CSC236 Worksheet 8 Solution

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Question 1

- **Part 1 (Building L_1 and L_2):**

L_1:

$$Q = \{E, O\}$$
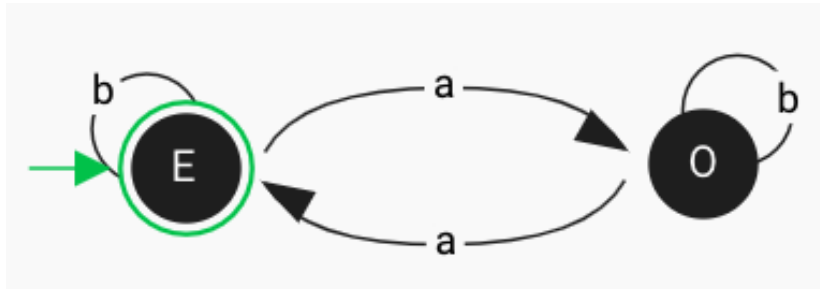$$\Sigma = \{a, b\}$$

$$\delta = \begin{array}{c|c|c} & a & b \\ \hline {}^*E & O & E \\ \hline O & E & O \end{array}$$

$$q_0 = E$$
$$F = \{E\}$$

Draw Diagram

L_2:

$$Q = \{0, 1, 2\}$$
$$\Sigma = \{a, b\}$$

$$\delta = \begin{array}{c|c|c} & a & b \\ \hline *0 & 1 & 1 \\ 1 & 2 & 2 \\ 2 & 0 & 0 \end{array}$$
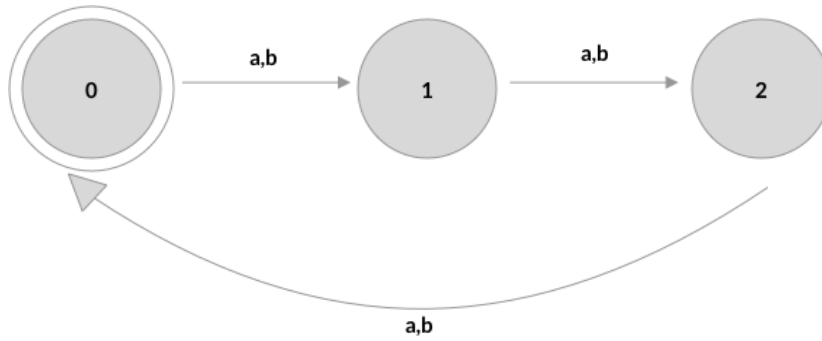
$$q_0 = 0$$
$$F = \{0\}$$

Draw Diagram

Part 1 (Building $L_1 \cap L_2$):

$$Q = \{(E, 0), (E, 1), (E, 2), (O, 0), (O, 1), (O, 2)\}$$
$$\Sigma = \{a, b\}$$

$$\delta = \begin{array}{c|c|c} & a & b \\ \hline *(E,0) & 1 & 1 \\ (E,1) & 2 & 2 \\ (E,2) & 0 & 0 \\ (O,0) & 1 & 1 \\ (O,1) & 2 & 2 \\ (O,2) & 0 & 0 \end{array}$$

$$q_0 = (E, 0)$$
$$F = \{(E, 0)\}$$

Correct Solution:

Part 1 (Building L_1 and L_2):

L_1:

$$Q = \{E, O\}$$
$$\Sigma = \{a, b\}$$

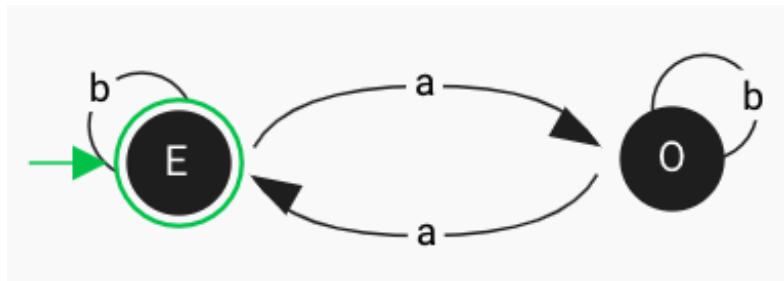$$\delta = \begin{array}{c|c|c} & a & b \\ \hline *E & O & E \\ \hline O & E & O \end{array}$$

$$q_0 = E$$
$$F = \{E\}$$

Draw Diagram

L_2:

$$Q = \{0, 1, 2\}$$
$$\Sigma = \{a, b\}$$

$$\delta = \begin{array}{c|c|c} & a & b \\ \hline *0 & 1 & 1 \\ \hline 1 & 2 & 2 \\ \hline 2 & 0 & 0 \end{array}$$

$$q_0 = 0$$
$$F = \{0\}$$

Draw Diagram
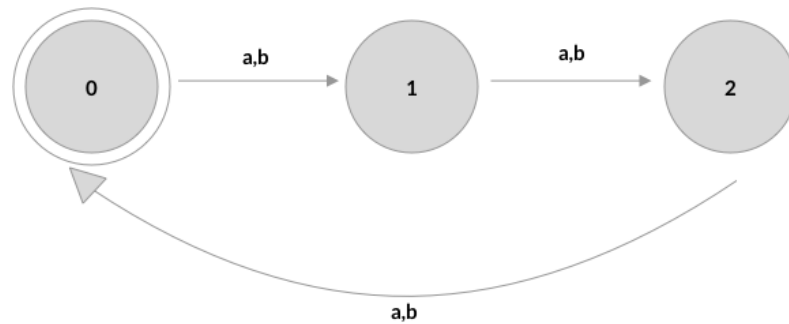
Part 1 (Building $L_1 \cap L_2$):

$$Q = \{(E,0), (E,1), (E,2), (O,0), (O,1), (O,2)\}$$
$$\Sigma = \{a, b\}$$

$$\delta = \begin{array}{c|c|c}
 & a & b \\
\hline
{}^*(E,0) & (O,1) & (E,1) \\
(E,1) & (O,2) & (E,2) \\
(E,2) & (O,0) & (E,0) \\
(O,0) & (E,1) & (O,1) \\
(O,1) & (E,2) & (O,2) \\
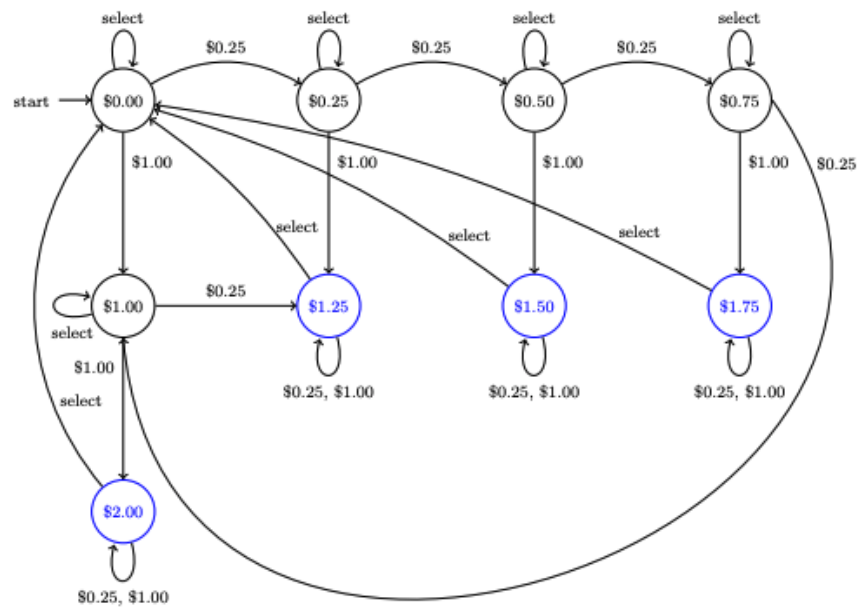(O,2) & (E,0) & (O,0) \\
\end{array}$$

$$q_0 = (E,0)$$
$$F = \{(E,0)\}$$

Notes:

- **Deterministic Finite State Automaton (DFSA)**: is a mathematical method of machine which, given any input string x, **accepts** or **rejects** x.

- Applications of DFSA

  1. Vending Machine

select
$0.25
select
$0.25
select
$0.25
select

start → $0.00 $0.25 $0.50 $0.75

$1.00 $1.00 $1.00 $1.00 $0.25

select   select   select

$1.00 $0.25 $1.25 $1.50 $1.75

select
$1.00
select

$0.25, $1.00 $0.25, $1.00 $0.25, $1.00

$2.00

$0.25, $1.00

2. Protocol analysis

3. Text parsing

4. Video game character behavior

Spot
Pac-Man

start → **Wander the Maze**   **Chase Pac-Man**

Lose
Pac-Man

Pac-Man Eats
Power Pellet

Reach
Central Base

Pac-Man Eats
Power Pellet

Power Pellet
Expires

Return to Base   **Flee Pac-Man**
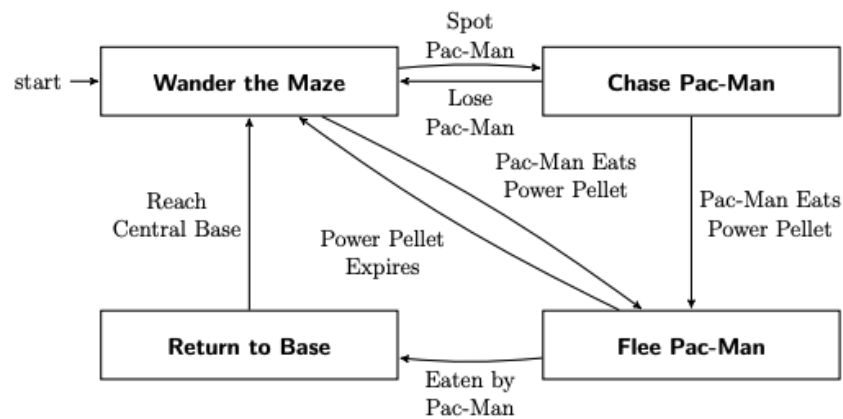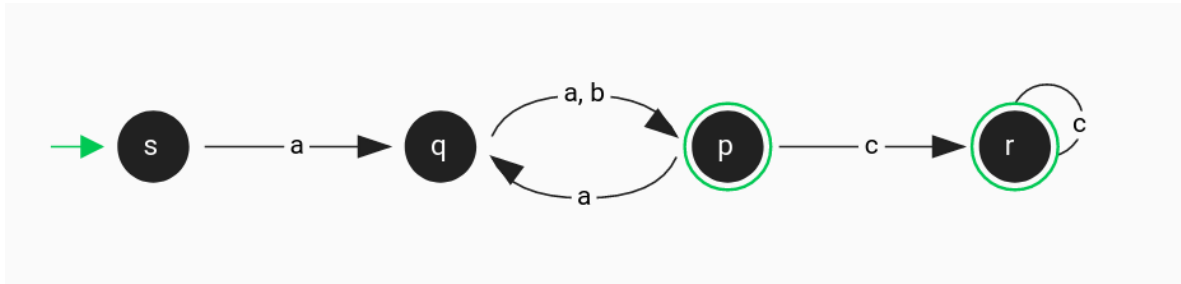
Eaten by
Pac-Man
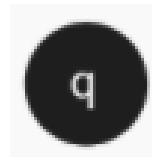
Figure 3: Behavior of a Pac-Man Ghost

5. Security Analysis

6. CPU control units (**)

7. Natural Language Processing (**)

8. Speech Recognition (**)
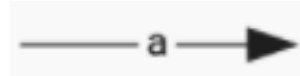
• Definitions and Syntax

- *DFSA M* is a quintuple $M = (Q, \Sigma, q_0, F, \delta)$, where
  * Q : a finite set of **states**.
    · Represents status of system
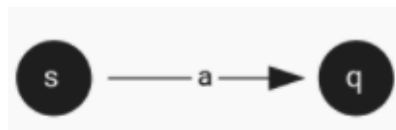   · Is represented by a black circle, i.e. s,q

    · i.e. automatic sliding door at walmart has two states: either close or open
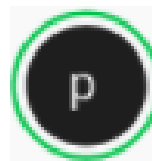   · i.e. traffic light has three states: red, yellow, green
  * Σ : a finite non-empty alphabet
    · is set of symbols in each transition, i.e. a, b, c

  * $q_0 \in Q$: the start or initial state
  * $\delta : Q \times \sigma \to Q$: a transition function
    · is a connection between two states.
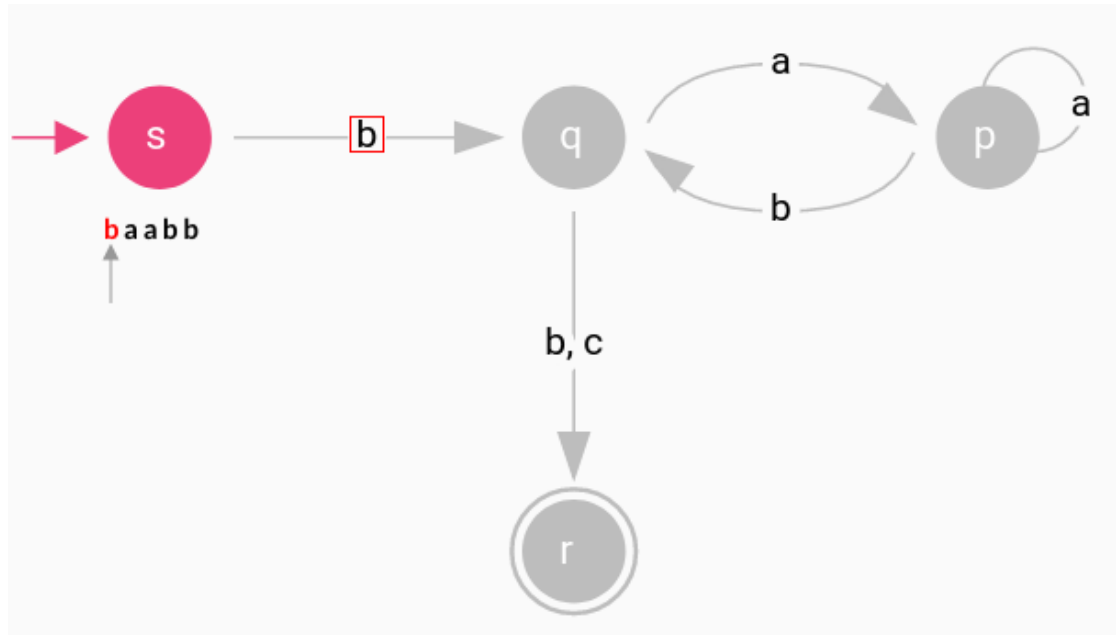   · is represented by an arrow

  * $F \subseteq Q$: the set of accepting or final states
    · Is represented by a double circle

    · Multiple accepting states may exists
    · Purpose: When processing ends, the output is either *accept* or *reject*
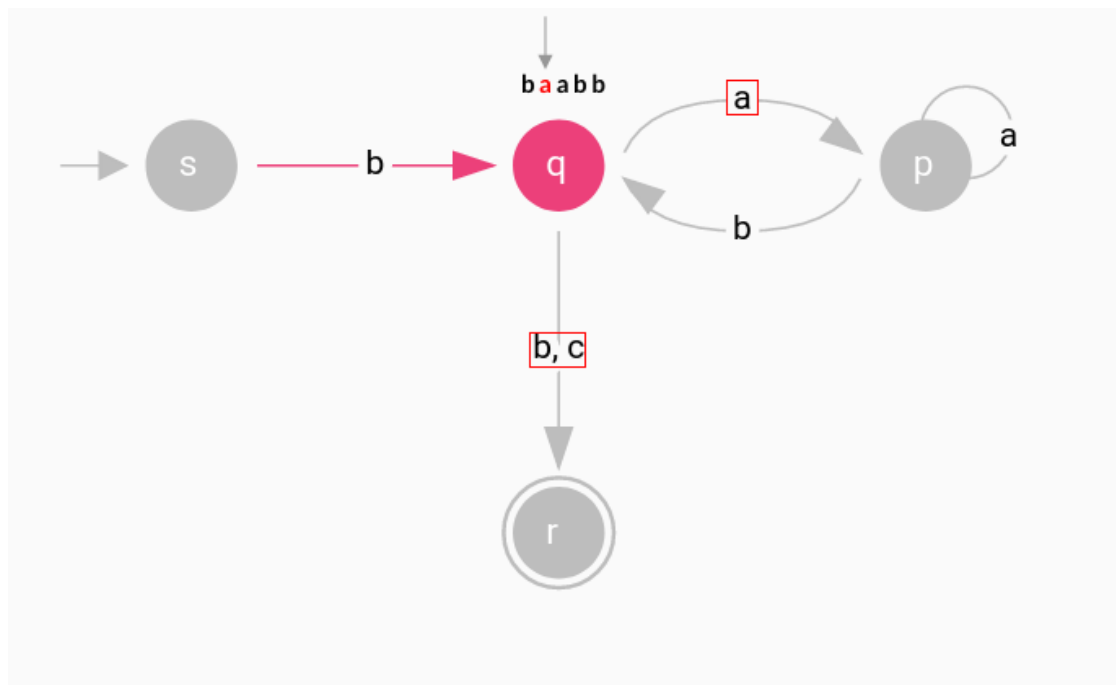
- Simple Example

1. First symbol of the input **baabb** is **b** and the current state is s.
2. Ask, is there any exiting transition from s that contains the symbol **b**?
3. The answer is yes, so move to q

– Step 2

1. Next symbol of the input **baabb** is **a** and the current state is q.
2. Ask, is there any exiting transition from q that contains the symbol **a** or **b,c**?

3. The answer is yes, and it's **a**. So move to p

– Step 3

1. Next symbol of the input **baabb** is **a** and the current state is p.
2. Ask, is there any exiting transition from p that contains the symbol **a** or **b**?
3. The answer is yes, and it's **a**. So move to p

– Step 4

1. Next symbol of the input **baabb** is **b** and the current state is p.
2. Ask, is there any exiting transition from p that contains the symbol **a** or **b**?
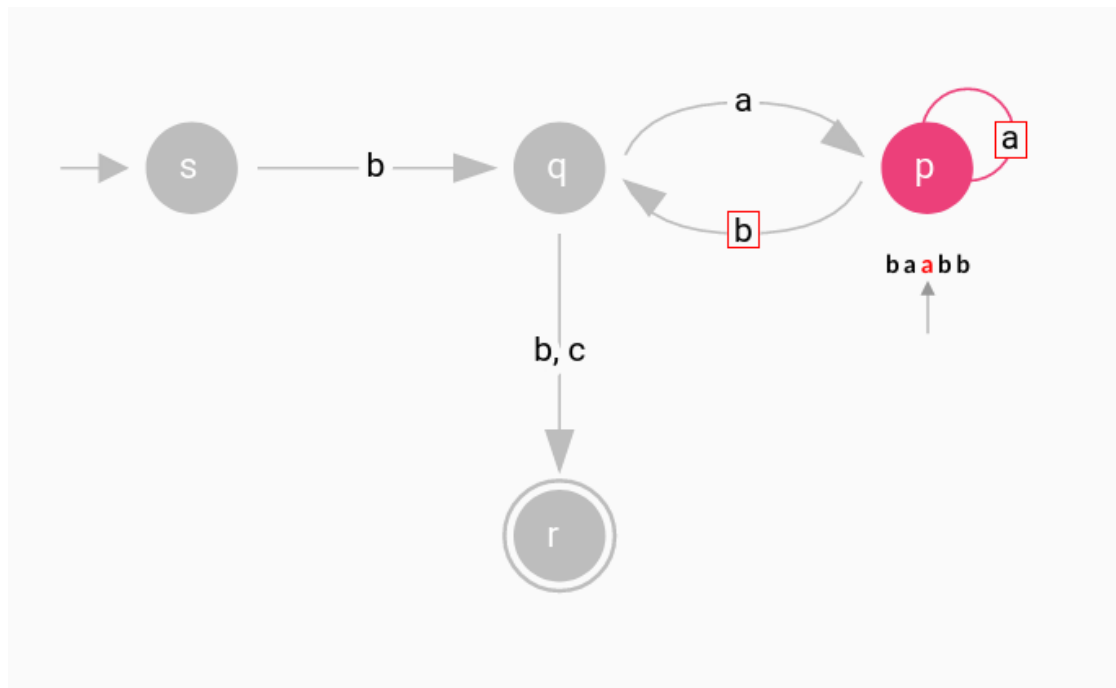3. The answer is yes, and it's **b**. So move to q

– Step 5

1. Next symbol of the input **baabb** is **b** and the current state is q.
2. Ask, is there any exiting transition from q that contains the symbol **a** or **b,c**?
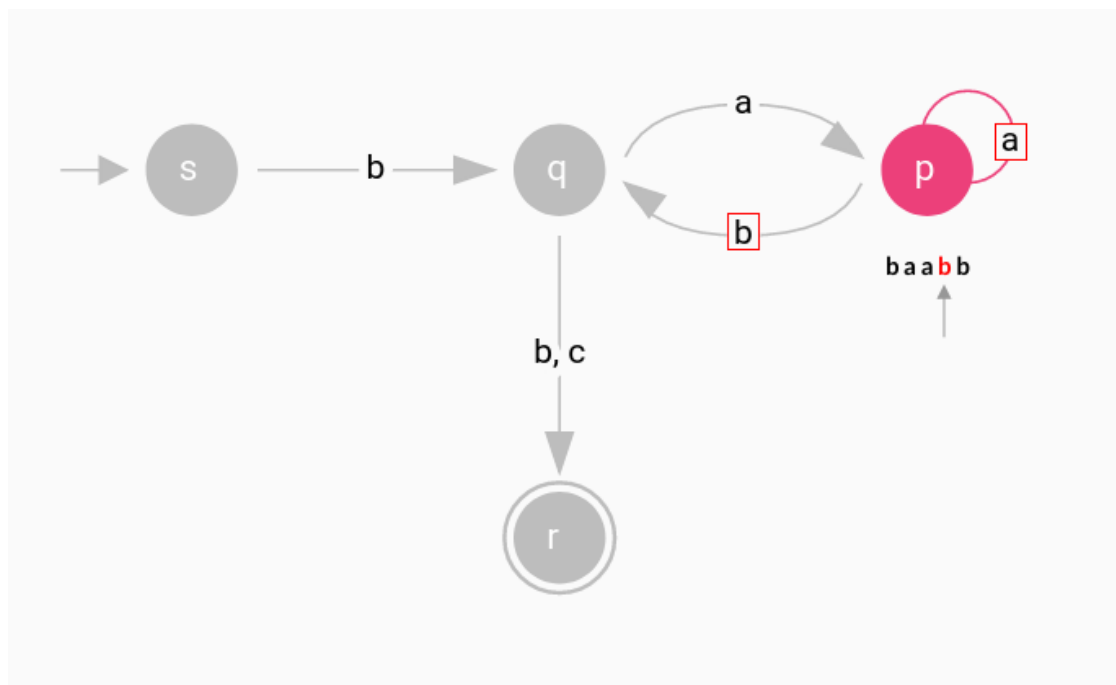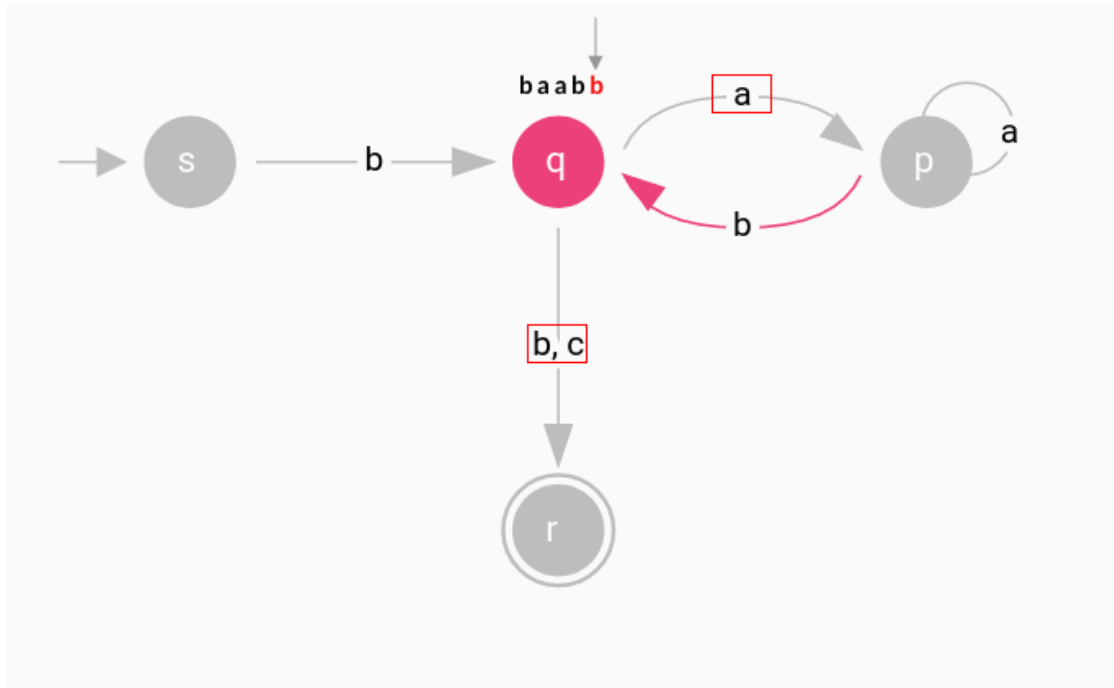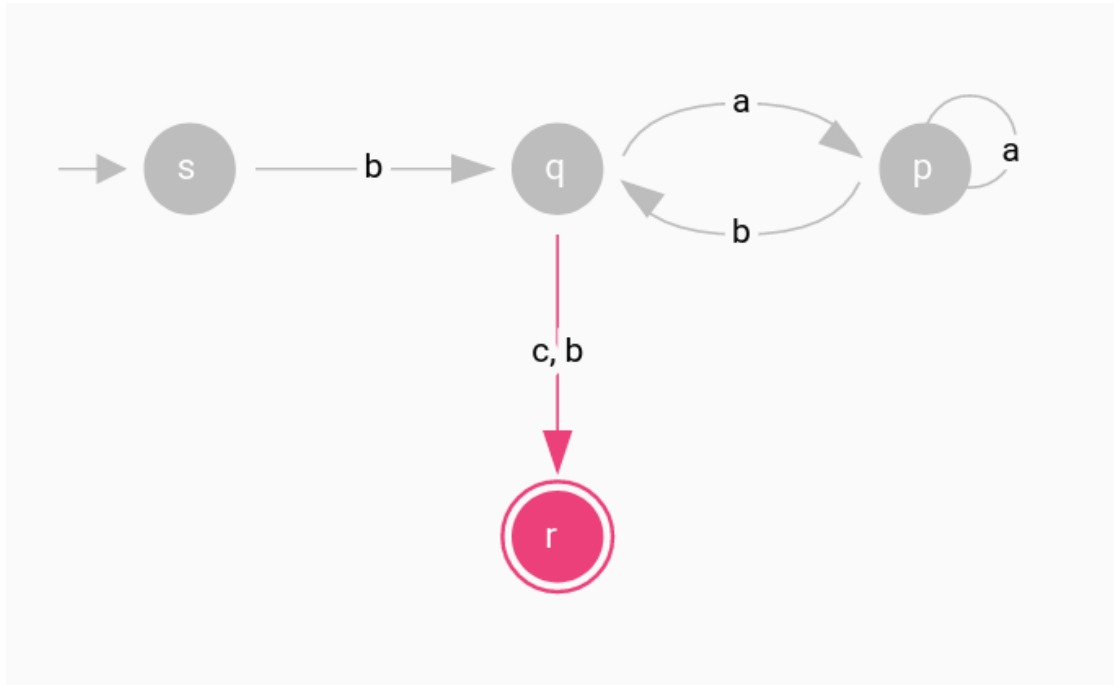3. The answer is yes, and it's **b**. So move to r

– Step 6

1. Next symbol of the input **baabb** is **b** and the current state is r.
2. Ask, if it satiesfies the accepting or final state (i.e, has the end of string been reached?). If so, the output is accept. Otherwise, it's reject.

- Formal Languages

  - is a <u>subset</u> of all possible words $\Sigma *$ formed by symbols of alphabet Σ.

    * $\Sigma *$ is set of all possible strings over the alphabet Σ.
    * i.e. $\Sigma = \{a, b\}$, $\Sigma * = \{a, b, aa, ab, ba, bb, aaa, aab, \cdots\}$

  - Example

    1. $L = \{w \mid w \text{ has at most seventeen 0's}\}$
    2. $L = \{w \mid w \text{ has equal number of 0's and 1's}\}$
    3. $L = \{x \in \{a, b\}^* \mid \text{the number of as in } x \text{ is even}\}$
       * $*$ in $\{a, b\}^*$ means all possible combinations
       * i.e. $\{a, b, aa, ab, ba, bb, aaa, baa, aba, \cdots\}$

- Tabular DFAs

  - Example

start

1

0

0

q_0 q_1 q_2 q_3

1   0

1

Σ
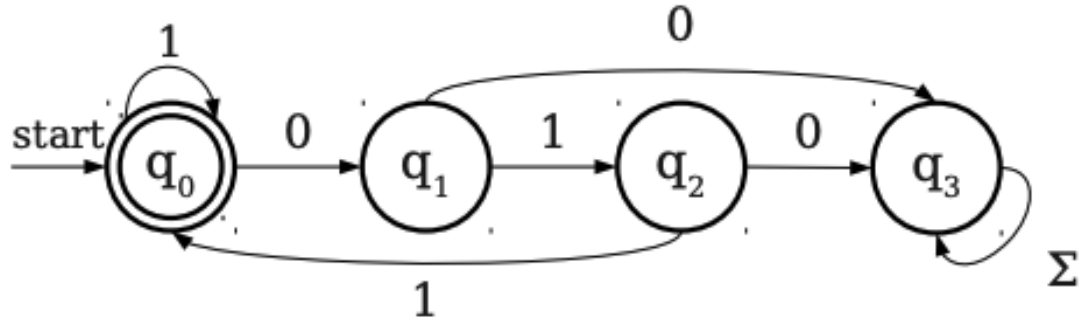
	0	1
$^*q_0$	q_1	q_0
$\delta = \quad q_1$	q_3	q_2
q_2	q_3	q_0
q_3	q_3	q_3

Note: * means it's an accepting state

Question 2

Partial Solution:

1. Prove that M_1 accepts L_1

   First, define Σ^* as the smallest set such that

   (a) $\varepsilon \in \Sigma^*$
   (b) $s \in \Sigma^* \Rightarrow sa \in \Sigma^* \land sb \in \Sigma^*$

   I will prove that M_1 accepts L_1.

   Define $P(s)$ as:

   $$P(s) : \delta^*(E, s) = \begin{cases} E & \text{if } s \text{ has an even number of } as \\ O & \text{if } s \text{ has an even number of } as \end{cases} \tag{1}$$

   I will prove $\forall s \in \Sigma^*$, $P(s)$ by structural induction.

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1. Basis Case

$|\varepsilon| = 0$, an even number, and $\delta^*(E, \varepsilon) = E$ so the implication in the first line of the invariant is true in this case. Also since $|\varepsilon|$ is not odd, the implication in the second line of the invariant is vacuously true. So $P(\varepsilon)$ holds.

2. Inductive Step

Let $s \in \Sigma^*$ and assume $P(s)$. I will show that $P(sa)$ and $P(sb)$ follow. There are two cases to consider:

1. Case sa

   Then,

$$\delta^*(E, sa) = \delta(\delta^*(E, s), a) = \begin{cases} \delta(E, a) & \text{if } s \text{ has even number of } as \\ \delta(O, a) & \text{if } s \text{ has odd number of } as \end{cases} \qquad [\text{By } P(s)]$$

$$(2)$$

$$= \begin{cases} O & \text{if } sa \text{ has odd number of } as \\ E & \text{if } sa \text{ has even number of } as \end{cases} \qquad [\text{One more } a]$$

$$(3)$$

2. Case sb (Let's first start with this)

   Then,

$$\delta^*(E, sb) = \delta(\delta^*(E, s), b) = \begin{cases} \delta(E, b) & \text{if } s \text{ has even number of } as \\ \delta(O, b) & \text{if } s \text{ has odd number of } as \end{cases} \qquad [\text{By } P(s)]$$

$$(4)$$

$$= \begin{cases} E & \text{if } sb \text{ has odd number of } as \\ O & \text{if } sb \text{ has even number of } as \end{cases} \qquad [\text{One more } b]$$

$$(5)$$

So $P(sa)$ and $P(sb)$ follow.

The first line of the invariant ensures that all strings with an even number as are accepted. The contrapositive of the second line of the invariant ensures that any string that does not drive the machine to state O does not have an odd number of as, in other words all strings that drive the machine to state E have an even number of as. So M_1 accepts L_1.

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2. Prove that M_2 accepts L_2

Define $P(s)$ as:

$$P(s) : \delta^*(0, s) = \begin{cases} 0 & \text{if } |s| \equiv 0 \mod (3) \\ 1 & \text{if } |s| \equiv 1 \mod (3) \\ 2 & \text{if } |s| \equiv 2 \mod (3) \end{cases} \tag{6}$$

I prove $\forall s \in \Sigma^*$, $P(s)$ by structural induction.

1. Basis Case (Let's give this a shot)

I need to prove $P(\varepsilon)$ holds.

Let $\varepsilon \in \Sigma^*$.

Then, since $|\varepsilon| = 0$, $\delta^*(0, \varepsilon) = 0$.

So, the implication of the first line is true in this case.

Also, since $|\varepsilon| \not\equiv 1 \mod 3$, the implication of second line is vacuously true in this case.

Also, since $|\varepsilon| \not\equiv 2 \mod 3$, the implication of third line is vacuously true in this case.

Thus, $P(\varepsilon)$ holds.

2. Inductive Step (Let's also try this)

Let $s \in \Sigma^*$ and assume $P(s)$. Let $@ \in \{a, b\}$. I will show $P(s@)$ follows.

Starting from $\delta^*(0, s@)$, we have

$$\delta^*(0, s@) = \begin{cases} \delta(0, @) & \text{if } |s| \equiv 0 \mod 3 \\ \delta(1, @) & \text{if } |s| \equiv 1 \mod 3 \\ \delta(2, @) & \text{if } |s| \equiv 2 \mod 3 \end{cases} \qquad \text{[By } P(s)\text{]}$$

(7)

$$= \begin{cases} 0 & \text{if } |s| \equiv 0 \mod 3 \\ 1 & \text{if } |s| \equiv 1 \mod 3 \\ 2 & \text{if } |s| \equiv 2 \mod 3 \end{cases} \qquad \text{[Since adding a char doesn't change inv. cond.]}$$

(8)

(9)

Thus, $P(s@)$ follows.

The first line of the invariant ensures that all strings with length equivalent to 0 mod 3 or multiples of 3 are accepted. The contrapositive of the second line of the invariant ensures that any string that does not drive the machine to state 1 does not have length equivalent to 1 mod 3. The contrapositive of the third line of the invariant ensures that any string that does not drive the machine to state 2 does not have length equivalent to 2 mod 3.

So, M_2 accepts L_2.

3. Prove that $M_{1\wedge2}$ accepts $L_1 \cap L_2$

Denote the states for M_1 as Q_1, the states for M_2 as Q_2, their respective transition functions as δ_1 and δ_2, and the transition function for $M_{1\wedge2}$ and the transition function for $M_{1\wedge2}$ as $\delta_{1\wedge2}$. Inspection of $\delta_{1\wedge2}$ shows that if $(q_1, q_2, c) \in Q_1 \times Q_2 \times \Sigma$, then $\delta_{1\wedge2}((q_1, q_2), c) = (\delta_1(q_1, c), \delta_2(q_2, c))$. Thus, the following invariant follows by simplying taking conjunctions of the invariants of the component machines, for any $s \in \Sigma^*$.

$$P(s) : \delta^*((E, 0), s) = \begin{cases} (E, 0) & \text{if } s \text{ has an even number of } as \wedge |s| \equiv 0 \mod 3 \\ (E, 1) & \text{if } s \text{ has an even number of } as \wedge |s| \equiv 1 \mod 3 \\ (E, 2) & \text{if } s \text{ has an even number of } as \wedge |s| \equiv 2 \mod 3 \\ (O, 0) & \text{if } s \text{ has an even number of } as \wedge |s| \equiv 0 \mod 3 \\ (O, 1) & \text{if } s \text{ has an even number of } as \wedge |s| \equiv 1 \mod 3 \\ (O, 2) & \text{if } s \text{ has an even number of } as \wedge |s| \equiv 2 \mod 3 \end{cases}$$

(10)

The implication on the first line ensures that all strings with an even number of as and a length that is a multiple of 3 end up in state $(e, 0)$. The contrapositive of the implications

on the other lines ensure that nay string that does not derive the machine to one of those 5 states must have an even number of as and a length that is a multiple of 3. Hence $M_{1\wedge2}$ accepts $L_1 \cap L_2$.

Notes:

- Hurrrayyyyy!!!!!!