

CSC373 Worksheet 4 Solution

August 6, 2020

1. • Calculating out-degree

Let $G = (V, E)$ be a directed graph. Let $[v_1, \dots, v_n]$ be a list of vertices in graph G .

I need to calculate the outdegree of every vertex using adjacency list.

We know that in addition to counting each v_i in adjacency list where $i = 1, \dots, n$, we are also counting $|Adj[v_i]|$ edges.

Since there are $|V| = n$ many vertices, we can write that the total count is $|V| + \sum_{i=1}^n |Adj[v_i]| = |V| + |E|$, which is $\mathcal{O}(|V| + |E|)$.

• Calculating In-degree

The outdegree of a vertex is indegree of another vertex.

Using this fact, we can conclude the running time of computing indegree of every vertex is $\mathcal{O}(|V| + |E|)$.

Notes:

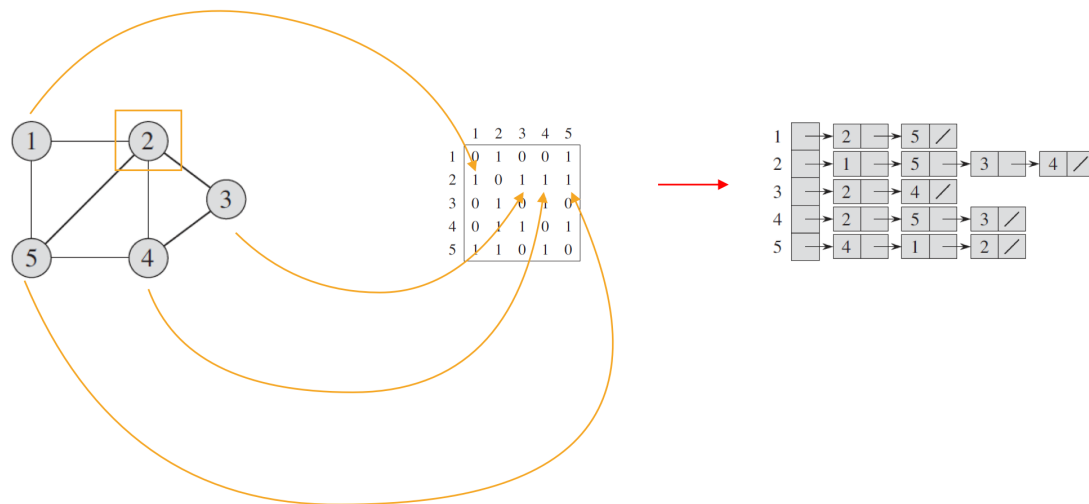
• Vertex

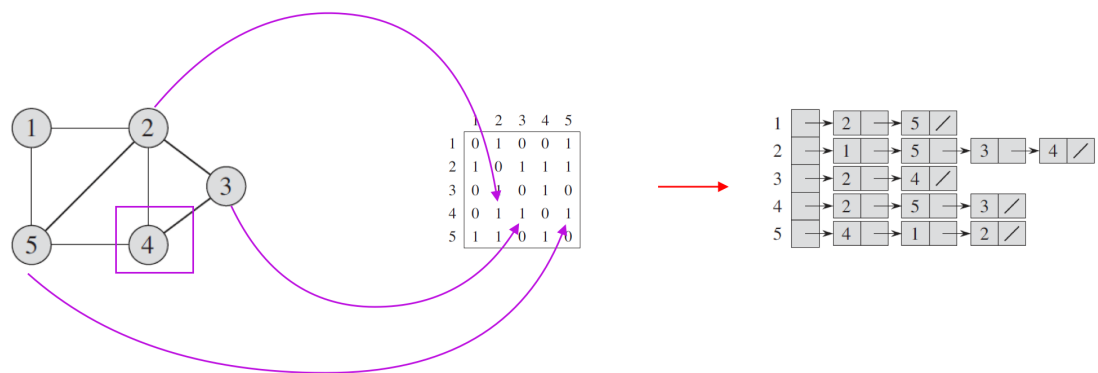
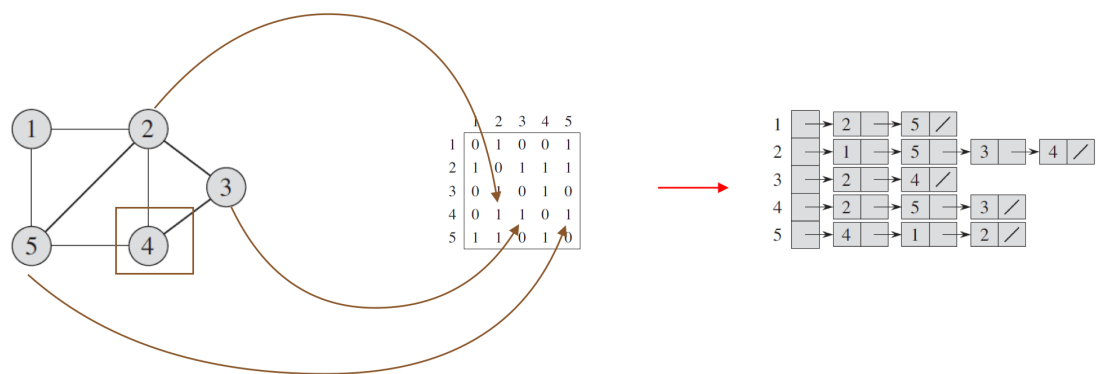
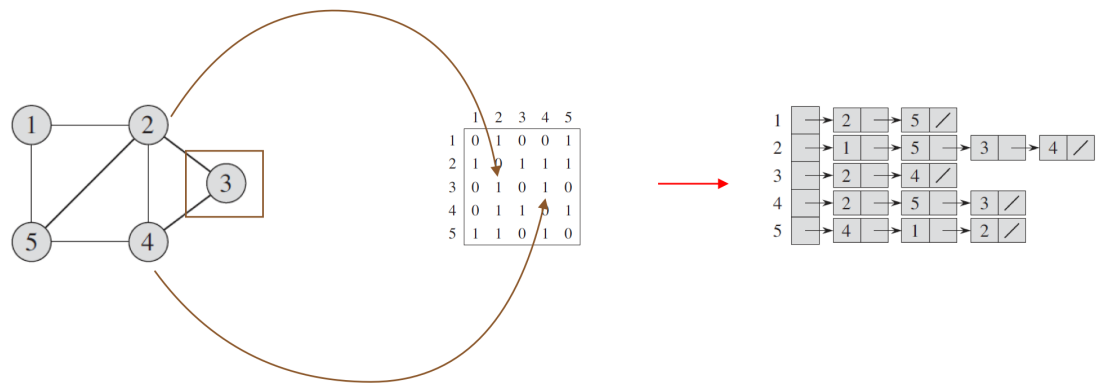
- Is a fundamental unit of which graphs are formed
- Also means node

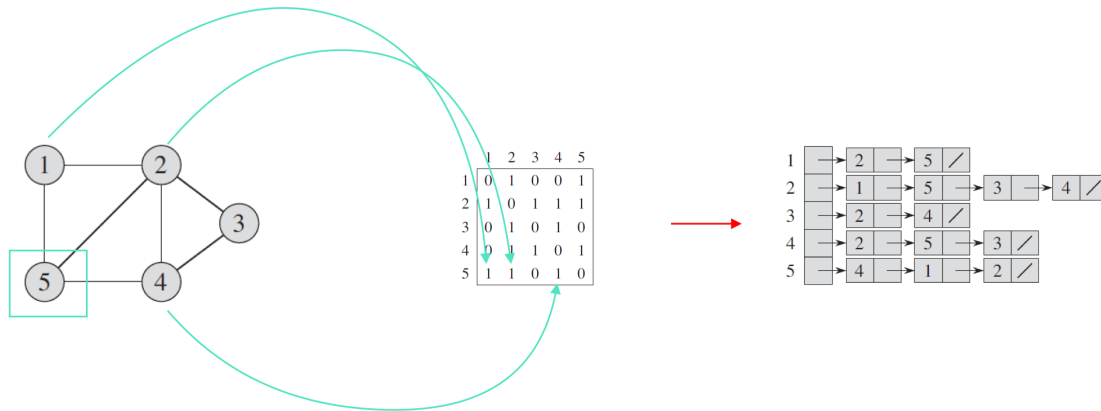


• Adjacency-list Representation

- Associates each vertex in a graph with the collection of its neighbouring vertices or edges
- Is represented by $Adj[v]$
 - * Means all vertices that are neighbour to vertex v
 - * In a directed graph, $Adj[v]$ are all out-degree vertices of vertex v
 - * $|Adj[v]|$ means the total number of outdegree of vertex v







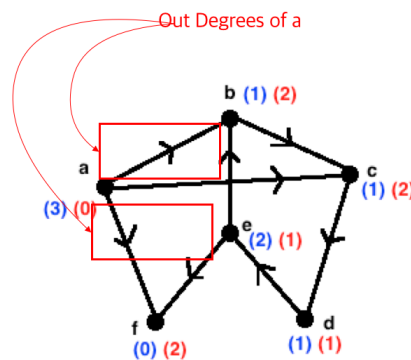
• Directed graph

- Is a graph that is made up of a set of vertices connected by edges, where the edges have a direction associated with them



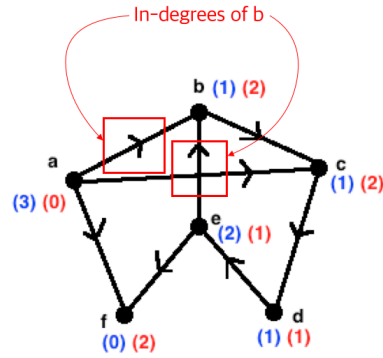
• Out-degrees

- For a directed graph $G = (V(G), E(G))$ and a vertex $x_1 \in V(G)$, the Out-Degree of x_1 refers to the number of arcs incident from x_1 . That is, the number of arcs directed away from the vertex x_1 .

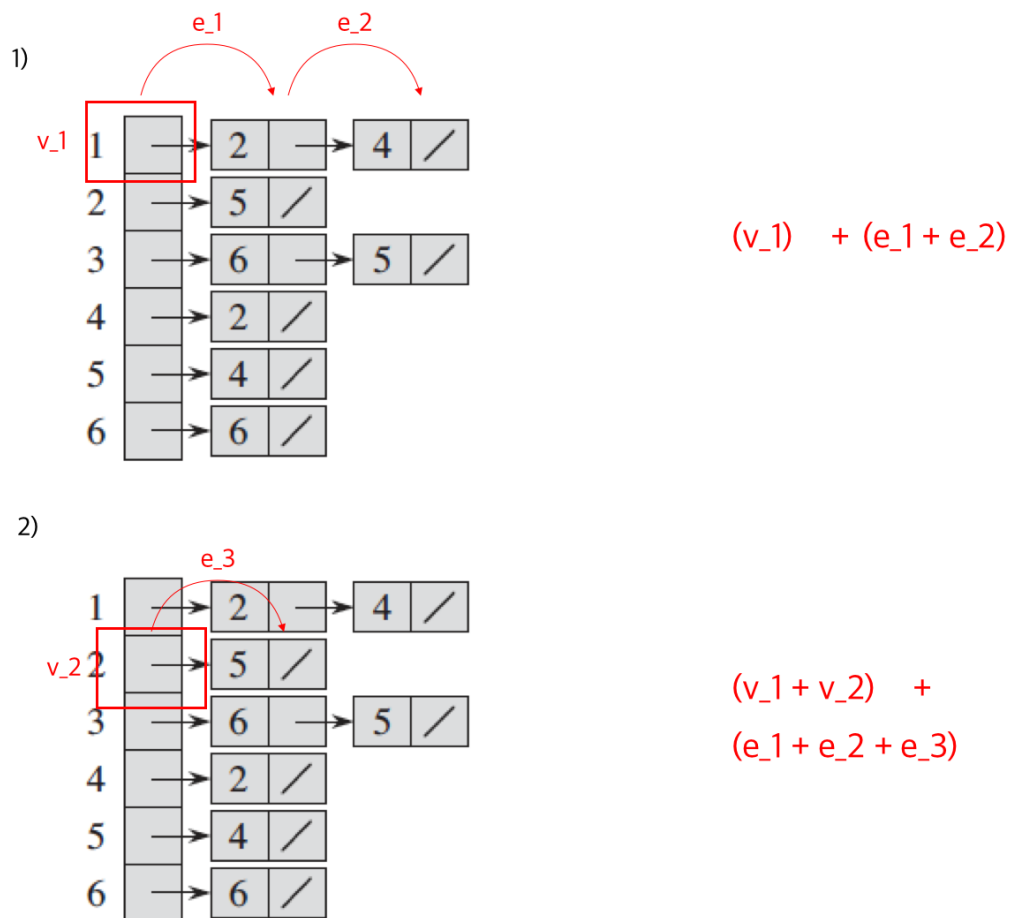


• In-degrees

- For a directed graph $G = (V(G), E(G))$ and a vertex $x_1 \in V(G)$, the In-Degree of x_1 refers to the number of arcs incident to x_1 . That is, the number of arcs directed towards the vertex x_1 .



- Computing the outdegree of every vertex using adjacency list



3)



$$(v_1 + v_2 + v_3) + (e_1 + e_2 + e_3 + e_4 + e_5)$$

3)



$$(v_1 + v_2 + v_3) + (e_1 + e_2 + e_3 + e_4 + e_5)$$

4)



$$(v_1 + v_2 + v_3 + v_4) + (e_1 + e_2 + e_3 + e_4 + e_5)$$

4)



$$(v_1 + v_2 + v_3 + v_4) + (e_1 + e_2 + e_3 + e_4 + e_5)$$

6)



$$(v_1 + v_2 + v_3 + v_4 + v_5 + v_6) + (e_1 + e_2 + e_3 + e_4 + e_5 + e_6 + e_7 + e_8)$$

So it has $\mathcal{O}(V + E)$

- Computing the outdegree of every vertex using adjacency list

The outdegree of a vertex is indegree of another vertex.

Using this fact, we can conclude the running time of computing indegree of every vertex is $\mathcal{O}(V + E)$.

- Computing G^T from G in Adjacency List



```

1  COMPUTE-G-TRANSPOSE-ADJ-MATRIX(Adj, V)
2      Let Adj' be a new adjacency list containing keys  $v_1 \dots v_n$ 
3
4      for i = 1 to |V|
5          for every vertex w in Adj[ $v_i$ ]
6              Insert(Adj'[w],  $v_i$ )
7
8      return Adj'
9

```

- Computing G^T from G in Adjacency-Matrix




```

1  COMPUTE-G-TRANSPOSE-ADJ-MATRIX(A,V)
2      Let A'[1..|V|, 1..|V|] be a new adjacency matrix
3
4      for i = 1 to |V|
5          for j = 1 to |V|
6              A'[j,i] = A[i,j]
7
8      return A'
9

```

Correct Solution:

- Computing G^T from G in Adjacency List



```

1  COMPUTE-G-TRANSPOSE-ADJ-MATRIX(Adj,V)
2      Let Adj' be a new adjacency list containing keys
3      v1...vn
4
5      for i = 1 to |V|
6          for every vertex w in Adj[vi]
7              Insert(Adj'[w], vi)
8
9      return Adj'

```

The running time is $\mathcal{O}(|V| + |E|)$

- Computing G^T from G in Adjacency-Matrix



```

1  COMPUTE-G-TRANSPOSE-ADJ-MATRIX(A,V)
2      Let A'[1..|V|, 1..|V|] be a new adjacency matrix
3
4      for i = 1 to |V|
5          for j = 1 to |V|
6              A'[j,i] = A[i,j]
7
8      return A'
9

```

The running time is $\mathcal{O}(|V|^2)$

```

31 Breadth-First-Search(V, v_i)
32     d = 0
33     for each v_i ∈ V
34         while performing BFS(V, v_i)
35             let w be the current node in BFS
36             if δ(v_i, w) > d
37                 d = δ(v_i, w)
38
39     return d
10

```

Finding Runtime of Algorithm

Since the graph iterates $\sum_{i=1}^n Adj[v_i] = |E|$ times for each $v_i \in V$, the algorithm iterates total of $|V| \cdot |E|$ times, which is $\mathcal{O}(|V||E|)$.

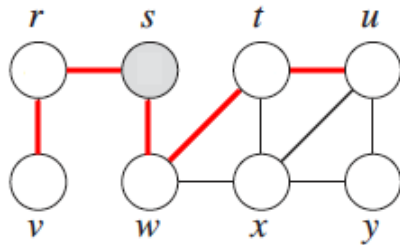
Notes:

- **Breadth First Search**

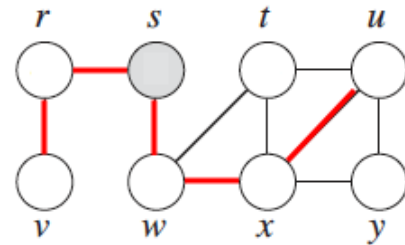
- Is an algorithm for searching or traversing a graph
- Is one of the simplest algorithm

- **Largest of All Shortest Path Distance**

- Means the shortest distance between two furthest apart nodes

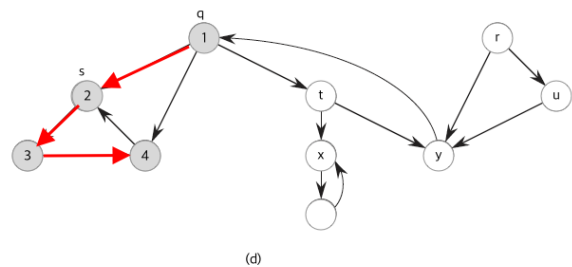
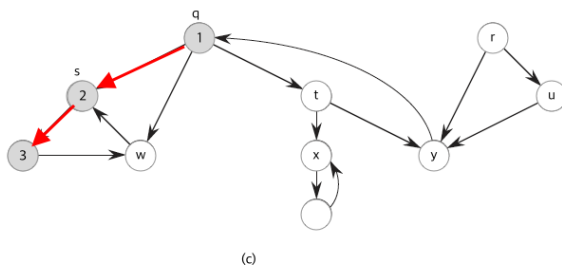
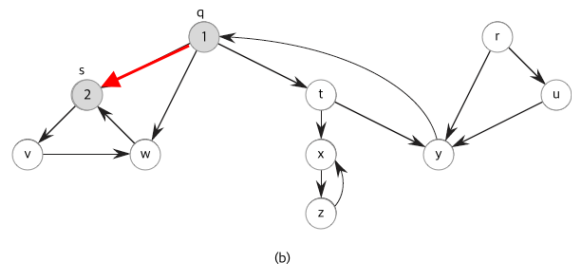


OR



References

1) McGill University, 308-360 Tutorial, link



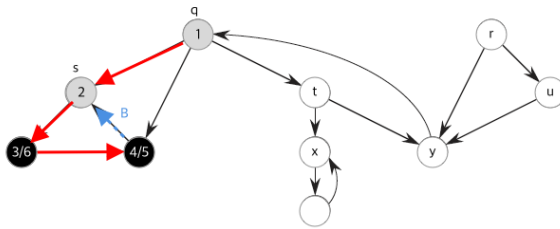
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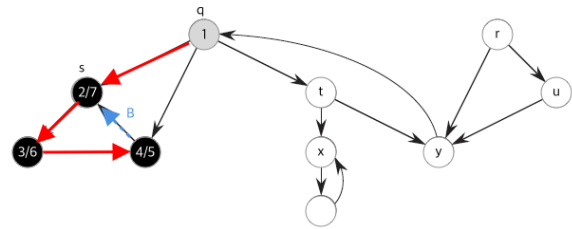
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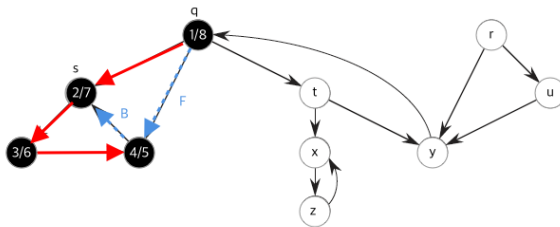
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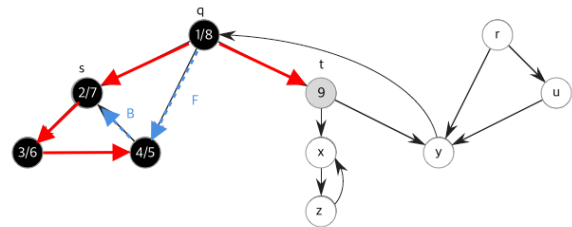
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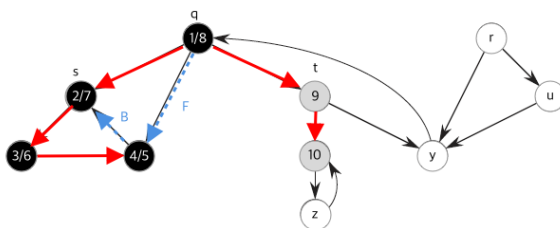
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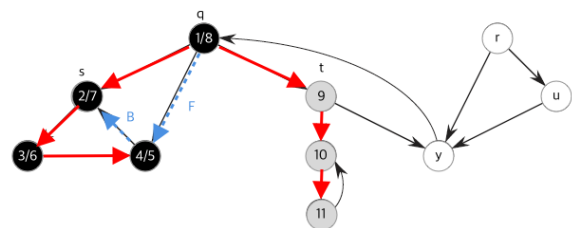
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(l)



(m)



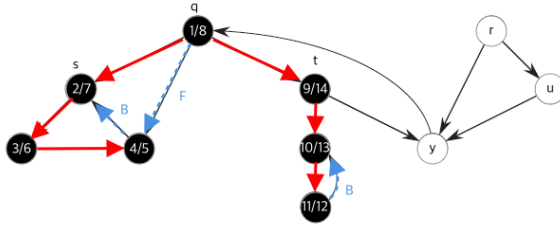
(n)



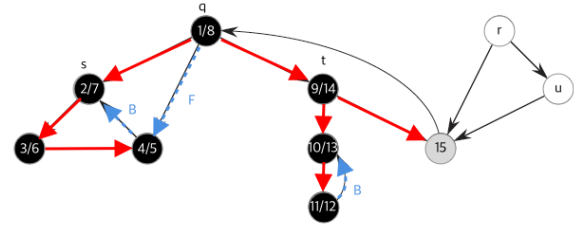
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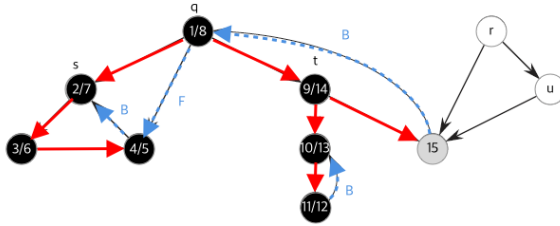
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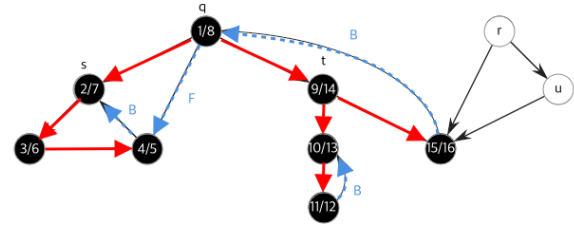
(q)



(r)



(s)



(t)



Notes:

- **Depth First Search**

- Searches deeper in the graph whenever possible

- **Forward Edge**

- Is an edge (u, v) such that v is descendant but not part of the DFS tree. Edge $1 \rightarrow 8$ is a forward edge



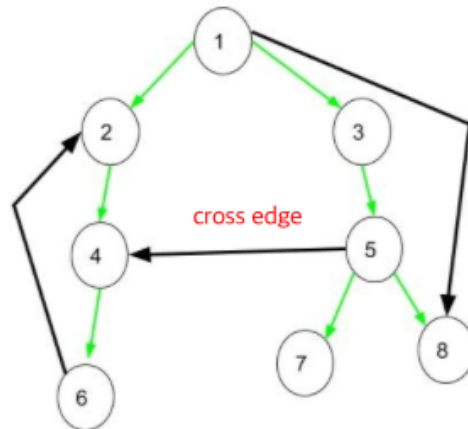
- **Back Edge**

- It is an edge (u, v) such that v is ancestor of edge u but not part of DFS tree. Edge from $6 \rightarrow 2$ is a back edge.
- Indicates a cycle in a graph



- **Cross Edge**

- It is a edge which connects two node such that they do not have any ancestor and a descendant relationship between them. Edge from node $5 \rightarrow 4$ is cross edge.



References

- 1) Geeks For Geeks, Tree, Back, Edge and Cross Edges in DFS of Graph, link
5. *Proof.* Let G be a connected graph. Let A be a subset of E that is always included in some minimum spanning tree for G .

Assume for the sake of contradiction that (u, v) is not contained in some minimum spanning tree of G .

We know from the minimum-spanning tree algorithm that a light edge crossings in a cut of G that respects A is always chosen (since this is a safe edge and the algorithm always picks the safe edge).

Since the edge (u, v) is not a part of minimum spanning tree, (u, v) is not chosen in a cut.

Since (u, v) is not chosen, the edge (u', v') with smaller weight is chosen.

Then this violates the assumption that (u, v) is the edge with the smallest weight.

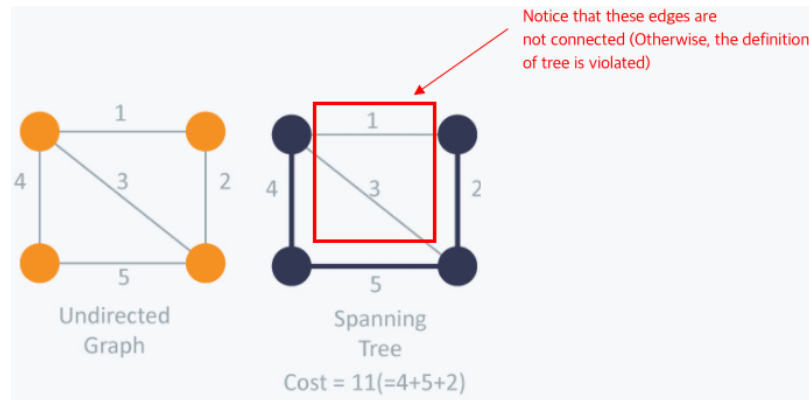
Thus, by contradiction, (u, v) belongs to some minimum spanning tree of G .

□

Notes:

• Spanning Tree

Given an undirected and connected graph $G = (V, E)$, a spanning tree of the graph G is a tree that spans G (that is, it includes every vertex of G) and is a subgraph of G (every edge in the tree belongs to G)



• Minimum Spanning Tree

- Is the spanning tree where the cost is minimum among all the spanning trees.
 - * The cost of the spanning tree is the sum of the weights of all the edges in the tree.
- There can be many minimum spanning trees.



- Is used in
 1. Network design (Telephone, electrical, hydraulic, TV cable, computer, road)
 2. Approximation algorithm for NP-hard problems
 3. Learning salient features for real-time face verification
 4. Reducing data storage in sequencing amino acids in a protein

1)



2)



3)



4)



5)



6)



• Cut

- A cut of an undirected graph $G = (V, E)$ is denoted $(S, V - S)$
- Is a partition of V



• Light Edge Crossing

- An edge is a light edge crossing if its weight is the minimum of any edge crossing the cut



- **Safe Edge**

- Is an edge (u, v) that may be added to A without violating the invariant that $A \cup (u, v)$ is a subset of some minimum spanning tree.

- **Cut respects a set A of edges**

- Means no edges in A crosses the crossing cut



References:

- 1) Princeton University, Minimum Spanning Tree, [link](#)
- 2) McGill University, 308-360 Tutorial, [link](#)
- 3) Hacker Earth, Minimum Spanning Tree, [link](#)

```

6.1  MST-PRIM-ADJACENCY-LIST(Adj, w, r)
      Let Q be an empty list
      Q = Extract node objects from Adj
      for i = 1 to |Q|

```

```

6      u.key =  $\infty$ 
7      u. $\pi$  = NIL
8
9      r.key = 0
10
11     while Q  $\neq$  0
12         u = EXTRACT-MIN(Q)
13         for each v  $\in$  G.Adj[u]
14             if v  $\in$  Q and w(u,v) < v.key
15                 v. $\pi$  = u
16                 v.key = w(u,v)
17

```

Notes:

- Kruskal Algorithm

- is a minimum-spanning-tree algorithm which finds an edge of the least possible weight that connects any two trees in the forest.

- Prim's Algorithm

- Is a greedy algorithm that finds a minimum spanning tree for a weighted undirected graph
- Is very similar to Dijkstra's algorithm for finding shortest path





References:

- 1) Wikipedia, Kruskal Algorithm, [link](#)

```

71  BELLMAN-FORD( $G, w, s$ )
72      INITIALIZE-SINGLE-SOURCE( $G, s$ )
73      active = true
74
75      while(active)
76          for each edge  $(u, v) \in G.E$ 
77              RELAX( $u, v, w$ )
78
79              if  $v.d > u.d + w(u, v)$ 

```

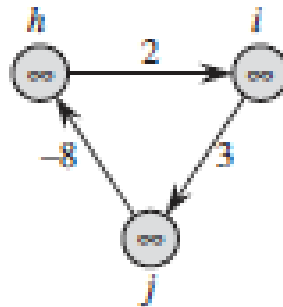
```

10         return False
11
12     return True
13
14
15     RELAX(u,v,w)
16         if v.d > u.d + w(u,v)
17             v.d = u.d + w(u,v)
18             v.π = u
19         else
20             active = false
21

```

Notes:

- m here represents the total number of edges in shortest path between u and v
- **Negative-Weight Cycle**
 - Is a cycle with weights that sum to a negative number



- **Bellman-Ford Algorithm**

- Solves the single-source shortest-paths problem in the general case in which edge weights may be negative.



- Returns TRUE if and only if the graph contains no negative-weight cycles that are reachable from the source

```

81 BELLMAN-FORD(G, w, s)
82     INITIALIZE-SINGLE-SOURCE(G, s)
83     for i = 1 to |G.V| - 1

```

```

4       for each edge  $(u,v) \in G.E$ 
5           RELAX( $u,v,w$ )
6
7       for each edge  $(u,v) \in G.E$ 
8           if  $v.d > u.d + w(u,v)$ 
9               return False
10
11       return True
12
13  RELAX( $u,v,w$ )
14      if  $v.d > u.d + w(u,v)$ 
15           $v.d = u.d + w(u,v)$ 
16           $v.\pi = u$ 
17      else
18           $v.d = -\infty$ 
19

```

9. I need to design algorithm that returns the most reliable path between two vertices.

```

1  DIJKSTRA( $G,w,s$ )
2      INITIALIZE-SINGLE-SOURCE( $G,s$ )
3       $S = \emptyset$ 
4       $Q = G.V$ 
5
6      while  $Q \neq \emptyset$ 
7           $u = \text{EXTRACT-MAX}(Q)$ 
8           $S = S \cup \{u\}$ 
9          for each vertex  $v \in G.\text{Adj}[u]$ 
10              RELAX( $u,v,w$ )
11
12  INITIALIZE-SINGLE-SOURCE( $G,s$ )
13      for each vertex  $v \in G.V$ 
14           $v.d = -\infty$ 
15           $v.\pi = \text{NIL}$ 
16       $s.d = 1$ 
17
18  RELAX( $u,v,w$ )
19      if  $u.d * w(u,v) > v.d$ 
20           $v.d = u.d * w(u,v)$ 
21           $v.\pi = u$ 
22

```

Notes:

- Dijkstra Algorithm
 - Finds shortest path from one node to all nodes
 - Solves the single-source shortest-paths problem on a weighted directed graph $G = (V, E)$

- Cannot have negative weight edges



10. The output of the Floyd-Warshall algorithm can detect the presence of a negative weight cycle by looking for a negative value in the matrix $D^{(n)}$.

Notes:

- Floyd-Warshall Algorithm
 - Is an algorithm for finding shortest paths between all pairs of vertices in a weighted graph.
 - Negative edges allowed
 - Cannot have negative weight cycles
 - Is useful in graphs with dense number of edges

```

FLOYD-WARSHALL( $W$ )
1   $n = W.rows$ 
2   $D^{(0)} = W$ 
3  for  $k = 1$  to  $n$ 
4    let  $D^{(k)} = (d_{ij}^{(k)})$  be a new  $n \times n$  matrix
5    for  $i = 1$  to  $n$ 
6      for  $j = 1$  to  $n$ 
7         $d_{ij}^{(k)} = \min(d_{ij}^{(k-1)}, d_{ik}^{(k-1)} + d_{kj}^{(k-1)})$ 
8  return  $D^{(n)}$ 

```

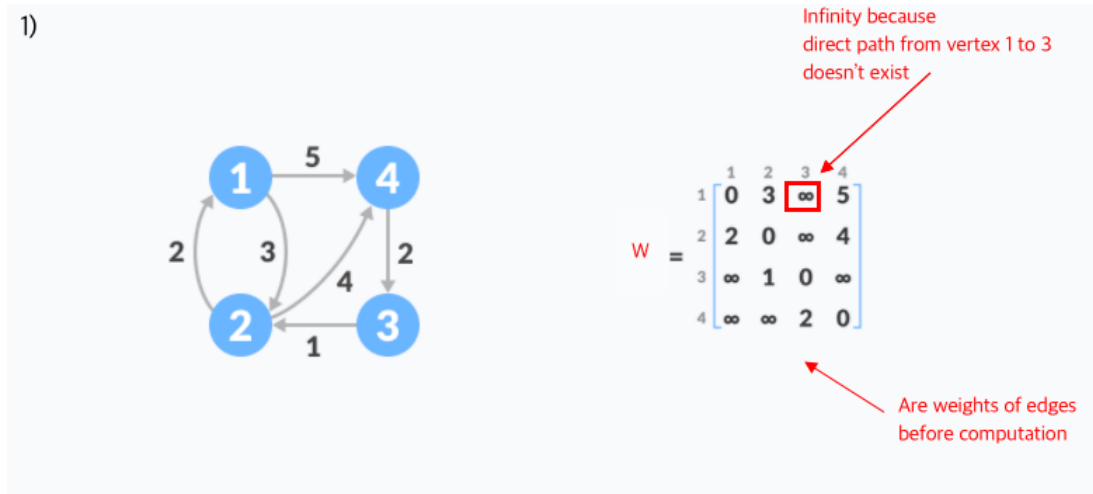
- * n : the number of vertices
- * W : the $n \times n$ matrix of edge weights of an n -vertex directed graph $G = (V, E)$

$$w_{ij} = \begin{cases} 0 & \text{if } i = j \\ \text{the weight of directed edge } (i, j) & \text{if } i \neq j \text{ and } (i, j) \in E \\ \infty & \text{if } i \neq j \text{ and } (i, j) \notin E \end{cases} \quad (1)$$

- * $D^{(k)}$: the $n \times n$ matrix of edge weights of shortest path where all intermediate vertices are in the set $1, 2, \dots, k$.

$$d_{ij} = \begin{cases} w_{ij} & \text{if } k = 0 \\ \min(d_{ij}^{(k-1)}, d_{ik}^{(k-1)} + d_{kj}^{(k-1)}) & \text{if } k \geq 1 \end{cases} \quad (2)$$

Example:



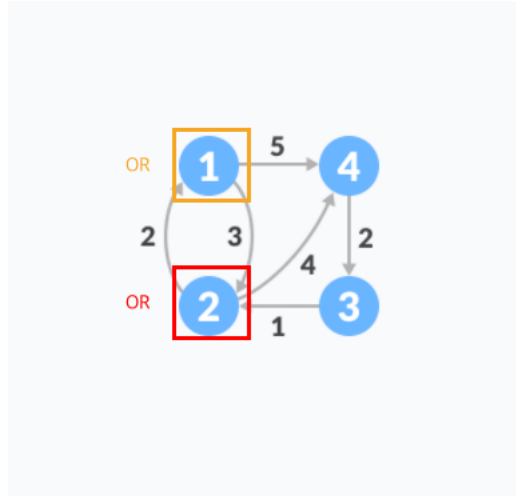
- * $k = 1$



$$D^1 = \begin{bmatrix} 0 & 3 & \infty & 5 \\ 2 & 0 & \infty & 4 \\ \infty & 1 & 0 & \infty \\ \infty & \infty & 2 & 0 \end{bmatrix} \quad (3)$$

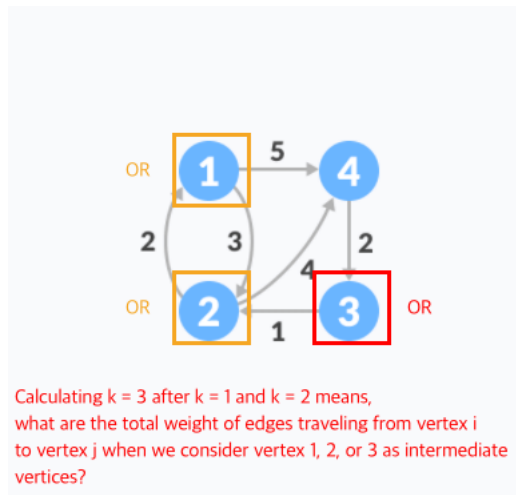
· Matrix calculated via line 3 to 7 in Floyd-Warshall algorithm

- * $k = 2$



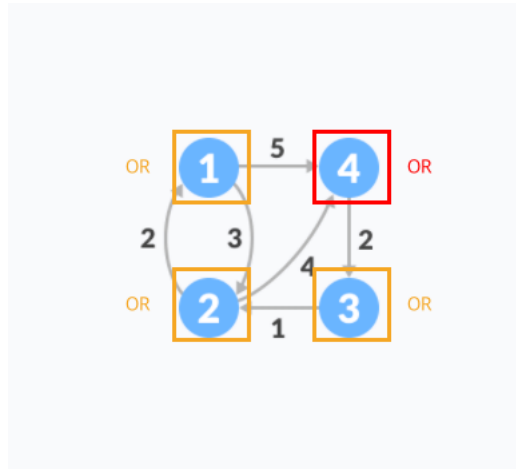
$$D^2 = \begin{bmatrix} 0 & 3 & \infty & 5 \\ 2 & 0 & \infty & 4 \\ \infty & 1 & 0 & 5 \\ \infty & \infty & 2 & 0 \end{bmatrix} \quad (4)$$

- 5 in D_{34}^2 is the only value that's different
- * $k = 3$



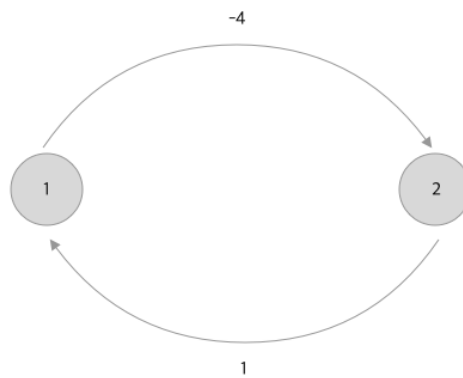
$$D^3 = \begin{bmatrix} 0 & 2 & \infty & 4 \\ 2 & 0 & \infty & 4 \\ \infty & 1 & 0 & 5 \\ 7 & 3 & 2 & 0 \end{bmatrix} \quad (5)$$

- * $k = 4$



$$D^4 = \begin{bmatrix} 0 & 3 & 7 & 5 \\ 2 & 0 & 4 & 4 \\ 12 & 1 & 0 & 5 \\ 7 & 3 & 2 & 0 \end{bmatrix} \quad (6)$$

Example 2:



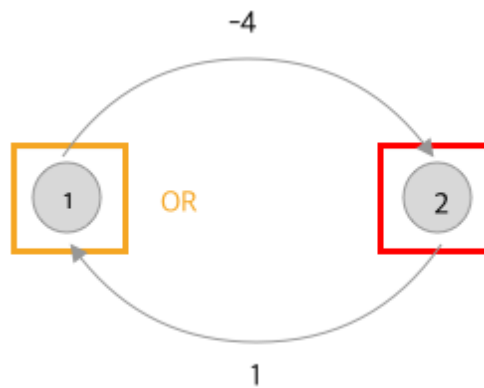
$$W = \begin{bmatrix} 0 & -4 \\ 1 & 0 \end{bmatrix} \quad (7)$$

* $k = 1$



$$W = \begin{bmatrix} 0 & -4 \\ 1 & 0 \end{bmatrix} \quad (8)$$

* $k = 2$



$$W = \begin{bmatrix} 0 & -4 \\ 1 & 0 \end{bmatrix} \quad (9)$$

11. Notes:

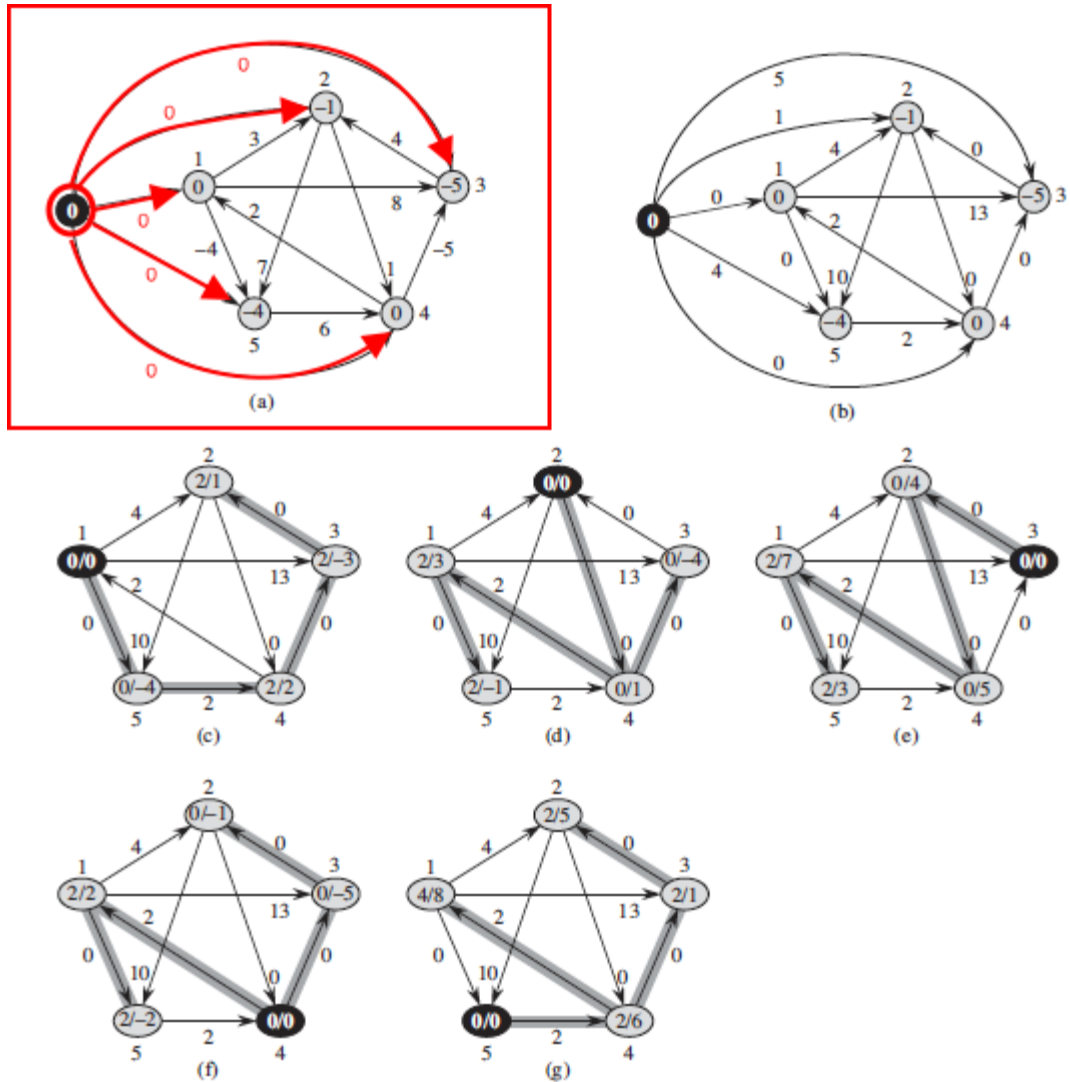
- Johnson's Algorithm

- Is a way to find the shortest paths between all pairs of vertices in an edge-weighted, directed graph
- Allows some edges to be negative-number
- No negative-cycles may exist
- Is similar to Floyd-Warshall Algorithm
- Is most effective in graph with sparse number of edges
- Works as a subroutine to Dijkstra and Bellman-Ford Algorithm

- Has running time of $O(nm + n(m + n \log n)) = O(nm + n \log n)$
- Reweighting
 - Is a technique used in Johnson's Algorithm
 - Is so that all edges have non-negative weights ^[1]

Steps:

1. Create a source vertex $V' = V \cup \{s\}$ for some $s \notin V$ and $E' = E \cup \{(s, v) : v \in V\}$
2. Extend the weight function w so that $w(s, v) = 0$ for all $v \in V$



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References

- 1) Columbia University, Johnson's Algorithm for All-Pairs Shortest Paths, link