 = \overline{X1} \cdot \overline{X2} \cdot \overline{X3} \cdot \overline{X4}$
2	$Z2 = \overline{X1} \cdot \overline{X2} \cdot \overline{X3} \cdot X4$
3	
4	
5	$Z2 = \overline{X1} \cdot X2 \cdot \overline{X3} \cdot \overline{X4}$
6	
7	
8	
9	$Z2 = X1 \cdot \overline{X2} \cdot \overline{X3} \cdot \overline{X4}$
10	$Z2 = X1 \cdot \overline{X2} \cdot \overline{X3} \cdot X4$
11	
12	
13	
14	$Z2 = X1 \cdot X2 \cdot \overline{X3} \cdot \overline{X4}$
15	$Z2 = X1 \cdot X2 \cdot \overline{X3} \cdot X4$
16	$Z2 = X1 \cdot X2 \cdot X3 \cdot X4$

Boolean Algebra Equation

First, only **outputs with Boolean 1** are highlighted (in yellow/green). Then, based on the inputs X1, X2, X3 and X4, the equation is created. An input of Boolean '0' for X1 will result in ' $\overline{X1}$ ', and an input of Boolean '1' will result in ' $X1$ '.

For example:

	X1	X2	X3	X4	Z1
6	0	1	0	1	1

↑
 $\overline{X1}$

↑
 $X2$

↑
 $\overline{X3}$

↑
 $X4$

↑
 $Z1$

The complete equation will then be $Z1 = \overline{X1} \cdot X2 \cdot \overline{X3} \cdot X4$.

Final Equation

For Table 2, there are 10 equations that result in Z1. To combine all the equations, the equations are simply summed up.

$$Z1 = (\overline{X1} \cdot \overline{X2} \cdot \overline{X3} \cdot \overline{X4}) + (\overline{X1} \cdot \overline{X2} \cdot X3 \cdot \overline{X4}) + (\overline{X1} \cdot \overline{X2} \cdot X3 \cdot X4) + (\overline{X1} \cdot X2 \cdot \overline{X3} \cdot X4) + (\overline{X1} \cdot X2 \cdot X3 \cdot \overline{X4}) + (\overline{X1} \cdot X2 \cdot X3 \cdot X4) + (X1 \cdot \overline{X2} \cdot \overline{X3} \cdot X4) + (X1 \cdot \overline{X2} \cdot X3 \cdot \overline{X4}) + (X1 \cdot \overline{X2} \cdot X3 \cdot X4) + (X1 \cdot X2 \cdot \overline{X3} \cdot \overline{X4})$$

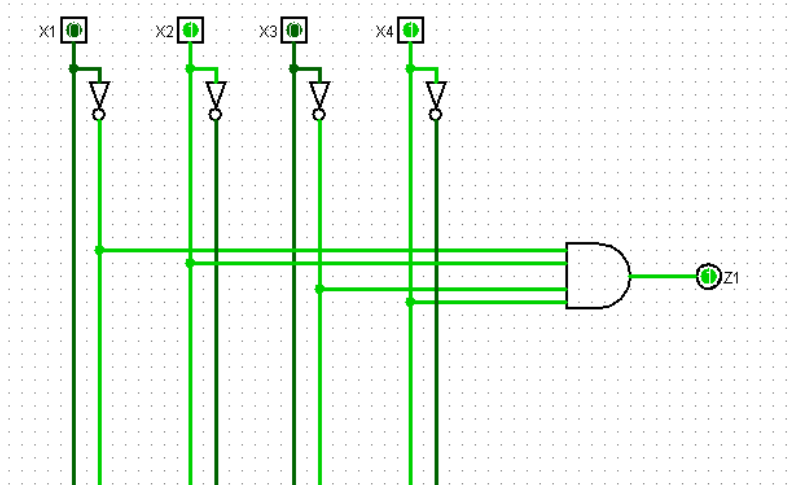
As for Z2, the same concept is applied to obtain the equation below:

$$Z2 = (\overline{X1} \cdot \overline{X2} \cdot \overline{X3} \cdot \overline{X4}) + (\overline{X1} \cdot \overline{X2} \cdot \overline{X3} \cdot X4) + (\overline{X1} \cdot X2 \cdot \overline{X3} \cdot \overline{X4}) + (X1 \cdot \overline{X2} \cdot \overline{X3} \cdot \overline{X4}) + (X1 \cdot \overline{X2} \cdot \overline{X3} \cdot X4) + (X1 \cdot X2 \cdot \overline{X3} \cdot \overline{X4}) + (X1 \cdot X2 \cdot \overline{X3} \cdot X4) + (X1 \cdot X2 \cdot X3 \cdot \overline{X4}) + (X1 \cdot X2 \cdot X3 \cdot X4)$$

Task 1.2

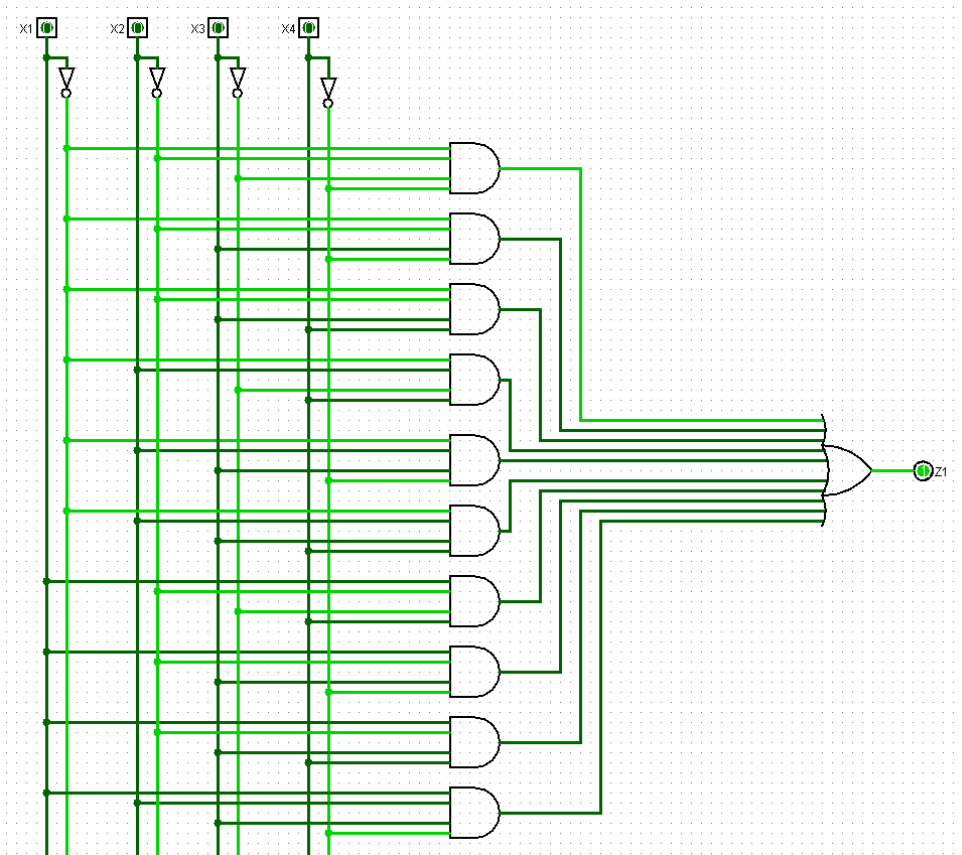
Logisim circuits are created based on the equations. A ‘.’ sign in the equation signifies an AND gate, and a ‘+’ sign signifies an OR gate. A NOT gate is used to express any input that is \bar{X} .

For example, $Z1 = \bar{X1} \cdot X2 \cdot \bar{X3} \cdot X4$ will result in Circuit1:

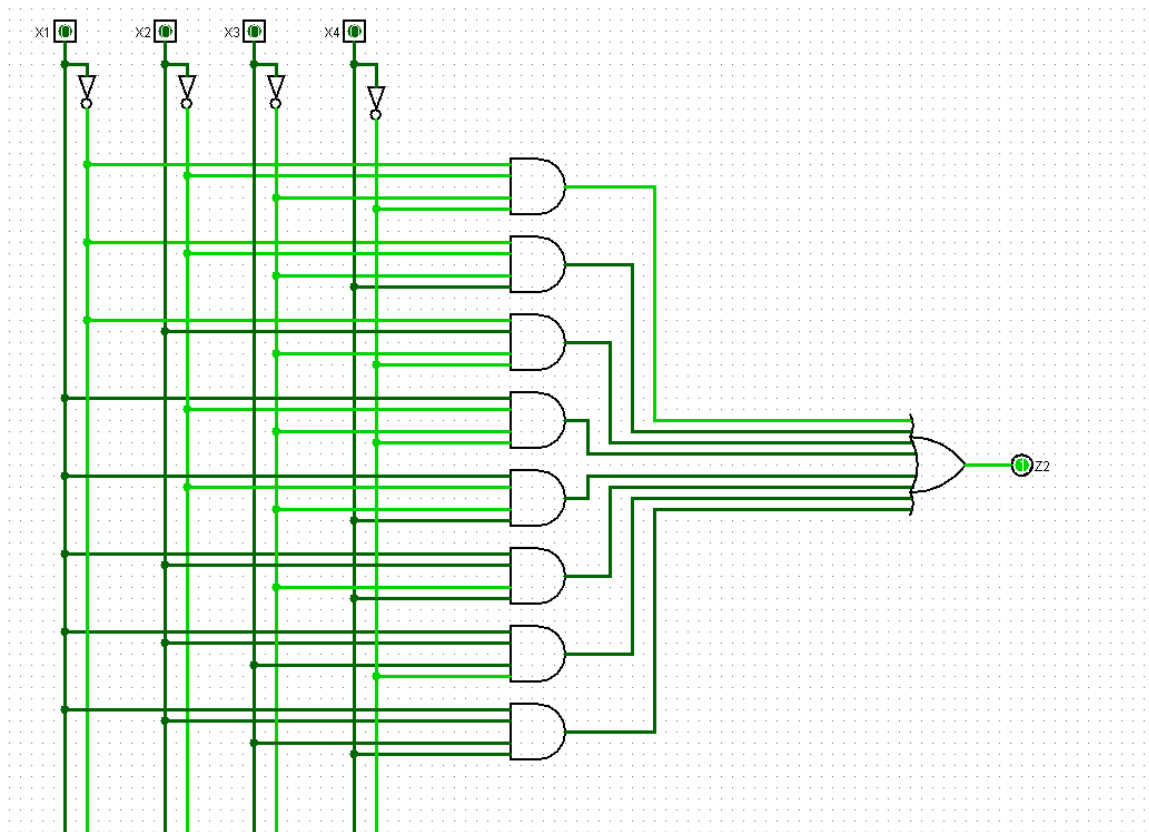


Circuit 1

For Z1, as there are ten equations, ten AND gates are used. After that, an OR gate is used to combine all the AND gates to give an accurate output of Z1. The same concept is applied for Z2. Below are the circuits for Z1(Circuit 2) and Z2(Circuit 3) respectively:

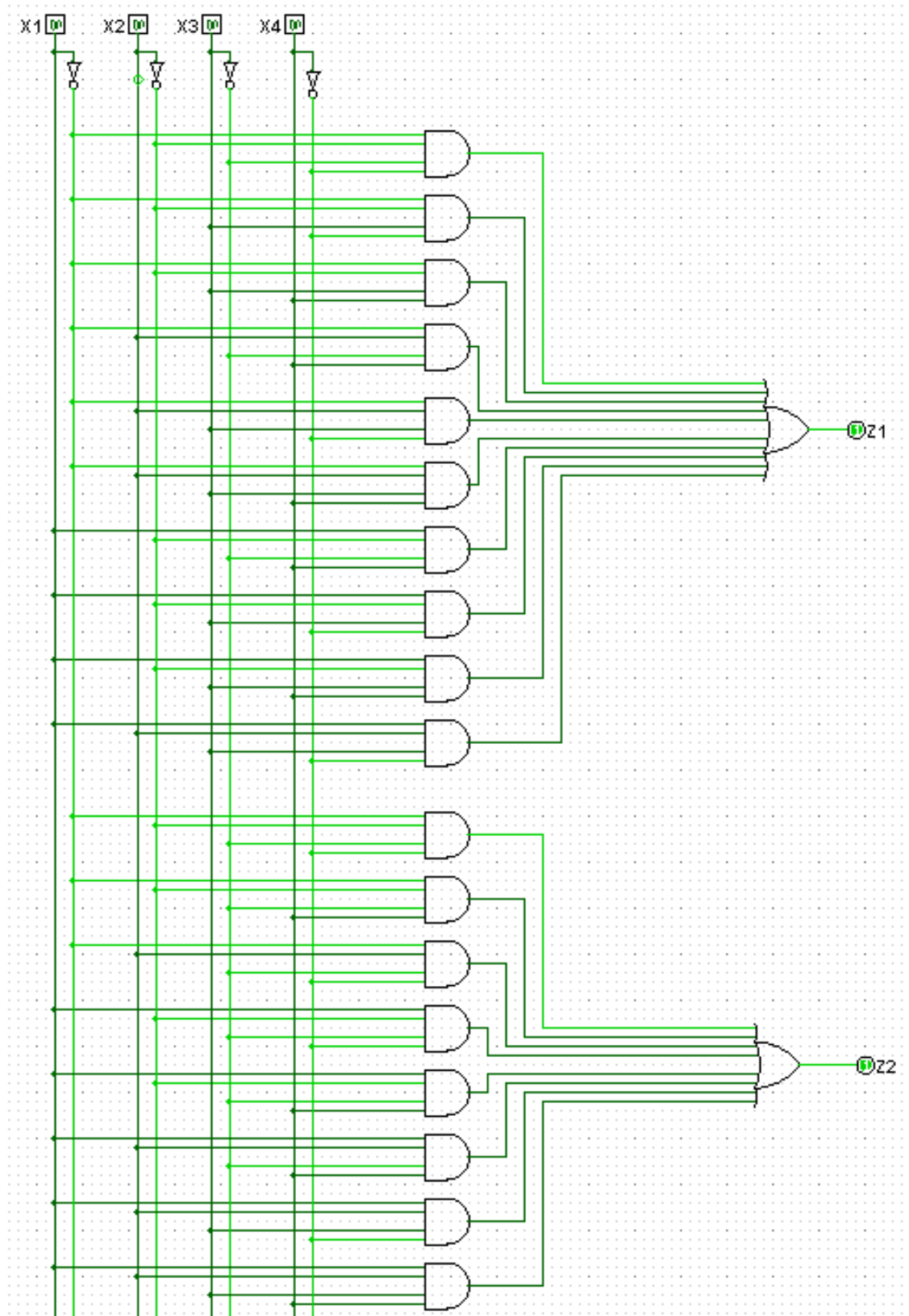


Circuit 2



Circuit 3

As Z1 AND Z2 have same input variables X1, X2, X3 and X4, the circuits can be combined to form the integrated circuit below containing outputs for both Z1 and Z2(Circuit 4):



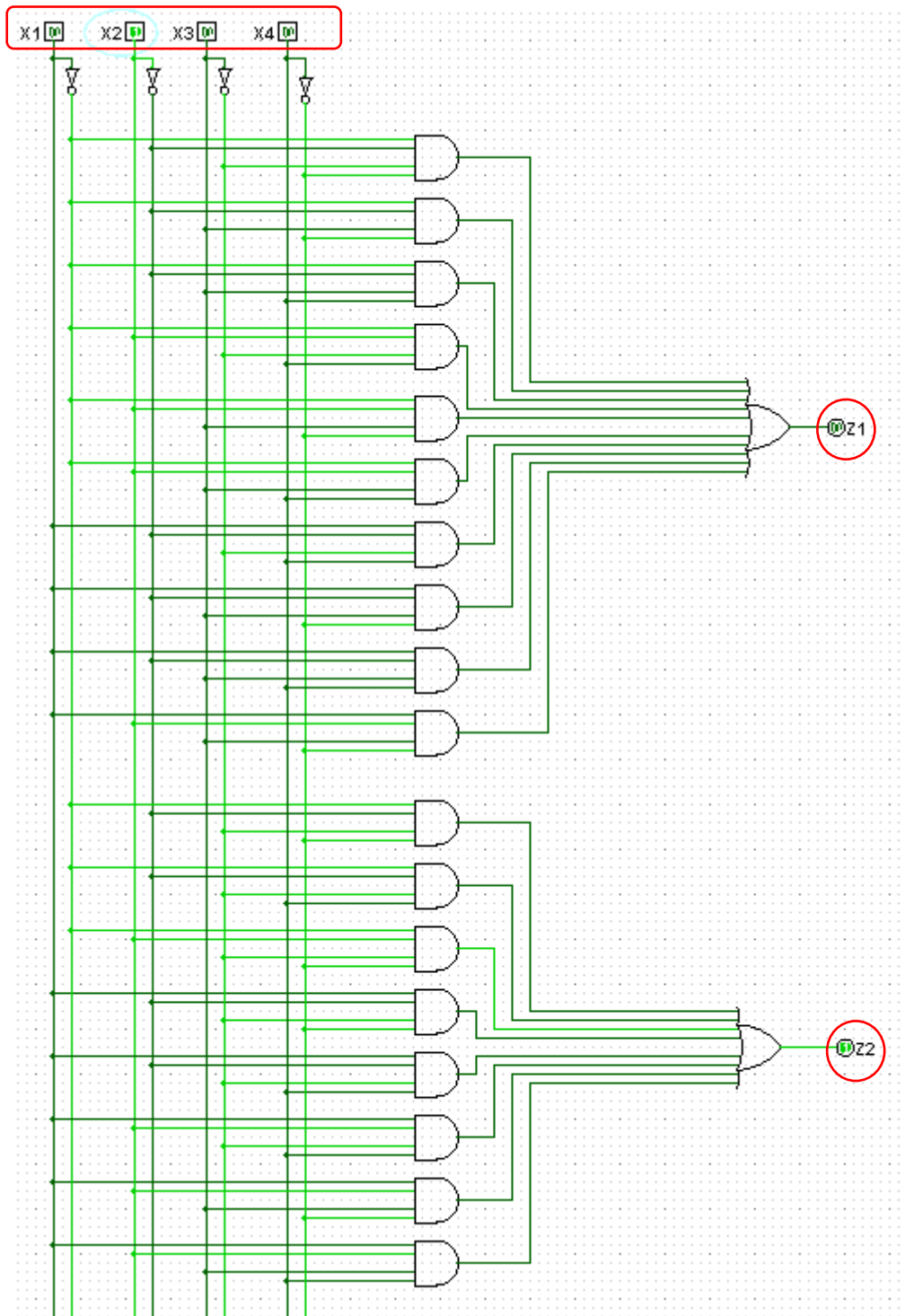
Circuit 4

Test Cases

3 test cases are documented to show that the integrated circuit gives the desired outputs. It is done by randomly selecting a given input and ensuring that it gives the respective outputs for Z1 and Z2.

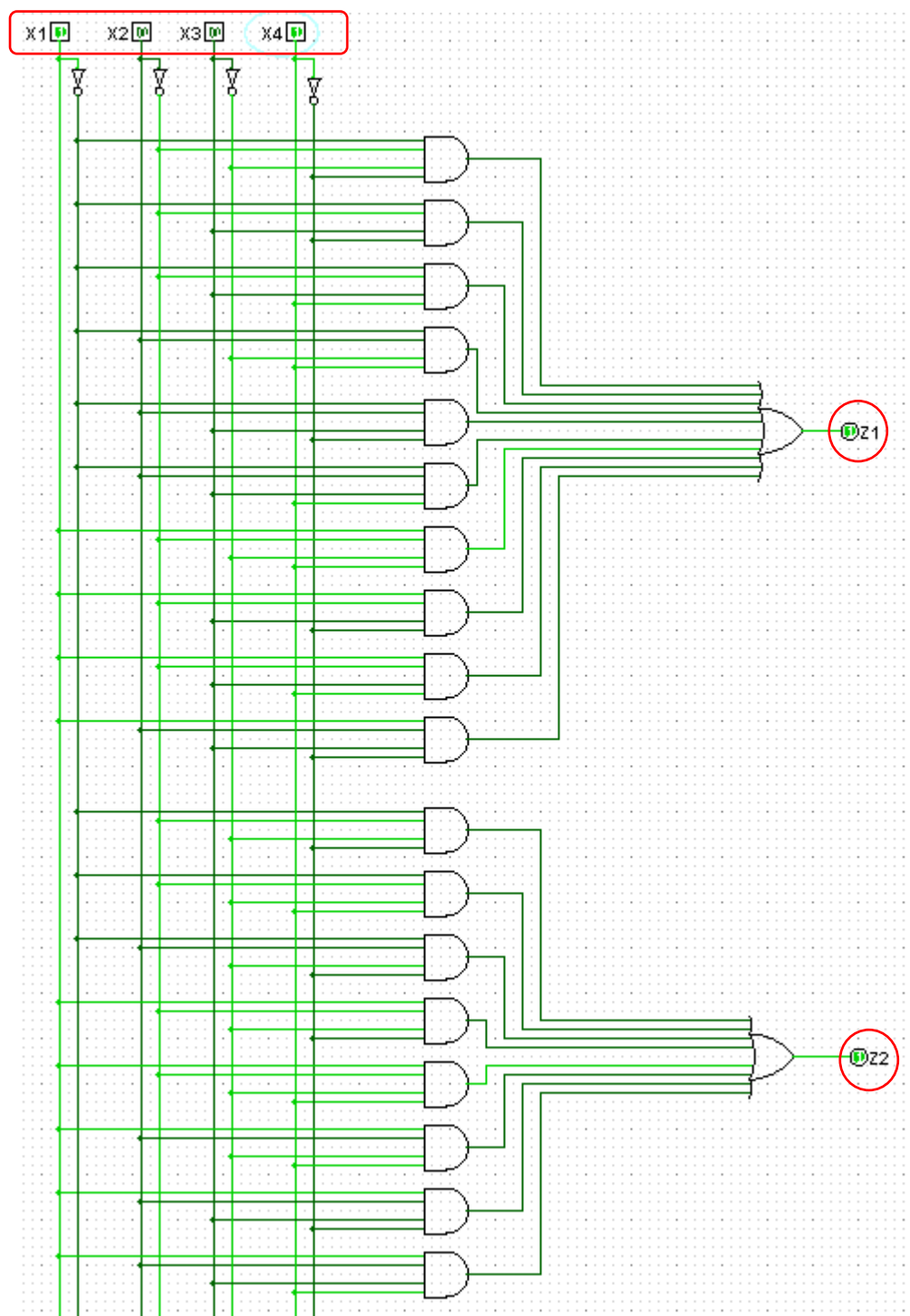
Test case 1:

	X1	X2	X3	X4	Z1	Z2
5	0	1	0	0	0	1



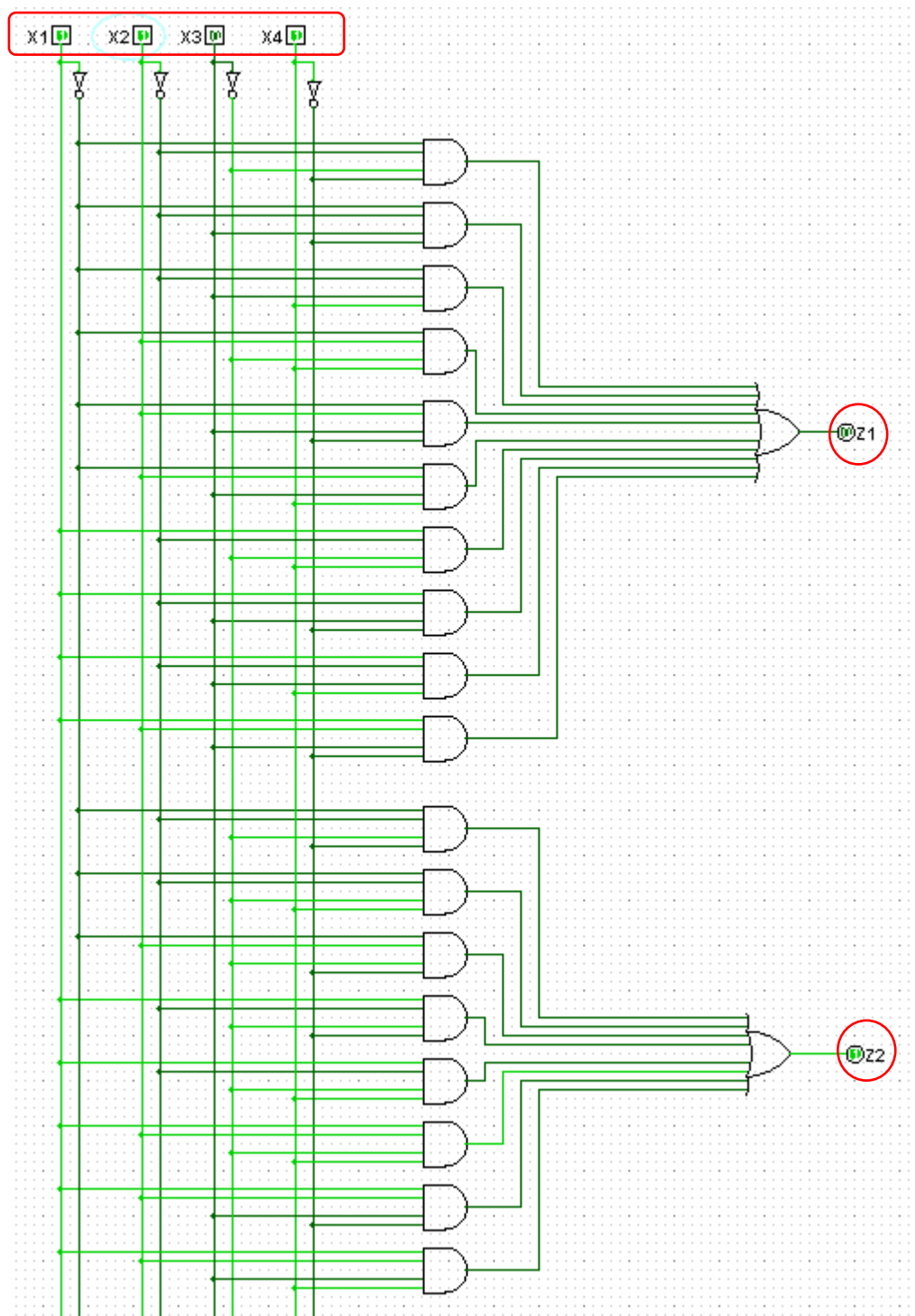
Test case 2:

	X1	X2	X3	X4	Z1	Z2
10	1	0	0	1	1	1



Test case 3:

	X1	X2	X3	X4	Z1	Z2
14	1	1	0	1	0	1



Task 1.3

The final equation for Z1 and Z2 is quite large, and therefore can be simplified using Karnaugh maps(K-maps). For Z1, a table is created with the axis containing inputs for X1X2 on one side and inputs for X3X4 on the other (K-map 1). After that, the table is filled with the outputs of Z1. The same is done for Z2 with K-map 2.

Some rules for K-maps:

1. **Only Ones** are grouped together
2. **All Ones** must be grouped
3. Groups cannot be diagonal.
4. Groups must contain 2^n of cells, where $n = 1, 2, 3..$
5. Groups must be as large as possible
6. Groups may overlap
7. Groups can wrap around the table

After the grouping for K-maps is complete, similar inputs from the same groups are analysed to build equations. Table 4.1 and 4.2 contains the equations for K-map 1 and K-map 2 respectively.

Eg:

We want to determine the equation of the **Red** group from **K-map 1**. It can be seen that for both cells, X1X2 are constant, as well as X4. However, X3 varies in the 2 cells, so it can be ignored. Hence, the equation for the red group is

$$Z1 = X1 \cdot \overline{X2} \cdot X4$$

		X3X4			
		00	01	11	10
X1X2	Z1	1	0	1	1
	00	1	0	1	1
	01	0	1	1	1
	11	0	0	0	1
	10	0	1	1	1

K-map 1

No.	Colour	Equation
1	Blue	$Z1 = \overline{X1} \cdot \overline{X2} \cdot \overline{X4}$
2	Yellow	$Z1 = X3 \cdot \overline{X4}$
3	Purple	$Z1 = \overline{X1} \cdot X2 \cdot X4$
4	Red	$Z1 = X1 \cdot \overline{X2} \cdot X4$
5	Green	$Z1 = \overline{X1} \cdot X3$

Table 4.1

		X3X4			
		00	01	11	10
X1X2	Z2	1	1	0	0
	00	1	1	0	0
	01	1	0	0	0
	11	0	1	1	1
	10	1	1	0	0

K-map 2

No.	Colour	Equation
1	Yellow	$Z2 = \overline{X1} \cdot \overline{X3} \cdot \overline{X4}$
2	Purple	$Z2 = \overline{X1} \cdot \overline{X2} \cdot \overline{X3}$
3	Red	$Z2 = X1 \cdot \overline{X2} \cdot \overline{X3}$
4	Green	$Z2 = X1 \cdot \overline{X3} \cdot X4$
5	Blue	$Z2 = X1 \cdot X2 \cdot X3$

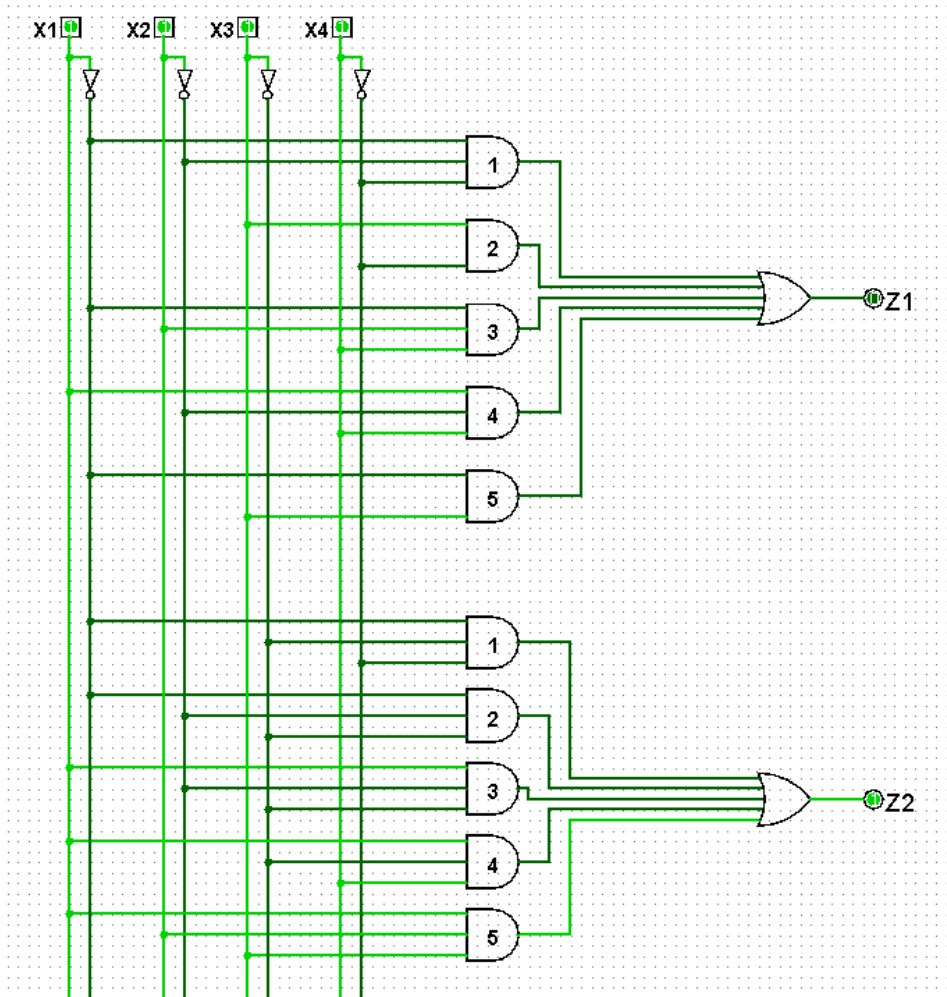
Table 5.1

Similar to previous explanation for combining equations, the final simplified equation for Z1 and Z2 is:

$$Z1 = (\overline{X1} \cdot \overline{X2} \cdot \overline{X4}) + (X3 \cdot \overline{X4}) + (\overline{X1} \cdot X2 \cdot X4) + (X1 \cdot \overline{X2} \cdot X4) + (\overline{X1} \cdot X3)$$

$$Z2 = (\overline{X1} \cdot \overline{X3} \cdot \overline{X4}) + (\overline{X1} \cdot \overline{X2} \cdot \overline{X3}) + (X1 \cdot \overline{X2} \cdot \overline{X3}) + (X1 \cdot \overline{X3} \cdot X4) + (X1 \cdot X2 \cdot X3)$$

A simplified circuit is created in Logisim using AND gates for ' \cdot ', OR gates for '+' and NOT gate for \overline{X} (refer to Circuit 5).

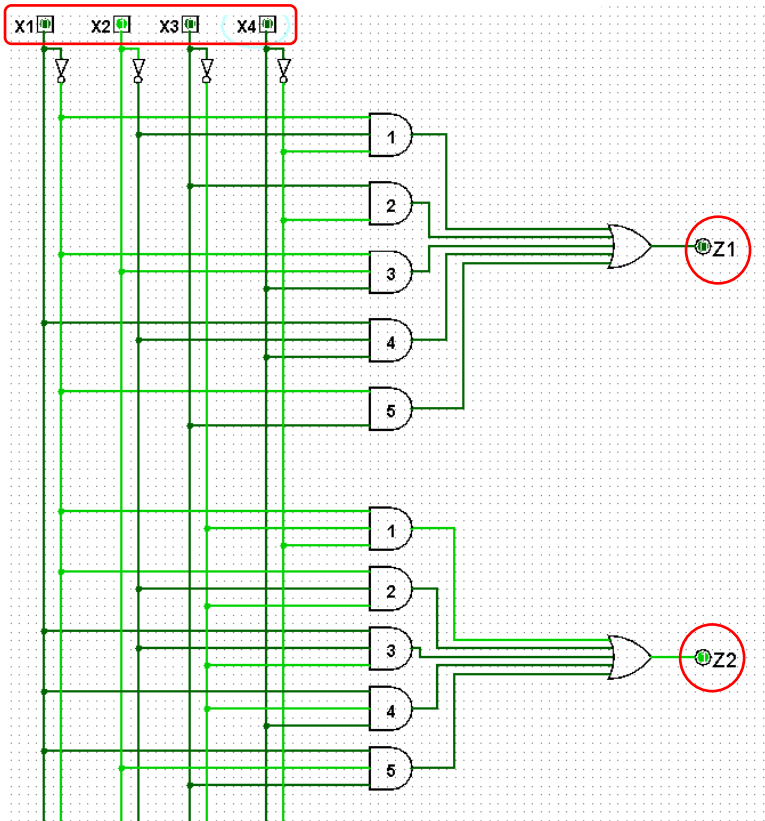


Circuit 5

The **same** test cases from Task 1.2 are again used to determine if the simplified circuits produce the correct output. The outputs for the simplified final circuit should be the same as the ones in Task 1.2.

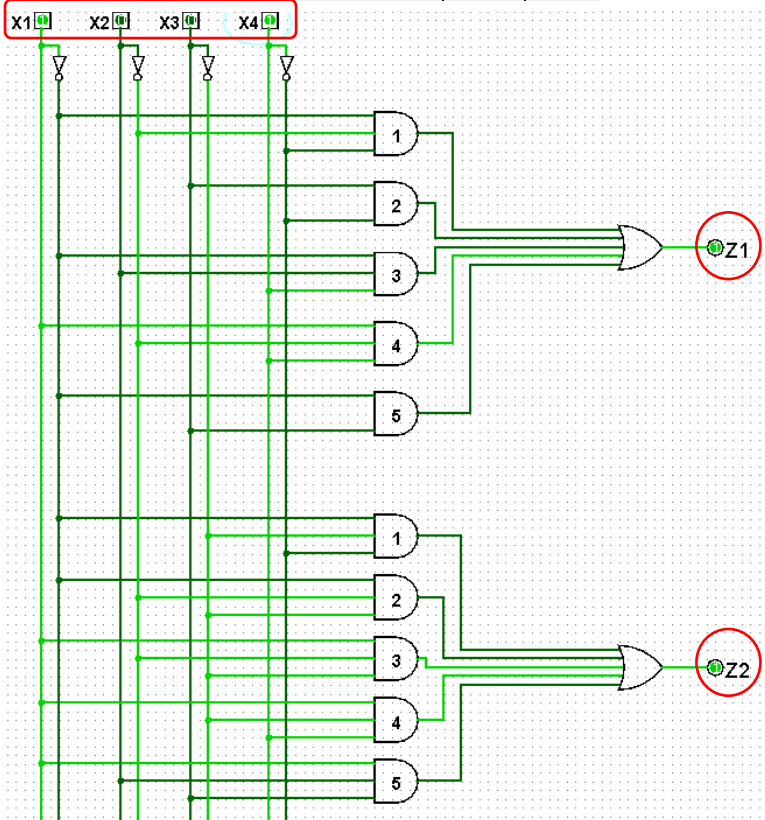
Test case 1:

	X1	X2	X3	X4	Z1	Z2
5	0	1	0	0	0	1



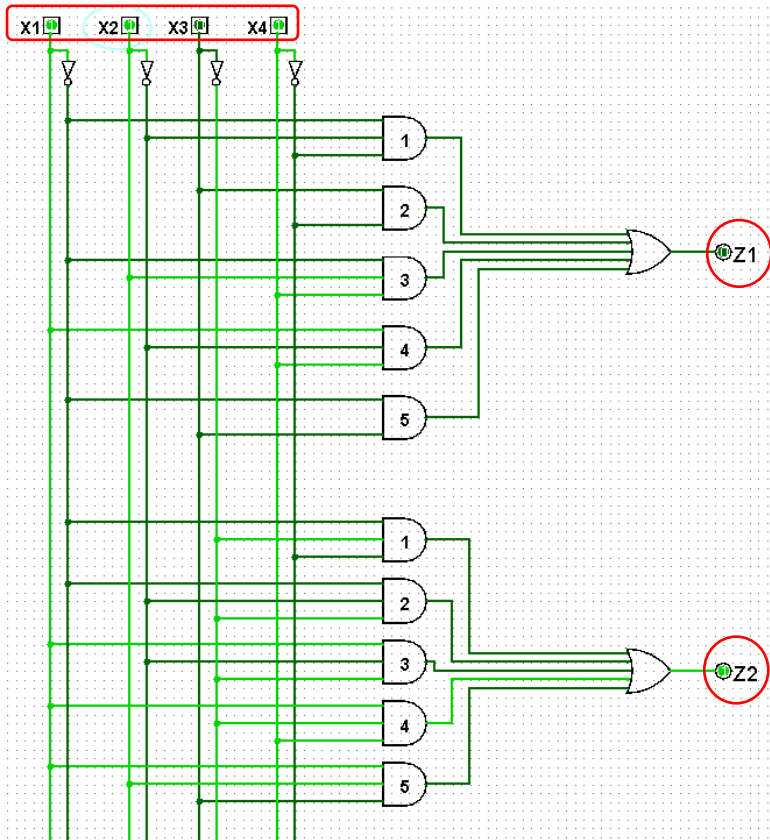
Test case 2:

	X1	X2	X3	X4	Z1	Z2
10	1	0	0	1	1	1



Test case 3:

	X1	X2	X3	X4	Z1	Z2
14	1	1	0	1	0	1



Task 2.1.1 – Your name as MARIE string

This task is to create a MARIE file that stores my name in memory. Below is a snapshot of the memory after the code is assembled with my name 'GLORIA(space)OOI'.

	+0	+1	+2	+3	+4	+5	+6	+7	+8	+9	+A	+B	+C	+D	+E	+F	
000	0047	004C	004F	0052	0049	0041	0020	004F	004F	0049	0000	0000	0000	0000	0000	0000	▲
010	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	■
020	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	
030	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	
040	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	
050	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	
060	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	
070	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	
080	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	
090	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	
0A0	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	
0B0	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	▼

Task 2.1.2 – Printing a String

This task requests the program to print out the name stored in memory. A loop is created to iterate each string (in ASCII) until it reaches the value HEX 000. Hence, when I stored my name, the last value I stored is HEX 000. For each iteration, the programme will check if the value is 0 and end loop.

Below is the output:

Output log	Register log	Watch list
<div>GLORIA OOI</div>		

Task 2.1.3 – A Subroutine for Printing a String

In Task 3, the address of the start of the string is 1st loaded and stored in *PrintString*. After that, the programme jumps into a subroutine called Print, iterates and prints each string found in memory until it reaches HEX 000, then exit the subroutine.

Output log	Register log	Watch list
		<div>GLORIA</div> <div>OOI</div>

This task requires the programme to ask for user input. As can be seen from the test cases below, memory will store the value entered by user.

[illegible][illegible]

Memory after input 1:

[illegible]

Test case 2:

The screenshot shows the 8086 simulator interface. The assembly code window displays the following code:

```

7  END, MAIN
8
9  USERINPUT,  HEX 0
10 LOOP,      LOAD BASE          /load address
11           ADD I              /add iterator
12           STORE ADDR         /store address+iterator in
13 INPUT      /request user input
14           SKIPCOND 400       /if input==0,
15           JUMP MAIN          /jump to MAIN
16           JUMPI USERINPUT    /exit subroutine
17 MAIN,      STOREI ADDR       /store input in ADDR's val
18           LOADI ADDR         /testing
19           OUTPUT

```

The memory dump window shows the following data:

Address	000	001	002	003	004	005	006	007	008	009	00A	00B	00C	00D	00E	00F
000	1014	2013	0004	7000	0003	1013	3017	2015	5000	8400	900C	C004	E015	D015	6000	1017
010	3016	2017	9005	001A	001A	001B	0001	0001	0000	0000	0047	0000	0000	0000	0000	0000
020	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000
030	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000
040	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000
050	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000
060	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000
070	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000
080	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000
090	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000
0A0	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000
0B0	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000

The CPU status window shows the following values:

Register	Value
MBR	001B
PC	009
IN	0047
OUT	0047

The OUTPUT MODE is set to ASCII.

An "Input Value" dialog box is open, prompting the user to input a value. The value "069" is entered, and the type is set to "Hexadecimal".

Memory after input 2:

[illegible]

Task 2.1.5 – Lower Case

This task is to convert any upper-case letters into lowercase and ignore those that are not alphabets, such as symbols.

Test case 1:

In test case one, the string stored is 'GLORIA(space)OOI', and the output is converted to 'gloria(space)ooi'. Here, the space is ignored as it is not an alphabet.

Assembly code:

Autosaved file

```
39 SIXFOUR,          DEC 64
40 THREETWO,         DEC 32
41 NINEONE,          DEC 91
42 ADDR,             HEX 0
43 ONE,              HEX 1
44 NUM,              DEC 43
45 PrintString,       HEX 0  /argument: address of string
46 I,                 HEX 0  /iterator
47 GLORIA_OOI,        HEX 047 /G
48                   HEX 04C /L
49                   HEX 04F /O
50                   HEX 052 /R
51                   HEX 049 /I
52                   HEX 041 /A
53                   HEX 020 /(space)
54                   HEX 04F /O
55                   HEX 04F /O
56                   HEX 049 /I
57                   HEX 000 /end
```

Machine halted normally.

Output log

Register log

Watch list

g
l
o
r
i
a

o
o
i

OUTPUT MODE: ASCII ▼

Test case 2:

In test case 2, the string stored is 'Hello(space)WoRlD'. It has a mixture of upper and lower case letters. The output is 'hello(space)world'.

Assembly code:

Modified file

```
40 THREETWO,          DEC 32
41 NINEONE,           DEC 91
42 ADDR,              HEX 0
43 ONE,               HEX 1
44 NUM,               DEC 43
45 PrintString,        HEX 0    /argument: address of string
46 I,                  HEX 0    /iterator
47 STRING,             HEX 048 /H
48                     HEX 065 /e
49                     HEX 04C /L
50                     HEX 06C /l
51                     HEX 04F /O
52                     HEX 020 /(space)
53                     HEX 057 /W
54                     HEX 06F /o
55                     HEX 052 /R
56                     HEX 06C /l
57                     HEX 044 /D
58                     HEX 000 /end
```

MARIE.js

Output log

Register log

Watch list

h
e
l
l
o

w
o
r
l
d

OUTPUT MODE: ASCII ▼

Task 2.1.6 – ROT13

This task require the programme to print out a string that is 13 places further in the alphabet, wrapping around from **z to a**. The programme will only make the conversion if the input is lower case alphabets, if not, the input will be left as it is.

Test case 1:

The string stored is my name 'gloria(space)ooi'. After performing ROT13, the output is 'tybevn(space)bbv'.

Assembly code:

Autosaved file

```
46 NINESIX,          DEC 96
47 ONETWOTWO,        DEC 122
48 ONETWOTHREE,       DEC 123
49 ADDR,             HEX 0
50 ONE,              HEX 1
51 NUM,              DEC 50
52 PrintString,       HEX 0  /argument: address of string
53 I,                 HEX 0  /iterator
54 GLORIA_OOI,        HEX 067 /g
55                   HEX 06C /l
56                   HEX 06F /o
57                   HEX 072 /r
58                   HEX 069 /i
59                   HEX 061 /a
60                   HEX 020 /(space)
61                   HEX 06F /o
62                   HEX 06F /o
63                   HEX 069 /i
64                   HEX 000 /end
```

Output log

Register log

Watch list

t
y
b
e
v
n

b
b
v

Test case 2:

The input for test case 2 is 'ap*G' and the output is 'nc*G', proving that only lower case alphabets are converted.

Assembly code:

Modified file

```
40      ADD ONE          /I+1
41      STORE I          /I = I+1
42      JUMP LOOP
43  CURRENTVALUE,        HEX 0
44  ONETHREE,            DEC 13
45  SIXFOUR,             DEC 64
46  NINESIX,            DEC 96
47  ONETWOTWO,          DEC 122
48  ONETWOTHREE,        DEC 123
49  ADDR,               HEX 0
50  ONE,                HEX 1
51  NUM,                DEC 50
52  PrintString,        HEX 0 /argument: address of string
53  I,                  HEX 0 /iterator
54  STRING,             HEX 061 /a
55                      HEX 070 /p
56                      HEX 02A /*
57                      HEX 047 /G
58                      HEX 000 /end
```

Machine halted normally.

Output log

Register log

Watch list

n
c
*
G

OUTPUT MODE: ASCII ▼

For this task, a user is to input a string. If the string is upper-case, the programme will convert it into lower-case, and then perform ROT13. Strings that are not alphabets will be ignored.

[illegible][illegible][illegible]

This is the output:

n	▲
\	
c	▼