

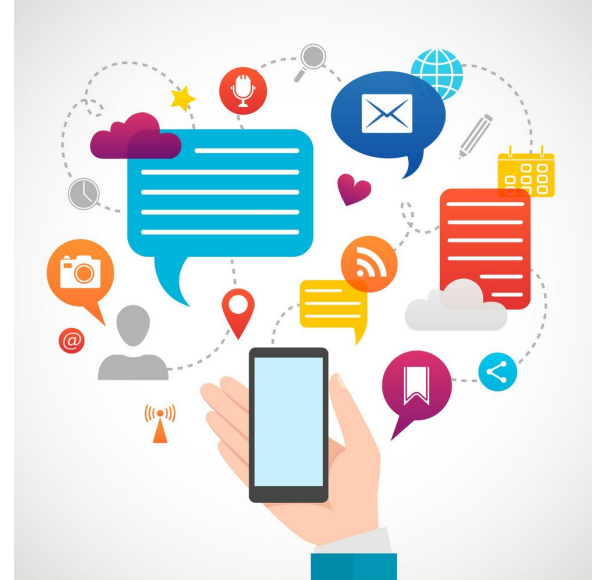


Parallel Betweenness Centrality Algorithm (ParBC)

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Introduction

With the popularization of Smart Devices and social applications, a huge quantity of vast social network are available



Problem Definition and Dataset

To measure how central is a node...

Betweenness Centrality is formulated:

$$Betweenness(k) = \sum_{i \neq k \neq j} \left(\frac{\sigma_{i,j}(k)}{\sigma_{i,j}} \right)$$



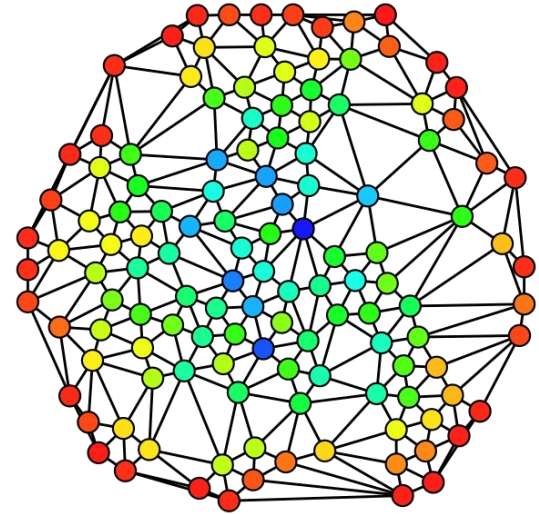
Algorithm Design

Originally proposed:

1. Calculate all pairwise shortest path length for all node pairs in G
2. Calculate betweenness Centrality. Inspired by Floyd-Warshall algorithm.

New proposed algorithm:

Direct calculations, easily implemented in pure parallel manner





Enhanced algorithm and parallel implementation

$$\sigma_{i,j} = \sum_{w \in \text{pred}_i(j)} \sigma_{i,w}$$

$$\delta_s(v) = \sum_{w \in \text{Pred}_s(v)} \sum_{t \in V} \delta_{s,t}(v, \{v, w\}) = \sum_{w \in \text{Pred}_s(v)} \left(\frac{\sigma_{s,v}}{\sigma_{s,w}} + \sum_{t \in V, t \neq w} \frac{\sigma_{s,v}}{\sigma_{s,w}} \frac{\sigma_{s,t}(w)}{\sigma_{s,t}} \right) = \sum_{w \in \text{Pred}_s(v)} \frac{\sigma_{s,v}}{\sigma_{s,w}} (1 + \delta_s(w))$$

$$\text{Betweenness}(v) = \sum_{s \in V} \delta_s(v)$$



Enhanced algorithm and parallel implementation

```
 $G \leftarrow$  the graph  
 $Betweenness \leftarrow \{n : 0 | \forall n \in G\}$   
for  $s$  in  $G$  do  
   $pred \leftarrow \{n : [ ] | \forall n \in G\}$ ;  $dist \leftarrow \{n : -1 | \forall n \in G\}$ ;  $sigma \leftarrow \{n : 0 | \forall n \in G\}$ ;  $S \leftarrow [ ]$   
   $dist[s] = 0$ ;  $sigma[s] = 1$   
   $queue \leftarrow [s]$   
  while  $queue$  is not empty do  
     $v = queue.get()$   
     $S.append(v)$   
    for  $w$  in  $G.neighbours(v)$  do  
      if  $dist[w] == -1$  then  
         $dist[w] = dist[v] + 1$   
         $queue.put(w)$   
      end if  
      if  $dist[w] == dist[v] + 1$  then  
         $sigma[w] += sigma[v]$   
         $pred[w].append(v)$   
      end if  
    end for  
   $delta \leftarrow \{n : 0 | \forall n \in G\}$   
  for  $w$  in  $reverse(S)$  do  
    for  $v$  in  $pred[w]$  do  
       $delta[v] += sigma[v] / sigma[w] * (1 + delta[w])$   
    if  $w \neq s$  then  
       $Betweenness[w] += delta[w] / 2$ 
```



Evaluation (correctness)

$$< 10^{-13}$$



Evaluation (efficiency)

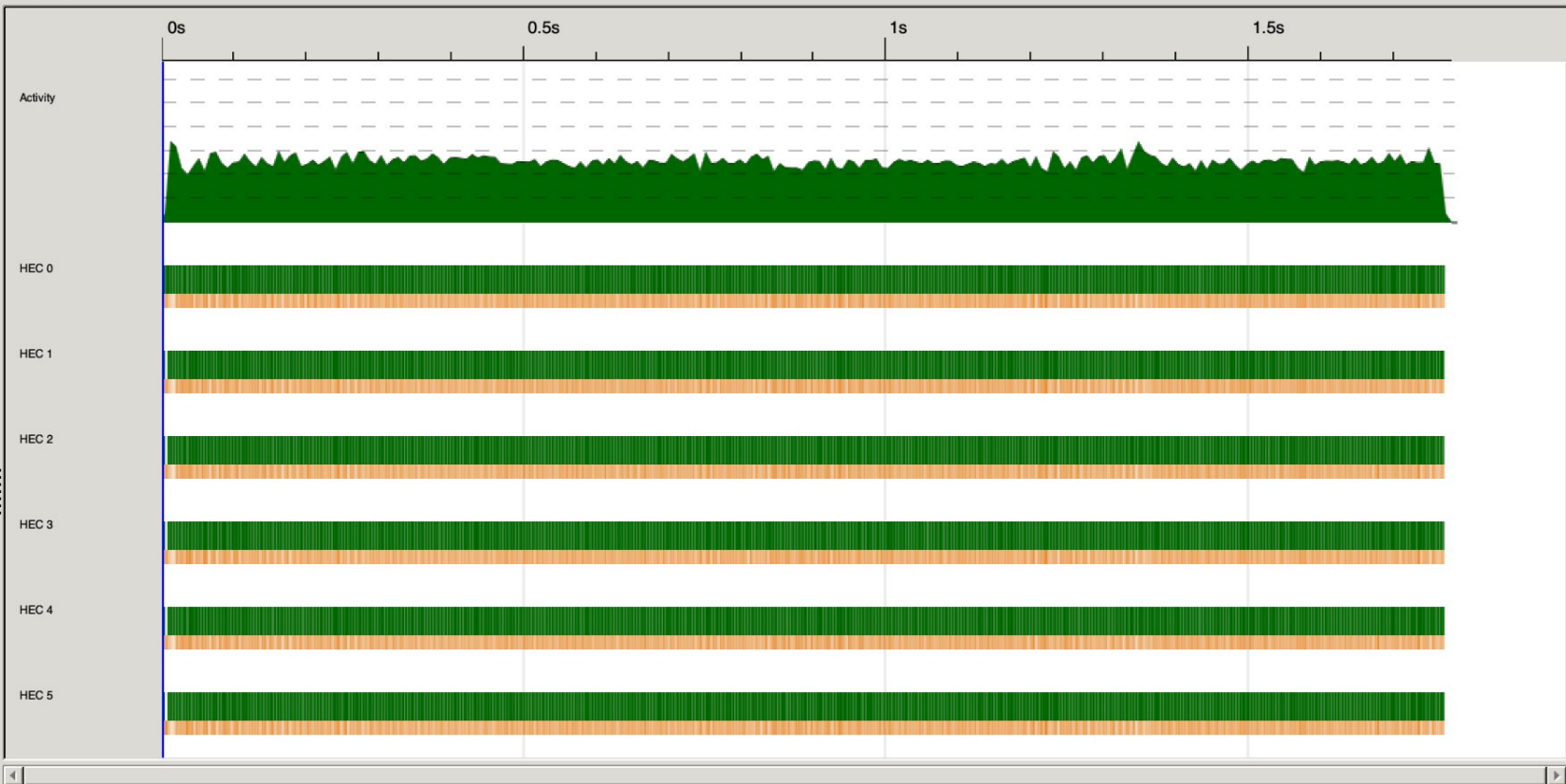
Table 3: Experiment Results for Simulate1000						
N	time(s)	converted	gc'd	fizzled	total	Speedup
seq	6.012	N/A	N/A	N/A	N/A	N/A
2	3.620	1998	1	1	2000	X1.66
3	2.590	1998	1	1	2000	X2.32
4	2.110	1998	1	1	2000	X2.85
5	1.920	1998	1	1	2000	X3.13
6	1.780	1998	1	1	2000	X3.38



Key Traces Bookmarks

Timeline

- running
- GC
- GC waiting
- create thread
- seq GC req
- par GC req
- migrate thread
- thread wakeup
- shutdown
- user message
- perf counter
- perf tracepoint
- create spark
- dud spark
- overflowed spark
- run spark
- fizzled spark
- GCed spark



Time Heap GC Spark stats Spark sizes Process info Raw events

Total time: 1.780s
Mutator time: 895.437ms
GC time: 884.992ms



Evaluation

Sounds inspiring!

Move to the social network of GitHub



Evaluation and Benchmark

Table 5: Experiment Results for SNAP		
Method	time	Speedup
seq	44h 43 mins	N/A
ParBC-N6	18h 38 mins	X2.40
NetworkX	16h 23 mins	X2.73





Thank you