

# CSC373 Worksheet 1

July 16, 2020

Source: [link](#)

1. **CLRS 4.2-2:** Write Pseudocode for Strassen's algorithm
2. **CLRS 4.2-4:** What is the largest  $k$  such that if you can multiply  $3 \times 3$  matrices using  $k$  multiplications (not assuming commutativity of multiplication), then you can multiply  $n \times n$  matrices in time  $O(n^{\lg 7})$ ? What would the running time of this algorithm be?
3. **CLRS 4.2-5:** V. Pan has discovered a way of multiplying  $68 \times 68$  matrices, using 132,464 multiplications, a way of multiplying  $70 \times 70$  matrices using 143,640 multiplications, and a way of multiplying  $72 \times 72$  matrices using 155,424 multiplications. Which method yields the best asymptotic running time when used in a divide-and-conquer matrix-multiplication algorithm? How does it compare to Strassen's algorithm?
4. **CLRS 4.2-7:** Show how to multiply the complex numbers  $a + bi$  and  $c + di$  using only three multiplications of real numbers. The algorithm should take  $a, b, c$  and  $d$  as input and produce the real component  $ac - bd$  and the imaginary component  $ad + bc$  separately.
5. **CLRS 4-1:** Give asymptotic upper and lower bounds for  $T(n)$  in each of the following recurrences. Assume that  $T(n)$  is constant for  $n \leq 2$ . Make your bounds as tight as possible, and justify your answers

a)  $T(n) = 2T(n/2) + n^4$

b)  $T(n) = T(7n/10) + n$

c)  $T(n) = 16T(n/4) + n^2$

d)  $T(n) = 7T(n/3) + n^2$

e)  $T(n) = 7T(n/2) + n^2$

f)  $T(n) = 2T(n/4) + \sqrt{n}$

g)  $T(n) = 2T(n - 2) + n^2$

6. **CLRS 33.4-2:** Show that it actually suffices to check only the points in the 5 array positions following each point in the array  $Y'$

7. **CLRS 33.4-4:** Give two points  $p_1$  and  $p_2$  in the plane, the  $L_\infty$ -distance between them is given by  $\max(|x_1 - x_2|, |y_1 - y_2|)$ . Modify the closest-pair algorithm to use the  $L_\infty$ -distance.
8. **CLRS 33.4-6:** Augment a change to the closest-pair algorithm that avoids presorting the  $Y$  array but leaves the running time as  $O(n \lg n)$ . (*Hint:* Merge sorted arrays  $Y_L$  and  $Y_R$  to form the sorted array  $Y$ ).