TRANSITION TABLES

PAST, PRESENT, AND FUTURE

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PAST

IT ALL STARTED SO INNOCENTLY

From: Kevin Grittner

To: "pgsql-hackers(at)postgresql(dot)org"

Subject: counting algorithm for incremental matview maintenance

Date: 2013-05-14 19:52:06

NOTE DATE

Date: 2013-05-14 19:52:06

Years pass

Kevin refines the concept through coding and conversing on pgsql-hackers.

Then

commit 8c48375e5f43ebd832f93c9166d1fe0e639ff806

Author: Kevin Grittner

Date: Fri Nov 4 10:49:50 2016 -0500

Implement syntax for transition tables in AFTER triggers.

This is infrastructure for the complete SQL standard feature. No support is included at this point for execution nodes or PLs. The intent is to add that soon.

As this patch leaves things, standard syntax can create tuplestores to contain old and/or new versions of rows affected by a statement. References to these tuplestores are in the TriggerData structure. C triggers can access the tuplestores directly, so they are usable, but they cannot yet be referenced within a SQL statement.

Nearly 5 months later...

commit 18ce3a4ab22d2984f8540ab480979c851dae5338

Author: Kevin Grittner

Date: Fri Mar 31 23:17:18 2017 -0500

Add infrastructure to support EphemeralNamedRelation references.

A QueryEnvironment concept is added, which allows new types of objects to be passed into queries from parsing on through execution. At this point, the only thing implemented is a collection of EphemeralNamedRelation objects -- relations which can be referenced by name in queries, but do not exist in the catalogs. The only type of ENR implemented is NamedTuplestore, but provision is made to add more types fairly easily.

PRESENT

A CURRENT EXAMPLE

changelog_trigger

THE ETERNAL QUESTION:

WHAT HAPPEN?

What was going on with client 187666 during February of last year?

What did tables postal, email, and phone look like?

Do you remember what we set that setting to last week?

...and what it had been before we did?

When did we delete that @%^# record?!?

COMMON ANSWER:

DUNNO!

COMMON REASON:

UM, HOW?

ANOTHER COMMON REASON:

TOO FINICKY TO SET UP!

YET ANOTHER COMMON REASON:

MAINTENANCE. UGH.

WHAT WOULD WE LIKE?

No manual steps after setup
Know who made the change
...when they made it
...down to the row and column level
Hardened against schema changes
Easy to query
In other words...
Magic!

OR IS IT?



- Triggers
 - Data changes per statement (New in 10!)
 - Event
- JSONB
- Partitioning

LOGGING DATA STRUCTURE

- Who
- When
- Where
- What

WHO

"user" TEXT DEFAULT CURRENT_USER

WHERE

table_schema TEXT NOT NULL table_name TEXT NOT NULL

WHEN

timestamp TIMESTAMP WITH TIME ZONE!!!

WHAT

- old_row (if applicable)
- new_row (if applicable)
- serialized as JSONB
 - Impervious to DDL
 - Good query tools
 - Index friendly

ROOT (DB) TABLE

```
CREATE TABLE IF NOT EXISTS the_log (
    "timestamp" timestamp with time zone DEFAULT now() NOT NULL,
    "user" text NOT NULL DEFAULT CURRENT_USER,
   action text NOT NULL,
    table_schema text NOT NULL,
    table_name text NOT NULL,
   old_row jsonb,
   new_row jsonb,
   CONSTRAINT the log_check CHECK (
        CASE action
            WHEN 'INSERT' THEN old_row IS NULL
            WHEN 'DELETE' THEN new_row IS NULL
        END
 PARTITION BY LIST(table_schema);
```

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BRANCH (SCHEMA) TABLE

```
CREATE TABLE IF NOT EXISTS public_log PARTITION OF the_log FOR VALUES IN ('public')
PARTITION BY LIST (table_name);
```

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LEAF (TABLE) TABLE

CREATE TABLE IF NOT EXISTS public_foo_log PARTITION OF public_log FOR VALUES IN ('foo');

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HOW DO WE GET STUFF IN THERE?

Per-statement triggers
INSERT
UPDATE
DELETE

INSERT TRIGGER

```
CREATE TRIGGER log_insert_public_foo
AFTER INSERT ON foo
REFERENCING NEW TABLE AS new_table
FOR EACH STATEMENT
EXECUTE PROCEDURE log();
```

DELETE TRIGGER

```
CREATE TRIGGER log_delete_public_foo
AFTER DELETE ON foo
REFERENCING OLD TABLE AS old_table
FOR EACH STATEMENT
EXECUTE PROCEDURE log();
```

UPDATE TRIGGER

(slightly wierd-looking)

```
CREATE TRIGGER log_update_public_foo
AFTER UPDATE ON foo
REFERENCING OLD TABLE AS old_table NEW TABLE AS new_table
FOR EACH STATEMENT
EXECUTE PROCEDURE log();
```

INSERT TRIGGER BODY

```
INSERT INTO the_log (
        action, table_schema, table_name, new_row
)
SELECT
      TG_OP, TG_TABLE_SCHEMA, TG_RELNAME, row_to_json(new_table)::
FROM
      new_table;
```

DELETE TRIGGER BODY

```
INSERT INTO the_log (
        action, table_schema, table_name, old_row
)
SELECT
        TG_OP, TG_TABLE_SCHEMA, TG_RELNAME, row_to_json(old_table)::
FROM
        old_table;
```

UPDATE TRIGGER BODY

```
INSERT INTO the_log (
    action, table_schema, table_name, old_row, new_row
)
SELECT
    TG_OP, TG_TABLE_SCHEMA, TG_RELNAME, old_row, new_row
FROM
    UNNEST(
        ARRAY(SELECT row_to_json(old_table)::jsonb FROM old_table
        ARRAY(SELECT row_to_json(new_table)::jsonb FROM new_table
        ) AS t(old_row, new_row)
```

THAT WAS TEDIOUS AND ERROR-PRONE

WHAT DO WE DO WHEN WE SEE TEDIOUS AND ERROR-PRONE?

AUTOMATE!

EVENT TRIGGER

```
CREATE EVENT TRIGGER add_logger
ON ddl_command_end
WHEN tag IN ('create table')
EXECUTE PROCEDURE add_logger();
```

EVENT TRIGGER FUNCTION

```
CREATE OR REPLACE FUNCTION add_logger()
RETURNS event_trigger
SECURITY DEFINER
LANGUAGE plpgsql
AS $$
DECLARE
    r RECORD;
BEGIN
    SELECT p.*, c.relname as table_name INTO STRICT r
    FROM
        pg_catalog.pg_event_trigger_ddl_commands() p
/* Call add_logging_items() somewhere in here */
END;
```

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FUNCTION THAT ADDS THE DDL

TODO

Propagate indexes like primary keys.
Separate TP from analytics with FDWs or something.
Less code/data on pub with logical decoding.
Stuff I haven't thought of with your help!
Please file bug reports!

FUTURE

MATERIALIZED VIEW MAINTENANCE

"SIMPLE" VIEWS (NO RECURSION OR AGGREGATION)

"Counting algorithm" built atop transition tables

RECURSIVE VIEWS

DRed (Delete and Rederive)
...also built atop transition tables
Heavier weight operation :(

AGGREGATES: THE GOOD

SUM COUNT Stored as-is Simple to update from with transition tables Yay!

AGGREGATES: THE BAD

STDDEV
Stored in a different form
Store AVG as SUM and COUNT, divide before display
PSTORE STDDEV as N, SUM, SUM(X^2), do some
February 7, 2018
COMPLEX arithmetic before display

AGGREGATES: THE UGLY

MEDIAN a.k.a. PERCENTILE_CONT ARRAY_AGG() Stored in a different form that's bulky Recompute entirely at every change. ...or do something too clever to reduce some of the load.

Questions? Comments? Snowballs?



СПАСИБО!