

# CSC343 Worksheet 13 Solution

July 4, 2020

1. a)

A	B	C	D	E
$a$	$b$	$c$	$d_1$	$e_1$
$a_1$	$b$	$c$	$d_1$	$e_2$
$a$	$b_1$	$c$	$d_1$	$e$

**Step 1 ( $B \rightarrow E$ ):**

A	B	C	D	E
$a$	$b$	$c$	$d_1$	$e_1$
$a_1$	$b$	$c$	$d_1$	$e_1$
$a$	$b_1$	$c$	$d_1$	$e$

**Step 2 ( $CE \rightarrow A$ ):**

A	B	C	D	E
$a$	$b$	$c$	$d_1$	$e_1$
$a$	$b$	$c$	$d_1$	$e_1$
$a$	$b_1$	$c$	$d_1$	$e$

So in this case, an example of an instance of  $R$  that is not lossless is:

Title	Year	Studio Name	President	President Address
Toy Story	2000	Pixar	Steve Jobs	123 ABC Street
Star Wars	1977	Fox	Lachlan Murdoch	Hollywood
Return of the Jedi	1983	Fox	Lachlan Murdoch	Hollywood

**Notes:**

- Decomposition: The good bad and ugly
  - 1) **Elimination of Anomalies** by decomposition as in Section 3
  - 2) **Recoverability of Information** Can we recover the original relation from the tuples in its decomposition?

- 3) **Preservation of Dependencies (lossless join):** Can we be sure that after reconstructing the original relation from the decompositions, the original FD's satisfy?

**BCNF:**  $\rightarrow$  satisfies 1) and 2) **Not good. NONO**

- The Chase Test for Lossless Join
  - Tests whether the decomposition is lossless

**Input:**

- A relation  $R$
- A decomposition of  $R$
- A set of functional dependencies

**Output:**

- Whether the decomposition is lossless or not
- $\Pi_{S_1}(R) \bowtie \Pi_{S_2}(R) \bowtie \cdots \Pi_{S_i}(R) = R$

**Three things to remember:**

1. The natural join is associate and commutative
2. Any tuple  $t$  in  $R$  is surely in  $\Pi_{S_1}(R) \bowtie \Pi_{S_2}(R) \bowtie \cdots \bowtie \Pi_{S_k}(R)$ .
3. We have to check to see any tuple in the  $\Pi_{S_1}(R) \bowtie \Pi_{S_2}(R) \bowtie \cdots \bowtie \Pi_{S_k}(R)$ .

**Example:**

$$S_1 = \{A, D\}, S_2 = \{B, C\}, S_3 = \{A, C\}$$

$$A \rightarrow B, B \rightarrow C, CD \rightarrow A$$

A	B	C	D
a	b <sub>1</sub>	c <sub>1</sub>	d
a	b <sub>2</sub>	c	d <sub>2</sub>
a <sub>3</sub>	b	c	d

a<sub>i</sub> represents arbitrary value

Represents S<sub>1</sub> = {A,D}

Represents S<sub>2</sub> = {A,C}

Represents S<sub>3</sub> = {B,C}

**Step 1:  $A \rightarrow B$**

Set the value  $b$  with the same value of  $a$  to be the same. (e.g.  $b_2 \rightarrow b_1$ )

1. The value of  $a$  is the same

2. Change the value of  $b_2$  to  $b_1$

$A$	$B$	$C$	$D$
$a$	$b_1$	$c_1$	$d$
$a$	$b_1$	$c$	$d_2$
$a_3$	$b$	$c$	$d$

### Step 2: $B \rightarrow C$

Set the value  $c$  with the same value of  $b$  to be the same. (e.g.  $b_2 \rightarrow b_1$ )

1. The value of  $b$  is the same

2. Change the value of  $c_1$  to  $c$

$A$	$B$	$C$	$D$
$a$	$b_1$	$c$	$d$
$a$	$b_1$	$c$	$d_2$
$a_3$	$b$	$c$	$d$

### Step 3: $CD \rightarrow A$

Set the value  $a$  with the same value of  $c$  and  $d$  to be the same. (e.g.  $a_3 \rightarrow a$ )

2. Change the value of  $a_3$  to  $a$  (e.g.  $a_3$  to  $a$ )

1. The value of  $c$  and  $d$  are the same

3. The value of  $a, b, c, d$  the same as  $(a, b, c, d)$  in  $R$ !

$A$	$B$	$C$	$D$
$a$	$b_1$	$c$	$d$
$a$	$b_1$	$c$	$d_2$
$a$	$b$	$c$	$d$

So, we can conclude the join is lossless.