

# CSC236 Worksheet 4 Solution

Hyungmo Gu

May 7, 2020

## Question 1

- Let  $n \in \mathbb{N}$  and assume that  $\exists k \in \mathbb{N}^+$ ,  $n = 3^k$ , so  $k = \log_3 n$ .

Then, since  $n = 3^k$  and  $3 \mid n$ , we have  $\lceil n/3 \rceil = n/3$ .

Then,

$$T(n) = 2n + T(n/3) \quad [\text{By def.}] \quad (1)$$

$$= 2(n/3) + (2(n/3) + T(n/3^2)) \quad [\text{By substituting } n/3 \text{ for } n \text{ in def.}] \quad (2)$$

$$= 2^2(n/3) + T(n/3^2) \quad (3)$$

$$= 2^3(n/3^2) + T(n/3^3) \quad [\text{By substituting } n/3 \text{ for } n \text{ in def.}] \quad (4)$$

$$\vdots \quad (5)$$

$$= 2^k(n/3^{k-1}) + T(n/3^k) \quad [\text{After } k \text{ applications}] \quad (6)$$

$$= 2^{\log_3 n}(n/3^{\log_3 n - 1}) + T(n/3^{\log_3 n}) \quad [\text{By replacing } k = \log_3 n] \quad (7)$$

$$= 2^{\log_3 n}(n(3)/n) + T(n/n) \quad (8)$$

$$= 3 \cdot 2^{\log_3 n} + T(1) \quad (9)$$

$$= 3 \cdot 2^{\log_3 n} + 2 \quad (10)$$

### Correct Solution:

Let  $n \in \mathbb{N}$  and assume that  $\exists k \in \mathbb{N}^+$ ,  $n = 3^k$ , so  $k = \log_3 n$ .

Then, since  $n = 3^k$  and  $3 \mid n$ , we have  $\lceil n/3 \rceil = n/3$ .

Then,

$$T(n) = 2n + T(n/3) \quad \text{[By def.]} \quad (11)$$

$$= 2n + 2(n/3) + T(n/3^2) \quad \text{[By substituting } n/3 \text{ for } n \text{ in def.]} \quad (12)$$

$$= 2n + 2(n/3) + 2(n/3^2) + T(n/3^3) \quad \text{[By substituting } n/3 \text{ for } n \text{ in def.]} \quad (13)$$

$$\vdots \quad (14)$$

$$= 2 \sum_{i=0}^{k-1} n/3^i + T(n/3^k) \quad (15)$$

$$= 2 \cdot 3^k \left( \frac{1 - (1/3)^k}{1 - 1/3} \right) + T(n/3^k) \quad \text{[By using geometric series]} \quad (16)$$

$$= 2 \cdot 3^k \cdot 3/2 \left( 1 - (1/3)^k \right) + T(n/n) \quad (17)$$

$$= 3(3^k - 1) + T(1) \quad (18)$$

$$= 3^{k+1} - 1 \quad (19)$$

### Notes:

#### • Repeated Substitution:

- Is a technique used to find a closed form formula
- **closed form formula** is a simple formula that allows evaluation of  $T(n)$  without the need to evaluate, say  $T(n/2)$

i.e. from

$$T(n) = \begin{cases} c & \text{if } n = 1 \\ 2T(n/2) + dn & \text{if } n > 1 \end{cases} \quad (20)$$

to

$$T(n) = cn + dn \log_2 n$$

### Example:

Consider the recurrence

$$T(n) = \begin{cases} c & \text{if } n = 1 \\ 2T(n/2) + dn & \text{if } n > 1 \end{cases} \quad (1)$$

Find closed form formula for  $T(n)$ , where  $n$  is an arbitrary power of 2. That is

$\exists k \in \mathbb{N}, n = 2^k$ .

Let  $n \in \mathbb{N}$  and assume that  $\exists k \in \mathbb{N}^+, n = 2^k$ , so  $k = \log_2 n$ .

Then,

$$T(n) = 2T(n/2) + dn \quad [\text{By 1}] \quad (2)$$

$$= 2\left(2T(n/2^2) + dn/2\right) + dn \quad [\text{By substituting } n/2 \text{ for } n \text{ in 1}] \quad (3)$$

$$= 2^2T(n/2^2) + 2dn \quad (4)$$

$$= 2^2\left(2T(n/2^3) + dn/2^2\right) + 2dn \quad [\text{By substituting } n/2^2 \text{ for } n \text{ in 1}] \quad (5)$$

$$= 2^3T(n/2^3) + 3dn \quad [\text{By substituting } n/2^2 \text{ for } n \text{ in 1}] \quad (6)$$

$$\vdots \quad (7)$$

$$= 2^kT(n/2^k) + kdn \quad [\text{After } k \text{ applications}] \quad (8)$$

$$= 2^{\log_2 n}T(n/2^{\log_2 n}) + (\log_2 n)dn \quad [\text{By replacing } k = \log_2 n] \quad (9)$$

$$= nT(1) + (\log_2 n)dn \quad (10)$$

$$= cn + (\log_2 n)dn \quad (11)$$

## Question 2