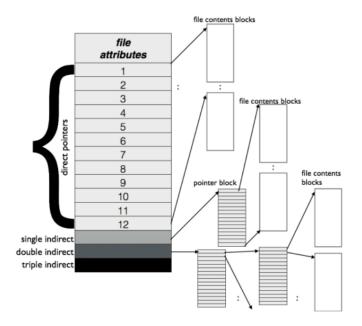
- 1. a) 1) 4 inode blocks. 1 for the file c, and 3 for the directdories /, a, b
 - 2) 3 directory blocks one for root /, one for a, the other for b
 - 3) 1 single indirect block as far as we know. The file definitely has more than 12 blocks (# of data blocks pointed by direct pounters), but less than 1036 (# of data blocks pointed by direct pointers and single indirect pointers). We are reading block 1034.
 - 4) 1 data block for file c
 - b) All of the above

Notes

• Inode



- Is short form of index node
- Describes a file system object such as file or data
- Contains all information about a file/directory, including
 - * File Type,
 - * Size
 - * Number of blocks allocated to it
 - * Protection information
 - * Time information (e.g time created, time modified)
 - * Location of data blocks residing on disk

References

- 1) Wikipedia, Inode, link
- 2) Machanick, Philip. (2016). Teaching Operating Systems: Just Enough Abstraction. 642. 10.1007/978-3-319-47680-3_10., link

c) Size, the location of data blocks that reside on disk

Notes

- I wonder what information about blocks inode has. Is it total number of blocks both inode and data, or just data?
- I struggled a bit on this one. I should find an easier way to remember which information inode has

d) Rough Work

• Creash Scenarios

- When only new data block is written to disk
 - * This is fine in system's point of view
 - * No inode points to it (it doesn't contain any information about file)
 - * No bitmap points to it
 - * Is as if write never occured
- When only the updated inode is written to disk
 - * There is no bitmap that's pointing to it
 - * There is new inode where existing inode is
 - * The data block Db hasn't been created
 - * Reading data where Db is will return garbage data
 - * there is a term for this. Is called **File-System inconsistency**
- When only inode bitmap is written to disk
 - * inode block pointed by bitmap is assumed to be allocated
 - * But there is no desired inode where it's pointing
 - * This is another example of File-System-Inconsistency
 - * If left as is, then space cannot be used for future use (inode leak)
- When only data bitmap is written to disk
 - * data block pointed by bitmap is assumed to be allocated
 - * But there is no desired inode where it's pointing
 - * This is another example of File-System-Inconsistency
 - * If left as is, then space cannot be used for future use (data leak)

Notes

- I wonder how system call for reading file/directory works in UNIX. Does it check for bitmap?
- I wonder how system call for deleting file/directory works in UNIX
- I wonder how system call for creatubg file/directory works in UNIX

• File API

- open (create)

- * Is a system call
- * Syntax:

```
int fd = open("foo". O_CREAT|O_WRONLY|O_TRUNC, S_IRUSR|S_IWUSR)
```

- · O_CREAT Creates file "foo" if does not exist
- · O_WRONLY Open file for writing only (default)
- · O_TRUNC Overwrites existing file Need example/Clarification
- · Can have multiple flags
- * Returns file descriptor or fd for short
 - · Is an integer
 - · Is used to access a file
 - · Is private per process
 - · Can be used to read() and write() files

```
File can also
                                             File can also
                                                                        File can also
                              File can be read
                                             be written by
                                                           be read by
                                                                        be read by
                                                           group
                                                                        others
                                             owner
        #include <fcntl.h>
        int fd;
        mode_t mode = S_IRUSR | S_IWUSR
                                              | S IRGRP
                                                            S IROTH;
        char *filename = "/tmp/file";
        fd = open(filename, O WRONLY | O CREAT | O TRUNC, mode);
        . . .
                                     Means
                                     1. File is Writable AND
                                     2. Create file if doesn't exist AND
                                    3. Overwrite file if exists
- (read)
  * Is a system call
  * Syntax:
     ssize_t read (int fd, void *buf, size_t count)
```

- · fd file descriptor (from open ())
- · buf container for the read data
- · count number of bytes to read
- * Returns number of bytes read, if successful

* Returns 0 if is at, or past the end of file

Example

```
char buf[4096];
int fd = open("/a/b/c", 0); // open in read-only mode
lseek(fd, 1034*4096, 0); // seek to position (1034*4096) from start of file
read(fd, buf, 4096); // read 4k of data from file
```

System Calls	Code	Offset		
fd = open("file", O_RDONLY);	3	0	_	read continues
read(fd, buffer, 100);	100	100	4	
read(fd, buffer, 100);	100	200		for each call
read(fd, buffer, 100);	100	300	_	
read(fd, buffer, 100);	0	300	◀——	returns 0
close(fd);	0	-	_	if at end

- write
 - * Is a system call
 - * Writes data out of a buffer
 - * Syntax:

```
ssize_t write (int fd, const void * buf, size_t nbytes)
```

- · fd file descriptor
- · buf A pointer to a buffer to write to file
- · nbytes number of bytes to write. If smaller than buffer, the output is truncated

```
#include <unistd.h>
#include <fcntl.h>

int main(void)
{
    int filedesc = open("testfile.txt", O_WRONLY | O_APPEND);

    if (filedesc < 0) {
        return -1;
    }

    if (write(filedesc, "This will be output to testfile.txt\n", 36) != 36) {
        write(2, "There was an error writing to testfile.txt\n", 43);
        return -1;
    }

    return 0;
}</pre>
```

- lseek
 - * Reads or write to a specific offset within a file
 - * Syntax:

```
off_t lseek (int fd, off_t offset, int whence)
```

- · fd file descriptor
- · offset the offset of pointer within file (in bytes)
- · whence the method of offset

```
SEEK_SET - offset from the start of file (absolute)
SEEK_CUR - offset from current location + offset bytes (relative)
SEEK_END - offset from the end of file
```

- * Returns offset amount (in bytes) from the beginning of file
- * Returns -1 if error

Example

System Calls	Return Code	Current Offset		move 200 bytes from the
fd = open("file", O-RDONLY);	3	0		start of file
lseek(fd, 200, SEEK_SET);	200	200	A	Start or file
read(fd, buffer, 50);	50	250] _	
close(fd);	0	-	-	
				read 50 bytes

- rename
 - * Is a system call
 - * Changes the name of file
 - * Syntax: int rename(const char *old, const char *new)
 - · old name of old file
 - · new name of new file
 - * Returns 0 if successful
 - * Returns -1 if error

- stat
 - * displays metadata of a certain file stored in **inode**
 - * Syntax: int stat(const char *path, struct stat *buf)
 - · path file descriptor of file that's being inquired
 - · buf A stat structure where data about the file will be stored (see below)

```
struct stat {
                       // ID of device containing file
  dev_t
           st_dev;
                       // inode number
 ino_t
           st_ino;
 mode t
           st_mode;
                       // protection
           st_nlink;
                       // number of hard links
  nlink_t
                       // user ID of owner
 uid t
           st uid:
  gid_t
           st_gid;
                       // group ID of owner
           st_rdev;
                       // device ID (if special file)
  dev_t
  off t
           st size;
                       // total size, in bytes
 blksize_t st_blksize; // blocksize for filesystem I/O
 blkcnt_t st_blocks; // number of blocks allocated
                       // time of last access
 time_t
           st_atime;
                       // time of last modification
 time_t
           st_mtime;
           st_ctime;
                       // time of last status change
  time_t
```

Figure 39.5: The stat structure.

```
#include <unistd.h>
      #include <stdio.h>
      #include <sys/stat.h>
      #include <sys/types.h>
4
      int main(int argc, char **argv)
6
          if (argc != 2)
              return 1;
          struct stat fileStat;
11
          if (stat(argv[1],&fileStat) < 0)</pre>
              return 1;
14
          printf("Information for %s\n", argv[1]);
          printf("----\n");
          printf("File Size: \t\t%d bytes\n", fileStat.st_size);
          printf("Number of Links: \t%d\n", fileStat.st_nlink);
          printf("File inode: \t\t%d\n", fileStat.st_ino);
19
          printf("File Permissions: \t");
21
          printf( (S_ISDIR(fileStat.st_mode)) ? "d" : "-");
22
          printf( (fileStat.st_mode & S_IRUSR) ? "r" : "-");
          printf( (fileStat.st_mode & S_IWUSR) ? "w" : "-");
          printf( (fileStat.st_mode & S_IXUSR) ? "x" : "-");
          printf( (fileStat.st_mode & S_IRGRP) ? "r" : "-");
26
          printf( (fileStat.st_mode & S_IWGRP) ? "w" : "-");
27
          printf( (fileStat.st mode & S IXGRP) ? "x" : "-");
28
```

```
printf( (fileStat.st_mode & S_IROTH) ? "r" : "-");
29
          printf( (fileStat.st_mode & S_IWOTH) ? "w" : "-");
30
          printf( (fileStat.st_mode & S_IXOTH) ? "x" : "-");
31
          printf("\n\n");
32
33
          printf("The file %s a symbolic link\n", (S_ISLNK(fileStat.
34
     st_mode)) ? "is" : "is not");
35
          return 0;
36
37
```

References

1) codewiki, stat, link