CSC 369 Worksheet 5 Solution

August 19, 2020

1. I need to run randomly-generated problems with two jobs and two queues using file mlfq.py with I/O turned off, and compute the MLFQ execution trace for each.

Using the command ./mlfq.py -s 1 -m 10 -n 2 -j 2 -M 0, we have

```
Job List:
Job 0: startTime 0 - runTime 2 - ioFreq 0
Job 1: startTime 0 - runTime 7 - ioFreq 0
```

with

- allotments for queue 1 is 1
- quantum length for queue 1 is 10
- allotments for queue 0 is 1
- quantum length for queue 0 is 10
- no priority boost

the exeuction trace is:

```
[time 0] Job begins by job 0
[time 0] Job begins by job 1
[time 0] Run job 0 at priority 1 [Ticks 9, Allotment 1, Time 1 (of 2)]

[time 1] Run job 0 at priority 1 [Ticks 8, Allotment 1, Time 0 (of 2)]

[time 2] Finished JOB 0
[time 2] Run job 1 at priority 1 [Ticks 9, Allotment 1, Time 6 (of 7)]
```

```
[time 3] Run job 1 at priority 1 [Ticks 8, Allotment 1, Time 5 (of 7)]

[time 4] Run job 1 at priority 1 [Ticks 7, Allotment 1, Time 4 (of 7)]

[time 5] Run job 1 at priority 1 [Ticks 6, Allotment 1, Time 3 (of 7)]

[time 6] Run job 1 at priority 1 [Ticks 5, Allotment 1, Time 2 (of 7)]

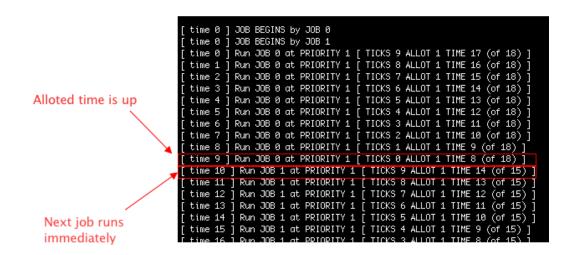
[time 7] Run job 1 at priority 1 [Ticks 4, Allotment 1, Time 1 (of 7)]

[time 8] Run job 1 at priority 1 [Ticks 3, Allotment 1, Time 0 (of 7)]

[time 9] Finished JOB 1
```

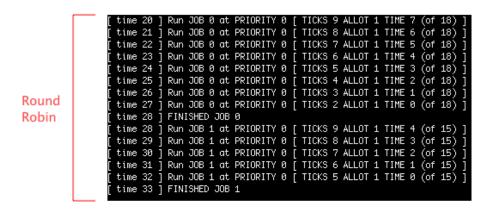
Notes

• Learned that when alloted time is up, the next job starts immediately (./mlfq.py -s 20 -m 20 -n 2 -j 2 -M 0 -c)



• Learned that when all jobs are at the bottom, without priority boost, jobs finishes by round robin

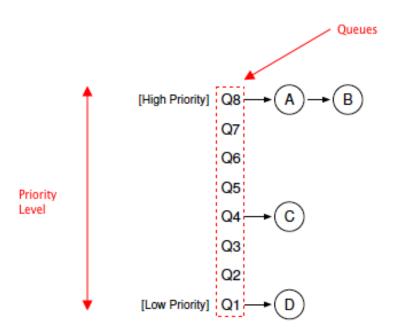
```
(./mlfq.py -s 20 -m 20 -n 2 -j 2 -M 0 -c)
```



- Learned that notification and subsequent job execution happen at the same time.
- The reason why round robin doesn't occur despite Priority(A) = Priority(B) is because allotment of queue is 1 (i.e. only one job can be in a queue)
- allotment means the amount of something allocated to a person/object (i.e. the size of queue)
- -m 10 sets the maximum runtime of a job to 10
- -M 0 turns off I/O in mlfq.py
- -n 2 sets number of queues to 2
- -j 2 sets number of jobs to 2
- Multi-level Feeback Queue (MLFQ):
 - Is one of the most well-known approaches to scheduling
 - Does two things:
 - a) Optimizes turnaround time
 - b) Minimizes response time
 - Uses **priority level** and **Queues** to achieve it's goal

• MLFQ Basic Rules:

- Jobs on same queue \rightarrow Same priority
- Rule 1: If Priority(A) > Priority(B), A runs (B doesn't)
- Rule 2: If Priority(A) = Priority(B), A & B run in RR

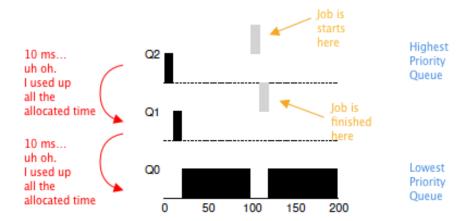


• Attemp #1: How to Change Priority

- Rule 3: When a job enters the system, it is placed at the <u>highest</u> priority (the topmost queue)
- Rule 4a: If a job uses up an entire time slice while running, its' priority is reduced (i.e. it moves down on queue).
- Rule 4b: If a job gives up the CPU before the time slice is up, it stays at the same priority level (e.g I/O Operation)
 - * Means that the shifting down of priority level only depends on CPU time

Example (Along Came a Short Job):

- 1) A job A enters system
- 2) Job is placed on highest Queue Q_2
- 3) After time-slice (e.g. 10 ms) in Q_2 , A is placed on lower queue Q_1
- 4) After time-slice in Q_1 , A is placed in lowest priority queue Q_0



• Attemp #2: The Priority Boost

- Rule 5: After some time period S, move all the jobs in the system to the topmost queue.
 - * This is to prevent starvation (i.e. a job never being run)
- Attempt #3: Better Accounting (Fix of Attempt # 1)
 - Is to prevent programmers from gaming (i.e tricking) the CPU so all programs get a fair share of allotment time
 - Rule 4: Once a job uses up its time allotment at a given level (regardless of how
 many times it has given up the CPU), its priority is reduced (it moves down one
 queue).
- 2. I need to run the scheduler (mlfq.py) to reproduce each of the examples in the chapter.
 - Example 1: A Single Long-Running Job

Here, the example has

- 3 queues
- 1 job
- 10ms as quantum length for queue 1
- 10ms as quantum length for queue 2
- 10ms as quantum length for queue 3
- 200ms as run time for job 1
- no priority boost

Combining together we have

$$./mlfq.py -1 0,200,0 -n 3 -j 1 -c$$

• Example 2: Along Came A Short Job

Here, the example has

- 3 queues
- -2 jobs
- 10ms as quantum length for queue 1
- 10ms as quantum length for queue 2
- 10ms as quantum length for queue 3
- 180ms as run time for job 1
- 20ms as run time for job 2
- 0ms as the starting time for job 1
- -100ms as the starting time for job 2
- no I/O operations for job 1
- no I/O operations for job 2
- no priority boost

Combining together we have

$$./mlfq.py -l 0,180,0:100,20,0 -n 3 -j 2 -c$$

• Example 3: What About I/O

Here, the example has

- 3 queues
- -2 jobs
- 10ms as quantum length for queue 1
- 10ms as quantum length for queue 2
- 10ms as quantum length for queue 3
- Oms as the starting time for job 1
- 185ms as run time for job 1
- no I/O operations for job 1
- 50ms as the starting time for job 2
- 15ms as run time for job 2
- 1ms as the frequency of I/O request (Start immediately)
- 9ms as the I/O time for job 2
- no priority boost
- Uses older rules (Rule 4.a, Rule 4.b)

Combining together we have

$$./mlfq.py -1 0,185,0:50,15,1 -i 9 -n 3 -j 2 -S -c$$

3. I need to answer the question how would .mlfq.py be configured to behave just like round-robin scheduler.

The round-robin scheduler occurs when jobs are at bottom most priority with no priority boost and I/O operations (i.e. jobs are at same priority until completion).

So, the configuration is:

$$./mlfq.py -l 0,100,0:0,100,0 -n 1 -j 2 -c$$

4. I need to craft a workload with two jobs and scheduler parameters so that one job takes of the older rules and obtain 99% of the CPU over a particular time interval.

The following configuration allows job 1 to occupy most of CPU time.

$$./mlfq.py -l 0,100,9:0,100,0 -i 1 -n 3 -j 2 -S -c$$

This occurs once job 2 moves down to lower priority