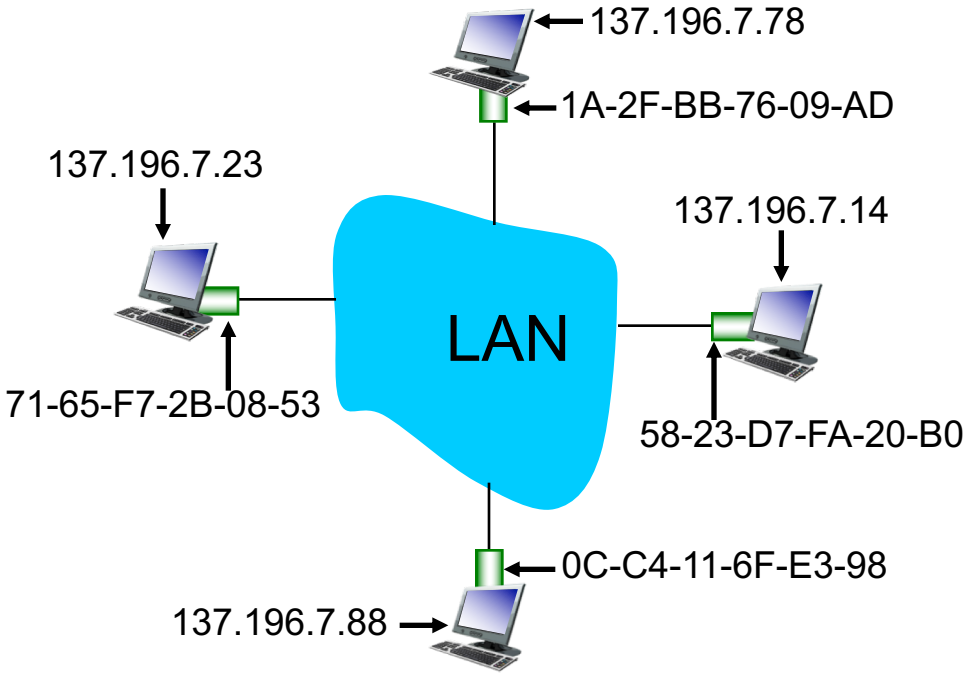
50.020 Network Security
Lab 0: ARP sniffing and spoofing

Summary of last lecture

- Definition of Internet
- Circuit switching vs. packet switching
- Encapsulation and layered network model
- Security issues in computer networking
- Review ARP protocol

ARP: address resolution protocol

Question: how to determine interface's MAC address, knowing its IP address?



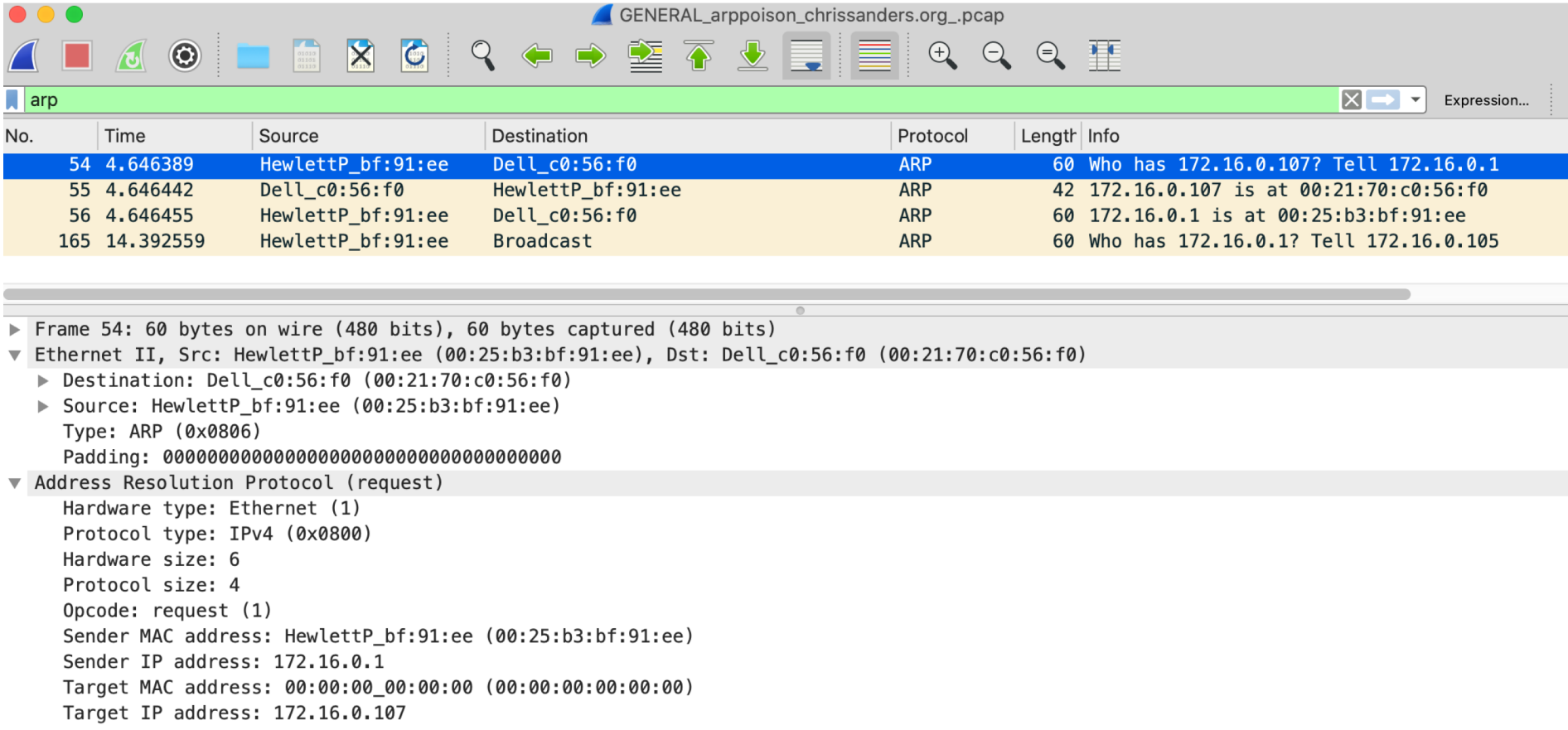
ARP table: each IP node (host, router) on LAN has table

- IP/MAC address mappings for some LAN nodes:
  < IP address; MAC address; TTL>
- TTL (Time To Live): time after which address mapping will be forgotten (typically 20 min)
- Try : arp –a

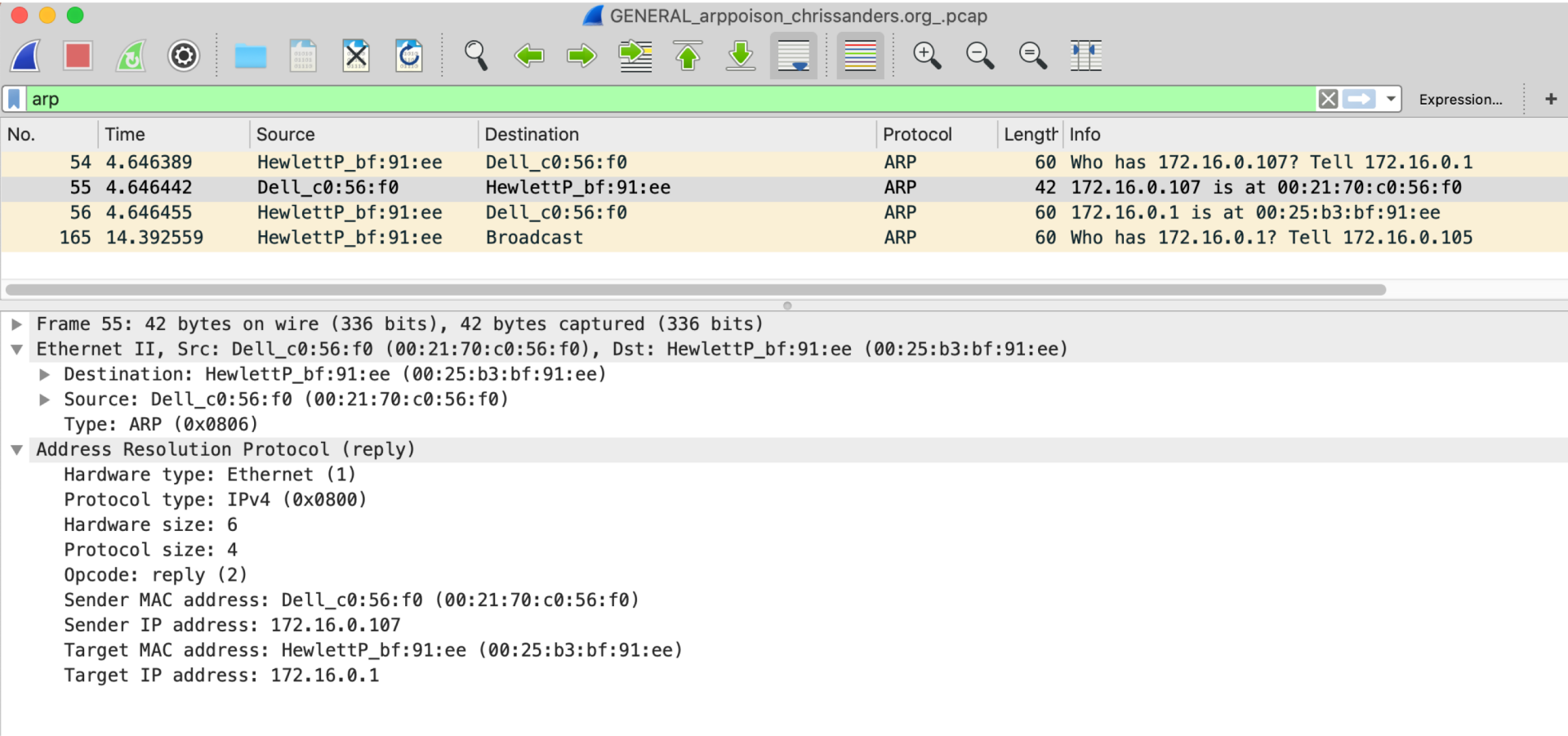
ARP protocol: same LAN

- A wants to send datagram to B
  - B's MAC address not in A's ARP table.
- A broadcasts ARP query packet, containing B's IP address
  - destination MAC address = FF-FF-FF-FF-FF-FF
  - all nodes on LAN receive ARP query
- B receives ARP packet, replies to A with its (B's) MAC address
  - frame sent to A's MAC address (unicast)
- A caches (saves) IP-to-MAC address pair in its ARP table until information becomes old (times out)
  - soft state: information that times out (goes away) unless refreshed
- ARP is "plug-and-play":
  - nodes create their ARP tables *without intervention from net administrator*

Example of an ARP request



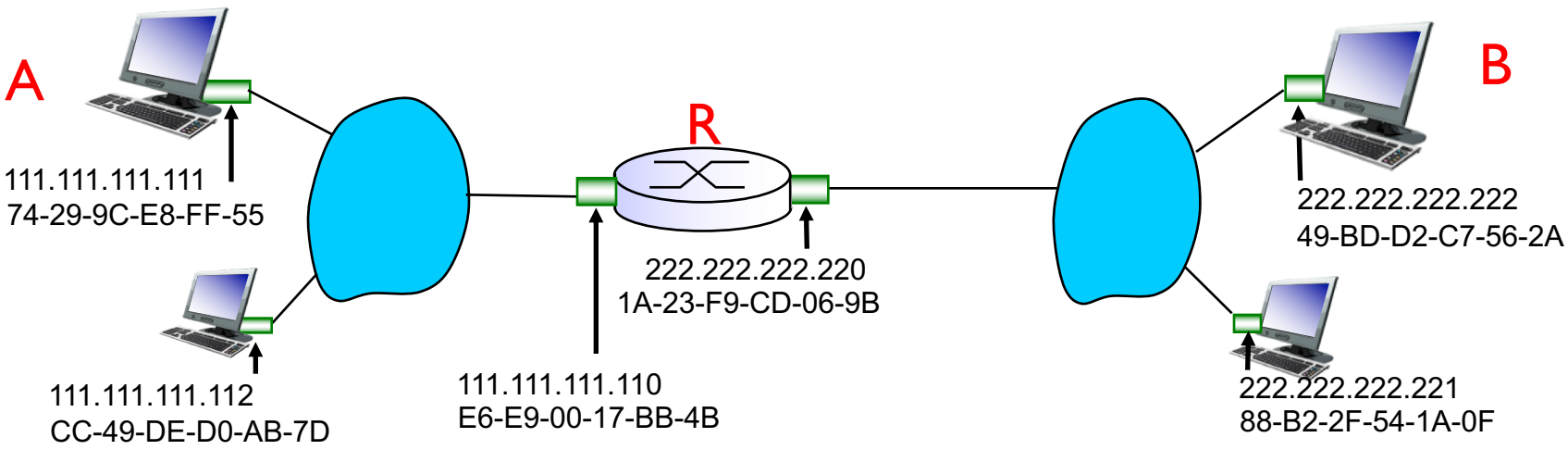
Example of an ARP reply



Addressing: routing to another LAN

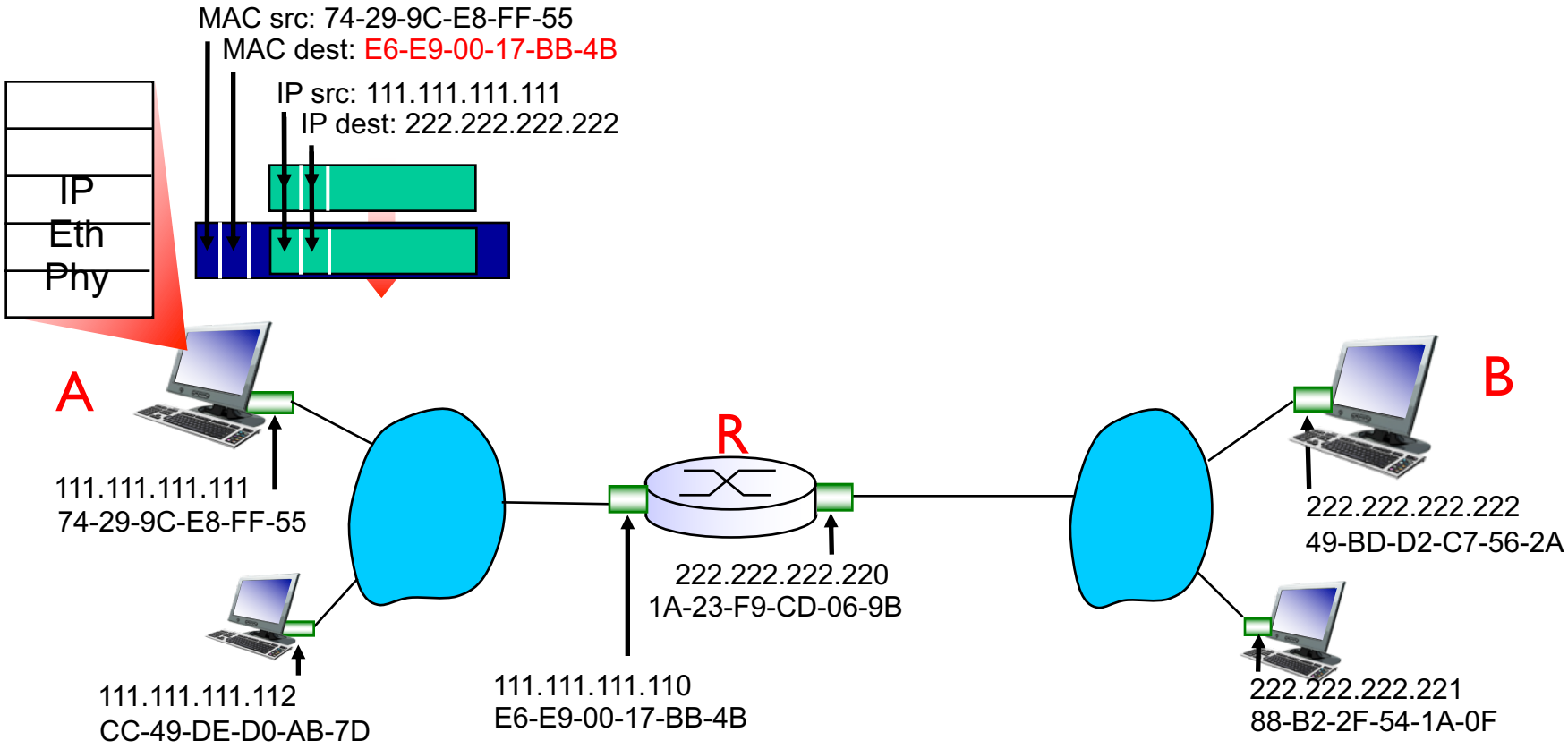
walkthrough: send datagram from A to B via R

- focus on addressing – at IP (datagram) and MAC layer (frame)
- assume A knows B's IP address
- assume A knows IP address of first hop router, R (how?)
- assume A knows R's MAC address (how?)



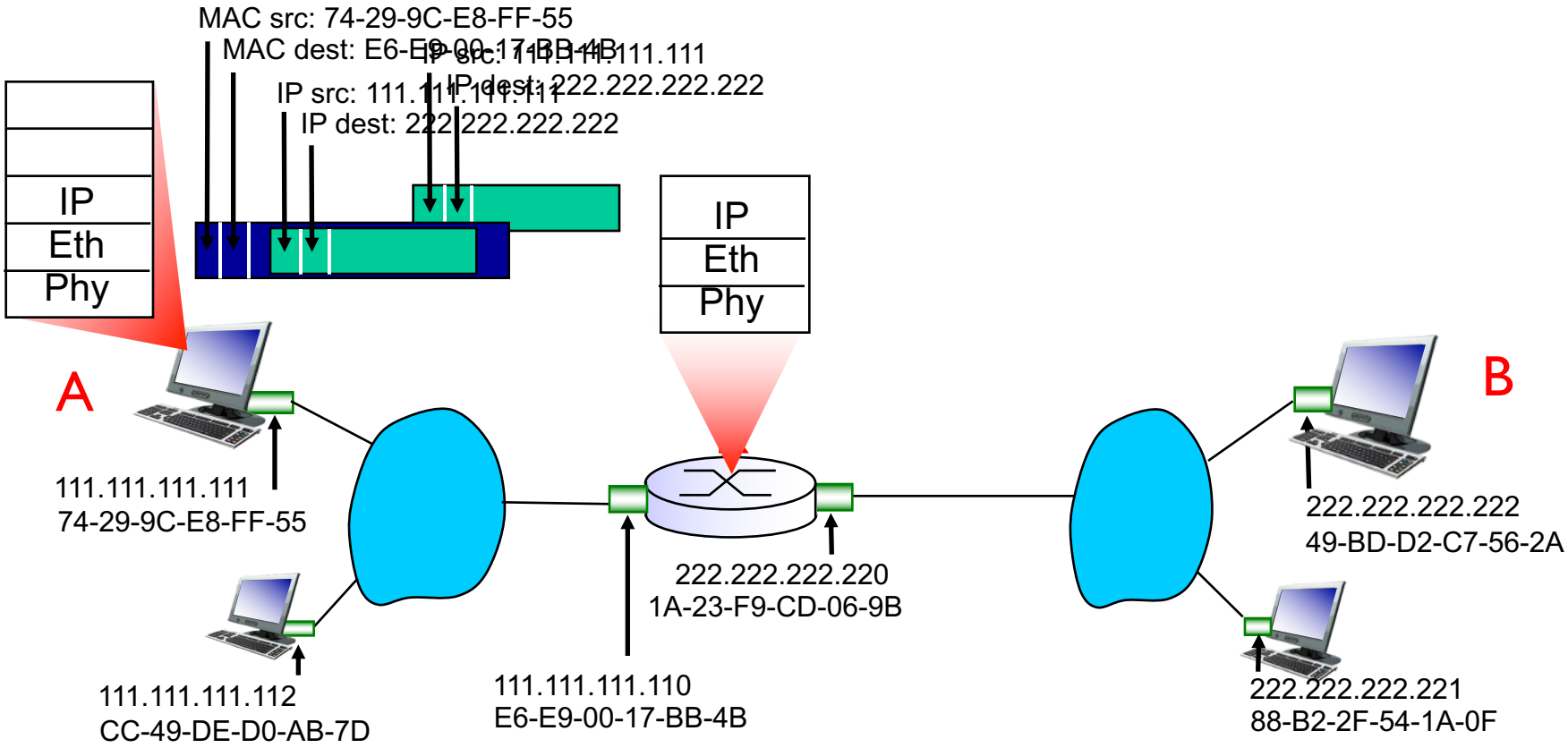
Addressing: routing to another LAN

- A creates IP datagram with IP source A, destination B
- A creates link-layer frame with R's MAC address as destination address, frame contains A-to-B IP datagram



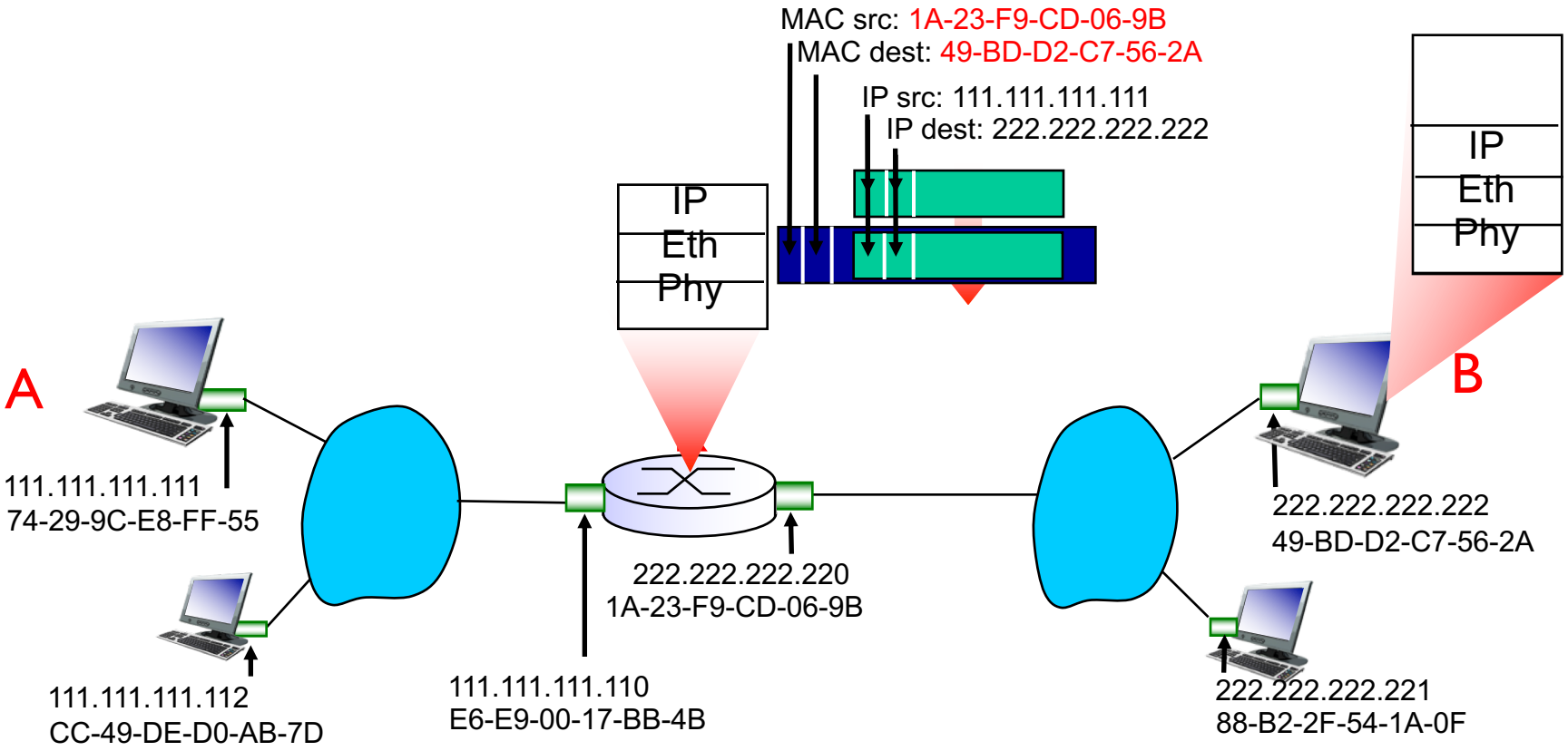
Addressing: routing to another LAN

- frame sent from A to R
- frame received at R, datagram removed, passed up to IP



Addressing: routing to another LAN

- R forwards datagram with IP source A, destination B
- R creates link-layer frame with B's MAC address as destination address, frame contains A-to-B IP datagram



Addressing: routing to another LAN

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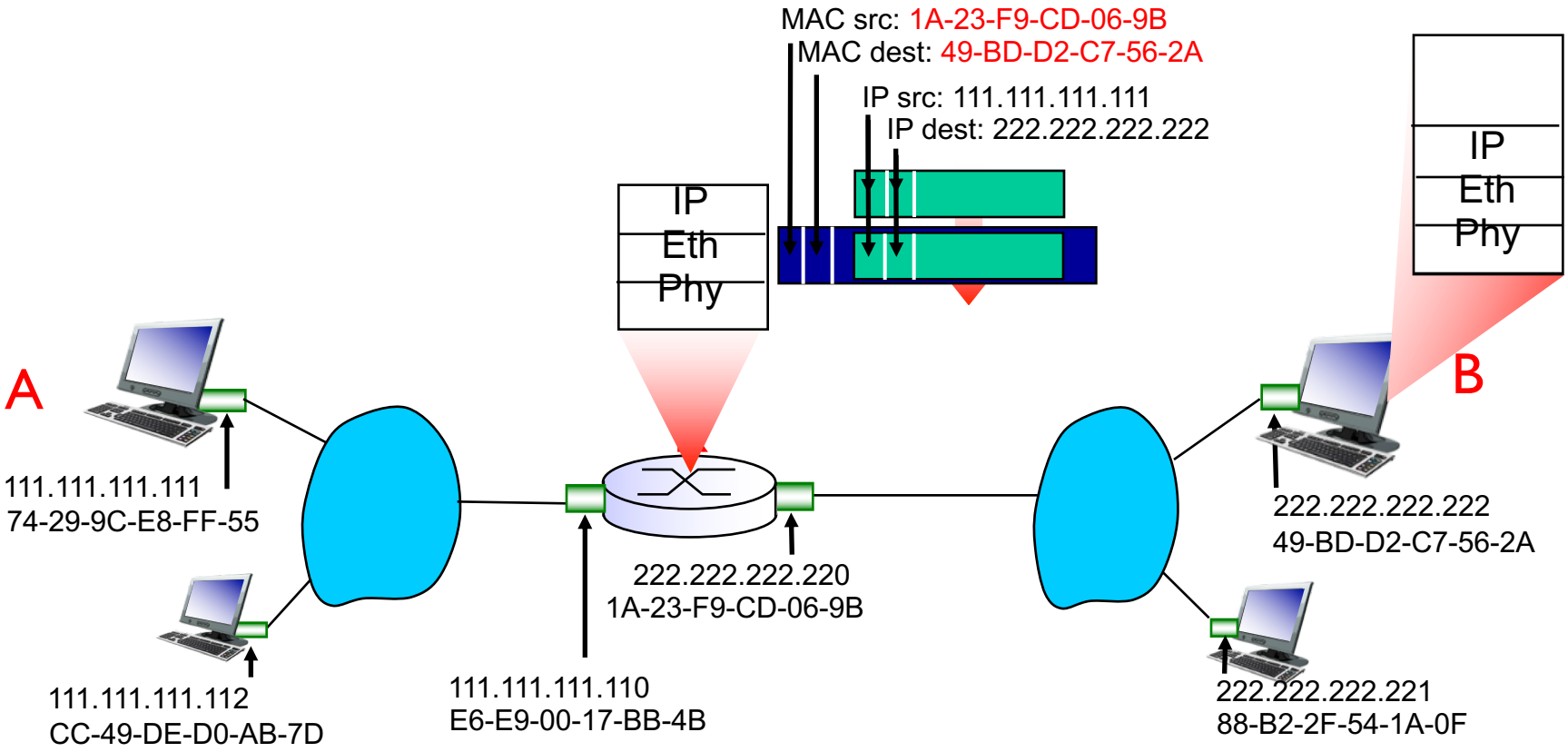
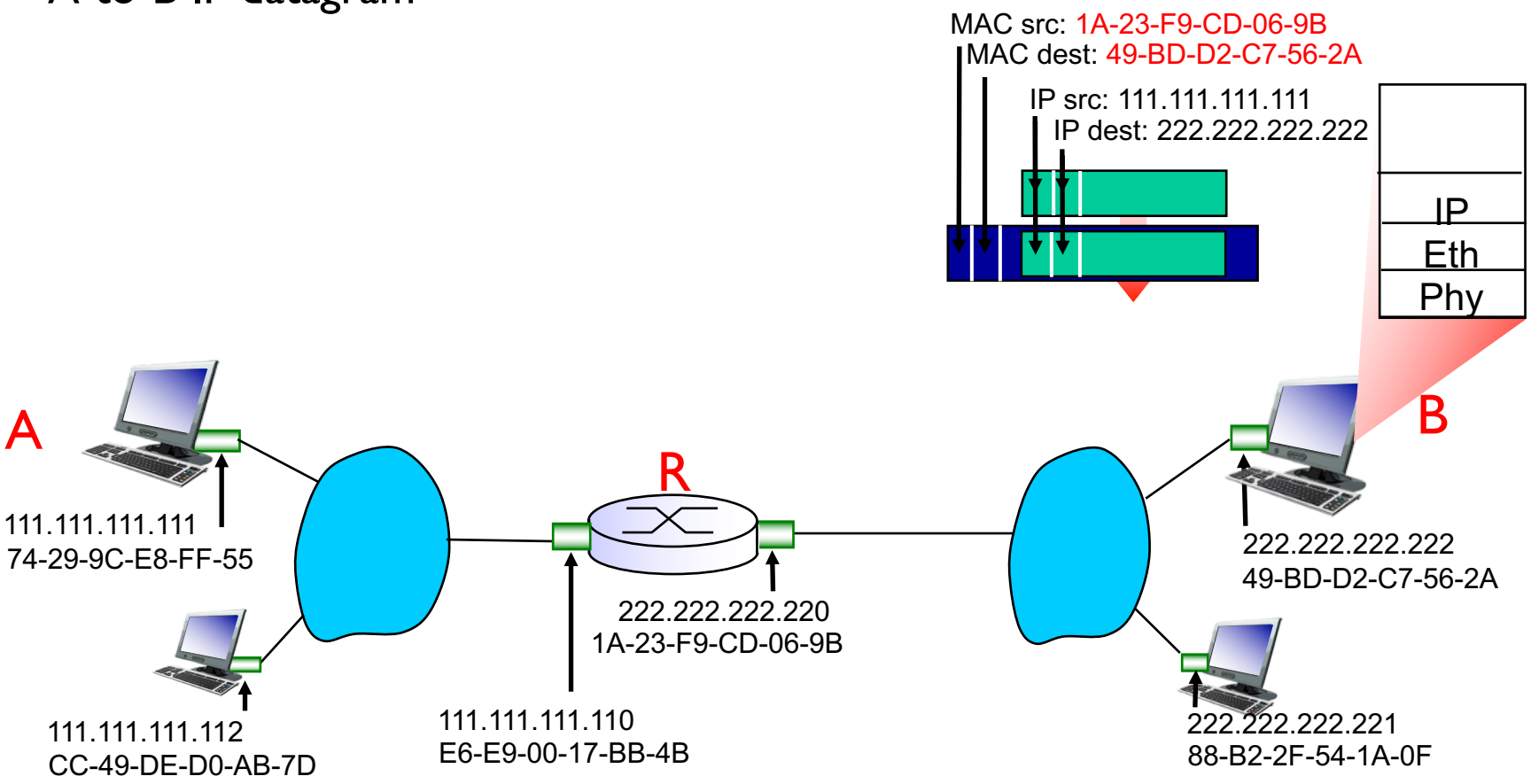
- R creates link-layer frame with B's MAC address as destination address, frame contains A-to-B IP datagram



Addressing: routing to another LAN

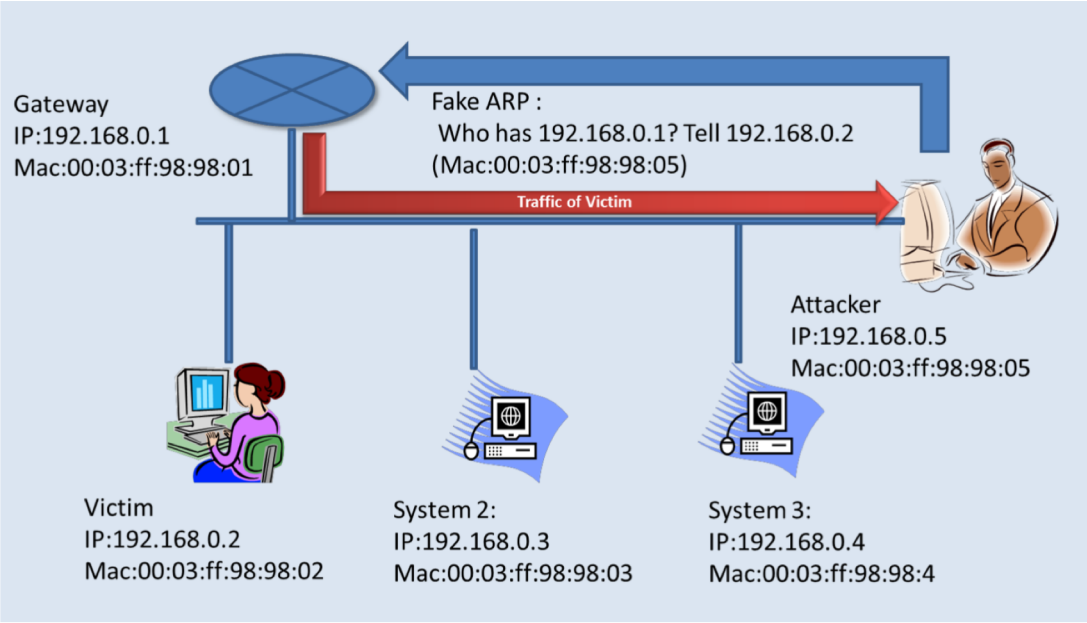
- R forwards datagram with IP source A, destination B
- R creates link-layer frame with B's MAC address as dest, frame contains A-to-B IP datagram



* Check out the online interactive exercises for more examples: http://gaia.cs.umass.edu/kurose_ross/interactive/

ARP Spoofing

- The ARP table is updated whenever an ARP response is received
- Requests are not tracked
- ARP announcements are not authenticated
- Machines trust each other
- A rogue machine can spoof other machines

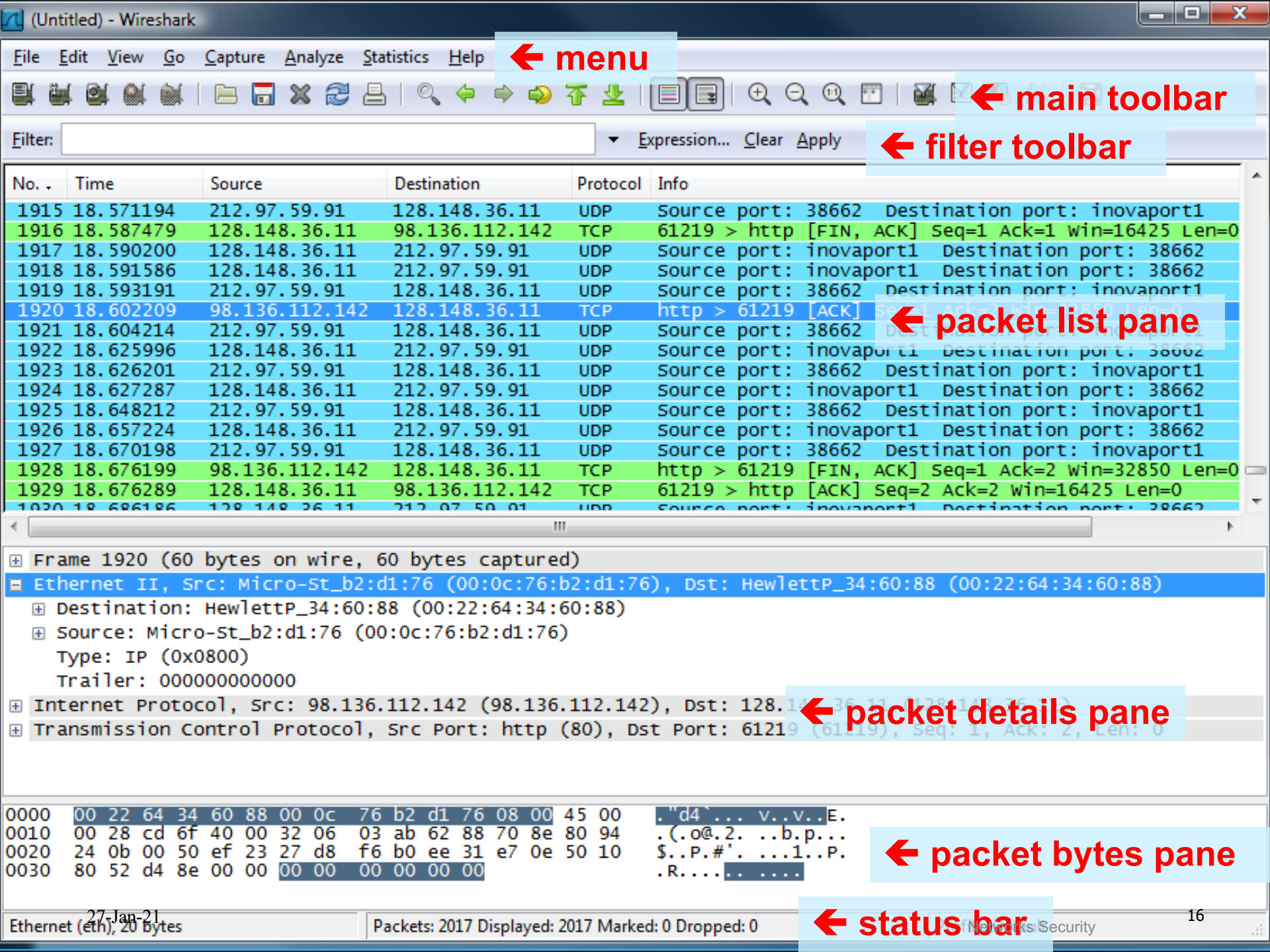


ARP Poisoning (ARP Spoofing)

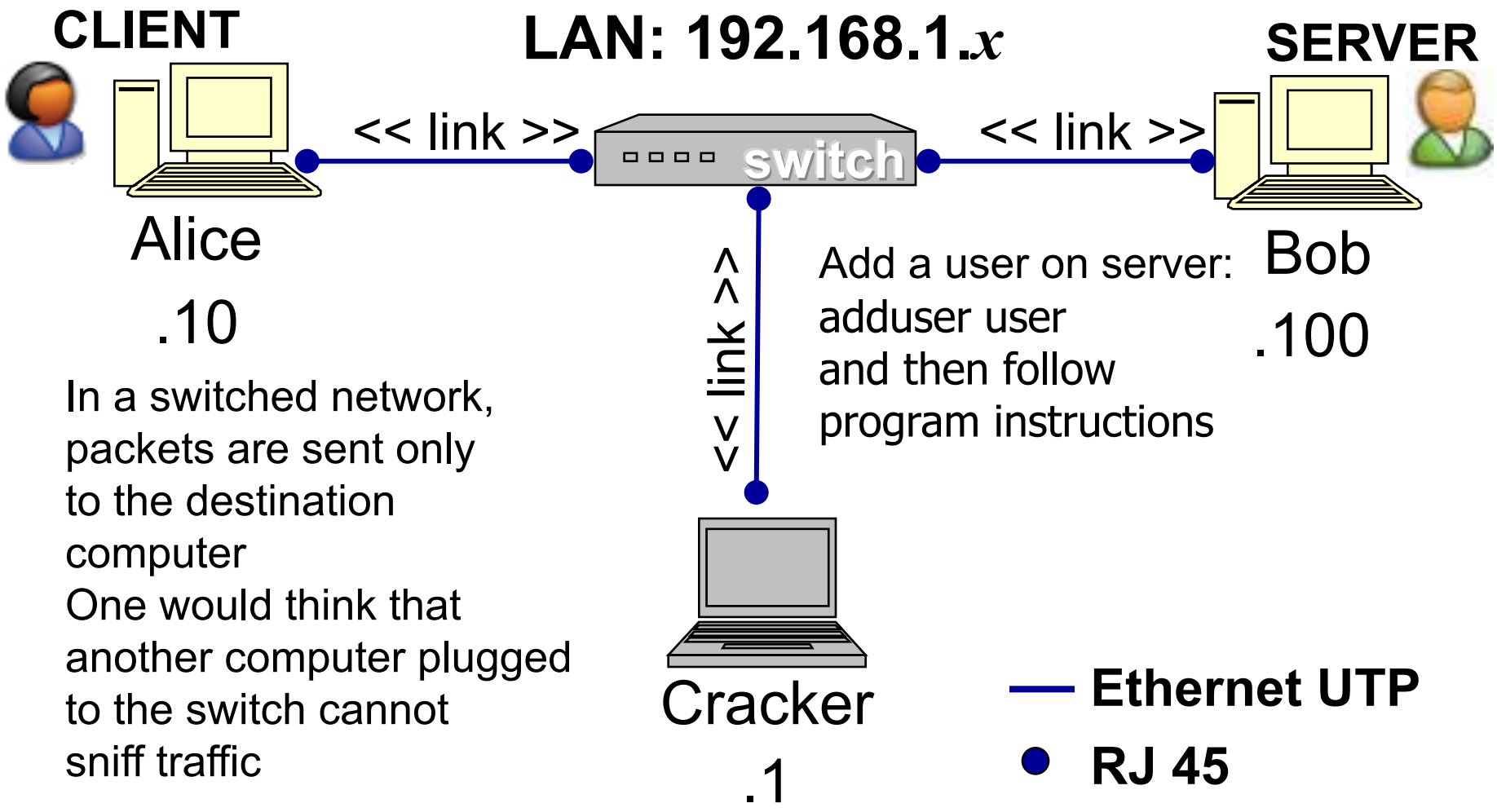
- According to the standard, almost all ARP implementations are stateless
- An arp cache updates every time that it receives an arp reply… even if it did not send any arp request!
- It is possible to "poison" an arp cache by sending gratuitous arp replies
- Using static entries solves the problem but it is almost impossible to manage!

Wireshark

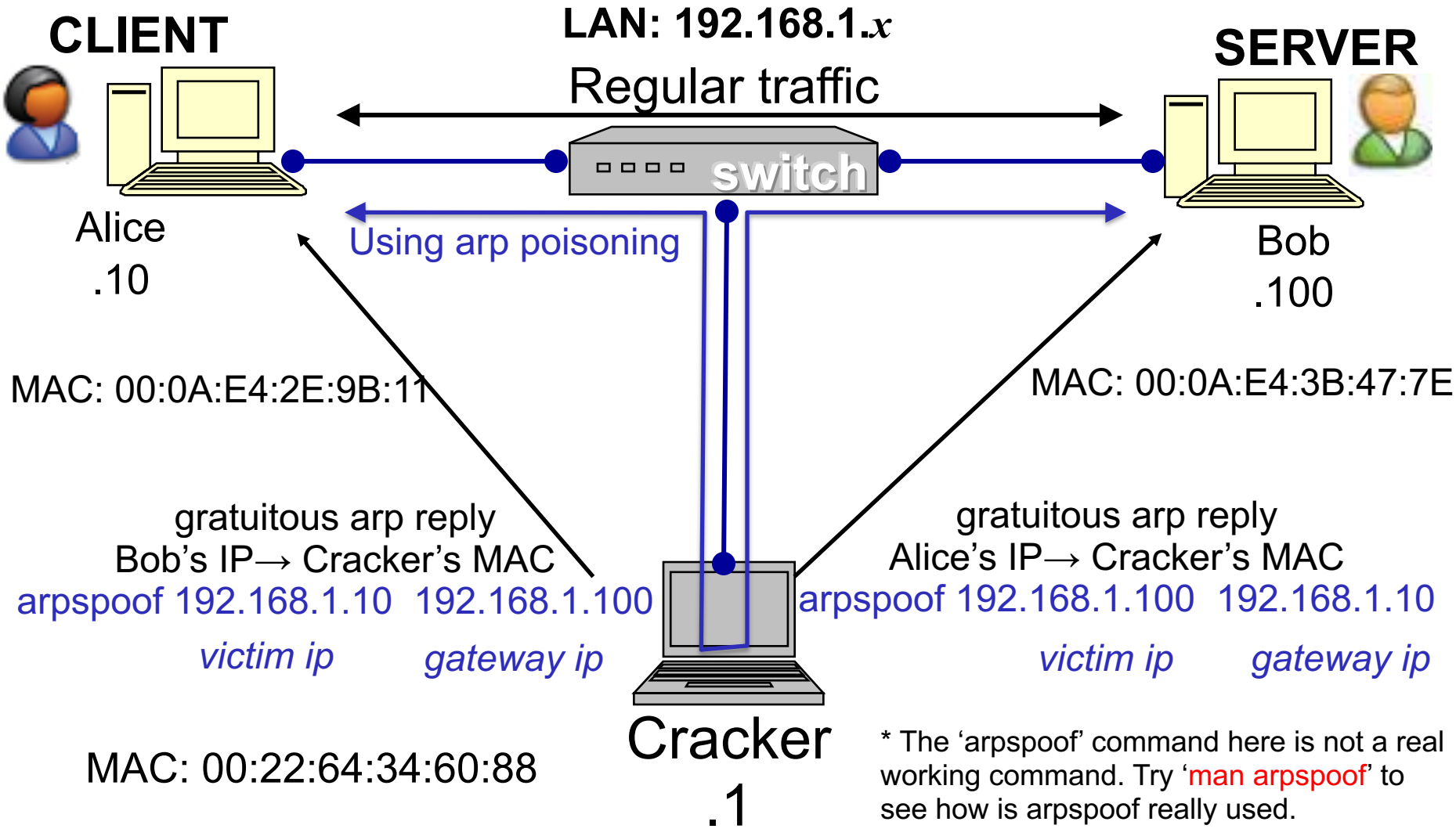
- Wireshark is a packet sniffer and protocol analyzer
  - Captures and analyzes frames
  - Supports plugins
- Usually required to run with administrator privileges
- Setting the network interface in promiscuous mode captures traffic across the entire LAN segment and not just frames addressed to the machine
- Freely available on www.wireshark.org


(Untitled) - Wireshark
File Edit View Go Capture Analyze Statistics Help
menu
main toolbar
Filter:
Expression... Clear Apply
filter toolbar
No. Time Source Destination Protocol Info
1915 18.571194 212.97.59.91 128.148.36.11 UDP Source port: 38662 Destination port: inovaport1
1916 18.587479 128.148.36.11 98.136.112.142 TCP 61219 > http [FIN, ACK] Seq=1 Ack=1 Win=16425 Len=0
1917 18.590200 128.148.36.11 212.97.59.91 UDP Source port: inovaport1 Destination port: 38662
1918 18.591586 128.148.36.11 212.97.59.91 UDP Source port: inovaport1 Destination port: 38662
1919 18.593191 212.97.59.91 128.148.36.11 UDP Source port: 38662
1920 18.602209 98.136.112.142 128.148.36.11 TCP http > 61219 [ACK]
1921 18.604214 212.97.59.91 128.148.36.11 UDP Source port: 38662
1922 18.625996 128.148.36.11 212.97.59.91 UDP Source port: inovaport1
1923 18.626201 212.97.59.91 128.148.36.11 UDP Source port: 38662 Destination port: inovaport1
1924 18.627287 128.148.36.11 212.97.59.91 UDP Source port: inovaport1 Destination port: 38662
1925 18.648212 212.97.59.91 128.148.36.11 UDP Source port: 38662 Destination port: inovaport1
1926 18.657224 128.148.36.11 212.97.59.91 UDP Source port: inovaport1 Destination port: 38662
1927 18.670198 212.97.59.91 128.148.36.11 UDP Source port: 38662 Destination port: inovaport1
1928 18.676199 98.136.112.142 128.148.36.11 TCP http > 61219 [FIN, ACK] Seq=1 Ack=2 Win=32850 Len=0
1929 18.676289 128.148.36.11 98.136.112.142 TCP 61219 > http [ACK] Seq=2 Ack=2 Win=16425 Len=0
packet list pane
Frame 1920 (60 bytes on wire, 60 bytes captured)
Ethernet II, Src: Micro-St_b2:d1:76 (00:0c:76:b2:d1:76), Dst: HewlettP_34:60:88 (00:22:64:34:60:88)
Destination: HewlettP_34:60:88 (00:22:64:34:60:88)
Source: Micro-St_b2:d1:76 (00:0c:76:b2:d1:76)
Type: IP (0x0800)
Trailer: 000000000000
Internet Protocol, Src: 98.136.112.142 (98.136.112.142), Dst: 128.
Transmission Control Protocol, Src Port: http (80), Dst Port: 6121
packet details pane
0000 00 22 64 34 60 88 00 0c 76 b2 d1 76 08 00 45 00 ."d4`... v..v..E.
0010 00 28 cd 6f 40 00 32 06 03 ab 62 88 70 8e 80 94 .(.o@.2. ..b.p...
0020 24 0b 00 50 ef 23 27 d8 f6 b0 ee 31 e7 0e 50 10 $..P.#'. ...1..P.
0030 80 52 d4 8e 00 00 00 00 00 00 00 00 .R...... ....
packet bytes pane
Ethernet (eth), 20 bytes
Packets: 2017 Displayed: 2017 Marked: 0 Dropped: 0
status bar

DEMO 1: Configuration using Telnet

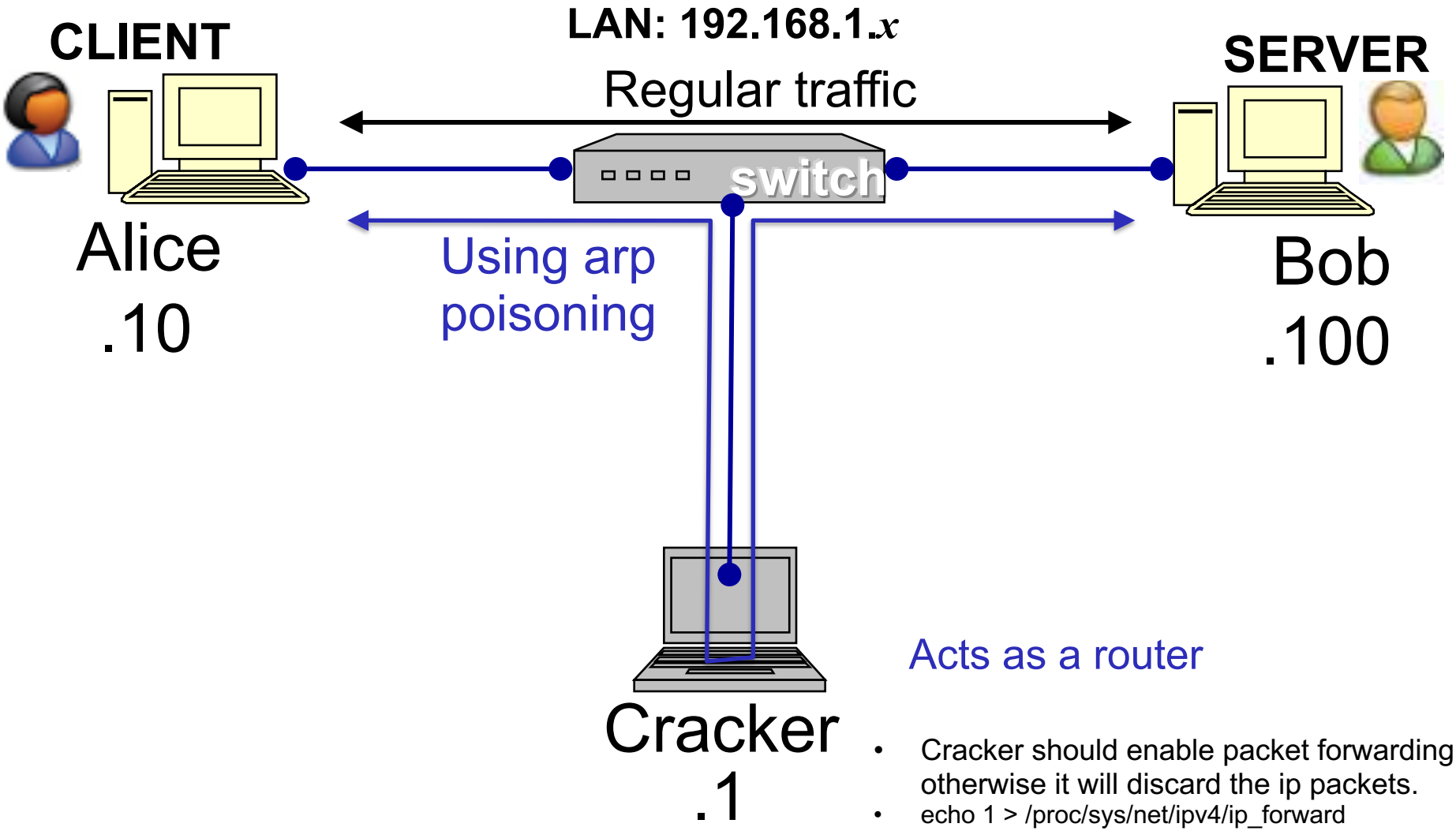


DEMO 1: ARP Spoofing



* The 'arpspoof' command here is not a real working command. Try 'man arpspoof' to see how is arpspoof really used.

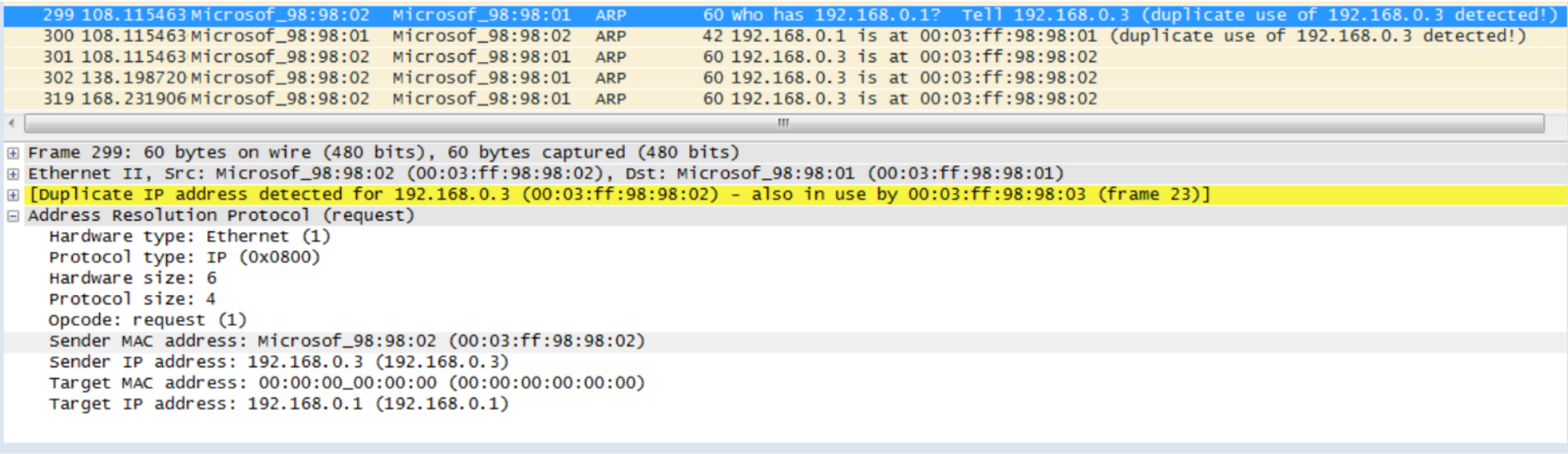
DEMO 1: catch telnet password



- Cracker should enable packet forwarding otherwise it will discard the ip packets.
- echo 1 > /proc/sys/net/ipv4/ip_forward

ARP Spoofing Detection

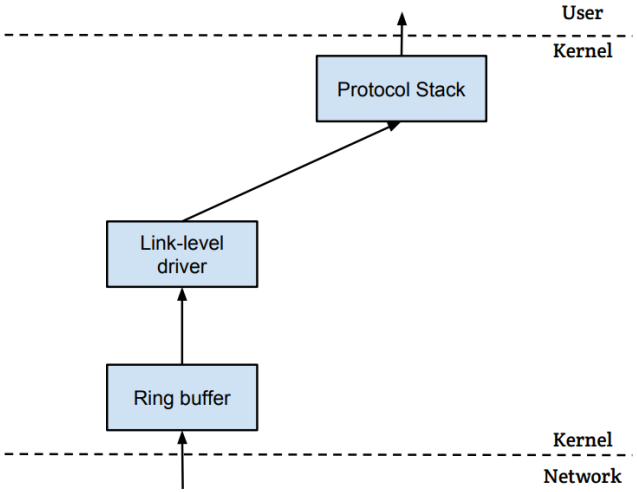
- Detected through duplicated MAC address



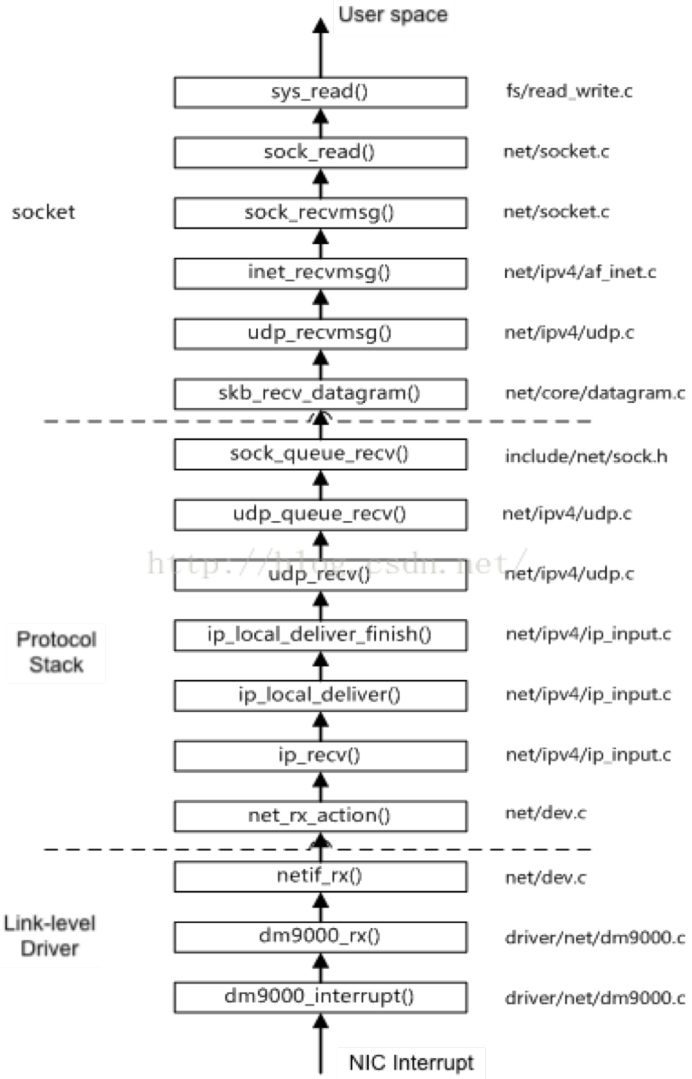
Packet Sniffing and Spoofing

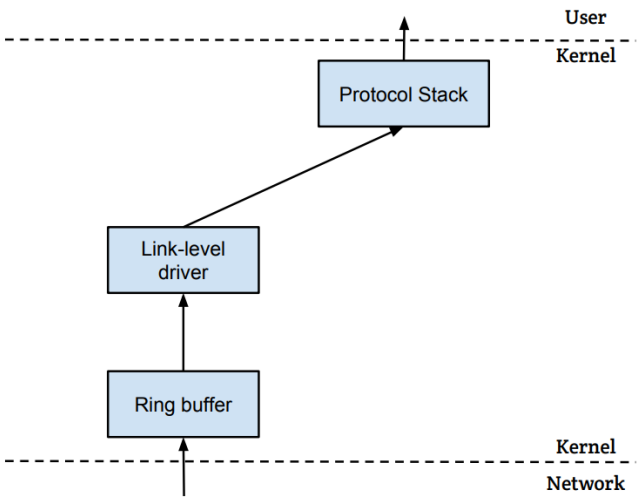
How Packets Are Received

- NIC (Network Interface Card) is a physical or logical link between a machine and a network
- Each NIC has a MAC address
- Every NIC on the network will hear all the frames on the wire
- NIC checks the destination address for every packet, if the address matches the cards MAC address, it is further copied into a buffer in the kernel
  - Copy through direct memory access (DMA)
  - Then NIC triggers an interrupt



What indeed happens inside a Linux kernel





(a) Packet flow

Q: where are ARP and ICMP processed?

Promiscuous Mode

- The frames that are not destined to a given NIC are discarded
  - i.e., if they do not match the destination MAC address
- When operating in promiscuous mode, NIC passes every frame received from the network to the kernel
  - promiscuous mode is a NIC mode!
- If a sniffer program is registered with the kernel, it will be able to see all the packets
  - Q: where are the packets now?
- In Wi-Fi, it is called Monitor Mode

BSD Packet Filter (BPF)

```
struct sock_filter code[] = {
  { 0x28,  0,  0, 0x0000000c }, { 0x15,  0,  8, 0x000086dd },
  { 0x30,  0,  0, 0x00000014 }, { 0x15,  2,  0, 0x00000084 },
  { 0x15,  1,  0, 0x00000006 }, { 0x15,  0, 17, 0x00000011 },
  { 0x28,  0,  0, 0x00000036 }, { 0x15, 14,  0, 0x00000016 },
  { 0x28,  0,  0, 0x00000038 }, { 0x15, 12, 13, 0x00000016 },
  { 0x15,  0, 12, 0x00000800 }, { 0x30,  0,  0, 0x00000017 },
  { 0x15,  2,  0, 0x00000084 }, { 0x15,  1,  0, 0x00000006 },
  { 0x15,  0,  8, 0x00000011 }, { 0x28,  0,  0, 0x00000014 },
  { 0x45,  6,  0, 0x00001fff }, { 0xb1,  0,  0, 0x0000000e },
  { 0x48,  0,  0, 0x0000000e }, { 0x15,  2,  0, 0x00000016 },
  { 0x48,  0,  0, 0x00000010 }, { 0x15,  0,  1, 0x00000016 },
  { 0x06,  0,  0, 0x0000ffff }, { 0x06,  0,  0, 0x00000000 },
};

struct sock_fprog bpf = {
   .len = ARRAY_SIZE(code),
   .filter = code,
};
```

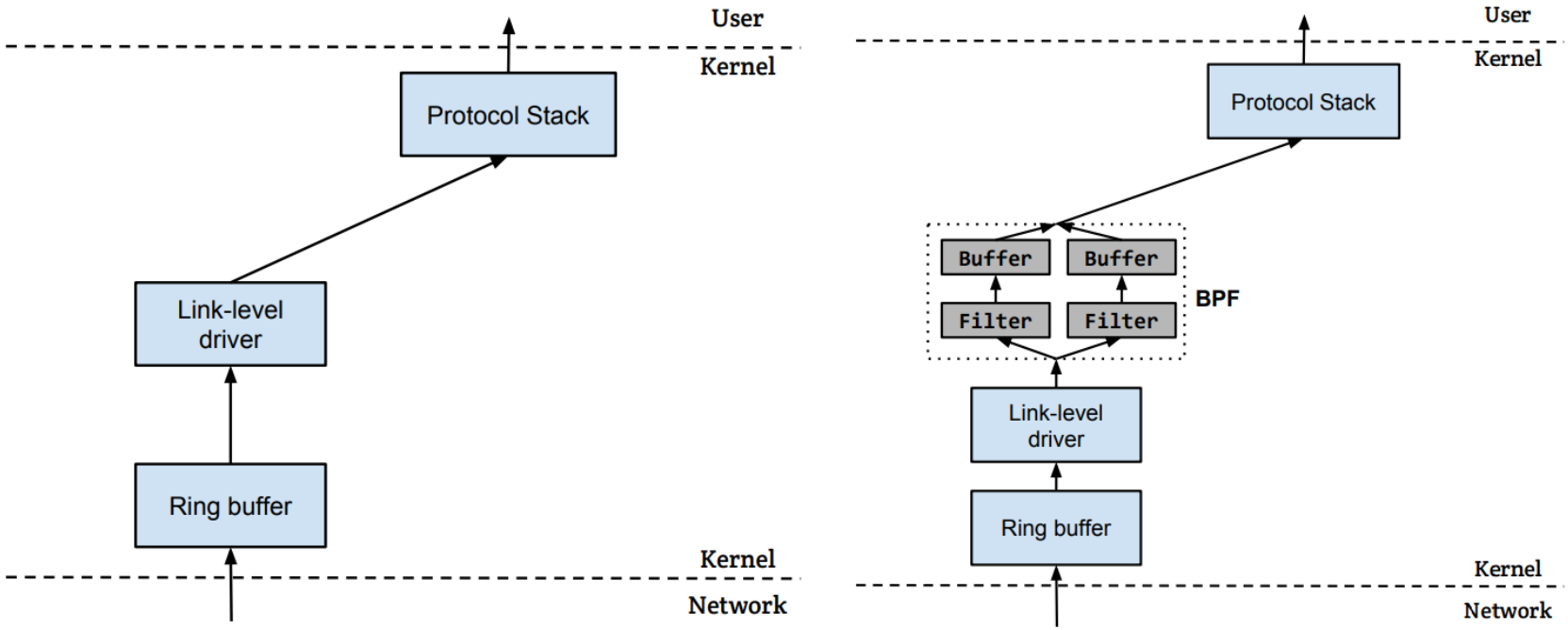
- BPF allows a user-program to attach a filter to the socket, which tells the kernel to discard unwanted packets.
- An example of the compiled BPF code is shown here.

BSD Packet Filter (BPF)

```
setsockopt(sock, SOL_SOCKET, SO_ATTACH_FILTER, &bpf, sizeof(bpf))
```

- A compiled BPF pseudo-code can be attached to a socket through `setsockopt()`
- When a packet is received by kernel, BPF will be invoked
- An accepted packet is pushed up the protocol stack. See the diagram on the following slide.

Packet Flow With/Without Filters



Packet Sniffing

Packet sniffing describes the process of capturing live data as they flow across a network

Let's first see how computers receive packets.

Receiving Packets Using Socket

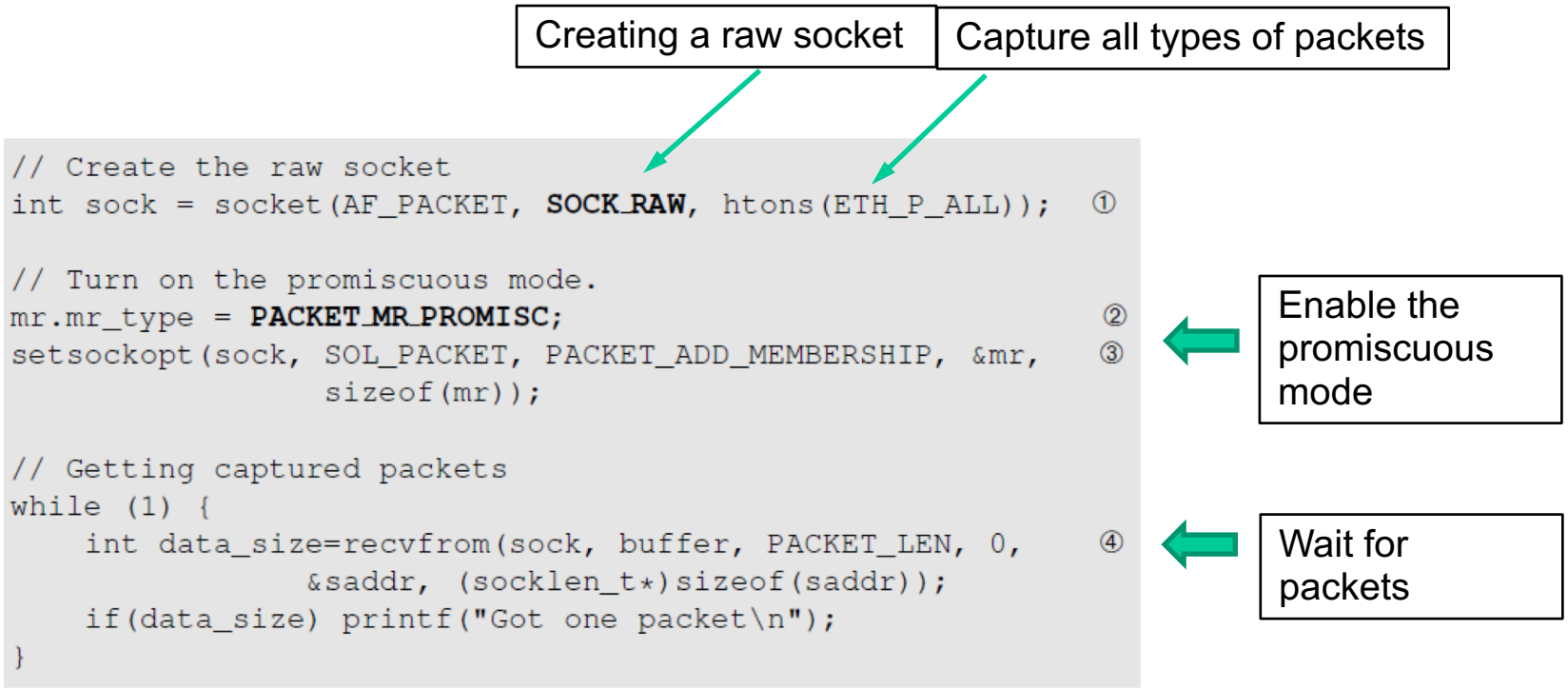
Create the socket

Provide information about server

Receive packets

```
// Step ①
int sock = socket(AF_INET, SOCK_DGRAM, IPPROTO_UDP);

// Step ②
memset((char *) &server, 0, sizeof(server));
server.sin_family = AF_INET;
server.sin_addr.s_addr = htonl(INADDR_ANY);
server.sin_port = htons(9090);

if (bind(sock, (struct sockaddr *) &server, sizeof(server)) < 0)
    error("ERROR on binding");

// Step ③
while (1) {
    bzero(buf, 1500);
    recvfrom(sock, buf, 1500-1, 0,
             (struct sockaddr *) &client, &clientlen);
    printf("%s\n", buf);
}
```

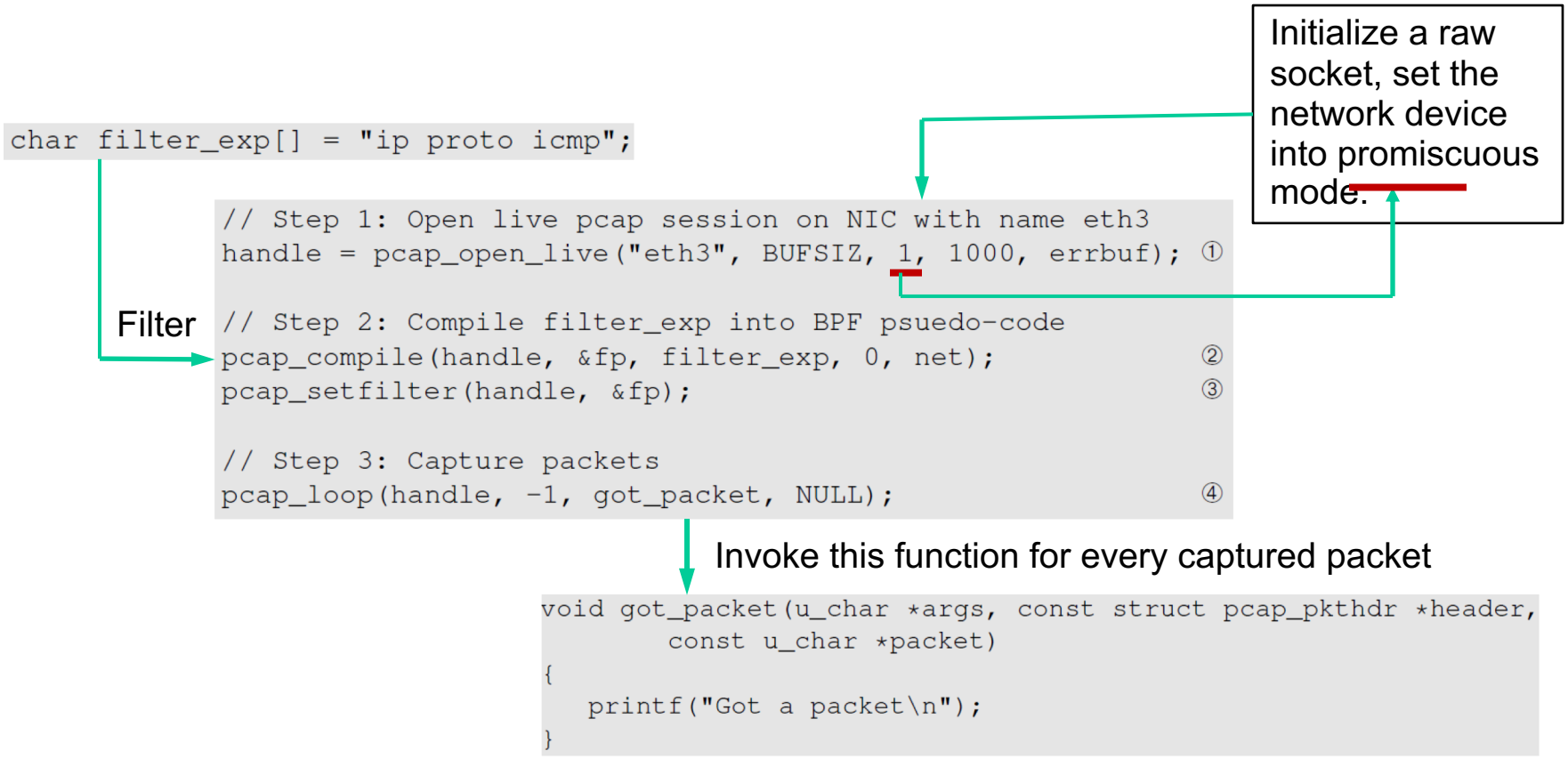
Receiving Packets Using Raw Socket



Limitation of the Approach

- This program is not portable across different operating systems.
- Setting filters is not easy.
- The program does not explore any optimization to improve performance.
- The PCAP library was thus created.
  - It still uses raw sockets internally, but its API is standard across all platforms. OS specifics are hidden by PCAP's implementation.
  - Allows programmers to specify filtering rules using human readable Boolean expressions.

Packet Sniffing Using the pcap API



Processing Captured Packet: Ethernet Header

```
/* Ethernet header */
struct ethheader {
  u_char  ether_dhost[ETHER_ADDR_LEN]; /* destination host address */
  u_char  ether_shost[ETHER_ADDR_LEN]; /* source host address */
  u_short ether_type;                  /* IP? ARP? RARP? etc */
};

void got_packet(u_char *args, const struct pcap_pkthdr *header,
                              const u_char *packet)
{
  struct ethheader *eth = (struct ethheader *)packet;
  if (ntohs(eth->ether_type) == 0x0800) { ... } // IP packet
  ...
}
```

The **packet** argument contains a copy of the packet, including the Ethernet header. We typecast it to the Ethernet header structure.

Now we can access the field of the structure

Processing Captured Packet: IP Header

```
void got_packet(u_char *args, const struct pcap_pkthdr *header,
                              const u_char *packet)
{
  struct ethheader *eth = (struct ethheader *)packet;

  if (ntohs(eth->ether_type) == 0x0800) { // 0x0800 is IP type
    struct ipheader * ip = (struct ipheader *)
                           (packet + sizeof(struct ethheader)); ①

    printf("       From: %s\n", inet_ntoa(ip->iph_sourceip));   ②
    printf("         To: %s\n", inet_ntoa(ip->iph_destip));     ③

    /* determine protocol */
    switch(ip->iph_protocol) {                                  ④
        case IPPROTO_TCP:
            printf("   Protocol: TCP\n");
            return;
        case IPPROTO_UDP:
            printf("   Protocol: UDP\n");
            return;
```

Find where the IP header starts, and typecast it to the IP Header structure.

Now we can easily access the fields in the IP header.

Further Processing Captured Packet

- If we want to further process the packet, such as printing out the header of the TCP, UDP and ICMP, we can use the similar technique.
  - We move the pointer to the beginning of the next header and type-cast
  - We need to use the header length field in the IP header to calculate the actual size of the IP header
- In the following example, if we know the next header is ICMP, we can get a pointer to the ICMP part by doing the following:

```
int ip_header_len = ip->iph_ihl * 4;
u_char *icmp = (struct icmpheader *)
               (packet + sizeof(struct ethheader) + ip_header_len);
```

Packet Spoofing

- When some critical information in the packet is forged, we refer to it as packet spoofing.
- Many network attacks rely on packet spoofing.
- Let's see how to send packets without spoofing.
- Code available on eDimension or http://github.com/kevin-w-du/BookCode

Sending Packets Without Spoofing

```
void main()
{
    struct sockaddr_in dest_info;
    char *data = "UDP message\n";

    // Step 1: Create a network socket
    int sock = socket(AF_INET, SOCK_DGRAM, IPPROTO_UDP);

    // Step 2: Provide information about destination.
    memset((char *) &dest_info, 0, sizeof(dest_info));
    dest_info.sin_family = AF_INET;
    dest_info.sin_addr.s_addr = inet_addr("10.0.2.5");
    dest_info.sin_port = htons(9090);

    // Step 3: Send out the packet.
    sendto(sock, data, strlen(data), 0,
                  (struct sockaddr *)&dest_info, sizeof(dest_info));
    close(sock);
}
```

Testing: Use the `netcat` (`nc`) command to run a UDP server on `10.0.2.5`. We then run the program on the left from another machine. We can see that the message has been delivered to the server machine:

```
seed@Server(10.0.2.5):$ nc -luv 9090
Connection from 10.0.2.6 port 9090 [udp/*] accepted
UDP message
```

Spoofing Packets Using Raw Sockets

There are two major steps in packet spoofing:

- Constructing the packet
- Sending the packet out

Spoofing Packets Using Raw Sockets

```
/*************************************************************
  Given an IP packet, send it out using a raw socket.
*************************************************************/
void send_raw_ip_packet(struct ipheader* ip)
{
    struct sockaddr_in dest_info;
    int enable = 1;

    // Step 1: Create a raw network socket.
    int sock = socket(AF_INET, SOCK_RAW, IPPROTO_RAW);

    // Step 2: Set socket option.
    setsockopt(sock, IPPROTO_IP, IP_HDRINCL,
                     &enable, sizeof(enable));

    // Step 3: Provide needed information about destination.
    dest_info.sin_family = AF_INET;
    dest_info.sin_addr = ip->iph_destip;

    // Step 4: Send the packet out.
    sendto(sock, ip, ntohs(ip->iph_len), 0,
           (struct sockaddr *)&dest_info, sizeof(dest_info));
    close(sock);
}
```

We use *setsockopt()* to enable *IP_HDRINCL* on the socket.

For raw socket programming, since the destination information is already included in the provided IP header, we do not need to fill all the fields

Since the socket type is raw socket, the system will send out the IP packet as is.

Spoofing Packets: Constructing the Packet

Fill in the ICMP Header

```
char buffer[1500];

memset(buffer, 0, 1500);

/*****************************************************
   Step 1: Fill in the ICMP header.
 *****************************************************/
struct icmpheader *icmp = (struct icmpheader *)
                          (buffer + sizeof(struct ipheader));
icmp->icmp_type = 8; //ICMP Type: 8 is request, 0 is reply.

// Calculate the checksum for integrity
icmp->icmp_chksum = 0;
icmp->icmp_chksum = in_cksum((unsigned short *)icmp,
                             sizeof(struct icmpheader));
```

Find the starting point of the ICMP header, and typecast it to the ICMP structure

Fill in the ICMP header fields

Spoofing Packets: Constructing the Packet

Fill in the IP Header

```
/*****************************************************
   Step 2: Fill in the IP header.
 *****************************************************/
struct ipheader *ip = (struct ipheader *) buffer;
ip->iph_ver = 4;
ip->iph_ihl = 5;
ip->iph_ttl = 20;
ip->iph_sourceip.s_addr = inet_addr("1.2.3.4");
ip->iph_destip.s_addr = inet_addr("10.0.2.5");
ip->iph_protocol = IPPROTO_ICMP;
ip->iph_len = htons(sizeof(struct ipheader) +
                    sizeof(struct icmpheader));
```

Typecast the buffer to the IP structure

Fill in the IP header fields

Finally, send out the packe

```
send_raw_ip_packet (ip);
```

Spoofing UDP Packets

```
memset(buffer, 0, 1500);
struct ipheader *ip = (struct ipheader *) buffer;
struct udpheader *udp = (struct udpheader *) (buffer +
                                        sizeof(struct ipheader));

/*********************************************************
   Step 1: Fill in the UDP data field.
 *********************************************************/
char *data = buffer + sizeof(struct ipheader) +
                      sizeof(struct udpheader);
const char *msg = "Hello Server!\n";
int data_len = strlen(msg);
strncpy (data, msg, data_len);

/*********************************************************
   Step 2: Fill in the UDP header.
 *********************************************************/
udp->udp_sport = htons(12345);
udp->udp_dport = htons(9090);
udp->udp_ulen = htons(sizeof(struct udpheader) + data_len);
udp->udp_sum =  0; /* Many OSes ignore this field, so we do not
                      calculate it. */
```

- Constructing UDP packets is similar, except that we need to include the payload data now.

Spoofing UDP Packets (continued)

```
/*************************************************************
   Step 3: Fill in the IP header.
 *************************************************************/

...... /* Code omitted here; same as that in Listing 12.6 */
ip->iph_protocol = IPPROTO_UDP; // The value is 17.
ip->iph_len = htons(sizeof(struct ipheader) +
                    sizeof(struct udpheader) + data_len);
```

Testing: Use the `nc` command to run a UDP server on `10.0.2.5`. We then spoof a UDP packet from another machine. We can see that the spoofed UDP packet was received by the server machine.

```
seed@Server(10.0.2.5):$ nc -luv 9090
Connection from 1.2.3.4 port 9090 [udp/*] accepted
Hello Server!
```

Sniffing and Then Spoofing

- In many situations, we need to capture packets first, and then spoof a response based on the captured packets.
- Procedure (using UDP as example)
  - Use PCAP API to capture the packets of interests
  - Make a copy from the captured packet
  - Replace the UDP data field with a new message and swap the source and destination fields
  - Send out the spoofed reply

UDP Packet

```
void spoof_reply(struct ipheader* ip)
{
    const char buffer[1500];
    int ip_header_len = ip->iph_ihl * 4;
    struct udpheader* udp = (struct udpheader *) ((u_char *)ip +
                                                  ip_header_len);
   if (ntohs(udp->udp_dport) != 9999) {
       // Only spoof UDP packet with destination port 9999
       return;
   }

   // Step 1: Make a copy from the original packet
   memset((char*)buffer, 0, 1500);
   memcpy((char*)buffer, ip, ntohs(ip->iph_len));
   struct ipheader  * newip  = (struct ipheader *) buffer;
   struct udpheader * newudp = (struct udpheader *) (buffer +
   ip_header_len);
   char *data = (char *)newudp + sizeof(struct udpheader);

   // Step 2: Construct the UDP payload, keep track of payload size
   const char *msg = "This is a spoofed reply!\n";
   int data_len = strlen(msg);
   strncpy (data, msg, data_len);
```

UDP Packet (Continued)

```
    // Step 3: Construct the UDP Header
    newudp->udp_sport = udp->udp_dport;
    newudp->udp_dport = udp->udp_sport;
    newudp->udp_ulen = htons(sizeof(struct udpheader) + data_len);
    newudp->udp_sum =  0;

    // Step 4: Construct the IP header (no change for other fields)
    newip->iph_sourceip = ip->iph_destip;
    newip->iph_destip = ip->iph_sourceip;
    newip->iph_ttl = 50; // Rest the TTL field
    newip->iph_len = htons(sizeof(struct ipheader) +
                           sizeof(struct udpheader) + data_len);

    // Step 5: Send out the spoofed IP packet
    send_raw_ip_packet(newip);
}
```

Packing Sniffing Using Scapy

```
#!/usr/bin/python3
from scapy.all import *

print("SNIFFING PACKETS.........")

def print_pkt(pkt):                                    ①
   print("Source IP:", pkt[IP].src)
   print("Destination IP:", pkt[IP].dst)
   print("Protocol:", pkt[IP].proto)
   print("\n")

pkt = sniff(filter='icmp',prn=print_pkt)               ②
```

Spoofing ICMP & UDP Using Scapy

```
#!/usr/bin/python3
from scapy.all import *

print("SENDING SPOOFED ICMP PACKET.........")
ip = IP(src="1.2.3.4", dst="93.184.216.34")  ①
icmp = ICMP()                                ②
pkt = ip/icmp                                ③
pkt.show()
send(pkt,verbose=0)                          ④
```

```
#!/usr/bin/python3
from scapy.all import *

print("SENDING SPOOFED UDP PACKET.........")
ip = IP(src="1.2.3.4", dst="10.0.2.69") # IP Layer
udp = UDP(sport=8888, dport=9090)       # UDP Layer
data = "Hello UDP!\n"                   # Payload
pkt = ip/udp/data      # Construct the complete packet
pkt.show()
send(pkt,verbose=0)
```

Sniffing and Then Spoofing Using Scapy

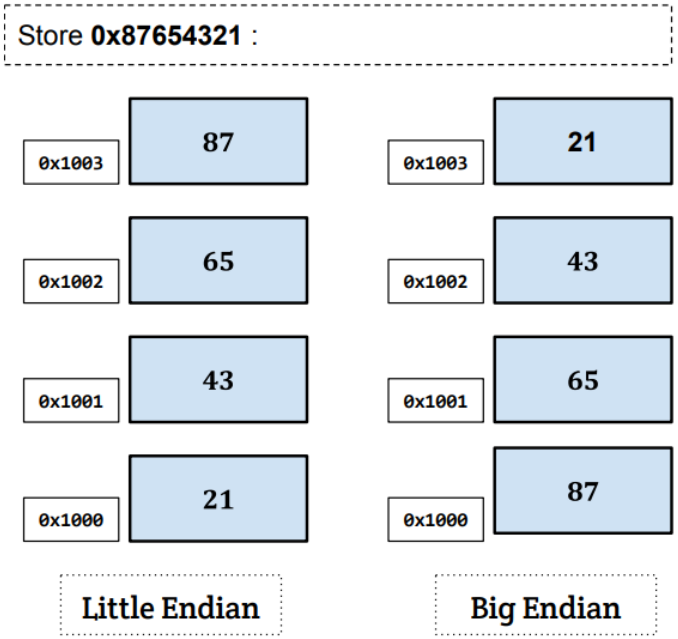
```
#!/usr/bin/python3
from scapy.all import *

def spoof_pkt(pkt):
  if ICMP in pkt and pkt[ICMP].type == 8:
     print("Original Packet.........")
     print("Source IP : ", pkt[IP].src)
     print("Destination IP :", pkt[IP].dst)

     ip = IP(src=pkt[IP].dst, dst=pkt[IP].src, ihl=pkt[IP].ihl)
     icmp = ICMP(type=0, id=pkt[ICMP].id, seq=pkt[ICMP].seq)
     data = pkt[Raw].load
     newpkt = ip/icmp/data

     print("Spoofed Packet.........")
     print("Source IP : ", newpkt[IP].src)
     print("Destination IP :", newpkt[IP].dst)
     send(newpkt,verbose=0)

pkt = sniff(filter='icmp and src host 10.0.2.69',prn=spoof_pkt)
```

Packet Spoofing: Scapy v.s C

- Python + Scapy
  - Pros: constructing packets is very simple
  - Cons: much slower than C code
- C Program (using raw socket)
  - Pros: much faster
  - Cons: constructing packets is complicated
- Hybrid Approach
  - Using Scapy to construct packets
  - Using C to slightly modify packets and then send packets

Endianness

- Endianness: a term that refers to the order in which a given multi-byte data item is stored in memory.
  - **Little Endian**: store the most significant byte of data at the highest address
  - **Big Endian**: store the most significant byte of data at the lowest address



Endianness In Network Communication

- Computers with different byte orders will "misunderstand" each other.
  - Solution: agree upon a common order for communication
  - This is called "network order", which is the same as big endian order
- All computers need to convert data between "host order" and "network order" .

Macro	Description
`htons()`	Convert unsigned short integer from host order to network order.
`htonl()`	Convert unsigned integer from host order to network order.
`ntohs()`	Convert unsigned short integer from network order to host order.
`ntohl()`	Convert unsigned integer from network order to host order.

Summary

- Packet sniffing
  - Using raw socket
  - Using PCAP APIs
- Packet spoofing using raw socket
- Sniffing and the spoofing
- Endianness