Bucket Sort

- uses hashing the values we are sorting are hashed
- Makes assumptions the data, like radix and counting sort
- because it makes assumptions, can sort in O(n) time
- performs best when hashed values of items being sorted are evenly distributed, so there aren't many collisions
 - Not in-place algorithm
- Stability will depend on sort algorithm used to sort the buckets ideally, we want a stable sort
 - to achieve O(n), must have only one item per bucket
- Insertion sort is often used to sort the buckets, because it is fast when the number of items is small

How it works

A generalization of Counting Sort

- Distribute the items into buckets based on their hashed values (scattering phase)
- 2. Sort the items each bucket
- 3. Merge the buckets can just concatenate them (gathering phase)

The value in bucket X must be greater than the values in bucket X - 1 and less than the values in bucket X + 1 — this means that the hash function we us must met this requirement

1) Values

54	46	83	66	95	92	43

2) Place items into Buckets (For example : Hash function uses the value of tens to hash them)

		46 ->	54	66	83	95 ->
		43				92

3) Sort items in Buckets

	43 ->	54	66	83	92 ->
	46				95

4) Merge Buckets

43	46	54	66	83	92	95