

CSC343 Worksheet 4 Solution

June 18, 2020

1. a) $[(1, 0, 1), (5, 4, 9), (1, 0, 1), (6, 4, 16), (7, 9, 16)]$
- b) $[(1, 0), (3, 3), (3, 4), (4, 3), (1, 1), (4, 3)]$
- c) $[(0, 1), (0, 1), (2, 3), (2, 4), (3, 4)]$

Notes:

- $\tau_L(R)$ sorts tuples in order indicated by L .
 - e.g.

$\tau_{C,B}(R)$ in $R(A, B, C)$ orders the tuples of R by their values of C , and tuples with the same C -value are ordered by their B value.

- d) $[(0, 1), (0, 2), (2, 4), (2, 5), (3, 4), (3, 4)]$
- e) $[(0, 1), (2, 4), (2, 5), (3, 4), (0, 2)]$

Notes:

- $\delta(R)$ converts a bag into a set
 - e.g.

Let $R = [(1, 2), (3, 4), (1, 2), (1, 2)]$

$\delta(R(A, B)) = [(1, 2), (3, 4)]$

- f) $[(0, 2), (2, 7), (3, 4)]$

Notes:

- $\gamma_L(R)$ is an operator that groups a relation and/or aggregate some columns.
 - L in $\gamma_L(R)$ is either
 1. **Grouping attribute** or an attribute by which R will be grouped.

2. **Aggregated attribute** or an attribute where an aggregation operator is applied to.

Example:

$\gamma_{starName, MIN(year) \rightarrow minYear, COUNT(title) \rightarrow ctTitle} (StarsIn)$

studioName
Disney
Disney
Disney
MGM
MGM

groups by studioName

Figure 5.4: A relation with imaginary division into groups

- g) $[(0, 1.5), (2, 4.5), (3, 4)]$
 h) $[(0, 1), (0, 1), (2, 3), (2, 4), (3, 4)]$
 i) $\gamma_{A, MAX(C)}([(2, 3, 4), (2, 3, 4)]) \rightarrow [(2, 4)]$
 j) $[(0, 1, \perp), (2, 3, 4), (2, 3, 4), (0, 1, \perp), (2, 4, \perp), (3, 4, \perp)]$

Notes:

- \bowtie is an outerjoin operator
 - \bowtie_L means Natural Left Outer Join
 - \bowtie_R means Natural Right Outer Join
 - \bowtie means Natural Full Outer Join
 - \perp means null
- e.g. $U \bowtie V$

<i>A</i>	<i>B</i>	<i>C</i>
1	2	3
4	5	6
7	8	9

(a) Relation U

<i>B</i>	<i>C</i>	<i>D</i>
2	3	10
2	3	11
6	7	12

(b) Relation V

<i>A</i>	<i>B</i>	<i>C</i>	<i>D</i>
1	2	3	10
1	2	3	11
4	5	6	\perp
7	8	9	\perp
\perp	6	7	12

(c) Result $U \bowtie V$

- k) $[(\perp, 0, 1), (\perp, 2, 4), (\perp, 2, 5), (2, 3, 4), (\perp, 0, 2), (2, 3, 4)]$
 l) $[(0, 1, \perp), (2, 3, 4), (2, 3, 4), (0, 1, \perp), (2, 4, \perp), (3, 4, \perp),$
 $(\perp, 0, 1), (\perp, 2, 4), (\perp, 2, 5), (2, 3, 4), (\perp, 0, 2), (2, 3, 4)]$
 m) $(0, 1) : \{(2, 4), (2, 5), (3, 4), (3, 4)\}$

But, $\{(2, 3), (2, 4), (3, 4)\}$ from R and $\{(0, 1), (0, 2)\}$ in S don't match. So,

$[(0, 1, 2, 4), (0, 1, 2, 5), (0, 1, 3, 4), (0, 1, 3, 4), (0, 1, 2, 4), (0, 1, 2, 5), (0, 1, 3, 4), (0, 1, 3, 4),$
 $(2, 3, \perp, \perp), (2, 4, \perp, \perp), (3, 4, \perp, \perp), (\perp, \perp, 0, 1), (\perp, \perp, 0, 2)]$

Notes:

- $R \bowtie_C S$ is equivalent form of $\sigma_C(R \times S)$ but instead of filtering, the unmatched tuples are filled with null.

2. a) SELECT model FROM PC WHERE speed > 3.0;
 b) SELECT DISTINCT maker FROM Products NATURAL JOIN Laptops WHERE hd >= 100;

```

c)      SELECT model, price FROM (
2         (SELECT model, price FROM PC NATURAL JOIN Products)
3         UNION
4         (SELECT model, price FROM Laptop NATURAL JOIN Products)
5         UNION
6         (SELECT model, price FROM Printer INNER JOIN Products ON
Printer.model = Product.model)
7     );
8

```

d) SELECT model FROM Printer WHERE color;

```
e) (SELECT DISTINCT makers FROM Products WHERE type='laptops') -
2  (SELECT DISTINCT makers FROM Products WHERE type='pc');
3
```

```
f) SELECT hd FROM PC WHERE EXISTS (
2     SELECT hd, COUNT(model) FROM PC GROUP BY hd
3     HAVING COUNT(model) > 2
4 );
5
```

3. a) SELECT class, country FROM classes WHERE bore >= 16;

b) SELECT * FROM Ships WHERE launched < 1921;

c) SELECT * FROM Outcomes WHERE result='sunk';

d) SELECT * FROM Classes NATURAL JOIN Ships WHERE displacement > 35000;

```
e) SELECT name, displacement, numGuns FROM Classes NATURAL JOIN (
2     SELECT * FROM Ships INNER JOIN Outcomes ON Ships.name =
Outcome.ship
3 );
4
```

```
f) (SELECT name FROM Ships)
2 UNION
3 (SELECT ship AS name FROM Outcomes);
4
```

```
g) SELECT class, COUNT(class) FROM Ships
2 GROUP BY Class
3 HAVING COUNT(class) = 1;
4
```

```
h) (SELECT countries FROM Classes WHERE type='bb')
2 INTERSECT
3 (SELECT countries FROM Classes WHERE type='bc');
4
```

i) Current attempt:

```
1 (SELECT Table1.name FROM Outcomes AS Table1 INNER JOIN Ships ON
Outcome.ship = Ships.name)
2
```

Took too much time. Omitted for now.

4. a) SELECT AVG(speed) FROM PC;

b) SELECT AVG(speed) FROM Laptop HAVING price > 1000;

c) SELECT AVG(price) FROM PC NATURAL JOIN Product HAVING maker = 'A';

d)

```

1  SELECT AVG(price) FROM (
2      (SELECT model, price FROM PC NATURAL JOIN Product WHERE maker
3      = 'D')
4      UNION
5      (SELECT model, price FROM Laptop NATURAL JOIN Product WHERE
6      maker = 'D')
7  );

```

e) SELECT speed, AVG(price) FROM PC GROUP BY speed;

f) SELECT maker, AVG(screen) FROM Laptop NATURAL JOIN Product GROUP BY maker;

g)

```

1  SELECT maker, COUNT(model) FROM Products GROUP BY maker HAVING
2  COUNT(model) >= 3;
3

```

h)

```

1  SELECT maker, MAX(price) FROM PC
2  NATURAL JOIN Products GROUP BY maker;
3

```

i)

```

1  SELECT speed, AVG(price) FROM PC GROUP BY speed
2  HAVING speed > 2.0;
3

```

5. a) SELECT COUNT(class) FROM Classes;

Correct Solution:

```
SELECT COUNT(class) FROM Classes HAVING type='bb';
```

b) SELECT class, AVG(numGuns) FROM Classes GROUP BY class;

Correct Solution:

```
SELECT class, AVG(numGuns) FROM Classes GROUP BY class
HAVING type='bb';
```

c) SELECT AVG(numGuns) FROM Classes HAVING type='bb';

d) SELECT class, MIN(launched) FROM Ships GROUP BY class;

e)

```

1  SELECT class, COUNT(Ships.name) FROM Ships INNER JOIN Outcomes
2  GROUP BY class ON Ships.name = Outcomes.ship
3  HAVING result='sunk';
4

```

6.

```

1  SELECT name, MIN(movieYear) FROM MovieStar INNER JOIN StarsIn ON
2  StarsIn.starName = MovieStar.name GROUP BY name HAVING COUNT(name)
3  >= 3;

```

7. Yes. It is possible

- Rename aggregate columns using ρ
- Use σ around $\rho(\gamma(\dots))$

8. a) INSERT INTO PC(model, speed, ram, hd, price) VALUES (1100, 3.2, 1024, 180, 2499);

INSERT INTO Product(maker, model, type) VALUES ('C', 1100, 'pc');

Notes:

- Insertion
 - **Syntax:** INSERT INTO R(A_1, A_2, \dots, A_n) VALUES (v_1, \dots, v_n)
 - Inserts a tuple with values into relation R
 - Values can be from select statement

```

1  \ Example 1
2  INSERT INTO StarsIn(movieTitle, movieYear, starName)
3  VALUES('The Maltese Fiction', 1942, 'Sydney Greenstreet');
4
5  \ Example 2
6  INSERT INTO Studio(name)
7  SELECT DISTINCT studioName
8  FROM Movies
9  WHERE studioName NOT IN
10     (SELECT name FROM Studio);
11

```

```

b) // SQL statement 1
2  INSERT INTO Laptop(model, speed, ram, hd, screen, price)
3  SELECT DISTINCT model + 110, speed, ram, hd, 17 AS screen,
4  price + 500
5  FROM PC;
6
7  // SQL statement 2
8  INSERT INTO Product(maker, model, type)
9  SELECT DISTINCT maker, model + 110, 'laptop' AS type
10 FROM Product WHERE type='pc';

```

```

c) DELETE FROM PC
2  WHERE hd < 100;
3

```

Notes:

- Deletion

- **Syntax:** DELETE FROM R WHERE $\langle \text{Condition} \rangle$;

Example:

```
1  DELETE FROM StarsIn
2  WHERE movieTitle = 'The Maltest Falcon' AND
3         movieYeeear = 1942 AND
4         starName = 'Sydney Greenstreet';
5
```

d)

```
DELETE FROM Product
WHERE maker = ALL(
    (SELECT maker FROM Product WHERE type='laptop')-
    (
        (SELECT maker FROM PRODUCT WHERE type='laptop')
        INTERSECT
        (SELECT maker FROM PRODUCT WHERE type='printer')
    )
)
```

e)

```
UPDATE Product
SET maker = 'A'
WHERE maker = 'B';
```

Notes:

- UPDATE
 - **Syntax:** UPDATE R SET $\langle \text{new-value assignments} \rangle$ WHERE $\langle \text{condition} \rangle$;
 - updates tuples in database

Example:

```
1  UPDATE MovieExec
2  SET name = 'Pres.' || name
3  WHERE cert# IN (SELECT presC# FROM Studio);
4
```