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CPSC 335 Algorithm Project 1 Written Report

This algorithm project was to determine the efficiency of (left-to-right) and lawnmower algorithm. Raymond and I had to decide what to implement and how to approach the problem. After timeless observation and thought process along with our pseudocode design we were able to complete the problem.

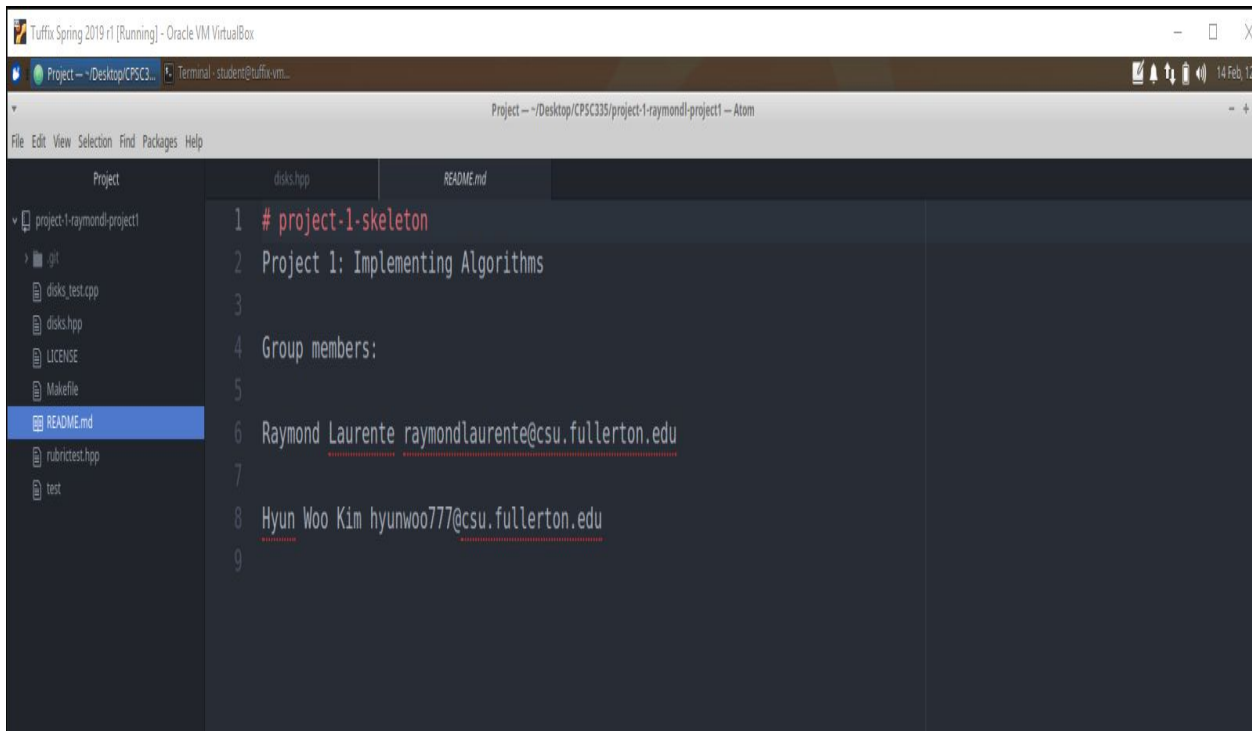
1. Left-to-right

- a. Simply left to right algorithm checks the first index of the array and checks the next index of the array and swap if necessary. Luckily, the array was in alternating order all we had to do was swap them.

2. Lawnmower

- a. This was the most difficult function to solve compare to the others
- b. We first divide the task in half
- c. Most outer loop will go half the size of the array
 - i. First inner (left) loop will iterate half way also swapping if necessary
 - ii. Second inner (right) loop will iterate half way also and swapping if necessary.
 - iii. We had to make sure that the second inner loop had to be pointing the last-2
 - iv. To compare the last-2 and last-1 to check statements

3. Screenshot of READ.ME file



```
1 # project-1-skeleton
2 Project 1: Implementing Algorithms
3
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```

4. Pseudocode

Left-to-right Algorithm

For l to end

 If(left(L) and Left+1(D) then swap

 Counter++

Lawnmower Algorithm

For l to end/2

 For left to end-1 and ++

 If(left(L) and Left+1(D) then swap

 Counter++

 For right = size-2 to right>0 and –

 If(right(L) and right+1(D) then swap

 Counter++

5. Mathematical Analysis

- Big O notation always consider the **worst-case scenario**
- left-to-right algorithm = **$O(n^2)$**
- lawnmower algorithm = **$O(n^3)$**
- proof and analysis are written at the bottom**

Left-to-right

For l = 1 to n do

 For j =1 to n do

 If(disk == D) && disk==L // 3

 Disk swap

 Swap++ //2 steps

 // no else therefore, 0

// the if statement step counter

3 + max(2,0) = 5 steps

// the inner for loop

[(n-1)/1] +1 * 5 = 5n

// outer for loop

$$[(n-1)/1] + 1 * 5n = 5n^2$$

therefore, the Big(O) notation for this algorithm is **$O(n^2)$**

Lawn Mower Algorithm

for l= 0 to n do

for left=0 to n - i;

if(disk(left) == L && disk(left+1) ==D

swap

counter++

for right = n-2 to right>0

if(disk(right) == L && disk(right+1) == D

swap

count++

// if- statement counter

$$3 + \max(2,0) = 5$$

// bottom inner loop

$$[(n-2) - 0/1] + 1 * 5 = 5n - 5$$

$$\sum 5n - 5 = 5n^2 - 5n$$

// upper inner loop

$$[(n-1)/1] + 1 * 5 = 5n$$

$$\sum 5n = 5n^2$$

// most outer for loop

$\lceil (n-1)/1 \rceil + 1 * \max(5n^2 - 5n, 5n^2)$

$n * 5n^2 = 5n^3$

therefore, the worst case scenario for lawnmower algorithm is $O(n^3)$