

## Real-World Linkage Convention (using a real stack frame)

A real-world linkage convention puts parameters onto the stack into a stack frame, and allows the subroutine to return a structured type (a struct or class object). There may be variations from this plan (e.g., functions with a variable number of parameters), but this holds all of the key components of a real stack frame calling convention in a compiled language like C. This convention, however, does not really deal with floating-point registers.

Caller prolog (done by the caller):

1. Push any registers \$t0-\$t9 that contain values that must be saved. Push the registers in ascending numerical order.
2. **Stack frame begins here.** If returning a structure (a struct or class object) from the subroutine, allocate space for the structure,  $\$sp = \$sp - \text{space\_for\_structure}$ .
3. Allocate space for parameters,  $\$sp = \$sp - \text{space\_for\_parameters}$ , and put argument values into parameters. Subtraction from  $\$sp$  may be combined with subtraction from step 2. **DO NOT** pass arguments in the \$a registers.
4. Call the subroutine using jal.

Subroutine Prolog (done by the subroutine):

5. Push \$ra (always).
6. Push the caller's frame pointer \$fp (always).
7. Push any registers \$s0-\$s7 that the subroutine might alter. Also push any registers \$a0-\$a3 you might alter. Push the registers in ascending numerical order.
8. Initialize the frame pointer:  $\$fp = \$sp - \text{space\_for\_variables}$ . The "space for variables" is normally four times the number of local (scalar) variables. (Remember that subtracting from  $\$sp$  grows the stack). **Stack frame ends here, with \$fp pointing to bottom of stack frame.** If no local variables, set  $\$fp = \$sp$  and skip the next step.
9. If not the same already, initialize the stack pointer:  $\$sp = \$fp$ .

Subroutine Body:

10. The subroutine may alter any \$t register, or any \$s or \$a register that it saved in the subroutine prolog.
11. The subroutine refers to structure return area, parameters and local variables using `disp($fp)`.
12. The subroutine may push and pop temporary variables and other values on the stack using  $\$sp$ .
13. If the subroutine calls another subroutine, then it does so by following these rules.

Subroutine Epilog (done at the end of the subroutine):

14. If returning a structure, copy it into the structure return area, otherwise, put return value in \$v0-\$v1 (a floating-point value may be returned in \$f0-\$f1) – DO NOT return a scalar on the stack.
15.  $\$sp = \$fp + \text{space\_for\_variables}$ .
16. Pop into \$a and \$s registers any values were previously saved in the stack frame, in reverse order.
17. Pop the caller's frame pointer into \$fp (always).
18. Pop \$ra (always).
19. Return to the caller using jr \$ra.

Regaining Control from a Subroutine (caller epilog):

20. Deallocate space for parameters,  $\$sp = \$sp + \text{space\_for\_parameters}$ .
21. If returning a structure, save to destination and then reclaim space from stack,  $\$sp = \$sp + \text{space\_for\_structure}$ . Or, returned structure may be used as a temporary variable, and de-allocated after use.
22. Pop any registers \$t0-\$t9 that the caller previously pushed, in reverse order.

The stack frame during the subroutine call looks like this:

saved \$t registers		\ (the \$t area is not part of the stack frame)
struct return area	\	<b>Caller-saved or allocated</b>
parameters	\	<b>/ (Caller-side)</b>
\$ra	<b>stack frame</b>	\ (Callee or subroutine side)
\$fp	/	\
saved \$s, \$a registers	/	<b>Callee-saved or allocated</b>
local variables	/	<b>/ &lt;- \$fp points to the bottom (top) here</b>