

# HYPERLINK

No Blocks, No Chains

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<https://hyperlink.com>

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Traditional DeFi platforms are shackled by legacy architectures, forcing reliance on Layer 2 networks and third-party providers that crumble under high volume. Hyperlink challenges this paradigm. Built on a directed acyclic graph (DAG)<sup>a</sup>, transactions are instantly broadcast into the network—no blocks, no chains, no waiting. While conventional networks slow down during peak usage, Hyperlink accelerates: higher volume means faster finality. Through native protocol-level financial primitives, built-in liquidity pools, and automated market making, Hyperlink delivers instantaneous DeFi operations without external dependencies. A sophisticated fee-per-byte system and reputation-based P2P propagation ensure network health, creating a truly scalable, self-regulating financial ecosystem.

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<sup>a</sup>[https://en.wikipedia.org/wiki/Directed\\_acyclic\\_graph](https://en.wikipedia.org/wiki/Directed_acyclic_graph)

## 1 Overview

Hyperlink represents a fundamental rethinking of decentralized networks, built from first principles to solve the core limitations of legacy architectures. Rather than layering solutions on top of an inherently constrained foundation, we've designed a system where high performance and rich functionality are intrinsic properties of the base protocol.

At its core, Hyperlink is built on three key innovations:

1. **DAG-based State Management:** Instead of forcing global consensus through periodically publishing held transactions, Hyperlink transactions directly reference and validate previous transactions as soon as they are broadcast. This creates a naturally parallel system where higher transaction volume accelerates rather than impedes network throughput.
2. **Protocol-level DeFi:** Financial primitives like token swaps, liquidity pools, and automated market making are implemented directly in the base protocol. This eliminates the complexity, cost, and security risks of smart contract layers while enabling atomic execution of complex financial operations.
3. **Reputation-based Network Control:** The system implements a dual reputation model that tracks both peer node behavior and wallet-level activity. This creates a self-regulating network that naturally resists spam and abuse while remaining permissionless and decentralized.

These architectural choices enable several unique capabilities:

- **Instant Finality:** Transactions are broadcast immediately into the network without waiting for transaction selection by miners or global consensus.
- **Atomic DeFi Operations:** Complex financial operations execute as single atomic transactions with protocol-level validation.
- **Natural Scalability:** Network performance improves with increased volume as more transactions create more parallel validation paths.
- **Economic Security:** A sophisticated fee-per-byte system combined with reputation tracking creates natural economic barriers to network abuse.

The following sections detail the technical implementation of these concepts, starting with the core network architecture and building up to the high-level financial capabilities. Each component is designed to work in concert, creating a cohesive system that maintains security and decentralization while delivering unprecedented performance and functionality.

## 2 Network Architecture

### 2.1 P2P Network

The network is built on libp2p<sup>1</sup> with both TCP and WebSocket transport support. Node discovery uses a Kademlia distributed hash table (DHT)<sup>2</sup> - a decentralized system that provides a lookup service similar to a hash

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<sup>1</sup><https://libp2p.io/>

<sup>2</sup>[https://en.wikipedia.org/wiki/Distributed\\_hash\\_table](https://en.wikipedia.org/wiki/Distributed_hash_table)

table, where the responsibility for maintaining key mappings is distributed among nodes in a way that minimizes disruption from node changes.

Each node maintains:

- A target of 8 peer connections
- Persistent connections to validator nodes
- Regular health checks on connections

Node initialization process:

1. Create or load P2P keypair from config directory
2. Initialize libp2p host with TCP and WebSocket transports
3. Set up protocol handlers for all supported protocols
4. Bootstrap DHT with validator nodes
5. Begin periodic peer discovery

## 2.2 Network Protocols

The system implements five core protocols:

### FullSyncProtocol

- Handles initial graph synchronization
- Uses batched transaction transfer
- Implements retry mechanism with exponential backoff
- Maximum 3 retry attempts per sync operation

### StateHashProtocol

- Quick state verification between peers
- Uses SHA3-256 of all transaction hashes
- Triggers full sync on mismatch

### TransactionProtocol

- Handles new transaction propagation
- Implements reputation-based forwarding
- Maintains connection pools for efficient broadcasting

### MissingTransactionsProtocol

- Retrieves specific transactions by hash
- Uses batched requests/responses
- Handles partial fulfillment of requests

### TransactionOrderProtocol

- Manages transaction ordering proposals
- Validator-specific protocol
- Ensures DAG convergence

## 2.3 Peer Selection and Reputation

The system implements a peer selection mechanism based on reputation scores. Each peer's reputation is tracked on a scale from 0.0 to 1.0, with several key thresholds:

- Initial peer reputation: 0.5
- Minimum threshold for transaction relay: 0.3
- Maximum peer reputation: 1.0
- Minimum peer reputation: 0.0

### 2.3.1 Reputation Scoring

Peer reputation is dynamically adjusted based on behavior:

- **Positive Actions (+0.05):**
  - Successfully propagating valid transactions
  - Maintaining stable connections
  - Providing valid sync responses
- **Negative Actions (-0.1):**
  - Propagating invalid transactions
  - Providing invalid sync data
  - Connection instability

The reputation system is implemented as:

```

1 func (n *Node) GetPeerReputation(peerID peer.ID)
2     float64 {
3     reputation, exists := n.reputations[peerID.String()
4     ]
5     if !exists {
6         return InitialPeerReputation // 0.5
7     }
8     return reputation
9 }
10
11 func (n *Node) UpdatePeerReputation(peerID peer.ID,
12     change float64) {
13     current, exists := n.reputations[peerID.String()]
14     if !exists {
15         current = InitialPeerReputation
16     }
17     current += change
18     if current > MaxPeerReputation {
19         current = MaxPeerReputation
20     } else if current < MinPeerReputation {
21         current = MinPeerReputation
22     }
23     n.reputations[peerID.String()] = current
24 }
```

### 2.3.2 Transaction Propagation Rules

The system implements different propagation rules based on transaction source and peer reputation:

#### 1. Priority Transactions:

- Mining rewards (from TokenGrantAddress)
- Transactions from wealthy accounts (balance  $\geq$  PovertyLine)
- Transactions with valid mining proofs
- These are propagated to all peers with reputation  $\geq 0.3$

#### 2. Standard Transactions:

- Must come from wallets with sufficient reputation
- Only propagated to peers with reputation  $\geq 0.3$
- Subject to additional validation checks

The propagation logic is implemented as:

```

1 func (n *Node) BroadcastTransaction(tx *Transaction) {
2     if tx.From == TokenGrantAddress ||
3         tx.MiningProof != nil ||
4         n.graph.State[tx.From] >= PovertyLine {
5         // Priority transaction - broadcast to all
6         // qualified peers
7         for _, peerID := range n.host.Network().Peers
8         () {
9             if n.GetPeerReputation(peerID) >=
10             MinimumPeerReputationThreshold {
11                 go n.SendTransactionToPeer(tx, peerID)
12             }
13         }
14     } else if n.GetWalletReputation(tx.From) >=
15     MinimumReputationThreshold {
16         // Standard transaction - broadcast only if
17         // sender has sufficient reputation
18         for _, peerID := range n.host.Network().Peers
19         () {
20             if n.GetPeerReputation(peerID) >=
21             MinimumPeerReputationThreshold {
22                 go n.SendTransactionToPeer(tx, peerID)
23             }
24         }
25     }
26 }

```

### 2.3.3 Connection Management

The system maintains a target of 8 peer connections, with special handling for validator nodes:

- Persistent connections to validator nodes
- Regular health checks every 30 seconds
- Automatic reconnection attempts on failure
- Connection pooling for efficient message distribution

This reputation-based peer selection system provides several benefits:

- Spam prevention through reputation requirements
- Natural isolation of misbehaving peers
- Efficient resource allocation to reliable peers
- Automatic network self-organization
- Resilience against network attacks

The combination of reputation tracking, dynamic adjustment, and differentiated propagation rules creates a robust peer-to-peer network that can efficiently distribute transactions while maintaining security against various forms of abuse.

## 3 Account System

### 3.1 Key Derivation (SLIP-0010)

Hyperlink implements the SLIP-0010<sup>3</sup> standard for universal private key derivation, providing a standardized and secure method for generating hierarchical deterministic wallets. The implementation focuses exclusively on ed25519 for signatures and key generation:

- Hardened-only derivation for ed25519
- Full compliance with SLIP-0010 test vectors
- Standardized BIP44 path structure

<sup>3</sup><https://github.com/satoshi-labs/slips/blob/master/slip-0010.md>

- Secure master key generation

The derivation process follows the SLIP-0010 specification:

1. Generate master key from seed using HMAC-SHA512 with key "ed25519 seed"
2. Derive child keys using standardized paths with hardened indices only
3. Implement proper key validation and regeneration

### 3.2 BIP44 Path Structure

The system uses standard BIP44 paths with enforced hardened derivation:

- Purpose: 44' (hardened)
- Coin type: 901' (hardened)
- Account: 0' (hardened)
- Change/Token type: x' (hardened)
- Address index: y' (hardened)

Example paths:

- Main account: m/44'/901'/0'/0'/0'
- Token account: m/44'/901'/0'/1'/0'
- NFT account: m/44'/901'/0'/2'/0'

### 3.3 Address Format

Addresses are derived directly from ed25519 public keys using:

- Base58 encoding for human readability
- 44-character fixed length format
- Direct derivation from Ed25519 public keys
- No checksum for simplicity

This standardized approach ensures:

- Cross-platform compatibility
- Secure key generation
- Deterministic recovery
- Hierarchical organization

## 4 Consensus Mechanism

### 4.1 DAG Structure

The system uses a DAG where each transaction must reference 1-4 previous transactions. The reference selection process follows these steps:

#### 4.1.1 Candidate Selection

- Scans recent transactions within 12-hour window
- Prioritizes newer transactions
- Excludes self-references

## 4.2 Full Sync Protocol

The sync process implements the following message structures:

```

1  type SyncRequest struct {
2      LastHash      string
3      LastTimestamp int64
4  }
5
6  type SyncResponse struct {
7      Transactions []Transaction
8      HasMore      bool
9      NextHash     string
10 }

```

### 4.2.1 Reference Selection Algorithm

---

#### Algorithm 1 Reference Selection

---

```

1: totalWeight ← 0
2: for i ← 0 to len(candidateTxes) - 1 do
3:     weight ← len(candidateTxes) - i
4:     totalWeight ← totalWeight + weight
5: end for
6: randomWeight ← random(0, totalWeight)

```

---

### 4.2.2 Reference Validation

- Ensures minimum reference count
- Verifies reference existence
- Checks reference age
- Prevents circular references

## 5 State Management

### 5.1 Graph State

The graph state represents the current network consensus, tracking:

1. Transaction History:
  - Complete DAG of all transactions
  - Reference relationships
  - Temporal ordering
2. Account Balances:
  - Current balance for all addresses
  - Atomic updates during transaction processing
  - Double-spend prevention
3. Validator Set:
  - Current validator nodes
  - Validator status tracking
  - Update history

### 5.2 Database Management

The persistence layer uses SQLite as its storage backend, implementing several critical optimizations and safety measures to ensure reliable operation at scale.

### 5.2.1 Performance Optimizations

The system implements several performance-focused features:

- **Batch Processing:**
  - Transactions are collected into batches of 100
  - Reduces disk I/O overhead
  - Improves throughput under high load
- **Prepared Statements:**
  - SQL statements are pre-compiled
  - Reduces parsing overhead
  - Prevents SQL injection vulnerabilities
- **Transaction Pooling:**
  - Reuses database connections
  - Minimizes connection overhead
  - Manages concurrent access efficiently

### 5.2.2 Reliability Mechanisms

The database layer implements multiple reliability features to ensure data integrity:

1. **Atomic Commits:**
  - Database transactions are all-or-nothing
  - Batch operations are atomic
  - Prevents partial updates during failures
2. **Rollback Protection:**
  - Automatic rollback on errors
  - Transaction boundary management
  - State consistency preservation
3. **Corruption Prevention:**
  - Safe file handling practices
  - Directory traversal protection
  - Proper permission management

The implementation ensures safe database operations:

```

1  func ProcessBatch() {
2      batchMutex.Lock()
3      batch := transactionBatch
4      transactionBatch = nil
5      batchMutex.Unlock()
6
7      db := GetDB()
8      dbTx, err := db.Beginx()
9      if err != nil {
10         LogError("Failed to begin database transaction
            : %v", err)
11         return
12     }
13
14     stmt, err := dbTx.Prepare('INSERT OR REPLACE INTO
            transactions (hash, data) VALUES (?, ?)')
15     if err != nil {
16         LogError("Failed to prepare statement: %v",
            err)
17         if err := dbTx.Rollback(); err != nil {
18             LogError("Failed to rollback database
            transaction: %v", err)
19         }
20         return
21     }
22     defer stmt.Close()
23 }

```

```

24 // Process batch with rollback protection
25 for _, tx := range batch {
26     txData, err := json.Marshal(tx)
27     if err != nil {
28         LogError("Failed to marshal transaction: %v", err)
29         continue
30     }
31
32     _, err = stmt.Exec(tx.Hash, string(txData))
33     if err != nil {
34         LogError("Failed to insert/update transaction: %v", err)
35         if err := dbTx.Rollback(); err != nil {
36             LogError("Failed to rollback database transaction: %v", err)
37         }
38         return
39     }
40 }
41
42 if err := dbTx.Commit(); err != nil {
43     LogError("Failed to commit database transaction: %v", err)
44     if err := dbTx.Rollback(); err != nil {
45         LogError("Failed to rollback database transaction: %v", err)
46     }
47     return
48 }
49 }

```

### 5.2.3 File System Safety

The system implements careful file system management:

#### 1. Path Validation:

- Sanitizes database file paths
- Prevents directory traversal attacks
- Ensures operations stay within config directory

#### 2. Permission Management:

- Database files created with 0600 permissions
- Directories created with 0700 permissions
- Prevents unauthorized access

#### 3. Resource Management:

- Proper file handle cleanup
- Deferred statement closing
- Memory leak prevention

The system also provides flexibility through configuration:

- Optional in-memory database for testing
- Configurable batch sizes
- Automatic database file creation

This comprehensive approach to database management ensures that:

- Transaction data is stored reliably
- System performance remains high under load
- Data integrity is maintained
- Recovery from errors is automatic
- Security is maintained at the storage layer

## 6 Transaction Model

Hyperlink implements a flexible transaction model where each transaction represents one or more state changes to the network. While the basic use case is token transfer, the system is designed to support various types of state changes:

### 6.1 Transaction Structure

Each transaction contains:

- Source address (From)
- One or more state changes
- Network fee
- References to previous transactions
- Cryptographic signature

### 6.2 State Changes

State changes can represent:

- Basic token transfers
- NFT minting and transfers
- Token program instructions
- Smart contract state updates

The change structure is designed to be extensible:

```

1 type TransactionChange struct {
2     To      string      // Target address
3     Amount uint64      // Value for token
4     transfers
5     Data interface{} // Program instructions or metadata
6 }

```

### 6.3 Program Instructions

Transactions can carry program instructions that define operations like:

- Creating new tokens
- Minting NFTs
- Updating token properties
- Executing smart contract functions

This flexible model allows Hyperlink to support a wide range of decentralized applications while maintaining the simplicity and efficiency of the core transaction system.

### 6.4 Transaction Processing

Transactions are processed in parallel with the following rules:

#### 6.4.1 Reference Resolution

- All referenced transactions must exist
- References must be within time window
- First transaction in empty graph exempted

### 6.4.2 State Application

- Atomic balance updates
- Double-spend prevention
- Fee processing
- Mining reward distribution

## 7 Token Generation Through Key Discovery

The Hyperlink token generation system implements a novel approach that is unique among cryptocurrency systems. Instead of competing to process transactions, participants generate Ed25519 key pairs and encode their public keys, searching for specific properties in the resulting public key hashes.

This approach serves multiple purposes:

1. Provides a deterministic and verifiable way to generate new tokens
2. Creates controlled token issuance through tunable difficulty requirements
3. Enables fully independent mining without coordination or competition

Unlike traditional mining systems, key discovery allows parallel token generation without requiring global consensus or competition between miners. Each discovered key meeting the difficulty requirement represents a provably rare artifact that can be independently verified by any participant.

### 7.1 Key Discovery Verification

The verification process ensures three critical aspects:

1. The winning key meets the validator's current difficulty requirement
2. The discoverer can sign with the winning key
3. The miner has properly claimed the reward using their HD wallet's primary account

This creates an unforgeable chain of custody from discovery to reward:

```

1 // Extract timestamp from proof message
2 parts := strings.Split(proofMessage, ":")
3 if len(parts) != 3 {
4     return fmt.Errorf("invalid proof message format")
5 }
6 timestamp, err := strconv.ParseInt(parts[2], 10, 64)
7 if err != nil {
8     return fmt.Errorf("invalid timestamp in proof message")
9 }
10
11 // Use the full proof message for validation
12 proofMessage := fmt.Sprintf("%s:%s:%d", proof.
13     WinningAddress, proof.MinerAddress, timestamp)
14
15 // Verify winning key signature
16 winningPubKey, err := AddressToPublicKey(proof.
17     WinningAddress)
18 if err != nil {
19     return fmt.Errorf("failed to get winning public
20     key: %v", err)
21 }
22 if !ed25519.Verify(winningPubKey, []byte(proofMessage),
23     proof.WinningSignature) {
24     return fmt.Errorf("invalid winning key signature")
25 }

```

```

21 }
22
23 // Verify miner signature using HD wallet's primary
24 // account
25 minerPubKey, err := AddressToPublicKey(proof.
26     MinerAddress)
27 if err != nil {
28     return fmt.Errorf("failed to get miner's public
29     key: %v", err)
30 }
31
32 if !ed25519.Verify(minerPubKey, []byte(proofMessage),
33     proof.MinerSignature) {
34     return fmt.Errorf("invalid miner signature")
35 }
36
37 // Verify difficulty meets validator's requirement
38 isWinner, difficulty := IsWinningKey(proof.
39     WinningAddress, validatorDifficulty)
40 if !isWinner {
41     return fmt.Errorf("winning address does not meet
42     difficulty requirement")
43 }

```

### 7.2 Mining Process Details

The mining process follows a simple but effective sequence:

1. Generate a new Ed25519 keypair
2. Take the public key and encode it using Base58:

$$address = base58(ed25519.publicKey) \quad (1)$$

3. Calculate SHA3-256 hash of the Base58-encoded public key:

$$hash = sha3.256(address) \quad (2)$$

4. Count the leading zeros in the resulting hash
5. If the number of leading zeros meets or exceeds the validator's difficulty requirement, a reward is issued proportional to the difficulty

This process creates a provably fair system where:

- The difficulty of finding a winning key is predictable
- The rarity of the key can be instantly verified by any participant
- Rewards scale exponentially with the difficulty of the discovered key
- No coordination between miners is required
- Difficulty is controlled by validator nodes

### 7.3 Mining Reward Calculation

The mining reward follows an exponential scaling formula:

$$reward = BaseReward \cdot 2^{(difficulty - MinDifficulty)} \quad (3)$$

Where:

- *BaseReward* is 1 unit (0.00000001 HYY)
- *MinDifficulty* is 6 leading zeros
- *difficulty* is the number of leading zeros in the hash

This creates an incentive structure where higher difficulties receive exponentially larger rewards:

- 6 zeros = 0.00000001 HYY
- 7 zeros = 0.00000002 HYY
- 8 zeros = 0.00000004 HYY



## 7.4 Mining Process Implementation

The mining process operates as a multi-threaded system that efficiently utilizes available CPU cores:

1. Generates Ed25519 key pairs
2. Checks if they meet the difficulty requirement
3. Submits proofs when winning keys are found
4. Maintains statistics on keys checked per second

```

1 func MiningWorker(node *Node) {
2     for {
3         // Generate new Ed25519 keypair
4         pubKey, privKey, err := ed25519.GenerateKey(
5             rand.Reader)
6         if err != nil {
7             continue
8         }
9
10        // Encode public key
11        pubKeyBase58 := base58.Encode(pubKey)
12
13        // Check if key meets difficulty requirement
14        hash := sha3.Sum256([]byte(pubKeyBase58))
15        difficulty := countLeadingZeros(hash)
16
17        if difficulty >= MinMinerDifficulty {
18            // Create and submit mining proof
19            proof := &MiningProof{
20                WinningAddress: pubKeyBase58,
21                MinerAddress:   node.wallet.Address
22            },
23            WinningSignature: nil,
24            MinerSignature:   nil,
25        }
26
27        // Sign proof with winning key
28        message := fmt.Sprintf("%s:%s", proof.
29            WinningAddress, proof.MinerAddress)
30        proof.WinningSignature = ed25519.Sign(
31            privKey, []byte(message))
32
33        // Sign with miner wallet
34        proof.MinerSignature = node.wallet.Sign([]
35            byte(message))
36
37        // Submit proof
38        node.SubmitMiningProof(proof)
39    }
40 }

```

## 7.5 Mining Reward Issuance

When a miner submits a proof, the validator processes it through several stages:

```

1 // First, verify the mining proof
2 proof := tx.MiningProof
3 proofMessage := fmt.Sprintf("%s:%s:%d", proof.
4     WinningAddress, proof.MinerAddress, timestamp)
5
6 // Verify winning key signature
7 winningPubKey, _ := AddressToPublicKey(proof.
8     WinningAddress)
9 if !ed25519.Verify(winningPubKey, []byte(proofMessage),
10     proof.WinningSignature) {
11     return false
12 }
13
14 // Verify miner key signature
15 minerPubKey, _ := AddressToPublicKey(proof.
16     MinerAddress)
17 if !ed25519.Verify(minerPubKey, []byte(proofMessage),
18     proof.MinerSignature) {
19     return false
20 }
21
22 isWinner, difficulty := IsWinningKey(proof.
23     WinningAddress, validatorDifficulty)
24 if !isWinner {
25     return false
26 }

```

```

21 // Calculate reward based on difficulty
22 reward := CalculateMinerReward(difficulty)
23
24 // Create mining reward instruction
25 miningInst := MiningRewardInstruction{
26     WinningAddress: proof.WinningAddress,
27     Difficulty:     difficulty,
28     ProofMessage:   proofMessage,
29 }
30
31 // Create and process the token issuance transaction
32 tx := Transaction{
33     From:      TokenGrantAddress,
34     Timestamp: timestamp,
35     Changes:   []TransactionChange{
36         {
37             To:      minerAddress,
38             Amount:   reward,
39             InstructionType: InstructionMiningReward,
40             InstructionData: instData,
41         },
42     },
43     Fee:      0,
44     MiningProof: proof,
45 }

```

The mining reward issuance process implements multiple security checks before creating new tokens:

1. **Proof Verification:** The system first verifies that both the winning key and miner's key have properly signed the proof message with timestamp, creating a cryptographic chain of custody.
2. **Difficulty Check:** The winning address is verified to meet the validator's current difficulty requirement by counting leading zeros in its hash.
3. **Reward Calculation:** The reward amount is calculated based on the achieved difficulty level, with higher difficulties earning exponentially larger rewards.
4. **Token Creation:** A special transaction is created from the TokenGrantAddress with the MiningReward instruction type.
5. **Fee Handling:** Mining transactions have zero fees to maximize rewards for miners.

This process ensures that token creation is:

- Cryptographically secure
- Independently verifiable
- Automatically scaled by difficulty
- Properly recorded in the network state
- Controlled by validator nodes

This process ensures that:

- The mining proof is valid and meets difficulty requirements
- The reward is calculated based on the achieved difficulty
- The token issuance is atomic with proof verification
- The miner pays the standard transaction fee from their reward
- The state change is properly recorded in the DAG

## 8 Transaction Processing

### 8.1 Transaction Creation and Validation

Each transaction must reference previous transactions, creating a web of confirmations. The system enforces several rules:

#### 8.1.1 Reference Requirements

- Minimum of 1 reference
- Maximum of 4 references
- References must be within a 12-hour window
- First transaction in network exempted from references

#### 8.1.2 Balance and Fee Rules

- All transactions must pay a base fee (10000 units)
- Total transaction amount must not exceed sender's balance
- Fees are burned (sent to BURN address)

The transaction structure is defined as:

```

1  type Transaction struct {
2      Timestamp    int64
3      Hash         string
4      From         string
5      Signature    string
6      Changes      []TransactionChange
7      References   []string
8      Fee          uint64
9      MiningProof *MiningProof
10 }
```

## 8.2 Transaction Propagation

The propagation system implements a reputation-based flooding protocol that balances network efficiency with spam prevention. The system uses multiple mechanisms to ensure reliable and secure transaction propagation:

### 8.2.1 Reputation Scoring

The system maintains two types of reputation scores:

#### 1. Wallet Reputation (0.0 - 1.0):

- Wealthy accounts (above poverty line) automatically receive maximum reputation
- New wallets start with a base reputation of 1.0
- Successful transactions increase reputation by 0.1
- Invalid transactions decrease reputation by 0.2
- Reputation is weighted by the wallet's balance relative to the poverty line

#### 2. Peer Reputation (0.0 - 1.0):

- New peers start with reputation of 0.5
- Propagating valid transactions increases reputation by 0.05
- Propagating invalid transactions decreases reputation by 0.1
- Minimum threshold of 0.3 required for transaction relay

### 8.2.2 Propagation Rules

Transactions are propagated differently based on their source:

#### 1. Priority Transactions:

- Mining reward transactions (from Token-GrantAddress)
- Transactions from accounts above the poverty line
- These are always propagated to all peers above minimum reputation

#### 2. Regular Transactions:

- Must come from wallets with reputation above 0.5
- Are only propagated to peers with reputation above 0.3
- First transaction from new wallets must exceed minimum amount threshold

### 8.2.3 Network Optimization

The system implements several optimizations for efficient propagation:

- Stream reuse for multiple transactions to the same peer
- Transaction deduplication using a broadcast cache
- Automatic stream health checks and recovery
- Batched transaction processing during sync

### 8.2.4 Gossip Protocol

A background gossip protocol runs every 10 seconds to ensure transaction propagation reliability:

- Checks for non-broadcasted transactions in the local graph
- Re-attempts propagation of previously failed broadcasts
- Helps network recovery after temporary partitions
- Maintains transaction availability across the network

This multi-layered approach creates a robust propagation system that:

- Prevents spam through reputation requirements
- Prioritizes transactions from proven participants
- Maintains network efficiency through peer scoring
- Ensures reliable transaction propagation
- Recovers automatically from network issues

### 8.3 Cryptographic Verification

Transactions must meet the following cryptographic requirements:

- Transactions must be signed by sender
- Hash must be correctly computed
- Mining proofs must be valid (for mining transactions)

```

1 // Hash verification
2 if tx.Hash != tx.ComputeHash() {
3     return fmt.Errorf("invalid transaction hash")
4 }
5
6 // Balance verification
7 totalAmount := tx.Fee
8 for _, change := range tx.Changes {
9     totalAmount += change.Amount
10 }
11 if g.State[tx.From] < totalAmount {
12     return fmt.Errorf("insufficient balance")
13 }
14
15 // Signature verification
16 senderPubKey, err := AddressToPublicKey(tx.From)
17 if !tx.Verify(senderPubKey) {
18     return fmt.Errorf("invalid signature")
19 }

```

### 8.4 Stream Management

The system implements stream management:

- Reuses existing healthy streams
- Implements timeout handling
- Provides acknowledgment system

```

1 // Check existing streams
2 for _, conn := range n.host.Network().ConnsToPeer(
3     peerID) {
4     for _, stream := range conn.GetStreams() {
5         if stream.Protocol() == protocol.ID(
6             TransactionProtocol) {
7             // Reuse existing stream if healthy
8             if err := stream.SetWriteDeadline(time.Now(
9                 ).Add(time.Second)); err == nil {
10                 if _, err := stream.Write([]byte{0});
11                 err == nil {
12                     stream.SetWriteDeadline(time.Time
13                         {}))
14                     existingStream = stream
15                     break
16                 }
17             }
18             stream.Reset()
19         }
20     }
21 }

```

### 8.5 State Updates

The Hyperlink system implements atomic state updates to ensure consistency and prevent race conditions when modifying account balances. The state update mechanism is a critical component that:

- Ensures atomic (all-or-nothing) balance updates
- Prevents negative balances
- Handles both regular transactions and mining rewards
- Maintains consistency during concurrent operations

#### 8.5.1 Update Process

The state update process follows several key steps:

1. **Balance Change Collection:** First, all balance changes from the transaction are collected and categorized:
  - Positive changes (credits to receiving accounts)
  - Negative changes (debits from sending accounts)
  - Fee deductions (sent to burn address)
2. **Special Case Handling:** Mining reward transactions from the TokenGrantAddress are processed differently:
  - No balance verification needed (allowed to create new tokens)
  - Direct credit to recipient addresses
  - Fee still collected and burned
3. **Balance Verification:** For regular transactions:
  - Total outgoing amount (including fee) is calculated
  - Sender's balance is verified to be sufficient
  - Prevents any possibility of overdraft
4. **Atomic Application:** Changes are applied atomically:
  - All balance updates happen together
  - If any part fails, entire transaction is rolled back
  - Maintains system consistency

The implementation ensures safe concurrent operation:

```

1 func (g *Graph) ApplyTransaction(tx Transaction) error
2 {
3     balanceChanges := make(map[string]uint64)
4     negativeChanges := make(map[string]uint64)
5
6     // Process changes
7     if tx.From == TokenGrantAddress {
8         // Handle mining rewards
9         for _, change := range tx.Changes {
10             balanceChanges[change.To] = SafeAdd(
11                 balanceChanges[change.To], change.Amount)
12         }
13     } else {
14         // Handle regular transactions
15         totalDeduction := tx.Fee
16         for _, change := range tx.Changes {
17             totalDeduction = SafeAdd(totalDeduction,
18                 change.Amount)
19             balanceChanges[change.To] = SafeAdd(
20                 balanceChanges[change.To], change.Amount)
21         }
22         negativeChanges[tx.From] = totalDeduction
23     }
24
25     // Apply changes atomically
26     for address, addition := range balanceChanges {
27         currentBalance := g.State[address]
28         subtraction := negativeChanges[address]
29
30         if subtraction > 0 {
31             if subtraction > currentBalance {
32                 return fmt.Errorf("negative balance
33                     would occur")
34             }
35             currentBalance -= subtraction
36         }
37         g.State[address] = SafeAdd(currentBalance,
38             addition)
39     }
40     return nil
41 }

```

### 8.5.2 Safety Mechanisms

The system implements several safety mechanisms to maintain integrity:

#### 1. Overflow Protection:

- All arithmetic operations are checked for overflow
- Uses safe addition function that panics on overflow
- Prevents balance corruption from integer overflow

#### 2. Validation Checks:

- Pre-validation of transaction amounts
- Balance sufficiency verification
- Fee verification

#### 3. Error Handling:

- Clear error messages for debugging
- Transaction rollback on any error
- Logging of all state changes

This comprehensive state management system ensures that the network maintains consistent and accurate account balances while preventing any potential exploitation through race conditions or arithmetic errors. The atomic nature of updates, combined with thorough validation and safety checks, provides a robust foundation for the network's financial operations.

## 8.6 Reputation System Implementation

The reputation system uses two distinct but interconnected reputation tracking mechanisms: wallet reputation and peer reputation. Each serves a different purpose in maintaining network health.

### 8.6.1 Wallet Reputation Calculation

The wallet reputation system implements a balance-weighted approach:

```

1 func (n *Node) GetWalletReputation(address string)
2     float64 {
3     balance := n.graph.State[address]
4     if balance >= PovertyLine {
5         return MaxReputation
6     }
7
8     reputation, exists := n.reputations[address]
9     if !exists {
10        balanceWeight := float64(balance) / float64(
11            PovertyLine)
12        return InitialReputation * (1 + balanceWeight)
13    }
14    return reputation
15 }
```

This implementation provides several key features:

#### 1. Wealth-Based Trust:

- Accounts with balances above the poverty line (10 HYY) automatically receive maximum reputation
- This reflects the assumption that accounts with significant stakes are less likely to spam
- Provides immediate trust for well-funded accounts

#### 2. Balance Weighting:

- For accounts below the poverty line, reputation is weighted by their balance
- The weighting factor is calculated as:  $balanceWeight = balance / PovertyLine$
- This creates a smooth progression of trust as accounts accumulate funds

#### 3. New Account Handling:

- New accounts start with a base reputation of 1.0
- This base value is then modified by their balance weight
- Allows new accounts to participate while maintaining spam protection

### 8.6.2 Peer Reputation Management

The peer reputation system focuses on network behavior:

```

1 func (n *Node) UpdatePeerReputation(peerID peer.ID,
2     change float64) {
3     current, exists := n.reputations[peerID.String()]
4     if !exists {
5         current = InitialPeerReputation
6     }
7
8     current += change
9
10    if current > MaxPeerReputation {
11        current = MaxPeerReputation
12    } else if current < MinPeerReputation {
13        current = MinPeerReputation
14    }
15
16    n.reputations[peerID.String()] = current
17 }
```

This implementation provides:

#### 1. Bounded Reputation:

- Reputation is capped between 0.0 and 1.0
- Prevents reputation inflation or underflow
- Ensures consistent behavior across the network

#### 2. Gradual Trust Building:

- New peers start at 0.5 reputation
- Small positive increments (+0.05) for good behavior
- Larger negative increments (-0.1) for bad behavior
- Creates bias toward conservative trust

#### 3. Persistent Memory:

- Reputation persists across connection sessions
- Allows network to remember reliable peers
- Helps maintain stable peer relationships

### 8.6.3 Reputation Impact on Network Operations

The reputation scores directly influence network behavior:

#### 1. Transaction Propagation:

- Transactions from high-reputation wallets are prioritized
- Only peers with reputation  $\geq 0.3$  receive transaction broadcasts

- Creates natural spam filtering at network level

## 2. Peer Selection:

- Higher reputation peers are preferred for connection maintenance
- Low reputation peers may be disconnected during network pruning
- Helps maintain optimal network topology

## 3. Resource Allocation:

- Connection slots are preferentially allocated to high-reputation peers
- Sync requests from high-reputation peers are prioritized
- Optimizes network resource utilization

This dual reputation system creates a self-regulating network where both account behavior and peer performance contribute to overall network health. The system is designed to be:

- **Self-balancing:** Poor behavior is naturally penalized while good behavior is rewarded
- **Attack-resistant:** Multiple reputation factors make network abuse expensive
- **Performance-oriented:** Resources are automatically allocated to reliable participants
- **Recovery-capable:** Reputation can be rebuilt through consistent good behavior

# 9 Security and Spam Prevention

The Hyperlink network implements a comprehensive security and spam prevention system that combines reputation tracking, transaction validation, and economic incentives. This multi-layered approach ensures network health while remaining permissionless and decentralized.

## 9.1 Reputation System Architecture

The system employs a dual-reputation model that separately tracks both wallet and peer behavior. This separation allows for fine-grained control over network participation while maintaining flexibility for different types of actors.

### 9.1.1 Wallet Reputation

Wallet reputation is calculated using a formula that considers both transaction history and economic stake:

$$R_w = \begin{cases} R_{max} & \text{if } B \geq P \\ R_i \cdot (1 + \frac{B}{P}) & \text{otherwise} \end{cases} \quad (4)$$

where:

- $R_w$  is the wallet's reputation score
- $R_{max}$  is the maximum reputation (1.0)
- $B$  is the wallet's current balance
- $P$  is the poverty line threshold (10 HYY)
- $R_i$  is the initial reputation score (1.0)

This formula creates several important security properties:

- Wealthy accounts automatically receive full trust, as they have economic stake
- New accounts receive a moderate initial trust level
- Trust scales with economic participation
- Malicious behavior results in reputation penalties

The implementation balances security with usability:

```
1 func (n *Node) GetWalletReputation(address string)
2   float64 {
3     balance := n.graph.State[address]
4     if balance >= PovertyLine {
5       return MaxReputation
6     }
7
8     reputation, exists := n.reputations[address]
9     if !exists {
10      balanceWeight := float64(balance) / float64(
11        PovertyLine)
12      return InitialReputation * (1 + balanceWeight)
13    }
14    return reputation
15  }
```

### 9.1.2 Transaction Validation Pipeline

Every transaction goes through a rigorous multi-stage validation process before being accepted into the network:

#### 1. Basic Structure Validation:

- Correct transaction format
- Required fields present
- Valid address formats
- Minimum fee requirements

#### 2. Cryptographic Validation:

- Transaction hash verification
- Signature authenticity
- Mining proof validation (if applicable)

#### 3. Economic Validation:

- Balance sufficiency
- Fee requirements
- Double-spend prevention

#### 4. Reputation-Based Validation:

- Sender reputation check
- Special rules for new accounts
- Minimum amount requirements for first transactions

The validation pipeline is implemented as a series of checks:

```
1 func (g *Graph) ValidateTransaction(tx Transaction)
2   error {
3     // Stage 1: Basic validation
4     if err := tx.ValidateBasic(); err != nil {
5       return err
6     }
7
8     // Stage 2: Cryptographic validation
9     if err := tx.ValidateCryptographic(); err != nil {
10      return err
11    }
12  }
```

```

12 // Stage 3: Reference validation
13 if err := g.ValidateReferences(tx); err != nil {
14     return err
15 }
16
17 // Stage 4: State validation
18 if err := g.ValidateState(tx); err != nil {
19     return err
20 }
21
22 return nil
23 }

```

### 9.1.3 Anti-Spam Mechanisms

The system implements multiple layers of spam prevention:

#### 1. Economic Barriers:

- Mandatory transaction fees (0.0001 HYY)
- All fees are burned, permanently reducing supply
- Higher minimum amounts for first transactions (0.0001 HYY)

#### 2. Reputation Requirements:

- Minimum reputation threshold for transaction propagation
- Reputation penalties for invalid transactions
- Gradual trust building through successful transactions

#### 3. Network-Level Protection:

- Peer reputation tracking
- Connection limits and pruning
- Bandwidth allocation based on peer reputation

### 9.1.4 Attack Resistance

The combination of economic incentives and reputation tracking creates strong resistance to common attack vectors:

#### 1. Sybil Attacks:

- New accounts have limited privileges
- Economic cost to create meaningful participation
- Reputation building requires consistent good behavior

#### 2. Spam Attacks:

- Fee burning makes sustained spam expensive
- Reputation loss for invalid transactions
- Network-level filtering of low-reputation actors

#### 3. Eclipse Attacks:

- Persistent connections to validator nodes
- Peer diversity requirements
- Regular peer rotation and health checks

This comprehensive security approach ensures that:

- The network remains open while resistant to abuse

- Economic incentives align with network health
- Bad actors face increasing costs and decreasing capabilities
- Legitimate users can easily participate and build trust
- The system can automatically adapt to changing threat patterns

## 9.2 Transaction Filtering

The Hyperlink network implements a transaction filtering system that acts as the first line of defense against spam and invalid transactions. This filtering happens before transactions enter the wider network, reducing unnecessary bandwidth usage and processing load.

### 9.2.1 Filtering Objectives

The transaction filtering system serves multiple purposes:

- Prevents obviously invalid transactions from propagating
- Enforces economic rules for network participation
- Protects against common spam patterns
- Provides special handling for new users
- Ensures efficient use of network resources

### 9.2.2 Core Filtering Rules

The system implements several key validation rules:

#### 1. Fee Validation:

- Every transaction must include a minimum base fee
- Fees are burned to create deflationary pressure

#### 2. Reference Requirements:

- Transactions must reference existing transactions
- Minimum reference count ensures DAG connectivity
- References must be within valid time window

#### 3. New User Protection:

- First-time transactions face stricter requirements.
- Mining proofs exempt from these restrictions

The filtering algorithm implements these rules in a systematic way:

### 9.2.3 Special Cases and Exceptions

The filtering system includes special handling for certain transaction types:

#### 1. Mining Rewards:

- Bypass normal reference requirements
- Must include valid mining proof
- Still subject to fee requirements

---

**Algorithm 2** Transaction Validation

```

1: if  $tx.Fee < BaseFee$  then return "fee too low"
2: end if
3: if  $len(tx.References) < MinReferences$  then return "insufficient references"
4: end if
5: if  $IsNewUser(tx.From)$  and  $tx.MiningProof = nil$  then
6:    $totalAmount \leftarrow 0$ 
7:   for each  $change$  in  $tx.Changes$  do
8:      $totalAmount \leftarrow totalAmount + change.Amount$ 
9:   end for
10:  if  $totalAmount < MinimumFirstTransactionAmount$  then return
    "first transaction amount too low"
11:  end if
12: end if

```

## 2. Validator Transactions:

- Priority processing
- Relaxed reference requirements
- Enhanced propagation rules

### 3. High-Value Accounts:

- Accounts above poverty line get preferential treatment
- Reduced filtering restrictions
- Faster propagation through network

### 9.2.4 Impact on Network Health

The filtering system provides several key benefits:

- **Resource Efficiency:**
  - Early rejection of invalid transactions
  - Reduced network bandwidth usage
  - Lower processing overhead on nodes
- **Spam Prevention:**
  - Economic barriers to entry
  - Progressive difficulty for new accounts
  - Automatic filtering of dust transactions
- **Network Stability:**
  - Controlled transaction flow
  - Predictable resource usage
  - Self-regulating participation rules

This filtering approach ensures that the network can maintain high performance and security while remaining accessible to legitimate users. The combination of economic incentives, progressive restrictions, and special case handling creates a robust first line of defense against network abuse.

## 10 The Hyyperlink Native Token

The Hyyperlink native token is the core utility token of the Hyyperlink network. It is used for transactions, mining rewards, and liquidity pool operations.

### 10.1 Initial Supply and Distribution

Hypperlink launches with an initial supply of 100 million tokens (100M HYY). While this provides the initial liquidity, the system implements a carefully designed mining mechanism to achieve targeted network growth and faster transaction finality:

- Initial supply of 100M HYY
- 2% annual inflation target through mining rewards
- Base mining reward of 1 atomic unit per successful mining proof
- Mining rewards increase with higher difficulty solutions

## 10.2 Mining and Network Growth

The mining system serves multiple purposes:

- Controls supply growth at 2% annually
- Mining volume accelerates transaction finality and graph convergence

Miners receive rewards based on difficulty:

- Base reward increases for each difficulty level above minimum
- Minimum difficulty requirement based on leading zeros in the hash of the winning mining proof pubkey
- Zero fees on mining transactions
- Difficulty is dictated by the validator nodes
- Rewards are automatically calculated and verified

The reward calculation follows:

$$reward = BaseReward \cdot 2^{(difficulty - MinDifficulty)} \quad (5)$$

Where:

- *BaseReward* is 1 unit (0.00000001 HYY)
- *MinDifficulty* is 6 leading zeros
- *difficulty* is the number of leading zeros in the hash

This creates an incentive structure where higher difficulties receive exponentially larger rewards:

- 6 zeros = 0.00000001 HYY
- 7 zeros = 0.00000002 HYY
- 8 zeros = 0.00000004 HYY

### 10.3 Special Addresses

The system utilizes two special-purpose addresses for token management:

### 10.3.1 Token Grant Address

The Token Grant address  
(111)  
serves as the source of protocol-controlled token operations:

- Initial supply distribution
- Protocol reward distribution
- System-level operations
- Special transaction handling

### 10.3.2 Burn Address

The Burn address (00) permanently removes tokens from circulation:

11.1

- Receives all transaction fees
- Cannot send tokens
- Tokens sent here are permanently removed from circulation
- Balance increases but never decreases

## 10.4 Token Units

Each Hyyperlink token (HYY) is divisible into 100,000,000 units:

- 1 HYY = 100,000,000 units
- Smallest unit = 0.00000001 HYY
- All internal calculations use 64bit unsigned integer units
- Display values are converted to decimal form

For example:

- 1.00000000 HYH = 100,000,000 units
- 0.00000001 HYH = 1 unit
- 123.45678900 HYH = 12,345,678,900 units

## 10.5 Supply Mechanics

The Hyyperlink token supply is governed by two opposing forces:

### 10.5.1 Protocol Operations

Assets flow through the system via protocol-level operations:

- Liquidity pool operations
- Swap fee collection and distribution
- Token creation, minting and burning
- NFT creation, minting and burning
- Protocol reward mechanisms
- System-level transactions

### 10.5.2 Supply Deflation

Assets are removed from circulation through:

- Transaction fees sent to Burn address
- Protocol operation fees
- Optional asset burning mechanisms
- Permanent removal from supply

This dual mechanism creates a dynamic balance:

- Protocol operations manage token distribution
- Transaction activity gradually decreases supply
- Network usage directly influences token economics
- Natural equilibrium between utility and scarcity

## 11 Native DeFi Capabilities

## 11.1 Protocol-Level Financial Primitives

Hyperlink implements core DeFi functionality directly in the protocol layer, eliminating the need for external smart contracts or Layer 2 solutions. This native implementation provides several key advantages:

- Direct token-to-token swaps without wrapping or bridging
- Atomic execution of complex DeFi operations
- Reduced transaction costs and latency
- Enhanced security through protocol-level validation
- Simplified user experience

## 11.2 Liquidity Pools and AMM

The system implements an Automated Market Maker (AMM) model with the following features:

- Native constant product market making formula
- Dynamic fee adjustment based on pool metrics
- Efficient price discovery mechanism
- Liquidity provider incentives
- Multi-token pool support

The AMM implementation follows the formula:

$$x \cdot y = k \quad (6)$$

Where:

- $x$  and  $y$  are the reserves of two tokens in a pool
- $k$  is the constant product maintained by the pool

### 11.3 Swap Mechanics

Token swaps are executed directly within the protocol through a series of atomic operations:

1. Pool reserve verification
2. Price impact calculation
3. Slippage protection check
4. Fee computation and distribution
5. Reserve update and token transfer

The swap amount calculation follows:

$$dy = \frac{y \cdot dx}{x + dx} \cdot (1 - f) \quad (7)$$

Where:

- $dx$  is the input token amount
- $dy$  is the output token amount
- $f$  is the pool fee rate



## 11.4 Liquidity Provider Operations

Liquidity providers can participate in pools through:

- Pool creation with initial liquidity
- Liquidity addition to existing pools
- Proportional liquidity removal
- Fee collection from swap operations

Provider shares are tracked using:

$$s = L \cdot \sqrt{\frac{x \cdot y}{X \cdot Y}} \quad (8)$$

Where:

- $s$  is the share tokens minted
- $L$  is the total existing shares
- $x, y$  are the deposited amounts
- $X, Y$  are the current pool reserves

## 12 Conclusion

The cryptocurrency landscape has long been defined by its limitations: networks that slow under load, DeFi constrained by smart contract overhead, and networks shackled by their own design. Hyperlink breaks free from these constraints by challenging their underlying assumptions.

By recognizing that financial primitives belong in the protocol layer, not in smart contracts, we've eliminated an entire class of complexity and risk. Our DAG architecture transforms high transaction volume from a liability into an asset—the more people use the network, the faster it becomes. The reputation system creates natural economic barriers to abuse while keeping the network permissionless, solving the volume overload problem without sacrificing decentralization.

These architectural choices compound to create emergent properties that transcend their individual benefits:

- When DeFi operations execute at the protocol level, and the network accelerates with usage, we achieve a system where peak demand improves performance instead of degrading it
- When reputation governs transaction propagation, and fees scale with operation complexity, we create an economic ecosystem that naturally resists abuse while rewarding legitimate usage
- When mining rewards scale exponentially with difficulty, and validators control difficulty requirements to maintain annual inflation targets, we establish a dynamic supply mechanism that adapts to network longevity needs.

The result is more than just another cryptocurrency—it's a fundamental rethinking of how decentralized networks should operate. Hyperlink demonstrates that we can build systems where security doesn't come at the cost of performance, where complexity doesn't require layers of abstraction, and where decentralization doesn't mean sacrificing efficiency.

This is the future of decentralized finance: a network that gets faster as more people join, a protocol

that executes complex operations without smart contract overhead, and an ecosystem that maintains security through economic incentives rather than artificial constraints. Hyperlink isn't just an improvement on existing systems—it's a blueprint for what they should have been all along.