Worksheet 13 Review

March 31, 2020

Question 1

a. Since the loop starts from i = 0 and ends at i = n - 1. The loop has

$$n - 1 - 0 + 1 = n \tag{1}$$

iterations.

Since each iteration runs 5 steps, the loop has total cost of

$$5 \cdot n = 5n \tag{2}$$

steps.

Because we know i = 0 at line 2 has cost of 1, we can conclude that the algorithm has total cost of 5n + 1 steps.

Correct Solution:

Because we know the loop starts from i = 0 and ends at i = n - 1 with i increasing by 5 per iteration, we can conclude the loop has

$$\left\lceil \frac{n}{5} \right\rceil \tag{3}$$

iterations.

Since each iteration takes constant time, the loop has runtime of $\Theta(n)$

Notes:

- How does professor begin a proof after 'We will prove that...' or at the beginning of each case/parts?
- Noticed professor doesn't provide a detailed explanation for the number of iterations.
- Realized the goal of this problem is to determine the exact cost and runtime of each loop.

There are
$$\lceil \frac{n}{5} \rceil$$
 iterations. \cdots

b. Because we know the loop starts at i = 4 and ends at i = n - 1 with i increasing by 1 per iteration, we can conclude that the loop has

$$\lceil n - 14 + 1 \rceil = n - 4 \tag{1}$$

iterations.

Since each iteration takes a constant time, we can conclude that the loop has runtime of $\mathcal{O}(n)$.

Correct Solution:

Because we know the loop starts at i = 4 and ends at i = n - 1 with i increasing by 1 per iteration, we can conclude that the loop has **at** most

$$\lceil n - 14 + 1 \rceil = n - 4 \tag{1}$$

iterations.

Since each iteration takes a constant time, we can conclude that the loop has runtime of $\Theta(n)$.

Notes:

- Noticed professor doesn't count i = 4 or i = 0 in question 1.a to the total cost of algorithm. But in later parts of the question, the two are considered. I wonder if the line with constant runtime should always be accounted for, or if it can be ignored in certain circumstances. If latter, when can the constants be ignored?
- c. Since the loop starts at i = 0 and ends at i = n 1 with i increasing by $\frac{n}{10}$ per iteration, we can conclude that the loop has at most

$$\left\lceil \frac{(n-1-0+1)}{\frac{n}{10}} \right\rceil = \left\lceil \frac{n \cdot 10}{n} \right\rceil$$

$$= 10$$
(1)

iterations.

Since each iteration takes constant time, we can conclude that the loop has runtime of $\Theta(1)$.

d. For the case where $n^2 \leq 20$, it's omitted because we are assuming the value of n is asymptotically large.

For the case where $n^2 > 20$, because we know the loop runs from i = 20 to $i = n^2 - 1$ with i increasing by 3 per iteration, we can conclude that the loop has at most

$$\left[\frac{n^2 - 1 - 20 + 1}{3} \right] = \left[\frac{n^2 - 20}{3} \right] \tag{1}$$

iterations.

Since the loop takes constant time per iteration, the loop has total runtime of $\Theta(n^2)$.

e. Since we are considering asymptotic runtime or the case where n is very large, the case where $n^2 \leq 20$ will be omitted.

For the case where $n^2 > 20$, because we know the first loop runs from i = 20 and ends at $i = n^2 - 1$ with i increasing by 3 per iteration, we can conclude that the first loop runs

$$\left[\frac{n^2 - 1 - 20 + 1}{3}\right] = \left[\frac{n^2 - 20}{3}\right] \tag{1}$$

iterations

For the second loop, because we know the loop runs from j = 0 to j = n-1 with j increasing by 0.01 per iteration, we can conclude that the second loop runs at most

$$\left\lceil \frac{n-1-0+1}{\frac{1}{100}} \right\rceil = \lceil 100 \cdot n \rceil \tag{2}$$

$$= 100 \cdot n \tag{3}$$

iterations

Since each iteration takes a constant time in both of the loops, the total runtime of the loops in this algorithm is

$$\Theta\left(\left\lceil \frac{n^2 - 20}{3} \right\rceil + 100 \cdot n\right) = \Theta(n^2) \tag{4}$$

Correct Solution:

For the case where $n^2 \leq 20$, it's omitted because we are assuming the value of n is asymptotically large.

For the case where $n^2 > 20$, because we know the loop runs from i = 20 to $i = n^2 - 1$ with i increasing by 3 per iteration, we can conclude that the loop has at most

$$\left[\frac{n^2 - 1 - 20 + 1}{3} \right] = \left[\frac{n^2 - 20}{3} \right] \tag{1}$$

iterations.

Since the loop takes constant time per iteration, the loop has total runtime of $\Theta(n^2)$.

Since we are considering asymptotic runtime or the case where n is very large, the case where $n^2 \leq 20$ will be omitted.

For the case where $n^2 > 20$, because we know the first loop runs from i = 20 and ends at $i = n^2 - 1$ with i increasing by 3 per iteration, we can conclude that the first loop runs

$$\left[\frac{n^2 - 1 - 20 + 1}{3}\right] = \left[\frac{n^2 - 20}{3}\right] \tag{1}$$

iterations

Since the first loop takes a constant time per iteration, we can conclude the first loop has runtime of $\Theta(n^2)$.

For the second loop, because we know the loop runs from j=0 to j=n-1 with j increasing by 0.01 per iteration, we can conclude that the second loop runs at most

$$\left\lceil \frac{n-1-0+1}{\frac{1}{100}} \right\rceil = \left\lceil 100 \cdot n \right\rceil \tag{2}$$

$$= 100 \cdot n \tag{3}$$

iterations

Since the first loop takes a constant time per iteration, we can conclude the second loop has runtime of $\Theta(n)$.

Since $n \in \Theta(n^2)$, the total runtime of algorithm is $\Theta(n^2)$.

Question 2

Question 3

Question 4