CSC373 Worksheet 5 Solution

August 8, 2020

1. **Rough Works:**

   Assume that a flow network $G = (V, E)$ violates the assumption that the network contains a path $s \rightsquigarrow v \rightsquigarrow t$ for all vertices $v \in V$. Let u be a vertex for which there is no path $s \rightsquigarrow u \rightsquigarrow t$.

   I must show such that there is no flow at vertex u. That is, there exists a maximum flow f in G such that $f(u, v) = f(v, u) = 0$ for all vertices $v \in V$.

   Assume for the sake of contradiction that there is some vertex u with flow f. That is, there exists some vertices $v \in V$ such that $f(u, v) > 0$ or $f(v, u) > 0$.

   I see that three cases follows, and I will prove each separately.

   1. **Cases 1:** $f(u, v) > 0$ and $f(v, u) = 0$
   2. **Cases 2:** $f(u, v) = 0$ and $f(v, u) > 0$
   3. **Cases 3:** $f(u, v) > 0$ and $f(v, u) > 0$

      Here, assume that $f(u, v) > 0$ and $f(v, u) > 0$.

      - Write that the path $s \rightsquigarrow u \rightsquigarrow t$ exists

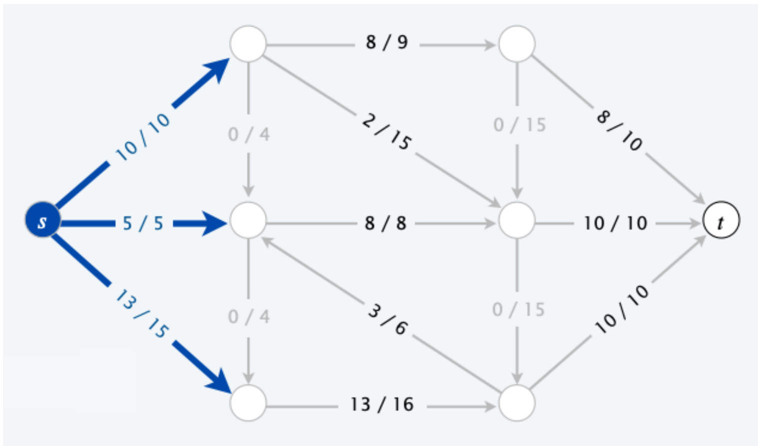
Since $s \rightsquigarrow v \rightsquigarrow t$ and u is connected by some vertices v, we can write $s \rightsquigarrow u \rightsquigarrow t$.

      - Write that this results in contradiction to the header that a vertex u has no path $s \rightsquigarrow u \rightsquigarrow t$.
      - Conclude that $f(u, v) = 0$ and $f(v, u) = 0$

   Notes

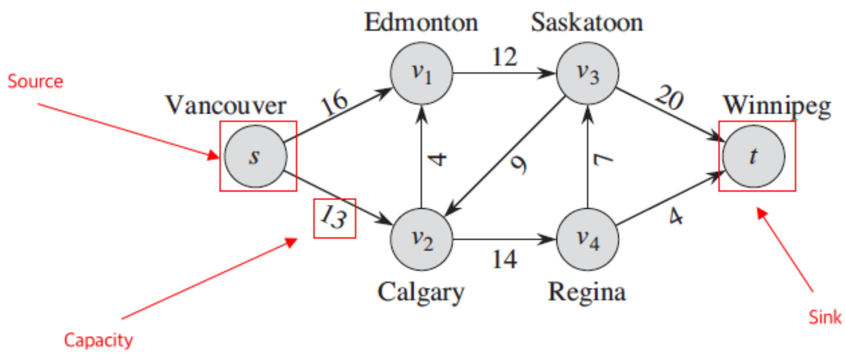
- **Maximum Flow:**
  - Finds a flow of maximum value [1]

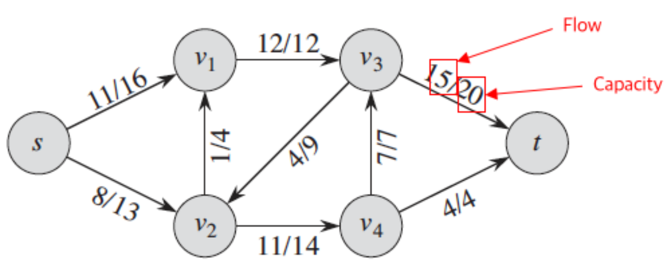
  Example



  Here, the maximum flow is $10 + 5 + 13 = 28$

- **Flow Network:**
  - $G = (V, E)$ is a directed graph in which each edge $(u, v) \in E$ has a nonnegative capacity $c(u, v) \geq 0$.
  - Two vertices must exist: **source** s and **sink** t
  - **path** from source s to vertax v to sink t is represented by $s \rightsquigarrow v \rightsquigarrow t$





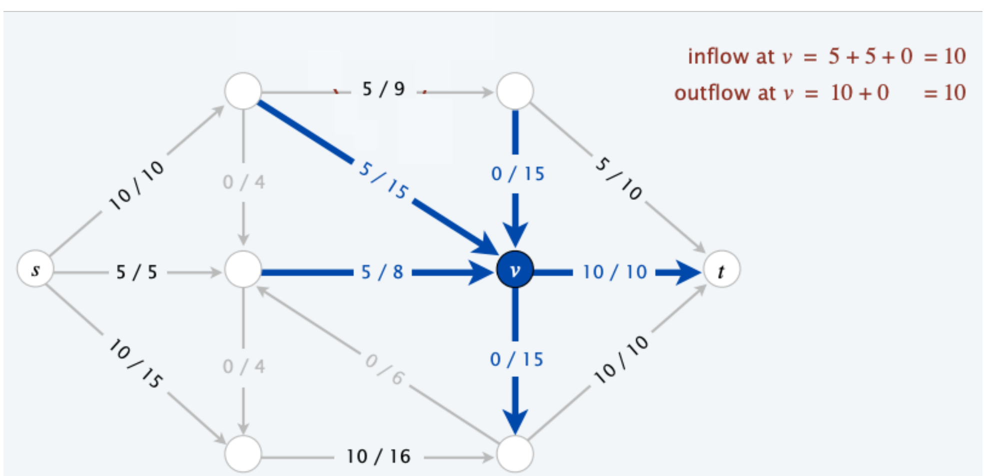
- **Capacity:**
  - Is a non-negative function $f : V \times V \to \mathbb{R}_{\geq 0}$
  - Has **capacity constraint** where for all $u, v \in V$ $0 \leq f(u, v) \leq c(u, v)$
    * Means flow cannot be above capacity constraint
- **Flow:**
  - Is a real valued function $f : V \times V \to \mathbb{R}$ in G
  - Satisfies **capacity constraint** (i.e for all $u, v \in V$, $0 \leq f(u, v) \leq c(u, v)$)
  - Satisfies **flow conservation**

    For all $u \in V - \{s, t\}$, we require

$$\sum_{v \in V} f(v, u) = \sum_{v \in V} f(u, v) \tag{1}$$

Means flow into vertex u is the same as flow going out of vertex u. [1]

Example:



References

1) Princeton University, Network Flow 1, link