- 1. a) 1) 4 inode blocks. 1 for the file c, and 3 for the directdories /, a, b
 - 2) 3 directory blocks one for root /, one for a, the other for b
 - 3) 1 single indirect block as far as we know. The file definitely has more than 12 blocks (# of data blocks pointed by direct pounters), but less than 1036 (# of data blocks pointed by direct pointers and single indirect pointers). We are reading block 1034.
 - 4) 1 data block for file c
 - b) All of the above

Notes

• Inode



- Is short form of index node
- Describes a file system object such as file or data
- Contains all information about a file/directory, including
 - * File Type,
 - * Size
 - * Number of blocks allocated to it
 - * Protection information
 - * Time information (e.g time created, time modified)
 - * Location of data blocks residing on disk

References

- 1) Wikipedia, Inode, link
- 2) Machanick, Philip. (2016). Teaching Operating Systems: Just Enough Abstraction. 642. 10.1007/978-3-319-47680-3_10., link

c) Size, the location of data blocks that reside on disk

Notes

- I wonder what information about blocks inode has. Is it total number of blocks both inode and data, or just data?
- I struggled a bit on this one. I should find an easier way to remember which information inode has
- d) Inode Bitmap and Data Block Bitmap
 - (b) Data Leak
 - (c) Inode Leak
 - New Directory Inode
 - (a) No inconsistency
 - Inode Bitmap, Data Block Bitmap, Existing Directory Data, New Directory Inode, and New Directory Data
 - (e) Inconsistent inode data
 - Inode Bitmap, and New Directory Inode
 - (c) Inode leak
 - (d) Multiple file paths may point to same inode
 - New Directory Inode, Existing Directory Inode, Existing Directory Data
 - (e) Inconsistent inode data
 - (f) Something points to garbage

Correct Solution

- Inode Bitmap and Data Block Bitmap
 - (b) Data Leak
 - (c) Inode Leak
- New Directory Inode
 - (a) No inconsistency
- Inode Bitmap, Data Block Bitmap, Existing Directory Data, New Directory Inode, and New Directory Data
 - (e) Inconsistent inode data
- Inode Bitmap, and New Directory Inode
 - (c) Inode leak
- New Directory Inode, Existing Directory Inode, Existing Directory Data

```
(d) - Multiple file paths may point to same inode(f) - Something points to garbage
```

Notes

- I wonder how system call for reading file/directory works in UNIX. Does it check for bitmap?
- I wonder how system call for deleting file/directory works in UNIX
- I wonder how system call for creatubg file/directory works in UNIX
- Learned that
 - Missing Inode Bitmap multiple file paths may point to same inode
- File API
 - open (create/access file)
 - * Is a system call
 - * Does three things on creation
 - 1) make structure (inode) that racks all relevant information about file
 - 2) link human readible name to the file, and put that link to a directory
 - 3) increment **reference count** in inode
 - * Syntax:

```
int fd = open("foo". O_CREAT|O_WRONLY|O_TRUNC, S_IRUSR|S_IWUSR)
```

- · O_CREAT Creates file "foo" if does not exist
- · O_WRONLY Open file for writing only (default)
- · O_TRUNC Overwrites existing file Need example/Clarification
- · Can have multiple flags
- * Returns file descriptor or fd for short
 - · Is an integer
 - · Is used to access a file
 - · Is private per process
 - · Can be used to read() and write() files

```
File can also
                                        File can also
                                                                    File can also
                       File can be read
                                        be written by
                                                      be read by
                                                                    be read by
                       by owner
                                        owner
                                                      group
                                                                    others
#include <fcntl.h>
int fd;
mode t mode = S IRUSR | S IWUSR
                                        S IRGRP
                                                       S IROTH
char *filename = "/tmp/file";
fd = open(filename, O WRONLY | O CREAT | O TRUNC, mode);
. . .
                              Means

    File is Writable AND

                              Create file if doesn't exist AND
                              3. Overwrite file if exists
```

- (read) (read file)
 - * Is a system call
 - * Syntax:

```
ssize_t read (int fd, void *buf, size_t count)
```

- · fd file descriptor (from open ())
- · buf container for the read data
- · count number of bytes to read
- * Returns number of bytes read, if successful
- * Returns 0 if is at, or past the end of file

Example

```
char buf[4096];
int fd = open("/a/b/c", 0); // open in read-only mode
lseek(fd, 1034*4096, 0); // seek to position (1034*4096) from start of file
read(fd, buf, 4096); // read 4k of data from file
```

```
Current
                                 Return
System Calls
                                  Code
                                           Offset
fd = open("file", O_RDONLY);
                                                              read continues
read(fd, buffer, 100);
                                    100
                                              100
                                                              for each call
                                              200
read(fd, buffer, 100);
                                    100
                                              300
read(fd, buffer, 100);
                                    100
read(fd, buffer, 100);
                                              300
                                      0
close(fd);
                                                              if at end
```

- write (write file)

- * Is a system call
- * Writes data out of a buffer
- * Syntax:

```
ssize_t write (int fd, const void * buf, size_t nbytes)
```

- · fd file descriptor
- · buf A pointer to a buffer to write to file
- · nbytes number of bytes to write. If smaller than buffer, the output is truncated

Example

```
#include <unistd.h>
#include <fcntl.h>

int main(void)
{
    int filedesc = open("testfile.txt", O_WRONLY | O_APPEND);

    if (filedesc < 0) {
        return -1;
    }

    if (write(filedesc, "This will be output to testfile.txt\n", 36) != 36) {
        write(2, "There was an error writing to testfile.txt\n", 43);
        return -1;
    }

    return 0;
}</pre>
```

- lseek
 - * Reads or write to a specific offset within a file
 - * Syntax:

```
off_t lseek (int fd, off_t offset, int whence)
```

- · fd file descriptor
- · offset the offset of pointer within file (in bytes)
- · whence the method of offset

```
SEEK_SET - offset from the start of file (absolute)
SEEK_CUR - offset from current location + offset bytes (relative)
SEEK_END - offset from the end of file
```

- * Returns offset amount (in bytes) from the beginning of file
- * Returns -1 if error

System Calls	Return Code	Current Offset	move 200 bytes from the
fd = open("file", O_RDONLY);	3	0	start of file
lseek(fd, 200, SEEK_SET);	200	200] A start of file
read(fd, buffer, 50);	50	250	
close(fd);	0	-	
			read 50 bytes

- rename (update file name)
 - * Is a system call
 - * Changes the name of file
 - * Is atomic (after crash, it will be either old or new, but not in-between)
 - * Syntax: int rename(const char *old, const char *new)
 - · old name of old file
 - · new name of new file
 - * Returns 0 if successful
 - * Returns -1 if error

- stat (get file info)
 - * displays metadata of a certain file stored in **inode**
 - * Syntax: int stat(const char *path, struct stat *buf)
 - · path file descriptor of file that's being inquired
 - · buf A stat structure where data about the file will be stored (see below)

```
struct stat {
  dev_t st_dev;
                          // ID of device containing file
            st_ino;
                         // inode number
  ino t
  mode_t
           st_mode;
                         // protection
                         // number of hard links
// user ID of owner
 nlink_t
            st_nlink;
 uid_t
            st_uid;
                         // group ID of owner
  gid_t
            st_gid;
                         // device ID (if special file)
// total size, in bytes
            st_rdev;
  dev_t
  off_t
            st_size;
 blksize_t st_blksize; // blocksize for filesystem I/O
 blkcnt_t st_blocks; // number of blocks allocated
            st_atime; // time of last access
st_mtime; // time of last modification
  time_t
 time_t
            st_ctime; // time of last status change
  time_t
);
```

Figure 39.5: The stat structure.

Example

```
#include <unistd.h>
#include <stdio.h>
#include <sys/stat.h>
#include <sys/types.h>
int main(int argc, char **argv)
{
    if(argc != 2)
        return 1;
    struct stat fileStat;
    if(stat(argv[1],&fileStat) < 0)</pre>
        return 1;
    printf("Information for %s\n",argv[1]);
    printf("----\n"):
    printf("File Size: \t\t%d bytes\n",fileStat.st_size);
printf("Number of Links: \t%d\n",fileStat.st_nlink);
    printf("File inode: \t\t%d\n",fileStat.st_ino);
    printf("File Permissions: \t");
    printf( (S_ISDIR(fileStat.st_mode)) ? "d" : "-");
    printf( (fileStat.st_mode & S_IRUSR) ? "r" : "-");
    printf( (fileStat.st_mode & S_IWUSR) ? "w" : "-");
    printf( (fileStat.st_mode & S_IXUSR) ? "x" : "-");
    printf( (fileStat.st_mode & S_IRGRP) ? "r" : "-");
    printf( (fileStat.st_mode & S_IWGRP) ? "w" : "-");
    printf( (fileStat.st_mode & S_IXGRP) ? "x" : "-");
    printf( (fileStat.st_mode & S_IROTH) ? "r" : "-");
printf( (fileStat.st_mode & S_IWOTH) ? "w" : "-");
    printf( (fileStat.st_mode & S_IXOTH) ? "x" : "-");
    printf("\n\n");
    printf("The file %s a symbolic link\n", (S_ISLNK(fileStat.st_mode)) ? "is" : "is not");
    return 0:
}
```

The result of above is:

```
$ ./testProgram testfile.sh

Information for testfile.sh

File Size: 36 bytes
Number of Links: 1
File inode: 180055
File Permissions: -rwxr-xr-x

The file is not a symbolic link
```

- unlink (removing file)
 - Is a system call
 - Removes a file (including symbolic link) from the system
 - Syntax: int unlink(const char *pathname)
 - * pathname path to file
 - Returns 0 if successful
 - Returns -1 if error

Example

```
#include <unistd.h>
char *path = "/modules/pass1";
int status;
...
status = unlink(path);
```

- mkdir (creating directory)
 - Is a system call
 - Syntax: int mkdir(const char *path, mode_t mode)
 - * path path of directory (including name)
 - * mode permission group
 - Returns 0 if successful
 - Returns -1 if error
 - directories can never be written directly
 - * directory is in format called File System Metadata
 - * directory can only be updated directly
 - creates two directories on creation . (current) and . . (parent)

```
#include <sys/types.h>
   #include <sys/stat.h>
   int status;
   status = mkdir("/home/cnd/mod1", S_IRWXU | S_IRWXG | S_IROTH | S_IXOTH);
• opendir, readdir, closedir (reading directory)
   - Are system calls
   - Are under <dirent.h> library
   - Requires struct dirent data structure
              struct dirent {
                char
                             d_name[256]; // filename
                            d_ino; // inode number
                ino_t
                                        // offset to the next dirent
                             d_off;
                off_t
                                         // length of this record
                unsigned short d_reclen;
               unsigned char d_type;
                                         // type of file
   - Syntax (opendir): DIR *opendir(const char *dirname)
     * dirname - directory path
     * Returns a pointer to the directory stream
     * The stream is positioned at the first entry in the directory.
   - Syntax (readdir): struct dirent *readdir(DIR *dirp);
     * dirp - directory stream
     * Returns a pointer to a direct structure representing the next directory entry
       in the directory stream
     * Returns NULL on reaching the end of the directory stream
   - Syntax (closedir): int closedir(DIR *dirp));
     * dirp - directory stream
     * Returns 0 if successful
     * Returns -1 otherwise
     Example
```

- rmdir (Deleting Directories)
 - Removes a directory whose name is given by path
 - Is performed only when directory is empty
 - Is included in <unistd.h> library
 - Fails if is symbolic link
 - Syntax: int rmdir(const char *path)
 - * path path of directory
 - Returns 0 if successful
 - Returns -1 if error

Example

```
#include <unistd.h>
int status;
...
status = rmdir("/home/cnd/mod1");
```

- unlink (Remove file)
 - Remove a link to a file
 - Is called **unlink** because it decrements **reference count** in inode
 - * Deletes file completely when reference count within the inode number is 0
 - Syntax:

```
#include <unistd.h>
int unlink(const char *pathname);
* pathname - pathname to file
```

- Returns 0 if successful
- Returns -1 if error
- Is used by linux command rm

#include <unistd.h>

```
char *path = "/modules/pass1";
int status;
...
status = unlink(path);
```

```
prompt> echo hello > file
prompt> stat file
... Inode: 67158084
                        Links: 1 ...
prompt> ln file file2
prompt> stat file
... Inode: 67158084
                        Links: 2 ...
prompt> stat file2
                        Links: 2 ...
... Inode: 67158084
prompt> ln file2 file3
prompt> stat file
... Inode: 67158084
                        Links: 3 ...
prompt> rm file
prompt> stat file2
                        Links: 2 ...
... Inode: 67158084
prompt> rm file2
prompt> stat file3
... Inode: 67158084
                        Links: 1 ...
prompt> rm file3
```

• Symbolic Link:

- Is directory entry containing "true" path to the file
- Is a shortcut that reference to a file instead of inode value [2]





• Hard Link:

- $-\,$ Is a direct reference to a file via its inode $^{[2]}$
- Is second directory entry identical to first





• Crash Consistency

- Inode before update

owner : remzi
permissions : read-write
size : 1
pointer : 4
pointer : null
pointer : null
pointer : null
i-number

- Inode after update

owner : remzi
permissions : read-write
size : 2
pointer : 4
pointer : 5
pointer : null
pointer : null

References

- 1) codewiki, stat, link
- 2) The Open Group Base Specification, unlink, link

e) a) Notes

- What is a sector? What is a sector address?
- inode
 - Inode block computation

block number = (inode #*sizeof(inode))/block size (1)

Example

Target: inode #32 Inode Size: 256 bytes Block Size: 4096 bytes

block number = (inode #*sizeof(inode))/block size (2)

$$=\frac{32*256}{4096}\tag{3}$$

$$=2 (4)$$

• inode bitmap

- uses bit to indicate whether the corres object/block is free
 - * 0 means free
 - * 1 means in use

• superblock

- Contains information about the following
 - * The number of inodes and data blocks in a particular file system
 - * The magic number of some knd to identify the file system type
 - * Where the inode table begins
- Is read first on mount before attaching to file system

• data bitmap