Worksheet 16 Review

April 3, 2020

Question 1

a. Let $k \in \mathbb{N}$.

Here, the minimum possible change occurs for the loop variable in a single iteration when $i = i + 1$.

The maximum possible change occurs for the loop variable in a single iteration when $i = i + 6$.

The exact upper bound of the variable after k iteration is

$$i_k \leq 6k \tag{1}$$

The exact lower bound of the variable after k iteration is

$$k \leq i_k \tag{2}$$

Using the fact that the termination occurs when $i_k = n$, we can calculate that for the upper bound, the loop terminates when

$$6k \geq n \tag{3}$$
$$k \geq \frac{n}{6} \tag{4}$$

Because we know $\frac{n}{6}$ may be a decimal, we can conclude the closest value at which the loop terminates is when

$$k = \left\lceil \frac{n}{6} \right\rceil \tag{5}$$

Using the same fact, we can calculate that for the lower bound, the loop terminates when

$$k \geq n \tag{6}$$

It follows from above that for the lower bound, the smallest value of k at which the loop termination occurs is when

$$k = n \tag{7}$$

Then, we can conclude the function has asymptotic lower bound of $\Omega(n)$, and asymptotic upper bound of $\mathcal{O}(n)$.

Then, since both Ω and $\mathcal{O}$ have the same value, $\Theta(n)$ is also true.

Correct Solution:

Here, the minimum possible change occurs for the loop variable in a single iteration when $i = i + 1$.

The maximum possible change occurs for the loop variable in a single iteration when $i = i + 6$.

The exact upper bound of the variable after k iteration is

$$i_k \leq 6k \tag{8}$$

The exact lower bound of the variable after k iteration is

$$k \leq i_k \tag{9}$$

Using the fact that the termination occurs when $i_k = n$, we can calculate that for the upper bound, the loop terminates when

$$6k \geq n \tag{10}$$

$$k \geq \frac{n}{6} \tag{11}$$

Because we know $\frac{n}{6}$ may be a decimal, we can conclude the closest value at which the loop terminates is when

$$k = \left\lceil \frac{n}{6} \right\rceil + 1 \tag{12}$$

Using the same fact, we can calculate that for the lower bound, the loop terminates when

$$k \geq n \tag{13}$$

It follows from above that for the lower bound, the smallest value of k at which the loop termination occurs is when

$$k = n + 1 \tag{14}$$

Then, we can conclude the function has asymptotic lower bound of $\Omega(n)$, and asymptotic upper bound of $\mathcal{O}(n)$.

Since both Ω and $\mathcal{O}$ have the same value, $\Theta(n)$ is also true.

Notes:

- where is $\mathbf{+1}$ coming from? Is it coming from the loop variable $i = 0$?

b. Let $k \in \mathbb{N}$.

Part 1 (Determining maximum and minimum possible change in a single iteration):

It follows from observation that the minimum possible change occurs when $i = i \cdot 2$, and the maximum possible change when $i = i \cdot 3$.

Part 2 (Determining lower bound and upper bound of loop iteration):

Because we know the smallest possible change occurs when $i = i \cdot 2$ occurs repeatedly, we can conclude that at k^{th} iteration i_k has the lower bound of 2^k.

Similarly, because we know largest possible change occurs when $i = i \cdot 3$ occurs repeatedly, we can conclude that at k^{th} iteration, i_k has the upper bound of 3^k.

Then, by putting together, we can conclude that

$$2^k \leq i_k \leq 3^k \tag{1}$$

Part 3 (Determining exact number of iterations for the lower bound and upper bound):

Because we know the loop runs until $i_k < n$, we can conclude that at lower bound, termination occurs when

$$i_k \geq n \tag{2}$$
$$2^k \geq n \tag{3}$$
$$\log_2 2^k \geq \log_2 n \tag{4}$$
$$k \geq \log_2 n \tag{5}$$

Using the fact that we are looking for smallest value of k, we can calculate that for lower bound

$$k = \lceil \log_2 n \rceil + 1 \tag{6}$$

Similarly, for the upper bound, loop terminates when

$$i_k \geq n \tag{7}$$
$$3^k \geq n \tag{8}$$
$$\log_3 3^k \geq \log_3 n \tag{9}$$
$$k \geq \log_3 n \tag{10}$$

Using the fact, we can calculate that for upper bound,

$$k = \lceil \log_3 n \rceil + 1 \tag{11}$$

Part 4 (Determining Big-Oh and Omega):

Because we know $\log_2 n$ dominates $\log_3 n$, we can conclude $\log_2 n$ is the asymptotic upper bound, and $\log_3 n$ is the asymptotic lower bound.

Then, we can conclude the algorithm has $\mathcal{O}(\log_2 n)$ and $\Omega(\log_3 n)$.

Notes:

- How come in solution, **+1** doesn't exist? What rules of thumb i can follow to better determine whether **+1** should be included?

Question 2

a. Because we know $n \in \Theta(n^2)$, we can conclude the algorithm has runtime of $\Theta(n^2)$.

> **Correct Solution:**
>
> Since **helper1** has cost of n and **helper2** has cost of n^2, we can conclude the algorithm has total cost of $n + n^2$.
>
> It follows from above the algorithm has runtime of $\Theta(n^2)$.

Notes:

- When is $\in$ in $n \in \Theta(n^2)$ used?

  Is $\in \Theta$ used when $\mathcal{O}$ and Ω exists with different values to choose which value works for both lower and upper bound of the algorithm?

- Noticed that professor evaluates total runtime before Theta

b. Because we know loop 1 starts at $i = 0$ and finishes at $i = n - 1$ with i increasing by 2 per iteration, we can conclude loop 1 has

$$\left\lceil \frac{n - 1 - 0 + 1}{2} \right\rceil = \left\lceil \frac{n}{2} \right\rceil \tag{1}$$

iterations.

Since each iteration in loop 1 takes n step, as required by **helper 1** function, we can conclude loop 1 has total cost of

$$n \cdot \left\lceil \frac{n}{2} \right\rceil \tag{2}$$

steps.

For loop 2, because we know it starts at $j = 0$ and finishes at $j = 9$, we can conclude loop 2 has

$$\lceil 9 - 0 + 1 \rceil = 10 \tag{3}$$

iterations.

Since each iteration in loop 2 takes n^2 step as required by **helper 2** function, we can conclude loop 2 has total of

$$10 \cdot n^2 \tag{4}$$

steps.

Since $i = 0$ and $j = 0$ have cost of 1 step each, the total cost of algorithm is

$$n \cdot \left\lceil \frac{n}{2} \right\rceil + 10n^2 + 2 \tag{5}$$

Then, we can conclude the algorithm has running time of $\Theta(n^2)$

Correct Solution:

Because we know loop 1 starts at $i = 0$ and finishes at $i = n - 1$ with i increasing by 2 per iteration, we can conclude loop 1 has

$$\left\lceil \frac{n-1-0+1}{2} \right\rceil = \left\lceil \frac{n}{2} \right\rceil \tag{1}$$

iterations.

Since each iteration in loop 1 takes n step, as required by **helper 1** function, we can conclude loop 1 has total cost of

$$n \cdot \left\lceil \frac{n}{2} \right\rceil \tag{2}$$

steps.

For loop 2, because we know it starts at $j = 0$ and finishes at $j = 9$, we can conclude loop 2 has

$$\lceil 9 - 0 + 1 \rceil = 10 \tag{3}$$

iterations.

Since each iteration in loop 2 takes n^2 step as required by **helper 2** function, we can conclude loop 2 has total of

$$10 \cdot n^2 \tag{4}$$

steps.

Combining together , the total cost of algorithm is

$$n \cdot \left\lceil \frac{n}{2} \right\rceil + 10n^2 \tag{5}$$

Then, we can conclude the algorithm has running time of $\Theta(n^2)$

Notes:

- Noticed professor doesn't count loop variables toward the total cost of algorithm.

  If other lines such as **return False** and **n = len(lst)** are included, would these count towards the total cost of the algorithm?

c. For loop 1, because we know it starts at $i = 0$ and finishes at $i = n - 1$ with each iteration having cost of i steps, we can conclude loop 1 has cost of

$$\sum_{i=0}^{n-1} i = \frac{n(n-1)}{2} \tag{1}$$

steps.

For loop 2, because we know it starts at $j = 0$ and finishes at $j = 9$ with each iteration costing j^2 steps, we can conclude loop 2 has

$$\begin{aligned} \sum_{j=0}^{9} j^2 &= \frac{9(9-1)(2(9)-1)}{6} && (2) \\ &= \frac{9 \cdot 8 \cdot 17}{6} && (3) \\ &= 204 && (4) \end{aligned}$$

steps.

Combining together, the total cost of algorithm is

$$\frac{n(n-1)}{2} + 204 \tag{5}$$

steps.

Then, we can conclude the running time of algorithm is $\Theta(n^2)$.

Question 3

a. **Predicate Logic:** $\forall x \in \mathbb{Z}^+,\ k \in \mathbb{N}$, (three iterations occur) $\Rightarrow x_{3(k+1)} \leq \frac{x_{3k}}{2}$

Proof. Let $k \in \mathbb{N}$.

Assume three iterations occur in loop.

We will show $x_{3(k+1)} \leq \frac{x_{3k}}{2}$ using proof by cases.

Case 1 (x divisible by 2 three times):

Assume x is divisible by 2 at least 3 times.

Because we know the if condition $x \mod 2 == 0$ is true in all iterations, we can conclude the line **x = x // 2** will run 3 times.

Then, it follows from above that the value of x at $3(k+1)^{th}$ iteration is

$$x_{3(k+1)} = \frac{x_{3k}}{2^3} \tag{6}$$

Then, because we know $\frac{1}{2^3} \leq \frac{1}{2}$, we can conclude

$$x_{3(k+1)} \leq \frac{1}{2} x_{3k} \tag{7}$$

Case 2 (x divisible by 2 two times):

Assume x is divisible by 2, 2 times.

Because we know the if condition $x \mod 2$ **== 0** is true in first two iterations, we can conclude the line **x = x // 2** will run twice.

Then, at the end of second iteration, we can conclude x will have the value of

$$\frac{x_{3k}}{2^2} \tag{8}$$

On the final iteration, because we know the if condition $x \mod 2$ **== 0** is false, we can conclude the line **x = 2 * x - 2** will run.

Then, using the above fact, we can calculate

$$\begin{align} x_{3(k+1)} &= \frac{x_{3k}}{2^2} \cdot 2 - 2 \tag{9} \\ &= \frac{x_{3k}}{2} - 2 \tag{10} \\ &\leq \frac{x_{3k}}{2} \tag{11} \end{align}$$

Case 3 (x divisible by 2 once):

Assume x is divisible by 2 once.

Because we know the if condition $x \mod 2$ **== 0** is true in first iteration, we can conclude the line **x = x // 2** will run.

Then, at the end of first iteration, we can conclude x will have the value of

$$\frac{x_{3k}}{2} \tag{12}$$

On the second iteration, because we know $x \mod 2 == \mathbf{0}$ is false, we can conclude the line **x = 2 * x - 2** will run.

Then, at the ned of second iteration, we can conclude x will have the value of

$$\frac{x_{3k}}{2} \cdot 2 - 2 = 3_{3k} - 2 \tag{13}$$

On the final iteration, because we know from assumption $2 \mid x_{3k}$ and $2 \mid -2$, we can conclude the if condition $x \mod 2 == \mathbf{0}$ will be satisfied and the line **x = x // 2** will run.

Then, at the end of final iteration, we can conclude $x_{3(k+1)}$ will have the value of

$$\begin{align} x_{3(k+1)} &= \frac{x_{3k} - 2}{2} \tag{14} \\ &\leq \frac{x_{3k}}{2} \tag{15} \end{align}$$

Case 4 (x is an odd number):

Assume x is an odd number.

Because we know $x \mod 2 == \mathbf{0}$ is false in first iteration, we can conclude the line **x = 2 * x - 2** will run.

Then, at the end of first iteration, we can conclude x will have the value of

$$x_{3k} \cdot 2 - 2 = 2 \cdot (x_{3k} - 1) \tag{16}$$

On the second iteration, because we know the if condition $x \mod 2 ==$ **0** is true, we can conclude the line **x = x // 2** will run.

Then, at the end of second iteration, we can conclude x will have the value of

$$\frac{2 \cdot (x_{3k} - 1)}{2} = x_{3k} - 1 \tag{17}$$

On the final iteration, because we know $x_{3k} - 1$ is an even number, we can conclude the if condition $x \mod 2 ==$ **0** is true, and the line **x = x // 2** will run.

Then, at the end of final iteration, we can conclude $x_{3(k+1)}$ will have the value of

$$\begin{aligned} x_{3(k+1)} &= \frac{x_{3k} - 1}{2} &(18) \\ &\leq \frac{x_{3k}}{2} &(19) \end{aligned}$$

$\square$

b. Let $k \in \mathbb{N}$.

We need to find the smallest value of k in terms of n at which the loop termination occurs.

The result in previous problem tells us that every 3 iterations, x decreases by half. So, initially we have

$$x_3 \leq \frac{x_0}{2} \tag{1}$$

Then, we can conclude that at $3k$ iterations,

$$x_3 \leq \frac{x_0}{2^k} \tag{2}$$

Then, since $x_0 = n$,

$$x_3 \leq \frac{n}{2^k} \tag{3}$$

Using the fact that the loop termination occurs when $x \leq 1$, we can calculate

$$2^k \geq n \tag{4}$$
$$k \geq \log n \tag{5}$$

Then, it follows from above the smallest value of k at which termination occurs is

$$\lceil \log n \rceil \tag{6}$$

Then, we can conclude the running time of algorithm is $\Theta(\log n)$

<u>Correct Solution:</u>

Let $k \in \mathbb{N}$.

We need to find the smallest value of k in terms of n at which the loop termination occurs.

The result in previous problem tells us that every 3 iterations, x decreases at least by half. So, initially we have

$$x_3 \leq \frac{x_0}{2} \tag{1}$$

Then, we can conclude that at $3k$ iterations,

$$x_{3k} \leq \frac{x_0}{2^k} \tag{2}$$

Then, since $x_0 = n$,

$$x_{3k} \leq \frac{n}{2^k} \tag{3}$$

Using the fact that the loop termination occurs when $x \leq 1$, we can calculate

$$2^k \geq n \tag{4}$$
$$k \geq \log n \tag{5}$$

Then, it follows from above the smallest value of k at which termination occurs is

$$\lceil \log n \rceil \tag{6}$$

Then, we can conclude the running time of algorithm is $\Theta(\log n)$