

## UNIT – II LINEAR DATA STRUCTURES

### Lecture NO. 8

## Queue

Queue is that type of list in which insertion can take place from one end called (Rear) and deletion can take place from another end called (Front).

Queue works on the principle of first in first out (FIFO) i.e

List come First Serve and Last in last out (LIFO).

### Daily Life Example

- i) The number of cars at police check
- ii) People purchase tickets from cinema make a queue.
- iii) People deposits their bills in bank make a queue.



### Computer Based Example

- i) When number of instructions are given for printing to computer, then that instructions will be executed first that was first assign to computer.

### Operations on Queue

Different operations can be performed at queue as

- 1) Create Queue  
By using this operation, the structure of the queue is generated.
- 2) Enqueue  
The insertion of element in the queue is called enqueue.
- 3) Dequeue  
The deletion of element from the queue is called dequeue.
- 4) Empty Queue

It returns Boolean value true, if queue is empty.

5) Full Queue

It returns Boolean value false, if queue is full.

6) Destroy Queue

This operation deletes all the elements of the queue and also deletes the structure of queue from the computer memory.

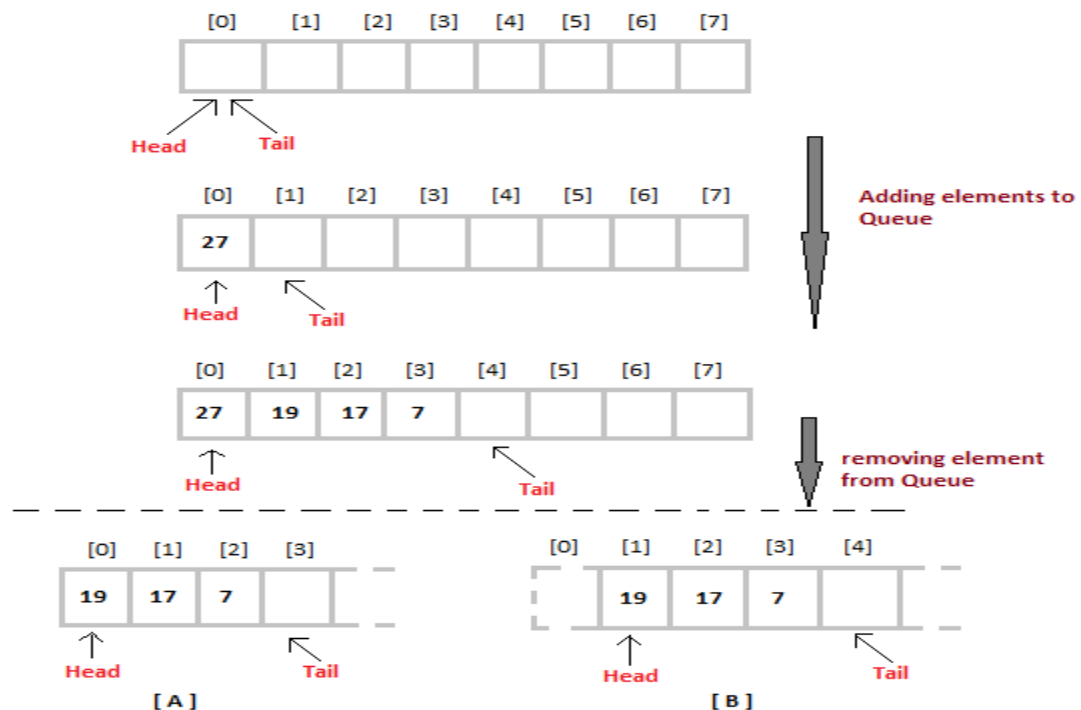
### Major Operation on Queue

1) Overflow Condition

When no more elements can be inserted in the queue then this condition is called overflow condition.

2) Underflow Condition

When queue is empty, and no more elements can be deleted from the queue then this condition is called underflow condition.



### ALGORITHM FOR THE INSERTION OF ELEMENT IN THE QUEUE

Algorithm enqueue (Q, N, item, Rear, Front)

[This algorithm is used to insert the element in the queue where, Q = Queue Array, N = Total number of elements, Item = which is to be inserted, Front, Rear = Deletion and Insertion pointers]

```
Step-1      [Check for overflow]
            If (Rear >= N) then
                Write (" Over Flow")
                Goto step-5
            End if

Step-2      [Increment the Rear Pointer]
            Rear = Rear + 1

Step-3      [Insert the element at Rear position]
            Q [Rear] = item

Step-4      [Reset Pointer]
            If (Front = 0) then
                Set Front = 1
            End if
            Goto step-1

Step-5      [Finished]
            Exit
```

#### ALGORITHM FOR THE DELETION OF ELEMENT FROM THE QUEUE

Algorithm Dequeue (Q, item, Front, Rear)

[ This algorithm is used for the deletion of element from the queue where Q = Array, Item = which is to be deleted]

```
Step-1      [Check for under flow]
            If ((Front = 0) or (Front > Rear)) then
                Write ("Under Flow")
                Goto step-4
            End if
```

Step-2 [Delete the item from queue]

Item= Q [Front]

Step-3 [Set Front Pointer]

If (Front = Rear) then

Front = Rear=0

Else

Front = Front + 1

End if

Goto step-1

Step-4 [Finished]

Exit

### Circular Queue

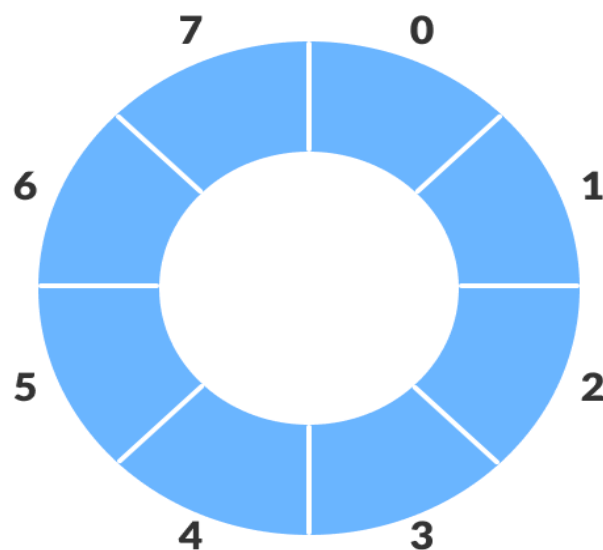
Circular queue provides flexible way for the insertion of element in a circular way.

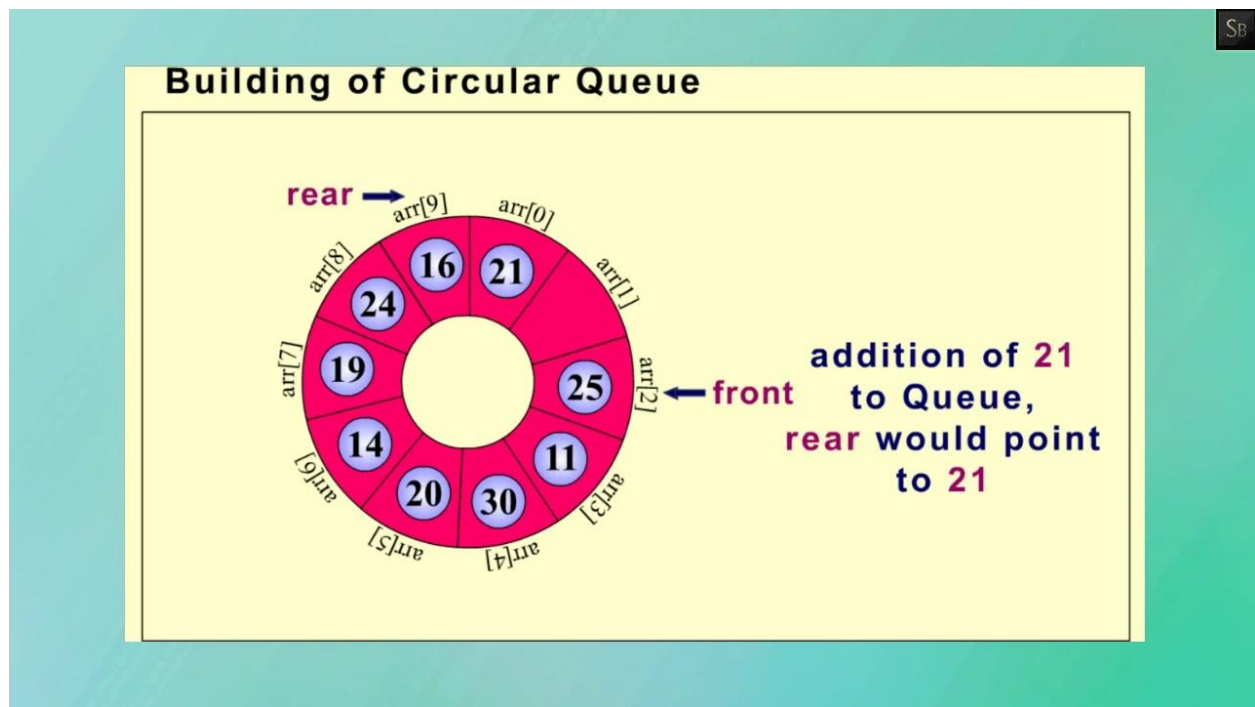
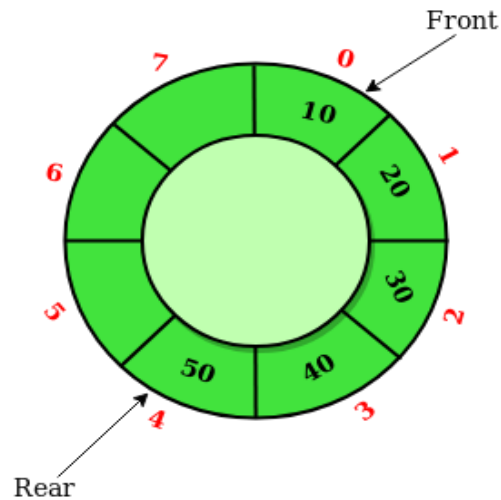
In queue we are not sure that when the rear pointer goes to an end i.e  $\text{Rear} = N$ , overflow condition occurs, maybe the same element from the front are deleted, it is the major drawback of queue.

This drawback can be overcome with the help of circular queue.

In circular queue the insertion can be taken place by using this sequence

1,2,3, ..... N





#### Check for Underflow

When the front and rear pointer become equal that is

$$\text{Front} = \text{Rear} = 0$$

Underflow condition may arise.

#### Check for Overflow

When  $\text{Front} = 1$  and  $\text{Rear} = N$  then overflow condition may arise. OR

When  $\text{Front} \geq \text{Rear} + 1$  then overflow condition may arise.

## ALGORITHM FOR THE INSERTION OF ELEMENT IN CIRCULAR QUEUE

Algorithm CircularInsertion (CQ, N, item, Rear, Front)

[This algorithm is used to insert the element in CQ, where

CQ = Array

N = Final Limit

Item = which is to be inserted

Front, Rear = Deletion and Insertion pointers ]

Step-1      [Check for overflow]

              If ((Front = 1) && (Rear = N) or (Front = Rear + 1) then

                  Write ("Over Flow")

                  Goto step-4

              End if

Step-2      [Reset Rear Pointer]

              If (Front = Rear = 0) then

                  Rear = Front = 1

              Else if (Rear = N) then

                  Rear = 1

              Else

                  Rear = Rear + 1

              End if

Step-3      [Insert the element]

              CQ [Rear] = item

              Goto step-1

Step-4            [Finished]  
Exit

#### ALGORITHM FOR THE DELETION OF ELEMENT FROM CIRCULAR QUEUE

Algorithm DeleteCQ (CQ, item, N, Front, Rear)

[This algorithm is used for the deletion of element from Circular queue where

CQ = Array

Item = which is to be deleted

N = Total number of element

Step-1            [Check for under flow]  
If ((Front = Rear = 0)) then  
    Write ("Under Flow")  
    Goto step-4  
End if

Step-2            [Delete the item from Circular queue]  
Item  $\leftarrow$  CQ [Front]

Step-3            [Set Front Pointer]  
If (Front = Rear) then  
    Front  $\leftarrow$  Rear  $\leftarrow$  0  
Else if (Front = N) then  
    Front = 1  
Else  
    Front  $\leftarrow$  Front + 1  
End if  
Goto step-1

Step-4

[Finished]

Exit