

Notes on Problems :-

Basic Maths:-

Problem 1: Given an integer N, write a program to count the number of digits in N.

Solution:

Approach 1: 1. store the integer in variable x. and initialize a counter variable to count the number of digits.
2. When we divide x by 10 it will result in an integer (given both the variables are integer). For example, 156/10 will result in 15 . similarly 5/10 will result in 0.
3. Using for loop and above observation keep on dividing x by 10. (x is input integer) and increment the count in every iteration when x equals 0 terminate the count and the count will have number of digits in N - $O(N)$

Approach 2:

- 1) Convert the integer N to string.
- 2) return the size of string. $O(1)$

Approach 3:

Use logarithm base 10 to count the number of digits. As

The number of digits in an integer
= the upper bound of $\log_{10}(n)$.

Problem 2: Reverse a given number and print it.

Solution:

Approach 1:

1. convert the number into string.
2. Reverse the string.

Approach 2:

The idea is to extract digits from the end of the given number and create a new number in reverse order.

How to extract digits from the end of a number?

\Rightarrow To extract the last digit, If you divide a number by 10, then the remainder will be last digit.

We can use the modulo(%) operator to do this,

$$\text{Ex. } 123 \% 10 = 3.$$

and to reduce the number by 1 digit, simply divide the number by 10.

To create a number from extracted digits :

⇒ The idea is to start with 0, and for every digit, multiply the number generated so far by 10, and add the digit to it.

Algorithm:

- Run a while loop until the given number N is equal to zero ($N!=0$).
- Initialize a variable reverse = 0;
- now in each step take the remainder of the given number N and store it as a variable digit,
 $digit = N \% 10$;
- Also divide the number by 10, $N=N/10$;
- in each step, variable reverse get updated as
 $reverse = reverse * 10 + digit$.

Problem 3: Given a number, check if it is a pallindrome .

Solution:

1) Reverse the number using algorithm given in Problem 2.

2) check if original & reversed numbers are equal or not.

Problem 4: Find the GCD of two numbers.

Greatest common divisor (GCD) and Highest common Factor (HCF) are same.

Solution:

Approach 1: Simply Traverse From 1 to $\min(a, b)$ and check for every number.

- Traverse from 1 to $\min(a, b)$.
- And check if i can divide by both a and b.
If yes store it in the answer.
- Find the maximum value of i which divides both a and b.

Approach 2:

Using Euclidean's theorem.

Intuition: GCD is the greatest number which can divide both a and b. If a number can divide both a and b. it should also divide both $(a-b)$ and b as well.

$$\gcd(a, b) = \gcd(b, a \% b) \quad a \geq b$$

Approach:

- Recursively call $\gcd(b, a \% b)$ function till the base condition is hit.

- $b=0$ is the base condition. When base condition is hit return a , as $\text{gcd}(a, 0)$ equal to a .

Problem 5: Given a number, check if it is a Armstrong number or not.

Armstrong numbers:

Armstrong numbers are the numbers having the sum of digits (individual digits of a number) raised to power no. of digits (in a number) is equal to a given number.

e.g. $1^3 + 5^3 + 3^3 = 153$

Number of digits $d=3$

$$1^4 + 6^4 + 3^4 + 4^4 = 1634.$$

Number of digits, $d=4$.

Solution:

1. count the number of digits in a number.
2. initialize variable $\text{sum} = 0$.
3. extract each digit, calculate $\text{digit}^{\text{no.of digits}}$. and add it to sum .
4. Repeat step 3 for each digit.
5. check if sum is equal to original number.

Problem 6: Print all divisors of a given number.

Given a number, print all the divisors of the number. A divisor is a number that gives the remainder as zero when divided.

Solution:

Approach 1: Brute Force.

start from 1 and check each number until n. if i divides n add it to list.

Approach 2:

If we take a closer look, we can notice that quotient of a division by one of the divisors is actually another divisor. Like, 4 divides 36.

The quotient is 9, and 9 also divides 36.

Another intuition is that root of a number actually acts as a splitting part of all the divisors of a number.

Also, the quotient of a division by any divisor gives an equivalent divisor on the other side. like 4 gives 9 while dividing 36.

→ From the intuition, we can come to conclusion that we don't need to traverse all the candidates

and if we found a single divisor, we can also find another divisor (Here, we need to be careful, if the given number is a perfect square, like 36, 6 also give 6 as a quotient. This is a corner case.)

- By keeping this in mind, it is enough if we traverse up to the root of a number. Whenever we find the divisor, we print it and also print the quotient. If the quotient is the same, we ignore it. Because the root of a perfect square will give the same quotient as itself.
- The quotients are the numbers that are on the other side of the root. so, it's okay if we stop traversing at the root.

Problem 7: check if a number is prime or not .

Given a number, check whether it is prime or not. A prime number is a natural number that is only divisible by 1 and itself.

Solution:

Brute Force :

check if number is divided by 2.

from 3 to n, check if it is divisible by any odd number.

We can optimise above solution, by only checking till the square root of a given number.
