H5 页面漏洞挖掘之路 - 混 淆篇 - SecPulse.COM | 安全 脉搏

66 复制代码混淆过后:

前言

针对上次我们提交漏洞之后,我们再次查看 JS 代码,定位加密函数和解密函数的位置,发现已经不是赤裸裸没有任何防护,而是已经进行的了 JS 混淆,接下来我们针对遇到 JS 混淆后,我们该如何破解 JS 混淆后的代码进行加解密,继续进行渗透测试。笔者在这里提供一个思路和方法。

前置知识

首先我们先了解下代码混淆的具体原理是什么? 其实很简单,就是去除代码中尽可能多的有意义的信息,比如注释、换行、空格、代码负号、变量重命名、属性重命名(允许的情况下)、无用代码的移除等等。因为代码是公开的,我们必须承认没有任何一种算法可以完全不被破解,所以,我们只能尽可能增加攻击者阅读代码的成本。

我将混淆类型分为两类:

变量名混淆

将变量名混淆成阅读比较难阅读的字符,增加代码阅读难度,而现在大部分厂商的混淆,都会将其混淆成 16 进制变量名。

效果如下:

```
`var test = 'helloworld';`
```

混淆后:

```
`var _0x7deb = 'helloworld';`
```

常量提取

将 JS 中的常量提取到数组中,调用的时候用数组下标的方式调用,这样的话直接读懂基本不可能了,要么反 AST 处理下,要么一步一步调试,工作量大增。

以上面的代码为例:

```
`var test = 'helloworld';`
```

复制代码混淆过后:

每个文件开头会有一个很长的字符数组,然后会有一段代码对这个数组进行加工,然后还有一个函数接收一个或两个参数输出一个字符串,这个字符串更接近原始的代码。将常量进行加密处理,上面的代码中,虽然已经是混淆过后的代码了,但是 helloworld 字符串还是以明文的形式出现在代码中,例如将关键字进行 Unicode16 进制编码。如下:

```
`var test = 'helloworld';`
```

结合常量提取得到混淆结果:

```
var _0x9d2b = ['\x68\x65\x6c\x6c\x6f'];

var _0xb7de = function (_0x4c7513) {
    __0x4c7513 = _0x4c7513 - 0x0;
    var _0x96ade5 = _0x9d2b[_0x4c7513];
    return _0x96ade5;
};

var test = _0xb7de('0x0');
```

第一部分: 变量名称存储数组

这里存储了一些在函数中用到的变量和字符串。

```
1 var _0x2ec2 = [
2 'UGtjczc=',
3 'dG9TdHJpbmc=',
4 'ZGVjcnlwdA==',
5 'c3RyaW5naWZ5',
6 'xxxx',
7 'bW9kZQ==',
8 'Q0JD',
9 'cGFk'
10 ];
```

第二部分 数组处理函数

```
1 /**
2 * params _0x167407: 上面的字符串数组
3 * params _0x353595: 计数个数
4 * 把前 _0x353595 +1 个元素放到数组末尾
5 */
6 (function (_0x167407, _0x353595) {
7 var _0x52a3ae = function (_0x3fbe47) {
8 while (--_0x3fbe47) {
9 _0x167407['push'](_0x167407['shift']());
10 }
11 };
12 _0x52a3ae(++_0x353595);
13 }(_0x2ec2, 312));
```

第三部分 数组字符串处理函数

```
// 这个是数组内容解码的函数,实际上第二个参数是没有用到的
var _0x523d = function (_0x4c10d0, _0x393bf7) {
 _0x4c10d0 = _0x4c10d0 - 0; // 这里第一个参数是通过字符串
 var _0x70d87b = _0x2ec2[_0x4c10d0]; // 这里 _0x70d87
 // 接下来判断有没有进行过初始化操作,如果没有的话,先初始化
 if (_0x523d['CuFQcU'] === undefined) {
  (function () {
 var _0x5b57a4 = function () {
  var _0x29e588;
  try {
    _0x29e588 = Function('return (function() ' + '{}.
  } catch (_0x4956c9) {
    _0x29e588 = window;
  return _0x29e588;
  };
 var _0x2b121a = _0x5b57a4(); // 这里实际上返回的是 Wind
 var _0x6c99b9 = 'ABCDEFGHIJKLMNOPQRSTUVWXYZabcdefghi
 // 下面这个是判断Window有没有atob这个函数,如果没有的话生成
 _0x2b121a['atob'] | (_0x2b121a['atob'] = function (
   var _0x901f5e = String(_0x13f6f4) ['replace'](/=+5
  for (var _0x240979 = 0, _0x43e3e8, _0x42ec25, _0x6e)
     _0x42ec25 = _0x6c99b9['index0f'](_0x42ec25);
   return _0x1c0a86;
```

```
});
  }());
  _0x523d['ZEesoG'] = function (_0x1de802) {
    var _0x216ff1 = atob(_0x1de802);
    var _0x42331f = [
    ];
    for (var _0x3a392f = 0, _0x2319db = _0x216ff1['le
     _0x42331f += '%' + ('00' + _0x216ff1['charCode/
    }
    return decodeURIComponent(_0x42331f);
   // 到这里完成初始化操作,置CuFQcU为true,添加VgXLDn属
    _0x523d['VgXLDn'] = {};
    _0x523d['CuFQcU'] = !![];
  }
  // 后面这段是先判断之前有没有对传入的参数进行解密过, 如果解释
  var _0x22ee7f = _0x523d['VqXLDn'][_0x4c10d0];
  if (_0x22ee7f === undefined) {
  _0x70d87b = _0x523d['ZEesoG'](_0x70d87b);
  _0x523d['VgXLDn'][_0x4c10d0] = _0x70d87b;
  } else {
     _0x70d87b = _0x22ee7f;
   return _0x70d87b;
};
```

第四部分 加解密函数

```
function encrypt(_0xd0a5dd) {
  var _0x2d682e = CryptoJS[_0x523d('0x0')][_0x523d('0x1
  var _0x2d053c = CryptoJS[_0x523d('0x0')][_0x523d('0x1
  var _0xa5c781 = CryptoJS[_0x523d('0x0')][_0x523d('0x1
  var _0x17d14e = CryptoJS[_0x523d('0x5')][_0x523d('0x6')]
  'iv': _0x2d053c,
  'mode': CryptoJS[_0x523d('0x7')][_0x523d('0x8')],
  'padding': CryptoJS[_0x523d('0x9')][_0x523d('0xa')]
  });
  return _0x17d14e[_0x523d('0xb')]();
}
function decrypt(_0x363945) {
  var _0x41412c = CryptoJS[_0x523d('0x0')][_0x523d('0x0')]
  var _0xf43728 = CryptoJS[_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('0x0')][_0x523d('
```

```
var _0x2f2c26 = CryptoJS[_0x523d('0x5')][_0x523d('0)
'iv': _0xf43728,
  'mode': CryptoJS[_0x523d('0x7')][_0x523d('0x8')],
  'padding': CryptoJS[_0x523d('0x9')][_0x523d('0xa')]
});
return CryptoJS[_0x523d('0x0')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x523d('0x1')][_0x52
```

当我们分析整个混淆后的代码后,我们可以手动断点调试,来看看具体的解密之后每参数是什么。我们首先将整个混淆后的 js 代码 copy 下来,定义 main()函数,调用加密 encrypt 和 decrypt 解密这两个函数,在浏览器下调试运行。

```
| Comparison | Com
```

代码完美运行,在第三部分数组字符串处理函数的位置我们手动断点 F10 进行调试。

密钥 key 成功拿到:

得知加密算法为 AES:

```
| Description |
```

AES 加密算法使用的填充方式: Pkcs7

总结

JS 混淆在安全对抗中必不可少,一是对保护前端页面的代码逻辑,二是对前端登陆的算法密钥和向量 Ⅳ 进行保护。而我们通过反混淆还原代码或者直接调用混淆后的 JS 代码进行调试,获取密钥和向量 Ⅳ,从而达到解密密文,篡改数据包继续进行漏洞挖掘。

参考

https://www.52pojie.cn/thread-1104122-1-1.html#2985657

全文完

本文由 简悦 SimpRead 优化,用以提升阅读体验

使用了全新的简悦词法分析引擎 beta, 点击查看详细说明



