Competitive Programming Library

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# Algorithms

#### 1.1 Count inversions

**Description:** Count the number of inversions when transforming the vector l in the vector r, which is also equivalent to the minimum number of swaps required.

**Usage:** If no r vector is provided it considers r as the sorted vector, if there is no such way to turn l into r using swaps then -1 is returned

**Time**:  $O(N \log N)$ 

```
#pragma once
#include "../Contest/template.cpp"
template <typename T>
ll countInversions(vector<T> l, vector<T> r = {}) {
    if (!len(r)) r = l, sort(all(r));
    int n = len(l);
    vi v(n), bit(n);
    vector<pair<T, int>> w;
    rep(i, 0, n) w.eb(r[i], i + 1);
    sort(all(w));
    rep(i, 0, n) {
        auto it = lower bound(all(w), make pair(l[i], 0));
        if (it == w.end() or it->first != \overline{\lfloor (i) \rfloor} return -1; // impossible
        v[i] = it -> second:
        it->second = -1;
    il ans = 0:
    rrep(i, n - 1, 0 - 1) {
        for (int j = v[i] - 1; j; j = j \& -j) ans += bit[j];
        for (int j = v[i]; j < n; j += j \& -j) bit[j]++;
    return ans;
```

# Ternary search (integer)

**Description:** Given a unimodal function f defined between the integers l and r finds an xsuch that f(x) is maximum/minimum.

**Usage:** Just pass the range l, r of the function you are interested, the function that receives an integer and if you want the maximum value use the cmp = greater < ll>(), otherwise less<ll>().

**Time**:  $O(\log r - l + 1)$ 

Memory: O(1)

```
#include "../Contest/template.cpp"
template <auto cmp = greater<ll>()>
ll ternary search(ll l, ll r, function<ll(ll)> f) {
   static const ll eps = 3:
   while (r - l >= eps) {
       ll m1 = l + (r - l) / 3;
        ll m2 = r - (r - l) / 3;
        if (cmp(f(m1), f(m2)))
            r = m2;
       else
```

```
l = m1:
rep(i, l, r + 1) if (cmp(f(i), f(l))) l = i;
return l;
```

# 1.3 Ternary search (real)

```
#include "../Contest/template.cpp"
template <auto cmp = greater<ld>()>
ld ternarySearch(ld l, ld r, function<ld(ld)> f, const ld eps = 1e-9) {
    while (r - l >= eps) {
        ld m1 = l + (r - l) / 3;
        ld m2 = r - (r - l) / 3;
        if (cmp(f(m1), f(m2)))
            r = m2:
        else
            l = m1:
    }
    return l;
```

### **Combinatorics**

# 2.1 Process all partitions of a set

```
Description: generate every distinct group of a set that contains elements from 0 to N-1,
and pass it to the given function "process". If N is 4 the sets generated would be:
[\{\{0,1,2,3\}\}] \ [\{\{0,1,2\},\{3\}\}] \ [\{\{0,1,3\},\{2\}\}] \ [\{\{0,1\},\{2,3\}\}] \ [\{\{0,1\},\{2\},\{3\}\}] \ [\{\{0,2,3\},\{1\}\}]]
[\{\{0,2\},\{1,3\}\}] \ [\{\{0,2\},\{1\},\{3\}\}] \ [\{\{0,3\},\{1,2\}\}] \ [\{\{0\},\{1,2,3\}\}] \ [\{\{0\},\{1,2\},\{3\}\}]]
[\{\{0,3\},\{1\},\{2\}\}] \ [\{\{0\},\{1,3\},\{2\}\}] \ [\{\{0\},\{1\},\{2,3\}\}] \ [\{\{0\},\{1\},\{2\},\{3\}\}]
Time: O(B(N)), Bell Number of N
Memory: O(N)
```

```
#include "../Contest/template.cpp"
void process all partitions of a set(
    const int N, const function<void(const vi2d &)> process) {
    vi2d groups;
    groups.reserve(N);
    function<void(int)> dfs = [&](int idx) {
        if (idx == N) {
            process(groups);
            return:
        rep(i, 0, len(groups)) {
            groups[i].eb(idx);
            dfs(idx + 1);
            groups[i].ppb();
        groups.pb({idx});
        dfs(idx + 1);
```

```
groups.ppb();
};
_dfs(0);
```

#### 3 Contest

### 3.1 bash config

```
#copy first argument to clipborad ! ONLY WORK ON XORG !
alias clip="xclip -sel clip"
# compile the $1 parameter, if a $2 is provided
# the name will be the the binary output, if
# none is provided the binary name will be
# 'a.out'
comp() {
  echo ">> COMPILING $1 <<" 1>&2
  if [ $# -gt 1 ]; then
    outfile="${2}"
  else
    outfile="a.out"
  time q++-std=c++20 \setminus
    -02 \
    -q3 \
    -Wall \
    -fsanitize=address,undefined \
    -fno-sanitize-recover \
    -D LOCAL \
    -o "${outfile}" \
  if [ $? -ne 0 ]; then
    echo ">> FAILED <<" 1>&2
    return 1
  fi
  echo ">> DONE << " 1>&2
# run the binary given in $1, if none is
# given it will try to run the 'a.out'
# binary
run() {
  to run=./a.out
  if [ -n "$1" ]; then
    to_run="$1"
  fi
  time $to_run
# just comp and run your cpp file
# accpets <in1 >out and everything else
comprun() {
  comp "$1" "a" && run ./a ${@:2}
testall() {
```

```
comp "$1" generator
  comp "$2" brute
  comp "$3" main
  input counter=1
  while true; do
    echo "$input counter"
    run ./generator >input
    run ./main <input >main output.txt
    run ./brute <input >brute output.txt
    diff brute output.txt main output.txt
    if [ $? -ne 0 ]; then
      echo "Outputs differ at input $input counter"
      echo "Brute file output:"
      cat brute output.txt
      echo "Main file output:"
      cat main output.txt
      echo "input used: "
      cat input
      break
    fi
    ((input counter++))
touch macro() {
  cat "$1"/template.cpp >"$2"
 cat "$1"/run.cpp >>"$2"
 cp "$1"/debug.cpp .
# Creates a contest with hame $2
# Copies the macro and debug file from $1
# Already creates files a...z .cpp and .py
prepare contest() {
  mkdir "$2"
  cd "$2"
  for i in {a..z}; do
    touch macro $1 $i.cpp
  done
get file hash() {
  local hash=$(cpp -dD -P -fpreprocessed "$1" | tr -d '[:space:]' | md5sum
      cut -c-6
  echo "$hash"
3.2 debug
template <typename T>
concept Printable = requires(T t) {
    { std::cout << t } -> std::same as<std::ostream &>;
template <Printable T>
void print(const T &x) {
```

```
cerr << x:
template <size t T>
void print(const bitset<T> &x) {
   cerr << x;
template <typename A, typename B>
void print(const pair<A, B> &p);
template <typename... A>
void print(const tuple<A...> &t);
template <typename T>
void print(stack<T> s);
template <typename T>
void print(queue<T> q);
template <typename T, typename... U>
void __print(priority_queue<T, U...> q);
template <typename A>
void print(const A &x) {
   bool first = true;
   cerr << '{';
   for (const auto &i : x) {
       cerr << (first ? "" : ","), print(i);</pre>
       first = false;
   cerr << '}';
template <typename A, typename B>
void print(const pair<A, B> &p) {
   cerr << '(';
     print(p.first);
   cerr << ',';
    print(p.second);
   cerr << ')';
template <typename... A>
void __print(const tuple<A...> &t) {
   bool first = true;
   cerr << '(';
   apply(
        [&first](const auto &...args) {
            ((cerr << (first ? "" : ","), __print(args), first = false),</pre>
   ...);
       t):
   cerr << ')';
template <typename T>
void print(stack<T> s) {
   vector<T> debugVector;
   while (!s.empty()) {
       T t = s.top();
       debugVector.push back(t);
       s.pop();
   reverse(debugVector.begin(), debugVector.end());
   __print(debugVector);
```

```
template <typename T>
void print(queue<T> q) {
    vector<T> debugVector;
    while (!q.empty()) {
        T t = q.front();
        debugVector.push back(t);
        q.pop();
    __print(debugVector);
template <typename T, typename... U>
void print(priority queue<T, U...> q) {
    vector<T> debugVector;
    while (!q.empty()) {
        T t = q.top();
        debugVector.push_back(t);
        q.pop();
    __print(debugVector);
void _print() { cerr << "]\n"; }</pre>
template <typename Head, typename... Tail>
void _print(const Head &H, const Tail &...T) {
    print(H);
    <u>if</u> (sizeof...(T)) cerr << ", ";
    _print(T...);
#define dbg(x...)
    cerr << "[" << #x << "] = ["; \
    print(x)
3.3 run
void run();
int32 t main() {
#ifndef LOCAL
    fastio;
#endif
    int T = 1;
    cin >> T;
    rep(t, 0, T) {
        dbg(t);
        run();
void run() {}
3.4 short-template
```

```
using namespace std;
#define fastio
    ios_base::sync_with_stdio(0);
    cin.tie(0);
void run() {}
int32_t main(void) {
    fastio;
    int t;
    t = 1;
    // cin >> t;
    while (t--) run();
}
```

### 3.5 template

```
#pragma once
#include <bits/stdc++.h>
using namespace std;
#ifdef LOCAL
#include "debug.cpp"
#else
#define dbq(...)
#endif
#define fastio
   ios base::sync with stdio(0); \
    cin.tie(0);
#define all(j) j.begin(), j.end()
#define rall(j) j.rbegin(), j.rend()
#define len(j) (int)j.size()
#define rep(i, a, b) \
    for (common type t < decltype(a), decltype(b) > i = (a); i < (b); i++)
#define rrep(i, a, b) \
    for (common type t < decltype(a), decltype(b) > i = (a); i > (b); i--)
#define trav(xi, xs) for (auto &xi : xs)
#define rtrav(xi, xs) for (auto &xi : ranges::views::reverse(xs))
using ll = long long;
#define endl '\n'
#define pb push_back
#define pf push front
#define ppb pop back
#define ppf pop front
#define eb emplace back
#define ef emplace back
#define lb lower bound
#define ub upper bound
#define fi first
#define se second
#define emp emplace
#define ins insert
\#define\ divc(a, b)\ ((a) + (b) - 1ll)\ /\ (b)
using str = string;
using ull = unsigned long long;
using ld = long double;
using vll = vector<ll>:
```

```
using pll = pair<ll, ll>;
using vll2d = vector<vll>:
using vi = vector<int>;
using vi2d = vector<vi>;
using pii = pair<int, int>;
using vpii = vector<pii>;
using vc = vector<char>;
using vs = vector<str>;
template <typename T, typename T2>
using umap = unordered_map<T, T2>;
template <tvpename T>
using pgmn = priority queue<T, vector<T>, greater<T>>;
template <typename T>
using pqmx = priority queue<T, vector<T>>;
template <typename T, typename U>
inline bool chmax(T &a, U const &b) {
    return (a < b ? a = b, 1 : 0);
template <typename T, typename U>
inline bool chmin(T &a, U const &b) {
    return (a > b ? a = b, 1 : 0);
template <typename T>
std::istream &operator>>(std::istream &is, std::vector<T> &vec) {
    for (auto &element : vec) {
        is >> element:
    return is;
template <typename T> // print vector
ostream &operator<<(ostream &os, vector<T> &xs) {
    rep(i, os.iword(0), xs.size()) os << xs[i] << (i == xs.size() ? "" : "
    ");
    os.iword(0) = 0;
    return os;
}
3.6 vim config
```

```
set sta nu rnu sc cindent noswapfile
set ts=2 sw=2
set bg=dark ruler clipboard=unnamed,unnamedplus, timeoutlen=100
colorscheme default
syntax on
" Takes the hash of the selected text and put
" in the vim clipboard
function! HashSelectedText()
    " Yank the selected text to the unnamed register
    normal! gvy
    " Use the system() function to call sha256sum with the yanked text
    let l:hash = system('echo' . shellescape(@@) . ' | sha256sum')
    " Yank the hash into Vim's unnamed register
    let @" = l:hash
endfunction
```

## 4 Data Structures

### 4.1 2D Segment Tree

#### 4.1.1 Point update query sum

```
#include "../../Contest/template.cpp"
template <typename T, auto op>
struct SeamentTree2D {
   int h. w:
   vector<vector<T>> t;
   SegmentTree2D(const vector<vector<T>> &a)
        : h(a.size()), w(a.back().size()), t(h * 4, vector<T>(w * 4)) {
       build x(1, 0, h - 1, a);
   void build_y(int vx, int lx, int rx, int vy, int ly, int ry,
                const vector<vector<T>> &a) {
       if (ly == ry) {
            if (lx == rx)
                t[vx][vy] = a[lx][ly];
                t[vx][vy] = op(t[vx * 2][vy], t[vx * 2 + 1][vy]);
       } else {
            int my = (ly + ry) / 2;
            build_y(vx, lx, rx, vy * 2, ly, my, a);
            build_y(vx, lx, rx, vy * 2 + 1, my + 1, ry, a);
            t[vx][vy] = op(t[vx][vy * 2], t[vx][vy * 2 + 1]);
       }
   void build x(int vx, int lx, int rx, const vector<vector<T>> &a) {
       if (lx != rx) {
            int mx = (lx + rx) / 2;
            build x(vx * 2, lx, mx, a);
            build x(vx * 2 + 1, mx + 1, rx, a);
       build y(vx, lx, rx, 1, 0, w - 1, a);
   T query y(int vx, int vy, int tly, int try , int ly, int ry) {
       if (ly > ry) return 0;
       if (ly == tly && try_ == ry) return t[vx][vy];
       int tmy = (tly + try) / 2;
       return op(query y(vx, vy * 2, tly, tmy, ly, min(ry, tmy)),
                  query y(vx, vy * 2 + 1, tmy + 1, try, max(ly, tmy + 1),
    ry));
   T query x(int vx, int tlx, int trx, int lx, int rx, int ly, int ry) {
       if (lx > rx) return 0;
       if (lx == tlx \&\& trx == rx) return query y(vx, 1, 0, w - 1, ly, ry)
   );
       int tmx = (tlx + trx) / 2;
        return op(
            query x(vx * 2, tlx, tmx, lx, min(rx, tmx), ly, ry),
            query_x(vx * 2 + 1, tmx + 1, trx, max(lx, tmx + 1), rx, ly, ry
   ));
```

```
void update y(int vx, int lx, int rx, int vy, int ly, int ry, int x,
                  int new val) {
        if (ly == ry) {
            if (lx == rx)
                t[vx][vy] = new val;
                t[vx][vy] = op(t[vx * 2][vy], t[vx * 2 + 1][vy]);
        } else {
            int my = (ly + ry) / 2;
            if (y \le my)
                update_y(vx, lx, rx, vy * 2, ly, my, x, y, new_val);
            else
                update y(vx, lx, rx, vy * 2 + 1, my + 1, ry, x, y, new val
   );
            t[vx][vy] = op(t[vx][vy * 2], t[vx][vy * 2 + 1]);
    }
    void update x(int vx, int lx, int rx, int x, int y, T new val) {
        if (lx != rx) {
            int mx = (lx + rx) / 2;
            if (x \le mx)
                update_x(vx * 2, lx, mx, x, y, new_val);
                update x(vx * 2 + 1, mx + 1, rx, x, y, new val);
        update_y(vx, lx, rx, 1, 0, w - 1, x, y, new_val);
    }
    T query(int lx, int rx, int ly, int ry) {
        return query x(1, 0, h - 1, lx, rx, ly, ry);
};
```

# 4.2 SQRT decomposition

#### 4.2.1 two-sequence-queries

```
r = numeric limits<int>::min();
        x = y = prod = sum as = sum bs = 0;
    };
};
int sqrtLen;
vector<t_sqrt> blocks;
vector<ll> as, bs:
SqrtDecomposition(const vector<ll> &as , const vector<ll> &bs ) {
    int n = as .size();
    sgrtLen = (int) sgrt(n + .0) + 1;
    blocks.resize(sqrtLen + 6.66);
    as = as_{;}
    bs = bs;
    for (int i = 0; i < n; i++) {
        auto &bi = blocks[i / sqrtLen];
        bi.l = min(bi.l, i);
        bi.r = max(bi.r, i);
        bi.sum as = sum(bi.sum as, as[i]);
        bi.sum bs = sum(bi.sum bs, bs[i]);
        bi.prod = sum(bi.prod, mul(as[i], bs[i]));
}
// adds x to a[i], and y to b[i], in range [l,
void update(int l, int r, ll x, ll y) {
    auto apply1 = [\&](int idx, ll x, ll y) -> void {
        auto &block = blocks[idx / sqrtLen];
        block.prod = sub(block.prod, mul(as[idx], bs[idx]));
        block.sum as = sub(block.sum as, as[idx]);
        block.sum bs = sub(block.sum bs, bs[idx]);
        as[idx] = sum(as[idx], x);
        bs[idx] = sum(bs[idx], y);
        block.prod = sum(block.prod, as[idx] * bs[idx]);
        block.sum as = sum(block.sum as, as[idx]);
        block.sum bs = sum(block.sum bs, bs[idx]);
    };
    auto apply2 = [&](int idx, ll x, ll y) -> void {
        blocks[idx].x = sum(blocks[idx].x, x);
        blocks[idx].y = sum(blocks[idx].y, y);
    };
    int cl = l / sqrtLen, cr = r / sqrtLen;
    if (cl == cr) {
        for (int i = l; i <= r; i++) {
            apply1(i, x, y);
    } else {
        for (int i = l; i \le (cl + 1) * sqrtLen - 1; i++) {
            apply1(i, x, y);
        for (int i = cl + 1; i \le cr - 1; i++) {
            apply2(i, x, y);
        }
```

```
for (int i = cr * sqrtLen; i <= r; i++) {
                apply1(i, x, y);
    }
    // sum of a[i]*b[i] in range [l r]
    ll query(int l, int r) {
        auto eval1 = [&](int idx) -> ll {
            auto &block = blocks[idx / sqrtLen];
            return mul(sum(as[idx], +block.x), sum(bs[idx], block.y));
        };
        auto eval2 = [\&](int idx) \rightarrow ll \{
            auto &block = blocks[idx];
            ll ret = 0:
            ret = sum(
                ret, mul(mul(block.x, block.y), sum(sub(block.r, block.l),
    1)));
            ret = sum(ret, block.prod);
            ret = sum(ret, block.y * block.sum as);
            ret = sum(ret, block.x * block.sum bs);
            return ret;
        };
        ll ret = 0;
        int cl = l / sqrtLen, cr = r / sqrtLen;
        if (cl == cr) {
            for (int i = l; i <= r; i++) {
                ret = sum(ret, eval1(i));
        } else {
            for (int i = l; i \le (cl + 1) * sqrtLen - 1; i++) {
                ret = sum(eval1(i), ret);
            for (int i = cl + 1; i \le cr - 1; i++) {
                ret = sum(ret, eval2(i));
            for (int i = cr * sqrtLen; i <= r; i++) {
                ret = sum(ret, eval1(i));
        }
        return ret;
};
```

# 4.3 Segment Tree Point Update Range Query (bottom-up)

#### 4.3.1 Query GCD

```
using ll = long long;
struct Node {
    ll value;
    bool undef;
    Node() : value(1), undef(1) {}; // Neutral element
```

```
Node(ll v) : value(v), undef(0) {};
};
inline Node combine(const Node &nl, const Node &nr) {
   if (nl.undef) return nr;
   if (nr.undef) return nl;
   Node m:
   m.value = gcd(nl.value, nr.value);
   m.undef = false:
   return m;
template <typename T = Node, auto F = combine>
struct SegTree {
   int n:
    vector<T> st;
   SegTree(int n) : n(n), st(n << 1) {}
   void assign(int p, const T &k) {
        for (st[p += n] = k; p >>= 1;) st[p] = F(st[p << 1], st[p << 1]
   1]);
   T query(int l, int r) {
       T ansl, ansr;
        for (l += n, r += n + 1; l < r; l >>= 1, r >>= 1) {
            if (l \& 1) ans l = F(ans l, st[l++]):
            if (r \& 1) ansr = F(st[--r], ansr);
        return F(ansl, ansr);
};
```

#### 4.3.2 Query Max Subarray Sum

```
#pragma once
#include "../../Contest/template.cpp"
#include "./Struct.cpp"
const ll oo = 1e9;
struct Node {
   ll tot, suf, pref, best;
    // Neutral element
   Node(): tot(-oo), suf(-oo), pref(-oo), best(-oo) {} // Neutral
   element
   // for assign
   Node(ll x) { tot = x, suf = x, pref = x, best = max(0ll, x); }
};
Node combine(Node &nl, Node &nr) {
   if (nl.tot == - oo) return nr;
   if (nr.tot == - oo) return nl;
   Node m:
   m.tot = nl.tot + nr.tot:
   m.pref = max({nl.pref, nl.tot + nr.pref});
   m.suf = max({nr.suf, nr.tot + nl.suf});
   m.best = max({nl.best, nr.best, nl.suf + nr.pref});
   return m:
```

```
using SegTreeMaxSubarraySum = SegTreeBottomUp<Node, Node(), combine>;
```

#### 4.3.3 Query min

#### 4.3.4 Query sum

```
#pragma once
#include "../../Contest/template.cpp"
#include "./Struct.cpp"
template <typename T>
using SegTreeBottomUpSumQuery =
    SegTreeBottomUp<T, T(0), [](T a, T b) { return a + b; }>;
```

#### 4.3.5 Struct

```
* @Description:
        merge should be function\langle T(T,T)\rangle, that
        makes the necessary operation between two
        nodes in the seament tree
 *
* */
#pragma once
#include "../../Contest/template.cpp"
template <typename T, T identity, auto merge>
struct SegTreeBottomUp {
    int size:
    vector<T> arr:
    SegTreeBottomUp(int n) {
        for (size = 1: size < n: size <<= 1):
        arr.resize(size << 1);</pre>
    void assign(int pos, const T &val) {
        for (arr[pos += size] = val; pos >>= 1;)
            arr[pos] = merge(arr[pos << 1], arr[pos << 1 | 1]);</pre>
    T query(int l, int r) {
        T ans l = identity, ans r = identity;
        for (\bar{l} += size, r += size + 1; l < r; l >>= 1, r >>= 1) {
            if (l & 1) ans l = merge(ans l, arr[l++]);
            if (r \& 1) ans r = merge(arr[--r], ans r);
```

```
return merge(ans l, ans r);
    SegTreeBottomUp(const vector<T> &vec) : SegTreeBottomUp(len(vec)) {
        copy(all(vec), begin(arr) + size);
        rrep(i, size - 1, 0) arr[i] = merge(arr[i << 1], arr[i << 1 | 1]);
};
4.4 Segment tree (dynamic)
4.4.1 Range Max Query Point Max Assignment
Description: Answers range queries in ranges until 10<sup>9</sup> (maybe more)
Time: Query and update O(n \cdot \log n)
struct node:
node *newNode():
struct node {
    node *left, *right;
    int lv, rv;
    ll val;
    node() : left(NULL), right(NULL), val(-oo) {}
    inline void init(int l, int r) {
        lv = l:
        rv = r;
    inline void extend() {
        if (!left) {
            int m = (lv + rv) / 2;
            left = newNode();
            right = newNode();
            left->init(lv, m);
            right->init(m + 1, rv);
    }
    ll query(int l, int r) {
        if (r < lv || rv < l) {</pre>
            return 0;
        if (l <= lv && rv <= r) {
            return val;
        extend();
        return max(left->query(l, r), right->query(l, r));
    void update(int p, ll newVal) {
        if (lv == rv) {
            val = max(val, newVal);
            return;
        }
        extend():
        (p <= left->rv ? left : right)->update(p, newVal);
        val = max(left->val, right->val);
```

```
};
const int BUFFSZ(1e7);
node *newNode() {
    static int bufSize = BUFFSZ;
    static node buf[(int)BUFFSZ];
    assert(bufSize);
    return &buf[--bufSize];
struct SegTree {
    int n;
    node *root;
    SegTree(int n) : n( n) {
        root = newNode():
        root->init(0, n);
    ll query(int l, int r) { return root->query(l, r); }
    void update(int p, ll v) { root->update(p, v); }
};
4.4.2 Range Sum Query Point Sum Update
Description: Answers range queries in ranges until 10<sup>9</sup> (maybe more)
Time: Query and update in O(n \cdot \log n)
struct node;
node *newNode();
struct node {
    node *left, *right;
    int lv. rv:
    ll val;
    node() : left(NULL), right(NULL), val(0) {}
    inline void init(int l, int r) {
        lv = l;
        rv = r;
    inline void extend() {
        if (!left) {
             int m = (rv - lv) / 2 + lv;
            left = newNode();
             right = newNode();
            left->init(lv, m);
             right->init(m + 1, rv);
        }
    ll query(int l, int r) {
        if (r < lv || rv < l) {</pre>
             return 0;
        if (l <= lv && rv <= r) {</pre>
             return val;
        extend();
```

```
return left->query(l, r) + right->query(l, r);
   }
   void update(int p, ll newVal) {
        if (lv == rv) {
            val += newVal:
            return:
        }
        extend():
        (p <= left->rv ? left : right)->update(p, newVal);
        val = left->val + right->val;
};
const int BUFFSZ(1.3e7);
node *newNode() {
   static int bufSize = BUFFSZ:
   static node buf((int)BUFFSZ1:
   // assert(bufSize):
   return &buf[--bufSizel:
struct SegTree {
   int n:
   node *root:
   SegTree(int _n) : n(_n) {
        root = newNode():
        root->init(0, n);
   ll query(int l, int r) { return root->query(l, r); }
   void update(int p, ll v) { root->update(p, v); }
};
```

## 4.5 Segment tree point update range query (top-down)

### 4.5.1 Query hash (top down)

```
#include "../../Contest/template.cpp"
const ll MOD = 1'000'000'009:
const ll P = 31;
const int MAXN = 2'000'000;
ll pows[MAXN + 1];
void computepows() {
    pows[0] = 1:
    for (int i = 1; i <= MAXN; i++) {
        pows[i] = (pows[i - 1] * P) % MOD:
struct Node {
   ll hash;
   Node(): hash(-1) {}; // Neutral element
   Node(ll v) : hash(v) {};
};
inline Node combine(Node &vl, Node &vr, int nl, int nr, int al, int ar) {
   if (vl.hash == -1) return vr;
   if (vr.hash == -1) return vl:
```

```
Node vm:
    int nm = midpoint(nl, nr);
    int lsize = min(nm, qr) - max(nl, ql) + 1;
    vm.hash = (vl.hash + ((vr.hash * pows[lsize]) % MOD)) % MOD;
    return vm:
template <typename T = Node, auto F = combine>
struct SegTree {
    int n:
    vector<T> st:
    SegTree(int n) : n(n), st(n \ll 2) {}
    void assign(int p, const T \&v) { assign(1, 0, n - 1, p, v); }
    void assign(int node, int l, int r, int p, const T &v) {
        if (l == r) {
            st[node] = v:
            return;
        int m = midpoint(l, r);
        if (p \ll m)
            assign(node << 1, l, m, p, v);
        else
            assign(node << 1 | 1, m + 1, r, p, v);
        st[node] = F(st[node << 1], st[node << 1 | 1], l, r, l, r);
    inline T query(int l, int r) { return query(1, 0, n - 1, l, r); }
    inline T query(int node, int nl, int nr, int l, int r) const {
        if (r < nl or nr < l) return T():
        if (l <= nl and nr <= r) return st[node];</pre>
        int m = midpoint(nl, nr);
        auto a = query(node << 1, nl, m, l, r);</pre>
        auto b = query(node << 1 | 1, m + 1, nr, l, r);
        return F(a, b, nl, nr, l, r);
};
```

## 4.6 Segment tree range update range query

#### 4.6.1 Arithmetic progression sum update, query sum

**Description**: Makes arithmetic progression updates in range and sum queries.

Usage: Considering PA(A, R) = [A + R, A + 2R, A + 3R, ...]

- $update_set(l, r, A, R)$ : sets [l, r] to PA(A, R)
- update\_add(l, r, A, R): sum PA(A, R) in [l, r]
- query(l, r): sum in range [l, r]

**Time**: build O(N), updates and queries O(log N)

```
const ll oo = 1e18;
struct SegTree {
    struct Data {
        ll sum;
        ll set_a, set_r, add_a, add_r;
        Data() : sum(0), set a(oo), set r(0), add a(0), add r(0) {}
```

```
};
int n;
vector<Data> seq;
SegTree(int n ) : n(n ), seg(vector<Data>(4 * n)) {}
void prop(int p, int l, int r) {
    int sz = r - l + 1;
    ll &sum = seg[p].sum, &set_a = seg[p].set_a, &set_r = seg[p].set_r
       &add a = seq[p].add a, &add r = seq[p].add r;
    if (set a != oo) {
        set a += add a, set r += add r;
        sum = set a * sz + set r * sz * (sz + 1) / 2;
        if (l != r) {
            int m = (l + r) / 2;
            seg[2 * p].set_a = set_a;
            seg[2 * p].set r = set r;
            seg[2 * p].add a = seg[2 * p].add r = 0;
            seq[2 * p + 1].set a = set a + set r * (m - l + 1);
            seq[2 * p + 1].set r = set r;
            seg[2 * p + 1].add a = seg[2 * p + 1].add r = 0;
        set a = oo, set r = 0;
        add a = add r = 0;
    } else \overline{i}f (add \overline{a} or add r) {
        sum += add a * sz + add r * sz * (sz + 1) / 2;
        if (l != r) {
            int m = (l + r) / 2;
            seg[2 * p].add a += add a;
            seg[2 * p].add r += add r;
            seq[2 * p + 1].add a += add a + add r * (m - l + 1);
            seq[2 * p + 1].add r += add r;
        add a = add r = 0;
int inter(pii a, pii b) {
    if (a.first > b.first) swap(a, b);
    return max(0, min(a.second, b.second) - b.first + 1);
ll set(int a, int b, ll aa, ll rr, int p, int l, int r) {
    prop(p, l, r);
    if (b < l or r < a) return seq[p].sum;</pre>
    if (a <= l and r <= b) {
        seg[p].set_a = aa;
        seg[p].set r = rr;
        prop(p, l, r);
        return seq[p].sum;
    int m = (l + r) / 2;
    int tam l = inter({l, m}, {a, b});
    return \overline{seg}[p].sum = set(a, b, aa, rr, 2 * p, l, m) +
                         set(a, b, aa + rr * tam l, rr, 2 * p + 1, m +
1, r);
}
```

```
void update set(int l, int r, ll aa, ll rr) {
        set(l, r, aa, rr, 1, 0, n - 1);
    ll add(int a, int b, ll aa, ll rr, int p, int l, int r) {
        prop(p, l, r);
        if (b < l or r < a) return seq[p].sum;</pre>
        if (a <= l and r <= b) {
            seq[p].add a += aa;
            seg[p].add r += rr;
            prop(p, l, r);
            return seg[p].sum;
        int m = (l + r) / 2:
        int tam_l = inter({l, m}, {a, b});
        return seg[p].sum = add(a, b, aa, rr, 2 * p, l, m) +
                            add(a, b, aa + rr * tam l, rr, 2 * p + 1, m +
   1, r);
    void update add(int l, int r, ll aa, ll rr) {
        add(l, r, aa, rr, 1, 0, n - 1);
    ll query(int a, int b, int p, int l, int r) {
        prop(p, l, r);
        if (b < l or r < a) return 0;</pre>
        if (a <= l and r <= b) return seq[p].sum;
        int m = (l + r) / 2;
        return query(a, b, 2 * p, l, m) + query(a, b, 2 * p + 1, m + 1, r)
    Il query(int l, int r) { return query(l, r, 1, 0, n - 1); }
};
4.6.2 Increment update query min & max (bottom up)
using SegT = ll;
    SegT mx, mn;
    OuervT()
```

```
struct QueryT {
        : mx(numeric limits<SeqT>::min()), mn(numeric limits<SeqT>::max())
    QueryT(SegT _v) : mx(_v), mn(_v) {}
};
inline QueryT combine(QueryT ln, QueryT rn, pii lr1, pii lr2) {
    chmax(ln.mx, rn.mx);
    chmin(ln.mn, rn.mn);
    return ln;
using LazyT = SeqT;
inline QueryT applyLazyInQuery(QueryT q, LazyT l, pii lr) {
    if (q.mx == QueryT().mx) q.mx = SeqT();
    if (q.mn == QueryT().mn) q.mn = SeqT();
    q.mx += l, q.mn += l;
    return q;
```

```
inline LazyT applyLazyInLazy(LazyT a, LazyT b) { return a + b; }
using UpdateT = SegT;
inline QueryT applyUpdateInQuery(QueryT q, UpdateT u, pii lr) {
   if (q.mx == QueryT().mx) q.mx = SeqT();
   if (q.mn == QueryT().mn) q.mn = SegT();
   q.mx += u, q.mn += u;
    return q;
inline LazyT applyUpdateInLazy(LazyT l, UpdateT u, pii lr) { return l + u;
template <typename Qt = QueryT, typename Lt = LazyT, typename Ut = UpdateT
          auto C = combine, auto ALQ = applyLazyInQuery,
          auto ALL = applyLazyInLazy, auto AUQ = applyUpdateInQuery,
          auto AUL = applyUpdateInLazy>
struct LazySegmentTree {
   int n, h;
   vector<Qt> ts;
   vector<Lt> ds;
   vector<pii> lrs;
   LazySegmentTree(int n)
        : n(_n),
          h(sizeof(int) * 8 - builtin clz(n)),
          ts(n \ll 1).
          ds(n),
          lrs(n << 1) {
        rep(i, 0, n) lrs[i + n] = {i, i};
        rrep(i, n - 1, 0) {
            lrs[i] = {lrs[i << 1].first, lrs[i << 1 | 1].second};</pre>
   LazySegmentTree(const vector<Qt> &xs) : LazySegmentTree(len(xs)) {
        copy(all(xs), ts.begin() + n);
        rep(i, 0, n) lrs[i + n] = {i, i};
        rrep(i, n - 1, 0) {
            ts[i] = C(ts[i << 1], ts[i << 1 | 1], lrs[i << 1], lrs[i << 1
   | 1]);
   void set(int p, Qt v) {
        ts[p + n] = v;
        build(p + n);
   void upd(int l, int r, Ut v) {
        l += n. r += n + 1:
        int 10 = 1, r0 = r:
        for (; l < r; l >>= 1, r >>= 1) {
            if (l & 1) apply(l++, v);
            if (r & 1) apply(--r, v);
        build(l0), build(r0 - 1);
   }
```

```
Qt qry(int l, int r) {
        l += n, r += n + 1:
        push(l), push(r - 1);
        Qt resl = Qt(), resr = Qt();
        pii lr1 = \{l, l\}, lr2 = \{r, r\};
        for (; l < r; l >>= 1, r >>= 1) {
            if (l & 1) resl = C(resl, ts[l], lr1, lrs[l]), l++;
            if (r & 1) r--, resr = C(ts[r], resr, lrs[r], lr2);
        return C(resl, resr, lr1, lr2);
    void build(int p) {
        while (p > 1) {
            p >>= 1;
            ts[p] =
                ALO(C(ts[p << 1], ts[p << 1 | 1], lrs[p << 1], lrs[p << 1]
    | 1]),
                    ds[p], lrs[p]);
    void push(int p) {
        rrep(s, h, 0) {
            int i = p \gg s;
            if (ds[i] != Lt()) {
                apply(i << 1, ds[i]), apply(i << 1 | 1, ds[i]);
                ds[i] = Lt();
    inline void apply(int p, Ut v) {
        ts[p] = AUQ(ts[p], v, lrs[p]);
        if (p < n) ds[p] = AUL(ds[p], v, lrs[p]);
};
4.6.3 Increment update sum query (top down)
```

```
struct Lnode {
    ll v;
    bool assign;
    Lnode() : v(), assign() {} // Neutral element
    Lnode(ll v, bool a = 0) : v(v), assign(a) {};
};
using Qnode = ll;
using Unode = Lnode;
struct LSegTree {
    int n, ql, qr;
    vector<0node> st:
    vector<Lnode> lz;
    Qnode merge(Qnode lv, Qnode rv, int nl, int nr) { return lv + rv; }
    void prop(int i, int l, int r) {
        if (lz[i].assign) {
```

```
st[i] = lz[i].v * (r - l + 1);
        if (l != r) lz[tol(i)] = lz[tor(i)] = lz[i];
    } else {
        st[i] += lz[i].v * (r - l + 1);
        if (l != r) lz[tol(i)].v += lz[i].v, lz[tor(i)].v += lz[i].v;
    lz[i] = Lnode();
void applyV(int i, Unode v) {
    if (v.assign) {
        lz[i] = v;
    } else {
        lz[i].v += v.v;
LSegTree() {}
LSegTree(int _n) : n(_n), st(_n << 2), lz(_n << 2) {}
bool disjoint(int l, int r) { return qr < l or r < ql; }</pre>
bool contains(int l, int r) { return ql <= l and r <= qr; }</pre>
int tol(int i) { return i << 1; }</pre>
int tor(int i) { return i << 1 | 1; }</pre>
void build(vector<Qnode> &v) { build(v, 1, 0, n - 1); }
void build(vector<Qnode> &v, int i, int l, int r) {
    if (l == r) {
        st[i] = v[l];
        return;
    int m = midpoint(l, r);
    build(v, tol(i), l, m);
    build(v, tor(i), m + 1, r);
    st[i] = merge(st[tol(i)], st[tor(i)], l, r);
void upd(int l, int r, Unode v) {
    al = l, ar = r:
    upd(1, 0, n - 1, v);
void upd(int i, int l, int r, Unode v) {
    prop(i, l, r);
    if (disjoint(l, r)) return;
    if (contains(l, r)) {
        applyV(i, v);
        prop(i, l, r);
        return;
    int m = midpoint(l, r);
    upd(tol(i), l, m, v);
    upd(tor(i), m + 1, r, v);
    st[i] = merge(st[tol(i)], st[tor(i)], l, r);
Qnode qry(int l, int r) {
    ql = l, qr = r;
    return gry(1, 0, n-1);
Qnode gry(int i, int l, int r) {
    prop(i, l, r);
```

```
if (disjoint(l, r)) return Qnode();
    if (contains(l, r)) return st[i];
    int m = midpoint(l, r);
    return merge(qry(tol(i), l, m), qry(tor(i), m + 1, r), l, r);
};
```

## 4.7 2D Sparse Table

```
const int N = 1001:
ll matrix[N][N];
ll M[1001][1001][10][10];
ll op(ll a, ll b) { return gcd(a, b); }
void SparseMatrix(int n, int m) {
           int i, j, x, y;
           for (i = 0; (1 << i) <= n; i++) {
                       for (j = 0; (1 << j) <= m; j++) {
                                  for (x = 0; (x + (1 << i) - 1) < n; x++) {
                                             for (y = 0; (y + (1 << j) - 1) < m; y++) {
                                                        if (i == 0 \&\& i == 0)
                                                                    M[x][y][i][i] = matrix[x][y];
                                                         else if (i == 0)
                                                                    M[x][y][i][j] = op(M[x][y][i][j-1],
                                                                                                                          M[x][y + (1 << (j - 1))][i][j -
             1]);
                                                         else if (i == 0)
                                                                   M[x][y][i][j] = op(M[x][y][i - 1][j],
                                                                                                                          M[x + (1 << (i - 1))][y][i -
          1][j]);
                                                         else {
                                                                    int tempa = op(M[x + (1 << (i - 1))][y][i - 1][i -
             1],
                                                                                                              M[x][y + (1 << (j - 1))][i - 1][j -
             1]);
                                                                    int tempb = op(M[x][y][i-1][j-1],
                                                                                                              M[x + (1 << (i - 1))][y 
          - 1))]
                                                                                                                 [i - 1][i - 1]);
                                                                   M[x][y][i][j] = op(tempa, tempb);
                                             }
           return;
int lg2(int x) { return sizeof(int) * 8 - builtin clz(x) - 1; }
ll query2d(int x, int y, int x1, int y1) {
           int k = lg2(x1 - x + 1);
           int l = lg2(y1 - y + 1);
           int tempa = op(M[x][y][k][l], M[x1 - (1 << k) + 1][y][k][l]);
           int tempb = op(M[x][y1 - (1 << l) + 1][k][l],
                                                     M[x1 - (1 \ll k) + 1][y1 - (1 \ll l) + 1][k][l]);
           return op(tempa, tempb);
```

```
}
```

#### 4.8 Bitree 2D

**Description**: Given a 2D array you can increment an arbitrary position, and also query the subsum of a subgrid

**Time**: Update and query in  $O(logN^2)$ 

```
struct Bit2d {
    int n;
    vll2d bit;
    Bit2d(int ni) : n(ni), bit(n + 1, vll(n + 1)) {}
   Bit2d(int ni, vll2d &xs) : n(ni), bit(n + 1, vll(n + 1)) {
        for (int i = 1: i <= n: i++) {
            for (int j = 1; j \le n; j++) {
                update(i, j, xs[i][j]);
        }
    void update(int x, int y, ll val) {
        for (; x \le n; x += (x \& (-x)))  {
            for (int i = y; i \le n; i += (i \& (-i))) {
                bit[x][i] += val;
        }
    il sum(int x, int y) {
        ll ans = 0:
        for (int i = x; i; i = (i & (-i))) {
            for (int j = y; j; j = (j \bar{k}'(-j))) {
                ans += bit[i][i];
            }
        return ans;
    il query(int x1, int y1, int x2, int y2) {
        return sum(x2, y2) - sum(x2, y1 - 1) - sum(x1 - 1, y2) +
               sum(x1 - 1, y1 - 1);
};
```

## 4.9 Convex Hull Trick / Line Container

**Description:** Container where you can add lines of the form mx + b, and query the maximum value at point x.

**Usage**:  $insert\_line(m, b)$  inserts the line  $m \cdot x + b$  in the container.

eval(x) find the highest value among all lines in the point x.

**Time**: Eval and insert in  $O(\log N)$ 

```
const ll LLINF = 1e18;
const ll is_query = -LLINF;
struct Line {
    ll m, b;
    mutable function<const Line *()> succ;
```

```
bool operator<(const Line &rhs) const {</pre>
        if (rhs.b != is query) return m < rhs.m;</pre>
        const Line *s = succ();
        if (!s) return 0:
        ll \dot{x} = rhs.m:
        return b - s -> b < (s -> m - m) * x;
};
struct Cht : public multiset<Line> { // maintain
                                        // max m*x+b
    bool bad(iterator y) {
        auto z = next(y);
        if (y == begin()) {
            if (z = end()) return 0;
            return y->m == z->m \&\& y->b <= z->b;
        auto x = prev(y);
        if (z == end()) return y->m == x->m && y->b <= x->b;
        return (ld)(x->b - v->b) * (z->m - v->m) >=
               (ld)(v->b-z->b) * (y->m-x->m);
    void insert line(ll m,
                     ll b) { // min -> insert (-m,-b) -> -eval()
        auto y = insert({m, b});
        y-succ = [=] { return next(y) == end() ? 0 : &*next(y); };
        if (bad(y)) {
            erase(y);
            return;
        while (next(y) != end() && bad(next(y))) erase(next(y));
        while (y != begin() && bad(prev(y))) erase(prev(y));
    il eval(ll x) {
        auto l = *lower bound((Line){x, is query});
        return l.m * x + l.b;
};
4.10 DSU (with rollback)
Description: Performs every operation a regular DSU does, but you can roll back to a
specific time.
Usage: int t = uf.time(); ...; uf.rollback(t); T
Time: O(\log(N))
struct RollbackUF {
    vi e:
    vector<pii> st;
    RollbackUF(int n) : e(n, -1) {}
    int size(int x) { return -e[find(x)]; }
    int find(int x) { return e[x] < 0 ? x : find(e[x]); }
    int time() { return len(st); }
    void rollback(int t) {
        for (int i = time(); i-- > t;) e[st[i].first] = st[i].second;
```

st.resize(t);

```
bool join(int a, int b) {
    a = find(a), b = find(b);
    if (a == b) return false;
    if (e[a] > e[b]) swap(a, b);
    st.push_back({a, e[a]});
    st.push_back({b, e[b]});
    e[a] += e[b];
    e[b] = a;
    return true;
}
```

## 4.11 DSU / UFDS

**Usage**: You may discomment the commented parts to find online which nodes belong to each set, it makes the  $union\_set$  method cost  $O(log^2)$  instead O(A)

```
struct DSU {
    vector<int> ps, sz;
    // vector<unordered set<int>> sts;
    DSU(int N)
        : ps(N + 1),
          sz(N, 1) /*, sts(N) */
        iota(ps.begin(), ps.end(), 0);
        // for (int i = 0; i < N; i++)
        // sts[i].insert(i):
    int find set(int x) { return ps[x] == x ? x : ps[x] = find set(ps[x]);
    int size(int u) { return sz[find set(u)]; }
    bool same set(int x, int y) { return find set(x) == find set(y); }
    void union set(int x, int y) {
        if (same_set(x, y)) return;
        int px = find set(x);
        int py = find set(y);
        if (sz[px] < sz[py]) swap(px, py);</pre>
        ps[py] = px;
        sz[px] += sz[py];
        // sts[px].merge(sts[py]);
};
```

# 4.12 Lichao Tree (dynamic)

**Description**: Lichao Tree that creates the nodes dynamically, allowing to query and update from range [MAXL, MAXR]

Usage:

- query(x): find the highest point among all lines in the structure
- add(a,b): add a line of form y = ax + b in the structure
- addSegment(a,b,l,r) : add a line segment of form y=ax+b which covers from range [l,r]

Time:  $O(\log N)$ 

```
template <typename T = ll, T MAXL = 0, T MAXR = 1 '000' 000'001>
struct LiChaoTree {
    static const T inf = -numeric_limits<T>::max() / 2;
    bool first_best(T a, T b) { return a > b; }
    T get best(T a, T b) { return first best(a, b) ? a : b; }
    struct line {
        T m, b;
        T operator()(T x) { return m * x + b; }
    struct node {
        line li;
        node *left. *right:
        node(line li = {0, inf}) : li( li), left(nullptr), right(nullptr)
    {}
        \simnode() {
            delete left:
            delete right;
    };
    node *root:
    LiChaoTree(line li = {0, inf}) : root(new node(li)) {}
    ~LiChaoTree() { delete root; }
    T query(T x, node *cur, T l, \bar{T} r) {
        if (cur == nullptr) return inf;
        if (x < l or x > r) return inf;
        T mid = midpoint(l, r);
        T ans = cur->li(x):
        ans = get best(ans, guery(x, cur->left, l, mid));
        ans = get best(ans, guery(x, cur->right, mid + 1, r));
        return ans:
    T query(T x) { return query(x, root, MAXL, MAXR); }
    void add(line li, node *&cur, T l, T r) {
        if (cur == nullptr) {
            cur = new node(li);
            return:
        T mid = midpoint(l, r);
        if (first best(li(mid), cur->li(mid))) swap(li, cur->li);
        if (first best(li(l), cur->li(l))) add(li, cur->left, l, mid);
        if (first best(li(r), cur->li(r))) add(li, cur->right, mid + 1, r)
    void add(T m, T b) { add({m, b}, root, MAXL, MAXR); }
    void addSegment(line li, node *&cur, T l, T r, T lseg, T rseg) {
        if (r < lseg || l > rseg) return;
        if (cur == nullptr) cur = new node;
        if (lseg \leftarrow 1 && r \leftarrow rseg) {
            add(li, cur, l, r);
            return;
        T mid = midpoint(l, r):
        if (l != r) {
            addSegment(li, cur->left, l, mid, lseg, rseg);
```

```
addSegment(li, cur->right, mid + 1, r, lseg, rseg);
}

void addSegment(T a, T b, T l, T r) {
    addSegment({a, b}, root, MAXL, MAXR, l, r);
};
```

#### 4.13 Merge sort tree

**Description**: Like a segment tree but each node stores the ordered subsegment it represents.

Usage:

• inrange(l, r, a, b): counts the number of positions  $i, l \le i \le r$  such that  $a \le x_i \le b$ . **Time**: Build  $O(N \log N^2)$ , inrange  $O(\log N^2)$ 

**Memory**:  $O(n \log N)$ 

```
template <class T>
struct MergeSortTree {
    int n:
    vector<vector<T>> st;
    MergeSortTree(vector<T> &xs) : n(len(xs)), st(n << 1) {</pre>
        rep(i, 0, n) st[i + n] = vector<T>({xs[i]});
        rrep(i, n - 1, 0) {
            st[i].resize(len(st[i << 1]) + len(st[i << 1 | 1]));
            merge(all(st[i << 1]), all(st[i << 1 | 1]), st[i].begin());
    int count(int i, T a, T b) {
        return upper bound(all(st[i]), b) - lower bound(all(st[i]), a);
    int inrange(int l, int r, T a, T b) {
        int ans = 0:
        for (l += n, r += n + 1; l < r; l >>= 1, r >>= 1) {
            if (l & 1) ans += count(l++, a, b);
            if (r \& 1) ans += count(--r, a, b);
        return ans;
};
```

# 4.14 Mex with update

**Description**: This DS allows you to mantain an array of elments, insert, and remove, and query the MEX at any time.

Ûsage:

- Mex(mxsz): Initialize the DS, mxsz must be the maximum number of elements that the structure may have.
- add(x): just adds one copy of x.
- rmv(x): just remove a copy of x.
- operator(): returns the MEX.

Time:

• Mex(mxsz):  $O(\log mxsz)$ 

- add(x):  $O(\log mxsz)$
- rmv(x):  $O(\log mxsz)$
- *operator()*: *O*(1)

```
struct Mex {
    int mx_sz;
    vi hs:
    set<int> st;
    Mex(int mx sz) : mx sz(mx sz), hs(mx sz + 1) {
        auto it = st.begin();
        rep(i, 0, mx sz + 1) it = st.insert(it, i);
    void add(int x) {
        if (x > mx sz) return;
        if (!hs[x]++) st.erase(x);
    void rmv(int x) {
        if (x > mx sz) return;
        if (!--hs[x]) st.emplace(x);
    int operator()() const { return *st.begin(); }
    /*
      Optional, you can just create with size
      len(xs) add N elements :D
    Mex(const vi &xs, int mx sz = -1) : Mex(\sim mx sz ? mx sz : len(xs)) {
        for (auto xi : xs) add(xi):
};
```

# 4.15 Orderd Set (GNU PBDS)

Usage: If you need an ordered multi set you may add an id to each value. Using greater equal, or less equal is considered undefined behavior.

- order\_of\_key (k): Number of items strictly smaller/greater than k.
- find\_by\_order(k): K-th element in a set (counting from zero).

**Time**: Both  $O(\log N)$ 

Warning: Is 2 or 3 times slower then a regular set/map.

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
template <typename T>
using ordered_set =
    tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
```

#### 4.16 Prefix Sum 2D

**Description**: Given an 2D array with N lines and M columns, find the sum of the subarray that have the left upper corner at (x1, y1) and right bottom corner at (x2, y2). **Time**: Build  $O(N \cdot M)$ , Query O(1).

```
template <typename T>
struct psum2d {
   vector<vector<T>> s;
   vector<vector<T>> psum;
   psum2d(vector<vector<T>> &grid, int n, int m)
        s(n + 1, vector < T > (m + 1)), psum(n + 1, vector < T > (m + 1)) {
        for (int i = 1; i \le n; i++)
            for (int j = 1; j <= m; j++) {
                s[i][i] = s[i][i-1] + qrid[i-1][i-1];
                psum[i][j] = psum[i - 1][j] + s[i][j];
   T query(int x1, int y1, int x2, int y2) \{
        T ans = psum[x2 + 1][y2 + 1] + psum[x1][y1];
        ans -= psum[x2 + 1][y1] + psum[x1][y2 + 1];
        return ans:
   }
};
```

# 4.17 Segment Tree Update Range Query (bottom-up)

```
* @Description:
        merge should be function<T(T,T)>, that
        makes the necessary operation between two
        nodes in the segment tree
 * */
#include "../../Contest/template.cpp"
template <typename T, T identity, auto merge>
struct SegTreeBottomUp {
   int size;
   vector<T> arr;
   SegTreeBottomUp(int n) {
        for (size = 1; size < n; size <<= 1);</pre>
        arr.resize(size << 1);
   void assign(int pos, const T &val) {
        for (arr[pos += size] = val; pos >>= 1;)
            arr[pos] = merge(arr[pos << 1], arr[pos << 1 | 1]);</pre>
   T query(int l, int r) {
       T ans l = identity, ans r = identity;
        for (l += size, r += size + 1; l < r; l >>= 1, r >>= 1) {
            if (l & 1) ans l = merge(ans l, arr[l++]);
            if (r \& 1) ans r = merge(arr[--r], ans r);
        return merge(ans_l, ans_r);
```

```
}
SegTreeBottomUp(const vector<T> &vec) : SegTreeBottomUp(len(vec)) {
    copy(all(vec), begin(arr) + size);
    rrep(i, size - 1, 0) arr[i] = merge(arr[i << 1], arr[i << 1 | 1]);
};
using SegTreeBottomUpSumQuery =
    SegTreeBottomUp<ll, Oll, [](ll a, ll b) { return a + b; }>;

4.18 Sparse table

template <typename T = ll,
</pre>
```

```
vi logs;
vector<vector<T>> st;
public:
   SparseTable(const vector<T> &v) : sz(len(v)), logs(sz + 1) {
      rep(i, 2, sz + 1) logs[i] = logs[i >> 1] + 1;
      st.resize(logs[sz] + 1, vector<T>(sz));
      rep(i, 0, sz) st[0][i] = v[i];
      for (int k = 1; (1 << k) <= sz; k++) {</pre>
```

for (int i = 0;  $i + (1 << k) <= sz; i++) {$ 

cmp(st[k-1][i], st[k-1][i+(1 << (k-1))], st[k][i])

```
}
}
T query(int l, int r) {
    r++;
    const int k = logs[r - l];
    T ret;
    cmp(st[k][l], st[k][r - (1 << k)], ret);
    return ret;
}</pre>
```

# 4.19 Static range queries

int sz:

};

```
}
T operator()(int l, int r) {
    T lv = (l ? acc[l - 1] : T());
    T ret;
    invop(acc[r], lv, ret);
    return ret;
}
```

#### 4.20 Venice Set

**Description**: A container that you can insert q copies of element e, increment every element in the container in x, query which is the best element and it's quantity and also remove k copies of the greatest element.

Time:

- add elment O(log N)
  remove O(log N)
  update: O(1)
- query O(1)

```
template <typename T = ll>
struct VeniceSet {
    using T2 = pair<T, ll>;
    priority queue<T2, vector<T2>, greater<T2>> pg;
    VeniceSet() : acc() {}
    void add element(const T &e, const ll q) { pq.emplace(e - acc, q); }
    void update all(const T &x) { acc += x; }
    T2 best() {
        auto ret = pg.top();
        ret.first += acc:
        return ret;
   void pop() { pq.pop(); }
    void pop k(int k) {
        auto [e, q] = pq.top();
        pq.pop();
        q -= k;
        if (q) pq.emplace(e, q);
};
```

# 4.21 Venice Set (complete)

**Description**: A container which you can insert elements update all at once and also make a few queries

Usage:

- $add\_element(e, q)$ : adds q copies of e, if no q is provided adds a single one
- $update\_all(x)$ : increment every value by x
- erase(e): removes every copy of e, and returns how much was removed.
- count(e): returns the number of e in the container

- high()/low(): returns the hightest/lowest element, and it's quantity
- $pop\_low(q)/pop\_high(q)$ : removes q copies of the lowest/highest elements if no q is provided removes all copies of the lowest/highest element.

You may answer which is the K-th value and it's quantity using an  $ordered\_set$ . Probably works with other operations

**Time:** Considering N the number of distinct numbers in the container

- $add\ element(e,q):\ O(\log(N))$
- $update \ all(x):O(1)$
- erase(e):  $O(\log(N))$
- count(e):  $O(\log(N))$
- high()/low(): O(1)
- $pop\_low(q)/pop\_high(q)$ : worst case is  $O(N \cdot \log{(N)})$  if you remove all elements and so on...

Warning: There is no error handling if you try to *pop* more elements than exists or related stuff

```
struct VeniceSet {
    set<pll> st;
    ll acc;
    VeniceSet() : acc() {}
    ll add element(ll e, ll q = 1) {
        q = erase(e);
        ė −= acc;
        st.emplace(e, q);
        return q;
    void update all(ll x) { acc += x; }
    ll erase(ll e) {
        e = acc;
        auto it = st.lb({e, LLONG MIN});
        if (it == end(st) | | (*it).first != e) return 0;
        ll ret = (*it).second;
        st.erase(it);
        return ret;
    ll count(ll x) {
        x \rightarrow acc;
        auto it = st.lb({x, LLONG MIN});
        if (it == end(st) || (*it).first != x) return 0;
        return (*it).second;
    pll high() { return *rbegin(st); }
    pll low() { return *begin(st); }
    void pop high(ll q = -1) {
        if (\overline{q} == -1) q = high().second;
        while (a) {
            auto [e, eq] = high();
            st.erase(prev(end(st)));
            if (eq > q) add_element(e, eq - q);
            q = max(0ll, q - eq);
        }
    void pop_low(ll q = -1) {
```

```
if (q == -1) q = low().second;
while (q) {
    auto [e, eq] = low();
    st.erase(st.begin());
    if (eq > q) add_element(e, eq - q);
    q = max(0ll, q - eq);
}
};
```

#### 4.22 Wavelet tree

```
using ll = long long;
template <typename T>
struct WaveletTree {
   struct Node {
       T lo, hi;
       int left child, right child;
       vector<int> pcnt;
       vector<ll> psum;
       Node(int lo_, int hi_)
            : lo(lo), hi(hi), left child(0), right child(0), pcnt(),
   psum() {}
   };
   vector<Node> nodes;
   WaveletTree(vector<T> v) {
       nodes.reserve(2 * v.size());
       auto [mn, mx] = minmax element(v.begin(), v.end());
        auto build = [&](auto &&self, Node &node, auto from, auto to) {
            if (node.lo == node.hi or from >= to) return;
            auto mid = midpoint(node.lo, node.hi);
            auto f = [&mid](T x) { return x <= mid; };</pre>
            node.pcnt.reserve(to - from + 1);
            node.pcnt.push back(0);
            node.psum.reserve(to - from + 1);
            node.psum.push_back(0);
            T left upper = node.lo, right lower = node.hi;
            for (auto it = from; it != to; it++) {
                auto value = f(*it);
                node.pcnt.push_back(node.pcnt.back() + value);
                node.psum.push back(node.psum.back() + *it);
                if (value)
                    left upper = max(left upper, *it);
                    right lower = min(right lower, *it);
            }
            auto pivot = stable partition(from, to, f);
            node.left child = make node(node.lo, left upper);
            self(self, nodes[node.left_child], from, pivot);
            node.right child = make node(right lower, node.hi);
            self(self, nodes[node.right_child], pivot, to);
        build(build, nodes[make_node(*mn, *mx)], v.begin(), v.end());
```

```
T kth element(int L, int R, int K) const {
        auto f = [&](auto &&self, const Node &node, int l, int r, int k)
   -> T {
            if (l > r) return 0;
            if (node.lo == node.hi) return node.lo;
            int lb = node.pcnt[l], rb = node.pcnt[r + 1], left size = rb -
    lb:
            return (left size > k
                         ? self(self, nodes[node.left child], lb, rb - 1, k
                         : self(self, nodes[node.right_child], l - lb, r -
    rb.
                                k - left size));
        return f(f, nodes[0], L, R, K);
    pair<int, ll> count and sum in range(int L, int R, T a, T b) const {
        auto f = [&](auto &&self, const Node &node, int l,
                     int r) -> pair<int, ll> {
            if (l > r \text{ or node.} lo > b \text{ or node.} hi < a) return {0, 0};
            if (a <= node.lo and node.hi <= b)</pre>
                return \{r - l + 1,
                         (node.lo == node.hi ? (r - l + 1ll) * node.lo
                                             : node.psum[r + 1] - node.psum
    [l])};
            int lb = node.pcnt[l], rb = node.pcnt[r + 1];
            auto [left cnt, left sum] =
                self(self, nodes[node.left child], lb, rb - 1);
            auto [right_cnt, right sum] =
                self(self, nodes[node.right_child], l - lb, r - rb);
            return {left cnt + right cnt, left sum + right sum};
        return f(f, nodes[0], L, R);
    inline int count in range(int L, int R, T a, T b) const {
        return count_and_sum_in_range(L, R, a, b).first;
    inline ll sum in range(int L, int R, T a, T b) const {
        return count and sum in range(L, R, a, b).second;
   private:
    int make_node(T lo, T hi) {
        int id = (int)nodes.size();
        nodes.emplace back(lo, hi);
        return id;
};
```

# 5 Dynamic Programming

## 5.1 Binary Knapsack (bottom up)

**Description**: Given the points each element have, and it repespective cost, computes the maximum points we can get if we can ignore/choose an element, in such way that the sum of costs don't exceed the maximum cost allowed.

Time: O(N\*W)

Warning: The vectors VS and WS starts at one, so it need an empty value at index 0.

```
const int MAXN(1'000), MAXCOST(1'000 * 20);
ll dp[MAXN + 1][MAXCOST + 1];
bool ps[MAXN + 1][MAXCOST + 1];
pair<ll, vi> knapsack(const vll &points, const vi &costs, int maxCost) {
   int n = len(points) - 1; // ELEMENTS START AT INDEX 1 !
    for (int m = 0; m <= maxCost; m++) {</pre>
        dp[0][m] = 0;
   for (int i = 1; i <= n; i++) {
        dp[i][0] = dp[i - 1][0] + (costs[i] == 0) * points[i];
        ps[i][0] = costs[i] == 0;
   for (int i = 1; i <= n; i++) {
        for (int m = 1; m <= maxCost; m++) {</pre>
            dp[i][m] = dp[i - 1][m], ps[i][m] = 0;
            int w = costs[i]:
            ll v = points[i];
            if (w \le m \text{ and } dp[i - 1][m - w] + v > dp[i][m]) {
                dp[i][m] = dp[i - 1][m - w] + v, ps[i][m] = 1;
        }
   }
   vi is;
   for (int i = n, m = maxCost; i >= 1; --i) {
        if (ps[i][m]) {
            is.emplace back(i);
            m -= costs[i];
    return {dp[n][maxCost], is};
```

### 5.2 Edit Distance

```
Time: O(N * M)
#include "../Contest/template.cpp"

ll edit_distance(const string &a, const string &b) {
   int n = a.size();
   int m = b.size();
   vll2d dp(n + 1, vi(m + 1, 0));
   const ll ADD = 1, DEL = 1, CHG = 1;
   for (int i = 0; i <= n; ++i) {</pre>
```

```
dp[i][0] = i * DEL;
}
for (int i = 1; i <= m; ++i) {
    dp[0][i] = ADD * i;
}

for (int i = 1; i <= n; ++i) {
    for (int j = 1; j <= m; ++j) {
        int add = dp[i][j - 1] + ADD;
        int del = dp[i - 1][j] + DEL;
        int chg = dp[i - 1][j - 1] + (a[i - 1] != b[j - 1]) * CHG;
        dp[i][j] = min({add, del, chg});
    }
}
return dp[n][m];</pre>
```

# 5.3 Knapsack

**Description**: Finds the maximum score you can achieve, given that you have N items, each item has a cost, a point and a quantity, you can spent at most maxcost and buy each item the maximum quantity it has.

Time:  $O(n \cdot maxcost \cdot \log maxqtd)$ 

**Memory**: O(maxcost).

```
ll knapsack(const vi &weight, const vll &value, const vi &qtd, int maxCost
    vi costs;
    vll values;
    for (int i = 0; i < len(weight); i++) {
         ll a = atd[i];
         for (ll x = 1; x \le q; q = x, x \le 1) {
             costs.eb(x * weight[i]);
             values.eb(x * value[i]);
        if (q) {
             costs.eb(q * weight[i]);
             values.eb(q * value[i]);
    vll dp(maxCost + 1):
    for (int i = 0; i < len(values); i++) {
   for (int j = maxCost; j > 0; j--) {
             if (i \ge costs[i]) dp[i] = max(dp[i], values[i] + dp[i - costs])
    [i]]);
    return dp[maxCost];
```

# 5.4 Longest Increasing Subsequence

**Description:** Find the pair (sz, psx) where sz is the size of the longest subsequence and psx is a vector where  $psx_i$  tells the size of the longest increase subsequence that ends at

position i.  $get_i dx$  just tells which indices could be in the longest increasing subsequence. **Time**:  $O(n \log n)$ 

```
#include "../Contest/template.cpp"
template <tvpename T>
pair<int, vi> lis(const vector<T> &xs, int n) {
   vector<T> dp(n + 1, numeric limits<T>::max());
   dp[0] = numeric limits<T>::min();
   int sz = 0:
   vi psx(n);
   rep(i, 0, n) {
       int pos = lower bound(all(dp), xs[i]) - dp.begin();
       sz = max(sz, pos);
       dp[pos] = xs[i];
       psx[i] = pos;
    return {sz, psx};
template <typename T>
vi get idx(vector<T> xs) {
   int n = xs.size();
   auto [sz1, psx1] = lis(xs, n);
   transform(rall(xs), xs.begin(), [](T x) { return -x; });
   auto [sz2, psx2] = lis(xs, n);
   vi ans;
   rep(i, 0, n) {
       int l = psx1[i];
       int r = psx2[n - i - 1];
       if (l + r - 1 == sz1) ans.eb(i);
    return ans;
```

# 5.5 Monery sum

**Description**: Find every possible sum using the given values only once.

```
}
}
for (int i = 0; i < _n; ++i)
    for (int j = 0; j <= _m; ++j)
        if (_dp[i][j]) _ans.insert(j);
return _ans;
}</pre>
```

#### 5.6 Steiner tree

```
template <typename T>
T steinerCost(const vector<vector<T>> &adj, const vi ks,
              T inf = numeric limits<T>::max()) {
    int k = len(ks), n = len(\overline{adj});
    vector<vector<T>> dp(n, vector<T>(1 << k, inf));</pre>
    vi inks(n):
    trav(ki, ks) inks[ki] = 1;
    trav(ki, ks) {
        rep(j, 0, n) {
            if (count(all(ks), j) == 0) {
                dp[j][1 << ki] = adj[ki][j];
        }
    rep(mask, 2, (1 << k)) {
        rep(i, 0, n) {
            if (inks[i]) continue;
            for (int mask2 = (mask - 1) & mask; mask2 >= 1;
                 mask2 = (mask2 - 1) \& mask) {
                int mask3 = mask ^ mask2;
                chmin(dp[i][mask], dp[i][mask2] + dp[i][mask3]);
            rep(j, 0, n) {
                if (inks[j]) continue;
                chmin(dp[i][mask], dp[i][mask] + adi[i][i]);
        }
    T ans = inf:
    rep(i, 0, n) chmin(ans, dp[i][(1 << k) - 1]);
    return ans;
```

#### 5.7 Sum of Subsets

**Description**: Allows you to find if some mask X is a super mask of any of the given masks **Usage**: Call *build* with the masks then it returns a vector of bool V where  $V_X$  says if X is a super mask of any of the initial maks

You can change it to count how many submasks of each mask exsists, by changing the bitwise or by a plus sign...

Time:  $O(LOG \cdot 2^{LOG})$ Memory:  $O(LOG^2 \cdot 2^{LOG})$  Warning: Remember to set LOG with the highest bit possible

```
const int LOG = 20;
vc build(const vi &masks) {
    vc ret(1 << LOG);
    trav(mi, masks) ret[mi] = 1;
    rep(b, 0, LOG) {
        rep(mask, 0, (1 << LOG)) {
            if (mask & (1 << b)) ret[mask] |= ret[mask ^ (1 << b)];
        }
    }
    return ret;
}</pre>
```

## 5.8 Travelling Salesman Problem

```
Time: O(N^2 \cdot 2^N)

Memory: O(N^2 \cdot 2^N)

vll2d dist;

vll memo;

int tsp(int i, int mask, int N) {

   if (mask == (1 << N) - 1) return dist[i][0];

   if (memo[i][mask] != -1) return memo[i][mask];

   int ans = INT_MAX << 1;

   for (int j = \overline{0}; j < N; ++j) {

      if (mask & (1 << j)) continue;

      auto t = tsp(j, mask | (1 << j), N) + dist[i][j];

      ans = min(ans, t);
   }

   return memo[i][mask] = ans;
}
```

## 6 Extras

## 6.1 Binary to gray

```
string binToGray(string bin) {
    string gray(bin.size(), '0');
    int n = bin.size() - 1;
    gray[0] = bin[0];
    for (int i = 1; i <= n; i++) {
        gray[i] = '0' + (bin[i - 1] == '1') ^ (bin[i] == '1');
    }
    return gray;
}</pre>
```

# 6.2 Get permutation cycles

**Description**: Receives a permutation [0, n-1] and return a vector 2D with each cycle.

```
vll2d getPermutationCicles(const vll &ps) {
    ll n = len(ps);
    vector<char> visited(n);
    vector<vll> cicles;
    rep(i, 0, n) {
        if (visited[i]) continue;
        vll cicle;
        ll pos = i;
        while (!visited[pos]) {
            cicle.pb(pos);
            visited[pos] = true;
            pos = ps[pos];
        }
        cicles.push_back(vll(all(cicle)));
    }
    return cicles;
}
```

#### 6.3 Max & Min Check

**Description**: Returns the min/max value in range [l, r] that satisfies the lambda function check, if there is no such value the 'nullopt' is returned.

Usage: check must be a function that receives an integer and return a boolean.

Time:  $O(\log r - l + 1)$ 

```
template <typename T>
optional<T> maxCheck(T l, T r, auto check) {
    optional<T> ret;
    while (l <= r) {
        T m = midpoint(l, r);
        if (check(m))
            ret ? chmax(ret, m) : ret = m, l = m + 1;
        else
            r = m - 1:
    return ret;
template <typename T>
optional<T> minCheck(T l, T r, auto check) {
    optional<T> ret;
    while (l <= r) {
        T m = midpoint(l, r);
        if (check(m))
            ret ? chmin(ret, m) : ret = m, r = m - 1;
        else
            l = m + 1;
    return ret;
```

## 6.4 Merge Interals

Time:  $(N \log N)$ 

```
Warning: It destroys the original array
#include "../Contest/template.cpp"
template <typename T>
vector<pair<T, T>> merge intervals(vector<pair<T, T>> &intervals) {
    if (!len(intervals)) return {};
    using Pt = pair<T, T>;
    sort(all(intervals)):
    vector<Pt> ret{intervals.front()};
    rep(i, 1, len(ret)) {
        auto &[pl, pr] = ret.back();
        auto &[l, r] = intervals[i];
        if (l <= pr)
            chmax(pr, r);
        else
            ret.eb(l, r);
    return ret;
```

# 6.5 Mo's algorithm

```
template <typename T, typename Tans>
struct Mo {
    struct Query {
        int l, r, idx, block;
        Ouery(int l, int _r, int _idx, int _block)
            : l(l), r(r), idx(idx), block(block) {}
        bool operator<(const Query &g) const {</pre>
            if (block != q.block) return block < q.block;</pre>
            return (block & 1 ? (r < q.r) : (r > q.r));
   };
   vector<T> vs:
   vector<Query> qs;
   const int block size;
   Mo(const vector<T> &a) : vs(a), block size((int)ceil(sqrt(a.size())))
   {}
   void add query(int l, int r) {
        qs.emplace_back(l, r, qs.size(), l / block_size);
   auto solve() {
        // get answer return type
        vector<Tans> answers(qs.size());
        sort(all(qs));
        int cur l = 0, cur r = -1;
        for (auto q : qs) {
            while (cur_l > q.l) add(--cur_l);
```

```
while (cur_r < q.r) add(++cur_r);
    while (cur_l < q.l) remove(cur_l++);
    while (cur_r > q.r) remove(cur_r--);
        answers[q.idx] = get_answer();
    }
    return answers;
}

private:
    // add value at idx from data structure
    inline void add(int idx) {}
    // remove value at idx from data structure
    inline void remove(int idx) {}
    // extract current answer of the data structure
    inline Tans get_answer() {}
};
```

#### 6.6 int128t stream

```
void print( int128 x) {
    if (x < 0) {
        cout << '-':
        X = -X;
    if (x > 9) print(x / 10);
    cout << (char)((x % 10) + '0');
int128 read() {
    string s;
    cin >> s:
    int128 x = 0;
    for (auto c : s) {
        if (c != '-') x += c - '0';
        x *= 10;
    x /= 10:
    if (s[0] == '-') x = -x;
    return x;
```

# 7 Geometry

# 7.1 All i know about 2D stuff

```
enum PointPosition { IN, ON, OUT };
template <class Point>
vector<Point> segInter(Point a, Point b, Point c, Point d);
template <typename T>
bool equals(T a, T b) {
    if (std::is floating_point<T>::value)
        return \overline{f}abs(a -\overline{b}) < EPS;
    else
        return a == b:
template <class T>
struct Point {
    typedef Point P;
    T x, y;
    explicit Point(T x = 0, T y = 0) : x(x), y(y) {}
    bool operator<(P p) { return tie(x, y) < tie(p.x, p.y); }</pre>
    bool operator>(P& rhs) { return rhs < *this; }</pre>
    bool operator==(P p) { return tie(x, y) == tie(p.x, p.y); }
    P operator+(P p) { return P(x + p.x, y + p.y); }
    P operator-(P p) { return P(x - p.x, y - p.y); }
    P operator*(T d) { return P(x * d, y * d); }
    P operator/(T d) { return P(x / d, y / d); }
    T dot(P p) \{ return x * p.x + y * p.y; \}
    T cross(P p) { return x * p.y - y * p.x; }
    T cross(P a, P b) { return (a - *this).cross(b - *this); }
    T dist2() { return x * x + y * y; }
    double dist() { return sqrt((double)dist2()); }
    // angle to x-axis in interval [-pi, pi]
    double angle() { return atan2(y, x); }
    P unit() { return *this / dist(); } // makes dist()=1
    P perp() \{ return P(-y, x); \}
                                         // rotates +90 degrees
    P normal() { return perp().unit(); }
    // returns point rotated 'a' radians ccw around
    // the origin
    P rotate(double a) {
        return P(x * cos(a) - y * sin(a), x * sin(a) + y * cos(a));
    pair<T, T> slope(Point<T>& o) {
        auto a = o.x - x;
        auto b = o.y - y;
        if (!is floating point<T>::value) {
            auto g = gcd(a, b);
            if (a) a /= a, b /= a:
        return {b, a};
    friend ostream& operator<<(ostream& os, P p) {</pre>
        return os << "(" << p,x << "," << p,v << ")";
    double distanceTo(Point<T>& other) {
        return hypot(other.x - x, other.y - y);
};
```

```
template <typename T>
struct Line {
    T a, b, c;
    Point<T> p1, p2;
    Line(T a = 0, T b = 0, T c = 0) : a(a), b(b), c(c) {
        if (a != 0) {
            double x = 0;
            double y = (-c) / b;
            p1 = Point < T > (x, v):
        if (b != 0) {
            double y = 0;
            double x = (-c) / a;
            p2 = Point < T > (x, y);
    Line(Point<T>& p, Point<T>& q) {
        a = p.y - q.y;
        b = q.x - p.x;
        c = p.cross(q);
        p1 = p, p2 = q;
    bool operator==(Line<T>& other) {
        return tie(a, b, c) == tie(other.a, other.b, other.c);
    // Less-than operator
    bool operator<(Line& rhs) {</pre>
        return tie(a, b, c) < tie(rhs.a, rhs.b, rhs.c);</pre>
    bool operator>(Line& rhs) { return rhs < *this; }</pre>
    Line<T> norm() {
        T d = a == 0 ? b : a;
        return Line(a / d, b / d, c / d);
    bool contains(PointT>& p) { return equals(a * p.x + b * p.y + c, (T)
    0); }
    bool parallel(Line<T>& r) {
        auto det = a * r.b - b * r.a;
        return equals(det, 0) and !(*this == r);
    bool orthogonal(Line<T>& r) { return equals(a * r.a + b * r.b, 0); }
    T direction(Point<T>& p3) { return p1.cross(p2, p3); }
    friend ostream& operator<<(ostream& os. Line 1) {</pre>
        return os << fixed << setprecision(6) << "(" << l.a << "," << l.b</pre>
    << ","
                  << l.c << ")":
    double distance(Point<T>& p) {
        return (a * p.x + b * p.y + c) / hypot(a, b);
    Point<T> closest(Point<T>& p) {
        auto den = (a * a + b * b):
```

```
auto x = (b * (b * p.x - a * p.y) - a * c) / den;
                                                                                    Point<T> c:
        auto v = (a * (-b * p.x + a * p.v) - b * c) / den:
                                                                                    Tr;
                                                                                    Circle(Point<T> c, T r) : c( c), r( r) {}
        return Point<T>{x, y};
                                                                                    Circle(T r) : Circle(\overline{Point} < T > (0, 0), r) {}
                                                                                    ld area() const { return PI * r * r; }
};
                                                                                    ld perimeter() const { return 2.0 * PI * r; }
                                                                                    ld arc(ld theta) const { return theta * r; }
   */
                                                                                    ld chord(ld theta) const { return 2.0 * r * sin(theta / 2.0); }
                                                                                    ld sector(ld theta) const { return (theta * r * r) / 2.0; }
template <typename T>
                                                                                    ld segment(ld theta) const { return ((theta - sin(theta)) * r * r) /
struct LineSeament {
    Point<T> p1, p2:
                                                                                    2.0:  }
                                                                                    PointPosition position(const Point<T>& p) const {
   LineSegment(Point<T> p, Point<T> q) { p1 = p, p2 = q; }
                                                                                        auto d = c.dist(p);
    LineSegment(T a, T b, T c, T d)
                                                                                        return equals(d, r) ? ON : (d < r ? IN : OUT);</pre>
        : LineSegment(Point<T>(a, b), Point<T>(c, d)) {}
    bool operator==(LineSegment<T>& other) {
                                                                                };
        return tie(p1, p2) == tie(other.p1, other.p2);
                                                                                /*
    // Less-than operator
    bool operator<(LineSegment& rhs) {</pre>
                                                                                template <typename T>
        return tie(p1, p2) < tie(rhs.p1, rhs.p2);</pre>
                                                                                struct Rectangle {
                                                                                    Point<T> P, 0;
    bool operator>(LineSegment& rhs) { return rhs < *this; }</pre>
                                                                                    T b, h;
   T direction(Point<T>& p3) { return p1.cross(p2, p3); }
                                                                                    Rectangle(const Point<T>& p. const Point<T>& a) : P(P), O(a) {
    friend ostream& operator<<(ostream& os, LineSegment l) {</pre>
                                                                                        assert(P != 0):
        return os << "(" << l.p1 << "," << l.p2 << ")";
                                                                                        b = \max(P.x, Q.x) - \min(P.x, Q.x);
                                                                                        h = max(P.y, Q.y) - min(P.y, Q.y);
    vector<Point<T>> intersection(LineSegment<T>& other) {
        return seaInter(pl. p2, other.p1, other.p2):
                                                                                    Rectangle(T base, T height)
                                                                                        : P(0, 0), Q(base, height), b(base), h(height) {}
                                                                                    T perimeter() const { return 2 * b + 2 * h; }
    // Verifica se o ponto P da reta r que écontm A e B pertence ao
                                                                                    T area() const { return b * h; }
   seamento
                                                                                    optional<Rectangle> intersection(const Rectangle& r) const {
    bool contains(Point<T>& P) {
                                                                                        using pt = pair<T, T>;
        return equals(p1.x, p2.x)
                   ? \min(p1.y, p2.y) \le P.y \text{ and } P.y \le \max(p1.y, p2.y)
                                                                                        auto i = pt(min(P.x, Q.x), max(P.x, Q.x));
                   : min(p1.x, p2.x) \le P.x \text{ and } P.x \le max(p1.x, p2.x);
                                                                                        auto u = pt(min(r.P.x, r.Q.x), max(r.P.x, r.Q.x));
                                                                                        auto a = max(i.first, u.first);
                                                                                        auto b = min(i.second, u.second);
    // Ponto mais óprximo de P no segmento AB
                                                                                        i = pt(min(P.y, Q.y), max(P.y, Q.y));
    Point<T> closest(Point<T>& P) {
                                                                                        u = pt(min(r.P.y, r.Q.y), max(r.P.y, r.Q.y));
        Line<T> r(p1, p2);
                                                                                        auto c = max(i.first, u.first);
        auto Q = r.closest(P);
                                                                                        auto d = max(i.second, u.first);
        if (this->contains(0)) return 0;
                                                                                        if (d < c or b < a) return nullopt:
        auto distp1 = P.distanceTo(p1);
                                                                                        return {{a, c}, {b, d}};
        auto distp2 = P.distanceTo(p2);
        if (distp1 <= distp2)</pre>
                                                                                };
            return p1:
        else
            return p2;
                                                                                template <typename T>
};
                                                                                struct Trapezium {
                                                                               ___ T B, b, h;
                                                                                    T area() const { return ((b + B) * h) / 2; }
                                                                                };
template <typename T>
                                                                               /*
struct Circle {
```

```
template <tvpename T>
struct Triangle {
   Point<T> A, B, C;
   enum SidesClass { EQUILATERAL, ISOCELES, SCALENE };
   SidesClass classification by sides() const {
        auto a = A.distanceTo(B);
        auto b = B.distanceTo(C);
        auto c = C.distanceTo(A):
        if (equals(a, b) && equals(b, c)) return EQUILATERAL;
        if (equals(a, b) or equals(a, c) or equals(b, c)) return ISOCELES:
        return SCALENE;
   enum AnglesClass { RIGHT, ACUTE, OBTUSE };
   AnglesClass classification by angles() const {
        auto a = dist(A, B);
        auto b = dist(B, C);
        auto c = dist(C, A):
        auto alpha = acos((a * a - b * b - c * c) / (-2 * b * c));
        auto beta = acos((b * b - a * a - c * c) / (-2 * a * c));
        auto gamma = acos((c * c - a * a - b * b) / (-2 * a * b));
        auto right = PI / 2.0:
        if (equals(alpha, right) || equals(beta, right) || equals(gamma,
   right))
            return RIGHT:
        if (alpha > right || beta > right || gamma > right) return OBTUSE;
        return ACUTE:
   }
   double perimeter() const {
        auto a = dist(A, B), b = dist(B, C), c = dist(C, A);
        return a + b + c;
   double area() const {
        Line<T> r(A, B);
        auto b = dist(A, B);
        auto h = r.distance(C);
        return (b * h) / 2;
};
template <typename T>
Point<T> triangleBarycenter(const Point<T>& a, const Point<T>& b,
                            const Point<T>& c) {
    return Point<T>((a.x + b.x + c.x) / 3.0, (a.y + b.y + c.y) / 3.0);
template <typename T>
Point<T> triangleOrthocenter(const Point<T>& a, const Point<T>& b,
                             const Point<T>& c) {
   LineT> r(a, b), s(a, c):
   Line<T> u\{r.b, -r.a, -(c.x * r.b - c.y * r.a)\};
   Line<T> v\{s.b, -s.a, -(b.x * s.b - b.y * s.a)\};
   auto det = u.a * v.b - u.b * v.a:
```

```
=== auto x = (-u.c * v.b + v.c * u.b) / det;
    auto y = (-v.c * u.a + u.c * v.a) / det;
    return {x, y};
template <typename T>
Point<double> triangleIncenter(const Point<T>& a, const Point<T>& b,
                               const Point<T>& c) {
    auto dab = distance(a, b):
    auto dbc = distance(b, c);
    auto dca = distance(c, a);
    auto p = dab + dbc + dca;
    auto \dot{x} = (a.x * dab + b.x * dbc + b.x * dca) / (p);
    auto y = (a.y * dab + b.y * dbc + b.y * dca) / (p);
    return Point<double>(x, v);
template <typename T>
Point<T> triangleCircumcenter(const Point<T>& A, const Point<T>& B,
                              const Point<T>& C) {
    auto D = 2 * (A.x * (B.v - C.v) + B.x * (C.v - A.v) + C.x * (A.v - B.v)
    auto A2 = A.x * A.x + A.v * A.v:
    auto B2 = B.x * B.x + B.y * B.y;
    auto C2 = C.x * C.x + C.y * C.y;
    auto x = (A2 * (B.y - C.y) + B2 * (C.y - A.y) + C2 * (A.y - B.y)) / D;
    auto y = (A2 * (C.x - B.x) + B2 * (A.x - C.x) + C2 * (B.x - A.x)) / D:
    return {x, v};
template <tvpename T>
Point<T> triangleCircumradius(const Point<T>& a, const Point<T>& b,
                              const Point<T>& c) {
    auto dab = distance(a, b);
    auto dbc = distance(b, c);
    auto dca = distance(c, a);
    return (dab + dbc + dca) / triangleArea(a, b, c);
}
template <class Point>
vector<Point> segInter(Point a, Point b, Point c, Point d) {
    auto oa = c.cross(d, a), ob = c.cross(d, b), oc = a.cross(b, c),
         od = a.cross(b, d);
    // Checks if intersection is single non-endpoint
    if (sgn(oa) * sgn(ob) < 0 \&\& sgn(oc) * sgn(od) < 0)
        return \{(a * ob - b * oa) / (ob - oa)\};
    set<Point> s:
    if (onSegment(c, d, a)) s.insert(a);
    if (onSegment(c, d, b)) s.insert(b);
    if (onSegment(a, b, c)) s.insert(c);
    if (onSegment(a, b, d)) s.insert(d);
    return {all(s)};
```

```
*/
                                                                                           st.pop back();
template <tvpename T>
double angle(const Point<T>& P, const Point<T>& Q, const Point<T>& R.
                                                                                       st.push back(pts[i]);
             const Point<T>& S) {
    auto ux = P.x - Q.x;
                                                                                   pts = st;
    auto uy = P.y - Q.y;
    auto vx = R.x - S.x:
    auto vy = R.y - S.y;
    auto num = ux * vx + uv * vv;
                                                                               template <tvpename T>
    auto den = hypot(ux, uy) * hypot(vx, vy);
   // Caso especial: se den == 0, algum dos vetores \acute{e} degenerado: os dois
    // pontos ãso iguais. Neste caso, o ângulo ãno áest definido
    return acos(num / den);
                                                                               template <tvpename T>
                                                                                   Point<T> b = C - A, c = B - A:
struct pt {
    double x, y;
    int id:
                                                                               template <typename T>
};
int orientation(pt a, pt b, pt c) {
    double v = a.x * (b.y - c.y) + b.x * (c.y - a.y) + c.x * (a.y - b.y);
                                                                                   Point<T> o = ps[0];
    if (v < 0) return -1; // clockwise</pre>
                                                                                   double r = 0, EPS = 1 + 1e-8;
   if (v > 0) return +1: // counter-clockwise
    return 0:
                                                                                       o = ps[i], r = 0:
                                                                                           o = (ps[i] + ps[i]) / 2;
bool cw(pt a, pt b, pt c, bool include collinear) {
                                                                                            r = (o - ps[i]).dist();
    int o = orientation(a, b, c);
    return o < 0 || (include collinear && o == 0);
bool collinear(pt a, pt b, pt c) { return orientation(a, b, c) == 0; }
void convex hull(vector<pt>& pts, bool include collinear = false) {
    pt p0 = *min element(all(pts), [](pt a, pt b) {
        return make pair(a.y, a.x) < make pair(b.y, b.x);</pre>
                                                                                   return {o, r};
    sort(all(pts), [&p0](const pt& a, const pt& b) {
        int o = orientation(p0, a, b);
        if (o == 0)
            return (p0.x - a.x) * (p0.x - a.x) + (p0.v - a.v) * (p0.v - a.x)
                                                                               template <tvpename T>
   y) <
                   (p0.x - b.x) * (p0.x - b.x) + (p0.y - b.y) * (p0.y - b.y)
                                                                                   auto a = 2 * (0.x - P.x);
                                                                                   auto b = 2 * (0.v - P.v);
        return o < 0;
                                                                                   return {a, b, c};
   if (include collinear) {
                                                                               }
        int i = len(pts) - 1;
        while (i \ge 0 \& collinear(p0, pts[i], pts.back())) i--;
        reverse(pts.begin() + i + 1, pts.end());
    vector<pt> st;
                                                                               ll polygonArea(vector<pll>& pts) {
    for (int i = 0; i < len(pts); i++) {
                                                                                   ll^{\tilde{a}ts} = 0;
        while (st.size() > 1 \&\&
```

```
!cw(st[len(st) - 2], st.back(), pts[i], include collinear))
double ccRadius(const Point<T>& A, const Point<T>& B, const Point<T>& C) {
    return (B - A).dist() * (C - B).dist() * (A - C).dist() /
          abs((B - A).cross(C - A)) / 2;
Point<T> ccCenter(const Point<T>& A, const Point<T>& B, const Point<T>& C)
    return A + (b * c.dist2() - c * b.dist2()).perp() / b.cross(c) / 2;
pair<Point<T>, double> mec(vector<Point<T>> ps) {
    shuffle(all(ps), mt19937(time(0)));
    rep(i, 0, len(ps)) if ((o - ps[i]).dist() > r * EPS) {
        rep(i, 0, i) if ((o - ps[i]).dist() > r * EPS) {
            rep(k, 0, j) if ((o - ps[k]).dist() > r * EPS) {
                o = ccCenter(ps[i], ps[i], ps[k]);
                r = (o - ps[i]).dist();
Line<T> perpendicular bisector(const Point<T>& P, const Point<T>& Q) {
    auto c = (P.x * P.x + P.y * P.y) - (Q.x * Q.x + Q.y * Q.y);
ll cross(ll x1, ll y1, ll x2, ll y2) { return x1 * y2 - x2 * y1; }
```

```
for (int i = 2; i < len(pts); i++)
        ats += cross(pts[i].first - pts[0].first, pts[i].second - pts[0].
    second,
                     pts[i - 1].first - pts[0].first,
                     pts[i - 1].second - pts[0].second);
    return abs(ats / 211):
ll boundary(vector<pll>& pts) {
    ll ats = pts.size();
    for (int i = 0; i < len(pts); i++) {</pre>
        il deltax = (pts[i].first - pts[(i + 1) % pts.size()].first);
ll deltay = (pts[i].second - pts[(i + 1) % pts.size()].second);
        ats += abs( gcd(deltax, deltay)) - 1;
    return ats;
pll latticePoints(vector<pll>& pts) {
    ll bounds = boundary(pts);
    ll area = polygonArea(pts);
    ll inside = area + 1ll - bounds / 2ll;
    return {inside, bounds};
            _____
template <typename T>
bool contains(const Point<T>& A, const Point<T>& B, const Point<T>& P) {
    // Verifica se P áest na ãregio retangular
    auto xmin = min(A.x, B.x);
    auto xmax = max(A.x, B.x);
    auto ymin = min(A.y, B.y);
    auto ymax = max(A.y, B.y);
    if (P.x < xmin \mid\mid P.x > xmax \mid\mid P.y < ymin \mid\mid P.y > ymax) return false
    // Verifica çãrelao de csemelhana no âtringulo
    return equals((P.y - A.y) * (B.x - A.x), (P.x - A.x) * (B.y - A.y);
}
// the polygon area of a intersection between a circle and a ccw polygon
template <typename T>
#define arg(p, q) atan2(p.cross(q), p.dot(q))
double circlePoly(Point<T> c, double r, vector<Point<T>> ps) {
    auto tri = [&](Point<T> p, Point<T> q) {
        auto r2 = r * r / 2;
        Point<T> d = q - p;
        auto a = d.dot(p) / d.dist2(), b = (p.dist2() - r * r) / d.dist2()
        auto det = a * a - b;
        if (det \le 0) return arg(p, q) * r2;
        auto s = max(0., -a - sqrt(det)), t = min(1., -a + sqrt(det));
        if (t < 0 | | 1 \le s) return arg(p, q) * r2;
        Point<T> u = p + d * s, v = p + d * t;
        return arg(p, u) * r2 + u.cross(v) / 2 + arg(v, q) * r2;
```

#### 7.2 Angle between three points

**Description**: Computes the angle apb in radians **Warning**: a is equal to b then the angle isn't defined.

```
#include "./template.cpp"
template <typename T>
ld angle(const Point<T>& p, const Point<T>& a, const Point<T>& b) {
    auto ux = p.x - a.x;
    auto uy = p.y - a.y;
    auto vx = p.x - b.x;
    auto vy = p.y - b.y;
    auto num = ux * vx + uy * vy;
    auto den = hypot(ux, uy) * hypot(vx, vy);
    return acos(num / den);
}
```

# 7.3 Area of union of rectangles

```
using SegT = ll;
const SeqT eSeq = 1e9;
struct QueryT {
    SeqT q, v;
    QueryT(): q(0), v(eSeq) {}
=== QueryT(SegT _v) : q(1), v(_v) {}
inline QueryT combine(QueryT ln, QueryT rn, pii lr1, pii lr2) {
    QueryT ret;
    if (ln.v < rn.v) ret = ln;</pre>
    if (rn.v < ln.v) ret = rn;</pre>
    if (rn.v == ln.v) {
        ret.v = ln.v;
        ret.q = ln.q + rn.q;
    return ret;
using LazvT = SegT:
inline OueryT applyLazyInOuery(QueryT q, LazyT l, pii lr) {
    if (l == LazyT()) return q;
    if (q.v == eSeq) q.v = 0, q.q = 1;
```

```
q.v += l;
   return q;
inline LazyT applyLazyInLazy(LazyT a, LazyT b) { return a + b; }
using UpdateT = SegT;
inline QueryT applyUpdateInQuery(QueryT q, UpdateT u, pii lr) {
    return applyLazyInQuery(q, u, lr);
inline LazyT applyUpdateInLazy(LazyT l, UpdateT u, pii lr) { return l + u;
template <typename Qt = QueryT, typename Lt = LazyT, typename Ut = UpdateT
          auto C = combine, auto ALQ = applyLazyInQuery,
          auto ALL = applyLazyInLazy, auto AUQ = applyUpdateInQuery,
          auto AUL = applyUpdateInLazy>
struct LazySegmentTree {
   int n, h;
   vector<Qt> ts;
   vector<Lt> ds;
   vector<pii> lrs;
   LazySegmentTree(int n)
        : n(_n),
         h(sizeof(int) * 8 - builtin clz(n)),
          ts(n \ll 1),
          ds(n),
          lrs(n << 1) {
        rep(i, 0, n) lrs[i + n] = {i, i};
        rrep(i, n - 1, 0) {
            lrs[i] = {lrs[i << 1].first, lrs[i << 1 | 1].second};</pre>
   LazySegmentTree(const vector<Qt> &xs) : LazySegmentTree(len(xs)) {
       copy(all(xs), ts.begin() + n);
        rep(i, 0, n) lrs[i + n] = \{i, i\};
        rrep(i, n - 1, 0) {
            ts[i] = C(ts[i << 1], ts[i << 1 | 1], lrs[i << 1], lrs[i << 1
   | 11);
   void set(int p, Qt v) {
       ts[p + n] = v;
       build(p + n);
   void upd(int l, int r, Ut v) {
       l += n. r += n + 1:
       int 10 = 1, r0 = r;
        for (; l < r; l >>= 1, r >>= 1) {
            if (l & 1) apply(l++, v);
            if (r & 1) apply(--r, v);
       build(l0), build(r0 - 1);
   Qt qry(int l, int r) {
```

```
l += n, r += n + 1:
        push(l), push(r - 1);
        Qt resl = Qt(), resr = Qt();
        pii lr1 = \{l, l\}, lr2 = \{r, r\};
        for (; l < r; l >>= 1, r >>= 1) {
            if (l & 1) resl = C(resl, ts[l], lr1, lrs[l]), l++;
            if (r \& 1) r--, resr = C(ts[r], resr, lrs[r], lr2);
        return C(resl, resr, lr1, lr2);
    }
    void build(int p) {
        while (p > 1) {
            p >>= 1;
            ts[p] =
                ALQ(C(ts[p << 1], ts[p << 1 | 1], lrs[p << 1], lrs[p << 1]
    | 1]),
                    ds[p], lrs[p]);
    }
    void push(int p) {
        rrep(s, h, 0) {
            int i = p \gg s;
            if (ds[i] != Lt()) {
                apply(i << 1, ds[i]), apply(i << 1 | 1, ds[i]);
                ds[i] = Lt();
        }
    inline void apply(int p, Ut v) {
        ts[p] = AUQ(ts[p], v, lrs[p]);
        if (p < n) ds[p] = AUL(ds[p], v, lrs[p]);
};
ll areaOfRectanglesUnion(
    const vector<pair<Point<int>, Point<int>>> &rectangles) {
    if (!size(rectangles)) return 0;
    int maxy = INT MIN;
    for (auto &[p1, p2] : rectangles) {
        assert(p1.x < p2.x && p1.y < p2.y);
        maxy = max(\{maxy, p1.y, p2.y\});
    vector<array<int, 4>> sl;
    sl.reserve(size(rectangles) * 2);
    for (auto &[p1, p2] : rectangles) {
        sl.push_back({p1.x, p1.y, p2.y - 1, 1});
        sl.push back(\{p2.x, p1.y, p2.y - 1, -1\});
    sort(sl.begin(), sl.end());
    vector<QueryT> aux(maxy, QueryT(0));
    LazySegmentTree seg(aux);
    // memset(seg vec, 0, sizeof(ll) * maxy);
    // seg::build(maxy, seg_vec);
    int prevx = get<0>(sl.front());
    ll ans = 0;
```

```
for (auto &[curx, ys, yf, inc] : sl) {
    auto [q, v] = seg.qry(0, maxy - 1);
    // auto [q, v] = seg::query(0, maxy - 1);
    ans += (ll)(curx - prevx) * (v ? maxy : maxy - q);
    seg.upd(ys, yf, inc);
    prevx = curx;
}
return ans;
```

### 7.4 Area: polygon

```
#include "./template.cpp"
template <typename T>
ld area(const vector<Point<T>>& pts) {
    ld a = 0.0;
    int n = size(pts);
    for (int i = 0; i < n; i++) {
        a += pts[i].x * pts[(i + 1) % n].y;
        a -= pts[i].y * pts[(i + 1) % n].x;
    }
    return fabs(a) / (ld)2;
}</pre>
```

## 7.5 Check if point belongs to line

## 7.6 Check if point belongs to segment

```
#include "./template.cpp"
template <class P>
bool segmentContainsPoint(const P& p, const P& a, const P& b) {
    auto xmin = min(a.x, b.x);
    auto xmax = max(a.x, b.x);
    auto ymin = min(a.y, b.y);
    auto ymax = max(a.y, b.y);
    if (p.x < xmin or p.x > xmax or p.y < ymin or p.y > ymax) return false;
    return equals((p.y - a.y) * (b.x - a.x), (p.x - a.x) * (b.y - a.y));
}
```

### 7.7 Check if point is inside polygon

**Description**: checks if the point p is inside the polygon with vertices in pts, works for both convex and concave polygons.

```
#pragma once
#include "./Angle between three points.cpp"
#include "./Check if point belongs to segment.cpp"
#include "./Determinant.cpp"
#include "./template.cpp"
template <tvpename T>
bool contains(const vector<Point<T>>& pts, const Point<T>& p) {
    int n = size(pts):
    if (n < 3) return false; // may treat it appart</pre>
    T sum = 0.0:
    for (int i = 0; i < n; i++) {
        auto d = determinant(p, pts[i], pts[(i + 1) % n]);
        auto a = angle(p, pts[i], pts[(i + 1) % n]);
        sum += d > 0 ? a : (d < 0 ? -a : 0):
    return equals(fabs(sum), 2 * PI);
// 0: outside, 1: inside, 2: boundary
template <class P>
int pointInPolygon(const vector<P>& pts, const P& p) {
    if (contains(pts, p)) return 1;
    int n = size(pts):
    for (int i = 0; i < n; i++) {
        if (segmentContainsPoint(p, pts[i], pts[(i + 1) % n])) {
            return 2:
    return 0;
```

#### 7.8 Convex hull

```
#include "../Contest/template.cpp"
#include "./Determinant.cpp"
#include "./template.cpp"

template <typename T>
vector<Point<T>> convexHull(vector<Point<T>> pts) {
    if (len(pts) <= 1) return pts;
        sort(all(pts));
        vector<Point<T>> h(len(pts) + 1);
    int s = 0, t = 0;
    for (int it = 2; it--; s = --t, reverse(all(pts)))
        for (Point<T> p: pts) {
             while (t >= s + 2 && determinant(h[t - 2], h[t - 1], p) <= 0)
        t--;
             h[t++] = p;
        }
}</pre>
```

```
return \{h.begin(), h.begin() + t - (t == 2 \& h[0] == h[1])\};
template <typename T>
vector<Point<T>> convexHull2(vector<Point<T>> pts) {
    int n = len(pts);
    sort(pts.begin(), pts.end());
   vector<Point<T>> l, u;
   for (int i = 0; i < n; i++) {
        while (len(l) >= 2 \&\&
               determinant(l[len(l) - 1], l[len(l) - 2], pts[i]) < 0) {
            l.pop back();
        l.push back(pts[i]);
   for (int i = n - 1; \simi; --i) {
        while (len(u) \geq 2 &&
               determinant(u[len(u) - 1], u[len(u) - 2], pts[i]) < 0)
            u.pop back();
        u.push back(pts[i]);
   u.pop back(), l.pop back();
   u.reserve(len(u) + len(l));
   u.insert(u.end(), all(l));
    return u:
```

### 7.9 Cross product between points

```
#pragma once
#include "./template.cpp"
template <typename T>
T cross(const Point<T>& p, const Point<T>& q) {
    return p.x * q.y - p.y * q.x;
}
```

# 7.10 Define line from two points

```
#pragma once
#include "./template.cpp"
template <typename T>
inline tuple<T, T, T> defineLine(const Point<T>& p, const Point<T>& q) {
    return {p.y - q.y, q.x - p.x, cross(p, q)};
}
```

### 7.11 Determinant

```
#pragma once
#include "./template.cpp"
template <typename T>
```

#### 7.12 Distance: point to point

```
#include "./template.cpp"
template <typename T>
T distance(const Point<T>& p, const Point<T>& q) {
   return hypot(p.x - q.x, p.y - q.y);
}
```

#### 7.13 Halfplane intersection

```
#pragma once
#include "./Point.cpp"
#include "./template.cpp"
// Basic half-plane struct.
struct Halfplane {
    // 'p' is a passing point of the line and 'pg' is the direction vector
    οf
    // the line.
    Point<ld> p, pq;
    long double angle;
    Halfplane() {}
    Halfplane(const Point<ld>& a, const Point<ld>& b) : p(a), pq(b - a) {
        angle = atan2l(pq.y, pq.x);
    // Check if point 'r' is outside this half-plane.
    // Every half-plane allows the region to the LEFT of its line.
    bool out(const Point<ld>& r) { return cross(pg, r - p) < -EPS; }
    // Intersection point of the lines of two half-planes. It is assumed
   they're
    // never parallel.
    friend Point<ld> inter(const Halfplane& s, const Halfplane& t) {
        long double alpha = cross((t.p - s.p), t.pq) / cross(s.pq, t.pq);
        return s.p + (s.pg * alpha);
};
// Actual algorithm
// receive it by reference if don't care messing with it
vector<Point<ld>> hp intersect(vector<Halfplane> H) {
    const ld inf = 2e6:
    Point<ld> box[4] = \{// Bounding box in CCW order
                        Point<ld>(inf, inf), Point<ld>(-inf, inf),
                        Point<ld>(-inf, -inf), Point<ld>(inf, -inf));
    for (int i = 0; i < 4; i++) { // Add bounding box half-planes.
        Halfplane aux(box[i], box[(i + 1) % 4]);
```

```
H.push back(aux);
}
// Sort by angle and start algorithm
sort(H.begin(), H.end(), [&](const Halfplane& a, const Halfplane& b) {
    return a.angle < b.angle:
deque<Halfplane> dg;
int len = 0:
for (int i = 0; i < int(H.size()); i++) {</pre>
    // Remove from the back of the deque while last half-plane is
redundant
    while (len > 1 && H[i].out(inter(dg[len - 1], dg[len - 2]))) {
        dq.pop back();
        --len:
    // Remove from the front of the deque while first half-plane is
    // redundant
    while (len > 1 \&\& H[i].out(inter(dg[0], dg[1])))  {
        dq.pop front();
        --len:
    }
    // Special case check: Parallel half-planes
    if (len > 0 \& fabsl(cross(H[i].pq, dq[len - 1].pq)) < EPS) {
        // Opposite parallel half-planes that ended up checked against
each
        if (dot(H[i].pq, dq[len - 1].pq) < 0.0) return vector<Point<ld
>>();
        // Same direction half-plane: keep only the leftmost half-
plane.
        if (H[i].out(dq[len - 1].p)) {
            dq.pop_back();
            --len:
        } else
            continue:
    // Add new half-plane
    dq.push back(H[i]);
    ++len:
// Final cleanup: Check half-planes at the front against the back and
while (len > 2 \& dq[0].out(inter(dq[len - 1], dq[len - 2]))) {
    dq.pop back();
    --len:
while (len > 2 && dq[len - 1].out(inter(dq[0], dq[1]))) {
    dq.pop front();
    --len:
// Report empty intersection if necessary
if (len < 3) return vector<Point<ld>>();
// Reconstruct the convex polygon from the remaining half-planes.
```

```
vector<Point<ld>> ret(len);
for (int i = 0; i + 1 < len; i++) {
    ret[i] = inter(dq[i], dq[i + 1]);
}
ret.back() = inter(dq[len - 1], dq[0]);
return ret;
}</pre>
```

#### 7.14 Lattice points

```
#pragma once
#include "../Contest/template.cpp"
#include "./Area: polygon.cpp"
#include "./template.cpp"
template <tvpename T>
pair<ll, ll> latticePoints(const vector<Point<T>> &pts) {
    ll bounds = pts.size();
    int n = pts.size();
    for (int i = 0; i < n; i++) {
        ll deltax = (pts[i].x - pts[(i + 1) % n].x);
        ll deltay = (pts[i].y - pts[(i + 1) % n].y);
        bounds += abs( gcd(deltax, deltay)) - 1;
    ll a = area(pts);
    ll inside = a + 1 - bounds / 2ll;
    return {inside, bounds};
}
```

# 7.15 Left of polygon cut

Warning: if some vertex lies exactly on the line A B, theese vertex will be included in teh answer

```
#include "./Determinant.cpp"
#include "./template.cpp"
template <tvpename T>
vector<Point<T>> leftOfPolygonCut(const vector<Point<T>>& vs, const Point<</pre>
   T>& A,
                                   const Point<T>& B) {
    // çãInterseo entre a reta AB e o segmento de reta PQ
    auto intersection = [&](const Point<T>& P, const Point<T>& Q,
                             const Point<T>& A, const Point<T>& B) -> Point
    <T> {
        auto a = B.y - A.y;
        auto b = A.x - B.x;
        auto c = B.x * A.y - A.x * B.y;
        auto u = fabs(a * P.x + b * P.y + c);
        auto v = fabs(a * 0.x + b * 0.v + c):
        // éMdia ponderada pelas âdistncias de P e O éat a reta AB
        return \{(P.x * v + Q.x * u) / (u + v), (P.y * v + Q.y * u) / (u + v)\}
   v)};
    };
    vector<Point<T>> points;
```

```
int n = size(vs);
for (int i = 0; i < n; ++i) {
    auto d1 = determinant(A, B, vs[i]);
    auto d2 = determinant(A, B, vs[(i + 1) % n]);

    // éVrtice à esquerda da reta
    if (d1 > -EPS) points.push_back(vs[i]);

    // A aresta cruza a reta
    if (d1 * d2 < -EPS)
        points.push_back(intersection(vs[i], vs[(i + 1) % n], A, B));
}
return points;</pre>
```

### 7.16 Perimeter: polygon

```
#include "./Distance: point to point.cpp"
#include "./template.cpp"

template <typename T>
T perimeter(const vector<Point<T>>& pts) {
    T p = 0.0;
    int n = size(pts);
    for (int i = 0; i < n; i++) {
        p += distance(pts[i], pts[(i + 1) % n]);
    }
    return p;
}</pre>
```

#### 7.17 Point

```
// Basic point/vector struct.
template <typename T>
struct Point {
   T x, y;
   Point(T x = 0, T y = 0) : x(x), y(y) {}
   // Addition, substraction, multiply by constant, dot product, cross
   product.
   friend Point<T> operator+(const Point<T>& p, const Point<T>& q) {
        return PointT>(p.x + q.x, p.y + q.y);
   }
   friend Point<T> operator-(const Point<T>& p, const Point<T>& q) {
        return Point<T>(p.x - q.x, p.y - q.y);
   template <typename T2>
   friend Point<T> operator*(const Point<T>& p, T2 k) {
        return Point<T>(p.x * k, p.y * k);
   friend T dot(const Point<T>& p, const Point<T>& q) {
        return p.x * q.x + p.y * q.y;
   friend T cross(const Point<T>& p, const Point<T>& q) {
```

```
return p.x * q.y - p.y * q.x; };
```

## 7.18 Polygon (regular): apothem

```
#include "./Distance: point to point.cpp"
#include "./template.cpp"
template <typename T>
ld apothem(const vector<Point<T>>& pts) {
    auto s = distance(pts[0], pts[1]);
    int n = size(pts);
    return (s / 2.0) * (1.0 / tan(PI / n));
}
```

## 7.19 Polygon (regular): circumradius

```
#include "./Distance: point to point.cpp"
#include "./template.cpp"

template <typename T>
ld circumradius(const vector<Point<T>>& pts) {
    auto s = distance(pts[0], pts[1]);
    int n = size(pts);
    return (s / 2.0) * (1.0 / sin(PI / (ld)n));
}
```

## 7.20 Polygon: check if is convex

```
#include "./Determinant.cpp"
#include "./template.cpp"

template <typename T>
bool checkIfPolygonIsConvex(vector<Point<T>>& pts) {
    int n = size(pts);
    if (n < 3) return false;
    int l, g, e;
    l = g = e = 0;
    for (int i = 0; i < n; i++) {
        auto d = determinant(pts[i], pts[(i + 1) % n], pts[(i + 2) % n]);
        d ? (d > 0 ? g++ : l++) : e++;
    }
    return l == n or g == n;
}
```

#### 7.21 Rectangle intersection

```
Assumes that the points P, Q that define
  a rectangle are the bottom-left and top-right
   corner, and also that the sides are parallel to the axis.
#pragma once
#include "../Contest/template.cpp"
#include "./Point.cpp"
template <typename T>
optional<pair<Point<T>, Point<T>>> rectangleIntersection(
    const pair<Point<T>, Point<T>> &r1, const pair<Point<T>, Point<T>> &r2
    assert(r1.first.x < r1.second.x && r1.first.y < r1.second.y);
   assert(r2.first.x < r2.second.x && r2.first.y < r2.second.y);</pre>
   T x1 = max(r1.first.x, r2.first.x);
   T \times 2 = min(r1.second.x, r2.second.x);
   T y1 = max(r1.first.y, r2.first.y);
   T y2 = min(r1.second.y, r2.second.y);
   if (x1 \ge x2 \text{ or } y1 \ge y2) \text{ return nullopt};
    return pair<Point<T>, Point<T>>{{x1, y1}, {x2, y2}};
```

#### 7.22 template

```
#pragma once
#include <bits/stdc++.h>
using namespace std;
using ld = long double;
template <tvpename T>
using Point = pair<T, T>:
#define x first
#define y second
const double EPS{1e-6};
const ld PI = acos(-1);
template <tvpename T>
bool equals(T a, T b) {
    if (std::is_floating_point<T>::value)
        return fabs(a - b) < EPS;</pre>
    else
        return a == b:
template <typename T>
bool equals(Point<T> a, Point<T> b) {
    if (std::is floating point<T>::value)
        return fabs(a,x - b,x) < EPS && fabs(a,v - b,v) < EPS:
   else
        return a == b;
}
```

# 8 Graphs

#### 8.1 Heavy-Light Decomposition (point update)

#### 8.1.1 Maximum number on path

```
struct Node {
    ll value:
    Node()
        : value(numeric limits<ll>::min()) {}; // Neutral
                                                // element
    Node(ll v) : value(v) {};
Node combine(Node l, Node r) {
    Node m;
    m.value = max(l.value, r.value);
    return m:
template <typename T = Node, auto F = combine>
struct SegTree {
    int n:
    vector<T> st;
    SegTree(int n) : n(n), st(n \ll 1) {}
    void set(int p, const T &k) {
        for (st[p += n] = k; p >>= 1;) st[p] = F(st[p << 1], st[p << 1])
   1]);
   T query(int l, int r) {
        T ansl. ansr:
        for (l += n, r += n + 1; l < r; l >>= 1, r >>= 1) {
            if (l & 1) ansl = F(ansl, st[l++]);
            if (r \& 1) ansr = F(st[--r], ansr);
        return F(ansl, ansr);
};
template <typename SegT = Node, auto SegOp = combine>
struct HeavyLightDecomposition {
    int n:
    vi ps, ds, sz, heavy, head, pos;
    SegTree<SegT, SegOp> seg;
    HeavyLightDecomposition(const vi2d &q, const vector<SegT> &v, int root
    = 0)
        : n(len(q)), seq(n) {
        ps = ds = sz = heavy = head = pos = vi(n, -1);
        auto dfs = [&](auto &&self, int u) -> void {
            sz[u] = 1;
            int mx = 0:
            for (auto x : q[u])
                if (x != ps[u]) {
                    ps[x] = u:
                    ds[x] = ds[u] + 1;
                    self(self, x);
                    sz[u] += sz[x];
```

```
if (sz[x] > mx) mx = sz[x], heavy[u] = x;
    };
    dfs(dfs, root);
    for (int i = 0, cur = 0; i < n; i++) {
        if (ps[i] == -1 \text{ or heavy}[ps[i]] != i)
            for (int j = i; j != -1; j = heavy[j]) {
                head[i] = i;
                pos[j] = cur++;
    }
    rep(i, 0, n) seq.set(pos[i], v[i]);
vector<pii> disjoint ranges(int u, int v) {
    vector<pii> ret:
    for (; head[u] != head[v]; v = ps[head[v]]) {
        if (ds[head[u]] > ds[head[v]]) swap(u, v);
        ret.eb(pos[head[v]], pos[v]);
    if (ds[u] > ds[v]) swap(u, v);
    ret.eb(pos[u], pos[v]);
    return ret;
SegT query path(int u, int v) {
    SegT res;
    for (auto [l, r] : disjoint ranges(u, v)) {
        res = SegOp(res, seg.query(l, r));
    return res;
SegT query subtree(int u) const {
    return seq.query(pos[u], pos[u] + sz[u] - 1);
void set(int u, SeqT x) { seq.set(pos[u], x); }
```

#### 8.2 2-SAT

};

**Description**: Calculates a valid assignment to boolean variables a, b, c,... to a 2-SAT problem, so that an expression of the type (a||b)&&(!a||c)&&(d||!b)&&... becomes true, or reports that it is unsatisfiable.

**Usage**: Negated variables are represented by bit-inversions  $(\tilde{x})$ .

Returns true iff it is solvable ts.values [0..N-1] holds the assigned values to the vars.

**Time**: O(N + E), where N is the number of boolean variables, and E is the number of clauses.

```
/
struct TwoSat {
    int N;
    vector<vi> gr;
    vi values; // 0 = false, 1 = true
    TwoSat(int n = 0) : N(n), gr(2 * n) {}
```

```
int addVar() { // (optional)
        ar.eb():
        gr.eb();
        return N++;
    void either(int f, int j) {
        f = max(2 * f, -1 - 2 * f);
        i = \max(2 * i, -1 - 2 * i);
        gr[f].pb(j ^ 1);
        gr[j].pb(f ^ 1);
    void setValue(int x) { either(x, x); }
    void implies(int f, int j) { either(\simf, j); } // (optional)
    void atMostOne(const vi &li) { // (optional)
        if (len(li) <= 1) return;</pre>
        int cur = \sim li[0];
        rep(i, 2, len(li)) {
            int next = addVar();
            either(cur, ~li[i]);
            either(cur, next);
            either(\simli[i], next);
            cur = \simnext;
        either(cur, \simli[1]);
    vi val, comp, z;
    int time = 0;
    int dfs(int i) {
        int low = val[i] = ++time, x;
        z.pb(i);
        for (int e : gr[i])
            if (!comp[e]) low = min(low, val[e] ?: dfs(e));
        if (low == val[i]) do {
                 x = z.back();
                 z.ppb();
                 comp[x] = low;
                 if (values[x >> 1] == -1) values[x >> 1] = x & 1;
            } while (x != i);
        return val[i] = low:
    bool solve() {
        values.assign(N, -1);
        val.assign(2 * N, 0);
        comp = val;
        rep(i, 0, 2 * N) if (!comp[i]) dfs(i);
        rep(i, 0, N) if (comp[2 * i] == comp[2 * i + 1]) return 0;
        return 1:
};
```

#### 8.3 BFS-01

**Description**: Similar to a Dijkstra given a weighted graph finds the distance from source s to every other node.

```
Warning: Applicable only when the weight of the edges \in \{0, x\}
vector<pair<ll, int>> adi[maxn];
ll dists[maxn];
int s, n;
void bfs 01() {
    fill(dists, dists + n, oo);
    dist[s] = 0;
    deque<int> q;
    g.emplace back(s);
    while (not q.empty()) {
        auto u = q.front();
        q.pop_front();
        for (auto [v, w] : adi[u]) {
            if (dist[v] <= dist[u] + w) continue;</pre>
            dist[v] = dist[u] + w;
            w ? q.emplace back(v) : q.emplace front(v);
```

#### 8.4 Bellman ford

**Description**: Find shortest path from a single source to all other nodes. Can detect negative cycles.

Time:  $O(V \cdot E)$ 

Time: O(V+E)

```
bool bellman ford(const vector<vector<pair<int, ll>>> &g, int s,
                  vector<ll> &dist) {
   int n = (int)q.size();
   dist.assign(n, LLONG MAX);
   vector<int> count(n);
   vector<char> in queue(n);
   queue<int> q;
   dist[s] = 0;
   q.push(s);
   in queue[s] = true;
   while (not q.empty()) {
        int cur = q.front();
        q.pop();
        in queue[cur] = false;
        for (auto [to, w] : g[cur]) {
            if (dist[cur] + w < dist[to]) {</pre>
                dist[to] = dist[cur] + w;
                if (not in queue[to]) {
                    q.push(to);
                    in queue[to] = true;
                    count[to]++:
                    if (count[to] > n) return false;
            }
       }
```

```
}
return true;
```

# 8.5 Bellman-Ford (find negative cycle)

**Description**: Given a directed graph find a negative cycle by running n iterations, and if the last one produces a relaxation than there is a cycle.

Time:  $O(V \cdot E)$ 

```
const ll oo = 2500 * 1e9;
using graph = vector<vector<pair<int, ll>>>;
vi negative cycle(graph &g, int n) {
    vll d(n, oo);
    vi p(n, -1);
    int x = -1:
    d[0] = 0;
    for (int i = 0; i < n; i++) {
        x = -1;
        for (int u = 0; u < n; u++) {
            for (auto &[v, l] : q[u]) {
                if (d[u] + l < d[v]) {
                    d[v] = d[u] + 1;
                    p[v] = u;
                     \dot{x} = v;
    }
    if (x == -1)
        return {};
    else {
        for (int i = 0; i < n; i++) x = p[x];
        vi cycle;
        for (int v = x; v = p[v]) {
            cvcle.eb(v);
            if (v == x and len(cycle) > 1) break;
        reverse(all(cycle));
        return cycle;
}
```

# 8.6 Biconnected Components

**Description**: Build a vector of vectors, where the i-th vector correspond to the nodes of the i-th, biconnected component, a biconnected component is a subset of nodes and edges in which there is no cut point, also exist at least two distinct routes in vertex between any two vertex in the same biconnected component.

**Time**: O(N+M)

```
const int maxn(5 '00' 000);
int tin[maxn], stck[maxn], bcc_cnt, n, top = 0, timer = 1;
vector<int> q[maxn], nodes[maxn];
```

```
int tarjan(int u, int p = -1) {
    int lowu = tin[u] = timer++;
    int son cnt = 0;
    stck[++top] = u;
    for (auto v : g[u]) {
        if (!tin[v]) {
            son cnt++;
            int lowx = tarjan(v, u);
            lowu = min(lowu, lowx);
            if (lowx >= tin[u]) {
                while (top != -1 \&\& stck[top + 1] != v)
                    nodes[bcc cnt].emplace back(stck[top--]);
                nodes[bcc cnt++] emplace back(u);
        } else {
            lowu = min(lowu, tin[v]);
    if (p == -1 \&\& son cnt == 0) {
        nodes[bcc cnt++].emplace back(u):
    return lowu:
void build bccs() {
    timer = 1:
    top = -1;
    memset(tin, 0, sizeof(int) * n);
    for (int i = 0; i < n; i++) nodes[i] = {};
    bcc cnt = 0;
    for (int u = 0; u < n; u++)
        if (!tin[u]) tarjan(u);
}
```

# 8.7 Binary Lifting/Jumping

**Description**: Given a function/successor grpah answers queries of the form which is the node after k moves starting from u.

**Time**: Build  $O(N \cdot MAXLOG2)$ , Query O(MAXLOG2).

```
const int MAXN(2e5), MAXLOG2(30);
int bl[MAXN][MAXLOG2 + 1];
int N;
int jump(int u, ll k) {
    for (int i = 0; i <= MAXLOG2; i++) {
        if (k & (1ll << i)) u = bl[u][i];
    }
    return u;
}
void build() {
    for (int i = 1; i <= MAXLOG2; i++) {
        for (int j = 0; j < N; j++) {
            bl[j][i] = bl[bl[j][i - 1]][i - 1];
        }
</pre>
```

```
}
```

# 8.8 Bipartite Graph

**Description**: Given a graph, find the 'left' and 'right' side if is a bipartite graph, if is not then a empty vi2d is returned

Time: O(N+M)

```
vi2d bipartite graph(vi2d &adj) {
    int n = len(adj);
    vi side(n, -1);
    vi2d ret(2);
    rep(u, 0, n) {
        if (side[u] == -1) {
            queue<int> q;
            a.emp(u):
            side[u] = 0;
            ret[0].eb(u);
            while (len(q)) {
                int u = q.front();
                q.pop();
                for (auto v : adj[u]) {
                     if (side[v] == -1) {
                         side[v] = side[u] ^ 1;
                         ret[side[v]].eb(v);
                         q.push(v);
                    } else if (side[u] == side[v])
                         return {};
        }
    return ret;
```

#### 8.9 Block-Cut Tree \* \*

**Description**: Builds the Block-Cut of a undirected graph. \*\*

Usage: isGraphCutpoint[u] answers how many connected components \* are created when the node u is removed from the graph, if \* isGraphCutpoint[u] is greater than 1, it means that u is a \* cutpoint. \* \*

Time:  $O(N + \hat{M})$  \* \* Memory: O(N) \* \*

Warning: Always careful with disconnected graphs ! you may end up having \* multiple trees. \* \*

```
#pragma once
#include "../Contest/template.cpp"
struct BlockCutTree {
   int n;
   vi idOnTree, tin, low, stk, isGraphCutpoint, isTreeCutpoint;
   vi2d comps, treeAdj;
```

```
BlockCutTree(vi2d &g)
    : n(len(q)), idOnTree(n), tin(n), low(n), isGraphCutpoint(n) {
    rep(i, 0, n) {
        if (!tin[i]) {
            int timer = 0;
            dfs(i, -1, timer, g);
    }
    buildTree();
void buildTree() {
    int node id = 0;
    rep(u, 0, n) {
        if (isGraphCutpoint[u]) {
            idOnTree[u] = node id++;
            isTreeCutpoint.eb(true);
            treeAdj.pb({});
        }
    }
    for (auto &comp : comps) {
        int node = node id++;
        treeAdj.pb({});
        isTreeCutpoint.eb(false);
        for (int u : comp) {
            if (!isGraphCutpoint[u]) {
                idOnTree[u] = node:
            } else {
                treeAdi[node].eb(idOnTree[u]),
                    treeAdj[idOnTree[u]].eb(node);
        }
void dfs(int u, int p, int &timer, vi2d &g) {
    tin[u] = low[u] = ++timer:
    stk.eb(u);
    for (auto v : q[u]) {
        if (v == p) continue;
        if (!tin[v]) {
            dfs(v, u, timer, g);
            chmin(low[u], low[v]);
            if (low[v] >= tin[u]) {
                isGraphCutpoint[u] += (tin[u] > 1 or tin[v] > 2);
                comps.pb({u}):
                while (comps.back().back() != v) {
                    comps.back().eb(stk.back());
                    stk.ppb();
        } else
            low[u] = min(low[u], tin[v]);
    }
int countMandatoryNodesOnPath(int startNode, int endNode);
```

};

#### 8.10 Centroid Decomposition

**Description**: Builds a vector fat where  $fat_i$  is who is the father of the node i in the centroid decomposed tree.

```
#pragma once
#include "../Contest/template.cpp"
vi centroidDecomposition(const vi2d &a) {
    int n = len(q);
    vi fat(n, -1), szt(n), tk(n);
    function<int(int, int)> calcsz = [&](int x, int f) {
        szt[x] = 1;
        for (auto y : g[x])
            if (y != f \&\& !tk[y]) szt[x] += calcsz(y, x);
        return szt[x]:
    function<void(int, int, int)> cdfs = [&](int x, int f, int sz) {
        if (sz < 0) sz = calcsz(x, -1);
        for (auto y : g[x])
            if (!tk[y] \&\& szt[y] * 2 >= sz) {
                szt[x] = 0;
                cdfs(y, f, sz);
                return:
        tk[x] = true;
        fat[x] = f;
        for (auto y : g[x])
            if (!tk[y]) cdfs(y, x, -1);
    cdfs(0, -1, -1);
    return fat;
}
```

# 8.11 Count mandatory nodes on a single path \* \*

**Description**: Given a startNode and an endNode, count the mandatory nodes \* in the path from startNode to endNode, that is the number of nodes such \* that are present in every possible such path. \* \*

Time: O(N+M) \* \* Memory: O(N) \* \*

Warning: The startNode and endNode is always included in the counting, \* ajust your final answer depending on the problem. Be careful with a \* disconnected graph where the path may not exist, treat it appart !. \* \*

```
#pragma once
#include "../Contest/template.cpp"
#include "./Block-Cut tree.cpp"

int BlockCutTree::countMandatoryNodesOnPath(int startNode, int endNode) {
    startNode = idOnTree[startNode], endNode = idOnTree[endNode];
    int ans = !isTreeCutpoint[startNode] + !isTreeCutpoint[endNode];
    int artPoints = 0;
```

```
function<void(int, int)> dfsCount = [&](int u, int p) {
    artPoints += isTreeCutpoint[u];
    if (u == endNode) ans += artPoints;
    for (auto v : treeAdj[u]) {
        if (v != p) {
            dfsCount(v, u);
        }
    }
    artPoints -= isTreeCutpoint[u];
};
dfsCount(startNode, -1);
return ans;
}
```

#### 8.12 DSU query

```
struct DSU {
    V<ii>p;
    V<int> s;
    int sum = 0;
    DSU(int n) : p(n, \{-1, -1\}), s(n, 1) \{\}
    int find(int x) {
        if (p[x].ff < 0) return x;
        return find(p[x].ff);
    void join(int x, int y, int w) {
        x = find(x);
        v = find(v);
        if (x == y) return;
        sum += w;
        if (s[x] < s[y]) swap(x, y);
        s[x] += s[y];
        p[y] = mp(x, w);
    int query(int x, int y) {
        int r = 0;
        while (x != y) {
            if (s[x] < s[y])
                r = max(r, p[x].ss), x = p[x].ff;
            else
                r = max(r, p[y].ss), y = p[y].ff;
        return r:
};
```

# 8.13 D'Escopo-Pape

**Description**: Is a single source shortest path that works faster than Dijkstra's algorithm and the Bellman-Ford algorithm in most cases, and will also work for negative edges. However not for negative cycles. There exists cases where it runs in exponential time.

Usage: Returns a pair containing two vectors, the first one with the distance from s to every other node, and another one with the ancestor of each node, note that the ancestor of s is -1

```
using Edge = pair<ll, int>;
using Adj = vector<vector<Edge>>;
pair<vll, vi> desopo pape(int s, int n, const Adj &adj) {
    vll ds(n, LLONG \overline{M}AX), ps(n, -1);
    ds[s] = 0;
    vi ms(n, 2);
    deque<int> q;
    q.eb(s);
    while (len(q)) {
        int u = q.front();
        q.pop_front();
        ms[u] = 0;
        for (auto [w, v] : adj[u]) {
            if (chmin(ds[v], w + ds[u])) {
                 ps[v] = u;
                 if (ms[v] == 2)
                     ms[v] = 1, q.pb(v);
                 else if (ms[v] == 0)
                     ms[v] = 1, q.pf(v);
        }
    return {ds, ps};
```

#### 8.14 Dijkstra

```
const int MAXN = 1'00'000:
const ll MAXW = 1'000'0001;
constexpr 11\ 00 = MAXW * MAXN + 1:
using Edge = pair<ll, int>; // { weigth, node}
using Adi = vector<vector<Edge>>;
template <tvpename T>
using min heap = priority queue<T, vector<T>, greater<T>>;
pair<vll, vi> dijkstra(const Adj &g, int s) {
    int n = len(q);
    min heap<Edge> pg;
    vll ds(n, 00);
    vi ps(n, -1);
    pq.emp(0, s);
    ds[s] = 0;
    while (len(pq)) {
        auto [du, u] = pq.top();
        pq.pop();
        if (ds[u] < du) continue;
        for (auto [w, v] : g[u]) {
            ll ndv = du + w;
            if (chmin(ds[v], ndv)) {
                ps[v] = u;
                pq.emp(ndv, v);
```

```
}
}
return {ds, ps};

// optional !
vi recover_path(int source, int ending, const vi &ps) {
    if (ps[ending] == -1) return {};
    int cur = ending;
    vi ans;
    while (cur != -1) {
        ans.eb(cur);
        cur = ps[cur];
    }
    reverse(all(ans));
    return ans;
}
```

#### 8.15 Dijkstra (K-shortest pahts)

```
const ll oo = 1e9 * 1e5 + 1;
using adj = vector<vector<pll>>>;
vector<priority queue<ll>> dijkstra(const vector<vector<pll>>> &g, int n,
   int s,
                                     int k) {
   priority queue<pll, vector<pll>, greater<pll>>> pq;
   vector<priority queue<ll>> dist(n);
   dist[0].emplace(0);
   pq.emplace(0, s);
   while (!pq.empty()) {
        auto [d1, v] = pq.top();
        pq.pop();
        if (not dist[v].empty() and dist[v].top() < d1) continue;</pre>
        for (auto [d2, u] : g[v]) {
            if (len(dist[u]) < k) {</pre>
                pq.emplace(d2 + d1, u);
                dist[u].emplace(d2 + d1);
            } else ·
                if (dist[u].top() > d1 + d2) {
                    dist[u].pop();
                    dist[u].emplace(d1 + d2);
                    pq.emplace(d2 + d1, u);
            }
        }
    return dist;
```

# 8.16 Extra Edges to Make Digraph Fully Strongly Connected

**Description**: Given a directed graph G find the necessary edges to add to make the graph a single strongly connected component.

```
Time: O(N+M)
Memory: O(N)
struct SCC {
    int n, num sccs;
    vi2d adj;
    vi scc id;
    SCC(int _n) : n(_n), num_sccs(0), adj(n), scc_id(n, -1) {}
    SCC(const vi2d & adj) : SCC(len( adj)) {
        adj = adj;
        find_sccs();
    void add edge(int u, int v) { adj[u].eb(v); }
    void find sccs() {
        int timer = 1;
        vi tin(n), st;
        st.reserve(n);
        function<int(int)> dfs = [&](int u) -> int {
            int low = tin[u] = timer++, siz = len(st);
            st.eb(u);
            for (int v : adj[u])
                 if (scc id[v] < 0) low = min(low, tin[v] ? tin[v] : dfs(v)
    );
            if (tin[u] == low) {
                 rep(i, siz, len(st)) scc_id[st[i]] = num_sccs;
                 st.resize(siz):
                 num sccs++;
            return low:
        };
        for (int i = 0; i < n; i++)
            if (!tin[i]) dfs(i);
};
vector<array<int, 2>> extra_edges(const vi2d &adj) {
    SCC scc(adi):
    auto scc_id = scc.scc_id;
    auto num sccs = scc.num sccs;
    if (num sccs == 1) return {};
    int n = len(adj);
    vi2d scc adj(num sccs);
    vi zero \overline{in}(num \ s\overline{ccs}, 1);
    rep(u, \overline{0}, n) {
        for (int v : adj[u]) {
            if (scc id[u] == scc id[v]) continue;
             scc adj[scc id[u]].eb(scc id[v]);
            zero in[scc id[v]] = 0;
    int random source = max element(all(zero in)) - zero in.begin();
    vi vis(num sccs);
    function\langle int(int) \rangle dfs = [&](int u) {
```

```
if (empty(scc adj[u])) return u;
    for (int v : scc adj[u])
        if (!vis[v]) -{
            vis[v] = 1;
            int zero_out = dfs(v);
            if (zero out != -1) return zero_out;
    return (int)-1;
};
vector<array<int, 2>> edges;
vi in unused;
rep(i, 0, num_sccs) {
    if (zero \overline{in}[i]) {
        vis[i] = 1;
        int zero out = dfs(i);
        if (zero out != -1)
            edges.push back({zero out, i});
            in unused.push back(i);
    }
}
rep(i, 1, len(edges)) { swap(edges[i][0], edges[i - 1][0]); }
rep(i, 0, num_sccs) {
    if (scc adj[i].empty() && !vis[i]) {
        if (!in unused.empty()) {
            edges.push back({i, in unused.back()});
            in unused.pop_back();
        } else {
            edges.push back({i, random source});
    }
for (int u : in unused) edges.push back({0, u});
vi to node(num sccs);
rep(i, 0, n) to node[scc id[i]] = i;
for (auto &[u, \overline{v}] : edges) u = to_node[u], v = to_node[v];
return edges;
```

#### 8.17 Find Articulation/Cut Points

**Description**: Given an **undirected** graph find it's articulation points.

Time: O(N+M)

}

Warning: A vertex u can be an articulation point if and only if has at least 2 adjascent vertex

```
const int MAXN(100);
int N;
vi2d G;
int timer;
int tin[MAXN], low[MAXN];
set<int> cpoints;
int dfs(int u, int p = -1) {
```

```
int cnt = 0:
    low[u] = tin[u] = timer++;
    for (auto v : G[u])
        if (not tin[v]) {
            cnt++;
            dfs(v, u);
            if (low[v] >= tin[u]) cpoints.insert(u);
            low[u] = min(low[u], low[v]);
        } else if (v != p)
            low[u] = min(low[u], tin[v]);
    return cnt;
void getCutPoints() {
    memset(low, 0, sizeof(low));
    memset(tin, 0, sizeof(tin));
    cpoints.clear();
    timer = 1:
    for (int i = 0; i < N; i++) {
        if (tin[i]) continue;
        int cnt = dfs(i);
        if (cnt == 1) cpoints.erase(i);
}
```

#### 8.18 Find Bridge-Tree components

Usage: label2CC(u, p) finds the 2-edge connected component of every node.

Time: O(n+m)

```
const int maxn(3 '00' 000);
int tin[maxn], compId[maxn], qtdComps;
vi q[maxn], stck;
int n;
int dfs(int u, int p = -1) {
    int low = tin[u] = len(stck);
    stck.emplace back(u);
    bool multEdge = false;
    for (auto v : q[u]) {
        if (v == p and !multEdge) {
            multEdge = 1;
            continue;
        low = min(low, tin[v] == -1 ? dfs(v, u) : tin[v]);
    if (low == tin[u]) {
        for (int i = tin[u]; i < len(stck); i++) compId[stck[i]] =</pre>
    atdComps:
        stck.resize(tin[u]);
        qtdComps++;
    }
    return low:
```

```
void label2CC() {
    memset(compId, -1, sizeof(int) * n);
    memset(tin, -1, sizeof(int) * n);
    stck.reserve(n);
    for (int i = 0; i < n; i++) {
        if (tin[i] == -1) dfs(i);
    }
}</pre>
```

#### 8.19 Find Bridges

**Description:** Find every bridge in a **undirected** connected graph.

Warning: Remember to read the graph as pair where the second is the id of the edge! @Time: O(N + M) const int MAXN(10000), MAXM(100000):

```
int N, M, clk, tin[MAXN], low[MAXN], isBridge[MAXM];
vector<pii> G[MAXN];
void dfs(int u, int p = -1) {
    tin[u] = low[u] = clk++;
    for (auto [v, i] : G[u]) {
        if (v == p) continue;
        if (tin[v]) {
            low[u] = min(low[u], tin[v]);
        } else {
            dfs(v, u);
            low[u] = min(low[u], low[v]);
            if (low[v] > tin[u]) {
                isBridge[i] = 1;
        }
   }
void findBridges() {
    fill(tin, tin + N, 0);
    fill(low, low + N, 0);
    fill(isBridge, isBridge + M, 0);
    clk = 1;
    for (int i = 0; i < N; i++) {
        if (!tin[i]) dfs(i);
}
```

#### 8.20 Find Centroid

**Description**: Given a tree (don't forget to make it 'undirected'), find it's centroids. @Time : O(V)

```
#pragma once
#include "../Contest/template.cpp"
void dfs(int u, int p, int n, vi2d &g, vi &sz, vi &centroid) {
    sz[u] = 1;
    bool iscentroid = true;
    for (auto v : g[u])
```

```
if (v != p) {
          dfs(v, u, n, g, sz, centroid);
          if (sz[v] > n / 2) iscentroid = false;
          sz[u] += sz[v];
     }
     if (n - sz[u] > n / 2) iscentroid = false;
     if (iscentroid) centroid.eb(u);
}
vi getCentroid(vi2d &g, int n) {
     vi centroid;
     vi sz(n);
     dfs(0, -1, n, g, sz, centroid);
     return centroid;
}
```

#### 8.21 Find bridges (online)

```
// O((n+m)*log(n))
struct BridgeFinder {
    // 2ecc = 2 edge conected component
    // cc = conected component
    vector<int> parent, dsu_2ecc, dsu_cc, dsu_cc_size;
    int bridges, lca iteration;
    vector<int> last visit;
    BridgeFinder(int n)
        : parent(n, -1),
          dsu 2ecc(n),
          dsu cc(n),
          dsu cc size(n, 1),
          bridges(0),
          lca iteration(0),
          last visit(n) {
        for (int i = 0; i < n; i++) {
            dsu \ 2ecc[i] = i;
            dsu cc[i] = i;
    int find 2ecc(int v) {
        if (\overline{v} == -1) return -1;
        return dsu 2ecc[v] == v ? v : dsu 2ecc[v] = find 2ecc(dsu 2ecc[v])
    int find cc(int v) {
        v = find 2ecc(v);
        return dsu cc[v] == v ? v : dsu cc[v] = find cc(dsu cc[v]);
    void make root(int v) {
        v = find 2ecc(v);
        int root = v;
        int child = -1;
        while (v != -1) {
            int p = find 2ecc(parent[v]);
```

```
parent[v] = child;
         dsu cc[v] = root;
         child = v;
         v = p;
    dsu cc size[root] = dsu cc size[child];
void merge_path(int a, int b) {
    ++lca iteration;
    vector<int> path a, path b;
    int lca = -1:
    while (lca == -1) {
         if (a != -1) {
             a = find 2ecc(a);
             path a.push back(a);
             if (last visit[a] == lca iteration) {
                 lca = a:
                 break:
             last visit[a] = lca iteration;
             a = parent[a];
         if (b != -1) {
             b = find 2ecc(b);
             path b.push back(b);
             if (last visit[b] == lca iteration) {
                 lca = b:
                 break:
             last visit[b] = lca iteration;
             b = parent[b];
    for (auto v : path a) {
         dsu \ 2ecc[v] = \overline{l}ca;
         if (v == lca) break;
         --bridges;
    for (auto v : path b) {
        dsu \ 2ecc[v] = \overline{l}ca;
         if (v == lca) break;
         --bridges;
    }
}
void add edge(int a, int b) {
    a = \overline{find} \ 2ecc(a);
    b = find 2ecc(b);
    if (a == b) return;
    int ca = find cc(a);
    int cb = find cc(b);
    if (ca != cb) {
         ++bridges;
         if (dsu_cc_size[ca] > dsu_cc_size[cb]) {
             swap(a, b);
```

```
swap(ca, cb);
}
make_root(a);
parent[a] = dsu_cc[a] = b;
dsu_cc_size[cb] += dsu_cc_size[a];
} else {
    merge_path(a, b);
}
};
```

#### 8.22 Floyd Warshall

**Description**: Simply finds the minimal distance for each node to every other node.  $O(V^3)$ 

#### 8.23 Functional/Successor Graph

**Description**: Given a functional graph find the vertice after k moves starting at u and also the distance between u and v, if it's impossible to reach v starting at u returns -1.

**Time**: build  $O(N \cdot MAXLOG2)$ , kth O(MAXLOG2), dist O(MAXLOG2)

```
const int MAXN(2 '000' 000), MAXLOG2(24);
int N;
vi2d succ(MAXN, vi(MAXLOG2 + 1));
vi dst(MAXN, 0);
int vis[MAXN];
void dfsbuild(int u) {
    if (vis[u]) return;
    vis[u] = 1;
    int v = succ[u][0];
    dfsbuild(v);
    dst[u] = dst[v] + 1;
void build() {
    for (int i = 0; i < N; i++) {
        if (not vis[i]) dfsbuild(i);
    for (int k = 1; k \le MAXLOG2; k++) {
        for (int i = 0; i < N; i++) {
            succ[i][k] = succ[succ[i][k-1]][k-1];
```

```
}
int kth(int u, ll k) {
    if (k <= 0) return u;
    for (int i = 0; i <= MAXL0G2; i++)
        if ((1ll << i) & k) u = succ[u][i];
    return u;
}
int dist(int u, int v) {
    int cu = kth(u, dst[u]);
    if (kth(u, dst[u] - dst[v]) == v)
        return dst[u] - dst[v];
    else if (kth(cu, dst[cu] - dst[v]) == v)
        return dst[u] + (dst[cu] - dst[v]);
    else
        return -1;
}
</pre>
```

# 8.24 Heavy light decomposition (supreme)

```
struct HLD {
   int V:
   int id:
   int nb_heavy_path;
   std::vector<std::vector<int>> g;
   std::vector<pair<int, int>> edges; // edges of the tree
                                        // par[i] = parent of
   std::vector<int> par;
                                        // vertex i (Default: -1)
   std::vector<int> depth;
                                        // depth[i] = distance between
   root
                                        // and vertex i
                                        // subtree sz[i] = size of
   std::vector<int> subtree sz;
                                        // subtree whose root is i
   std::vector<int> heavy child;
                                        // heavy child[i] = child of
                                        // vertex i on heavy path
                                        // (Default: -1)
   std::vector<int> tree id;
                                        // tree id[i] = id of tree vertex
                                        // i belongs to
   std::vector<int> aligned id,
       aligned id inv;
                                     // aligned id[i] = aligned
                                     // id for vertex i
                                     // (consecutive on heavy
                                     // edges)
   std::vector<int> head;
                                     // head[i] = id of vertex on heavy
                                     // path of vertex i, nearest to root
   std::vector<int> head ids;
                                     // consist of head vertex id's
   std::vector<int> heavy path id; // heavy path id[i] =
                                     // heavy_path_id for vertex
                                     // [i]
   HLD(const std::vector<std::vector<int>> &e, vector<int> roots = {0})
        : HLD((int)e.size()) {
       q = e;
       build(roots);
   HLD(int sz = 0)
        : V(sz),
```

```
id(0),
      nb heavy path(0),
      g(sz),
      par(sz),
      depth(sz),
      subtree_sz(sz),
      heavy child(sz),
      tree id(sz, -1),
      aligned id(sz),
      aligned id inv(sz),
      head(sz),
      heavy path id(sz, -1) {}
void add edge(int u, int v) {
    edges.emplace back(u, v);
    q[u].emplace \overline{b}ack(v);
    g[v].emplace back(u);
void build dfs(int root) {
    std::stack<std::pair<int, int>> st;
    par[root] = -1;
    depth[root] = 0;
    st.emplace(root, 0);
    while (!st.empty()) {
        int now = st.top().first;
        int &i = st.top().second;
        if (i < (int)g[now] size()) {</pre>
             int nxt = q[now][i++];
             if (nxt == par[now]) continue;
             par[nxt] = now;
             depth[nxt] = depth[now] + 1;
             st.emplace(nxt, 0);
        } else {
             st.pop();
            int max_sub_sz = 0;
             subtree sz[now] = 1;
             heavy c\overline{h}ild[now] = -1;
             for (auto nxt : q[now]) {
                 if (nxt == par[now]) continue;
                 subtree sz[now] += subtree sz[nxt];
                 if (max sub sz < subtree sz[nxt])</pre>
                     max sub sz = subtree sz[nxt], heavy child[now] =
nxt;
        }
    }
void build bfs(int root, int tree id now) {
    std::queue<int> q({root});
    while (!q.empty()) {
        int h = q.front();
        q.pop();
        head ids.emplace back(h);
         for (int now = h; now != -1; now = heavy child[now]) {
             tree id[now] = tree id now;
```

```
aligned id[now] = id++;
            aligned id inv[aligned id[now]] = now;
            heavy path id[now] = nb heavy path;
            head[now] = h:
            for (int nxt : q[now])
                if (nxt != par[now] and nxt != heavy child[now])
                     a.push(nxt):
        nb heavy path++;
void build(std::vector<int> roots = {0}) {
    int tree id now = 0;
    for (auto r : roots) build_dfs(r), _build_bfs(r, tree_id_now++);
// data[i] = value of vertex i
template <class T>
std::vector<T> segtree rearrange(const std::vector<T> &data) const {
    assert(int(data.size()) == V);
    std::vector<T> ret;
    ret.reserve(V);
    for (int i = 0; i < V; i++) ret.emplace back(data[aligned id inv[i</pre>
]]);
    return ret;
// data[i] = weight of edge[i]
template <class T>
std::vector<T> segtree rearrange weighted(
    const std::vector<T> &data) const {
    assert(data.size() == edges.size());
    vector<T> ret(V);
    for (int i = 0; i < (int)edges.size(); i++) {</pre>
        auto [u, v] = edges[i];
        if (depth[u] > depth[v]) swap(u, v);
        ret[aligned_id[v]] = data[i];
    return ret;
int segtree edge index(int i) const {
    auto [u, v] = edges[i];
    if (depth[u] > depth[v]) swap(u, v);
    return aligned id[v];
// query for vertices on path [u, v] (INCLUSIVE)
void for each vertex(int u, int v, const auto &f) const {
    static assert(std::is invocable r v<void, decltype(f), int, int>);
    assert(tree id[u] == tree id[v] and tree id[u] >= 0);
    while (true) {
        if (aligned id[u] > aligned id[v]) std::swap(u, v);
        f(std::max(\overline{aligned}_id[head[\overline{v}]), aligned_id[u]), aligned_id[v])
        if (head[u] == head[v]) break;
        v = par[head[v]];
}
```

```
void for each vertex noncommutative(int from, int to, const auto &fup,
                                     const auto &fdown) const {
    static assert(std::is invocable r v<void, decltype(fup), int, int
>);
    static assert(std::is invocable r v<void, decltype(fdown), int,
int>);
    assert(tree id[from] == tree id[to] and tree id[from] >= 0);
    int u = from, v = to:
    const int lca = lowest common ancestor(u, v), dlca = depth[lca];
    while (u >= 0 and depth[u] > dlca) {
        const int p = (depth[head[u]] > dlca ? head[u] : lca);
        fup(aligned id[p] + (p == lca), aligned id[u]), u = par[p];
    static std::vector<std::pair<int, int>> lrs;
    int sz = 0:
    while (v >= 0 \text{ and } depth[v] >= dlca) {
        const int p = (depth[head[v]] >= dlca ? head[v] : lca);
        if (int(lrs.size()) == sz) lrs.emplace back(0, 0);
        lrs.at(sz++) = \{p, v\}, v = par.at(p);
    while (sz--)
        fdown(aligned id[lrs.at(sz).first], aligned id[lrs.at(sz).
second]);
// query for edges on path [u, v]
void for_each_edge(int u, int v, const auto &f) const {
    static assert(std::is invocable r v<void, decltype(f), int, int>);
    assert(tree_id[u] == tree_id[v] and tree_id[u] >= 0);
    while (true) {
        if (aligned_id[u] > aligned_id[v]) std::swap(u, v);
        if (head[u] != head[v]) {
            f(aligned id[head[v]], aligned id[v]);
            v = par[head[v]];
        } else {
            if (u != v) f(aligned id[u] + 1, aligned id[v]);
            break:
    }
// lowest common ancestor: O(log V)
int lowest common ancestor(int u, int v) const {
    assert(tree id[u] == tree id[v] and tree id[u] >= 0);
    while (true) {
        if (aligned_id[u] > aligned_id[v]) std::swap(u, v);
        if (head[u] == head[v]) return u;
        v = par[head[v]];
int distance(int u, int v) const {
    assert(tree id[u] == tree id[v] and tree id[u] >= 0);
    return depth[u] + depth[v] - 2 * depth[lowest common ancestor(u, v
)];
// Level ancestor, O(log V)
```

```
// if k-th parent is out of range, return -1
    int kth parent(int v, int k) const {
        if (k < 0) return -1;
        while (v \ge 0) {
            int h = head.at(v), len = depth.at(v) - depth.at(h);
            if (k <= len) return aligned id inv.at(aligned id.at(v) - k);</pre>
            k = len + 1, v = par.at(h);
        return -1;
   }
    // Jump on tree, O(log V)
    int s to t by k steps(int s, int t, int k) const {
        if (\overline{k} < 0) return -1;
        if (k == 0) return s:
        int lca = lowest common ancestor(s, t);
        if (k <= depth.at(s) - depth.at(lca)) return kth parent(s, k);</pre>
        return kth parent(t, depth.at(s) + depth.at(t) - depth.at(lca) * 2
    - k);
};
```

#### 8.25 Kruskal

**Description**: Find the minimum spanning tree of a graph. **Time**:  $O(E \log E)$ 

```
#include "./../Data Structures/DSU.cpp"
vector<tuple<ll, int, int>> kruskal(int n, vector<tuple<ll, int, int>> &
    edges) {
    DSU dsu(n);
    vector<tuple<ll, int, int>> ans;
    sort(all(edges));
    for (auto [a, b, c] : edges) {
        if (dsu.same_set(b, c)) continue;
        ans.emplace_back(a, b, c);
        dsu.union_set(b, c);
    }
    return ans;
}
```

#### 8.26 Lowest Common Ancestor

**Description**: Given two nodes of a tree find their lowest common ancestor, or their distance

```
#pragma once
#include "../Contest/template.cpp"
template <typename T>
struct SparseTable {
   vector<T> v;
   int n;
   static const int b = 30;
   vi mask, t;
```

```
int op(int x, int y) { return v[x] < v[y] ? x : y; }
    int msb(int x) { return builtin clz(1) - builtin clz(x); }
    SparseTable() {}
    SparseTable(const vector<T> &v ) : v(v ), n(v.size()), mask(n), t(n) {
        for (int i = 0, at = 0; i < n; \max \overline{k}[i++] = at |= 1) {
            at = (at << 1) & ((1 << b) - 1);
            while (at and op(i, i - msb(at & -at)) == i) at ^= at & -at;
        for (int i = 0; i < n / b; i++)
            t[i] = b * i + b - 1 - msb(mask[b * i + b - 1]);
        for (int j = 1; (1 << j) <= n / b; j++)
            for (int i = 0; i + (1 << j) <= n / b; i++)
                t[n / b * i + i] = op(t[n / b * (j - 1) + i],
                                       t[n / b * (j - 1) + i + (1 << (j -
   1))]);
    int small(int r, int sz = b) { return r - msb(mask[r] \& ((1 << sz) -
   1)): }
    T query(int l, int r) {
        if (r - l + 1 \le b) return small(r, r - l + 1);
        int ans = op(small(l + b - 1), small(r));
        int x = l / b + 1, y = r / b - 1;
        if (x \le y) {
            int j = msb(y - x + 1);
                op(ans, op(t[n / b * j + x], t[n / b * j + y - (1 << j) +
   1]));
        return ans:
};
struct LCA {
    SparseTable<int> st;
    int n;
    vi v, pos, dep;
    vll wdep;
    LCA(const Graph \&g, int root) : n(len(g)), pos(n), wdep(n) {
        dfs(root, 0, -1, q);
        st = SparseTable<int>(vector<int>(all(dep)));
    void dfs(int i, int d, int p, const Graph &g) {
        v.eb(len(dep)) = i, pos[i] = len(dep), dep.eb(d);
        for (auto [w, j] : g[i])
            if (j != p) {
                wdep[j] = wdep[i] + w;
                dfs(i, d + 1, i, q);
                v.eb(len(dep)) = i, dep.eb(d);
    int lca(int a, int b) {
        int l = min(pos[a], pos[b]);
        int r = max(pos[a], pos[b]);
        return v[st.query(l, r)];
    il dist(int a, int b) { return wdep[a] + wdep[b] - 2ll * wdep[lca(a, b
```

```
)]; }
};
```

#### 8.27 Lowest Common Ancestor (Binary Lifting)

**Description**: Given a directed tree, finds the LCA between two nodes using binary lifting, and answer a few queries with it.

Usage:

- lca: returns the LCA between the two given nodes
- on path: fids if c is in the path from a to b

**Time**: build  $O(N \cdot MAXLOG_2)$ , all queries  $O(MAXLOG_2)$ 

```
struct LCA {
   int n;
    const int maxlog;
   vector<vector<int>> up;
   vector<int> depth:
   LCA(const vector<vector<int>> &tree)
        : n(tree.size()),
          maxlog(ceil(log2(n))),
          up(n, vector<int>(maxlog + 1)),
          depth(n, -1) {
        for (int i = 0; i < n; i++) {
            if (depth[i] == -1) {
                depth[i] = 0;
                dfs(i, -1, tree);
            }
        }
   void dfs(int u, int p, const vector<vector<int>> &tree) {
        if (p != -1) {
            depth[u] = depth[p] + 1;
            up[u][0] = p;
            for (int i = 1; i <= maxlog; i++) {
                up[u][i] = up[up[u][i - 1]][i - 1];
        for (int v : tree[u]) {
            if (v == p) continue;
            dfs(v, u, tree);
   int kth jump(int u, int k) {
        for (int i = maxlog; i \ge 0; i--) {
            if ((1 << i) & k) {
                u = up[u][i];
        return u;
   int lca(int u, int v) {
        if (depth[u] < depth[v]) swap(u, v);</pre>
        int diff = depth[u] - depth[v];
        u = kth jump(u, diff);
```

```
if (u == v) return u;
for (int i = maxlog; i >= 0; i--) {
        if (up[u][i] != up[v][i]) {
            u = up[u][i];
            v = up[v][i];
        }
    }
    return up[u][0];
}
bool on_path(int u, int v, int s) {
        int uv = lca(u, v), us = lca(u, s), vs = lca(v, s);
        return (uv == s or (us == uv and vs == s) or (vs == uv and us == s));
}
int dist(int u, int v) {
        return depth[u] + depth[v] - 2 * depth[lca(u, v)];
}
};
```

#### 8.28 Maximum flow (Dinic)

**Description:** Finds the **maximum flow** in a graph network, given the **source** s and the **sink** t. Add edge from a to b with capcity c.

**Time**: In general  $O(E \cdot V^2)$ , if every capacity is 1, and every vertice has in degree equal 1 or out degree equal 1 then  $O(E \cdot \sqrt{V})$ ,

Warning: Suffle the edges list for every vertice may take you out of the worst case

```
struct Dinic {
    struct Edge {
        int to, rev;
        ll c, oc;
        Il flow() { return max(oc - c, OLL); } // if you need flows
    };
    vi lvl, ptr, q;
    vector<vector<Edge>> adj;
    Dinic(int n) : lvl(n), ptr(n), q(n), adj(n) {}
    void addEdge(int a, int b, ll c, ll rcap = 0) {
        adj[a].pb({b, len(adj[b]), c, c});
        adj[b].pb({a, len(adj[a]) - 1, rcap, rcap});
    ll dfs(int v, int t, ll f) {
        if (v == t || !f) return f;
        for (int &i = ptr[v]; i < len(adj[v]); i++) {</pre>
            Edge &e = adj[v][i];
            if (lvl[e.to] == lvl[v] + 1)
                if (ll p = dfs(e.to, t, min(f, e.c))) {
                    e.c -= p, adj[e.to][e.rev].c += p;
                     return p;
        return 0:
```

```
ll maxFlow(int s, int t) {
        Il flow = 0;
        q[0] = s;
        rep(L, 0, 31) {
            do { // 'int L=30' maybe faster for random
                  // data
                lvl = ptr = vi(len(q));
                int qi = 0, qe = lvl[s] = 1;
                while (qi < qe && !lvl[t]) {</pre>
                    int v = q[qi++];
                    for (Edge e : adi[v])
                        if (!lvl[e.to] && e.c >> (30 - L))
                             q[qe++] = e.to, lvl[e.to] = lvl[v] + 1;
                while (ll p = dfs(s, t, LLONG MAX)) flow += p;
            } while (lvl[t]);
        return flow;
    bool leftOfMinCut(int a) { return lvl[a] != 0; }
};
```

#### 8.29 Minimum Cost Flow

**Description**: Given a network find the minimum cost to achieve a flow of at most f. Works with **directed** and **undirected** graphs **Usage**:

- add(u, v, c, w): adds an edge from u to v with capacity c and cost w.
- flow(f): return a pair (flow, cost) with the maximum flow until f with source at s and sink at t, with the minimum cost possible.

**Time**:  $O(N \cdot M + f \cdot m \log n)$ 

```
template <typename T>
struct MinCostFlow {
    struct Edge {
        int to;
        ll c, rc; // capcity, residual capacity
                   // cost
       T w:
    };
   <u>int</u> n, s, t;
   vector<Edge> edges;
   vi2d q;
   vector<T> dist;
   vi pre;
   MinCostFlow() {}
   MinCostFlow(int n_, int _s, int _t) : n(n_), s(_s), t(_t), g(n) \{ \}
   void addEdge(int u, int v, ll c, T w) {
        g[u].pb(len(edges));
        edges.eb(v, c, 0, w);
        g[v].pb(len(edges));
        edges.eb(u, 0, 0, -w);
   // {flow, cost}
   pair<ll, T> flow(ll flow limit = LLONG MAX) {
```

```
II flow = 0:
     T cost = 0:
     while (flow < flow_limit and dijkstra(s, t)) {</pre>
         ll aug = LLONG MAX;
         for (int i = t; i != s; i = edges[pre[i] ^ 1].to) {
             aug = min({flow limit - flow, aug, edges[pre[i]].c});
         for (int i = t; i != s; i = edges[pre[i] ^ 1].to) {
             edges[pre[i]].c -= aug;
             edges[pre[i] ^1].c += aug;
             edges[pre[i]].rc += aug;
             edges[pre[i] ^ 1].rc -= aug;
         flow += aug;
         cost += (T)aug * dist[t];
     return {flow, cost};
 // Needs to be called after flow method
 vi2d paths() {
     vi2d p;
     for (;;) {
         int cur = s;
         auto &res = p.eb();
         res.pb(cur);
         while (cur != t) {
             bool found = false;
             for (auto i : q[cur]) {
                 auto &[v, , c, cost] = edges[i];
                 if (c > 0) {
                     --c:
                     res.pb(cur = v);
                     found = true:
                     break:
                 }
             if (!found) break;
         if (cur != t) {
             p.ppb();
             break;
     }
     return p;
private:
 bool bellman ford(int s, int t) {
     dist.assign(n, numeric limits<T>::max());
     pre.assign(n, -1);
     vc ing(n, false);
     queue<int> q;
     dist[s] = 0;
     q.push(s);
```

```
while (len(q)) {
        int u = q.front();
        q.pop();
        inq[u] = false;
        for (int i : q[u]) {
            auto [v, c, w, _] = edges[i];
             auto new dist = dist[u] + w;
             if (c > \overline{0} \text{ and } dist[v] > \text{new dist})  {
                 dist[v] = new dist;
                 pre[v] = i;
                 if (not ing[v]) {
                     ina[v] = true:
                     q.push(v);
                 }
            }
        }
    return dist[t] != numeric limits<T>::max();
bool dijkstra(int s, int t) {
    dist.assign(n, numeric limits<T>::max());
    pre.assign(n, -1);
    dist[s] = 0;
    using PQ = pair<T, int>;
    pgmn<PQ> pg;
    pq.emp(0, s);
    while (len(pg)) {
        auto [cost, u] = pq.top();
        pq.pop();
        if (cost != dist[u]) continue;
        for (int i : g[u]) {
             auto [v, c, , w] = edges[i];
             auto new dist = dist[u] + w;
             if (c > \overline{0} \text{ and } dist[v] > new_dist) {
                 dist[v] = new dist;
                 pre[v] = i;
                 pq.emp(new dist, v);
        }
    return dist[t] != numeric limits<T>::max();
```

#### 8.30 Minimum Vertex Cover (already divided)

};

**Description**: Given a bipartite graph g with n vertices at left and m vertices at right, where g[i] are the possible right side matches of vertex i from left side, find a minimum vertex cover. The size is the same as the size of the maximum matching, and the complement is a maximum independent set.

```
vector<int> min_vertex_cover(vector<vector<int>> &g, int n, int m) {
   vector<int> match(m, -1), vis;
```

```
auto find = [&](auto &&self, int j) -> bool {
    if (match[i] == -1) return 1;
    vis[i] = 1;
    int di = match[j];
    for (int e : q[di])
        if (!vis[e] and self(self, e)) {
            match[e] = di;
            return 1;
    return 0;
};
for (int i = 0; i < (int)q.size(); i++) {
    vis.assign(match.size(), 0);
    for (int j : q[i]) {
        if (find(find, j)) {
            match[j] = i;
            break:
        }
    }
int res = (int)match.size() - (int)count(match.begin(), match.end(),
vector<char> lfound(n, true), seen(m);
for (int it : match)
    if (it != -1) lfound[it] = false;
vector<int> q, cover;
for (int i = 0; i < n; i++)
    if (lfound[i]) q.push back(i);
while (!q.empty()) {
    int i = q.back();
    g.pop back();
    lfound[i] = 1;
    for (int e : g[i])
        if (!seen[e] and match[e] !=-1) {
            seen[e] = true;
            q.push back(match[e]);
for (int i = 0; i < n; i++)
    if (!lfound[i]) cover.push back(i);
for (int i = 0; i < m; i++)
    if (seen[i]) cover.push back(n + i);
assert((int)size(cover) == res);
return cover;
```

# 8.31 Prim (MST)

**Description**: Given a graph with N vertex finds the minimum spanning tree, if there is no such three returns inf, it starts using the edges that connect with each  $s_i \in s$ , if none is provided than it starts with the edges of node 0.

Time:  $O(V \log E)$ 

}

```
#include "../Contest/template.cpp"
const int MAXN(1'00'000);
int N:
vector<pair<ll, int>> G[MAXN];
ll prim(vi s = vi(1, 0)) {
   priority queue<pair<ll, int>, vector<pair<ll, int>>, greater<pair<ll,</pre>
   int>>>
        pq;
   vector<char> ingraph(MAXN);
   int ingraphcnt(0);
   for (auto si : s) {
        ingraphcnt++;
        ingraph[si] = true;
        for (auto &[w, v] : G[si]) pg.emplace(w, v);
   ll\ mstcost = 0;
   while (ingraphent < N and !pg.empty()) {</pre>
        ll w:
        int v;
        do {
            tie(w, v) = pq.top();
            pq.pop();
        } while (not pg.empty() and ingraph[v]);
        mstcost += w, ingraph[v] = true, ingraphcnt++;
        for (auto &[w2, v2] : G[v]) {
            pq.emplace(w2, v2);
    return ingraphent == N ? mstcost : oo:
```

#### 8.32 Reachability Tree

```
#include <bits/stdc++.h>
using namespace std;
const int MAXN = 20000 + 100;
int dsu[MAXN];
int n:
const int MAXM = 100000:
int U[MAXM], V[MAXM];
vector<int> adi[MAXN];
int getRoot(int u) {
   if (u == dsu[u]) return u;
   return dsu[u] = getRoot(dsu[u]);
void addEdge(int u, int v) {
   u = aetRoot(u):
   v = getRoot(v);
   dsu[n] = n:
   dsu[u] = dsu[v] = n;
```

```
adj[n].push_back(u);
    if (u != v) adj[n].push_back(v);
    ++n;
}

void build() {
    for (int i = 0; i < n; ++i) dsu[i] = i;
    for (int i = 0; i < m; ++i) addEdge(U[i], V[i]);
}
int32_t main() {
    ios_base::sync_with_stdio(0);
    cin.tie(0);
}</pre>
```

#### 8.33 Shortest Path With K-edges

**Description**: Given an adjacency matrix of a graph, and a number K computes the shortest path between all nodes that uses exactly K edges, so for  $0 \le i, j \le N - 1$  ans[i][j] = "the shortest path between i and j that uses exactly K edges, remember to initialize the adjacency matrix with  $\infty$ .

Time:  $O(N^3 \cdot \log K)$ 

```
template <typename T>
vector<vector<T>> prod(vector<vector<T>> &a, vector<vector<T>> &b) {
    const T oo = numeric limits<T>::max();
    int n = a.size();
    vector<vector<T>> c(n, vector<T>(n, _oo));
    for (int i = 0; i < n; i++)
        for (int j = 0; j < n; j++)
            for (int k = 0; k < n; k++)
                if (a[i][k] != oo and b[k][j] != oo)
                    c[i][i] = m\overline{i}n(c[i][j], a[i][k] + b[k][j]);
    return c;
template <typename T>
vector<vector<T>> shortest with k moves(vector<vector<T>> adj, long long k
   ) {
    if (k == 1) return adj;
    auto ans = adj;
    k--;
    while (k) {
        if (k & 1) ans = prod(ans, adj);
        k >>= 1:
        adj = prod(adj, adj);
    return ans;
```

# 8.34 Strongly Connected Components (struct)

**Description**: Find the connected component for each edge (already in a topological order), some additional functions are also provided.

```
Time: Build: O(V+E)
struct SCC {
   int n, num_sccs;
   vi2d adi;
   vi scc_id;
```

```
SCC(int n) : n(n), num sccs(0), adj(n), scc id(n, -1) {}
   void add edge(int u, int v) { adj[u].eb(v); }
   void find sccs() {
        int timer = 1:
        vi tin(n), st:
        st.reserve(n);
        function<int(int)> dfs = [&](int u) -> int {
            int low = tin[u] = timer++, siz = len(st);
            st.eb(u);
            for (int v : adi[u])
                if (scc id[v] < 0) low = min(low, tin[v] ? tin[v] : dfs(v)
   );
            if (tin[u] == low) {
                rep(i, siz, len(st)) scc id[st[i]] = num sccs;
                st.resize(siz):
                num_sccs++;
            return low;
        };
        for (int i = 0; i < n; i++)
            if (!tin[i]) dfs(i);
   vector<set<int>> build qscc() {
        vector<set<int>> qscc;
        for (int i = 0; i < len(adj); ++i)</pre>
            for (auto j : adj[i])
                if (scc_id[i] != scc_id[j]) gscc[scc_id[i]].emplace(scc_id
   [j]);
        return gscc;
   vi2d per comp() {
        vi2d ret(num sccs);
        rep(i, 0, n) ret[scc id[i]].eb(i);
        reverse(all(ret)); // already in topological order ;)
        return ret:
};
      Topological Sorting (Kahn)
```

**Description:** Finds the topological sorting in a **DAG**, if the given graph is not a **DAG** than an empty vector is returned, need to 'initialize' the **INCNT** as you build the graph. Time: O(V+E)

```
const int MAXN(2 '00' 000);
int INCNT[MAXN];
vi2d GOUT(MAXN);
int N:
```

```
vi toposort() {
    vi order;
    queue<int> q;
    for (int i = 0; i < N; i++)
        if (!INCNT[i]) g.emplace(i);
    while (!q.empty()) {
        auto u = q.front();
        q.pop();
        order emplace back(u);
        for (auto v : GOUT[u]) {
            INCNT[v]--;
            if (INCNT[v] == 0) g.emplace(v);
    }
    return len(order) == N ? order : vi();
```

# 8.36 Topological Sorting (Tarjan)

**Description:** Finds a the topological order for the graph, if there is no such order it means the graph is cyclic, then it returns an empty vector

Time: O(V+E)

```
const int maxn(1 '00' 000);
int n, m;
vi q[maxn];
int not found = 0, found = 1, processed = 2;
int state[maxn];
bool dfs(int u, vi &order) {
    if (state[u] == processed) return true;
    if (state[u] == found) return false;
    state[u] = found;
    for (auto v : q[u]) {
        if (not dfs(v, order)) return false;
    state[u] = processed:
    order.emplace back(u);
    return true;
vi topo sort() {
    vi order:
    memset(state, 0, sizeof state);
    for (int u = 0: u < n: u++) {
        if (state[u] == not found and not dfs(u, order)) return {};
    reverse(all(order));
    return order;
```

#### 8.37 Tree Isomorphism (not rooted)

**Description**: Two trees are considered **isomorphic** if the hash given by thash() is the same. **Time**:  $O(V \cdot \log V)$ 

```
map<vi, int> mphash;
struct Tree {
    int n;
    vi2d q;
    vi sz, cs;
   Tree(int n): n(n), g(n), sz(n) {}
    void add edge(int u, int v) {
        g[u].emplace back(v);
        g[v].emplace back(u);
    void dfs centroid(int v, int p) {
        sz[v] = 1:
        bool cent = true;
        for (int u : q[v])
            if (u != p) {
                dfs centroid(u, v);
                sz[v] += sz[u];
                cent &= not(sz[u] > n / 2);
        if (cent and n - sz[v] <= n / 2) cs.push back(v);</pre>
    int fhash(int v, int p) {
        vi h:
        for (int u : g[v])
            if (u != p) h.push back(fhash(u, v));
        sort(all(h));
        if (!mphash.count(h)) mphash[h] = mphash.size();
        return mphash[h];
    ll thash() {
        cs.clear();
        dfs centroid(0, -1);
        if (cs.size() == 1) return fhash(cs[0], -1);
        ll h1 = fhash(cs[0], cs[1]), h2 = fhash(cs[1], cs[0]);
        return (min(h1, h2) \ll 30ll) + max(h1, h2);
};
```

#### 8.38 Tree Isomorphism (rooted)

**Description**: Given a rooted tree find the hash of each subtree, if two roots of two distinct trees have the same hash they are considered isomorphic

**Time**: hash first time in  $O(\log N_v \cdot N_v)$  where  $(N_v)$  is the of the subtree of v

```
map<vi, int> hasher;
int hs = 0;
struct RootedTreeIso {
   int n;
```

```
vi2d adj;
vi hashes;
RootedTreeIso(int _n) : n(_n), adj(_n), hashes(_n, -1) {};
void add_edge(int u, int v) {
    adj[u].emplace_back(v);
    adj[v].emplace_back(u);
}
int hash(int u, int p = -1) {
    if (hashes[u] != -1) return hashes[u];
    vi children;
    for (auto v : adj[u])
        if (v != p) children.emplace_back(hash(v, u));
    sort(all(children));
    if (!hasher.count(children)) hasher[children] = hs++;
    return hashes[u] = hasher[children];
}
};
```

# 8.39 Tree diameter (DP)

```
const int MAXN(1 '000' 000);
int N;
vi G[MAXN];
int diameter, toLeaf[MAXN];
void calcDiameter(int u = 0, int p = -1) {
   int d1, d2;
   d1 = d2 = -1;
   for (auto v : G[u]) {
        if (v != p) {
            calcDiameter(v, u);
            d1 = max(d1, toLeaf[v]);
            tie(d1, d2) = minmax({d1, d2});
        }
   }
   toLeaf[u] = d2 + 1;
   diameter = max(diameter, d1 + d2 + 2);
}
```

#### 8.40 Tree edge queries

```
h(n, -1),
      par(n, vector<int>(LOG + 1, root)),
      ed(n, vector < T > (LOG + 1, E)) {
    h[root] = 0, dfs(root, q);
void dfs(int u, const Graph& g) {
    for (auto& [w, v] : g[u]) {
        if (h[v] == -1)
            h[v] = h[u] + 1, par[v][0] = u, ed[v][0] = w;
            for (int k = 0, p; k < LOG; k++) {
                p = par[v][k];
                par[v][k + 1] = par[p][k];
                ed[v][k + 1] = F(ed[v][k], ed[p][k]);
            dfs(v, g);
        }
    }
pair<int, T> up(int u, int dis) {
    T res = E;
    for (int k = 0; k \le LOG; k++) {
        if (dis & (1 << k)) {
            res = F(res, ed[u][k]);
            u = par[u][k];
        }
    return {u, res};
pair<int, T> lca(int u, int v) {
    if (h[u] > h[v]) swap(u, v);
    T res = E:
    tie(v, res) = up(v, h[v] - h[u]);
    if (v == u) return {v, res};
    for (int k = LOG; \sim k; k--) {
        if (par[u][k] != par[v][k]) {
            res = F(res, ed[v][k]):
            res = F(res, ed[u][k]);
            u = par[u][k], v = par[v][k];
        }
    res = F(res, ed[v][0]);
    res = F(res, ed[u][0]);
    return {par[v][0], res};
```

#### 8.41 Virtual Tree

};

```
#pragma once
#include "../Contest/template.cpp"
#include "./Lowest common ancestor (sparse table).cpp"
struct VTree {
   int n;
   LCA lca;
```

```
VTree(const Graph\& q, int root = 0) : n(len(q)), lca(q, root) {}
    pair<vector<tuple<li, int, int>>, int> vtree(vector<int> vs) {
        sort(vs.begin(), vs.end(),
             [&](int u, int v) { return lca.pos[u] < lca.pos[v]; });</pre>
        for (int i = 0, n = size(vs); i + 1 < n; i++) {
            vs.eb(lca.lca(vs[i], vs[i + 1]));
        sort(vs.begin(), vs.end(),
             [&](int u, int v) { return lca.pos[u] < lca.pos[v]; });
        vs.erase(unique(all(vs)), vs.end());
        vi st{vs.front()};
        vector<tuple<ll, int, int>> ret;
        for (int i = 1; i < len(vs); i++) {
            int v = vs[i];
            while (len(st) >= 2 \& lca.lca(v, st.back()) != st.back()) {
                int a = end(st)[-2];
                int b = st.back();
                ll c = lca.dist(a, b);
                ret.eb(c, a, b);
                st.pop back();
            st.pb(v);
        while (len(st) \geq 2) {
            int a = end(st)[-2]:
            int b = st.back();
            ll c = lca.dist(a, b);
            ret.eb(c, a, b);
            st.pop back();
        return {ret, st.back()};
};
```

# 9 Linear Algebra

# 9.1 Matrix (primitive)

```
#include "../Contest/template.cpp"
template <typename T>
struct Matrix {
   int n, m;
   valarray<valarray<T>> v;
   Matrix(int _n, int _m, int id = 0) : n(_n), m(_m), v(valarray<T>(m), n
) {
    if (id) {
       rep(i, 0, n) v[i][i] = 1;
    }
}
valarray<T> &operator[](int x) { return v[x]; }
Matrix transpose() {
    Matrix newv(m, n);
    rep(i, 0, n) rep(j, 0, m) newv[j][i] = (*this)[i][j];
```

```
return newv;
Matrix operator+(Matrix &b) {
    Matrix ret(*this);
    return ret.v += b.v;
Matrix & operator += (Matrix &b) { return v += b.v; }
Matrix operator*(Matrix b) {
    Matrix ret(n, b.m);
    rep(i, 0, n) rep(i, 0, m) rep(k, 0, b.m) ret[i][k] +=
        v[i][j] * b.v[j][k];
    return ret:
Matrix & operator *= (Matrix b) { return *this = *this * b; }
Matrix power(ll exp) {
    Matrix in = *this;
    Matrix ret(n, n, 1);
    while (exp) {
        if (exp & 1) ret *= in:
        in *= in:
        \exp >>= 1;
    return ret;
 * Alters current matrix.
 * Does gaussian elimination and puts matrix in
 * upper echelon form (possibly reduced).
 * Returns the determinant of the square matrix
 * with side equal to the number of rows of the
 * original matrix.
T gaussjordanize(int reduced = 0) {
    T \det = T(1);
    int line = 0:
    rep(col, 0, m) {
        int pivot = line:
        while (pivot < n && v[pivot][col] == T(0)) pivot++;
        if (pivot >= n) continue;
        swap(v[line], v[pivot]);
        if (line != pivot) det *= T(-1);
        det *= v[line][line]:
        v[line] /= T(v[line][col]);
        if (reduced) rep(i, 0, line) {
                v[i] = T(v[i][col]) * v[line];
        rep(i, line + 1, n) { v[i] -= T(v[i][col]) * v[line]; }
        line++:
    return det * (line == n);
```

```
* Needs to be called in a square matrix that
* represents a system of linear equations. Returns {possible solution
* number of solutions (2 if infinite solutions)}
pair<vector<T>, int> solve system(vector<T> results) {
    Matrix aux(n, m + 1);
    rep(i, 0, n) {
        rep(i, 0, m) aux[i][i] = v[i][i];
        aux[i][m] = results[i]:
   T det = aux.gaussjordanize(1);
    int ret = 1 + (det == T(0));
    int n = results.size():
    rrep(i, n - 1, 0 - 1) {
        int pivot = 0;
        while (pivot < n && aux[i][pivot] == T(0)) pivot++;</pre>
        if (pivot == n) {
            if (aux[i][m] != T(0)) ret = 0;
        } else
            swap(aux[i], aux[pivot]);
    rrep(i, n - 1, 0 - 1) rep(j, i + 1, n) aux[i][m] =
        aux[i][i] * aux[i][m];
    rep(i, 0, n) results[i] = aux[i][m];
    rep(i, 0, n) rep(j, 0, m) v[i][j] = aux[i][j];
    return {results, ret}:
/* Does not alter current matrix. Returns {inverse matrix, is curent
* matrix invertable} */
pair<Matrix<T>, bool> find inverse() {
    int n = v.size();
    Matrix<T> aug(n, 2 * n);
    rep(i, 0, n) rep(i, 0, n) aug[i][i] = v[i][i];
    rep(i, 0, n) aug[i][n + i] = 1;
   T det = aug.gaussjordanize(1);
   Matrix<T> ret(n, n):
    rep(i, 0, n) ret[i] = valarray < T > (aug[i][slice(n, n, 1)]);
    return {ret, det != T(0)};
/* Returns rank of matrix. does not alter it. */
int get rank() const {
   if (m == 0) return 0;
   Matrix<T> aux(*this);
    aux.gaussjordanize();
    int resp = 0:
    rep(i, 0, n) resp += (aux[i] != valarray<T>(m)).sum();
    return resp;
```

```
};
    Math
10.1 Arithmetic Progression Sum
Usage:
  • s: first term
  • d: common difference
  • n: number of terms
ll arithmeticProgressionSum(ll s, ll d, ll n) {
    return (s + (s + d * (n - 1))) * n / 2ll;
10.2 Binomial
Time: O(N \cdot K)
Memory: O(K)
ll binom(ll n, ll k) {
    if (k > n) return 0;
    vll dp(k + 1, 0);
    dp[0] = 1;
    for (ll i = 1; i <= n; i++)
        for (ll i = k; i > 0; i--) dp[i] = dp[i] + dp[i - 1];
    return dp[k];
}
10.3 Binomial MOD
Description: find \binom{n}{k} (mod MOD)
  • precompute: on first call it takes O(MAXNBIN) to precompute the factorials
  • query: O(1).
Memory: O(MAXNBIN)
Warning: Remember to set MAXNBIN properly!
const ll MOD = 998244353:
inline ll binom(ll n, ll k) {
    static const int BINMAX = 2'000'000;
    static vll FAC(BINMAX + 1), FINV(BINMAX + 1);
    static bool done = false;
    if (!done) {
        vll INV(BINMAX + 1);
        FAC[0] = FAC[1] = INV[1] = FINV[0] = FINV[1] = 1;
        for (int i = 2; i <= BINMAX; i++) {</pre>
            FAC[i] = FAC[i - 1] * i % MOD;
            INV[i] = MOD - MOD / i * INV[MOD % i] % MOD;
```

FINV[i] = FINV[i - 1] \* INV[i] % MOD;

done = true:

```
if (n < k \mid | n < 0 \mid | k < 0) return 0;
    return FAC[n] * FINV[k] % MOD * FINV[n - k] % MOD;
10.4 Chinese Remainder Theorem
Description: Find the solution X to the N modular equations.
                                 x \equiv a_1(modm_1)
                                x \equiv a_n(modm_n)
The m_i don't need to be coprime, if there is no solution then it returns -1.
Warning: Make sure all M_i are coprime!
#include "../Contest/template.cpp"
tuple<ll, ll, ll> ext gcd(ll a, ll b) {
    if (!a) return {b, 0, 1};
    auto [g, x, y] = ext_gcd(b % a, a);
    return \{q, y - b / a * x, x\};
template <typename T = ll>
struct crt {
    T a, m;
    crt() : a(0), m(1) {}
    crt(T a_, T m_) : a(a_), m(m_) {}
    crt operator*(crt C) {
        auto [g, x, y] = ext gcd(m, C.m);
        if ((a - C.a) \% q != 0) a = -1;
        if (a == -1 \text{ or } C.a == -1) \text{ return } crt(-1, 0);
        T lcm = m / g * C.m;
        T ans = a + (x * (C.a - a) / g % (C.m / g)) * m;
        return crt((ans % lcm + lcm) % lcm, lcm);
};
template <typename T = ll>
struct Congruence {
    T a, m;
};
template <typename T = ll>
T chinese remainder theorem(const vector<Congruence<T>> &equations) {
    crt<T> ans:
```

for (auto &[a , m ] : equations) {

return ans.a;

ans = ans \* crt<T>(a , m );

(1)

}

# 10.5 Derangement / Matching Problem

```
D_N = N! \left(1 - \frac{1}{1!} + \frac{1}{2!} - \frac{1}{3!} + \dots + (-1)^N \frac{1}{N!}\right)
Time: O(N)
#warning Remember to call precompute !
const l\bar{l} MOD = le9 + 7:
const int MAXN(1 '000' 000);
ll fats[MAXN + 1];
void precompute() {
    fats[0] = 1;
    for (ll i = 1; i <= MAXN; i++) {
         fats[i] = (fats[i - 1] * i) % MOD;
}
ll fastpow(ll a, ll p, ll m) {
    ll ret = 1:
    while (p) {
         if (p & 1) ret = (ret * a) % MOD;
         p >>= 1;
         a = (a * a) % MOD;
    return ret;
ll divmod(ll a, ll b) { return (a * fastpow(b, MOD - 2, MOD)) % MOD; }
ll derangement(const ll n) {
    ll \ ans = fats[n]:
    for (ll i = 1; i <= n; i++) {
         ll k = divmod(fats[n], fats[i]);
         if (i & 1) {
             ans = (ans - k + MOD) % MOD;
         } else {
             ans = (ans + k) % MOD;
    return ans;
```

**Description**: Computes the derangement of N, which is given by the formula:

#### 10.6 Euler Phi

**Description**: Computes the number of positive integers less than N that are coprimes with N, in  $O(\sqrt{N})$ .

```
int phi(int n) {
    if (n == 1) return 1;
    auto fs = factorization(n);  // a vctor of pair or a map
    auto res = n;
    for (auto [p, k] : fs) {
        res /= p;
        res *= (p - 1);
    }
    return res;
}
```

#### 10.7 Euler phi (in range)

**Description**: Computes the number of positive integers less than n that are coprimes with N, in the range [1, N], in  $O(N \log N)$ .

# 10.8 Extended Euclidian algorithm

**Description:** Finds the gcd between a and b and x and y such that ax + by = g**Time:**  $O(\log min(a,b))$ 

Warning: If a = b = 0 then there is infity solutions, but 0 is returned. Be careful about overflow.

```
#pragma once
#include "../Contest/template.cpp"
template <typename T>
tuple<T, T, T> extGcd(T a, T b) {
    if (!b) return {a, 1, 0};
    auto [d, x1, y1] = extGcd(b, a % b);
    T x = y1, y = x1 - y1 * (a / b);
    return {d, x, y};
}
```

#### 10.9 FFT convolution and exponentiation

```
const ld PI = acos(-1);
/* change the ld to doulbe may increase
 * performance =D */
struct num {
    ld a{0.0}, b{0.0};
    num() {}
    num(ld na) : a{na} {}
```

```
num(ld na, ld nb) : a{na}, b{nb} {}
   const num operator+(const num &c) const { return num(a + c.a, b + c.b)
   const num operator-(const num &c) const { return num(a - c.a. b - c.b)
    const num operator*(const num &c) const {
        return num(a * c.a - b * c.b, a * c.b + b * c.a);
    const num operator/(const ll &c) const { return num(a / c, b / c); }
};
void fft(vector<num> &a, bool invert) {
   int n = len(a);
   for (int i = 1, j = 0; i < n; i++) {
        int bit = n \gg 1;
        for (; j & bit; bit >>= 1) j ^= bit;
        i ^= bit;
        if (i < j) swap(a[i], a[j]);</pre>
   for (int sz = 2; sz <= n; sz <<= 1) {
        ld ang = 2 * PI / sz * (invert ? -1 : 1);
        num wsz(cos(ang), sin(ang));
        for (int i = 0; i < n; i += sz) {
            num w(1);
            rep(j, 0, sz / 2)  {
                num u = a[i + j], v = a[i + j + sz / 2] * w;
                a[i + i] = u + v;
                a[i + j + sz / 2] = u - v;
                W = W * WSZ:
   if (invert)
        for (num \&x : a) x = x / n:
vi conv(vi const a, vi const b) {
   vector<num> fa(all(a));
   vector<num> fb(all(b));
   int n = 1;
    while (n < len(a) + len(b)) n <<= 1;
   fa.resize(n);
   fb.resize(n):
   fft(fa, false);
   fft(fb, false);
    rep(i, 0, n) fa[i] = fa[i] * fb[i];
   fft(fa, true):
   vi result(n);
   rep(i, 0, n) result[i] = round(fa[i].a);
   while (len(result) and result.back() == 0) result.pop back();
   /* Unconment this line if you want a boolean
    * convolution*/
   // for (auto &xi : result) xi = min(xi, 1ll);
   return result:
vll poly exp(vll &ps, int k) {
   vll ret(len(ps));
```

```
auto base = ps;
ret[0] = 1;
while (k) {
    if (k & 1) ret = conv(ret, base);
    k >>= 1;
    base = conv(base, base);
}
return ret;
}
```

#### 10.10 Factorial Factorization

**Description**: Computes the factorization of N! in  $\varphi(N) * \log N$ **Time**:  $O(\varphi(N) \cdot \log N)$ 

```
ll E(ll n, ll p) {
    ll k = 0, b = p;
    while (b <= n) {
        k += n / b;
        b *= p;
    }
    return k;
}
map<ll, ll> factorial_factorization(ll n, const vll &primes) {
        map<ll, ll> fs;
        for (const auto &p : primes) {
             if (p > n) break;
             fs[p] = E(n, p);
        }
        return fs;
}
```

#### 10.11 Factorization

**Description**: Computes the factorization of N.

Time:  $O(\sqrt{n})$ .

```
map<ll, ll> factorization(ll n) {
    map<ll, ll> ans;
    for (ll i = 2; i * i <= n; i++) {
        ll count = 0;
        for (; n % i == 0; count++, n /= i);
        if (count) ans[i] = count;
    }
    if (n > 1) ans[n]++;
    return ans;
}
```

# 10.12 Factorization (Pollard's Rho)

**Description**: Factorizes a number into its prime factors.

**Time**:  $O(N^{(\frac{1}{4})} * \log(N))$ .

```
ll mul(ll a, ll b, ll m) {
    ll ret = a * b - (ll)((ld)1 / m * a * b + 0.5) * m;
    return ret < 0 ? ret + m : ret;
ll pow(ll a, ll b, ll m) {
   ll\ ans = 1:
    for (; b > 0; b /= 2ll, a = mul(a, a, m)) {
        if (b % 2ll == 1) ans = mul(ans, a, m);
    return ans;
bool prime(ll n) {
   if (n < 2) return 0:
    if (n <= 3) return 1;
    if (n % 2 == 0) return 0;
    ll r = builtin ctzll(n - 1), d = n >> r;
    for (int a: {2, 325, 9375, 28178, 450775, 9780504, 795265022}) {
        ll x = pow(a, d, n);
        if (x == 1 \text{ or } x == n - 1 \text{ or a } % n == 0) continue;
        for (int j = 0; j < r - 1; j++) {
            x = mul(x, x, n);
            if (x == n - 1) break;
        if (x != n - 1) return 0;
    return 1;
ll rho(ll n) {
    if (n == 1 or prime(n)) return n;
    auto f = [n](ll x) \{ return mul(x, x, n) + 1; \};
    ll x = 0, y = 0, t = 30, prd = 2, x0 = 1, q;
    while (t % 40 != 0 or qcd(prd, n) == 1) {
        if (x == y) x = ++x0, y = f(x);
        q = mul(prd, abs(x - y), n);
        if (q != 0) prd = q;
        x = f(x), y = f(f(y)), t++;
    return gcd(prd, n);
vector<ll> fact(ll n) {
    if (n == 1) return {};
    if (prime(n)) return {n};
    ll d = rho(n);
    vector<ll> l = fact(d), r = fact(n / d);
    l.insert(l.end(), r.begin(), r.end());
    return l:
10.13 Fast Pow
```

**Description:** Computes  $a^b \pmod{m}$ **Time:**  $O(\log B)$ .

```
ll fpow(ll a, ll b, ll m) {
```

```
ll ret = 1;
while (b) {
    if (b & 1) ret = (ret * a) % m;
    b >>= 1;
    a = (a * a) % m;
}
return ret;
}
```

#### 10.14 Find linear recurrence (Berlekamp-Massey)

**Description**: Given the first N terms of a linear recurrence finds the smallest recurrence that matches the sequence.

Time:  $O(N^2)$ 

**Warning**: Works faster if the *mod* is const but can be also be a parameter.

Absolute magic!

```
const ll mod = 998244353;
ll modpow(ll b, ll e) {
    ll ans = 1:
    for (; e; b = b * b % mod, e /= 2)
        if (e & 1) ans = ans * b % mod;
    return ans:
vl berlekampMassey(vll s) {
    int n = len(s), L = 0, m = 0;
    if (!n) return {};
    vll C(n), B(n), T;
    C[0] = B[0] = 1;
    ll b = 1;
    rep(i, 0, n) {
        ++m:
        ll d = s[i] \% mod;
        rep(j, 1, L + 1) d = (d + C[j] * s[i - j]) % mod;
        if (!d) continue;
        T = C;
        ll coef = d * modpow(b, mod - 2) % mod;
        rep(j, m, n) C[j] = (C[j] - coef * B[j - m]) % mod;
        if (2 * L > i) continue;
        L = i + 1 - L;
        B = T;
        b = d;
        m = 0:
    C.resize(L + 1):
    C.erase(C.begin());
    for (ll \&x : C) x = (mod - x) % mod;
    return C:
```

#### 10.15 Find multiplicatinve inverse

```
ll inv(ll a, ll m) { return a > 1ll ? m - inv(m % a, a) * m / a : 1ll; }
```

#### 10.16 Floor division

```
template <typename T1, typename T2>
constexpr typename std::common_type<T1, T2>::type floor_div(T1 x, T2 y) {
    assert(y != 0);
    if (y < 0) x = -x, y = -y;
    return x < 0 ? (x - y + 1) / y : x / y;
}</pre>
```

#### 10.17 GCD

```
template <typename T>
T gcd(T a, T b) {
    return b ? gcd(b, a % b) : a;
}
```

#### 10.18 Gauss XOR elimination / XOR-SAT

**Description:** Execute gaussian elimination with xor over the system Ax = b in. The add method must receive a bitset indicating which variables are present in the equation, and the solution of the equation.

Time:  $O(\frac{nm^2}{64})$ 

```
const int MAXXI = 2009;
using Equation = bitset<MAXXI>;
struct GaussXor {
   vector<char> B;
   vector<Equation> A:
   void add(const Equation &ai, bool bi) {
        A.push back(ai):
        B.push back(bi):
   pair<bool, Equation> solution() {
        int cnt = 0. n = A.size():
        Equation vis:
        vis.set();
        Equation x;
        for (int j = MAXXI - 1, i; j >= 0; j--) {
            for (i = cnt; i < n; i++) {
                if (A[i][j]) break;
            if (i == n) continue;
            swap(A[i], A[cnt]), swap(B[i], B[cnt]);
            i = cnt++;
            vis[j] = 0;
            for (int k = 0; k < n; k++) {
                if (i == k || !A[k][j]) continue;
                A[k] ^= A[i]:
                B[k] ^= B[i];
            }
        x = vis:
```

```
for (int i = 0; i < n; i++) {
    int acum = 0;
    for (int j = 0; j < MAXXI; j++) {
        if (!A[i][j]) continue;
        if (!vis[j]) {
            vis[j] = 1;
            x[j] = acum ^ B[i];
        }
        acum ^= x[j];
    }
    if (acum != B[i]) return {false, Equation()};
}
return {true, x};
}</pre>
```

#### 10.19 Guess K-th (Berlekamp-Massey)

```
/* Berlekamp-Massey algorithm
 * Given the first n terms of a linear recurrence relation, this algorithm
 * finds the shortest linear recurrence relation that generates the given
 * sequence.
  Note: mod needs to have inverse
 * Time complexity: O(n^2)
template <typename T>
vector<T> berlekamp massey(const vector<T> &s) {
    vector<T> cur, best;
    int lf, ld;
    for (int i = 0; i < (int)s.size(); i++) {
        T delta = 0;
        for (int j = 0; j < (int)cur.size(); j++)</pre>
            delta += s[i - j - 1] * cur[j];
        if (delta == s[i]) continue;
        if (cur.empty()) {
            cur.resize(i + 1);
            lf = i;
            ld = (int)(delta - s[i]).value();
            continue:
        T coef = -(s[i] - delta) / ld;
        vector<T> c(i - lf - 1);
        c.push back(coef);
        for (auto &x : best) c.push back(-x * coef);
        if (c.size() < cur.size()) c.resize(cur.size());</pre>
        for (int j = 0; j < (int)cur.size(); j++) c[j] += cur[j];</pre>
        if (i - lf + (int)best.size() >= (int)cur.size())
            best = cur, lf = i, ld = (int)(delta - s[i]).value();
        cur = c;
    return cur;
template <typename T>
T get kth(const vector<T> &rec, const vector<T> &dp, ll k) {
```

```
int n = (int)rec.size():
    assert(rec.size() <= dp.size());
    // use fft to speed up
    auto mul = [&](const vector<T> &a, const vector<T> &b) {
        vector<T> res(2 * n);
        for (int i = 0; i < n; i++)
            for (int j = 0; j < n; j++) res[i + j] += a[i] * b[j];
        for (int i = 2 * n - 1; i >= n; i--)
            for (int j = 1; j \le n; j++) res[i - j] += res[i] * rec[j - j]
    11;
        res.resize(n);
        return res:
    vector<T> a(n), x(n);
    x[0] = 1;
    if (n != 1)
        a[1] = 1;
        a[0] = rec[0]:
    while (k) {
        if (k \& 1) x = mul(x, a);
        a = mul(a, a):
        k >>= 1:
    T res = 0:
    for (int i = 0; i < n; i++) res += x[i] * dp[i];
    return res;
template <typename T>
T guess kth term(const vector<T> &s, ll k) {
    if (k < (int)s.size()) return s[k];</pre>
    auto coef = berlekamp massey(s);
    if (coef.empty()) return 0;
    return get kth(coef, s, k);
```

#### 10.20 Integer partition

**Description**: Find the total of ways to partition a given number N in such way that none of the parts is greater than K.

```
Time: O(N \cdot min(N, K))
```

Memory: O(N)

**Warning:** Remember to memset everything to -1 before using it

```
const ll MOD = 10000000007;
const int MAXN(100);
ll memo[MAXN + 1];
ll dp(ll n, ll k = oo) {
    if (n == 0) return 1;
    ll &ans = memo[n];
    if (ans != -1) return ans;
    ans = 0;
    for (int i = 1; i <= min(n, k); i++) {
        ans = (ans + dp(n - i, k)) % MOD;
    }
    return ans;</pre>
```

#### 10.21 LCM

}

```
ll gcd(ll a, ll b) { return b ? gcd(b, a % b) : a; }
ll lcm(ll a, ll b) { return a / gcd(a, b) * b; }
```

#### 10.22 Linear Recurrence

**Description**: Find the n-th term of a linear recurrence, given the recurrence rec and the first K values of the recurrence, remember that first\_k[i] is the value of f(i), considering 0-indexing.

**Usage**: Suppose you want the N-th term of Fibonacci the first k should be 1,1, and the rec should be 0.1,1.1.

Time:  $O(K^3 \log N)$ 

```
template <typename T>
vector<vector<T>> prod(vector<vector<T>> &a, vector<vector<T>> &b,
                       const ll mod) -
    assert(a.back().size() == b.size());
    int n = a.size();
    int m = a.back().size();
    vector<vector<T>> c(n, vector<T>(m));
    for (int i = 0: i < n: i++) {
        for (int i = 0; i < m; i++) {
            for (int k = 0; k < m; k++) {
                c[i][j] = (c[i][j] + ((a[i][k] * b[k][j]) % mod)) % mod;
    return c;
template <tvpename T>
vector<vector<T>> fpow(vector<vector<T>> &xs, ll p, ll mod) {
    vector<vector<T>> ans(xs.size(), vector<T>(xs.size()));
    for (int i = 0; i < (int)xs.size(); i++) ans[i][i] = 1;</pre>
    for (auto b = xs; p; p >>= 1, b = prod(b, b, mod))
        if (p \& 1) ans = prod(ans, b, mod);
    return ans:
ll linear reg(vector<vector<ll>> rec, vector<ll> first k, ll n, const ll
   mod) {
    int k = first k.size();
    if (n <= k) return first k[n - 1];</pre>
    ll n2 = n - k:
    rec = fpow(rec, n2, mod);
    ll ret = 0:
    for (int i = 0; i < k; i++)
        ret = (ret + (rec.back()[i] * first_k[i]) % mod) % mod;
    return ret;
```

#### 10.23 Linear diophantine equation (count)

Description:

```
Time: O(\log min(a,b))
#pragma once
#include "../Contest/template.cpp"
#include "./Extended Euclidian algorithm.cpp"
#include "./Linear diophantine equation (solve).cpp"
template <typename T>
T countSolutionsInRange(T a, T b, T c, T minX, T maxX, T minY, T maxY) {
   auto ss = [\&](T \&x, T \&y, T a, T b, T cnt) { x += cnt * b, y -= cnt *}
   assert(a and b):
   auto sol = diophantineEquationSolution(a, b, c);
   if (!sol) return 0;
   auto [x, y] = *sol;
   auto g = get<0>(extGcd(a, b));
   a /= q:
   b /= a:
   int signA = a > 0 ? +1 : -1;
   int signB = b > 0 ? +1 : -1;
   ss(x, y, a, b, (minX - x) / b);
   if (x < minX) ss(x, y, a, b, signB);
   if (x > maxX) return 0;
   int lx1 = x;
    ss(x, y, a, b, (maxX - x) / b);
   if (x > maxX) ss(x, y, a, b, -signB);
   int rx1 = x:
    ss(x, y, a, b, -(minY - y) / a);
   if (y < minY) ss(x, y, a, b, -signA);
   if (y > maxY) return 0;
   int lx2 = x:
   ss(x, y, a, b, -(maxY - y) / a);
   if (y > maxY) ss(x, y, a, b, signA);
   int rx2 = x;
   if (lx2 > rx2) swap(lx2, rx2);
   int lx = max(lx1, lx2);
   int rx = min(rx1, rx2);
   if (lx > rx) return 0;
   return (rx - lx) / abs(b) + 1;
```

# 10.24 Linear diophantine equation (solve)

**Description:** Finds a solution for ax + by = c, where a, b, c, are given and x and y unknown.

Time:  $O(\log min(a, b))$ 

```
#pragma once
#include "../Contest/template.cpp"
#include "./Extended Euclidian algorithm.cpp"
```

```
template <typename T>
optional<pair<T, T>> diophantineEquationSolution(T a, T b, T c) {
    if (a == 0 \text{ and } b == 0) {
        if (c)
            return nullopt;
        else
            return pair<T, T>{(T)0, (T)0};
    }
    auto [g, x0, y0] = extGcd(a < 0 ? a * -1 : a, b < 0 ? b * -1 : b);
    if (c % g) return nullopt;
    x0 *= c / q, y0 *= c / q;
    if (a < 0) x0 = -x0:
    if (b < 0) v0 = -v0;
    pair<T, T> ret;
    ret.first = x0, ret.second = y0;
    return ret;
```

#### 10.25 List N elements choose K

**Description**: Process every possible combination of K elements from N elements, thoose index marked as 1 in the index vector says which elements are choosed at that moment.

Time:  $O(\binom{N}{K} \cdot O(process))$ 

```
void process(vi &index) {
    for (int i = 0; i < len(index); i++) {
        if (index[i]) cout << i << " \n"[i == len(index) - 1];
    }
}
void n_choose_k(int n, in k) {
    vi index(n);
    fill(index.end() - k, index.end(), 1);
    do {
        process(index);
    } while (next_permutation(all(index)));
}</pre>
```

# 10.26 List primes (Sieve of Eratosthenes)

```
const ll MAXN = 2e5;
vll list_primes(ll n = MAXN) {
    vll ps;
    bitset<MAXN + 1> sieve;
    sieve.set();
    sieve.reset(1);
    for (ll i = 2; i <= n; ++i) {
        if (sieve[i]) ps.push_back(i);
        for (ll j = i * 2; j <= n; j += i) {
            sieve.reset(j);
        }
    }
}</pre>
```

```
return ps;
}
```

# 10.27 Matrix exponentiation

```
const ll MOD = 1 '000' 000'007;
template <typename T>
vector<vector<T>> prod(vector<vector<T>> &a, vector<vector<T>> &b) {
   int n = len(a):
   vector<vector<T>> c(n, vector<T>(n));
   for (int i = 0; i < n; i++) {
        for (int j = 0; j < n; j++) {
            for (int k = 0; k < n; k++) {
                c[i][j] = (c[i][j] + ((a[i][k] * b[k][j]) % MOD)) % MOD;
       }
   }
    return c;
template <typename T>
vector<vector<T>> fpow(vector<vector<T>> &xs, ll p) {
   vector<vector<T>> ans(len(xs), vector<T>(len(xs)));
   for (int i = 0; i < len(xs); i++) ans[i][i] = 1;
   auto b = xs;
   while (p) {
        if (p \& 1) ans = prod(ans, b);
        p >>= 1;
        b = prod(b, b);
    return ans;
}
```

#### 10.28 NTT integer convolution and exponentiation

#### Time:

- Convolution  $O(N \cdot \log N)$ ,
- Exponentiation:  $O(\log K \cdot N \cdot \log N)$

```
mint(ll v ) {
        if (v_ >= _mod or v_ <= -_mod) v_ %= _mod;
        if (v_{-} < 0) v_{-} += \underline{mod}; v_{-} += \underline{mod};
    m &operator+=(const m &a) {
        v += a.v;
        if (v \ge mod) v = mod;
        return *this:
    m &operator-=(const m &a) {
        \dot{v} -= a.v:
        if (v < 0) v += mod;
        return *this;
    m &operator*=(const m &a) {
        v = v * ll(a.v) % _mod;
        return *this;
    m &operator/=(const m &a) {
        v = v * inv(a.v) % mod;
        return *this;
    m operator-() { return m(-v); }
    m & operator^=(ll e) {
        if (e < 0) {
            v = inv(v):
             e = -e;
        v = expo(v, e);
        // possivel otimizacao:
        // cuidado com 0^0
        // v = \exp(v, e^{(p-1)});
        return *this;
    bool operator==(const m &a) { return v == a.v; }
    bool operator!=(const m &a) { return v != a.v; }
    friend istream &operator>>(istream &in, m &a) {
        ll val:
        in >> val;
        a = m(val);
        return in;
    friend ostream &operator<<(ostream &out, m a) { return out << a.v; }</pre>
    friend m operator+(m a, m b) { return a += b; }
    friend m operator-(m a, m b) { return a -= b;
    friend m operator*(m a, m b) { return a *= b; }
    friend m operator/(m a, m b) { return a /= b; }
    friend m operator^(m a, ll e) { return a ^= e; }
};
const ll MOD1 = 998244353;
const ll MOD2 = 754974721;
const ll MOD3 = 167772161;
template <int mod>
void ntt(vector<mint<_mod>> &a, bool rev) {
    int n = len(a);
```

```
auto b = a:
   assert(!(n \& (n - 1)));
   mint < mod > g = 1;
   while ((g \land (mod / 2)) == 1) g += 1;
   if (rev) q = 1 / q;
   for (int step = n / 2; step; step /= 2) {
        mint < mod > w = q ^ (mod / (n / step)), wn = 1;
        for (int i = 0; i < n / 2; i += step) {
            for (int j = 0; j < step; j++) {
                auto u = a[2 * i + j], v = wn * a[2 * i + j + step];
                b[i + j] = u + v;
                b[i + n / 2 + i] = u - v:
            \dot{w}n = wn * w;
        swap(a, b);
   if (rev) {
        auto n1 = mint< mod>(1) / n;
        for (auto &x : \overline{a}) x *= n1;
template <ll mod>
vector<mint< mod>> convolution(const vector<mint< mod>> &a,
                                const vector<mint< mod>> &b) {
   vector<mint< mod>> l(all(a)), r(all(b));
   int N = len(\overline{l}) + len(r) - 1, n = 1;
   while (n \le N) n *= 2;
   l.resize(n), r.resize(n);
   ntt(l, false), ntt(r, false);
   for (int i = 0; i < n; i++) l[i] *= r[i];
   ntt(l, true);
   l.resize(N);
   // Uncommnent for a boolean convolution :)
    for (auto& li : l) {
     li.v = min(li.v, 1ll);
    return l;
template <ll mod>
vector<mint< mod>> poly exp(vector<mint< mod>> &ps, int k) {
   vector<mint< mod>> ret(len(ps));
   auto base = \overline{ps}:
   ret[0] = 1;
   while (k) {
        if (k & 1) ret = convolution(ret, base);
        k >>= 1:
        base = convolution(base, base);
    return ret:
```

# 10.29 NTT integer convolution and exponentiation (2 mods) modules)

**Description**: Computes the convolution between the two polynomials and.

Time:  $O(N \log N)$ 

Warning: This is pure magic!

```
template <int mod>
struct mint {
    ll expo(ll b, ll e) {
        ll ret = 1;
        while (e) {
            if (e % 2) ret = ret * b % mod;
            e /= 2, b = b * b % mod;
        return ret;
    ll inv(ll b) { return expo(b, mod - 2); }
    usina m = mint:
    ll v:
    mint() : v(0) {}
    mint(ll v ) {
        if (v_ >= _mod or v_ <= -_mod) v_ %= _mod;
        if (v_{-} < 0) v_{-} += _{mod};
    m &operator+=(const m &a) {
        \dot{v} += a.v:
        if (v >= mod) v -= mod;
        return *this;
    m &operator-=(const m &a) {
        \dot{v} -= a.v;
        if (v < 0) v += mod;
        return *this;
    m &operator*=(const m &a) {
        v = v * ll(a.v) % mod;
        return *this;
    m &operator/=(const m &a) {
        v = v * inv(a.v) % mod;
        return *this:
    m operator-() { return m(-v); }
    m &operator^=(ll e) {
        if (e < 0) {
            v = inv(v);
            e = -e;
        v = expo(v, e);
        // possivel otimizacao:
        // cuidado com 0^0
        // v = \exp(v, e^{(p-1)});
        return *this;
```

```
bool operator==(const m &a) { return v == a.v; }
    bool operator!=(const m &a) { return v != a.v; }
    friend istream &operator>>(istream &in, m &a) {
        ll val:
        in >> val;
        a = m(val);
        return in;
    friend ostream &operator<<(ostream &out, m a) { return out << a.v; }</pre>
    friend m operator+(m a, m b) { return a += b; }
    friend m operator-(m a, m b) { return a -= b; }
    friend m operator*(m a, m b) { return a *= b; }
    friend m operator/(m a, m b) { return a /= b; }
    friend m operator^(m a, ll e) { return a ^= e; }
};
const ll MOD1 = 998244353;
const ll MOD2 = 754974721;
const ll MOD3 = 167772161;
template <int mod>
void ntt(vector<mint<_mod>> &a, bool rev) {
    int n = len(a):
    auto b = a;
    assert(!(n \& (n - 1)));
    mint < mod > q = 1;
    while ((q ^ (mod / 2)) == 1) q += 1;
   if (rev) q = 1 / q;
    for (int step = n / 2; step; step /= 2) {
        mint < mod > w = g ^ (mod / (n / step)), wn = 1;
        for (int i = 0; i < n / 2; i += step) {
            for (int j = 0; j < step; j++) {
                auto u = a[2 * i + j], v = wn * a[2 * i + j + step];
                b[i + j] = u + v;
                b[i + n / 2 + i] = u - v;
            \dot{w}n = wn * w;
        swap(a, b);
    if (rev) {
        auto n1 = mint<_mod>(1) / n;
        for (auto \&x : \overline{a}) x *= n1;
tuple<ll, ll, ll> ext gcd(ll a, ll b) {
   if (!a) return {b, 0, 1};
    auto [g, x, y] = ext gcd(b % a, a);
    return \{g, y - b / a * x, x\};
template <typename T = ll>
struct crt {
   T a, m;
    crt() : a(0), m(1) {}
    crt(T a_, T m_) : a(a_), m(m_) {}
    crt operator*(crt C) {
```

```
auto [q, x, y] = ext qcd(m, C.m);
        if ((a - C.a) % q != 0) a = -1;
        if (a == -1 \text{ or } C.a == -1) \text{ return } crt(-1, 0);
        T lcm = m / g * C.m;
        T ans = a + (x * (C.a - a) / g % (C.m / g)) * m;
        return crt((ans % lcm + lcm) % lcm, lcm);
};
template <typename T = ll>
struct Congruence {
    T a, m;
};
template <typename T = ll>
T chinese remainder_theorem(const vector<Congruence<T>> &equations) {
    crt<T> ans:
    for (auto &[a_, m_] : equations) {
        ans = ans * crt<T>(a , m );
    return ans.a;
#define int long long
template <ll m1, ll m2>
vll merge_two_mods(const vector<mint<ml>> &a, const vector<mint<m2>> &b) {
    int n = len(a):
    vll ans(n);
    for (int i = 0; i < n; i++) {
        auto cur = crt<ll>():
        auto ai = a[i].v;
        auto bi = b[i].v;
        cur = cur * crt<ll>(ai, m1);
        cur = cur * crt < ll > (bi, m2);
        ans[i] = cur.a:
    }
    return ans;
vll convolution 2mods(const vll &a, const vll &b) {
    vector<mint<MOD1>> l(all(a)), r(all(b));
    int N = len(l) + len(r) - 1, n = 1;
    while (n \le N) n *= 2;
    l.resize(n), r.resize(n);
    ntt(l, false), ntt(r, false);
    for (int i = 0; i < n; i++) l[i] *= r[i];</pre>
    ntt(l, true);
    l.resize(N);
    vector<mint<MOD2>> l2(all(a)), r2(all(b));
    l2.resize(n), r2.resize(n);
    ntt(l2, false), ntt(r2, false);
    rep(i, 0, n) l2[i] *= r2[i];
    ntt(l2, true);
    l2.resize(N);
    return merge_two_mods(l, l2);
vll poly exp(const vll &xs, ll k) {
    vll ret(len(xs));
```

```
ret[0] = 1;
auto base = xs;
while (k) {
    if (k & 1) ret = convolution_2mods(ret, base);
    k >>= 1;
    base = convolution_2mods(base, base);
}
return ret;
```

#### 10.30 Polynomial Taylor Shift

```
using C = complex<double>;
const ll mod = 998244353;
void fft(vector<C> &a) {
   int n = len(a), L = 31 - \_builtin\_clz(n);
   static vector<complex<long double>> R(2, 1);
    static vector<C> rt(2, 1);
   for (static int k = 2; k < n; k *= 2) {
        R.resize(n);
        rt.resize(n);
        auto x = polar(1.0L, acos(-1.0L) / k);
        for (int i = k; i < 2 * k; i++)
            rt[i] = R[i] = i \& 1 ? R[i / 2] * x : R[i / 2];
   vector<int> rev(n):
   for (int i = 0; i < n; i++) rev[i] = (rev[i / 2] | (i \& 1) << L) / 2;
   for (int i = 0; i < n; i++)
        if (i < rev[i]) swap(a[i], a[rev[i]]);</pre>
   for (int k = 1; k < n; k *= 2) {
        for (int i = 0; i < n; i += 2 * k)
            for (int j = 0; j < k; j++) {
                auto x = (double *)&rt[j + k], y = (double *)&a[i + j + k]
   ];
                C z(x[0] * y[0] - x[1] * y[1], x[0] * y[1] + x[1] * y[0]);
                a[i + j + k] = a[i + j] - z;
                a[i + j] += z;
            }
   }
vector<double> conv(const vector<double> &a, const vector<double> &b) {
   if (a.empty() || b.empty()) return {};
   vector<double> res(len(a) + len(b) - 1);
   int L = 32 - builtin clz(len(res)), n = 1 \ll L;
   vector<C> in(\overline{n}), out(\overline{n});
    copy(a.begin(), a.end(), begin(in));
   for (int i = 0; i < len(b); i++) in[i].imag(b[i]);
   fft(in):
    for (C &x : in) x *= x;
    for (int i = 0; i < n; i++) {
        out[i] = in[-i \& (n - 1)] - conj(in[i]);
   fft(out);
   for (int i = 0; i < len(res); i++) {
```

```
res[i] = imag(out[i]) / (4 * n);
    return res;
template <ll M>
vector<ll> convMod(const vector<ll> &a, const vector<ll> &b) {
    if (a.empty() || b.empty()) return {};
    vector<ll> res(len(a) + len(b) + 1);
    int B = 32 - builtin clz(len(res)), n = 1 \ll B, cut = int(sqrt(M));
    vector<C> L(n), R(n), outs(n), outl(n);
    for (int i = 0; i < len(a); i++) {
        L[i] = C((int)a[i] / cut, (int)a[i] % cut);
    for (int i = 0; i < len(b); i++) {
        R[i] = C((int)b[i] / cut, (int)b[i] % cut);
    fft(L), fft(R);
    for (int i = 0; i < n; i++) {
        int j = -i \& (n - 1);
        outl[j] = (L[i] + conj(L[j])) * R[i] / (2.0 * n);
        outs[j] = (L[i] - conj(L[j])) * R[i] / (2.0 * n) / 1i;
    fft(outl), fft(outs);
    for (int i = 0; i < len(res); i++) {
        ll av = ll(real(outl[i]) + .5), cv = ll(imag(outs[i]) + .5);
        ll bv = ll(imag(outl[i]) + .5) + ll(real(outs[i]) + .5);
        res[i] = ((av % M * cut + bv) % M * cut + cv) % M;
    }
    return res;
ll fexp(ll b, ll e) {
    ll res = 1;
    while (e > 0) {
        if (e & 1) res = res * b % mod;
        b = b * b % mod;
        e >>= 1:
    return res;
ll inv(ll n) { return fexp(n, mod - 2); }
vector<ll> shift(vector<ll> &a, ll v) {
    int n = len(a) - 1;
    vector<ll> f(n + 1), g(n + 1);
    vector<ll> i fact(n + 1);
    f[0] = a[0];
    q[n] = 1;
    i fact[0] = 1;
    l\bar{l} fact = 1, potk = 1;
    for (int i = 1; i < n + 1; i++) {
        fact = fact * i % mod;
        f[i] = fact * a[i] % mod;
        potk = (potk * v % mod + mod) % mod;
        q[n-i] = ((potk * inv(fact)) % mod + mod) % mod;
        i fact[i] = inv(fact);
    auto p = convMod<mod>(f, g);
```

```
vector<ll> res(n + 1);
for (int i = 0; i < n + 1; i++) {
    res[i] = (p[i + n] * i_fact[i] % mod + mod) % mod;
}
return res;
}</pre>
```

# 10.31 Polyominoes

**Usage**: buildPolyominoes(x) creates every polyomino until size x, and put it in polyominoes[x], access polyomino.v to find the vector of pairs representing the coordinates of each piece, considering that the polyomino was 'rooted' in coordinate (0,0).

Warning: note that when accessing polyominoes[x] only the first x coordinates are valid.

```
#include "../Contest/template.cpp"
const int MAXP = 10;
using pii = pair<int, int>;
// This implementation considers the rotations as
// distinct
                 0, 10, 10+9, 10+9+8...
int pos[11] = \{0, 10, 19, 27, 34, 40, 45, 49, 52, 54, 55\};
struct Polyominoes {
    pii v[MAXP];
   ll id:
   int n;
   Polyominoes() {
        n = 1;
        v[0] = \{0, 0\};
        normalize();
   pii &operator[](int i) { return v[i]; }
   bool add(int a, int b) {
        for (int i = 0; i < n; i++)
            if (v[i].first == a and v[i].second == b) return false;
        v[n++] = pii(a, b);
        normalize();
        return true:
   void normalize() {
        int mnx = 100, mny = 100;
        for (int i = 0; i < n; i++)
            mnx = min(mnx, v[i].first), mny = min(mny, v[i].second);
        id = 0:
        for (int i = 0; i < n; i++) {
            v[i].first -= mnx, v[i].second -= mny;
            id |= (1LL << (pos[v[i].first] + v[i].second));</pre>
   }
vector<Polyominoes> polyominoes[MAXP + 1];
void buildPolyominoes(int mxN = 10) {
   vector<pair<int, int>> dt({{1, 0}, {-1, 0}, {0, -1}, {0, 1}});
    for (int i = 0; i <= mxN; i++) polyominoes[i].clear();</pre>
   Polyominoes init;
   queue<Polyominoes> q;
```

```
unordered set<int64 t> used;
g.push(init);
used.insert(init.id);
while (!q.empty()) {
    Polyominoes u = q.front();
    q.pop();
    polyominoes[u.n].push back(u);
    if (u n == mxN) continue;
    for (int i = 0; i < u.n; i++) {
        for (auto [dx, dy] : dt) {
            Polyominoes to = u;
            bool ok = to.add(to[i].first + dx, to[i].second + dy);
            if (ok and !used.count(to.id)) {
                q.push(to);
                used.insert(to.id);
   }
```

#### 11 Primitives

#### 11.1 Bigint

```
const int maxn = 1e2 + 14, lg = 15;
const int base = 10000000000:
const int base digits = 9;
struct bigint ₹
    vi a;
    int sign;
    int size() {
        if (a.empty()) return 0;
        int ans = (a.size() - 1) * base digits;
        int ca = a.back();
        while (ca) ans++, ca /= 10;
        return ans;
    bigint operator^(const bigint &v) {
        bigint ans = 1, a = *this, b = v;
        while (!b.isZero()) {
            if (b \% 2) ans *= a;
            a *= a, b /= 2;
        return ans;
    string to_string() {
        stringstream ss;
        ss << *this;
        string s;
        ss >> s;
        return s;
    int sumof() {
        string s = to string();
```

```
int ans = 0:
    for (auto c : s) ans += c - '0';
    return ans;
/*</arpa>*/
bigint() : sign(1) {}
bigint(long long v) { *this = v; }
bigint(const string &s) { read(s); }
void operator=(const bigint &v) {
    sign = v.sign;
    a = v.a:
void operator=(long long v) {
    sign = 1;
    a.clear();
    if (v < 0) sign = -1, v = -v;
    for (; v > 0; v = v / base) a.push back(v % base);
bigint operator+(const bigint &v) const {
    if (sign == v.sign) {
        bigint res = v;
        for (int i = 0, carry = 0;
             i < (int)max(a.size(), v.a.size()) || carry; ++i) {</pre>
            if (i == (int)res.a.size()) res.a.push back(0);
            res.a[i] += carry + (i < (int)a.size() ? a[i] : 0);
            carry = res.a[i] >= base;
            if (carry) res.a[i] -= base;
        return res;
    return *this - (-v);
bigint operator-(const bigint &v) const {
    if (sign == v.sign) {
        if (abs() >= v.abs()) {
            bigint res = *this;
            for (int i = 0, carry = 0; i < (int)v.a.size() || carry;
++i) {
                res.a[i] \rightarrow carry + (i < (int)v.a.size() ? v.a[i] : 0)
                carrv = res.a[i] < 0:
                if (carry) res.a[i] += base;
            res.trim():
            return res;
        return -(v - *this);
    return *this + (-v);
void operator*=(int v) {
    if (v < 0) sign = -sign, v = -v;
    for (int i = 0, carry = 0; i < (int)a.size() || carry; ++i) {
        if (i == (int)a.size()) a.push back(0);
```

```
long long cur = a[i] * (long long)v + carry;
        carry = (int)(cur / base);
        a[i] = (int)(cur % base);
        // asm("divl %%ecx" : "=a"(carry),
        // "=d"(a[i]) : "A"(cur), "c"(base));
    trim();
bigint operator*(int v) const {
    bigint res = *this:
    res *= v;
    return res;
void operator*=(long long v) {
    if (v < 0) sign = -sign, v = -v;
    if (v > base) {
        *this = *this * (v / base) * base + *this * (v % base);
    for (int i = 0, carry = 0; i < (int)a.size() || carry; ++i) {</pre>
        if (i == (int)a.size()) a.push back(0);
        long long cur = a[i] * (long long)v + carry;
        carry = (int)(cur / base);
        a[i] = (int)(cur % base);
        // asm("divl %%ecx" : "=a"(carry),
        // "=d"(a[i]) : "A"(cur), "c"(base));
    trim():
bigint operator*(long long v) const {
    bigint res = *this;
    res *= v;
    return res:
friend pair<br/>
bigint, bigint> divmod(const bigint &a1, const bigint &b1)
    int norm = base / (b1.a.back() + 1);
    bigint a = a1.abs() * norm;
    bigint b = b1.abs() * norm;
    biaint a, r:
    q.a.resize(a.a.size());
    for (int i = a.a.size() - 1; i \ge 0; i--) {
        r *= base;
        r += a.a[i];
        int s1 = r.a.size() <= b.a.size() ? 0 : r.a[b.a.size()];</pre>
        int s2 = r.a.size() \le b.a.size() - 1 ? 0 : r.a[b.a.size() -
1];
        int d = ((long long)base * s1 + s2) / b.a.back();
        r -= b * d:
        while (r < 0) r += b, --d;
        q.a[i] = d;
    q.sign = a1.sign * b1.sign;
    r.sign = a1.sign;
    q.trim();
```

```
r.trim();
    return make pair(q, r / norm);
bigint operator/(const bigint &v) const { return divmod(*this, v).
bigint operator%(const bigint &v) const { return divmod(*this, v).
second; }
void operator/=(int v) {
    if (v < 0) sign = -sign, v = -v;
    for (int i = (int)a.size() - 1, rem = 0; i >= 0; --i) {
        long long cur = a[i] + rem * (long long)base;
        a[i] = (int)(cur / v);
        rem = (int)(cur % v);
    trim();
bigint operator/(int v) const {
    bigint res = *this;
    res /= v;
    return res;
int operator%(int v) const {
    if (v < 0) v = -v;
    int m = 0:
    for (int i = a.size() - 1; i >= 0; --i)
        m = (a[i] + m * (long long)base) % v;
    return m * sign;
void operator+=(const bigint &v) { *this = *this + v; }
void operator==(const bigint &v) { *this = *this - v; }
void operator*=(const bigint &v) { *this = *this * v; }
void operator/=(const bigint &v) { *this = *this / v; }
bool operator<(const bigint &v) const {</pre>
    if (sign != v.sign) return sign < v.sign;</pre>
    if (a.size() != v.a.size())
        return a.size() * sign < v.a.size() * v.sign;</pre>
    for (int i = a.size() - 1; i >= 0; i--)
        if (a[i] != v.a[i]) return a[i] * sign < v.a[i] * sign;</pre>
    return false:
bool operator>(const bigint &v) const { return v < *this; }</pre>
bool operator<=(const bigint &v) const { return !(v < *this); }</pre>
bool operator>=(const bigint &v) const { return !(*this < v); }</pre>
bool operator==(const bigint &v) const {
    return !(*this < v) && !(v < *this);
bool operator!=(const bigint &v) const { return *this < v || v < *this
; }
void trim() {
    while (!a.empty() && !a.back()) a.pop back();
    if (a.empty()) sign = 1;
bool isZero() const { return a.empty() || (a.size() == 1 \&\& !a[0]); }
```

```
bigint operator-() const {
    bigint res = *this:
    res.sign = -sign;
    return res;
bigint abs() const {
    bigint res = *this;
    res.sian *= res.sian:
    return res;
long longValue() const {
    long long res = 0;
    for (int i = a.size() - 1; i \ge 0; i--) res = res * base + a[i];
    return res * sign:
friend bigint gcd(const bigint &a, const bigint &b) {
    return b.isZero() ? a : gcd(b, a % b);
friend bigint lcm(const bigint &a, const bigint &b) {
    return a / \gcd(a, b) * b;
void read(const string &s) {
    sign = 1;
    a.clear();
    int pos = 0;
    while (pos < (int)s.size() && (s[pos] == '-' || s[pos] == '+')) {
        if (s[pos] == '-') sign = -sign:
        ++pos;
    for (int i = s.size() - 1; i \ge pos; i -= base digits) {
        int x = 0:
        for (int j = max(pos, i - base digits + 1); j <= i; j++)
            x = x * 10 + s[i] - '0';
        a.push back(x);
    trim();
friend istream &operator>>(istream &stream, bigint &v) {
    string s:
    stream >> s:
    v.read(s);
    return stream;
friend ostream &operator<<(ostream &stream, const bigint &v) {</pre>
    if (v.sign == -1) stream << '-';
    stream << (v.a.empty() ? 0 : v.a.back());</pre>
    for (int i = (int)v.a.size() - 2; i >= 0; --i)
        stream << setw(base digits) << setfill('0') << v.a[i];</pre>
    return stream:
static vector<int> convert_base(const vector<int> &a, int old digits,
                                 int new digits) {
    vector<long long> p(max(old digits, new digits) + 1);
    p[0] = 1;
```

```
for (int i = 1; i < (int)p.size(); i++)p[i] = p[i-1] * 10;
    vector<int> res:
    long long cur = 0;
    int cur digits = 0;
    for (int i = 0; i < (int)a.size(); i++) {
        cur += a[i] * p[cur digits];
        cur digits += old digits;
        while (cur digits >= new digits) {
            res.push back(int(cur % p[new digits]));
            cur /= p[new digits]:
            cur digits -= new_digits;
        }
    res.push back((int)cur);
    while (!res.empty() && !res.back()) res.pop back();
    return res;
typedef vector<long long> vll;
static vll karatsubaMultiply(const vll &a, const vll &b) {
    int n = a.size();
    vll res(n + n):
    if (n <= 32) {
        for (int i = 0; i < n; i++)
            for (int i = 0; i < n; i++) res[i + i] += a[i] * b[i];
        return res:
    }
    int k = n \gg 1;
    vll a1(a.begin(), a.begin() + k);
    vll a2(a.begin() + k, a.end());
    vll b1(b.begin(), b.begin() + k);
    vll b2(b.beain() + k.b.end()):
    vll alb1 = karatsubaMultiply(a1, b1);
    vll a2b2 = karatsubaMultiply(a2, b2);
    for (int i = 0; i < k; i++) a2[i] += a1[i];
    for (int i = 0; i < k; i++) b2[i] += b1[i];
    vll r = karatsubaMultiply(a2, b2);
    for (int i = 0; i < (int)alb1.size(); i++) r[i] -= alb1[i];
    for (int i = 0; i < (int)a2b2.size(); i++) r[i] -= a2b2[i];
    for (int i = 0; i < (int)r.size(); i++) res[i + k] += r[i];
    for (int i = 0; i < (int)alb1.size(); i++) res[i] += alb1[i];
    for (int i = 0; i < (int)a2b2.size(); i++) res[i + n] += a2b2[i];
    return res;
bigint operator*(const bigint &v) const {
    vector<int> a6 = convert base(this->a, base digits, 6);
    vector<int> b6 = convert base(v.a, base digits, 6);
    vll a(a6.begin(), a6.end\overline{()});
    vll b(b6.begin(), b6.end());
    while (a.size() < b.size()) a.push_back(0);</pre>
    while (b.size() < a.size()) b.push back(0);</pre>
    while (a.size() \& (a.size() - 1)) \overline{a.push back(0)}, b.push back(0);
    vll c = karatsubaMultiply(a, b);
    bigint res;
```

```
res.sign = sign * v.sign;
for (int i = 0, carry = 0; i < (int)c.size(); i++) {
    long long cur = c[i] + carry;
    res.a.push_back((int)(cur % 1000000));
    carry = (int)(cur / 1000000);
}
res.a = convert_base(res.a, 6, base_digits);
res.trim();
return res;
}
};</pre>
```

#### 11.2 Integer Mod

```
#include "../Contest/template.cpp"
template <ll m>
struct mod int {
    ll x;
    mod int(ll v = 0)  {
        x = v % m;
        if (x < 0) v += m;
    mod int &operator+=(mod int const &b) {
        x += b.x;
        if (x \ge m) x = m:
        return *this;
    mod int &operator=(mod int const &b) {
        x \rightarrow b.x:
        if (x < 0) x += m:
        return *this;
    mod int &operator*=(mod int const &b) {
        x = (ll)x * b.x % m;
        return *this;
    friend mod int mpow(mod int a, ll e) {
        mod int res = 1:
        while (e) {
            if (e & 1) res *= a;
            a ∗= a;
            e >>= 1:
        return res;
    friend mod int inverse(mod int a) { return mpow(a, m - 2); }
    mod int &operator/=(mod int const &b) { return *this *= inverse(b); }
    friend mod int operator+(mod int a, mod int const b) { return a += b;
    mod_int operator++(int) { return this->x = (this->x + 1) % m; }
    mod int operator++() { return this->x = (this->x + 1) % m; }
    friend mod int operator-(mod int a, mod int const b) { return a -= b;
```

```
friend mod_int operator-(mod_int const a) { return 0 - a; }
mod_int operator--(int) { return this->x = (this->x - 1 + m) % m; }
mod_int operator--() { return this->x = (this->x - 1 + m) % m; }
friend mod_int operator*(mod_int a, mod_int const b) { return a *= b; }
friend mod_int operator/(mod_int a, mod_int const b) { return a /= b; }
friend ostream & operator<<(ostream & os, mod_int const & a) {
    return os << a.x; }
friend bool operator==(mod_int const & a, mod_int const & b) {
    return a.x == b.x; }
friend bool operator!=(mod_int const & a, mod_int const & b) {
    return a.x != b.x; }
};</pre>
```

## 11.3 Integer Mod (complete)

```
#include "../Contest/template.cpp"
template <ll Mod>
struct modint {
    static constexpr ll mod = Mod;
   ll v;
   modint() : v(0) \{ \}
   template <ll Mod2>
   modint(const modint<Mod2> &x) : v(x.value()) {}
   modint(ll x) : v(safe_mod(x)) {}
   ll &value() { return \overline{v}; }
   const ll &value() const { return v; }
   static ll safe mod(ll x) {
        return x \ge 0 ? x \% \mod : ((x \% \mod) + \mod) \% \mod;
   template <typename T>
   explicit operator T() const {
        return (T)v;
    bool operator==(const modint rhs) const noexcept { return v == rhs.v;
    bool operator!=(const modint rhs) const noexcept { return v != rhs.v;
   bool operator<(const modint rhs) const noexcept { return v < rhs.v; }</pre>
   bool operator<=(const modint rhs) const noexcept { return v <= rhs.v;</pre>
   bool operator>(const modint rhs) const noexcept { return v > rhs.v; }
   bool operator>=(const modint rhs) const noexcept { return v >= rhs.v;
   modint operator++(int) {
        modint res = *this:
        *this += 1;
        return res;
   modint operator--(int) {
```

```
modint res = *this:
    *this -= 1:
    return res;
modint &operator++() { return *this += 1; }
modint &operator--() { return *this -= 1; }
modint operator+() const { return modint(*this); }
modint operator-() const { return mod - modint(*this): }
friend modint operator+(const modint lhs, const modint rhs) noexcept {
    return modint(lhs) += rhs;
friend modint operator-(const modint lhs, const modint rhs) noexcept {
    return modint(lhs) -= rhs:
friend modint operator*(const modint lhs, const modint rhs) noexcept {
    return modint(lhs) *= rhs;
friend modint operator/(const modint lhs, const modint rhs) noexcept {
    return modint(lhs) /= rhs;
modint &operator+=(const modint rhs) {
    v += rhs.v:
    if (v \ge mod) v = mod;
    return *this;
modint & operator = (const modint rhs) {
    if (v < rhs.v) v += mod:
    v = rhs.v:
    return *this;
modint & operator *= (const modint rhs) {
    v = v * rhs.v % mod;
    return *this;
modint &operator/=(modint rhs) { return *this *= rhs.inv(); }
modint pow(ll p) const {
    static assert(mod < static cast<ll>(1) << 32,
                  "Modulus must be less than 2**32");
    modint res = 1, a = *this;
    while (p) {
        if (p \& 1) res *= a;
        a *= a;
        p >>= 1;
    return res;
modint inv() const { return pow(mod - 2); }
modint sqrt() const {
    modint b = 1;
    while (b.pow((mod - 1) >> 1) == 1) b += 1;
    ll m = mod - 1. e = 0:
    while (\sim m \& 1) m >>= 1. e++:
    auto x = pow((m - 1) >> 1);
    auto v = *this * x * x:
    x *= *this;
    auto z = b.pow(m);
    while (y != 1) {
```

```
ll i = 0:
             for (modint t = y; t != 1; t *= t, ++j);
            z.pow(111 << (e - j - 1));
            X *= Z;
Z *= Z;
            V *= Z;
            e = i;
        return x;
    friend ostream &operator<<(ostream &s, const modint &x) {</pre>
        s << x.value():
        return s:
    friend istream &operator>>(istream &s, modint &x) {
        ll value:
        s >> value:
        x = \{value\};
        return s;
};
```

#### 11.4 Matrix

```
template <tvpename T>
struct Matrix {
   vector<vector<T>> d;
   Matrix() : Matrix(0) {}
   Matrix(int n) : Matrix(n, n) {}
   Matrix(int n, int m) : Matrix(vector<vector<T>>(n, vector<T>(m))) {}
   Matrix(const vector<vector<T>> &v) : d(v) {}
   constexpr int n() const { return (int)d.size(); }
   constexpr int m() const { return n() ? (int)d[0].size() : 0; }
   void rotate() { *this = rotated(); }
   Matrix<T> rotated() const {
       Matrix<T> res(m(), n());
       for (int i = 0; i < m(); i++) {
            for (int j = 0; j < n(); j++) {
                res[i][i] = d[n() - i - 1][i];
        return res;
   Matrix<T> pow(int power) const {
       assert(n() == m()):
       auto res = Matrix<T>::identity(n());
       auto b = *this;
       while (power) {
            if (power & 1) res *= b;
            b *= b:
            power >>= 1;
        return res:
```

```
Matrix<T> submatrix(int start i, int start j, int rows = INT MAX,
                    int cols = INT MAX) const {
    rows = min(rows, n() - start i);
    cols = min(cols, m() - start i):
    if (rows \leq 0 or cols \leq 0) return \{\}:
    Matrix<T> res(rows, cols);
    for (int i = 0: i < rows: i++)
        for (int j = 0; j < cols; j++)
            res[i][i] = d[i + start i][j + start_j];
    return res;
}
Matrix<T> translated(int x, int y) const {
    Matrix<T> res(n(), m());
    for (int i = 0; i < n(); i++) {
        for (int j = 0; j < m(); j++) {
            if (i + x < 0 \text{ or } i + x >= n() \text{ or } j + y < 0 \text{ or } j + y >= m()
                 continue:
            res[i + x][j + y] = d[i][j];
    return res;
static Matrix<T> identity(int n) {
    Matrix<T> res(n):
    for (int i = 0; i < n; i++) res[i][i] = 1;
    return res;
vector<T> &operator[](int i) { return d[i]; }
const vector<T> &operator[](int i) const { return d[i]; }
Matrix<T> & operator += (T value) {
    for (auto &row : d) {
        for (auto &x : row) x += value;
    return *this:
Matrix<T> operator+(T value) const {
    auto res = *this;
    for (auto &row : res) {
        for (auto &x : row) x = x + value;
    return res;
Matrix<T> &operator==(T value) {
    for (auto &row : d) {
        for (auto &x : row) x -= value;
    return *this:
Matrix<T> operator-(T value) const {
    auto res = *this;
    for (auto &row : res) {
        for (auto \&x : row) x = x - value:
    return res:
```

```
Matrix<T> &operator*=(T value) {
    for (auto &row : d) {
        for (auto &x : row) x *= value;
    return *this:
Matrix<T> operator*(T value) const {
    auto res = *this:
    for (auto &row : res) {
        for (auto &x : row) x = x * value;
    return res;
Matrix<T> &operator/=(T value) {
    for (auto &row : d) {
        for (auto &x : row) x /= value;
    return *this;
Matrix<T> operator/(T value) const {
    auto res = *this;
    for (auto &row : res) {
        for (auto &x : row) x = x / value;
    return res;
Matrix<T> & operator += (const Matrix<T> &o) {
    assert(n() == o.n() and m() == o.m());
    for (int i = 0; i < n(); i++) {
        for (int i = 0; i < m(); i++) {
            d[i][i] += o[i][i];
    return *this:
Matrix<T> operator+(const Matrix<T> &o) const {
    assert(n() == o.n() and m() == o.m());
    auto res = *this;
    for (int i = 0; i < n(); i++) {
        for (int j = 0; j < m(); j++) {
            res[i][j] = res[i][j] + o[i][j];
    return res;
Matrix<T> &operator=(const Matrix<T> &o) {
    assert(n() == o.n() and m() == o.m()):
    for (int i = 0; i < n(); i++) {
        for (int j = 0; j < m(); j++) {
            d[i][i] -= o[i][i];
    return *this;
Matrix<T> operator-(const Matrix<T> &o) const {
    assert(n() == o.n() and m() == o.m());
    auto res = *this;
```

```
for (int i = 0; i < n(); i++) {
            for (int i = 0; i < m(); i++) {
                res[i][j] = res[i][j] - o[i][j];
        return res;
    Matrix<T> & operator*=(const Matrix<T> &o) {
        *this = *this * 0:
        return *this:
    Matrix<T> operator*(const Matrix<T> &o) const {
        assert(m() == o.n()):
        Matrix<T> res(n(), o.m());
        for (int i = 0; i < res.n(); i++) {
            for (int j = 0; j < res.m(); j++) {
                auto &x = res[i][i];
                for (int k = 0; k < m(); k++) {
                    x += (d[i][k] * o[k][i]):
            }
        return res:
    friend istream &operator>>(istream &is. Matrix<T> &mat) {
        for (auto &row : mat)
            for (auto &x : row) is >> x;
        return is:
    friend ostream &operator<<(ostream &os, const Matrix<T> &mat) {
        bool frow = 1:
        for (auto &row : mat) {
            if (not frow) os << '\n';</pre>
            bool first = 1:
            for (auto &x : row) {
                if (not first) os << ' ':
                os << x:
                first = 0:
            frow = 0:
        return os;
    auto begin() { return d.begin(); }
    auto end() { return d.end(); }
    auto rbegin() { return d.rbegin(); }
    auto rend() { return d.rend(); }
    auto begin() const { return d.begin(); }
    auto end() const { return d.end(); }
    auto rbegin() const { return d.rbegin(); }
    auto rend() const { return d.rend(); }
};
```

## 12 Problems

## 12.1 2081 - Fixed-Lenght Paths II

```
#include <bits/stdc++.h>
using namespace std;
const int MAXN = 2'00'000;
int N, K1, K2;
vector<int> ADJ[MAXN];
int64 t ans = 0;
int sz[MAXN], removed[MAXN];
void calcSize(int u, int p = -1) {
    sz[u] = 1:
    for (int v : ADJ[u]) {
        if (v != p and !removed[v]) {
            calcSize(v, u);
            sz[u] += sz[v];
   }
}
int findCentroid(int u, int mxSz, int p = -1) {
    for (int v : ADJ[u]) {
        if (!removed[v] \text{ and } v != p \text{ and } sz[v] * 2 >= mxSz)
            return findCentroid(v, mxSz, u);
    return u;
int64 t cnt[MAXN], totCnt[MAXN], initialSum;
int mxD;
void dfs(int u, int p, int d) {
   if (d > K2) return;
    cnt[d]++:
    mxD = max(mxD, d);
    if (K1 - 1 \le d \text{ and } d \le K2 - 1) initialSum++;
    for (int v : ADJ[u]) {
        if (v != p and !removed[v]) {
            dfs(v, u, d + 1);
void solve(int curRoot) {
    calcSize(curRoot);
    int centroid = findCentroid(curRoot, sz[curRoot]);
    removed[centroid] = true;
    int totMxD = 0:
    initialSum = (K1 == 1);
    // cerr << "centroid: " << centroid << '\n';</pre>
    for (int v : ADJ[centroid]) {
        if (!removed[v]) {
            // cerr << "v: " << v << '\n';
            m \times D = 0:
            int64 t curSum = initialSum;
            dfs(v. centroid, 1):
            totMxD = max(totMxD, mxD);
```

```
for (int d = 1; d <= mxD; d++) {</pre>
                 // cerr << "d : " << d << " curSum: " << curSum << '\n';
                 ans += (curSum * cnt[d]);
                 int pl = max(0, K1 - d) - 1;
                 if (pl >= 0) curSum += totCnt[pl];
                 int pr = K2 - d:
                 curSum -= totCnt[pr];
            for (int d = 1; d <= mxD; d++) totCnt[d] += cnt[d];</pre>
            fill(\&cnt[1], \&cnt[1] + mxD + 1, 0);
    // cerr << "centroid: " << centroid
    //<< " ans: " << ans << '\n';
    for (int d = 1; d \le totMxD; d++) totCnt[d] = 0;
    for (int v : ADJ[centroid])
        if (!removed[v]) solve(v);
int32_t main() {
    ios base::sync with stdio(!cin.tie(0));
    tot\overline{C}nt[0] = 1;
    cin >> N >> K1 >> K2;
    for (int i = 0; i < N - 1; i++) {
        int u. v:
        cin >> u >> v:
        u--, v--;
        ADJ[u] emplace back(v);
        ADJ[v].emplace_back(u):
    solve(0):
    cout << ans << '\n':
// AC, centroid decomposition
```

## 12.2 Fixed lenght pahts I

```
#include <bits/stdc++.h>
using namespace std;

const int MAXN = 2'00'000;
int N, K;
vector<int> ADJ[MAXN];
int64_t ans;
bool removed[MAXN];
int cnt[MAXN];
int sz[MAXN];
void calcSize(int u, int p = -1) {
    sz[u] = 1;
    for (int v : ADJ[u]) {
        if (v != p and !removed[v]) {
            calcSize(v, u);
            sz[u] += sz[v];
```

```
}
}
int getCentroid(int mxSz, int u, int p = -1) {
    for (int v : ADJ[u]) {
        if (v != p and !removed[v] and sz[v] >= mxSz)
            return getCentroid(mxSz, v, u);
    return u;
}
int mxd:
void dfs(int u, int p, bool upd, int d = 1) {
    if (d > K) return:
    mxd = max(mxd, d);
    upd ? cnt[d]++ : ans += cnt[K - d];
    for (int v : ADJ[u]) {
        if (v != p \text{ and } !removed[v]) dfs(v, u, upd, d + 1);
}
void solve(int u) {
    calcSize(u);
    int c = getCentroid(sz[u] >> 1, u);
    removed[c] = true;
    mxd = 0:
    cnt[0] = 1;
    for (int v : ADJ[c]) {
        if (!removed[v]) {
            dfs(v, c, false);
            dfs(v, c, true);
    for (int i = 0; i \le mxd; i++) cnt[i] = 0;
    for (int v : ADJ[c]) {
        if (!removed[v]) solve(v);
int32 t main() {
    ios base::sync with stdio(0);
    cin.tie(0);
    cin >> N >> K;
    for (int i = 0; i < N - 1; i++) {
        int u, v;
        cin >> u >> v;
        u--, v--;
        ADJ[u].emplace back(v);
        ADJ[v].emplace back(u);
    }
    solve(0);
    cout << ans << '\n';
    return 0;
```

### 12.3 Fixed length paths II

```
#include <bits/stdc++.h>
using namespace std;
const int MAXN = 2'00'000;
int N, K1, K2;
vector<int> ADJ[MAXN];
int64 t ans = 0;
int sz[MAXN], removed[MAXN];
void calcSize(int u, int p = -1) {
    sz[u] = 1:
    for (int v : ADJ[u]) {
        if (v != p and !removed[v]) {
            calcSize(v, u);
            sz[u] += sz[v];
int findCentroid(int u, int mxSz, int p = -1) {
    for (int v : ADJ[u]) {
        if (!removed[v] \text{ and } v != p \text{ and } sz[v] * 2 >= mxSz)
            return findCentroid(v, mxSz, u);
    return u;
int64 t cnt[MAXN], totCnt[MAXN], initialSum;
int mxD:
void dfs(int u, int p, int d) {
    if (d > K2) return;
    cnt[d]++:
    mxD = max(mxD, d):
    if (K1 - 1 \le d \text{ and } d \le K2 - 1) initialSum++;
    for (int v : ADJ[u]) {
        if (v != p and !removed[v]) {
            dfs(v, u, d + 1);
void solve(int curRoot) {
    calcSize(curRoot);
    int centroid = findCentroid(curRoot, sz[curRoot]);
    removed[centroid] = true;
    int totMxD = 0:
    initialSum = (K1 == 1);
    // cerr << "centroid: " << centroid << '\n';</pre>
    for (int v : ADJ[centroid]) {
        if (!removed[v]) {
            // cerr << "v: " << v << '\n':
            mxD = 0:
            int64_t curSum = initialSum;
            dfs(v, centroid, 1);
            totMxD = max(totMxD, mxD);
            for (int d = 1; d <= mxD; d++) {</pre>
                 // cerr << "d : " << d << " curSum: " << curSum << '\n';
```

```
ans += (curSum * cnt[d]);
                int pl = \max(0, K1 - d) - 1;
                if (pl >= 0) curSum += totCnt[pl];
                int pr = K2 - d:
                curSum -= totCnt[pr];
            }
            for (int d = 1: d \le mxD: d++) totCnt[d] += cnt[d]:
            fill(\&cnt[1], \&cnt[1] + mxD + 1, 0):
    // cerr << "centroid: " << centroid
    //<< " ans: " << ans << '\n';
    for (int d = 1; d \le totMxD; d++) totCnt[d] = 0;
    for (int v : ADJ[centroid])
        if (!removed[v]) solve(v);
int32 t main() {
    ios base::sync with stdio(!cin.tie(0));
    totCnt[0] = 1:
    cin >> N >> K1 >> K2:
    for (int i = 0; i < N - 1; i++) {
        int u, v;
        cin >> u >> v;
        u--, v--;
        ADJ[u].emplace back(v);
        ADJ[v].emplace back(u);
    solve(0);
    cout << ans << '\n';
// AC, centroid decomposition
```

# 13 Strings

#### 13.1 Z-Function

## 13.1.1 Z-function building

**Description**: The Z-function is an algorithm used to compute the Z-array of a given string. For a string s, Z[i] represents the length of the longest common prefix between the string s and the suffix of s starting from the index i.

Usage: The function z\_function\_build(s) takes a single argument s, which is a string (or any container-like structure), and returns a vector of integers representing the Z-function of the input.

```
string s = "abacaba";
vector<int> result = z_function_build(s);
// result = [0, 0, 1, 0, 3, 0, 1]
vector<int> v = {1, 2, 3, 1, 2, 3};
vector<int> result_v = z_function_build(v);
// result v = [0, 1, 2, 3, 0, 1]
```

```
Time: O(n)
Memory: O(n)
Warning: By definition Z[0] = 0, remember to treat it appart.
```

```
#pragma once
#include "../../Contest/template.cpp"
template <typename Seq>
vector<int> z_function_build(const Seq& s) {
   int n = int(s.size());
   vector<int> z(n);
   for (int i = 1, l = 0, r = 0; i < n; ++i) {
      if (i <= r) z[i] = min(r - i + 1, z[i - l]);
      while (i + z[i] < n && s[z[i]] == s[i + z[i]]) ++z[i];
      if (i + z[i] - 1 > r) l = i, r = i + z[i] - 1;
   }
   return z;
}
```

## 13.2 Count distinct anagrams

```
const ll\ MOD = 1e9 + 7;
const int maxn = 1e6;
vll fs(maxn + 1):
void precompute() {
    fs[0] = 1;
    for (ll i = 1; i <= maxn; i++) {
        fs[i] = (fs[i - 1] * i) % MOD;
ll fpow(ll a, int n, ll mod = LLONG MAX) {
    if (n == 0) return 1:
    if (n == 1) return a:
    ll x = fpow(a, n / 2, mod) % mod;
    return ((x * x) % mod * (n & 1 ? a : 111)) % mod;
ll distinctAnagrams(const string &s) {
    precompute();
    vi hist('z' - 'a' + 1, 0);
    for (auto &c : s) hist[c - 'a']++;
    ll ans = fs[len(s)];
    for (auto &q : hist) {
        ans = (ans * fpow(fs[q], MOD - 2, MOD)) % MOD;
    return ans;
```

## 13.3 Double hash range query

```
#include "../Contest/template.cpp"
using ll = long long;
using vll = vector<ll>;
using pll = pair<ll, ll>;
```

```
const int MAXN(1'000'000);
const ll MOD = 1000027957;
const ll MOD2 = 1000015187;
const ll P = 31;
ll p[MAXN + 1], p2[MAXN + 1];
void precompute() {
    p[0] = p2[0] = 1;
    for (int i = 1; i <= MAXN; i++)
        p[i] = (P * p[i - 1]) % MOD, p2[i] = (P * p2[i - 1]) % MOD2;
struct Hash {
   int n;
   vll h, h2, hi, hi2;
   Hash() {}
   Hash(const string \&s) : n(s.size()), h(n), h2(n), hi(n), hi2(n) {
        h[0] = h2[0] = s[0];
        for (int i = 1; i < n; i++)
            h[i] = (s[i] + h[i - 1] * P) % MOD,
            h2[i] = (s[i] + h2[i - 1] * P) % MOD2;
        hi[n - 1] = hi2[n - 1] = s[n - 1]:
        for (int i = n - 2; i >= 0; i--)
            hi[i] = (s[i] + hi[i + 1] * P) % MOD,
            hi2[i] = (s[i] + hi2[i + 1] * P) % MOD2;
   pll query(int l, int r) {
        ll\ hash = (h[r] - (l?h[l-1]*p[r-l+1]%MOD:0));
        ll\ hash2 = (h2[r] - (l? h2[l-1] * p2[r-l+1] % MOD2 : 0));
        return {(hash < 0 ? hash + MOD : hash),</pre>
                (hash2 < 0 ? hash2 + MOD2 : hash2);
   pll query inv(int l, int r) {
        ll\ ha\bar{s}h = (hi[l] - (r + 1 < n ? hi[r + 1] * p[r - l + 1] % MOD :
   0));
        ll\ hash2 =
            (hi2[l] - (r + 1 < n ? hi2[r + 1] * p2[r - l + 1] % MOD2 : 0))
        return {(hash < 0 ? hash + MOD : hash),</pre>
                (hash2 < 0 ? hash2 + MOD2 : hash2);
};
```

## 13.4 Hash range query

```
#include "../Contest/template.cpp"
const ll P = 31;
const ll MOD = 1e9 + 9;
const int MAXN(1e6);
ll ppow[MAXN + 1];
void pre_calc() {
    ppow[0] = 1;
    for (int i = 1; i <= MAXN; i++) ppow[i] = (ppow[i - 1] * P) % MOD;
}
struct Hash {</pre>
```

```
int n;
    vll h, hi;
    Hash(const string &s) : n(s.size()), h(n), hi(n) {
        h[0] = s[0];
        hi[n-1] = s[n-1];
        for (int i = 1; i < n; i++) {
            h[i] = (s[i] + h[i - 1] * P) % MOD:
            hi[n - i - 1] = (s[n - i - 1] + hi[n - i - 1] * P) % MOD;
    }
    ll gry(int l, int r) {
        ll hash = (h[r] - (l ? h[l - 1] * ppow[r - l + 1] % MOD : 0));
        return hash < 0 ? hash + MOD : hash;</pre>
    ll gry inv(int l, int r) {
        ll^{-} hash = (hi[l] - (r + 1 < n ? hi[r + 1] * ppow[r - l + 1] % MOD
        return hash < 0 ? hash + MOD : hash:
};
```

## 13.5 Hash unsigned long long $2^{64} - 1$

**Description**: Arithmetic mod  $2^{64} - 1$ . 2x slower than mod  $2^{64}$  and more code, but works on evil test data (e.g. Thue-Morse, where ABBA... and BAAB... of length  $2^{10}$  hash the same mod  $2^{64}$ ).

"typedef ull H;" instead if you think test data is random.

```
#include "../Contest/template.cpp"
struct H {
    ull x;
    H(ull x = 0) : x(x) {}
    H operator+(H o) { return x + o.x + (x + o.x < x); }
    H operator-(H o) { return *this + \sim0.x; }
    H operator*(H o) {
        auto m = (uint128 t)x * o.x;
        return H((\overline{ull})m) + (\overline{ull})(m >> 64):
    ull get() const { return x + !\sim x; }
    bool operator==(H o) const { return get() == o.get(); }
    bool operator<(H o) const { return get() < o.get(); }</pre>
};
static const H C = (long long)le11 + 3; // (order ~ 3e9; random also ok)
struct Hash {
    int n;
    vector<H> ha, pw;
    Hash(string \&str) : n(str.size()), ha((int)str.size() + 1), pw(ha) {
        pw[0] = 1:
        for (int i = 0; i < (int)str.size(); i++)</pre>
             ha[i + 1] = ha[i] * C + str[i], pw[i + 1] = pw[i] * C;
    H query(int a, int b) { // hash [a, b]
        return ha[b] - ha[a] * pw[b - a];
```

```
vector<H> getHashes(string &str, int length) {
    if ((int)str.size() < length) return {};</pre>
    H h = 0, pw = 1;
    for (int i = 0; i < length; i++) h = h * C + str[i], pw = pw * C;
    vector<H> ret = {h}:
    for (int i = length; i < (int)str.size(); i++)</pre>
        ret.push back(h = h * C + str[i] - pw * str[i - length]);
    return ret:
}
H hashString(string &s) {
    H h{};
    for (char c : s) h = h * C + c;
    return h:
13.6 K-th digit in digit string
Description: Find the k-th digit in a digit string, only works for 1 \le k \le 10^{18}!
```

**Time**: precompute O(1), query O(1)

```
using vull = vector<ull>:
vull pow10;
vector<array<ull, 4>> memo;
void precompute(int maxpow = 18) {
    ull qtd = 1;
    ull start = 1;
    ull end = 9;
    ull curlenght = 9;
    ull startstr = 1:
    ull endstr = 9:
    for (ull i = 0, j = 111; (int) i < maxpow; i++, j *= 1011) pow10.eb(j);
    for (ull i = 0; i < maxpow - 1ull; i++) {
        memo.push back({start, end, startstr, endstr});
        start = end + 1ll:
        end = end + (9ll * pow10[qtd]);
        curlenght = end - start + 1ull:
        atd++;
        startstr = endstr + 1ull;
        endstr = (endstr + 1ull) + (curlenght)*gtd - 1ull;
char kthDigit(ull k) {
    int qtd = 1;
    for (auto [s, e, ss, es] : memo) {
        if (k \ge ss and k \le ss) {
            ull pos = k - ss:
            ull index = pos / qtd;
            ull nmr = s + index;
            int i = k - ss - qtd * index;
            return ((nmr / pow10[qtd - i - 1]) % 10) + '0';
        }
```

```
atd++:
return 'X';
```

#### 13.7 KMP

```
/**
* Author: Johan Sannemo
 * Date: 2016-12-15
 * License: CC0
 * Description: pi[x] computes the length of the longest prefix of s that
 * at x, other than s[0...x] itself (abacaba \rightarrow 0010123). Can be used to
 * all occurrences of a string. Time: O(n) Status: Tested on
 * kattis:stringmatching
 */
 * @Title: Prefix function - Knuth-Morris-Pratt
 * @Description: Given a string $S$ builds an array $A$ such that
 * $A i$ is the longest suffix that ends in $i$ and is also a prefix
 * of $S$.
*
*/
#pragma once
vi pi(const string& s) {
    vi p(sz(s));
    rep(i, 1, sz(s)) {
        int g = p[i - 1];
        while (g \&\& s[i] != s[g]) g = p[g - 1];
        p[i] = g + (s[i] == s[g]);
    return p;
vi match(const string& s, const string& pat) {
    vi p = pi(pat + '\0' + s), res;
    rep(i, sz(p) - sz(s), sz(p)) if (p[i] == sz(pat))
        res.push back(i - 2 * sz(pat));
    return res:
```

# 13.8 Longest Palindrome Substring (Manacher)

**Description**: Finds the longest palindrome substring, manacher returns a vector where the i-th position is how much is possible to grow the string to the left and the right of i and keep it a palindrome.

Time: O(N)

```
vi manacher(const string &s) {
    int n = len(s) - 2:
    vi p(n + 2);
    int l = 1. r = 1:
```

```
for (int i = 1; i <= n; i++) {
        p[i] = max(0, min(r - i, p[l + (r - i)]));
        while (s[i - p[i]] == s[i + p[i]]) p[i]++;
        if (i + p[i] > r) l = i - p[i], r = i + p[i];
        p[i]--;
    return p;
string longest palindrome(const string &s) {
    string t("$#");
    for (auto c : s) t.push back(c), t.push back('#');
   t.push back('^');
   vi xs = manacher(t);
   int mpos = max element(all(xs)) - xs.begin();
    for (int k = xs[mpos], i = mpos - k; i \le mpos + k; i++)
        if (t[i] != '#') p.push back(t[i]);
    return p;
}
```

## 13.9 Longest palindrome

```
string longest palindrome(const string &s) {
    int n = (int)s.size();
    vector<array<int, 2>> dp(n);
    pii odd(0, -1), even(0, -1);
    pii ans:
    for (int i = 0; i < n; i++) {
        int k = 0;
        if (i > odd.second)
            k = 1;
        else
            k = min(dp[odd.first + odd.second - i][0], odd.second - i + 1)
        while (i - k) = 0 and i + k < n and s[i - k] = s[i + k] + k + k
        dp[i][0] = k--;
        if (i + k > odd.second) odd = \{i - k, i + k\};
        if (2 * dp[i][0] - 1 > ans.second) ans = \{i - k, 2 * dp[i][0] - k\}
    1};
        k = 0:
        if (i <= even.second)</pre>
            k = min(dp[even.first + even.second - i + 1][1],
                    even.second -i + 1:
        while (i - k - 1) = 0 and i + k < n and s[i - k - 1] == s[i + k]
        dp[i][1] = k--;
        if (i + k > even.second) even = \{i - k - 1, i + k\};
        if (2 * dp[i][1] > ans.second) ans = \{i - k - 1, 2 * dp[i][1]\};
    return s.substr(ans.first, ans.second);
}
```

## 13.10 Lyndon factorization

```
vi lyndon_factorization(string S) {
    auto sa = suffix_array(S);
    vi ans;
    vi mex(len(S) + 1, 0);
    int p = 0;
    rtrav(si, sa) {
        if (si == p) {
            ans.eb(si);
        }
        mex[si] = 1;
        while (mex[p]) p++;
    }
    ans.eb(len(S));
    return ans;
}
```

## 13.11 Rabin-Karp

```
size t rabin karp(const string &s, const string &p) {
    if (s.size() < p.size()) return 0;</pre>
    auto n = s.size(), m = p.size();
    const ll p1 = 31, p2 = 29, q1 = 1e9 + 7, q2 = 1e9 + 9;
    const ll p1_1 = fpow(p1, q1 - 2, q1), p1_2 = fpow(p1, m - 1, q1);
    const ll p2 1 = fpow(p2, q2 - 2, q2), p2 2 = fpow(p2, m - 1, q2);
    pair<ll. ll> hs. hp:
    for (int i = (int)m - 1; \simi; --i) {
        hs.first = (hs.first * p1) % q1;
        hs.first = (hs.first + (s[i] - 'a' + 1)) % q1;
        hs.second = (hs.second * p2) % q2;
        hs.second = (hs.second + (s[i] - 'a' + 1)) % q2;
        hp.first = (hp.first * p1) % q1;
        hp.first = (hp.first + (p[i] - 'a' + 1)) % q1;
        hp.second = (hp.second * p2) % q2;
        hp.second = (hp.second + (p[i] - 'a' + 1)) % q2;
    size t occ = 0;
    for \overline{\text{(size t i = 0; i < n - m; i++)}}
        occ += (hs == hp);
        int fi = s[i] - 'a' + 1;
        int fm = s[i + m] - 'a' + 1;
        hs.first = (hs.first - fi + q1) % q1;
        hs.first = (hs.first * p1 1) % q1;
        hs.first = (hs.first + fm * p1 2) % q1;
        hs.second = (hs.second - fi + q2) % q2;
        hs.second = (hs.second * p2 1) % q2;
        hs.second = (hs.second + fm * p2 2) % q2;
    occ += hs == hp;
    return occ:
}
```

#### 13.12 Suffix Automaton

**Description:** A suffix automaton A for a string s is a minimal finite automaton that recognizes the suffixes of s.

```
#pragma once
#include "../Contest/template.cpp"
struct SuffixAutomaton {
    int n;
    vi suffix_link, max_len;
    vi2d transition;
    SuffixAutomaton(const string &s, int alphabet_size='z'-'a'+1, int norm
    = 'a') : n(len(s), suffix_link(n<<1), max_len(n<<1) , transition(n
    <<1, vi(alphabet_size, -1){
        int last = 0;
        trav(c, s) {
            int max_len_cur = max_len[last] + 1;
            while
        }
    }
};</pre>
```

### 13.13 Suffix array

```
#include <bits/stdc++.h>
using namespace std;
#ifdef LOCAL
#include "debug.cpp"
#else
#define dbg(...)
#endif
#define endl '\n'
#define fastio
   ios_base::sync with stdio(0); \'
    cin.tie(0);
#define int long long
#define all(j) j.begin(), j.end()
#define rall(j) j.rbegin(), j.rend()
#define len(j) (int)j.size()
#define rep(i, a, b) \
    for (common type t<decltype(a), decltype(b) > i = (a); i < (b); i++)
#define rrep(i, a, b) \
    for (common type t<decltype(a), decltype(b)> i = (a); i > (b); i--
#define trav(xi, xs) for (auto &xi : xs)
#define rtrav(xi, xs) for (auto &xi : ranges::views::reverse(xs))
#define pb push back
#define pf push front
#define ppb pop back
#define ppf pop front
#define eb emplace back
#define lb lower bound
#define ub upper bound
#define fi first
#define se second
```

```
#define emp emplace
#define ins insert
#define divc(a, b) ((a) + (b) - 1ll) / (b)
using str = string;
using ll = long long;
using ull = unsigned long long;
using ld = long double;
using vll = vector<ll>;
using pll = pair<ll, ll>;
using vll2d = vector<vll>;
using vi = vector<int>;
using vi2d = vector<vi>;
using pii = pair<int, int>;
using vpii = vector<pii>;
using vc = vector<char>;
using vs = vector<str>;
template <typename T, typename T2>
using umap = unordered map<T, T2>;
template <tvpename T>
using pgmn = priority queue<T, vector<T>, greater<T>>;
template <typename T>
using pqmx = priority_queue<T, vector<T>>;
template <typename T, typename U>
inline bool chmax(T &a, U const &b) {
    return (a < b ? a = b, 1 : 0);
template <typename T, typename U>
inline bool chmin(T &a, U const &b) {
    return (a > b ? a = b, 1 : 0);
vector<int> sort cyclic shifts(string const &s) {
    int n = s.size();
    const int alphabet = 128;
    vector<int> p(n), c(n), cnt(max(alphabet, n), 0);
    for (int i = 0; i < n; i++) cnt[s[i]]++;
    for (int i = 1; i < alphabet; i++) cnt[i] += cnt[i - 1];
    for (int i = 0; i < n; i++) p[--cnt[s[i]]] = i;
    c[p[0]] = 0;
    int classes = 1;
    for (int i = 1; i < n; i++) {
        if (s[p[i]] != s[p[i - 1]]) classes++;
        c[p[i]] = classes - 1;
    }
    vector<int> pn(n), cn(n);
    for (int h = 0; (1 << h) < n; ++h) {
        for (int i = 0; i < n; i++) {
            pn[i] = p[i] - (1 << h);
            if (pn[i] < 0) pn[i] += n;
        fill(cnt.begin(), cnt.begin() + classes, 0);
        for (int i = 0; i < n; i++) cnt[c[pn[i]]]++;
        for (int i = 1; i < classes; i++) cnt[i] += cnt[i - 1];
        for (int i = n - 1; i \ge 0; i--) p[--cnt[c[pn[i]]]] = pn[i];
        cn[p[0]] = 0:
        classes = 1:
```

```
for (int i = 1; i < n; i++) {
            pair<int, int> cur = {c[p[i]], c[(p[i] + (1 << h)) % n]};
            pair<int, int> prev = \{c[p[i-1]], c[(p[i-1] + (1 << h)) \%
   n]};
            if (cur != prev) ++classes;
            cn[p[i]] = classes - 1;
        c.swap(cn);
    return p;
vector<int> suffix array(string s) {
    s += "$";
    vector<int> p = sort_cyclic_shifts(s);
    p.erase(p.begin());
    return p;
vector<int> longestCommonPrefix(const string &s, const vector<int> &suf) {
    int n = s.size();
    vector<int> isuf(n), res(n - 1);
    for (int i = 0; i < n; ++i) isuf[suf[i]] = i;
    int k = 0:
    for (; isuf[k] != n - 1; ++k) {
        int cmp i = suf[isuf[k] + 1];
        int r = k == 0 ? 0 : max(res[isuf[k - 1]] - 1, (int)0);
        while (k + r < n \&\& cmp i + r < n \&\& s[k + r] == s[cmp i + r]) ++r
        res[isuf[k]] = r;
    ++k;
    for (int i = k; i < n; ++i) {
        int cmp i = suf[isuf[i] + 1];
        int r = i == k ? 0 : max(res[isuf[i - 1]] - 1, (int)0);
        while (i + r < n \&\& cmp i + r < n \&\& s[i + r] == s[cmp i + r]) ++r
        res[isuf[i]] = r;
    return res;
}
ll distinct substrings(const string &s, const vi &sa) {
   int n = len(s):
    vi lcp = longestCommonPrefix(s, sa);
    ll ans = n - sa[0]:
    rep(i, 1, n) \{ ans += n - sa[i] - lcp[i - 1]; \}
    return ans;
void run();
int32 t main() {
#ifndef LOCAL
    fastio:
#endif
    int T = 1:
    /*cin >> T;*/
    rep(t, 0, T) {
        dbg(t);
```

```
run():
void run() {
    string S;
    cin >> S;
    auto sa = suffix array(S);
    cout << distinct substrings(S, sa) << endl;</pre>
13.14 Suffix array (supreme)
template <typename T = ll,
          auto cmp = [](T \& src1, T \& src2, T \& dst) \{ dst = min(src1, src2); \}
class SparseTable {
   private:
    int sz:
    vi logs;
    vector<vector<T>> st:
   public:
    SparseTable() {}
    SparseTable(const vector<T> &v) : sz(len(v)), logs(sz + 1) {
        rep(i, 2, sz + 1) logs[i] = logs[i >> 1] + 1;
        st.resize(logs[sz] + 1, vector<T>(sz));
        rep(i, 0, sz) st[0][i] = v[i];
        for (int k = 1: (1 << k) <= sz: k++) {
            for (int i = 0; i + (1 << k) <= sz; i++) {
                cmp(st[k-1][i], st[k-1][i+(1 << (k-1))], st[k][i])
    T query(int l, int r) {
        const int k = logs[r - l];
        cmp(st[k][l], st[k][r - (1 << k)], ret);
        return ret:
};
template <typename T>
using RMQ = SparseTable<T, [](T &a, T &b, T &c) { c = min(a, b); }>;
// éCrditos: ShahjalalShohaq
// O(N)
struct SA {
    string s;
    int n:
    vector<int> sa, lcp, pos;
    RMQ<int> rmq;
    void induced sort(vector<int> &vec, int val, vector<int> &sa,
                      vector<bool> &sl, vector<int> &lms) {
        vector<int> l(val), r(val);
        for (int c : vec) {
```

```
if (c + 1 < val) l[c + 1] ++;
                      r[c]++;
           partial sum(l.begin(), l.end(), l.begin());
           partial sum(r.begin(), r.end(), r.begin());
           fill(sa.begin(), sa.end(), -1);
           for (int i = lms.size() - 1; i >= 0; i--) sa[--r[vec[lms[i]]]] =
lms[i];
           for (int i : sa) {
                      if (i >= 1 \&\& sl[i - 1]) sa[l[vec[i - 1]]++] = i - 1;
           fill(r.begin(), r.end(), 0);
           for (int c : vec) r[c]++;
           partial sum(r.begin(), r.end(), r.begin());
           for (int k = sa.size() - 1, i = sa[k]; k >= 1; --k, i = sa[k]) {
                      if (i >= 1 \&\& !sl[i - 1]) sa[--r[vec[i - 1]]] = i - 1;
}
vector<int> build sa(vector<int> &vec, int val) {
           int n = vec.size();
           vector<int> sa(n), lms;
           vector<bool> sl(n):
           sl[n-1] = false;
           for (int i = n - 2; i \ge 0; i--) {
                      sl[i] =
                                 (\text{vec}[i] > \text{vec}[i + 1] \mid | (\text{vec}[i] == \text{vec}[i + 1] \& sl[i + 1])
1]));
                      if (sl[i] \&\& !sl[i + 1]) lms.push back(i + 1);
           reverse(lms.begin(), lms.end());
           induced sort(vec, val, sa, sl, lms);
           vector<int> new lms(lms.size()), lms vec(lms.size());
           for (int i = 0, k = 0; i < n; i++) {
                      if (!sl[sa[i]] \&\& sa[i] >= 1 \&\& sl[sa[i] - 1]) new lms[k++] =
sa[i];
           int cur = 0;
           sa[n-1] = cur;
           for (int k = 1; k < (int) new lms.size(); k++) {
                      int i = new lms[k - 1], \bar{j} = new lms[k];
                      if (vec[i] != vec[i]) {
                                 sa[j] = ++cur;
                                 continue:
                      bool flag = false;
                      for (int a = i + 1, b = j + 1;; ++a, ++b) {
                                 if (vec[a] != vec[b]) {
                                           flag = true;
                                           break;
                                if ((!sl[a] && sl[a - 1]) || (!sl[b] && sl[b - 1])) {
                                           flag = !((!sl[a] \&\& sl[a - 1]) \&\& (!sl[b] \&\& sl[b - 1]) \&\& (!sl[b] \&\&
1]));
                                           break:
                                }
                      }
```

```
sa[i] = (flag ? ++cur : cur);
    for (int i = 0; i < (int)lms.size(); i++) lms vec[i] = sa[lms[i]];</pre>
    if (cur + 1 < (int)lms.size()) {</pre>
        auto lms sa = build sa(lms vec, cur + 1);
        for (int i = 0; i < (int)lms.size(); i++)</pre>
            new lms[i] = lms[lms sa[i]];
    induced sort(vec, val, sa, sl, new lms);
    return sa:
vector<int> suffix array() {
    vector < int > vec(n + 1);
    copy(begin(s), end(s), begin(vec));
    vec.back() = '$';
    auto sa = build sa(vec, 256);
    sa.erase(sa.begin());
    return sa;
vector<int> build_lcp() {
    int n = (int)s.size(), k = 0;
    vector<int> rank(n), lcp(n);
    for (int i = 0; i < n; i++) rank[sa[i]] = i;
    for (int i = 0; i < n; i++, k -= !!k) {
        if (rank[i] == n - 1) {
            k = 0;
            continue:
        int j = sa[rank[i] + 1];
        while (i + k < n \&\& j + k < n \&\& s[i + k] == s[j + k]) k++;
        lcp[rank[i]] = k;
    return lcp;
SA() {}
SA(string _s) : s(_s), n(len(s)), pos(n) {
    sa = suffix array();
    lcp = build lcp();
    rmq = RMQ < \overline{int} > (lcp);
    for (int i = 0; i < n; i++) pos[sa[i]] = i;
}
int get lcp(int i,
            int j) { // lcp na çãposio i, indica o lcp
                       // das çõposies i e i+1 do sa
    if (i == i) return n - i;
    int l = pos[i], r = pos[j];
    if (l > r) swap(l, r);
    return rmq.query(l, r);
// string s = a + '+' + b;
tuple<int, int, int> lcs(int n) { // m é o tamanho da string a
    int m = len(s) - n - 1:
    int best len = 0;
```

```
int index s = 0;
        int index t = 0;
        for (int \bar{i} = 0; i < n + m; ++i) {
            if ((sa[i] < n \&\& sa[i + 1] >= n + 1) ||
                (sa[i] >= n + 1 \&\& sa[i + 1] < n)) {
                if (lcp[i] > best len) {
                    best len = lcp[i];
                    index s = min(sa[i], sa[i + 1]);
                    index t = \max(sa[i], sa[i + 1]) - n - 1;
            }
        /*int maior = 0, pos = -1;*/
        /*for (int i = 2; i < n; i++) {*/}
        /* if ((sa[i] < n) != (sa[i - 1] < n)) {*/
        /* if (lcp[i-1] > maior)*/
                maior = lcp[i - 1], pos = sa[i];*/
        /* }*/
        /*}*/
        /*return {maior, pos};*/
        return {best len, index s, index t};
   ll distinct subs() { // n*(n+1)/2 - sum(lcp[i])
        ll resp = (ll)n * ((ll)n + 1) / 2;
        return resp - accumulate(lcp.begin(), lcp.end(), OLL);
};
```

#### 13.15 Suffix automaton

```
#include <bits/stdc++.h>
using namespace std;
#ifdef LOCAL
#include "debug.cpp"
#else
#define dbg(...)
#endif
#define endl '\n'
#define fastio
    ios base::sync with stdio(0); \
    cin.tie(0):
#define int long long
#define all(j) j.begin(), j.end()
#define rall(j) j.rbegin(), j.rend()
#define len(j) (int)j.size()
#define rep(i, a, b) \
    for (common_type_t<decltype(a), decltype(b)> i = (a); i < (b); i++)</pre>
#define rrep(i, a, b) \
    for (common_type_t < decltype(a), decltype(b) > i = (a); i > (b); i--)
#define trav(xi, xs) for (auto &xi : xs)
#define rtrav(xi, xs) for (auto &xi : ranges::views::reverse(xs))
#define pb push back
#define pf push front
#define ppb pop back
```

```
#define ppf pop_front
#define eb emplace back
#define lb lower bound
#define ub upper bound
#define fi first
#define se second
#define emp emplace
#define ins insert
#define divc(a, b) ((a) + (b) - 111) / (b)
using str = string;
using ll = long long;
using ull = unsigned long long;
using ld = long double;
using vll = vector<ll>;
using pll = pair<ll, ll>;
using vll2d = vector<vll>;
using vi = vector<int>;
using vi2d = vector<vi>;
using pii = pair<int, int>;
using vpii = vector<pii>;
using vc = vector<char>;
using vs = vector<str>;
template <typename T, typename T2>
using umap = unordered map<T, T2>;
template <typename T>
using pqmn = priority queue<T, vector<T>, greater<T>>;
template <typename T>
using pqmx = priority queue<T, vector<T>>;
template <typename T, typename U>
inline bool chmax(T &a, U const &b) {
    return (a < b ? a = b, 1 : 0);
template <typename T, typename U>
inline bool chmin(T &a, U const &b) {
    return (a > b ? a = b, 1 : 0);
struct SuffixAutomaton {
    struct state {
        int len, link, cnt, firstpos;
        // this can be optimized using a vector with
        // the alphabet size
        map<char, int> next;
        vi inv link;
    };
    vector<state> st;
    int sz = 0;
    int last;
    vc cloned:
    SuffixAutomaton(const string &s, int maxlen)
        : st(maxlen * 2), cloned(maxlen * 2) {
        st[0].len = 0;
        st[0].link = -1:
        SZ++;
        last = 0:
        for (auto &c : s) add char(c);
```

```
// precompute for count occurences
    for (int i = 1: i < sz: i++) {
        st[i].cnt = !cloned[i];
    vector<pair<state, int>> aux;
    for (int i = 0; i < sz; i++) {
        aux.push back({st[i], i});
    sort(all(aux),
         [](const pair<state, int> &a, const pair<state, int> &b) {
             return a.fi.len > b.fi.len:
    for (auto &[stt, id] : aux) {
        if (stt.link !=-1) {
            st[stt.link].cnt += st[id].cnt;
    // for find every occurende position
    for (int v = 1; v < sz; v++) {
        st[st[v].link].inv_link.push_back(v);
}
void add char(char c) {
    int cur = sz++;
    st[cur].len = st[last].len + 1;
    st[cur].firstpos = st[cur].len - 1;
    int p = last;
    // follow the suffix link until find a
    // transition to c
    while (p != -1 \text{ and } !st[p].next.count(c))  {
        st[p].next[c] = cur:
        p = st[p].link;
    // there was no transition to c so create and
    // leave
    if (p == -1) {
        st[cur].link = 0;
        last = cur;
        return:
    int q = st[p].next[c];
    if (st[p].len + 1 == st[q].len) {
        st[cur].link = q;
    } else {
        int clone = sz++:
        cloned[clone] = true:
        st[clone].len = st[p].len + 1;
        st[clone].next = st[q].next;
        st[clone].link = st[q].link;
        st[clone].firstpos = st[q].firstpos;
        while (p != -1 \text{ and } st[p].next[c] == q) {
            st[p].next[c] = clone;
            p = st[p].link;
        st[q].link = st[cur].link = clone;
```

```
last = cur;
    bool checkOccurrence(const string &t) { // O(len(t))
        int cur = 0;
        for (auto &c : t) {
            if (!st[cur].next.count(c)) return false;
            cur = st[cur].next[c];
        return true;
    Íl totalSubstrings() { // distinct, O(len(s))
        ll tot = 0:
        for (int i = 1; i < sz; i++) {
            tot += st[i].len - st[st[i].link].len;
        return tot;
    // count occurences of a given string t
    int countOccurences(const string &t) {
        int cur = 0:
        for (auto &c : t) {
            if (!st[cur].next.count(c)) return 0;
            cur = st[cur].next[c];
        return st[cur].cnt;
    // find the first index where t appears a
    // substring O(len(t))
    int firstOccurence(const string &t) {
        int cur = 0:
        for (auto c : t) {
            if (!st[cur].next.count(c)) return -1;
            cur = st[cur].next[c];
        return st[cur].firstpos - len(t) + 1;
    vi everyOccurence(const string &t) {
        int cur = 0;
        for (auto c : t) {
            if (!st[cur].next.count(c)) return {};
            cur = st[cur].next[c];
        vi ans;
        qetEveryOccurence(cur, len(t), ans);
        return ans;
    void getEveryOccurence(int v, int P length, vi &ans) {
        if (!cloned[v]) ans.pb(st[v].firstpos - P length + 1);
        for (int u : st[v].inv link) getEveryOccurence(u, P length, ans);
    }
};
void run();
int32 t main() {
#ifndef LOCAL
```

```
fastio;
#endif
   int T = 1;
   /*cin >> T;*/
   rep(t, 0, T) {
       dbg(t);
       run();
   }
}

void run() {
   string S;
   cin >> S;
   SuffixAutomaton sa(S, len(S));
   cout << sa.totalSubstrings() << endl;
}</pre>
```

## 13.16 Suffix-Tree (Ukkonen's Algorithm)

```
/**
 * Author: Unknown
 * Date: 2017-05-15
 * Source: https://e-maxx.ru/algo/ukkonen
 * Description: Ukkonen's algorithm for online suffix tree construction.
 * Each node contains indices [l, r) into the string, and a list of child
 * nodes. Suffixes are given by traversals of this tree, joining [l, r)
 * substrings. The root is 0 (has l = -1, r = 0), non-existent children
   are -1.
 * To get a complete tree, append a dummy symbol -- otherwise it may
   an incomplete path (still useful for substring matching, though).
 * Time: $0(26N)$
 * Status: stress-tested a bit
#pragma once
#include "../Contest/template.cpp"
struct SuffixTree {
    enum { N = 200010, ALPHA = 26 }; // N \sim 2*maxlen+10
   int toi(char c) { return c - 'a'; }
   string a; // v = cur node, g = cur position
   int t[N][ALPHA], l[N], r[N], p[N], s[N], v = 0, q = 0, m = 2;
   void ukkadd(int i, int c) {
   suff:
        if (r[v] <= q) {
            if (t[v][c] == -1) {
                t[v][c] = m;
                l[m] = i;
                p[m++] = v:
                v = s[v];
                q = r[v];
                goto suff;
            v = t[v][c]:
            q = l[v];
        }
```

```
if (q == -1 || c == toi(a[q]))
        q++;
    else {
        l[m + 1] = i;
        p[m + 1] = m;
        l[m] = l[v];
        r[m] = q;
        [v]q = [m]q
        t[m][c] = m + 1;
        t[m][toi(a[q])] = v;
        l[v] = q;
        p[v] = m;
        t[p[m]][toi(a[l[m]])] = m;
        v = s[p[m]];
        a = l[m];
        while (q < r[m]) {
            v = t[v][toi(a[q])];
            a += r[v] - l[v]:
        if (q == r[m])
            s[m] = v;
        else
            s[m] = m + 2:
        q = r[v] - (q - r[m]);
        m += 2;
        qoto suff:
SuffixTree(string a) : a(a) {
    fill(r. r + N. len(a)):
    memset(s, 0, sizeof s);
    memset(t, -1, sizeof t);
    fill(t[1], t[1] + ALPHA, 0);
    s[0] = 1;
    l[0] = l[1] = -1;
    r[0] = r[1] = p[0] = p[1] = 0;
    rep(i, 0, len(a)) ukkadd(i, toi(a[i]));
// example: find longest common substring (uses ALPHA = 28)
int lcs(int node, int i1, int i2, int olen) {
    if (l[node] <= i1 && i1 < r[node]) return 1;</pre>
    if (l[node] <= i2 && i2 < r[node]) return 2;</pre>
    int mask = 0, len = node ? olen + (r[node] - l[node]) : 0;
    rep(c, 0, ALPHA) if (t[node][c] != -1) mask |=
        lcs(t[node][c], i1, i2, len);
    if (mask == 3) best = max(best, {len, r[node] - len});
    return mask;
static pii LCS(string s, string t) {
    SuffixTree st(s + (char)('z' + 1) + t + (char)('z' + 2));
    st.lcs(0, len(s), len(s) + 1 + len(t), 0);
    return st.best;
```

};

#### 13.17 Trie

#### Description:

- build with the size of the alphabet (sigma) and the first char (norm)
- insert(s) insert the string in the trie O(|s| \* sigma)
- erase(s) remove the string from the trie O(|s|)
- find(s) return the last node from the string s, 0 if not found O(|s|)

```
#include "../Contest/template.cpp"
struct Trie {
    vi2d to:
    vi end, pref;
    int sigma:
    char norm:
    Trie(int sigma = 'z' - 'a' + 1, char norm = 'a')
        : sigma(sigma ), norm(norm ) {
        to = {vector<int>(sigma)};
        end = \{0\}, pref = \{0\}:
    int next(int node, char key) { return to[node][key - norm]; }
    void insert(const string &s) {
        int x = 0;
        for (auto c : s) {
            int \&nxt = to[x][c - norm];
            if (!nxt) {
                nxt = len(to);
                to.push back(vi(sigma));
                end.emplace back(0), pref.emplace back(0);
            x = nxt, pref[x]++;
        end[x]++, pref[0]++;
    void erase(const string &s) {
        int x = 0:
        for (char c : s) {
            int &nxt = to[x][c - norm];
            x = nxt, pref[x]--;
            if (!pref[x]) nxt = 0;
        end[x]--, pref[0]--;
    int find(const string &s) {
        int x = 0:
        for (auto c : s) {
            x = to[x][c - norm];
            if (!x) return 0;
        return x;
};
```

#### 13.18 Z-function get occurrence positions

```
Time: O(len(s) + len(p))
#include "../Contest/template.cpp"
vi getOccPos(string& s, string& p) {
    // Z-function
    char delim = '#':
    string t{p + delim + s};
    vi zs(len(t));
    for (int i = 1, l = 0, r = 0; i < len(t); i++) {
        if (i \le r) zs[i] = min(zs[i - l], r - i + 1);
        while (zs[i] + i < len(t)) and t[zs[i]] == t[i + zs[i]]) zs[i]++;
        if (r < i + zs[i] - 1) l = i, r = i + zs[i] - 1;
    // Iterate over the results of Z-function to get
    // ranges
    vi ans:
    int start = len(p) + 1 + 1 - 1;
    for (int i = start; i < len(zs); i++) {</pre>
        if (zs[i] == len(p)) {
            int l = i - start;
            ans.emplace back(l);
    return ans:
template <class T>
std::vector<int> z algorithm(const std::vector<T>& s) {
    int n = int(s.\overline{size}()):
    if (n == 0) return {};
    std::vector<int> z(n);
    for (int i = 1, j = 0; i < n; i++) {
        int & k = z[i];
        k = (i + z[i] \le i) ? 0 : std::min(i + z[i] - i, z[i - i]);
        while (i + k < n \&\& s[k] == s[i + k]) k++;
        if (j + z[j] < i + z[i]) j = i;
    z[0] = n;
    return z;
std::vector<int> z algorithm(const std::string& s) {
    int n = int(s.size());
    std::vector<int> s2(n);
    for (int i = 0; i < n; i++) {
        s2[i] = s[i];
    return z algorithm(s2);
// Pattern matching com um erro ??
// prefixo do apadro dando certo
// depois o sufixo continua dando certo...
//
// Pattern matching com x erros
//
```

```
// únmero de substrings distintas...
//
// únmero de bustrings distintas
// oline ??
//
```

```
// achar o íperodo
//
// kmp é áintercambivel com z
```