

CSC343 Worksheet 12 Solution

June 30, 2020

1.
 - Keys
 - {id of molecule}
 - {x position, y position, z position}
 - Functional Dependencies
 - 1. id of molecule \rightarrow x position, y position, z position, x velocity, y velocity, z velocity
 - 2. x position, y position, z position \rightarrow id of molecule, x velocity, y velocity, z velocity

Notes:

- Function Dependencies
 - *Functional Dependency* is a relationship between two attributes typically between the key and other non-key attributes within a table.

Example:

SIN \rightarrow Name, Address, Birthdate

Example 2:

ISBN \rightarrow Title

- Key of Relations
 - One or more attributes $\{A_1, A_2, \dots, A_n\}$ is a key for a relation R if
 1. Those attributes functionally determine all other attributes of the relation
 2. No proper subset of $\{A_1, A_2, \dots, A_n\}$ functionally determines all other attributes of R

Example:

Given relation

$R = \text{Movies1}(\text{title, year, length, genre, studioName, starName})$

- i. {title, year, starName } form a key for the relation **Movies1**

- ii. $\{ \text{year}, \text{starName} \}$ is not a key. Same star can be in multiple movies per year
- Superkeys
 - * Means a set of attributes that contains a key
 - * Don't need to be minimal

Example:

Given relation

 $R = \text{Movies1}(\text{title}, \text{year}, \text{length}, \text{genre}, \text{studioName}, \text{starName})$

- $\{ \text{title}, \text{year}, \text{starName} \}$ is a key and superkey
- $\{ \text{title}, \text{year}, \text{starName}, \text{title}, \text{year}, \text{length} \}$ is a superkey

References:

- 1) OpenTextBC, Chapter 11 Functional Dependencies, link
2. a)
 1. $AB \rightarrow C$
 2. $AB \rightarrow D$
 3. $C \rightarrow A$
 4. $C \rightarrow B$
 5. $D \rightarrow B$
 6. $D \rightarrow C$
 7. $C \rightarrow D$
 8. $D \rightarrow A$

Second Attempt:

$\{A, B\}^+ = \{A, B, C, D\}$, so the following non-trivial FDs follows: $AB \rightarrow C$ and $AB \rightarrow D$.

$\{C\}^+ = \{D, A\}$, so the following non-trivial FDs follows $C \rightarrow D$ and $C \rightarrow A$.

$\{D\}^+ = \{A\}$, so the following non-trivial FDs follows: $D \rightarrow A$.

Notes:

- The Splitting / Combining Rule
 - Combining Rule
 - * $A_1, A_2, \dots, A_n \rightarrow B_i$ for $i = 1, 2, \dots, m$
to
 $A_1, A_2, \dots, A_n \rightarrow B_1, B_2, \dots, B_m$

Example:

Given

title year \rightarrow length
 title year \rightarrow genre
 title year \rightarrow studioName

it's combined form is

title year \rightarrow length genre studioName

– Splitting Rule

*

* $A_1, A_2, \dots, A_n \rightarrow B_1, B_2, \dots, B_m$

to

$A_1, A_2, \dots, A_n \rightarrow B_i$ for $i = 1, 2, \dots, m$

Example:

Given

title year \rightarrow length

It's splitted form is

title \rightarrow length
 year \rightarrow length

• Trivial Functional Dependencies

- A functional dependency $FD : X \rightarrow Y$ is **trivial** if Y is a subset of X

Exmample:

title year \rightarrow title

Example 2:

title \rightarrow title

• Non-trivial Functional Dependencies

- is a case where some but not all of the attributes on the R.H.S of an FD are also on L.H.S

Example:

title year \rightarrow title movieLength

- Can be simplified using **trivial-dependency rule**
 - * The FD $A_1A_2 \cdots A_n \rightarrow B_1B_2 \cdots B_m$ is equivalent to $A_1A_2 \cdots A_n \rightarrow C_1C_2 \cdots C_k$

where C 's are all those B 's that are not in A 's.



Figure 3.3: The trivial-dependency rule

- Computing the Closure of Attributes
 - Closure of attribute set $\{X\}$ is denoted as $\{X\}^+$.
 - The closure means a given set of attributes A satisfying FD, are a sets of all attributes B such that $A \rightarrow B$

Example:

Given attributes A, B, C, D, E, F and FDs $AB \rightarrow C$, $BC \rightarrow AD$, $D \rightarrow E$ and $CF \rightarrow B$, What is the closure of $\{A, B\}$ or $\{A, B\}^+$

1. Start with $\{A, B\}$.
2. Split $BC \rightarrow AD$
 - * We have $BC \rightarrow A$ and $BC \rightarrow D$
 - * Since A is in $\{A, B\}$, this is not included
 - * Since D is not in $\{A, B\}$, this IS included

So, we have $\{A, B, D\}$

3. Since C in $AB \rightarrow C$ is NOT in $\{A, B, C, D\}$, C is included and we have $\{A, B, C, D\}$
4. Since A in $BC \rightarrow A$ is in $\{A, B, C, D\}$, this is skipped
5. Since E is not in $D \rightarrow E$, E is included and we have $\{A, B, C, D, E\}$ as our solution

- Why the Closure Algorithm Works
- Transitive Rule
 - Definition

If $A_1A_2 \cdots A_n \rightarrow B_1B_2 \cdots B_m$ and $B_1B_2 \cdots B_m \rightarrow C_1C_2 \cdots C_k$ hold in relation R , $A_1A_2 \cdots A_n \rightarrow C_1C_2 \cdots C_k$ also holds in R .

Example:

Given

title year \rightarrow studioName
 studioName \rightarrow studioAddr

Transitive rule says the above is equal to the following

title year \rightarrow studioAddr

- Inference Rules
 - Is also called **Armstrong's Axioms**
 - Has 3 axioms
 1. *Reflexivity*
 - * If $\{B_1, B_2, \dots, B_n\} \subseteq \{A_1, A_2, \dots, A_n\}$ then $A_1A_2 \cdots A_n \rightarrow B_1B_2 \cdots B_m$
 - * also called **trivial FDs**
 2. *Augmentation*
 - * If $A_1A_2 \cdots A_n \rightarrow B_1B_2 \cdots B_m$ then $A_1A_2 \cdots A_nC_1C_2 \cdots C_k \rightarrow B_1B_2 \cdots B_mC_1C_2 \cdots C_k$
 - * $C_1C_2 \cdots C_k$ are any set of attributes
 3. *Transitivity*
 - * If $A_1A_2 \cdots A_n \rightarrow B_1B_2 \cdots B_m$ and $B_1B_2 \cdots B_m \rightarrow C_1C_2 \cdots C_k$ then $A_1A_2 \cdots A_n \rightarrow C_1C_2 \cdots C_k$

b) A, B is the only key of R .

Notes:

- Key of Attributes
 - **Definition:** A set of attributes $\{A_1, A_2, \dots, A_n\}$ is a key for a relation R if
 1. Those attributes functionally determine all other attributes

2. No proper subset of $\{A_1, A_2, \dots, A_n\}$ functionally determines all other attributes of R .
- c) The superkeys that are not keys are: $\{A, B, C\}$, $\{A, B, D\}$, $\{A, B, C, D\}$