

CSC343 Worksheet 13 Solution

July 4, 2020

1. a)

A	B	C	D	E
a	b	c	d_1	e_1
a_1	b	c	d_1	e_2
a	b_1	c	d_1	e

Step 1 ($B \rightarrow E$):

A	B	C	D	E
a	b	c	d_1	e_1
a_1	b	c	d_1	e_1
a	b_1	c	d_1	e

Step 2 ($CE \rightarrow A$):

A	B	C	D	E
a	b	c	d_1	e_1
a	b	c	d_1	e_1
a	b_1	c	d_1	e

So in this case, an example of an instance of R that is not lossless is:

Title	Studio Name	President	Year	President Address
Toy Story	Pixar	Steve Jobs	2000	123 ABC Street
Star Wars	Fox	Lachlan Murdoch	1977	Hollywood
Return of the Jedi	Fox	Lachlan Murdoch	1983	Hollywood

- $S_1 = \{A, B, C\}$

Title	Studio Name	President
Toy Story	Pixar	Steve Jobs
Star Wars	Fox	Lachlan Murdoch
Return of the Jedi	Fox	Lachlan Murdoch

- $S_2 = \{C, D, E\}$

President	Year	President Address
Steve Jobs	2000	123 ABC Street
Lachlan Murdoch	1977	Hollywood
Lachlan Murdoch	1983	Hollywood

- $S_3 = \{C, E, A\}$

Title	President	President Address
Toy Story	Steve Jobs	123 ABC Street
Star Wars	Lachlan Murdoch	Hollywood
Return of the Jedi	Lachlan Murdoch	Hollywood

-

Notes:

- Decomposition: The good bad and ugly
 - 1) **Elimination of Anomalies** by decomposition as in Section 3
 - 2) **Recoverability of Information** Can we recover the original relation from the tuples in its decomposition?
 - 3) **Preservation of Dependencies (lossless join):** Can we be sure that after reconstructing the original relation from the decompositions, the original FD's satisfy?

BCNF: \rightarrow satisfies 1) and 2) **Not good. NONO**

- The Chase Test for Lossless Join
 - Tests whether the decomposition is lossless

Input:

- A relation R
- A decomposition of R
- A set of functional dependencies

Output:

- Whether the decomposition is lossless or not
- $\Pi_{S_1}(R) \bowtie \Pi_{S_2}(R) \bowtie \cdots \bowtie \Pi_{S_i}(R) = R$

Three things to remember:

1. The natural join is associative and commutative
2. Any tuple t in R is surely in $\pi_{S_1}(R) \bowtie \pi_{S_2}(R) \bowtie \cdots \bowtie \pi_{S_k}(R)$.
3. We have to check to see any tuple in the $\pi_{S_1}(R) \bowtie \pi_{S_2}(R) \bowtie \cdots \bowtie \pi_{S_k}(R)$.

Example:

$$S_1 = \{A, D\}, S_2 = \{B, C\}, S_3 = \{A, C\}$$

$$A \rightarrow B, B \rightarrow C, CD \rightarrow A$$

a_i represents arbitrary value

A	B	C	D
a	b ₁	c ₁	d
a	b ₂	c	d ₂
a ₃	b	c	d

← Represents S₁ = {A,D}
 ← Represents S₂ = {A,C}
 ← Represents S₃ = {B,C}

Step 1: $A \rightarrow B$

Set the value b with the same value of a to be the same. (e.g. $b_2 \rightarrow b_1$)

1. The value of a is the same
 2. Change the value of b_2 to b_1

A	B	C	D
a	b ₁	c ₁	d
a	b ₁	c	d ₂
a ₃	b	c	d

Step 2: $B \rightarrow C$

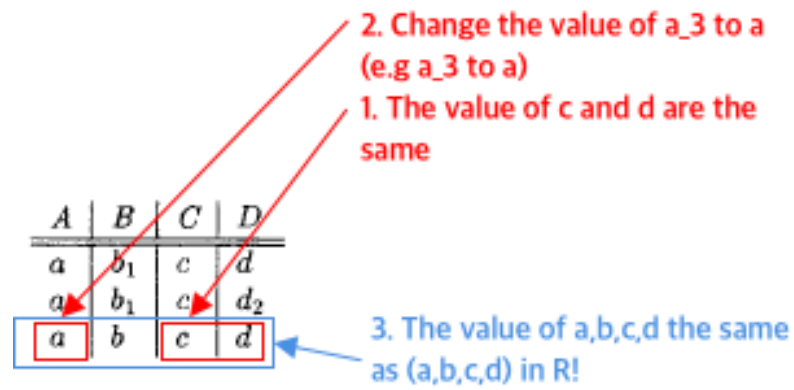
Set the value c with the same value of b to be the same. (e.g. $b_2 \rightarrow b_1$)

1. The value of b is the same
 2. Change the value of c_1 to c

A	B	C	D
a	b ₁	c	d
a	b ₁	c	d ₂
a ₃	b	c	d

Step 3: $CD \rightarrow A$

Set the value a with the same value of c and d to be the same. (e.g. $a_3 \rightarrow a$)



So, we can conclude the join is lossless.