1. a) Trap instruction is run in user mode, and privileged operation is run in kernel mode

Notes

• Previliged Instructions

- Is the instruction that can run only in **kernel mode**
- Attempt at execution in **user mode** \rightarrow treated as an illegal operation & will not run.

• Trap

- Is a special hardware instruction
- Is a software generated interrupt ^[4]
- Is a type of synchronous interrupt ^[1]
- Is caused by an exceptional condition ^[1]
 - 1. Division by zero [1]
 - 2. Invalid memory access (segmentation fault) [1]
 - 3. Previleged instruction by **user mode** code ^[2]
- Usually results in a switch to **kernel mode** \rightarrow Operating system performs action \rightarrow Returns control to original process

• Trap Instruction

 Is executed when a user wants to invoke a service from the operating system (i.e. reading hard drive) in user mode

• User Mode

- Executing code has no ability to directly access hardware or reference memory
- Crashes are always recoverable ^[3]
- Is where most of the code on our computer are executed [3]

• Kernel Mode

- Executing code has complete and unrestricted access to the underlying hardware
 [3]
- Is generally reserved for the lowest-level, most trusted functions of the operating system [3]
- Is fatal to crash; it will halt the entire PC (i.e the blue screen of death) [3]

References

- 1) Wikipedia, Trap (computing), link
- 2) University of Utah, CS5460: Operating Systems Lecture 3 OS Organization, link
- 3) Coding Horror, Understanding User and Kernel Mode, link
- 4) ETH Zurich, Programming in Systems, link
- b) No. Lock uses a variable with binary states 0 (acquired) and 1 (available), where as semaphore uses counter variable that can have value greater than 1 to keep track of the amount of resource remaining.

Notes

• Locks

```
- Is a variable with two boolean states
  * 1 - (available/unlock/free)
  * 0 - (acquired/locked/held)

    Has two operations

 1. acquire()
                boolean test and set(boolean *lock)
                 {
                       boolean old = *lock;
                       *lock = True;
                       return old;
                 boolean lock;
                 void acquire(boolean *lock) {
                         while (test and set(lock));
                 }
 2. release()
                void release (boolean *lock) {
                         *lock = false;
                 }
```

 Is put around critical section to ensure critical section executes as if it's a single atomic instruction

```
lock_t mutex; // some globally-allocated lock 'mutex'
lock(&mutex);
balance = balance + 1;
unlock(&mutex);
```

- Can only be released by the thread that acquired it
- Is used to protect shared resource (e.g. from race condition in files and data structure) [2]

Semaphore

Is an abstract data types suitable for synchronization problems [2]

- Has variable count that allows arbitrary resource count [1]
- Has two atomic operations
 - 1. (wait/P/decrement) block until count > 0 then decrement variable

```
wait(semaphore *s) {
     while (s->count == 0);
     s->count -= 1;
}
```

2. (signal/V/increment) - increment count, unblock a waiting thread

```
signal(semaphore *s) {
    s->count += 1;
    ..... //unblock one waiter
}
```

- Can be signaled by any thread ^[2]

References

- 1) Wikipedia, Semaphore (programming), link
- 2) Stack Overflow, Difference between binary semaphore and mutex, link
- c) All shared variable of threads are located in heap memory, and if multiple threads access the variable at the same time, a race condition would occur.

Notes

• What is concurrency error? Where and when does it occur?

• Race condition

 if multiple threads of execution enter the critical section roughly at the same time; both attempt to update the shared data structure, leading to a supring result

Concurrency

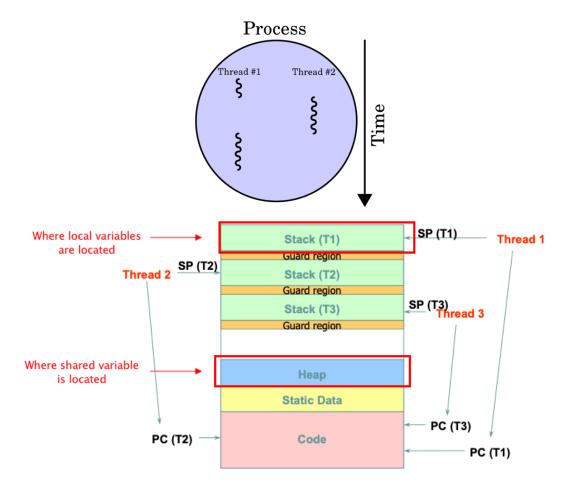
 Is the ability of different parts or units of a program, algorithm, or problem to be executed out of order, without affecting the final outcome. ^[1]

• Concurrency Error

- Two types of concurrency errors

• Thread

– Is the smallest sequence of programmed instructions that can be managed independently by a schdeduler ^[2]



- A thread is bound to a single process
- A process can have multiple threads

References

- 1) Wikipedia, Concurrency (computer science), link
- 2) Wikipedia, Thread, link