# 动态规划 Dynamic Programming



#### 本节内容

- 1. 递推! 而递归+记忆化是少数情况
- 2. 状态的定义 fib[n]
- 3. 最优子结构 opt[n] = best\_of(opt[n-1], opt[n-2], ...)
- 4. 状态转移方程 (DP方程)



#### 要点

- 1. 递推! 而递归+记忆化是少数情况
- 2. 状态的定义 fib[n]
- 3. 最优子结构 opt[n] = best\_of(opt[n-1], opt[n-2], ...)
- 4. 状态转移方程 (DP方程)



• Fibonacci array

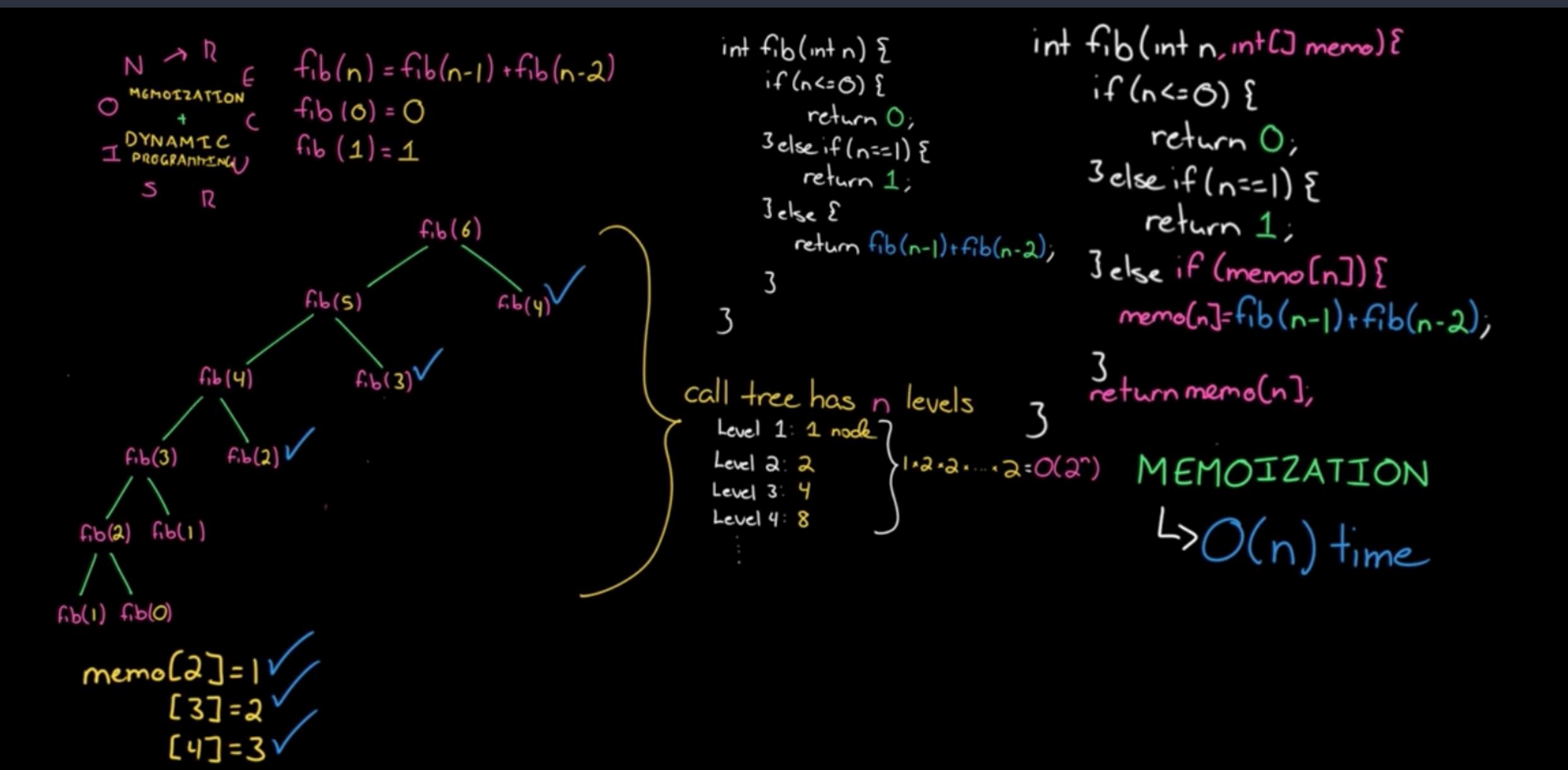
• 0, 1, 1, 2, 3, 5, 8, 13, 21, ...

• F[n] = F[n-1] + F[n-2]

```
int fib (int n) {
                      fib(n) = fib(n-1) + fib(n-2)
                                                                   if (n<=0) {
      MEMOTZATION
   0
                      fib (0) = 0
                                                                       return 0;
     DYNAMIC
                                                                  3 else if (n==1) {
                      fib (1) = 1
   I PROGRAMMENGU
                                                                      return 1;
       S
                                                                  Jelse E
                                  t.p(6)
                                                                      return fib (n-1)+fib(n-2);
                      fib(5)
                                         F. P(A)
                                                f.b(2)
            fib (4)
                                        fib(3)
                            f(9)
                                                           call tree has n levels
                                                              Level 1: 1 node 7
     f.b(3)
               f.b(2)
                                                                               11222 = O(2")
                                                              Level 2: 2
                                                              Level 3: 4
                                   t.P(1) t.P(0)
                                                              Level 4: 8
fib(2) fib(1) fib(1) fib(0)
(P(1) tP(0)
```

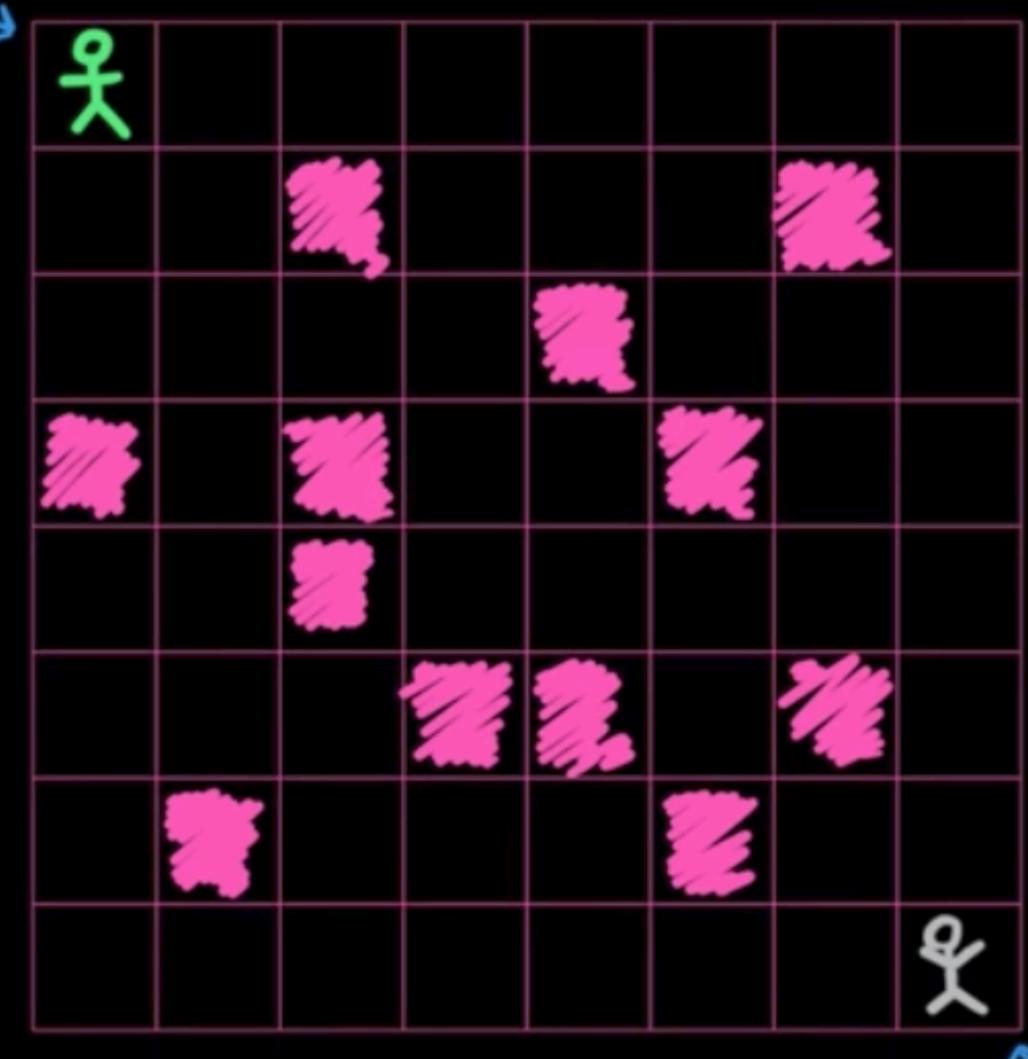
```
int fib(int n) {
return n <= 1 ? n : fib(n - 1) + fib(n - 2);</li>
```

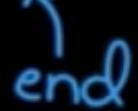
```
int fib (int n) {
                      fib(n) = fib(n-1) + fib(n-2)
                                                                   if (n<=0) {
      MEMOTZATION
   0
                      fib (0) = 0
                                                                       return 0;
     DYNAMIC
                                                                  3 else if (n==1) {
                      fib (1) = 1
   I PROGRAMMENGU
                                                                      return 1;
       S
                                                                  Jelse E
                                  t.p(6)
                                                                      return fib (n-1)+fib(n-2);
                      fib(5)
                                         F. P(A)
                                                f.b(2)
            fib (4)
                                        fib(3)
                            f(9)
                                                           call tree has n levels
                                                              Level 1: 1 node 7
     f.b(3)
               f.b(2)
                                                                               11222 = O(2")
                                                              Level 2: 2
                                                              Level 3: 4
                                   t.P(1) t.P(0)
                                                              Level 4: 8
fib(2) fib(1) fib(1) fib(0)
(P(1) tP(0)
```



[5]=5

# COUNT THE PATHS



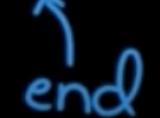




```
COUNT THE PATHS
        MGMOTZATTON
       DYNAMIC
               paths(start, end) =
                      + paths(B,ena)
   paths(A,en&)
paths(D,end)+paths(C,end) paths(C,end)+paths(E,end)
         int countPaths(boolean[][] grid, int row, int col) {
           if (!validSquare(grid, row, col)) return 0;
           if (isAtEnd(grid, row, col)) return 1;
           return countPaths(grid, row+1, col) + countPaths(grid, row, col+1);
```

# COUNT THE PATHS columns (x) start

rows





# COUNT THE PATHS columns (x)

	A	<b>?</b>	Bol	E					
rows	D	•							
(g)									
									ı
		١	ı	1	1	1	1	ı	<b>Y</b>





$$opt[i, j] = opt[i + 1, j] + opt[i, j + 1]$$

\_\_\_\_\_\_

$$opt[i, j] = opt[i + 1, j] + opt[i, j + 1]$$

else: // 石头

$$opt[i, j] = 0$$

# COUNT THE PATHS columns (x)

웃	Bot	• E					
A옷-	c <sup>*</sup>						
D							
		3	a	ı		2	ı
١	١	3	١	١	1	ı	4

rows





### COUNT THE PATHS

۵			L
0	О	r	Т
		u	
		\	

	,27						
옷	17	12	12	7	4	١	ı
10	5		5	3	3		1
5	S	2	2		3	3	١
	3		2	١		2	ı
7	3		ı	١	ı	ı	1
4	3	3			0		ı
		3	a	ı		2	١,
١	ı	ı	١	1	1	ı	ا ا





#### 动态规划 Dynamic Programming

- 1. 递推! (recursion)
- 2. 状态的定义 fib[n]
- 3. 最优子结构 optimized = fib[n]
- 4. 状态转移方程(DP方程: fib[n] = fib[n-1] + fib[n-2])



#### 动态规划 Dynamic Programming

- 1. 打破自己的思维惯性, 形成机器思维
- 2. 理解复杂逻辑的关键
- 3. 也是职业进阶的要点要领



#### 01背包问题

假设山洞里共有a,b,c,d,e这5件宝物(不是5种宝物),它们的重量分别是2,2,6,5,4,它们的价值分别是6,3,5,4,6,现在给你个承重为10的背包,怎么装背包,可以才能带走最多的财富。

01背包的状态转换方程

 $f[i,j] = Max\{ f[i-1,j-Wi]+Pi(j >= Wi), f[i-1,j] \}$ 

f[i,j]表示在前i件物品中选择若干件放在承重为 j 的背包中,可以取得的最大价值。

Pi表示第i件物品的价值。

决策:为了背包中物品总价值最大化,第i件物品应该放入背包中吗?



## 预习题目 (上课考察)

- https://leetcode.com/problems/best-time-to-buy-and-sellstock/#/description
- 2. https://leetcode.com/problems/climbing-stairs/description/
- 3. https://leetcode.com/problems/triangle/description/
- 4. https://leetcode.com/problems/maximum-product-subarray/description/
- 5. https://leetcode.com/problems/coin-change/description/



## 预习题目 (上课考察)

- 1. https://leetcode.com/problems/house-robber/
- 2. https://leetcode.com/problems/house-robber-ii/description/
- 3. https://leetcode.com/problems/best-time-to-buy-and-sell-stock/#/description
- 4. https://leetcode.com/problems/best-time-to-buy-and-sell-stock-ii/
- 5. https://leetcode.com/problems/best-time-to-buy-and-sell-stock-iii/
- 6. https://leetcode.com/problems/best-time-to-buy-and-sell-stock-with-cooldown/
- 7. https://leetcode.com/problems/best-time-to-buy-and-sell-stock-iv/
- 8. https://leetcode.com/problems/best-time-to-buy-and-sell-stock-with-transaction-fee/



# 实战题目

- 1. <a href="https://leetcode-cn.com/problems/perfect-squares/">https://leetcode-cn.com/problems/perfect-squares/</a>
- 2. https://leetcode-cn.com/problems/burst-balloons/
- 3. https://leetcode-cn.com/problems/jump-game/
- 4. https://leetcode-cn.com/problems/jump-game-ii/
- 5. https://leetcode-cn.com/problems/unique-paths/
- 6. https://leetcode-cn.com/problems/unique-paths-ii/
- 7. https://leetcode-cn.com/problems/unique-paths-iii/
- 8. https://leetcode-cn.com/problems/coin-change/
- 9. https://leetcode-cn.com/problems/coin-change-2/
- 10.



# 实战题目

- 1. https://leetcode-cn.com/problems/longest-valid-parentheses/
- 2. https://leetcode-cn.com/problems/minimum-path-sum/
- 3. https://leetcode-cn.com/problems/edit-distance/
- 4. https://leetcode-cn.com/problems/decode-ways
- 5. https://leetcode-cn.com/problems/maximal-square/
- 6. https://leetcode-cn.com/problems/max-sum-of-rectangle-no-larger-than-k/
- 7. https://leetcode-cn.com/problems/frog-jump/
- 8. https://leetcode.com/problems/split-array-largest-sum
- 9. https://leetcode-cn.com/problems/student-attendance-record-ii/
- 10.https://leetcode-cn.com/problems/task-scheduler/
- 11.https://leetcode-cn.com/problems/palindromic-substrings/
- 12.https://leetcode-cn.com/problems/minimum-window-substring/



#