# CSC369 Week 2 Notes

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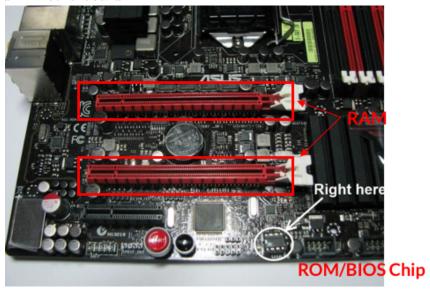
## 1 System Calls

- Bootstraping
  - Bootstraping

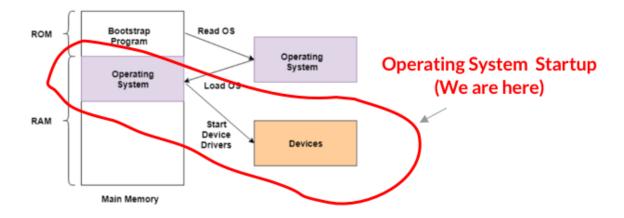


- \* executes **Bootstrap Program** 
  - · is the first code that runs when the computer system is started
- \* Entire operating system depnds on the bootstrap program to work correctly
- \* Locates and loads kernel (code of operating system) onto RAM
  - $\cdot$  kernel = code of the operating system
  - $\cdot$  kernel is in HDD
- \* Bootstrap program is in ROM
- ROM
  - \* is called **read-only-memory**
  - \* Is also called **BIOS chip** (Basic Input/Output System)

- \* is non-volatile
- \* is stored in motherboard



• Operating System Startup

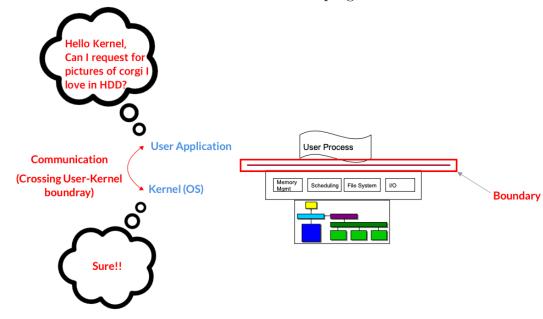


- Initializes OS
  - \* Initialize internal data structures
  - \* Create first process
  - \* Switch mode to user and start runing first process
  - \* Wait for something to happen
- Requesting OS Services
  - Some services offered by OS are:
    - \* Program execution
      - · Loading program to memory and executing program

- \* I/O operations
  - · Keyboard, mouse, speaker
- \* File system manipulation
  - · Reading and writing files and directories
- \* Error Detection
  - · Error that pops when printer ink is empty

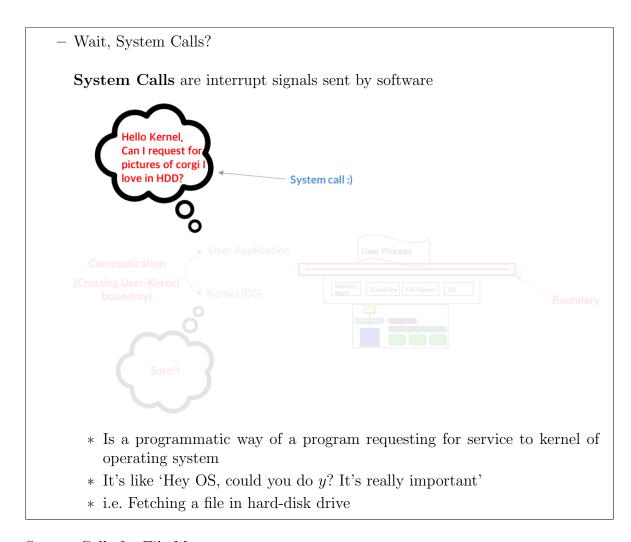


- Operating system and user programs are isolated
- How do they communicate?
- Boundary Crossings
  - Boundary
    - \* Is the line between user applications and kernel
    - \* Data is difficult to move back and forth between this line
  - Boundary Corssings
    - \* Is the communication that occurs between a program and kernel



- \* Communication occurs by sending data from one program into kernel, and then back
- More can be found here
- System Calls for Process Management
  - Major system calls

Call	Description
pid = fork()	Create a child process identical to the parent
pid = waitpid(pid, &statloc, options)	Wait for a child to terminate
s = execve(name argv, environp)	Replace a process' core image
exit(status)	Terminate process execution and return status



- System Calls for File Management
  - Major system calls

Call	Description
fd = open(file, how,)	Open a file for reading, writing, or both
s = close(fd)	Close an open file
n = read(fd, buffer, nbytes)	Read data from a file into a buffer
n = write(fd, buffer, nbytes)	Write data from a buffer into a file
position = lseek(fd, offset, whence)	Move the file pointer
s = stat(name, &buf)	get a file's status information

### • Systemm Call Interface

#### - Interface

\* Is a point where two systems, subjects, organizations, etc. meet and interact. (Definition)

#### - System Call Interface

\* Is the point where user mode and kernel mode meet and interact

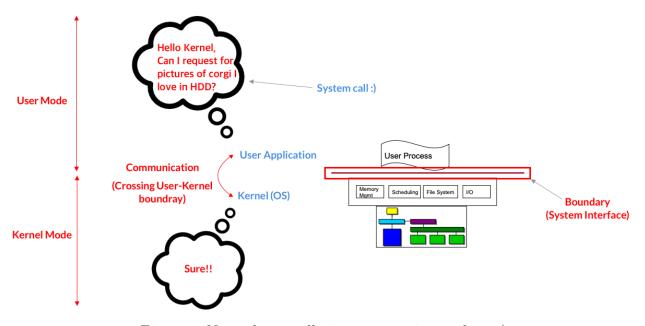


Figure 1: Now, there really is a party going on here:)

\* Most of the system call interface are hidden from the programmer by API

#### • System Call Operation

- Register is a local storage device present in CPU
- Register may include the address of the memory location instead of real data
- CPU takes data that has to be processed from register

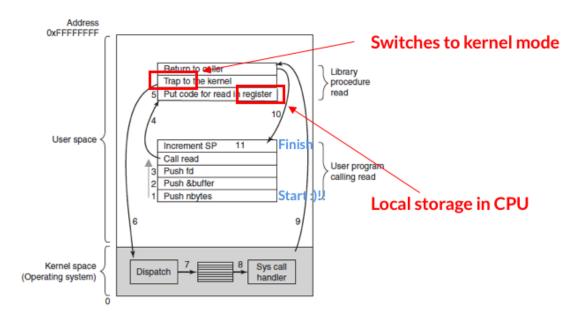


Figure 1-17. The 11 steps in making the system call read(fd, buffer, nbytes).