Intro to DB

CHAPTER 4 INTERMEDIATE SQL

Chapter 4: Intermediate SQL

- Join Expressions
- Views
- Transactions
- Integrity Constraints
- SQL Data Types and Schemas
- Authorization

Joined Relations

Join operation

- takes two relations and return as a result another relation.
- a Cartesian product which requires that tuples in the two relations match (under some condition).
- also specifies the attributes that are present in the result of the join
- The join operations are typically used as subquery expressions in the from clause

Join Operation

Join matches tuples with the same values specified in the on condition.

select *
from instructor join teaches on (instructor.ID=teaches.ID);

	ID	na	me	dept_name	salary		ID	cou	ırse_id	sec_id	semester	year
	10101 12121 15151 22222	Wu Mo Ein	zart stein	Comp. Sci Finance Music Physics	90000 40000 95000		10101 10101 10101 12121	CS CS	-101 -315 -347 N-201	1 1 1 1	Fall Spring Fall Spring	2009 2010 2009 2010
	32343		ID	пате	dept_name	salary	course	_id	sec_id	semester	year	2010 2009
			10101 10101 10101 12121 15151 22222 32343 45565 45565	Srinivasan Srinivasan Srinivasan Wu Mozart Einstein El Said Katz Katz	Comp. Sci.	65000 65000 90000 40000 95000 60000 75000	CS-31 CS-34 FIN-2 MU-1 PHY- HIS-3 CS-10	5 7 01 99 101 51	1 1 1 1 1 1 1 1	Fall Spring Fall Spring Spring Fall Spring Spring Spring	2009 2010 2009 2010 2010 2009 2010 2010	
dar	shan		76766 76766	Crick Crick	Biology Biology	72000 72000	BIO-1	.01	1 1	Summer		

Join Operation (cont.)

• Let $r_1(A_1, ..., A_k)$, $r_2(B_1, ..., B_n)$ select * from r_1 join r_2 on condition;

is equivalent to

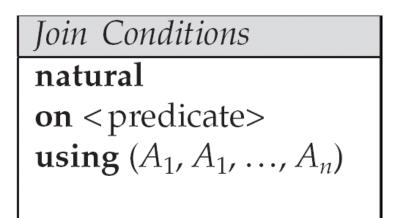
You can use expressions (that evaluate to relations) in from clauses
 from E₁, ..., E_n

select distinct *name, title* **from** *instructor* **join** *teaches* **on** (instructor.ID=teaches.ID), course where teaches.course_id=course.course_id

Joined Relations

- take two relations and return as a result another relation.
- typically used as subquery expressions in the from clause
- Join condition defines which tuples in the two relations match, and what attributes are present in the result of the join.
- Join type defines how tuples in each relation that do not match any tuple in the other relation (based on the join condition) are treated.

Join types
inner join
left outer join
right outer join
full outer join



Inner join

- Inner join = join (default)
 select * from instructor join teaches on (instructor.ID=teaches.ID);
 = select * from instructor inner join teaches on (instructor.ID=teaches.ID);
- ON clause
 - an explicit join clause (acts like a where clause)
- USING clause
 - specifies which columns to test for equality

Above example is equivalent to

select * **from** *instructor* **inner join** *teaches* **using** (*ID*);

Natural Join

- (Section 3.3.3)
- Natural join matches tuples with the same values for all common attributes, and retains only one copy of each common column

select *
from instructor natural join teaches;

Let r₁(A₁, ..., A_k, C₁, ..., C_m), r₂(B₁, ..., B_n, C₁, ..., C_m)
 select *
 from r₁ natural join r₂

is equivalent to

select
$$A_1, ..., A_k, r_1.C_1, ..., r_1.C_m, B_1, ..., B_n$$

Natural Join (cont.)

- Danger in natural join: beware of unrelated attributes with same name which get equated incorrectly
- List the names of instructors along with the titles of courses that they teach
 - Incorrect version (makes course.dept_name = instructor.dept_name)
 - select distinct name, title
 from instructor natural join teaches natural join course;
 - Correct version
 - select distinct name, title
 from instructor natural join teaches, course
 where teaches.course_id = course.course_id;
 - Another correct version
 - select distinct name, title from (instructor natural join teaches)

join course using (course_id);

Natural Join (cont.)

 Find the IDs of all students who were taught by an instructor named Einstein in building 301 (remove duplicates in the result).

select distinct takes.ID

from takes, section, teaches, instructor
where takes.course_id=section.course_id and takes.sec_id=section.sec_id and
takes.semester=section.semester and takes.year=section.year and
section.course_id=teaches.course_id and section.sec_id=teaches.sec_id and
section.year=teaches.year and section.year=teaches.year and
teaches.ID=instructor.ID and
instructor.name='Einstein' and building='301'

Join operation – Example

course

course_id	title	dept_name	credits
BIO-301	Genetics	Biology	4
CS-190	Game Design	Comp. Sci.	4
CS-315	Robotics	Comp. Sci.	3

prereq

course_id	prereq_id
BIO-301	BIO-101
CS-190	CS-101
CS-347	CS-101

course natural join prereq

course_id	title	dept_name	credits	prere_id	course_id
BIO-301	Genetics	Biology	4	BIO-101	BIO-301
CS-190	Game Design	Comp. Sci.	4	CS-101	CS-190

Note: prereq information missing for CS-315 and course information missing for CS-437.

Outer Join

- An extension of the join operation that avoids loss of information.
- Computes the join and then adds tuples form one relation that does not match tuples in the other relation to the result of the join.
- Uses null values.

Left Outer Join

course natural left outer join prereq

course_id	title	dept_name	credits	prere_id
BIO-301	Genetics	Biology	4	BIO-101
CS-190	Game Design	Comp. Sci.	4	CS-101
CS-315	Robotics	Comp. Sci.	3	null

- = course left outer join prereq on course.course_id = prereq.course_id
- = course left outer join prereq using (course_id)

course

course_id	title	dept_name	credits
BIO-301	Genetics	Biology	4
CS-190	Game Design	Comp. Sci.	4
CS-315	Robotics	Comp. Sci.	3

prereq

course_id	prereg_id
BIO-301	BIO-101
CS-190	CS-101
CS-347	CS-101

Right Outer Join

course natural right outer join prereq

course_id	title	dept_name	credits	prere_id
BIO-301	Genetics	Biology	4	BIO-101
CS-190	Game Design	Comp. Sci.	4	CS-101
CS-347	null	null	null	CS-101

course

course_id	title	dept_name	credits
BIO-301	Genetics	Biology	4
CS-190	Game Design	Comp. Sci.	4
CS-315	Robotics	Comp. Sci.	3

prereq

course_id	prereg_id
BIO-301	BIO-101
CS-190	CS-101
CS-347	CS-101

Full Outer Join

course natural full outer join prereq

course_id	title	dept_name	credits	prere_id
BIO-301	Genetics	Biology	4	BIO-101
CS-190	Game Design	Comp. Sci.	4	CS-101
CS-315	Robotics	Comp. Sci.	3	null
CS-347	null	null	null	CS-101

course

course_id	title	dept_name	credits
BIO-301	Genetics	Biology	4
CS-190	Game Design	And the second s	4
CS-315	Robotics	Comp. Sci.	3

prereq

course_id	prereg_id
BIO-301	BIO-101
CS-190	CS-101
CS-347	CS-101

Joined Relations – Examples

course inner join prereq on course.course_id = prereq.course_id

course_id	title	dept_name	credits	prere_id	course_id
BIO-301	Genetics	Biology	4	BIO-101	BIO-301
CS-190	Game Design	Comp. Sci.	4	CS-101	CS-190

course left outer join prereq on course.course_id = prereq.course_id

course_id	title	dept_name	credits	prere_id	course_id
	CHARLES AND	Biology	1923/8	BIO-101	BIO-301
CS-190	Game Design	Comp. Sci.	4	CS-101	CS-190
CS-315	Robotics	Comp. Sci.	3	null	null

Integrity Constraints

- Integrity constraints guard against accidental damage to the database
 - Examples

- A checking account must have a balance greater than \$10,000.00
- A salary of a bank employee must be at least \$4.00 an hour
- A customer must have a (non-null) phone number
- Constraints on a single relation
 - not null
 - unique
 - primary key
 - check (P), where P is a predicate
- Referential integrity

Not Null Constraint

Declare name for budget to be not null name varchar(20) not null budget numeric(12,2) not null

Exa	mple					

The Unique & Primary Key Constraints

```
• unique ( A<sub>1</sub>, A<sub>2</sub>, ..., A<sub>m</sub>)
```

- primary key (A₁, A₂, ..., A_m)
 - $A_1, A_2, \dots A_m$ form the primary key for the table

The check clause

- check (P), where P is a predicate
- Example:
 - ensure that semester is one of fall, winter, spring or summer

```
create table section (
     course_id varchar (8),
     sec_id varchar (8),
     semester varchar (6),
     year numeric (4,0),
     building varchar (15),
     room_number varchar (7),
     time slot id varchar (4),
     primary key (course_id, sec_id, semester, year),
     check (semester in ('Fall', 'Winter', 'Spring', 'Summer'))
```

Referential Integrity

- "A value that appears in one relation for a given set of attributes also appears for a certain set of attributes in another relation."
 - Example:
 - If "Biology" is a department name appearing in one of the tuples in the course relation,
 - then there exists a tuple in the *department* relation for "Biology".
- specified as part of **create table** statement
- **Foreign Key**
 - the attributes that comprise the foreign key, and
 - the name of the relation referenced by the foreign key

	course_id	title	dept_name	credits
	BIO-101	Intro. to Biology	Biology	4
	BIO-301	Genetics	Biology	4
	BIO-399	Computational Biology	Biology	3
	CS-101	Intro. to Computer Science	Comp. Sci.	4
	CS-190	Game Design	Comp. Sci.	4
	CS-315	Robotics	Comp. Sci.	3
	CS-319	Image Processing	Comp. Sci.	3
Original Slide	CS-347	Database System Concepts	Comp. Sci.	3
© Silberscha EE-181 Intro. to Digital Systems		Elec. Eng.	3	
	EINI 201	Investment Realing	Finance	2

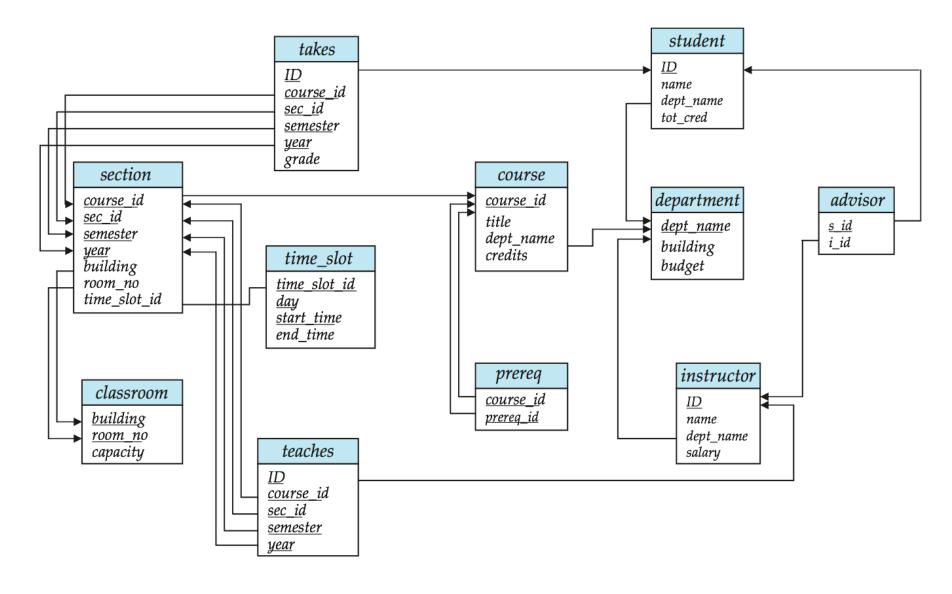
dept_name	building	budget
Biology	Watson	90000
Comp. Sci.	Taylor	100000
Elec. Eng.	Taylor	85000
Finance	Painter	120000
History	Painter	50000
Music	Packard	80000
Physics	Watson	70000

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Cascading Actions in Referential Integrity

```
create table course (
  course_id char(5) primary key,
             varchar(20),
  title
  dept_name varchar(20) references department
create table course (
  dept_name varchar(20),
  foreign key (dept_name) references department
         on delete cascade
         on update cascade,
  . . .
```

Primary & Foreign Keys in Schema Diagram



Authorization Specification in SQL

- The grant statement is used to confer authorization
 - grant <privilege list>
 - on <relation name or view name>
 - to <user list>
- <user list> is:
 - a user-id
 - public, which allows all valid users the privilege granted
 - a role (explained later)

Privileges in SQL

- select: allows read access to relation
 - (or the ability to query using the view)
 - Example: grant users U_1 , U_2 , and U_3 select authorization on the *branch* relation:



- **insert**: the ability to insert tuples
- update: the ability to update using the SQL update statement
- delete: the ability to delete tuples.
- all privileges: used as a short form for all the allowable privileges
- index: allows creation and deletion of indices.
- resources: allows creation of new relations.
- alteration: allows addition or deletion of attributes in a relation.
- drop: allows deletion of relations.

Revoking Authorization in SQL

The revoke statement is used to revoke authorization.

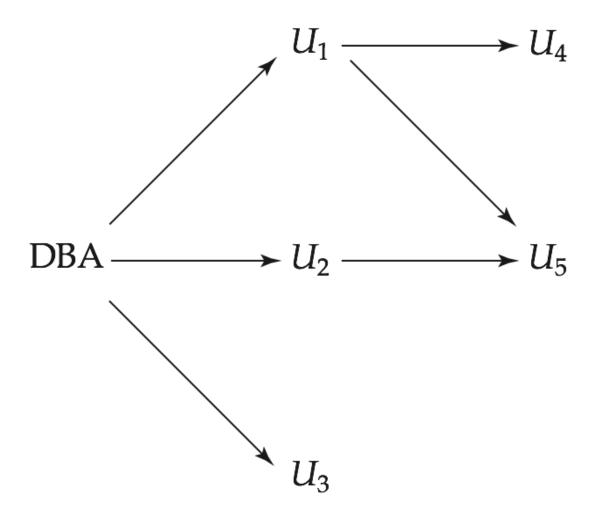
```
revoke <privilege list>
on <relation name or view name>
from <user list>
```

Example:

revoke select on branch from U_1 , U_2 , U_3

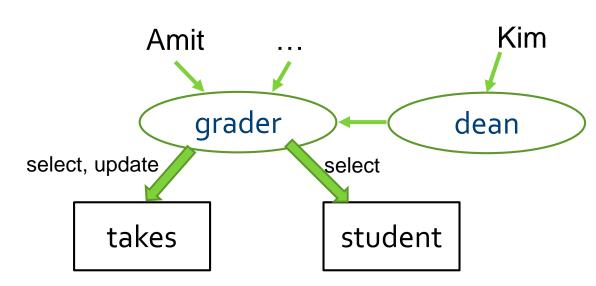
- <privilege-list> may be all to revoke all privileges the revokee may hold.
- All privileges that depend on the privilege being revoked are also revoked.
- If the same privilege was granted twice to the same user by different grantees, the user may retain the privilege after the revocation.

Authorization-Grant Graph



Roles

- Roles are used to represent a group of users and their privileges
 - create role grader,
- Privileges can be granted to roles:
 - grant select on student to grader
 - grant select, update on takes to grader
- Roles can be granted to users, as well as to other roles
 - grant grader to Amit;
 - create role dean;
 - grant grader to dean;
 - grant dean to Kim;



END OF CHAPTER 4