

# Preliminaries

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## Mathematical Preliminaries

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### Sets

集合根据元素数量可分为 finite set 和 infinite set

Universal set: all possible elements

集合的操作:  $\cup$  (Union),  $\cap$  (Intersection),  $-$  (Difference),  $\overline{A}$  (Complement)

De Morgan's Law:

$$\begin{aligned}\overline{A \cup B} &= \overline{A} \cap \overline{B} \\ \overline{A \cap B} &= \overline{A} \cup \overline{B}\end{aligned}$$

Null set:  $\emptyset$

Subset:  $A \subseteq B$ , 如果  $A \neq B$  即为 proper subset, 记为  $A \subset B$

Disjoint Sets:  $A \cap B = \emptyset$

集合中元素的数量称为集合的势 (Cardinality), 记为  $|A|$

**Power sets:** a set of sets, 是所有子集的集合, 记为  $P(S)$ ,  $2^S$ , 有  $|2^S| = 2^{|S|}$

Cartesian Product: 笛卡尔积,  $|A \times B| = |A| \cdot |B|$

### Functions

给定集合  $A, B$ , 函数  $f$ , 函数将每个  $A$  中的元素映射到至多一个  $B$  中的元素, 记为

$$f: A \rightarrow B$$

total function:  $A = \text{domain}$

injective function:  $\forall a, a' \in A, a \neq a' \rightarrow f(a) \neq f(a')$

surjective function:  $\forall b \in B, \exists a \in A, f(a) = b$

**bijective function:** total, injective and surjective

Big O notation: 参见 [Asymptotic](#)

## Relations

Given two sets,  $A, B$ , a **relation**  $R$  is any subset of  $A \times B$

Equivalence Relations: Reflexive, Symmetric and Transitive

对于一个等价关系  $R$ , 可定义 Equivalence Class

$$[x]_R = \{y : xRy\}$$

两个等价类之间的关系只有相等或 disjoint

对于  $A$  上的关系  $R$ , 有

- Partial order: Reflexive, Transitive and **Antisymmetric**
- Total order: partial order and  $\forall a, b \in A$ , either  $aRb$  or  $bRa$ , also called **linear order**

## Graphs

a directive graph  $G = \langle V, E \rangle$

- Walk: a sequence of adjacent edges
- Path: a walk where no edge is repeated
- Simple path: a path where no node is repeated
- Cycle: a walk from a node to itself
- Simple cycle: only the base node is repeated

A **tree** is a directed graph that has no **cycle**.

## Proof Techniques

归纳法 or 归谬法

鸽笼原理

## String and Languages

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### Alphabet and string

Alphabet: 符号的有限集合, 记为  $\Sigma$

String: alphabet 中的符号组成的序列 (有穷? )

空串即没有符号的 string, 记为  $\lambda$  或是  $\epsilon$

$\Sigma^*$ : the set of all possible strings from alphabet  $\Sigma$  (including  $\lambda$ )

$\Sigma^+$ : the set of all possible strings from alphabet  $\Sigma$  except  $\lambda$ ,  $\Sigma^+ = \Sigma^* - \{\lambda\}$

string 有几种运算, 对于 string  $w = a_1 a_2 \dots a_n, v = b_1 b_2 \dots b_m$

- concatenation:  $wv = a_1 a_2 \dots a_n b_1 b_2 \dots b_m$
- reverse:  $w^R = a_n \dots a_2 a_1$
- length: The length of a string  $x$  is the number of symbols contained in the string  $x$ , denoted by  $|x|$ .  $|w| = n, |wv| = |w| + |v|, |\lambda| = 0$
- substring:  $s$  is a substring of  $x$  if there exist strings  $y$  and  $z$  such that  $x = ysz$  ( $y, z$  can be empty string)
- prefix: when  $x = sz$  ( $z = \epsilon$ ),  $s$  is called a prefix of  $x$
- suffix: when  $x = ys$  ( $y = \epsilon$ ),  $s$  is called a suffix of  $x$
- $w^n = \underbrace{ww \dots w}_n, w^0 = \lambda$

解决字符串相关的问题时, 一般根据其**长度**进行分情况讨论

## Languages

Language 是 string 的集合, 即  $\Sigma^*$  的子集

■ e.g.  $L = \{a^n b^n : n \geq 0\}$ , 这个 Language 不能通过正则描述

由于 Language 本质是集合, 故集合的操作同样适用于 Language (Union, Intersection, Difference, Complement),  $\bar{L} = \Sigma^* - L$

除此之外还有一些 Language 独有的操作

- reverse:  $L^R = \{w^R : w \in L\}$
- concatenation:  $L_1 L_2 = \{xy : x \in L_1, y \in L_2\}$
- $L^n = \underbrace{LL \dots L}_n, L^0 = \{\lambda\}$
- Kleene Closure:  $L^* = \bigcup_{i=0}^{\infty} L^i$
- Positive Closure:  $L^+ = \bigcup_{i=1}^{\infty} L^i = L^* - \{\lambda\}$