CSC236 Worksheet 4 Solution

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Question 1

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Notes:

- Repeated Subtitution:
 - Is a technique used to find a closed form formula
 - closed form formula is a simple formula that allows evaluation of T(n) without the need to evaluate, say T(n/2)

i.e. from

$$T(n) = \begin{cases} c & \text{if } n = 1\\ 2T(n/2) + dn & \text{if } n > 1 \end{cases}$$
 (1)

to

$$T(n) = cn + dn \log_2 n$$

Example:

Consider the recurrence

$$T(n) = \begin{cases} c & \text{if } n = 1\\ 2T(n/2) + dn & \text{if } n > 1 \end{cases}$$
 (2)

Find closed form formula for T(n), where n is an arbitrary power of 2. That is $\exists k \in \mathbb{N}, n = 2^k$.

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Let n \in \mathbb{N} and assume that \exists k \in \mathbb{N}^+, n = 2^k, so k = \log_2 n.
Then,
   T(n) = 2T(n/2) + dn
                                                                                     [By 1]
                                                                                                  (3)
          = 2\Big(2T(n/2^2) + dn/2\Big) + dn
                                               [By subtituting n/2 for n in 1]
                                                                                                  (4)
          = 2^{2}T(n/2^{2}) + 2dn
                                                                                                  (5)
          =2^{2}\left(2T(n/2^{3})+dn/2^{2}\right)+2dn \quad \text{[By subtituting } n/2^{2} \text{ for } n \text{ in } 1\text{]}
                                                                                                  (6)
          = 2^3 T(n/2^3) + 3dn
                                                     [By subtituting n/2^2 for n in 1]
                                                                                                  (7)
                                                                                                  (8)
          =2^kT(n/2^k)+kdn
                                                                  [After k applications]
                                                                                                  (9)
          = 2^{\log_2 n} T(n/2^{\log_2 n}) + (\log_2 n) dn
                                                              [By replacing k = \log_2 n]
                                                                                                 (10)
          = nT(1) + (\log_2 n)dn
                                                                                                 (11)
          = cn + (\log_2 n) dn
                                                                                                (12)
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Question 2