September 1, 2020

1. a) 0

Notes

- a) is 0 because (i >> 1 + j >> 1 = i >> 10 >> 1 = 0)
- Bitwise Shift Operators
 - has lower precedence than arithematic operators

Example:

```
i << 2 + 1 \text{ means } i << (2+1) \text{ and not } (i << 2) + 1
```

- << : Left Shift
- >> : Right Shift
- Tip: Always shift only on unsigned numbers for portability

Example

- >> = / << = : Are bitwise shift equivalent of + = (

b) 0

Notes

- i is 1111111111111111
- i is 0000000000000000
- so i & i = 0
- : Bitwise complement (NOT)

a	\sim a
0	1
1	0

Example:

```
1 0 1 1 1 //<- this is 7
2 -------
3 1 0 0 0 //<- this is 8
4
5 so, ~ 7 = 8
```

• &: Bitwise and

a	b	a & b
0	0	0
0	1	1
1	0	0
1	1	1

Example:

```
0 1 1 1 //<- this is 7
0 1 0 0 //<- this is 4

------
0 1 0 0 //<- this is 4

so, 7 & 4 = 4
```

- ullet : Bitwise exclusive or
- \bullet |: Bitwise inclusive or
- c) 1

Notes

- i is 111111111111110
- j is 000000000000000
- $\bullet\,$ i & j is 0000000000000000 or 1
- i & j ^ k is 1

• ^: Bitwise XOR

a	b	a ^ b
0	0	0
0	1	1
1	0	1
1	1	0

Example:

d) 0

Example

- i is 000000000000111
- j is 000000000001000
- \bullet i ^ j is 0000000000000000 or 0
- k is 000000000001001
- i ^ j & k is 0000000000000000 or 0

Correct Solution

15

\underline{Notes}

• There is a precendence to the order of operations



e) • toggling from 0 to 1

```
i = 0x0000;
i |= 0x0001;
or
i |= 1 << 0; where i = 0x0000;
• toggling from 1 to 0

i = 0x0001;
i &= ~0x0001;
or
i &= ~(1 << 0); where i = 0x0001;</pre>
```

Correct Solution

• toggling from 0 to 1 of 4th bit

```
i = 0x0010;
i ^= 0x0000;
or
i ^= 1 << 4; where i = 0x0000;
• toggling from 1 to 0 of 4th bit

i = 0x0010;
i ^= 0x0010;
or
i ^= (1 << 4); where i = 0x0010;</pre>
```

Notes

- Toggling can be done using bitwise XOR
- Setting a bit
 - Is done using | or bitwise OR

- The idiom of above is $i = 1 \ll j$

- Clearing a bit
 - Is done using | or bitwise AND

- The idiom of above is i &= \sim (i << j)
- 2. It swaps the elements between x and y.

Notes

• Preprocessor performs operations of statements in order from left to right

#define
$$M(x,y)$$
 ((x)^=(y), (y)^=(x), (x)^=(y))

New value of y, using x from 1

New value of x, using y from 2, x from 1

3. #define MK_COLOR(r,g,b) (long) ((b | (g << 8)) | (b | (r << 16)))

Rough Work

- 1. store b in bit 0
 - b
- 2. store g in bit 8
 - b | g << 8
- 3. store r in bit 16
 - b | r << 16

Notes

• First Byte is furthest from 0x and first byte is closest to 0x



4. • GET_RED

```
#define GET_RED(c) (long) (c & 0x007)
```

• GET_GREEN

```
#define GET_GREEN(c) (long) ((c >> 8) & 0x007)
```

• GET_BLUE

```
#define GET_BLUE(c) (long) ((c >> 16) & 0x007)
```

Notes

- 0x0007 in binary is 0x000000000001111
- c >> 4 shifts c to right by 4 bits and return overlapping value between c >> 4 and 0x00000000001111 (0x007)
- Test code is below

```
#include <stdio.h>
#include <stdlib.h>

#define MK_COLOR(r,g,b) (long) ( (r | (g << 8)) | (r | (b << 16))

#define GET_RED(c) (long) (c & 0x007)
#define GET_GREEN(c) (long) ((c >> 8) & 0x007)
#define GET_BLUE(c) (long) ((c >> 16) & 0x007)

#define GET_BLUE(c) (long) ((c >> 16) & 0x007)

int main() {
    long i, r = 4, g = 5, b = 6, r2, g2, b2;

i = MK_COLOR(r,g,b);
```

```
5. a)
         unsigned short swap_bytes(unsigned short i) {
              unsigned short j, k;
    3
              j = i \& 0x007; // extract first byte
    4
   5
              i = i >> 4;
   6
              k = i \& 0x007; // extract second byte
   8
              i = i >> 4; // shift down layter two bytes
   9
   10
              i \mid = j \ll 8; // add first byte to position of fourth byte
   11
              i \mid= k << 12; // add second byte to position of third byte
   12
   13
  b)
         unsigned short swap_bytes(unsigned short i) {
```

```
b) unsigned short swap_bytes(unsigned short i) {
    i = i >> 8 | i << 8;
    return i;
    }
}</pre>
```

Rough Works

1. Extract first two bytes

```
j = i & 0x0007

i = i >> 4

k = i & 0x0007
```

2. Shift later two bytes down

```
i = i >> 4
```

3. Add first two bytes to last two bytes

```
i |= j << 8;
i |= k << 12;
```

```
unsigned int rotate_left(unsigned int i, int n) {
    return i >> 28 | i << n;
}

unsigned int rotate_right(unsigned int i, int n) {
    return i << 28 | i >> n;
}
```

- 7. a) Is a binary with n many 1s from the first bit
 - b) Extracts last n bits in i

Correct Solution

- a) Is a binary with n many 1s from the first bit
- b) Extracts last n bits starting from position m in i

```
8. a)     unsigned int count_ones(unsigned char *ch)
{
        unsigned char *p;
        unsigned int count;

        for (p = ch; *p != '\0'; p++){
            if (*p == '1') {
                count++;
            }
        }
        }

return count;
}
```

```
b)
       unsigned int count_ones(unsigned char ch)
       {
 2
 3
           int sum = 0;
           sum += (i >> 0) & 1;
 5
           sum += (i >> 1) & 1;
 6
           sum += (i >> 2) & 1;
           sum += (i >> 3) & 1;
 9
           sum += (i >> 4) & 1;
           sum += (i >> 5) & 1;
10
           sum += (i >> 6) & 1;
```

```
sum += (i >> 7) & 1;
return count;
}
```

Notes

- Unsigned char goes from 0 (00000000) to 255 (111111111)
- I am having trouble how to convert from loop to without loop: '(. I need help
- Example

100010101 - Here there are 4 1s.

```
9_1
       unsigned int reverse_bits (unsigned int n)
 2
           int m = 15, p = 0;
3
           unsigned int byte_left_end, byte_right_end, res = 0;
 5
           while (m > p) {
 6
                byte_left_end = n >> m & 1;
                byte_right_end = n >> p & 1;
9
               res |= byte_left_end << (15 - m) | byte_right_end << (15 - p)
10
      );
11
12
               m--;
13
               p++;
           }
14
           return res;
15
```

Rough Work

- Start at n = 15, m = 0
- Swap bit at n with m
- Repeat until m > n

<u>Notes</u>

- unsigned int has 4 bytes or (0x0000) or (000000000000000, 16 bits)
- 10. The precedence of &, ^, and is lower than the precedence of the relational and equality operators.

As a result, the expression will first evaluate

```
(SHIFT_BIT | CTRL_BIT | ALT_BIT) == 0
```

first.

The work around is to put parenthesis around $key_code \& (SHIFT_BIT | CTRL_BIT | ALT_BIT)$.

11. The function tries to insert low_byte to high_byte after shifting high_byte by 8 bits to left.

It doesn't work because + sign takes precedence over <<.

To fix this issue, a parenthesis is required around high_byte << 8.

12. It reduces the value of n by 1, then uses bitwise operator (&) to extract the common bits between n and n-1.