# CSC373 Worksheet 6 Solution

## August 13, 2020

## 1. 1. Multiply objective function by - 1

Maximize

$$-2x_1 - 7x_2 - x_3$$

Subject to

$$x_1 - x_3 = 7$$
$$3x_1 + x_2 \ge 7$$
$$x_2 \ge 0$$
$$x_3 \le 0$$

2. Replace non-nonnegative constraints  $x_1$ 

Maximize

$$-2x_1' + 2x_1'' - 7x_2 - x_3$$

Subject to

$$x'_{1} - x''_{1} - x_{3} = 7$$

$$3x'_{1} - 3x''_{1} + x_{2} \ge 7$$

$$x'_{1}, x''_{1}, x_{2} \ge 0$$

$$x_{3} \le 0$$

3. Replace non-nonnegative constraints  $x_3$ 

Maximize

$$-2x_1' + 2x_1'' - 7x_2 - x_3' + x_3''$$

Subject to

$$x'_1 - x''_1 - x'_3 + x''_3 = 7$$
$$3x'_1 - 3x''_1 + x_2 \ge 7$$
$$x'_1, x''_1, x_2, x'_3, x''_3 \ge 0$$

4. Replace equality constraints with  $\geq$  and  $\leq$ 

Maximize

$$-2x_1' + 2x_1'' - 7x_2 - x_3' + x_3''$$

Subject to

$$\begin{aligned} x_1' - x_1'' - x_3' + x_3'' &\leq 7 \\ x_1' - x_1'' - x_3' + x_3'' &\geq 7 \\ 3x_1' - 3x_1'' + x_2 &\geq 7 \\ x_1', x_1'', x_2, x_3', x_3'' &\geq 0 \end{aligned}$$

5. Correct greater-than-or-equal-to inequality constraints

Maximize

$$-2x_1' + 2x_1'' - 7x_2 - x_3' + x_3''$$

Subject to

$$x'_1 - x''_1 - x'_3 + x''_3 \le 7$$

$$-x'_1 + x''_1 + x'_3 - x''_3 \le -7$$

$$-3x'_1 + 3x''_1 - x_2 \le 7$$

$$x'_1, x''_1, x_2, x'_3, x''_3 \ge 0$$

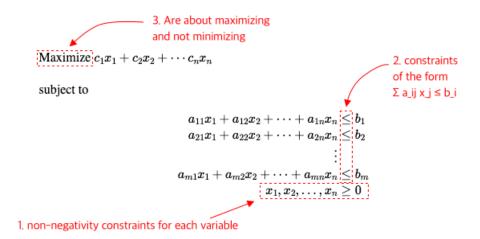
#### Notes:

### • Linear Programming

- Is a method to achieve the best outcome (such as maximum profit or lowest cost) in a mathematical model whose requirements are represented by linear relationships. [1]
- Is named to make it sound cool for government funding
  - \* Like dynamic programming
- Applications
  - \* Microeconomics (maximize profits, minimize costs)
  - \* Company management

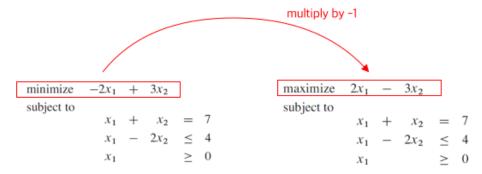
#### • Standard Form

- Is a form of linear programming
- Are about maximizing, not minimizing [2]
- All have a positivity constraint for each variable [2]
- All other constraints are all of the form "linear combination of variables  $\leq$  constant".  $^{[2]}$

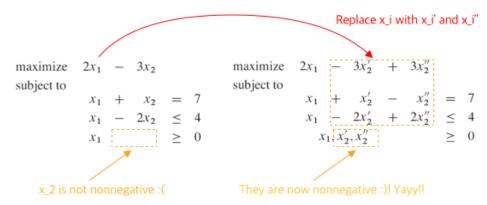


#### • Converting Linear Programming to Standard Form

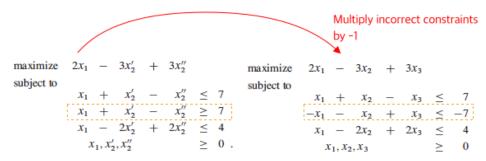
- 1) The objective function might be a minimization rather than a maximization
  - Negate coefficients of the objective function



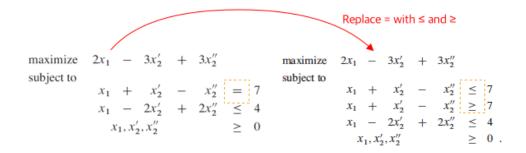
- 2) There might be variables without nonnegativity constraints
  - Replace each non-nonnegative variable  $x_i$  with  $x_i'$  and  $x_i''$
  - Modify linear program



- 3) There might be **equality constraints**, which have an equal sign rather than a less-than-or-equal-to sign
  - Replace equality constraint  $f(x_1, x_2, ..., x_n) = b$  with  $f(x_1, x_2, ..., x_n) \le b$  and  $f(x_1, x_2, ..., x_n) \ge b$



- 4) There might be **inequality constraints**, but instead of having a less-than-or-equal-to-sign
  - Multiply incorrect inequality constraints by -1



#### References:

- 1) Wikipedia, Linear Programming, link
- 2) Instituto de Mathematicas, Standard form for Linear Programs, link

2.

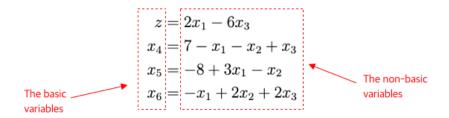
$$z = 2x_1 - 6x_3$$

$$x_4 = 7 - x_1 - x_2 + x_3$$

$$x_5 = -8 + 3x_1 - x_2$$

$$x_6 = -x_1 + 2x_2 + 2x_3$$

The basic variables are variables on the lhs (i.e B = 4, 5, 6), and the non-basic variables are the variables on the rhs of the expressions (i.e. N = 1, 2, 3).



#### Notes:

#### • Slack Form

- Is a form of linear programming
- Is for efficient solving of liner programming problem using simplex algorithm

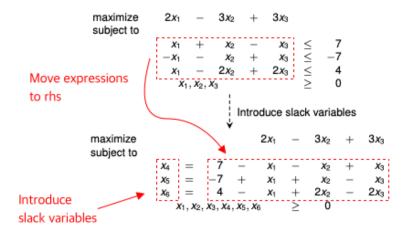


## • Converting Linear Programs into Slack Form

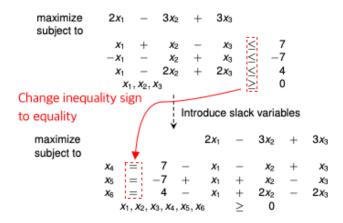
- 1) Start from the standard form of linear programming
- 2) Shift objective functions to right



3) Introduce slack variable  $x_i$  to lhs and move expressions  $\sum_{j=1}^n a_{ij}x_j$  to rhs



4) Change inequalities in linear programming to equality



5) Use Variable z to denote objective function



6) Omit the nonnegativivty constraints

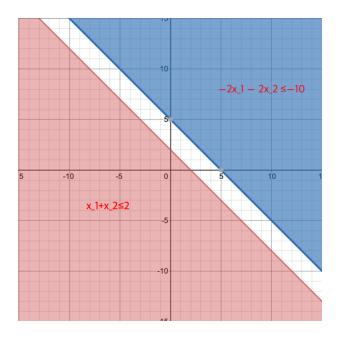


## References:

- 1) Cambridge University, Linear Programming, link
- 3. Multiplying the first expression (under subject to) by 2, and summing the inequality constraints, we have

$$0 \le -6 \tag{1}$$

which is impossible.



## **Notes**

- I noticed that infeasible solution has non-overlapping region
- Infeasible
  - A Linear Program is infeasible if there is no solution that satisfies all of the constraints
- 4. Let  $x_1 = 3r$  and  $x_2 = r$  where  $r \ge 0$ . Then the inequality constraints become

$$-2(3r) + (r) = -5r \le -1$$
$$-(3r) - 2(r) = -5r \le -2$$
$$3r, r > 0$$

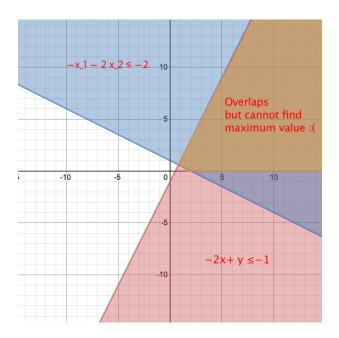
and are valid.

Now, looking at the objective functions, with  $x_1 = 3r$  and  $x_2 = r$ , it becomes

$$3r - r = 2r$$

which increases without bound.

Thus, there is no maximum, and the linear program is unboudned.



## Notes

- I learned that to show an LP is unbounded, I first have to subtitute  $x_i$  with a common variable r (e.g.  $x_1 = 3r$ ,  $x_2 = r$ ), check inequality constraints, and then look at objective functions and see if I can get max/min.
- Unbounded
  - A Linear Program is unbounded if it has some feasible solutions but does not have a finite optimal objective value

## References:

- 1) CLRS Solutions, 29.1 Standard and slack forms, link
- 5. Notes

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