I21: Introduction à l'algorithmique Cours 4: Algorithmes de tris

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Problème du tri

Problème : Tri

 $\textit{Entrée}: \mathsf{tableau} \ \mathsf{d'entiers} \ T \ \mathsf{de} \ \mathsf{taille} \ n$

 $\it Sortie$: une permutation des éléments de $\it T$ telle que

$$T[1] \le T[2] \le \cdots \le T[n].$$

Problème du tri

- Trier des données est la base de nombreux algorithmes :
 - recherche d'élément
 - plus proche paire
 - unicité
 - enveloppe convexe
- Les algorithmes de tris font appels à des idées réutilisables dans de nombreux autres contextes.
- C'est historiquement le problème le plus étudié en informatique.

Les tris standards

Trois algorithmes classiques :

- tri par sélection
- tri par propagation
- tri par insertion

Caractéristiques des tris standards

Ces tris partagent 2 grandes caractéristiques :

Tris Comparatifs

Ils reposent exlusivement sur l'opération de comparaison des éléments.

Tris in situ

Seul un nombre constant d'éléments est stocké hors du tableau à trié.

Caractéristiques des tris standards

L'opération principale est l'échange de deux valeurs du tableau.

```
algorithme Swap(T,i,j):

DONNEES

T: tableau d entiers
i,j: entiers

VARIABLES:
aux: entier

DEBUT

aux \leftarrow T[i]
T[i] \leftarrow T[j]
T[j] \leftarrow aux

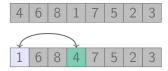
FIN
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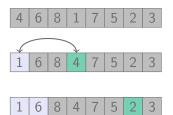
Idée générale :

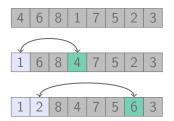
- Chercher le plus petit élément du tableau;
- l'échanger avec l'élément d'indice 1;
- chercher le deuxième plus petit élément;
- l'échanger avec l'élément d'indice 2;
- etc

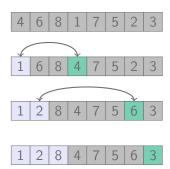
4 | 6 | 8 | 1 | 7 | 5 | 2 | 3

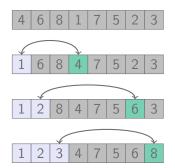


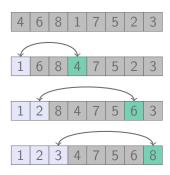




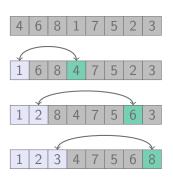


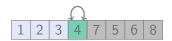


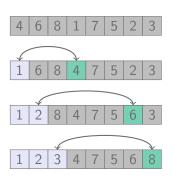




1 2 3 4 7 5 6 8

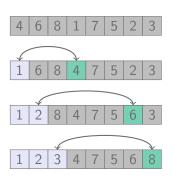


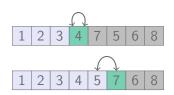


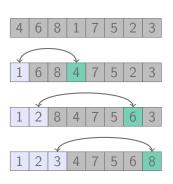


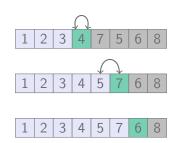


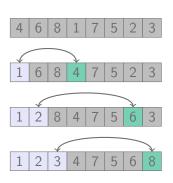


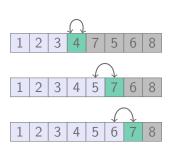


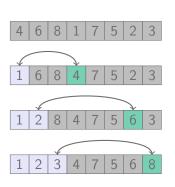


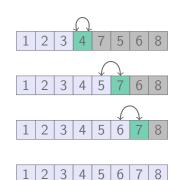












```
ALGORITHME TriSelection(T):
   DONNEES
      T: tableau d'entiers de taille n
   VARIABLES
    imin,i,j: entiers
   DEBUT
     i \leftarrow 1
      TQ i < n-1 FAIRE
        imin←i
9
        i \leftarrow i + 1
10
        TQ j < n FAIRE
11
           SI T[j] < T[imin] ALORS
12
              imin←j
13
           FSI
14
          j \leftarrow j + 1
15
        FTQ
16
         Swap(T, i, imin)
17
        i \leftarrow i + 1
18
      FTQ
19
    FIN
20
```

```
ALGORITHME TriSelection (T):
    DONNEES
       To tableau d'entiers de taille n
    VARIABLES
       imin, i, j: entiers
    DEBUT
      i \leftarrow 1
       TQ i < n-1 FAIRE
          imin ← i
10
         i \leftarrow i + 1
         TQ \mid < n FAIRE
11
             SI T[j] < T[imin] ALORS
12
               imin \leftarrow i
13
             FSI
14
            i \leftarrow i + 1
15
         FTQ
16
          Swap(T, i, imin)
17
          i \leftarrow i + 1
18
       FTQ
19
    FIN
20
```

- Arrêt : deux boucles imbriquées à incrément constant
- Validité : (T[1:i-1]est trié et $T[i-1] \le \min(T[i:n])$) est un invariant de boucle
- Complexité : $\Theta(n^2)$

Idée générale :

- Faire remonter les élément les plus grands en échangeant les éléments contigu mal arrangés;
- effectuer autant de passes que nécessaire pour que tous les éléments soient à la bonne place.

6 4 8 1 7 5 2 3

6 4 8 1 7 5 2 3



4 6 8 1 7 5 2 3

4 6 8 1 7 5 2 3



4 6 1 8 7 5 2 3



4 6 1 7 8 5 2 3



4 6 1 7 5 8 2 3

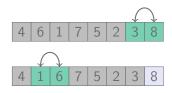


4 6 1 7 5 2 8 3



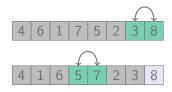




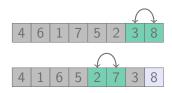




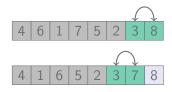


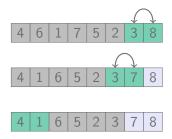


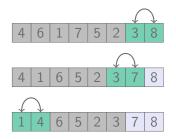


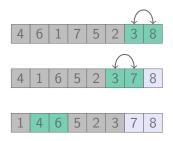


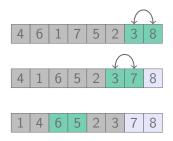


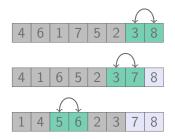


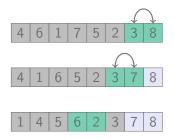


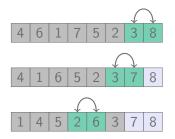


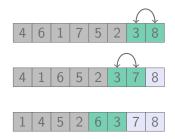


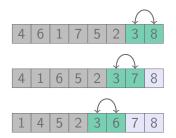


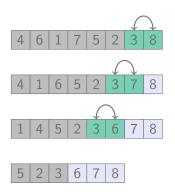


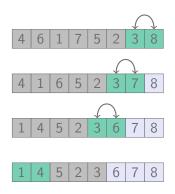


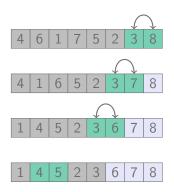


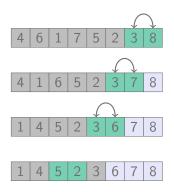


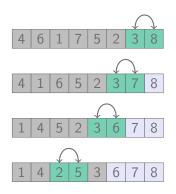


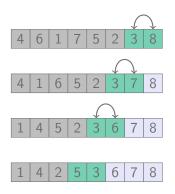


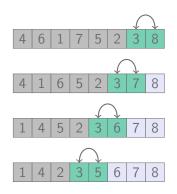




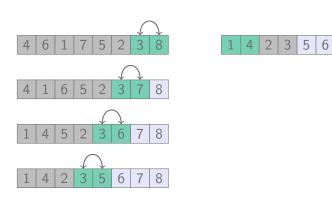




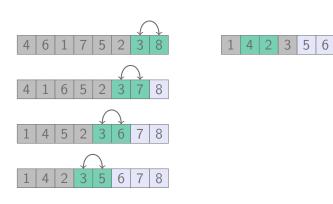


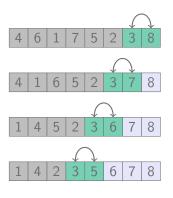


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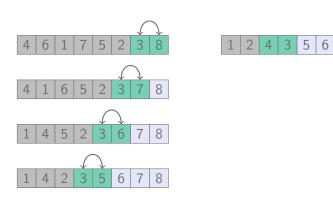
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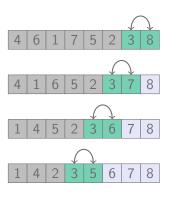




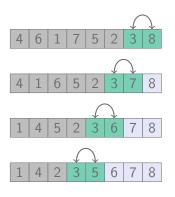


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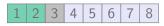


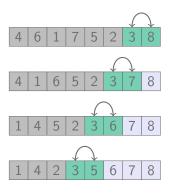


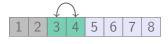


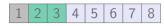


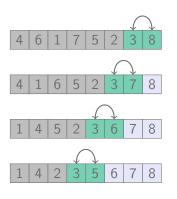






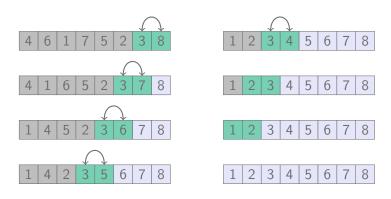












```
ALGORITHME TriBulles (T):
    DONNEES
      T: tableau d'entier de taille n
    VARIABLES:
      d.i: entiers
    DEBUT
      d\leftarrow n
      TQ d > 1 FAIRE
        i \leftarrow 1
         TQ i < d FAIRE
10
            SI T[i] > T[i+1] ALORS
11
               Swap(T, i, i+1)
12
            FSI
13
            i \leftarrow i + 1
14
         FTQ
15
       d \leftarrow d - 1
16
      FTQ
17
    FIN
18
```

- Arrêt : deux boucles imbriquées à incrément constant
- Validité : (T[d+1:n]est trié et t[d+1] >max(T[1:d])) est un invariant
- Complexité : $\Theta(n^2)$

- Figure 1 Si aucun échange n'est effectué alors le tableau est trié.
- C'est un moyen simple pour savoir quand arrêter le travail de l'algorithme.
- On améliore ainsi la complexité dans le meilleur cas.

```
ALGORITHME TriBulles (T):
    DEBUT
       d \leftarrow n
       echange←VRAI
       TQ echnage = VRAI FAIRE
          i \leftarrow 1
6
          echange \leftarrow FAUX
         TQ i < d FAIRE
            SI T[i] > T[i+1] ALORS
9
               Swap(T, i, i+1)
10
               echange←VRAI
11
            FSI
12
            i \leftarrow i + 1
13
         FTQ
14
       d\leftarrow d-1
15
       FTQ
16
    FIN
17
```

```
ALGORITHME TriBulles (T):
    DEBUT
       d \leftarrow n
       echange←VRAI
      TQ echnage = VRAI FAIRE
         i \leftarrow 1
6
         echange←FAUX
         TQ i < d FAIRE
            SI T[i] > T[i+1] ALORS
9
               Swap(T, i, i+1)
               echange←VRAI
11
            FSI
12
            i \leftarrow i + 1
13
         FTQ
14
       d \leftarrow d-1
15
      FTQ
16
    FIN
17
```

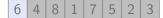
```
Meilleur cas : \Theta(n)
Pire cas : \Theta(n^2)
  Complexité : O(n^2)
```

Idée générale :

- Tri utiliser naturellement pour ranger des cartes;
- on parcours le tableau de droite à gauche et on range chaque élément à sa place parmis les éléments précédents.

5 3

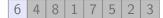
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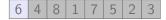
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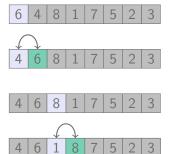


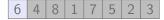








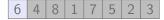


























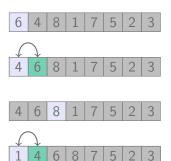
























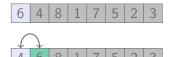






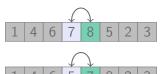


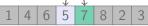










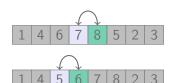






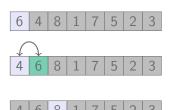


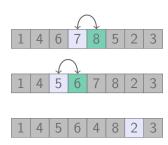


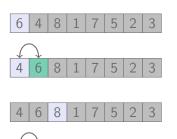


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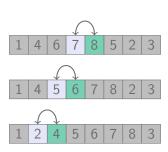


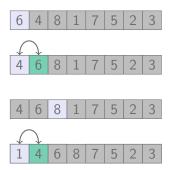


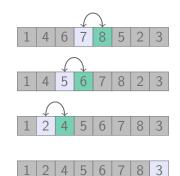


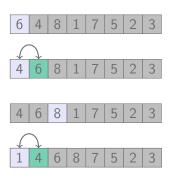
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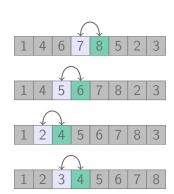
5











```
ALGORITHME TriInsertion(T):
   DONNEES
     T: tableau d'entier de taille n
   VARIABLES:
   i, j: entiers
   DEBUT
    i ←2
     TQ i < n FAIRE
       j \leftarrow i
        TQ j > 1 ET T[j-1] > T[j] FAIRE
10
        Swap(T, j, j-1)
11
         j \leftarrow j - 1
12
     FTQ
13
     i \leftarrow i + 1
14
   FTQ
15
   FIN
16
```

```
ALGORITHME TriInsertion(T):
   DONNEES
      To tableau d'entier de taille n
   VARIABLES:
    i, j: entiers
   DEBUT
     i ← 2
      TQ i < n FAIRE
        i \leftarrow i
        TQ j > 1 ET T[j-1]>T[j] FAIRE
           Swap(T, j, j-1)
11
           j \leftarrow j - 1
12
        FTQ
13
        i \leftarrow i + 1
14
      FTQ
15
    FIN
16
```

- Arrêt : deux boucles imbriquées à incrément constant
- ▶ Validité : (T[1:i-1]est trié) est un invariant
- Complexité : $O(n^2)$