## CSC373 Worksheet 0

## July 18, 2020

- 1. CLRS 4.3-1: Show that the solution of T(n) = T(n-1) + n is  $\mathcal{O}(n^2)$ .
- 2. CLRS 4.3-2: Show that the solution of  $T(n) = T(\lceil n/2 \rceil) + 1$  is  $\mathcal{O}(\lg n)$ .
- 3. **CLRS 4.3-3:** We saw that the solution of  $T(n) = 2T(\lfloor n/2 \rfloor) + n$  is  $O(n \lg n)$ . Show that the solution of this recurrence is also  $\Omega(n \lg n)$ . Conclude that the solution is  $\Theta(n \lg n)$ .
- 4. CLRS 4.3-5: Show that  $\Theta(n \lg n)$  is the solution to the "exact" recurrence (4.3) for merge sort.
- 5. **CLRS 4.3-6:** Show that the solution to T(n) = 2T(|n/2| + 17) + n is  $O(n \lg n)$ .
- 6. CLRS 4.3-7: Using the master method in Section 4.5, you can show that the solution to the recurrence T(n) = 4T(n/3) + n is  $T(n) = \Theta(n^{\log_3 4})$ . Show that a substitution proof with the assumption  $T(n) \leq c n^{\log_3 4}$  fails. Then show how to subtract off a lower-order term to make a substitution proof work.
- 7. CLRS 4.3-8: Using the master method in Section 4.5, you can show that the solution to the recurrence T(n) = 4T(n/2) + n is  $T(n) = \Theta(n^2)$ . Show that a substitution proof with the assumption  $T(n) \le cn^2$  fails. Then show how to subtract off a lower-order term to make a substitution proof work.
- 8. CLRS 4.4-1: Use a recursion tree to determine a good asymptotic upper bound on the recurrence  $T(n) = 3T(\lfloor n/2 \rfloor) + n$ . Use the subtitution method to verify your answer.