

1. a) 1) 4 - inode blocks. 1 for the file c, and 3 for the directories /, a, b
 - 2) 3 - directory blocks - one for root /, one for a, the other for b
 - 3) 1 - single indirect block as far as we know. The file definitely has more than 12 blocks (# of data blocks pointed by direct pointers), but less than 1036 (# of data blocks pointed by direct pointers and single indirect pointers). We are reading block 1034.
 - 4) 1 - data block for file c
- b) All of the above

Notes

- **Inode**



- Is short form of **index node**
- Describes a file system object such as file or data
- Contains all information about a file/directory, including
 - * File Type,
 - * Size
 - * Number of blocks allocated to it
 - * Protection information
 - * Time information (e.g time created, time modified)
 - * Location of data blocks residing on disk

References

- 1) Wikipedia, Inode, link
- 2) Machanick, Philip. (2016). Teaching Operating Systems: Just Enough Abstraction. 642. 10.1007/978-3-319-47680-3_10., link

- c) Size, the location of data blocks that reside on disk

Notes

- I wonder what information about blocks inode has. Is it total number of blocks both inode and data, or just data?
- I struggled a bit on this one. I should find an easier way to remember which information inode has

d) Rough Work

• **Creash Scenarios**

- When only new data block is written to disk
 - * This is fine in system's point of view
 - * No inode points to it (it doesn't contain any information about file)
 - * No bitmap points to it
 - * Is as if write never occurred
- When only the updated inode is written to disk
 - * There is no bitmap that's pointing to it
 - * There is new inode where existing inode is
 - * The data block Db hasn't been created
 - * Reading data where Db is will return garbage data
 - * there is a term for this. Is called **File-System inconsistency**
- When only inode bitmap is written to disk
 - * inode block pointed by bitmap is assumed to be allocated
 - * But there is no desired inode where it's pointing
 - * This is another example of **File-System-Inconsistency**
 - * If left as is, then space cannot be used for future use (**inode leak**)
- When only data bitmap is written to disk
 - * data block pointed by bitmap is assumed to be allocated
 - * But there is no desired inode where it's pointing
 - * This is another example of **File-System-Inconsistency**
 - * If left as is, then space cannot be used for future use (**data leak**)

Notes

- I wonder how system call for reading file/directory works in UNIX. Does it check for bitmap?
- I wonder how system call for deleting file/directory works in UNIX
- I wonder how system call for creatubg file/directory works in UNIX
- **File API**
 - open (create)

* Is a system call

* **Syntax:**

```
int fd = open("foo", O_CREAT|O_WRONLY|O_TRUNC, S_IRUSR|S_IWUSR)
```

- O_CREAT - Creates file "foo" if does not exist
- O_WRONLY - Open file for writing only (default)
- O_TRUNC - Overwrites existing file **Need example/Clarification**
- Can have multiple flags

* Returns **file descriptor** or fd for short

- Is an integer
- Is used to access a file
- Is private per process
- Can be used to read() and write() files

Example

```
#include <fcntl.h>
...
int fd;
mode_t mode = S_IRUSR | S_IWUSR | S_IRGRP | S_IROTH;
char *filename = "/tmp/file";
...
fd = open(filename, O_WRONLY | O_CREAT | O_TRUNC, mode);
...
```

File can be read by owner

File can also be written by owner

File can also be read by group

File can also be read by others

Means

1. File is Writable AND
2. Create file if doesn't exist AND
3. Overwrite file if exists

– (read)

* Is a system call

* **Syntax:**

```
ssize_t read (int fd, void *buf, size_t count)
```

- fd - file descriptor (from open())
- buf - container for the read data
- count - number of bytes to read

* Returns number of bytes read, if successful

- * Returns 0 if is at, or past the end of file

Example

```
char buf[4096];
int fd = open("/a/b/c", 0); // open in read-only mode
lseek(fd, 1034*4096, 0);    // seek to position (1034*4096) from start of file
read(fd, buf, 4096);        // read 4k of data from file
```

System Calls	Return Code	Current Offset	
fd = open("file", O_RDONLY);	3	0	
read(fd, buffer, 100);	100	100	← read continues for each call
read(fd, buffer, 100);	100	200	
read(fd, buffer, 100);	100	300	
read(fd, buffer, 100);	0	300	← returns 0 if at end
read(fd, buffer, 100);	0	-	
close(fd);			

– write

- * Is a system call
- * Writes data out of a buffer
- * **Syntax:**

```
ssize_t write (int fd, const void * buf, size_t nbytes)
```

- fd - file descriptor
- buf - A pointer to a buffer to write to file
- nbytes - number of bytes to write. If smaller than buffer, the output is truncated

Example

```
#include <unistd.h>
#include <fcntl.h>

int main(void)
{
    int filedesc = open("testfile.txt", O_WRONLY | O_APPEND);

    if (filedesc < 0) {
        return -1;
    }

    if (write(filedesc, "This will be output to testfile.txt\n", 36) != 36) {
        write(2, "There was an error writing to testfile.txt\n", 43);
        return -1;
    }

    return 0;
}
```

– lseek

* Reads or write to a specific offset within a file

* **Syntax:**

```
off_t lseek (int fd, off_t offset, int whence)
```

· fd - file descriptor

· offset - the offset of pointer within file (in bytes)

· whence - the method of offset

SEEK_SET - offset from the start of file (absolute)

SEEK_CUR - offset from current location + offset bytes (relative)

SEEK_END - offset from the end of file

* Returns offset amount (in bytes) from the beginning of file

* Returns -1 if error

Example

System Calls	Return Code	Current Offset
fd = open("file", O_RDONLY);	3	0
lseek(fd, 200, SEEK_SET);	200	200
read(fd, buffer, 50);	50	250
close(fd);	0	-

move 200 bytes from the start of file

read 50 bytes

– rename

* Is a system call

* Changes the name of file

* **Syntax:** int rename(const char *old, const char *new)

· old - name of old file

· new - name of new file

* Returns 0 if successful

* Returns -1 if error

Example

```
int fd = open("foo.txt.tmp", O_WRONLY|O_CREAT|O_TRUNC,
              S_IRUSR|S_IWUSR);
write(fd, buffer, size); // write out new version of file
fsync(fd);
close(fd);
rename("foo.txt.tmp", "foo.txt");
```

e.g. "hello\n"

- stat
 - * displays metadata of a certain file stored in **inode**
 - * **Syntax:** `int stat(const char *path, struct stat *buf)`
 - path - file descriptor of file that's being inquired
 - buf - A stat structure where data about the file will be stored (see below)

```

struct stat {
    dev_t     st_dev;      // ID of device containing file
    ino_t     st_ino;      // inode number
    mode_t    st_mode;     // protection
    nlink_t   st_nlink;    // number of hard links
    uid_t     st_uid;      // user ID of owner
    gid_t     st_gid;      // group ID of owner
    dev_t     st_rdev;     // device ID (if special file)
    off_t     st_size;     // total size, in bytes
    blksize_t st_blksize;  // blocksize for filesystem I/O
    blkcnt_t  st_blocks;   // number of blocks allocated
    time_t    st_atime;    // time of last access
    time_t    st_mtime;    // time of last modification
    time_t    st_ctime;    // time of last status change
};

```

Figure 39.5: The **stat** structure.

Example

```

1  #include <unistd.h>
2  #include <stdio.h>
3  #include <sys/stat.h>
4  #include <sys/types.h>
5
6  int main(int argc, char **argv)
7  {
8      if(argc != 2)
9          return 1;
10
11     struct stat fileStat;
12     if(stat(argv[1], &fileStat) < 0)
13         return 1;
14
15     printf("Information for %s\n", argv[1]);
16     printf("-----\n");
17     printf("File Size: \t\t%d bytes\n", fileStat.st_size);
18     printf("Number of Links: \t%d\n", fileStat.st_nlink);
19     printf("File inode: \t\t%d\n", fileStat.st_ino);
20
21     printf("File Permissions: \t");
22     printf( (S_ISDIR(fileStat.st_mode)) ? "d" : "-");
23     printf( (fileStat.st_mode & S_IRUSR) ? "r" : "-");
24     printf( (fileStat.st_mode & S_IWUSR) ? "w" : "-");
25     printf( (fileStat.st_mode & S_IXUSR) ? "x" : "-");
26     printf( (fileStat.st_mode & S_IRGRP) ? "r" : "-");
27     printf( (fileStat.st_mode & S_IWGRP) ? "w" : "-");
28     printf( (fileStat.st_mode & S_IXGRP) ? "x" : "-");

```

```

29     printf( (fileStat.st_mode & S_IROTH) ? "r" : "-");
30     printf( (fileStat.st_mode & S_IWOTH) ? "w" : "-");
31     printf( (fileStat.st_mode & S_IXOTH) ? "x" : "-");
32     printf("\n\n");
33
34     printf("The file %s a symbolic link\n", (S_ISLNK(fileStat.
35 st_mode)) ? "is" : "is not");
36
37     return 0;
38 }

```

The result of above is:

```

$ ./testProgram testfile.sh

Information for testfile.sh
-----
File Size:          36 bytes
Number of Links:    1
File inode:         180055
File Permissions:   -rwxr-xr-x

The file is not a symbolic link

```

- `unlink` (remove)
 - Is a system call
 - Removes a file (including symbolic link) from the system
 - **Syntax:** `int unlink(const char *pathname)`
 * `pathname` - path to file
 - Returns 0 if successful
 - Returns -1 if error

Example

```

#include <unistd.h>

char *path = "/modules/pass1";
int  status;
...
status = unlink(path);

```

- `mkdir`
 - Is a system call
 - **Syntax:** `int mkdir(const char *path, mode_t mode)`
 * `path` - path of directory (including name)

- * mode - permission group
- Returns 0 if successful
- Returns -1 if error
- directories can never be written directly
 - * directory is in format called **File System Metadata**
 - * directory can only be updated directly
- creates two directories on creation . (current) and .. (parent)

References

- 1) codewiki, stat, link
- 2) The Open Group Base Specification, unlink, link