# CSC343 Worksheet 6 Solution

June 22, 2020

## 1. Exercise 6.6.1:

```
a) SET TRANSACTION READONLY;
BEGIN TRANSACTION;
SELECT model, price FROM PC
WHERE speed = speed AND
ram=ram
COMMIT;
```

## Notes:

- Transactions
  - is a collection of one or more operations that must be executed atomically
  - COMMIT causes the transaction to end successfully
  - ROLLBACK causes the transaction to abort. Any changes are undone
  - SET TRANSACTION READ ONLY
    - \* tells the database that it will not be modified
    - \* Must be declared before transaction
    - \* Is useful when one user is running multiple queries while other is updating the same table

## Example:

```
BEGIN TRANSACTION;

UPDATE accounts

SET balance = balance - 1000

WHERE account_no = 100;

UPDATE accounts

SET balance = balance + 1000

WHERE account_no = 200;

INSERT INTO account_changes(account_no,flag,amount,changed_at)
```

```
13
               COMMIT;
        14
        15
               // Example - SET TRANSACTION READONLY
        16
               SET TRANSACTION READONLY;
        17
               BEGIN TRANSACTION;
        18
        19
               COMMIT;
        20
        21
b)
       BEGIN TRANSACTION;
       DELETE FROM PC
 2
       WHERE model = < model number >
 3
 4
       DELETE FROM Product
 5
       WHERE model = < model number >
 6
       COMMIT;
 8
 9
c)
       BEGIN TRANSACTION;
       UPDATE PC
 3
       SET price=price - 100
 4
       WHERE model = < model number >
 5
 6
       COMMIT;
d)
       BEGIN TRANSACTION;
 2
       IF (<model> IN (
 3
            SELECT <model > FROM Product
            NATURAL JOIN PC)
 5
 6
           PRINT 'Error occured';
 7
       ELSE
 8
           INSERT INTO PC
 9
            VALUES (<model>, <speed>, <ram>, <hd>, <price>)
10
11
           INSERT INTO Product
12
            VALUES (<maker>, <model>, <type>)
13
       COMMIT;
14
15
```

VALUES(100, '-', 1000, datetime('now'));

## 2. Exercise **6.6.2**:

For all cases, when system crashes, the operations in transaction are aborted and database is reverted back to pre-transaction state.

#### 3. Exercise 6.6.3:

The following would be observed

- Reading data modified by another transaction (Dirty Read)
- Repeated retrieval of rows resulting in different values (Non-repeatible read)
- Insertion/deletion of data (Phantom)

## Notes:

- Dirty Reads
  - A dirty read occurs when a transaction is allowed to read data from a row that
    has been modified by another running transaction and not yet committed.



- Non-repeatible Reads
  - A non-repeatable read occurs when, during the course of a transaction, a row is retrieved twice and the values within the row differ between reads.



#### • Phantom Reads

 A phantom read occurs when, in the course of a transaction, new rows are added or removed by another transaction to the records being read.



#### • Isolation Levels

- SET TRANSACTION ISOLATION LEVEL READ UNCOMMITTED;
  - \* is the lowest isolation level
  - \* allows to read a transaction that's not yet committed
  - \* transactions are not isolated from each other
- SET TRANSACTION ISOLATION LEVEL READ COMMITTED:
  - \* Does not allow to read a transaction that's not yet committed
  - \* Prevents other transactions from reading, updating or deleting while commit
- SET TRANSACTION ISOLATION LEVEL REPEATIBLE READ;
  - \* Is the higher level of isolation
  - \* Guarentees everything of READ COMMITTED level
  - \* Can read unchanged data in subsequent reads
- SET TRANSACTION ISOLATIO LEVEL SERIALIZABLE
  - \* Is the highest level of isolation
  - \* Guarentees everything of READ REPEATIBLE READ;
  - \* No new data can be seen by a subsequent read.

Isolation Level	Dirty Reads	Nonrepeat- able Reads	Phantoms
Read Uncommitted	Allowed	Allowed	Allowed
Read Committed	Not Allowed	Allowed	Allowed
Repeatable Read	Not Allowed	Not Allowed	Allowed
Serializable	Not Allowed	Not Allowed	Not Allowed

## References:

\* Stack Overflow: Difference between 'read committed' and 'repeatable read', link

\* Wikipedia: Isolation (database systems), link

## 4. Exercise 8.1.1:

```
a) CREATE VIEW RichExec AS

SELECT * FROM MovieExec

WHERE netWorth >= 10000000;
```

## Notes:

- Virtual Views
  - Syntax: CREATE VIEW < view-name > AS < view-definition >
  - Contrasts to database that exists in physical storage
  - Exists in RAM
  - Is created using query
  - can be used like a relation

# Notes:

```
CREATE VIEW ParamountMovies AS

SELECT title, year

FROM Movies

WHERE studioName = 'Paramount';
```

```
SELECT * FROM MovieExec
NATURAL JOIN MovieStar;
```

## 5. Exericse 8.1.2:

```
a) SELECT name, gender FROM ExecutiveStar;
b) SELECT name FROM RichExec WHERE netWorth > 10000000;
```

```
C) SELECT name FROM StudioPres

NATURAL JOIN ExecutiveStar

WHERE netWorth > 50000000
```

#### 6. Exericse 8.2.1:

*RichExec* is updatable.

#### Notes:

- Updatable View Conditions
  - The WHERE cluase in CREATE VIEW must not be a subquery
  - The FROM clause has only one occurrence of R
  - The SELECT clause must include enough attributes
  - NOT NULL attributes must have default values
    - \* A solution to this is by including the attribute without default value in CREATE VIEW

# Example:

```
Movies(title, year, length, genre, studioName, producerC#)
Suppose studioName is NOT NULL but has no default value.
Then, a fix is:

CREATE VIEW Paramount AS
SELECT studioName, title, year
FROM Movies
WHERE studioName = 'Paramount';
```

#### 7. Exericse 8.2.2:

- a) No. It is not updatable. Since,
  - 1. studioName attribute in Movies is NOT NULL without default value

## Notes:

- Using Trigger in VIEW
  - Uses INSTEAD OF in place of BEFORE or AFTER
  - When event causes the trigger, the trigger is done instead of the event

## Example:

```
CREATE VIEW ParamountMovies AS

SELECT title, year
FROM Movies
WHERE studioName = 'paramount';

CREATE TRIGGER ParamountInsert
INSTEAD OF INSERT ON ParamountMovies
REFERENCING NEW ROW AS NewRow
FOR EACH ROW
INSERT INTO Movies(title, year, studioName)
VALUES(NewRow.title, NewRow.year, 'Paramount');
```

```
C) CREATE TRIGGER DisneyComediesInsert

INSTEAD OF INSERT ON DisneyComedies

REFERENCING

NEW ROW AS NewTuple

OLD ROW AS OldTuple

FOR EACH ROW

UPDATE Movies

SET length=NewTuple.length

WHERE title=OldTuple.title AND year=OldTuple.year;
```

#### 8. Exercise **8.2.3**

a) No. the view is not updatable. Because for it to be updatable, only one relation must exist in FROM

```
b)
       CREATE TRIGGER NewPCInsert
       INSTEAD OF INSERT ON NewPC
       REFERENCING
 3
           NEW ROW AS NewTuple
           OLD ROW AS OldTuple
       FOR EACH ROW
 6
       INSERT INTO PC(model speed, ram, hd ,price)
       VALUES (NewTuple.model, NewTuple.speed, NewTuple.ram, NewTuple.hd
 8
      , NewTuple.price);
 9
       INSERT INTO Product(maker, model, type)
 10
       VALUES (NewTuple.maker, NewTuple.model, 'pc');
11
 12
```

```
c)
       CREATE TRIGGER NewPCUpdate
       INSTEAD OF INSERT ON NewPC
 2
       REFERENCING
 3
           NEW ROW AS NewTuple
 4
       FOR EACH ROW
 5
       UPDATE PC
 6
       SET model=NewTuple.model
 7
           speed=NewTuple.speed,
 8
           ram = NewTuple.ram ,
 9
10
           hd=NewTuple.hd,
           price=NewTuple.price;
11
12
       UPDATE Product
13
       SET maker=NewTuple.maker,
14
           model=NewTuple.model,
15
           type='pc';
16
```

```
Correct Solution:
    CREATE TRIGGER NewPCUpdate
    INSTEAD OF UPDATE ON NewPC
    REFERENCING
        NEW ROW AS NewTuple
    FOR EACH ROW
    UPDATE PC
    SET model=NewTuple.model
        speed=NewTuple.speed,
        ram = NewTuple.ram ,
        hd=NewTuple.hd,
        price=NewTuple.price;
    UPDATE Product
    SET maker=NewTuple.maker,
        model=NewTuple.model,
        type='pc';
```

```
d
       CREATE TRIGGER NewPCDelete
       INSTEAD OF DELETE ON NewPC
       REFERENCING
 3
           NEW ROW AS NewTuple
 4
       FOR EACH ROW
 5
       DELETE FROM PC
 6
       WHERE model=NewTuple.model;
 8
       DELETE FROM Product
 9
       WHERE model=NewTuple.model;
10
```

```
9. a) CREATE INDEX studioNameIndex Studio(name)
```

## Notes:

- Indexes
  - Syntax (Create Index):CREATE INDEX < index-name > R(< attributes >)
  - Syntax (Drop Index):DROP INDEX < index-name >
  - Used to find tuples in a very large database
    - \* Is efficient
  - Can be thought as (key, value) pair in a binary search tree
  - e.g. Declaring Index

```
CREATE INDEX KeyIndex ON Movies(title, year);
```

- e.g. Dropping index

```
CREATE INDEX KeyIndex ON Movies(title, year);
```

- CREATE INDEX movieKeyIndex Movies(genre, length)

## 10. Exercise 8.4.1:

Action	No Index	Star Index	Movie Index	Both Indexes
$\overline{Q_1}$	100	4	100	4
$\overline{Q_2}$	100	100	4	4
$\overline{I}$	2	4	4	6
Average	$2 + 100p_1 + 100p_2$	$4 + 96p_2$	$4 + 96p_1$	$6 - 2p_1 - 2p_2$

## Notes:

- Database Tuning
  - Index speeds up queries that can use it
  - Index should NOT be created when modifications are the frequent choice of action

## 11. Exercise **8.4.2**:

Omitted for the time being

#### 12. Exercise 8.5.1:

```
UPDATE MovieProd
SET name='New Name'
WHERE (title, year) IN

(
SELECT title, year FROM Movies
INNER JOIN MovieExecs
ON Movies.productC# = MovieExec.cert#
WHERE cert# = '4567'
);
```

## Notes:

- Materialized Views
  - Is also known as a summary
  - Is also known as black-box abstraction
  - Stores view in physical storage
  - Useful when storing expensive operation like AVG or COUNT

## 13. Exercise 8.5.3

The following modifications to base tables require the modification of the materialized view

- PC: Updates(model, speed, ram, hd, price), Delete
- Product: Updates(maker, model, type), Delete

## Notes:

• Materialized view of NewPC

```
CREATE MATIERLIZED VIEW NewPC AS

SELECT maker, model, speed, ram, hd, price

FROM Product, PC

WHERE Product.model = PC.model AND type = 'pc';
```

- 14. Exercise 8.5.3
- 15. Exercise 8.5.4