

Sep 23: Change to ELB LTL

Recall: (P)  $\inf_{x \in X} f(x)$

(4) Quadratic Programming (QP)

$$f(x) = \sum_{i,j=1}^n Q_{ij} x_i x_j = \underbrace{\bar{x}}_{1 \times n} \underbrace{Q}_{n \times n} \underbrace{x}_{n \times 1}, \quad Q = [Q_{ij}] \in \mathbb{R}^{n \times n}$$

homogeneous quadratic:  $f(\alpha x) = \alpha^2 f(x)$

$X$ : Same as LP

\* Compare quadratic constrained QP (QCQP)

$$\begin{aligned} \inf \quad & \bar{x}^T Q x \\ \text{s.t.} \quad & \bar{x}^T A_i x \leq b_i, \quad i=1, \dots, m. \end{aligned}$$

\* Note:  $Q$  can be assumed WLOG to be symmetric when defining  $f$ . Indeed, observe

$$\bar{x}^T Q x = \bar{x}^T \left( \underbrace{\frac{Q+Q^T}{2}}_{\text{symmetric}} \right) x$$

$$\begin{aligned} & \bar{x}^T \frac{Q}{2} x + \bar{x}^T \frac{Q^T}{2} x \\ &= \frac{1}{2} \bar{x}^T Q x + \frac{1}{2} (\bar{x}^T Q^T x)^T \\ &= \bar{x}^T Q x \end{aligned}$$

(5) Semidefinite Programming (SDP)

Definition: Let  $Q \in \underline{S}^n$ . Then, the following are equivalent:

Set of  $n \times n$  real  
Sym. matrices

- (1)  $Q$  is positive semidefinite (psd)
- (2)  $\forall x \in \mathbb{R}^n : x^T Q x \geq 0$
- (3) All eigenvalues of  $Q$  are non-negative.

Let  $C, A_1, \dots, A_m \in \underline{S}^n$ ;  $b_1, \dots, b_m \in \mathbb{R}$  be given.

$$\begin{aligned} \inf \quad & b^T y \\ \text{(SDP)} \quad \text{s.t.} \quad & C - \sum_{i=1}^m y_i A_i \succeq 0 \quad (*) \quad (\text{linear matrix} \end{aligned}$$

$$(SDP) \quad \text{s.t.} \quad C - \sum_{i=1}^m y_i A_i \succeq 0 \quad (*) \quad \text{(linear matrix inequality (LMI))}$$

$y \in \mathbb{R}^m$   
is psd

### Remarks

(1) (\*) is equivalent to  $\underbrace{-\sum_{i=1}^m y_i A_i}_{\triangleq M(y)} \succeq -C$

$\mathbb{R}^m \ni y \mapsto M(y) \in S^n$

Note:  $M(\cdot)$  is a linear map:  $\forall \alpha, \beta \in \mathbb{R}, y, z \in \mathbb{R}^m$ ,  
 $M(\alpha y + \beta z) = \alpha M(y) + \beta M(z)$  (exercise)

(2) Suppose that  $C, A_1, \dots, A_m$  are diagonal. Then, (\*) becomes

$$\begin{bmatrix} c_1 & & \\ & \ddots & \\ & & c_n \end{bmatrix} - \sum_{i=1}^m y_i \begin{bmatrix} a_{i1} & & \\ & \ddots & \\ & & a_{in} \end{bmatrix} \succeq 0$$

$$\Rightarrow \begin{bmatrix} c_1 - \sum_i y_i a_{i1} & & \\ & \ddots & \\ & & c_n - \sum_i y_i a_{in} \end{bmatrix} \succeq 0$$

$$\Leftrightarrow c_j - \sum_i y_i a_{ij} \geq 0 \quad \text{for } j=1, \dots, n. \quad \text{linear inequalities}$$

(3) How about more LMIs in the constraint?

e.g.

$$\begin{aligned} C_1 - \sum_{i=1}^m y_i A_i \succeq 0 \\ C_2 - \sum_{i=1}^m y_i B_i \succeq 0 \end{aligned} \quad \text{(exercise)} \quad \Leftrightarrow \begin{bmatrix} C_1 \\ C_2 \end{bmatrix} - \sum_{i=1}^m y_i \begin{bmatrix} A_i \\ B_i \end{bmatrix} \succeq 0$$

(Hint: Let  $A = \begin{bmatrix} A_1 & 0 \\ 0 & A_2 \end{bmatrix} \in S^n$ . Then,  $A \succeq 0$  iff  $A_1, A_2 \succeq 0$  (why?))

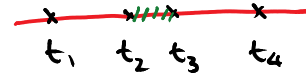
### Reformulation Examples

(1) Air Traffic Control

- n planes arriving
- $i^{\text{th}}$  plane arrives within  $[a_i, b_i]$
- assume that the planes land in order

- Let  $t_i$  be the assigned landing time of  $i^{\text{th}}$  plane
- For Safety, define the shortest metering time as

$$\min_{1 \leq i \leq n-1} \{t_{i+1} - t_i\}$$



- Goal: maximize shortest metering time



$$\max \quad \min_{1 \leq i \leq n-1} \{t_{i+1} - t_i\}$$

$$\text{s.t.} \quad \left. \begin{array}{ll} a_i \leq t_i \leq b_i & i=1, \dots, n \\ t_i \leq t_{i+1} & i=1, \dots, n-1 \end{array} \right\} \begin{array}{l} \text{linear} \\ \text{constraints} \end{array}$$

$\Leftrightarrow$

$$\max \quad z$$

$$\text{s.t.} \quad a_i \leq t_i \leq b_i, \quad t_i \leq t_{i+1}$$

$$z \stackrel{\text{can be replaced by } z}{=} \min_{1 \leq i \leq n-1} \{t_{i+1} - t_i\}$$

At optimality, equality is achieved (why?)

$\Leftrightarrow$

$$\max \quad z$$

$$\text{s.t.} \quad a_i \leq t_i \leq b_i, \quad t_i \leq t_{i+1}$$

$$z \leq t_{i+1} - t_i$$

$\left. \begin{array}{l} \text{LP} \\ \text{in variables} \\ t_1, \dots, t_n, z \end{array} \right\}$

## (2) Data Fitting

data points  $(a_i, b_i) \in \mathbb{R}^n \times \mathbb{R}; \quad i=1, \dots, m.$



Need to restrict the class of functions for the black box

Typical choice: affine function

$$f(y) = \underbrace{\bar{x}^T y}_{\text{linear}} + \underbrace{t}_{\text{constant}} \quad \begin{array}{l} x, y \in \mathbb{R}^n \\ t \in \mathbb{R} \end{array}$$

Goal: Find  $(x, t)$  s.t.  $f(a_i) \cong b_i \quad \forall i$

e.g.  $n=1$

