

Fall 2019

# System Software Experiment 2

# Thread Synchronization

---

Prof. Jinkyu Jeong ([jinkyu@skku.edu](mailto:jinkyu@skku.edu))

TA – Gyusun Lee ([gyusun.lee@csi.skku.edu](mailto:gyusun.lee@csi.skku.edu))

TA – Jiwon Woo ([jiwon.woo@csi.skku.edu](mailto:jiwon.woo@csi.skku.edu))

Computer Systems and Intelligence Laboratory (<http://csi.skku.edu>)

Sungkyunkwan University

# Example 1

- Make program
  - 5 Reader + 1 Writer
  - Writer updates value 1,000,000 times
  - Each reader reads value 10,000,000 times
- Use pthread\_mutex

## Example 2

- Implement same thing using pthread\_spinlock
  - pthread\_spinlock\_t s
  - pthread\_spin\_init(&s, PTHREAD\_PROCESS\_PRIVATE)
  - pthread\_spin\_[un]lock(&s)
  - pthread\_spin\_destroy(&s)
- What is the difference?

# Readers-Writer Lock

- Reader blocks other reader
  - Do we need this?
- Implementation
  - Using two mutexes
  - Using a condition variable and a mutex

# Readers-Writer Lock(1)

## ■ Two mutex

### Begin Read

- Lock  $r$ .
- Increment  $b$ .
- If  $b = 1$ , lock  $g$ .
- Unlock  $r$ .

### End Read

- Lock  $r$ .
- Decrement  $b$ .
- If  $b = 0$ , unlock  $g$ .
- Unlock  $r$ .

### Begin Write

- Lock  $g$ .

### End Write

- Unlock  $g$ .

# Readers-Writer Lock(2)

- A condition variable and a mutex
  - Lock for read

- Input: mutex  $m$ , condition variable  $c$ , integer  $r$  (number of readers waiting), flag  $w$  (writer waiting).
- Lock  $m$  (blocking).
- While  $w$ :
  - wait  $c, m^{[a]}$
- Increment  $r$ .
- Unlock  $m$ .

- Lock for write

- Lock  $m$  (blocking).
- While  $w$ :
  - wait  $c, m$
- Set  $w$  to true.
- While  $r > 0$ :
  - wait  $c, m$
- Unlock  $m$ .

# Example 3-4

- Example 3
  - Implement program with readers-writer lock(Using mutex)
- Example 4
  - Implement program with readers-writer lock(Using condition variables)
- What is the difference?