

# CSC343 Worksheet 4 Solution

June 18, 2020

1. a)  $[(1, 0, 1), (5, 4, 9), (1, 0, 1), (6, 4, 16), (7, 9, 16)]$
- b)  $[(1, 0), (3, 3), (3, 4), (4, 3), (1, 1), (4, 3)]$
- c)  $[(0, 1), (0, 1), (2, 3), (2, 4), (3, 4)]$

## Notes:

- $\tau_L(R)$  sorts tuples in order indicated by  $L$ .
  - e.g.

$\tau_{C,B}(R)$  in  $R(A, B, C)$  orders the tuples of  $R$  by their values of  $C$ , and tuples with the same  $C$ -value are ordered by their  $B$  value.

- d)  $[(0, 1), (0, 2), (2, 4), (2, 5), (3, 4), (3, 4)]$
- e)  $[(0, 1), (2, 4), (2, 5), (3, 4), (0, 2)]$

## Notes:

- $\delta(R)$  converts a bag into a set
  - e.g.

Let  $R = [(1, 2), (3, 4), (1, 2), (1, 2)]$

$\delta(R(A, B)) = [(1, 2), (3, 4)]$

- f)  $[(0, 2), (2, 7), (3, 4)]$

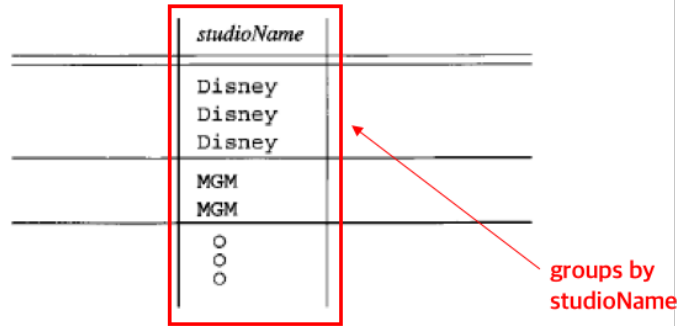
## Notes:

- $\gamma_L(R)$  is an operator that groups a relation and/or aggregate some columns.
  - $L$  in  $\gamma_L(R)$  is either
    1. **Grouping attribute** or an attribute by which  $R$  will be grouped.

2. **Aggregated attribute** or an attribute where an aggregation operator is applied to.

**Example:**

$\gamma_{starName, MIN(year) \rightarrow minYear, COUNT(title) \rightarrow ctTitle} (StarsIn)$



| studioName |
|------------|
| Disney     |
| Disney     |
| Disney     |
| MGM        |
| MGM        |
| ○          |
| ○          |
| ○          |

Figure 5.4: A relation with imaginary division into groups

- g)  $[(0, 1.5), (2, 4.5), (3, 4)]$   
 h)  $[(0, 1), (0, 1), (2, 3), (2, 4), (3, 4)]$   
 i)  $\gamma_{A, MAX(C)}([(2, 3, 4), (2, 3, 4)]) \rightarrow [(2, 4)]$   
 j)  $[(0, 1, \perp), (2, 3, 4), (2, 3, 4), (0, 1, \perp), (2, 4, \perp), (3, 4, \perp)]$

**Notes:**

- $\bowtie$  is an outerjoin operator
  - $\bowtie_L$  means Natural Left Outer Join
  - $\bowtie_R$  means Natural Right Outer Join
  - $\bowtie$  means Natural Full Outer Join
  - $\perp$  means null
- e.g.  $U \bowtie V$

| <i>A</i> | <i>B</i> | <i>C</i> |
|----------|----------|----------|
| 1        | 2        | 3        |
| 4        | 5        | 6        |
| 7        | 8        | 9        |

(a) Relation  $U$ 

| <i>B</i> | <i>C</i> | <i>D</i> |
|----------|----------|----------|
| 2        | 3        | 10       |
| 2        | 3        | 11       |
| 6        | 7        | 12       |

(b) Relation  $V$ 

| <i>A</i> | <i>B</i> | <i>C</i> | <i>D</i> |
|----------|----------|----------|----------|
| 1        | 2        | 3        | 10       |
| 1        | 2        | 3        | 11       |
| 4        | 5        | 6        | $\perp$  |
| 7        | 8        | 9        | $\perp$  |
| $\perp$  | 6        | 7        | 12       |

(c) Result  $U \bowtie V$ 

- k)  $[(\perp, 0, 1), (\perp, 2, 4), (\perp, 2, 5), (2, 3, 4), (\perp, 0, 2), (2, 3, 4)]$   
 l)  $[(0, 1, \perp), (2, 3, 4), (2, 3, 4), (0, 1, \perp), (2, 4, \perp), (3, 4, \perp),$   
 $(\perp, 0, 1), (\perp, 2, 4), (\perp, 2, 5), (2, 3, 4), (\perp, 0, 2), (2, 3, 4)]$   
 m)  $(0, 1) : \{(2, 4), (2, 5), (3, 4), (3, 4)\}$

But,  $\{(2, 3), (2, 4), (3, 4)\}$  from  $R$  and  $\{(0, 1), (0, 2)\}$  in  $S$  don't match. So,

$[(0, 1, 2, 4), (0, 1, 2, 5), (0, 1, 3, 4), (0, 1, 3, 4), (0, 1, 2, 4), (0, 1, 2, 5), (0, 1, 3, 4), (0, 1, 3, 4),$   
 $(2, 3, \perp, \perp), (2, 4, \perp, \perp), (3, 4, \perp, \perp), (\perp, \perp, 0, 1), (\perp, \perp, 0, 2)]$

### Notes:

- $R \bowtie_C S$  is equivalent form of  $\sigma_C(R \times S)$  but instead of filtering, the unmatched tuples are filled with null.

2. a) SELECT model FROM PC WHERE speed > 3.0;  
 b) SELECT DISTINCT maker FROM Products NATURAL JOIN Laptops WHERE hd >= 100;

```

c)      SELECT model, price FROM (
2         (SELECT model, price FROM PC NATURAL JOIN Products)
3         UNION
4         (SELECT model, price FROM Laptop NATURAL JOIN Products)
5         UNION
6         (SELECT model, price FROM Printer INNER JOIN Products ON
Printer.model = Product.model)
7     );
8

```

d) SELECT model FROM Printer WHERE color;

```
e) (SELECT DISTINCT makers FROM Products WHERE type='laptops') -
2 (SELECT DISTINCT makers FROM Products WHERE type='pc');
3
```

```
f) SELECT hd FROM PC WHERE EXISTS (
2 SELECT hd, COUNT(model) FROM PC GROUP BY hd
3 HAVING COUNT(model) > 2
4 );
5
```

3. a) SELECT class, country FROM classes WHERE bore >= 16;

b) SELECT \* FROM Ships WHERE launched < 1921;

c) SELECT \* FROM Outcomes WHERE result='sunk';

d) SELECT \* FROM Classes NATURAL JOIN Ships WHERE displacement > 35000;

```
e) SELECT name, displacement, numGuns FROM Classes NATURAL JOIN (
2 SELECT * FROM Ships INNER JOIN Outcomes ON Ships.name =
Outcome.ship
3 );
4
```

```
f) (SELECT name FROM Ships)
2 UNION
3 (SELECT ship AS name FROM Outcomes);
4
```

```
g) SELECT class, COUNT(class) FROM Ships
2 GROUP BY Class
3 HAVING COUNT(class) = 1;
4
```

```
h) (SELECT countries FROM Classes WHERE type='bb')
2 INTERSECT
3 (SELECT countries FROM Classes WHERE type='bc');
4
```

i) Current attempt:

```
1 (SELECT Table1.name FROM Outcomes AS Table1 INNER JOIN Ships ON
Outcome.ship = Ships.name)
2
```

Took too much time. Omitted for now.

4. a) SELECT AVG(speed) FROM PC;

b) SELECT AVG(speed) FROM Laptop HAVING price > 1000;

c) SELECT AVG(price) FROM PC NATURAL JOIN Product HAVING maker = 'A';

d)

```

1  SELECT AVG(price) FROM (
2      (SELECT model, price FROM PC NATURAL JOIN Product WHERE maker
3      = 'D')
4      UNION
5      (SELECT model, price FROM Laptop NATURAL JOIN Product WHERE
6      maker = 'D')
7  );

```

e) SELECT speed, AVG(price) FROM PC GROUP BY speed;

f) SELECT maker, AVG(screen) FROM Laptop NATURAL JOIN Product GROUP BY maker;

g)

```

1  SELECT maker, COUNT(model) FROM Products GROUP BY maker HAVING
2  COUNT(model) >= 3;
3

```

h)

```

1  SELECT maker, MAX(price) FROM PC
2  NATURAL JOIN Products GROUP BY maker;
3

```

i)

```

1  SELECT speed, AVG(price) FROM PC GROUP BY speed
2  HAVING speed > 2.0;
3

```

5. a) SELECT COUNT(class) FROM Classes;

**Correct Solution:**

```
SELECT COUNT(class) FROM Classes HAVING type='bb';
```

b) SELECT class, AVG(numGuns) FROM Classes GROUP BY class;

**Correct Solution:**

```
SELECT class, AVG(numGuns) FROM Classes GROUP BY class
HAVING type='bb';
```

c) SELECT AVG(numGuns) FROM Classes HAVING type='bb';

d) SELECT class, MIN(launched) FROM Ships GROUP BY class;

e)

```

1  SELECT class, COUNT(Ships.name) FROM Ships INNER JOIN Outcomes
2  GROUP BY class ON Ships.name = Outcomes.ship
3  HAVING result='sunk';
4

```

6.

```

1  SELECT name, MIN(movieYear) FROM MovieStar INNER JOIN StarsIn ON
2  StarsIn.starName = MovieStar.name GROUP BY name HAVING COUNT(name)
3  >= 3;

```

7. Yes. It is possible

- Rename aggregate columns using  $\rho$
- Use  $\sigma$  around  $\rho(\gamma(\dots))$

8. a) INSERT INTO PC(model, speed, ram, hd, price) VALUES (1100, 3.2, 1024, 180, 2499);

INSERT INTO Product(maker, model, type) VALUES ('C', 1100, 'pc');

### Notes:

- Insertion
  - **Syntax:** INSERT INTO R( $A_1, A_2, \dots, A_n$ ) VALUES ( $v_1, \dots, v_n$ )
  - Inserts a tuple with values into relation R
  - Values can be from select statement

```

1  \ Example 1
2  INSERT INTO StarsIn(movieTitle, movieYear, starName)
3  VALUES('The Maltese Fiction', 1942, 'Sydney Greenstreet');
4
5  \ Example 2
6  INSERT INTO Studio(name)
7  SELECT DISTINCT studioName
8  FROM Movies
9  WHERE studioName NOT IN
10     (SELECT name FROM Studio);
11

```

```

b) // SQL statement 1
2  INSERT INTO Laptop(model, speed, ram, hd, screen, price)
3  SELECT DISTINCT model + 110, speed, ram, hd, 17 AS screen,
4  price + 500
5  FROM PC;
6
7  // SQL statement 2
8  INSERT INTO Product(maker, model, type)
9  SELECT DISTINCT maker, model + 110, 'laptop' AS type
10 FROM Product WHERE type='pc';

```

```

c) DELETE FROM PC
2  WHERE hd < 100;
3

```

### Notes:

- Deletion

- **Syntax:** DELETE FROM  $R$  WHERE  $\langle \text{Condition} \rangle$ ;

**Example:**

```

1  DELETE FROM StarsIn
2  WHERE movieTitle = 'The Maltest Falcon' AND
3         movieYeeear = 1942 AND
4         starName = 'Sydney Greenstreet';
5

```

d)

```

1  DELETE FROM Product
2  WHERE maker = ALL(
3      (SELECT maker FROM Product WHERE type='laptop')-
4      (
5          (SELECT maker FROM PRODUCT WHERE type='laptop')
6          INTERSECT
7          (SELECT maker FROM PRODUCT WHERE type='printer')
8      )
9  )
10

```

e)

```

1  UPDATE Product
2  SET maker = 'A'
3  WHERE maker = 'B';
4

```

**Notes:**

- UPDATE

- **Syntax:** UPDATE  $R$  SET  $\langle \text{new-value assignments} \rangle$  WHERE  $\langle \text{condition} \rangle$ ;
- updates tuples in database

**Example:**

```

1  UPDATE MovieExec
2  SET name = 'Pres.' || name
3  WHERE cert# IN (SELECT presC# FROM Studio);
4

```

f) UPDATE PC SET ram = 2 \* ram, hd = 60 + hd;