Problem Set 4 Solution

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Question 1

a. **Statement:** $\forall f, g : \mathbb{N} \to \mathbb{R}^+,\ b \in \mathbb{R}^+,\ (g(n) \in \Theta(f(n))) \wedge (n_0 \in \mathbb{N},\ n \geq n_0 \Rightarrow f(n) \geq b \wedge g(n) \geq b) \wedge (b > 1) \Rightarrow \log_b(g(n)) \in \Theta(\log_b(f(n)))$

Statement Expanded: $\forall f, g : \mathbb{N} \to \mathbb{R}^+,\ b \in \mathbb{R}^+,\ \Big(\exists c_1, c_2, n_0 \in \mathbb{R}^+,\ \forall n \in \mathbb{N},\ n \geq n_0 \Rightarrow c_1 \cdot g(n) \leq f(n) \leq c_2 \cdot g(n)\Big) \wedge \Big(\exists n_1 \in \mathbb{N},\ n \geq n_1 \Rightarrow f(n) \geq b \wedge g(n) \geq b\Big) \wedge \Big(b > 1\Big) \Rightarrow \Big(\exists d_1, d_2, n_2 \in \mathbb{R}^+,\ \forall n \in \mathbb{N},\ n \geq n_2 \Rightarrow d_1 \cdot \log_b(g(n)) \leq \log_b(f(n)) \leq d_2 \cdot \log_b(g(n)\Big)$

Proof. Let $f, g : \mathbb{N} \to \mathbb{R}^+$, and $b \in \mathbb{R}^+$. Assume $c_1 = 1$, $c_2 = b$, and $n_0 = 1$, and $n \in \mathbb{N}$ such that $n \geq n_0$ and $c_1 \cdot g(n) \leq f(n) \leq c_2 \cdot g(n)$. Assume $f(n)$ and $g(n)$ are eventually $\geq b$. Assume $b > 1$. Let $d_1 = 1$, $d_2 = 2$, and $n_2 = n_0$. Assume $n \geq n_2$.

We need to show $d_1 \cdot \log_b g(n) \leq \log_b f(n) \leq d_2 \cdot \log_b g(n)$.

We will do so in two parts. One for $(d_1 \cdot \log_b g(n) \leq \log_b f(n))$ and the other for $(\log_b f(n) \leq d_2 \cdot \log_b g(n))$.

Part 1 ($d_1 \cdot \log_b g(n) \leq \log_b f(n)$):

The assumption tell us

$$c_1 \cdot g(n) \leq f(n) \tag{1}$$

Then, it follows from the fact $\forall x, y \in \mathbb{R}^+, x \geq y \Leftrightarrow \log x \geq \log y$

$$\log(c_1 \cdot g(n)) \leq \log(f(n)) \tag{2}$$

Then, using the fact $b > 1$, we can calculate

$$\frac{\log(c_1 \cdot g(n))}{\log b} \leq \frac{\log(f(n))}{\log b} \tag{3}$$
$$\frac{\log(c_1) + \log(g(n))}{\log b} \leq \frac{\log(f(n))}{\log b} \tag{4}$$

Then,

$$\frac{\log(g(n))}{\log b} \leq \frac{\log(f(n))}{\log b} \tag{5}$$

by the fact $c_1 = 1$ and $\log c_1 = 0$.

Then, since $\frac{\log f(x)}{\log b} = \log_b f(x)$,

$$\log_b(g(n)) \leq \log_b(f(n)) \tag{6}$$

Then, because we know $d_1 = 1$, we can conclude

$$\log_b(g(n)) \leq d_1 \cdot \log_b(f(n)) \tag{7}$$

Part 2 ($\log_b f(n) \leq d_2 \cdot \log_b g(n)$):

The assumption tells us

$$f(n) \leq c_2 \cdot g(n) \tag{8}$$

Then, it follows from the fact $\forall x, y \in \mathbb{R}^+, x \geq y \Leftrightarrow \log x \geq \log y$

$$\log(f(n)) \leq \log(c_2 \cdot g(n)) \tag{9}$$

Then, using the fact $b > 1$, we can calculate

$$\frac{\log(f(n))}{\log b} \leq \frac{\log(c_2 \cdot g(n))}{\log b} \tag{10}$$

$$\frac{\log(f(n))}{\log b} \leq \frac{\log(c_2) + \log(g(n))}{\log b} \tag{11}$$

Then, since $c_2 = b$,

$$\frac{\log(f(n))}{\log b} \leq \frac{\log(b) + \log(g(n))}{\log b} \tag{12}$$

Then, using the fact $g(n)$ is eventually $\geq b$, we can write

$$\frac{\log(f(n))}{\log b} \leq \frac{\log(g(n)) + \log(g(n))}{\log b} \tag{13}$$

$$\frac{\log(f(n))}{\log b} \leq \frac{2 \cdot \log(g(n))}{\log b} \tag{14}$$

Then, since $\frac{\log f(x)}{\log b} = \log_b f(x)$,

$$\log_b(f(n)) \leq 2 \cdot \log_b(g(n)) \tag{15}$$

Then, because we know $d_2 = 2$, we can conclude

$$\log_b(f(n)) \leq d_2 \cdot \log_b(g(n)) \tag{16}$$

$\square$

Notes:

- $\forall x, y \in \mathbb{R}^+, x \geq y \Leftrightarrow \log x \geq \log y$
- $\exists c_1, c_2, n_0 \in \mathbb{R}^+$, $\forall n \in \mathbb{N}, n \geq n_0 \Rightarrow c_1 \cdot g(n) \leq f(n) \leq c2 \cdot g(n)$
- **Definition of Eventually:** $\exists n_0 \in \mathbb{N}, n \geq n_0 \Rightarrow P$, where $P : \mathbb{N} \to \{\text{True}, \text{False}\}$

b. *Proof.* Let $k \in \mathbb{N}$.

First, we will analyze the cost of loop 2 over iteration of loop 1.

The code tells us loop 2 starts at $j_k = 1$ with j_k increasing by a factor of 3 per iteration until $j_k \geq 1$.

Using these facts, we can calculate that the terminating condition occurs when

$$
\begin{aligned}
3^k &\geq i && (1)\\
k &\geq \log_3 i && (2)
\end{aligned}
$$

Because we know the number of iterations is the smallest value of k satisfying the above inequality, we can conclude loop 2 has

$$\lceil \log_3 i \rceil \tag{3}$$

iterations.

Next, we need to determine the total number of iterations of loop 2 over all iterations of loop 1.

The code tells us loop 1 starts at $i = 1$ and ends at $i = n$ with each i increasing by 1 per iteration.

Using these facts, we can conclude loop 2 has total of

$$\lceil \log_3 1 \rceil + \lceil \log_3 2 \rceil + \cdots + \lceil \log_3 n \rceil = \sum_{i=1}^{n} \lceil \log_3 i \rceil \tag{4}$$

iterations. □

c. After scratching head and looking at solution many times, I realized that there are many things I do not yet understand, and it's the best to write what I have and learn from the solution. Here is my best attempt :).

Proof. Let $n \in \mathbb{N}$.

The previous answer tells us the exact cost of the algorithm is

$$\sum_{i=1}^{n} \lceil \log_3 i \rceil \tag{1}$$

Then, it follows by changing the variable i to $i' = \log_3 i$ we can write

$$\sum_{i'=0}^{\lceil \log_3 n \rceil} i' \tag{2}$$

Then, because we know $\sum\limits_{i=0}^{n} i = \frac{n(n+1)}{2}$, we can conclude

$$\sum_{i'=0}^{\lceil \log_3 n \rceil} i' = \frac{(\lceil \log_3 n \rceil)(\lceil \log_3 n \rceil + 1)}{2} \tag{3}$$

$$= \frac{\lceil \log_3 n \rceil^2 + \lceil \log_3 n \rceil}{2} \tag{4}$$

Then, we can conclude the runtime of the algorithm is $\Theta(\log_3^2 n)$. □

Correct Solution:

We need to determne Θ of the algorithm.

We will prove that the Θ of the algorithm is $\Theta(n \log n)$.

The answer to previous question tells us the total exact cost of the algorithm is

$$\sum_{i=1}^{n} \lceil \log_3 i \rceil \tag{5}$$

Then, by using fact 1 $\forall x \in \mathbb{R}, x \leq \lceil x \rceil \leq x + 1$, we can calculate

$$\sum_{i=1}^{n} \log_3 i \leq \sum_{i=1}^{n} \lceil \log_3 i \rceil \leq \sum_{i=1}^{n} \Big(\log_3 i + 1 \Big) \tag{6}$$

$$\sum_{i=1}^{n} \log_3 i \leq \sum_{i=1}^{n} \lceil \log_3 i \rceil \leq \Big(\sum_{i=1}^{n} \log_3 i + \sum_{i=1}^{n} 1 \Big) \tag{7}$$

$$\sum_{i=1}^{n} \log_3 i \leq \sum_{i=1}^{n} \lceil \log_3 i \rceil \leq \sum_{i=1}^{n} \log_3 i + n \tag{8}$$

Then,

$$\log_3 \Big(\prod_{i=1}^{n} i \Big) \leq \sum_{i=1}^{n} \lceil \log_3 i \rceil \leq \log_3 \Big(\prod_{i=1}^{n} i \Big) + n \tag{9}$$

$$\log_3(n!) \leq \sum_{i=1}^{n} \lceil \log_3 i \rceil \leq \log_3(n!) + n \tag{10}$$

by the fact $\forall a, b \in \mathbb{R}^+, \log(a) + \log(b) = \log(ab)$.

Then,

$$\frac{\ln n!}{\ln 3} \leq \sum_{i=1}^{n} \lceil \log_3 i \rceil \leq \frac{\ln(n!)}{\ln 3} + n \tag{11}$$

by changing the base to e using the formula $\log_3 n! = \frac{\log_e n!}{\log_e 3} = \frac{\ln n!}{\ln 3}$.

Now, the fact 2 tells us $n! \in \Theta(e^{n \ln n - n + \frac{1}{2} \ln n})$.

Because we know from fact 3 that $n \ln n - n + \frac{1}{2} \ln n$ is eventually ≥ 1, we can conclude $e^{n \ln n - n \frac{1}{2} \ln n}$ is eventually $\geq e$.

Since $n!$ is also eventually $\geq e$, by using solution to problem 1.a with $g(n) = n!$ and $f(n) = e^{n \ln n - n + \frac{1}{2} \ln}$ and $b = e$, we can write

$$\ln(n!) \in \Theta(\ln(e^{n \ln n - n + \frac{1}{2} \ln n})) \tag{12}$$

$$\ln(n!) \in \Theta(n \ln n - n + \frac{1}{2} \ln n) \tag{13}$$

Then,

$$\ln(n!) \in \Theta(n \ln n) \tag{14}$$

by the fact $n \ln n - n + \frac{1}{2} \ln n \in \Theta(n \ln n)$.

So, since the algorithm runs at least $\frac{\ln n!}{\ln 3}$, we can conclude it has asymptotic lower bound of $\Omega(n \ln n)$, and since the algorithm runs at most $\frac{\ln n!}{\ln 3} + n$, we can conclude it has upper bound running time of $\mathcal{O}(n \ln n)$.

Since the value of Ω and $\mathcal{O}$ are the same, we can conclude the algorithm has running time of $\Theta(n \ln n)$ or $\Theta(n \log n)$.

Notes:

- In a main flow of proof, when there is a huge interruption like showing $\ln(n!) \in \Theta(n \ln n)$, how can a sentence be started to tell the audience we are working on another major idea?
- When an interruption in proof has been occured for another major part of a proof, how can a sentence be started to combine parts together?
- How can a sentence be written to say condition x_1, x_2, and x_3 are satisfied, so a statement y can be used to an equation or an idea?

d. We need to evaluate tight asymptotic upper bound.

We will prove that the tight asymptotic upper bound of the algorithm is $\mathcal{O}(n^2)$.

First, we need to analyze the number of iterations of loop 2 per iteration of loop 1.

The code tells us loop 2 starts at $j = 0$ and ends at most$j = i - 1$ with j increasing by 1 per iteration.

Then, using these facts, we can conclude loop 2 has at most

$$\left\lceil \frac{i-1-0+1}{1} \right\rceil = i \tag{1}$$

iterations.

Next, we need to determine the total number of iterations of loop 2 over all iterations of loop 1.

The code tells us that loop 1 starts at $i = n$ and ends at most $i = 0$ with i decreasing by 1 per iteration.

Because we know each iteration of loop 1 takes i iterations by loop 2, using these facts, we can conclude the total number of iterations of loop 2 is at most

$$\begin{align} n + (n-1) + (n-2) + \cdots + 0 &= \sum_{i=1}^{n} \tag{2} \\ &= \frac{n(n+1)}{2} \tag{3} \end{align}$$

iterations, or $\mathcal{O}(n^2)$.

Question 2

Question 3

Question 4