Competitive Programming Library

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Algorithms

1.1 Count inversions

Description: Count the number of inversions when transforming the vector l in the vector r, which is also equivalent to the minimum number of swaps required.

Usage: If no r vector is provided it considers r as the sorted vector, if there is no such way to turn l into r using swaps then -1 is returned

Time: $O(N \log N)$

```
#pragma once
#include "../Contest/template.cpp"
template <typename T>
ll countInversions(vector<T> l, vector<T> r = {}) {
    if (!len(r)) r = l, sort(all(r));
    int n = len(l);
    vi v(n), bit(n);
    vector<pair<T, int>> w;
    rep(i, 0, n) w.eb(r[i], i + 1);
    sort(all(w));
    rep(i, 0, n) {
        auto it = lower bound(all(w), make pair(l[i], 0));
        if (it == w.end() or it->first != \overline{\lfloor (i) \rfloor} return -1; // impossible
        v[i] = it -> second:
        it->second = -1;
    il ans = 0:
    rrep(i, n - 1, 0 - 1) {
        for (int j = v[i] - 1; j; j = j \& -j) ans += bit[j];
        for (int j = v[i]; j < n; j += j \& -j) bit[j]++;
    return ans;
```

Ternary search (integer)

Description: Given a unimodal function f defined between the integers l and r finds an xsuch that f(x) is maximum/minimum.

Usage: Just pass the range l, r of the function you are interested, the function that receives an integer and if you want the maximum value use the cmp = greater < ll>(), otherwise less<ll>().

Time: $O(\log r - l + 1)$

Memory: O(1)

```
#include "../Contest/template.cpp"
template <auto cmp = greater<ll>()>
ll ternary search(ll l, ll r, function<ll(ll)> f) {
   static const ll eps = 3:
   while (r - l >= eps) {
       ll m1 = l + (r - l) / 3;
        ll m2 = r - (r - l) / 3;
        if (cmp(f(m1), f(m2)))
            r = m2;
       else
```

```
l = m1:
rep(i, l, r + 1) if (cmp(f(i), f(l))) l = i;
return l;
```

1.3 Ternary search (real)

```
#include "../Contest/template.cpp"
template <auto cmp = greater<ld>()>
ld ternarySearch(ld l, ld r, function<ld(ld)> f, const ld eps = 1e-9) {
    while (r - l >= eps) {
        ld m1 = l + (r - l) / 3;
        ld m2 = r - (r - l) / 3;
        if (cmp(f(m1), f(m2)))
            r = m2:
        else
            l = m1:
    }
    return l;
```

Combinatorics

2.1 Process all partitions of a set

```
Description: generate every distinct group of a set that contains elements from 0 to N-1,
and pass it to the given function "process". If N is 4 the sets generated would be:
[\{\{0,1,2,3\}\}] \ [\{\{0,1,2\},\{3\}\}] \ [\{\{0,1,3\},\{2\}\}] \ [\{\{0,1\},\{2,3\}\}] \ [\{\{0,1\},\{2\},\{3\}\}] \ [\{\{0,2,3\},\{1\}\}]]
[\{\{0,2\},\{1,3\}\}] \ [\{\{0,2\},\{1\},\{3\}\}] \ [\{\{0,3\},\{1,2\}\}] \ [\{\{0\},\{1,2,3\}\}] \ [\{\{0\},\{1,2\},\{3\}\}]]
[\{\{0,3\},\{1\},\{2\}\}] \ [\{\{0\},\{1,3\},\{2\}\}] \ [\{\{0\},\{1\},\{2,3\}\}] \ [\{\{0\},\{1\},\{2\},\{3\}\}]
Time: O(B(N)), Bell Number of N
Memory: O(N)
```

```
#include "../Contest/template.cpp"
void process all partitions of a set(
    const int N, const function<void(const vi2d &)> process) {
    vi2d groups;
    groups.reserve(N);
    function<void(int)> dfs = [&](int idx) {
        if (idx == N) {
            process(groups);
            return:
        rep(i, 0, len(groups)) {
            groups[i].eb(idx);
            dfs(idx + 1);
            groups[i].ppb();
        groups.pb({idx});
        dfs(idx + 1);
```

```
groups.ppb();
};
_dfs(0);
```

3 Contest

3.1 bash config

```
#copy first argument to clipborad ! ONLY WORK ON XORG !
alias clip="xclip -sel clip"
# compile the $1 parameter, if a $2 is provided
# the name will be the the binary output, if
# none is provided the binary name will be
# 'a.out'
comp() {
  echo ">> COMPILING $1 <<" 1>&2
  if [ $# -gt 1 ]; then
    outfile="${2}"
  else
    outfile="a.out"
  time q++-std=c++20 \setminus
    -02 \
    -q3 \
    -Wall \
    -fsanitize=address,undefined \
    -fno-sanitize-recover \
    -D LOCAL \
    -o "${outfile}" \
  if [ $? -ne 0 ]; then
    echo ">> FAILED <<" 1>&2
    return 1
  fi
  echo ">> DONE << " 1>&2
# run the binary given in $1, if none is
# given it will try to run the 'a.out'
# binary
run() {
  to run=./a.out
  if [ -n "$1" ]; then
    to_run="$1"
  fi
  time $to_run
# just comp and run your cpp file
# accpets <in1 >out and everything else
comprun() {
  comp "$1" "a" && run ./a ${@:2}
testall() {
```

```
comp "$1" generator
  comp "$2" brute
  comp "$3" main
  input counter=1
  while true; do
    echo "$input counter"
    run ./generator >input
    run ./main <input >main output.txt
    run ./brute <input >brute output.txt
    diff brute output.txt main output.txt
    if [ $? -ne 0 ]; then
      echo "Outputs differ at input $input counter"
      echo "Brute file output:"
      cat brute output.txt
      echo "Main file output:"
      cat main output.txt
      echo "input used: "
      cat input
      break
    fi
    ((input counter++))
touch macro() {
  cat "$1"/template.cpp >"$2"
 cat "$1"/run.cpp >>"$2"
 cp "$1"/debug.cpp .
# Creates a contest with hame $2
# Copies the macro and debug file from $1
# Already creates files a...z .cpp and .py
prepare contest() {
  mkdir "$2"
  cd "$2"
  for i in {a..z}; do
    touch macro $1 $i.cpp
  done
get file hash() {
  local hash=$(cpp -dD -P -fpreprocessed "$1" | tr -d '[:space:]' | md5sum
      cut -c-6
  echo "$hash"
3.2 debug
template <typename T>
concept Printable = requires(T t) {
    { std::cout << t } -> std::same as<std::ostream &>;
template <Printable T>
void print(const T &x) {
```

```
cerr << x:
template <size t T>
void print(const bitset<T> &x) {
   cerr << x;
template <typename A, typename B>
void print(const pair<A, B> &p);
template <typename... A>
void print(const tuple<A...> &t);
template <typename T>
void print(stack<T> s);
template <typename T>
void print(queue<T> q);
template <typename T, typename... U>
void __print(priority_queue<T, U...> q);
template <typename A>
void print(const A &x) {
   bool first = true;
   cerr << '{';
   for (const auto &i : x) {
       cerr << (first ? "" : ","), print(i);</pre>
       first = false;
   cerr << '}';
template <typename A, typename B>
void print(const pair<A, B> &p) {
   cerr << '(';
     print(p.first);
   cerr << ',';
    print(p.second);
   cerr << ')';
template <typename... A>
void __print(const tuple<A...> &t) {
   bool first = true;
   cerr << '(';
   apply(
        [&first](const auto &...args) {
            ((cerr << (first ? "" : ","), __print(args), first = false),</pre>
   ...);
       t):
   cerr << ')';
template <typename T>
void print(stack<T> s) {
   vector<T> debugVector;
   while (!s.empty()) {
       T t = s.top();
       debugVector.push back(t);
       s.pop();
   reverse(debugVector.begin(), debugVector.end());
   __print(debugVector);
```

```
template <typename T>
void print(queue<T> q) {
    vector<T> debugVector;
    while (!q.empty()) {
        T t = q.front();
        debugVector.push back(t);
        q.pop();
    __print(debugVector);
template <typename T, typename... U>
void print(priority queue<T, U...> q) {
    vector<T> debugVector;
    while (!q.empty()) {
        T t = q.top();
        debugVector.push_back(t);
        q.pop();
    __print(debugVector);
void _print() { cerr << "]\n"; }</pre>
template <typename Head, typename... Tail>
void _print(const Head &H, const Tail &...T) {
    print(H);
    <u>if</u> (sizeof...(T)) cerr << ", ";
    _print(T...);
#define dbg(x...)
    cerr << "[" << #x << "] = ["; \
    print(x)
3.3 run
void run();
int32 t main() {
#ifndef LOCAL
    fastio;
#endif
    int T = 1;
    cin >> T;
    rep(t, 0, T) {
        dbg(t);
        run();
void run() {}
3.4 short-template
```

```
using namespace std;
#define fastio
    ios_base::sync_with_stdio(0);
    cin.tie(0);
void run() {}
int32_t main(void) {
    fastio;
    int t;
    t = 1;
    // cin >> t;
    while (t--) run();
}
```

3.5 template

```
#pragma once
#include <bits/stdc++.h>
using namespace std;
#ifdef LOCAL
#include "debug.cpp"
#else
#define dbq(...)
#endif
#define fastio
   ios base::sync with stdio(0); \
    cin.tie(0);
#define all(j) j.begin(), j.end()
#define rall(j) j.rbegin(), j.rend()
#define len(j) (int)j.size()
#define rep(i, a, b) \
    for (common type t < decltype(a), decltype(b) > i = (a); i < (b); i++)
#define rrep(i, a, b) \
    for (common type t < decltype(a), decltype(b) > i = (a); i > (b); i--)
#define trav(xi, xs) for (auto &xi : xs)
#define rtrav(xi, xs) for (auto &xi : ranges::views::reverse(xs))
using ll = long long;
#define endl '\n'
#define pb push_back
#define pf push front
#define ppb pop back
#define ppf pop front
#define eb emplace back
#define ef emplace back
#define lb lower bound
#define ub upper bound
#define fi first
#define se second
#define emp emplace
#define ins insert
\#define\ divc(a, b)\ ((a) + (b) - 1ll)\ /\ (b)
using str = string;
using ull = unsigned long long;
using ld = long double;
using vll = vector<ll>:
```

```
using pll = pair<ll, ll>;
using vll2d = vector<vll>:
using vi = vector<int>;
using vi2d = vector<vi>;
using pii = pair<int, int>;
using vpii = vector<pii>;
using vc = vector<char>;
using vs = vector<str>;
template <typename T, typename T2>
using umap = unordered_map<T, T2>;
template <tvpename T>
using pgmn = priority queue<T, vector<T>, greater<T>>;
template <typename T>
using pqmx = priority queue<T, vector<T>>;
template <typename T, typename U>
inline bool chmax(T &a, U const &b) {
    return (a < b ? a = b, 1 : 0);
template <typename T, typename U>
inline bool chmin(T &a, U const &b) {
    return (a > b ? a = b, 1 : 0);
template <typename T>
std::istream &operator>>(std::istream &is, std::vector<T> &vec) {
    for (auto &element : vec) {
        is >> element:
    return is;
template <typename T> // print vector
ostream &operator<<(ostream &os, vector<T> &xs) {
    rep(i, os.iword(0), xs.size()) os << xs[i] << (i == xs.size() ? "" : "
    ");
    os.iword(0) = 0;
    return os;
}
3.6 vim config
```

```
set sta nu rnu sc cindent noswapfile
set ts=2 sw=2
set bg=dark ruler clipboard=unnamed,unnamedplus, timeoutlen=100
colorscheme default
syntax on
" Takes the hash of the selected text and put
" in the vim clipboard
function! HashSelectedText()
    " Yank the selected text to the unnamed register
    normal! gvy
    " Use the system() function to call sha256sum with the yanked text
    let l:hash = system('echo' . shellescape(@@) . ' | sha256sum')
    " Yank the hash into Vim's unnamed register
    let @" = l:hash
endfunction
```

4 Data Structures

4.1 2D Segment Tree

4.1.1 Point update query sum

```
#include "../../Contest/template.cpp"
template <typename T, auto op>
struct SeamentTree2D {
   int h. w:
   vector<vector<T>> t;
   SegmentTree2D(const vector<vector<T>> &a)
        : h(a.size()), w(a.back().size()), t(h * 4, vector<T>(w * 4)) {
       build x(1, 0, h - 1, a);
   void build_y(int vx, int lx, int rx, int vy, int ly, int ry,
                const vector<vector<T>> &a) {
       if (ly == ry) {
            if (lx == rx)
                t[vx][vy] = a[lx][ly];
                t[vx][vy] = op(t[vx * 2][vy], t[vx * 2 + 1][vy]);
       } else {
            int my = (ly + ry) / 2;
            build_y(vx, lx, rx, vy * 2, ly, my, a);
            build_y(vx, lx, rx, vy * 2 + 1, my + 1, ry, a);
            t[vx][vy] = op(t[vx][vy * 2], t[vx][vy * 2 + 1]);
       }
   void build x(int vx, int lx, int rx, const vector<vector<T>> &a) {
       if (lx != rx) {
            int mx = (lx + rx) / 2;
            build x(vx * 2, lx, mx, a);
            build x(vx * 2 + 1, mx + 1, rx, a);
       build y(vx, lx, rx, 1, 0, w - 1, a);
   T query y(int vx, int vy, int tly, int try , int ly, int ry) {
       if (ly > ry) return 0;
       if (ly == tly && try_ == ry) return t[vx][vy];
       int tmy = (tly + try) / 2;
       return op(query y(vx, vy * 2, tly, tmy, ly, min(ry, tmy)),
                  query y(vx, vy * 2 + 1, tmy + 1, try, max(ly, tmy + 1),
    ry));
   T query x(int vx, int tlx, int trx, int lx, int rx, int ly, int ry) {
       if (lx > rx) return 0;
       if (lx == tlx \&\& trx == rx) return query y(vx, 1, 0, w - 1, ly, ry)
   );
       int tmx = (tlx + trx) / 2;
        return op(
            query x(vx * 2, tlx, tmx, lx, min(rx, tmx), ly, ry),
            query_x(vx * 2 + 1, tmx + 1, trx, max(lx, tmx + 1), rx, ly, ry
   ));
```

```
void update y(int vx, int lx, int rx, int vy, int ly, int ry, int x,
                  int new val) {
        if (ly == ry) {
            if (lx == rx)
                t[vx][vy] = new val;
                t[vx][vy] = op(t[vx * 2][vy], t[vx * 2 + 1][vy]);
        } else {
            int my = (ly + ry) / 2;
            if (y \le my)
                update_y(vx, lx, rx, vy * 2, ly, my, x, y, new_val);
            else
                update y(vx, lx, rx, vy * 2 + 1, my + 1, ry, x, y, new val
   );
            t[vx][vy] = op(t[vx][vy * 2], t[vx][vy * 2 + 1]);
    }
    void update x(int vx, int lx, int rx, int x, int y, T new val) {
        if (lx != rx) {
            int mx = (lx + rx) / 2;
            if (x \le mx)
                update_x(vx * 2, lx, mx, x, y, new_val);
                update x(vx * 2 + 1, mx + 1, rx, x, y, new val);
        update_y(vx, lx, rx, 1, 0, w - 1, x, y, new_val);
    }
    T query(int lx, int rx, int ly, int ry) {
        return query x(1, 0, h - 1, lx, rx, ly, ry);
};
```

4.2 SQRT decomposition

4.2.1 two-sequence-queries

```
r = numeric limits<int>::min();
        x = y = prod = sum as = sum bs = 0;
    };
};
int sqrtLen;
vector<t_sqrt> blocks;
vector<ll> as, bs:
SqrtDecomposition(const vector<ll> &as , const vector<ll> &bs ) {
    int n = as .size();
    sgrtLen = (int) sgrt(n + .0) + 1;
    blocks.resize(sqrtLen + 6.66);
    as = as_{;}
    bs = bs;
    for (int i = 0; i < n; i++) {
        auto &bi = blocks[i / sqrtLen];
        bi.l = min(bi.l, i);
        bi.r = max(bi.r, i);
        bi.sum as = sum(bi.sum as, as[i]);
        bi.sum bs = sum(bi.sum bs, bs[i]);
        bi.prod = sum(bi.prod, mul(as[i], bs[i]));
}
// adds x to a[i], and y to b[i], in range [l,
void update(int l, int r, ll x, ll y) {
    auto apply1 = [\&](int idx, ll x, ll y) -> void {
        auto &block = blocks[idx / sqrtLen];
        block.prod = sub(block.prod, mul(as[idx], bs[idx]));
        block.sum as = sub(block.sum as, as[idx]);
        block.sum bs = sub(block.sum bs, bs[idx]);
        as[idx] = sum(as[idx], x);
        bs[idx] = sum(bs[idx], y);
        block.prod = sum(block.prod, as[idx] * bs[idx]);
        block.sum as = sum(block.sum as, as[idx]);
        block.sum bs = sum(block.sum bs, bs[idx]);
    };
    auto apply2 = [&](int idx, ll x, ll y) -> void {
        blocks[idx].x = sum(blocks[idx].x, x);
        blocks[idx].y = sum(blocks[idx].y, y);
    };
    int cl = l / sqrtLen, cr = r / sqrtLen;
    if (cl == cr) {
        for (int i = l; i <= r; i++) {
            apply1(i, x, y);
    } else {
        for (int i = l; i \le (cl + 1) * sqrtLen - 1; i++) {
            apply1(i, x, y);
        for (int i = cl + 1; i \le cr - 1; i++) {
            apply2(i, x, y);
        }
```

```
for (int i = cr * sqrtLen; i <= r; i++) {
                apply1(i, x, y);
    }
    // sum of a[i]*b[i] in range [l r]
    ll query(int l, int r) {
        auto eval1 = [&](int idx) -> ll {
            auto &block = blocks[idx / sqrtLen];
            return mul(sum(as[idx], +block.x), sum(bs[idx], block.y));
        };
        auto eval2 = [\&](int idx) \rightarrow ll \{
            auto &block = blocks[idx];
            ll ret = 0:
            ret = sum(
                ret, mul(mul(block.x, block.y), sum(sub(block.r, block.l),
    1)));
            ret = sum(ret, block.prod);
            ret = sum(ret, block.y * block.sum as);
            ret = sum(ret, block.x * block.sum bs);
            return ret;
        };
        ll ret = 0;
        int cl = l / sqrtLen, cr = r / sqrtLen;
        if (cl == cr) {
            for (int i = l; i <= r; i++) {
                ret = sum(ret, eval1(i));
        } else {
            for (int i = l; i \le (cl + 1) * sqrtLen - 1; i++) {
                ret = sum(eval1(i), ret);
            for (int i = cl + 1; i \le cr - 1; i++) {
                ret = sum(ret, eval2(i));
            for (int i = cr * sqrtLen; i <= r; i++) {
                ret = sum(ret, eval1(i));
        }
        return ret;
};
```

4.3 Segment Tree Point Update Range Query (bottom-up)

4.3.1 Query GCD

```
using ll = long long;
struct Node {
    ll value;
    bool undef;
    Node() : value(1), undef(1) {}; // Neutral element
```

```
Node(ll v) : value(v), undef(0) {};
};
inline Node combine(const Node &nl, const Node &nr) {
   if (nl.undef) return nr;
   if (nr.undef) return nl;
   Node m:
   m.value = gcd(nl.value, nr.value);
   m.undef = false:
   return m;
template <typename T = Node, auto F = combine>
struct SegTree {
   int n:
    vector<T> st;
   SegTree(int n) : n(n), st(n << 1) {}
   void assign(int p, const T &k) {
        for (st[p += n] = k; p >>= 1;) st[p] = F(st[p << 1], st[p << 1]
   1]);
   T query(int l, int r) {
       T ansl, ansr;
        for (l += n, r += n + 1; l < r; l >>= 1, r >>= 1) {
            if (l \& 1) ans l = F(ans l, st[l++]):
            if (r \& 1) ansr = F(st[--r], ansr);
        return F(ansl, ansr);
};
```

4.3.2 Query Max Subarray Sum

```
#pragma once
#include "../../Contest/template.cpp"
#include "./Struct.cpp"
const ll oo = 1e9;
struct Node {
   ll tot, suf, pref, best;
    // Neutral element
   Node(): tot(-oo), suf(-oo), pref(-oo), best(-oo) {} // Neutral
   element
   // for assign
   Node(ll x) { tot = x, suf = x, pref = x, best = max(0ll, x); }
};
Node combine(Node &nl, Node &nr) {
   if (nl.tot == - oo) return nr;
   if (nr.tot == - oo) return nl;
   Node m:
   m.tot = nl.tot + nr.tot:
   m.pref = max({nl.pref, nl.tot + nr.pref});
   m.suf = max({nr.suf, nr.tot + nl.suf});
   m.best = max({nl.best, nr.best, nl.suf + nr.pref});
   return m:
```

```
using SegTreeMaxSubarraySum = SegTreeBottomUp<Node, Node(), combine>;
```

4.3.3 Query min

4.3.4 Query sum

```
#pragma once
#include "../../Contest/template.cpp"
#include "./Struct.cpp"
template <typename T>
using SegTreeBottomUpSumQuery =
    SegTreeBottomUp<T, T(0), [](T a, T b) { return a + b; }>;
```

4.3.5 Struct

```
* @Description:
        merge should be function\langle T(T,T)\rangle, that
        makes the necessary operation between two
        nodes in the seament tree
 *
* */
#pragma once
#include "../../Contest/template.cpp"
template <typename T, T identity, auto merge>
struct SegTreeBottomUp {
    int size:
    vector<T> arr:
    SegTreeBottomUp(int n) {
        for (size = 1: size < n: size <<= 1):
        arr.resize(size << 1);</pre>
    void assign(int pos, const T &val) {
        for (arr[pos += size] = val; pos >>= 1;)
            arr[pos] = merge(arr[pos << 1], arr[pos << 1 | 1]);</pre>
    T query(int l, int r) {
        T ans l = identity, ans r = identity;
        for (\bar{l} += size, r += size + 1; l < r; l >>= 1, r >>= 1) {
            if (l & 1) ans l = merge(ans l, arr[l++]);
            if (r \& 1) ans r = merge(arr[--r], ans r);
```

```
return merge(ans l, ans r);
    SegTreeBottomUp(const vector<T> &vec) : SegTreeBottomUp(len(vec)) {
        copy(all(vec), begin(arr) + size);
        rrep(i, size - 1, 0) arr[i] = merge(arr[i << 1], arr[i << 1 | 1]);
};
4.4 Segment tree (dynamic)
4.4.1 Range Max Query Point Max Assignment
Description: Answers range queries in ranges until 10<sup>9</sup> (maybe more)
Time: Query and update O(n \cdot \log n)
struct node:
node *newNode():
struct node {
    node *left, *right;
    int lv, rv;
    ll val;
    node() : left(NULL), right(NULL), val(-oo) {}
    inline void init(int l, int r) {
        lv = l:
        rv = r;
    inline void extend() {
        if (!left) {
            int m = (lv + rv) / 2;
            left = newNode();
            right = newNode();
            left->init(lv, m);
            right->init(m + 1, rv);
    }
    ll query(int l, int r) {
        if (r < lv || rv < l) {</pre>
            return 0;
        if (l <= lv && rv <= r) {
            return val;
        extend();
        return max(left->query(l, r), right->query(l, r));
    void update(int p, ll newVal) {
        if (lv == rv) {
            val = max(val, newVal);
            return;
        }
        extend():
        (p <= left->rv ? left : right)->update(p, newVal);
        val = max(left->val, right->val);
```

```
};
const int BUFFSZ(1e7);
node *newNode() {
    static int bufSize = BUFFSZ;
    static node buf[(int)BUFFSZ];
    assert(bufSize);
    return &buf[--bufSize];
struct SegTree {
    int n;
    node *root;
    SegTree(int n) : n( n) {
        root = newNode():
        root->init(0, n);
    ll query(int l, int r) { return root->query(l, r); }
    void update(int p, ll v) { root->update(p, v); }
};
4.4.2 Range Sum Query Point Sum Update
Description: Answers range queries in ranges until 10<sup>9</sup> (maybe more)
Time: Query and update in O(n \cdot \log n)
struct node;
node *newNode();
struct node {
    node *left, *right;
    int lv. rv:
    ll val;
    node() : left(NULL), right(NULL), val(0) {}
    inline void init(int l, int r) {
        lv = l;
        rv = r;
    inline void extend() {
        if (!left) {
             int m = (rv - lv) / 2 + lv;
            left = newNode();
             right = newNode();
            left->init(lv, m);
             right->init(m + 1, rv);
        }
    ll query(int l, int r) {
        if (r < lv || rv < l) {</pre>
             return 0;
        if (l <= lv && rv <= r) {</pre>
             return val;
        extend();
```

```
return left->query(l, r) + right->query(l, r);
   }
   void update(int p, ll newVal) {
        if (lv == rv) {
            val += newVal:
            return:
        }
        extend():
        (p <= left->rv ? left : right)->update(p, newVal);
        val = left->val + right->val;
};
const int BUFFSZ(1.3e7);
node *newNode() {
   static int bufSize = BUFFSZ:
   static node buf((int)BUFFSZ1:
   // assert(bufSize):
   return &buf[--bufSizel:
struct SegTree {
   int n:
   node *root:
   SegTree(int _n) : n(_n) {
        root = newNode():
        root->init(0, n);
   ll query(int l, int r) { return root->query(l, r); }
   void update(int p, ll v) { root->update(p, v); }
};
```

4.5 Segment tree point update range query (top-down)

4.5.1 Query hash (top down)

```
#include "../../Contest/template.cpp"
const ll MOD = 1'000'000'009:
const ll P = 31;
const int MAXN = 2'000'000;
ll pows[MAXN + 1];
void computepows() {
    pows[0] = 1:
    for (int i = 1; i <= MAXN; i++) {
        pows[i] = (pows[i - 1] * P) % MOD:
struct Node {
   ll hash;
   Node(): hash(-1) {}; // Neutral element
   Node(ll v) : hash(v) {};
};
inline Node combine(Node &vl, Node &vr, int nl, int nr, int al, int ar) {
   if (vl.hash == -1) return vr;
   if (vr.hash == -1) return vl:
```

```
Node vm:
    int nm = midpoint(nl, nr);
    int lsize = min(nm, qr) - max(nl, ql) + 1;
    vm.hash = (vl.hash + ((vr.hash * pows[lsize]) % MOD)) % MOD;
    return vm:
template <typename T = Node, auto F = combine>
struct SegTree {
    int n:
    vector<T> st:
    SegTree(int n) : n(n), st(n \ll 2) {}
    void assign(int p, const T \&v) { assign(1, 0, n - 1, p, v); }
    void assign(int node, int l, int r, int p, const T &v) {
        if (l == r) {
            st[node] = v:
            return;
        int m = midpoint(l, r);
        if (p \ll m)
            assign(node << 1, l, m, p, v);
        else
            assign(node << 1 | 1, m + 1, r, p, v);
        st[node] = F(st[node << 1], st[node << 1 | 1], l, r, l, r);
    inline T query(int l, int r) { return query(1, 0, n - 1, l, r); }
    inline T query(int node, int nl, int nr, int l, int r) const {
        if (r < nl or nr < l) return T():
        if (l <= nl and nr <= r) return st[node];</pre>
        int m = midpoint(nl, nr);
        auto a = query(node << 1, nl, m, l, r);</pre>
        auto b = query(node << 1 | 1, m + 1, nr, l, r);
        return F(a, b, nl, nr, l, r);
};
```

4.6 Segment tree range update range query

4.6.1 Arithmetic progression sum update, query sum

Description: Makes arithmetic progression updates in range and sum queries.

Usage: Considering PA(A, R) = [A + R, A + 2R, A + 3R, ...]

- $update_set(l, r, A, R)$: sets [l, r] to PA(A, R)
- update_add(l, r, A, R): sum PA(A, R) in [l, r]
- query(l, r): sum in range [l, r]

Time: build O(N), updates and queries O(log N)

```
const ll oo = 1e18;
struct SegTree {
    struct Data {
        ll sum;
        ll set_a, set_r, add_a, add_r;
        Data() : sum(0), set a(oo), set r(0), add a(0), add r(0) {}
```

```
};
int n;
vector<Data> seq;
SegTree(int n ) : n(n ), seg(vector<Data>(4 * n)) {}
void prop(int p, int l, int r) {
    int sz = r - l + 1;
    ll &sum = seg[p].sum, &set_a = seg[p].set_a, &set_r = seg[p].set_r
       &add a = seq[p].add a, &add r = seq[p].add r;
    if (set a != oo) {
        set a += add a, set r += add r;
        sum = set a * sz + set r * sz * (sz + 1) / 2;
        if (l != r) {
            int m = (l + r) / 2;
            seg[2 * p].set_a = set_a;
            seg[2 * p].set r = set r;
            seg[2 * p].add a = seg[2 * p].add r = 0;
            seq[2 * p + 1].set a = set a + set r * (m - l + 1);
            seq[2 * p + 1].set r = set r;
            seg[2 * p + 1].add a = seg[2 * p + 1].add r = 0;
        set a = oo, set r = 0;
        add a = add r = 0;
    } else \overline{i}f (add \overline{a} or add r) {
        sum += add a * sz + add r * sz * (sz + 1) / 2;
        if (l != r) {
            int m = (l + r) / 2;
            seg[2 * p].add a += add a;
            seg[2 * p].add r += add r;
            seq[2 * p + 1].add a += add a + add r * (m - l + 1);
            seq[2 * p + 1].add r += add r;
        add a = add r = 0;
int inter(pii a, pii b) {
    if (a.first > b.first) swap(a, b);
    return max(0, min(a.second, b.second) - b.first + 1);
ll set(int a, int b, ll aa, ll rr, int p, int l, int r) {
    prop(p, l, r);
    if (b < l or r < a) return seq[p].sum;</pre>
    if (a <= l and r <= b) {
        seg[p].set_a = aa;
        seg[p].set r = rr;
        prop(p, l, r);
        return seq[p].sum;
    int m = (l + r) / 2;
    int tam l = inter({l, m}, {a, b});
    return \overline{seg}[p].sum = set(a, b, aa, rr, 2 * p, l, m) +
                         set(a, b, aa + rr * tam l, rr, 2 * p + 1, m +
1, r);
}
```

```
void update set(int l, int r, ll aa, ll rr) {
        set(l, r, aa, rr, 1, 0, n - 1);
    ll add(int a, int b, ll aa, ll rr, int p, int l, int r) {
        prop(p, l, r);
        if (b < l or r < a) return seq[p].sum;</pre>
        if (a <= l and r <= b) {
            seq[p].add a += aa;
            seg[p].add r += rr;
            prop(p, l, r);
            return seg[p].sum;
        int m = (l + r) / 2:
        int tam_l = inter({l, m}, {a, b});
        return seg[p].sum = add(a, b, aa, rr, 2 * p, l, m) +
                            add(a, b, aa + rr * tam l, rr, 2 * p + 1, m +
   1, r);
    void update add(int l, int r, ll aa, ll rr) {
        add(l, r, aa, rr, 1, 0, n - 1);
    ll query(int a, int b, int p, int l, int r) {
        prop(p, l, r);
        if (b < l or r < a) return 0;</pre>
        if (a <= l and r <= b) return seq[p].sum;
        int m = (l + r) / 2;
        return query(a, b, 2 * p, l, m) + query(a, b, 2 * p + 1, m + 1, r)
    Il query(int l, int r) { return query(l, r, 1, 0, n - 1); }
};
4.6.2 Increment update query min & max (bottom up)
using SegT = ll;
    SegT mx, mn;
    OuervT()
```

```
struct QueryT {
        : mx(numeric limits<SeqT>::min()), mn(numeric limits<SeqT>::max())
    QueryT(SegT _v) : mx(_v), mn(_v) {}
};
inline QueryT combine(QueryT ln, QueryT rn, pii lr1, pii lr2) {
    chmax(ln.mx, rn.mx);
    chmin(ln.mn, rn.mn);
    return ln;
using LazyT = SeqT;
inline QueryT applyLazyInQuery(QueryT q, LazyT l, pii lr) {
    if (q.mx == QueryT().mx) q.mx = SeqT();
    if (q.mn == QueryT().mn) q.mn = SeqT();
    q.mx += l, q.mn += l;
    return q;
```

```
inline LazyT applyLazyInLazy(LazyT a, LazyT b) { return a + b; }
using UpdateT = SegT;
inline QueryT applyUpdateInQuery(QueryT q, UpdateT u, pii lr) {
   if (q.mx == QueryT().mx) q.mx = SeqT();
   if (q.mn == QueryT().mn) q.mn = SegT();
   q.mx += u, q.mn += u;
    return q;
inline LazyT applyUpdateInLazy(LazyT l, UpdateT u, pii lr) { return l + u;
template <typename Qt = QueryT, typename Lt = LazyT, typename Ut = UpdateT
          auto C = combine, auto ALQ = applyLazyInQuery,
          auto ALL = applyLazyInLazy, auto AUQ = applyUpdateInQuery,
          auto AUL = applyUpdateInLazy>
struct LazySegmentTree {
   int n, h;
   vector<Qt> ts;
   vector<Lt> ds;
   vector<pii> lrs;
   LazySegmentTree(int n)
        : n(_n),
          h(sizeof(int) * 8 - builtin clz(n)),
          ts(n \ll 1).
          ds(n),
          lrs(n << 1) {
        rep(i, 0, n) lrs[i + n] = {i, i};
        rrep(i, n - 1, 0) {
            lrs[i] = {lrs[i << 1].first, lrs[i << 1 | 1].second};</pre>
   LazySegmentTree(const vector<Qt> &xs) : LazySegmentTree(len(xs)) {
        copy(all(xs), ts.begin() + n);
        rep(i, 0, n) lrs[i + n] = {i, i};
        rrep(i, n - 1, 0) {
            ts[i] = C(ts[i << 1], ts[i << 1 | 1], lrs[i << 1], lrs[i << 1
   | 1]);
   void set(int p, Qt v) {
        ts[p + n] = v;
        build(p + n);
   void upd(int l, int r, Ut v) {
        l += n. r += n + 1:
        int 10 = 1, r0 = r:
        for (; l < r; l >>= 1, r >>= 1) {
            if (l & 1) apply(l++, v);
            if (r & 1) apply(--r, v);
        build(l0), build(r0 - 1);
   }
```

```
Qt qry(int l, int r) {
        l += n, r += n + 1:
        push(l), push(r - 1);
        Qt resl = Qt(), resr = Qt();
        pii lr1 = \{l, l\}, lr2 = \{r, r\};
        for (; l < r; l >>= 1, r >>= 1) {
            if (l & 1) resl = C(resl, ts[l], lr1, lrs[l]), l++;
            if (r & 1) r--, resr = C(ts[r], resr, lrs[r], lr2);
        return C(resl, resr, lr1, lr2);
    void build(int p) {
        while (p > 1) {
            p >>= 1;
            ts[p] =
                ALO(C(ts[p << 1], ts[p << 1 | 1], lrs[p << 1], lrs[p << 1]
    | 1]),
                    ds[p], lrs[p]);
    void push(int p) {
        rrep(s, h, 0) {
            int i = p \gg s;
            if (ds[i] != Lt()) {
                apply(i << 1, ds[i]), apply(i << 1 | 1, ds[i]);
                ds[i] = Lt();
    inline void apply(int p, Ut v) {
        ts[p] = AUQ(ts[p], v, lrs[p]);
        if (p < n) ds[p] = AUL(ds[p], v, lrs[p]);
};
4.6.3 Increment update sum query (top down)
```

```
struct Lnode {
    ll v;
    bool assign;
    Lnode() : v(), assign() {} // Neutral element
    Lnode(ll v, bool a = 0) : v(v), assign(a) {};
};
using Qnode = ll;
using Unode = Lnode;
struct LSegTree {
    int n, ql, qr;
    vector<0node> st:
    vector<Lnode> lz;
    Qnode merge(Qnode lv, Qnode rv, int nl, int nr) { return lv + rv; }
    void prop(int i, int l, int r) {
        if (lz[i].assign) {
```

```
st[i] = lz[i].v * (r - l + 1);
        if (l != r) lz[tol(i)] = lz[tor(i)] = lz[i];
    } else {
        st[i] += lz[i].v * (r - l + 1);
        if (l != r) lz[tol(i)].v += lz[i].v, lz[tor(i)].v += lz[i].v;
    lz[i] = Lnode();
void applyV(int i, Unode v) {
    if (v.assign) {
        lz[i] = v;
    } else {
        lz[i].v += v.v;
LSegTree() {}
LSegTree(int _n) : n(_n), st(_n << 2), lz(_n << 2) {}
bool disjoint(int l, int r) { return qr < l or r < ql; }</pre>
bool contains(int l, int r) { return ql <= l and r <= qr; }</pre>
int tol(int i) { return i << 1; }</pre>
int tor(int i) { return i << 1 | 1; }</pre>
void build(vector<Qnode> &v) { build(v, 1, 0, n - 1); }
void build(vector<Qnode> &v, int i, int l, int r) {
    if (l == r) {
        st[i] = v[l];
        return;
    int m = midpoint(l, r);
    build(v, tol(i), l, m);
    build(v, tor(i), m + 1, r);
    st[i] = merge(st[tol(i)], st[tor(i)], l, r);
void upd(int l, int r, Unode v) {
    al = l, ar = r:
    upd(1, 0, n - 1, v);
void upd(int i, int l, int r, Unode v) {
    prop(i, l, r);
    if (disjoint(l, r)) return;
    if (contains(l, r)) {
        applyV(i, v);
        prop(i, l, r);
        return;
    int m = midpoint(l, r);
    upd(tol(i), l, m, v);
    upd(tor(i), m + 1, r, v);
    st[i] = merge(st[tol(i)], st[tor(i)], l, r);
Qnode qry(int l, int r) {
    ql = l, qr = r;
    return gry(1, 0, n-1);
Qnode gry(int i, int l, int r) {
    prop(i, l, r);
```

```
if (disjoint(l, r)) return Qnode();
    if (contains(l, r)) return st[i];
    int m = midpoint(l, r);
    return merge(qry(tol(i), l, m), qry(tor(i), m + 1, r), l, r);
};
```

4.7 2D Sparse Table

```
const int N = 1001:
ll matrix[N][N];
ll M[1001][1001][10][10];
ll op(ll a, ll b) { return gcd(a, b); }
void SparseMatrix(int n, int m) {
           int i, j, x, y;
           for (i = 0; (1 << i) <= n; i++) {
                       for (j = 0; (1 << j) <= m; j++) {
                                  for (x = 0; (x + (1 << i) - 1) < n; x++) {
                                             for (y = 0; (y + (1 << j) - 1) < m; y++) {
                                                        if (i == 0 \&\& i == 0)
                                                                    M[x][y][i][i] = matrix[x][y];
                                                         else if (i == 0)
                                                                    M[x][y][i][j] = op(M[x][y][i][j-1],
                                                                                                                          M[x][y + (1 << (j - 1))][i][j -
             1]);
                                                         else if (i == 0)
                                                                   M[x][y][i][j] = op(M[x][y][i - 1][j],
                                                                                                                          M[x + (1 << (i - 1))][y][i -
          1][j]);
                                                         else {
                                                                    int tempa = op(M[x + (1 << (i - 1))][y][i - 1][i -
             1],
                                                                                                              M[x][y + (1 << (j - 1))][i - 1][j -
             1]);
                                                                    int tempb = op(M[x][y][i-1][j-1],
                                                                                                              M[x + (1 << (i - 1))][y 
          - 1))]
                                                                                                                 [i - 1][i - 1]);
                                                                   M[x][y][i][j] = op(tempa, tempb);
                                             }
           return;
int lg2(int x) { return sizeof(int) * 8 - builtin clz(x) - 1; }
ll query2d(int x, int y, int x1, int y1) {
           int k = lg2(x1 - x + 1);
           int l = lg2(y1 - y + 1);
           int tempa = op(M[x][y][k][l], M[x1 - (1 << k) + 1][y][k][l]);
           int tempb = op(M[x][y1 - (1 << l) + 1][k][l],
                                                     M[x1 - (1 \ll k) + 1][y1 - (1 \ll l) + 1][k][l]);
           return op(tempa, tempb);
```

```
}
```

4.8 Bitree 2D

Description: Given a 2D array you can increment an arbitrary position, and also query the subsum of a subgrid

Time: Update and query in $O(logN^2)$

```
struct Bit2d {
    int n;
    vll2d bit;
    Bit2d(int ni) : n(ni), bit(n + 1, vll(n + 1)) {}
   Bit2d(int ni, vll2d &xs) : n(ni), bit(n + 1, vll(n + 1)) {
        for (int i = 1: i <= n: i++) {
            for (int j = 1; j \le n; j++) {
                update(i, j, xs[i][j]);
        }
    void update(int x, int y, ll val) {
        for (; x \le n; x += (x \& (-x)))  {
            for (int i = y; i \le n; i += (i \& (-i))) {
                bit[x][i] += val;
        }
    il sum(int x, int y) {
        ll ans = 0:
        for (int i = x; i; i = (i & (-i))) {
            for (int j = y; j; j = (j \bar{k}'(-j))) {
                ans += bit[i][i];
            }
        return ans;
    il query(int x1, int y1, int x2, int y2) {
        return sum(x2, y2) - sum(x2, y1 - 1) - sum(x1 - 1, y2) +
               sum(x1 - 1, y1 - 1);
};
```

4.9 Convex Hull Trick / Line Container

Description: Container where you can add lines of the form mx + b, and query the maximum value at point x.

Usage: $insert_line(m, b)$ inserts the line $m \cdot x + b$ in the container.

eval(x) find the highest value among all lines in the point x.

Time: Eval and insert in $O(\log N)$

```
const ll LLINF = 1e18;
const ll is_query = -LLINF;
struct Line {
    ll m, b;
    mutable function<const Line *()> succ;
```

```
bool operator<(const Line &rhs) const {</pre>
        if (rhs.b != is query) return m < rhs.m;</pre>
        const Line *s = succ();
        if (!s) return 0:
        ll \dot{x} = rhs.m:
        return b - s -> b < (s -> m - m) * x;
};
struct Cht : public multiset<Line> { // maintain
                                        // max m*x+b
    bool bad(iterator y) {
        auto z = next(y);
        if (y == begin()) {
            if (z = end()) return 0;
            return y->m == z->m \&\& y->b <= z->b;
        auto x = prev(y);
        if (z == end()) return y->m == x->m && y->b <= x->b;
        return (ld)(x->b - v->b) * (z->m - v->m) >=
               (ld)(v->b-z->b) * (y->m-x->m);
    void insert line(ll m,
                     ll b) { // min -> insert (-m,-b) -> -eval()
        auto y = insert({m, b});
        y-succ = [=] { return next(y) == end() ? 0 : &*next(y); };
        if (bad(y)) {
            erase(y);
            return;
        while (next(y) != end() && bad(next(y))) erase(next(y));
        while (y != begin() && bad(prev(y))) erase(prev(y));
    il eval(ll x) {
        auto l = *lower bound((Line){x, is query});
        return l.m * x + l.b;
};
4.10 DSU (with rollback)
Description: Performs every operation a regular DSU does, but you can roll back to a
specific time.
Usage: int t = uf.time(); ...; uf.rollback(t); T
Time: O(\log(N))
struct RollbackUF {
    vi e:
    vector<pii> st;
    RollbackUF(int n) : e(n, -1) {}
    int size(int x) { return -e[find(x)]; }
    int find(int x) { return e[x] < 0 ? x : find(e[x]); }
    int time() { return len(st); }
    void rollback(int t) {
        for (int i = time(); i-- > t;) e[st[i].first] = st[i].second;
```

st.resize(t);

```
bool join(int a, int b) {
    a = find(a), b = find(b);
    if (a == b) return false;
    if (e[a] > e[b]) swap(a, b);
    st.push_back({a, e[a]});
    st.push_back({b, e[b]});
    e[a] += e[b];
    e[b] = a;
    return true;
}
```

4.11 DSU / UFDS

Usage: You may discomment the commented parts to find online which nodes belong to each set, it makes the $union_set$ method cost $O(log^2)$ instead O(A)

```
struct DSU {
    vector<int> ps, sz;
    // vector<unordered set<int>> sts;
    DSU(int N)
        : ps(N + 1),
          sz(N, 1) /*, sts(N) */
        iota(ps.begin(), ps.end(), 0);
        // for (int i = 0; i < N; i++)
        // sts[i].insert(i):
    int find set(int x) { return ps[x] == x ? x : ps[x] = find set(ps[x]);
    int size(int u) { return sz[find set(u)]; }
    bool same set(int x, int y) { return find set(x) == find set(y); }
    void union set(int x, int y) {
        if (same_set(x, y)) return;
        int px = find set(x);
        int py = find set(y);
        if (sz[px] < sz[py]) swap(px, py);</pre>
        ps[py] = px;
        sz[px] += sz[py];
        // sts[px].merge(sts[py]);
};
```

4.12 Lichao Tree (dynamic)

Description: Lichao Tree that creates the nodes dynamically, allowing to query and update from range [MAXL, MAXR]

Usage:

- query(x): find the highest point among all lines in the structure
- add(a,b): add a line of form y = ax + b in the structure
- addSegment(a,b,l,r) : add a line segment of form y=ax+b which covers from range [l,r]

Time: $O(\log N)$

```
template <typename T = ll, T MAXL = 0, T MAXR = 1 '000' 000'001>
struct LiChaoTree {
    static const T inf = -numeric_limits<T>::max() / 2;
    bool first_best(T a, T b) { return a > b; }
    T get best(T a, T b) { return first best(a, b) ? a : b; }
    struct line {
        T m, b;
        T operator()(T x) { return m * x + b; }
    struct node {
        line li;
        node *left. *right:
        node(line li = {0, inf}) : li( li), left(nullptr), right(nullptr)
    {}
        \simnode() {
            delete left:
            delete right;
    };
    node *root:
    LiChaoTree(line li = {0, inf}) : root(new node(li)) {}
    ~LiChaoTree() { delete root; }
    T query(T x, node *cur, T l, \bar{T} r) {
        if (cur == nullptr) return inf;
        if (x < l or x > r) return inf;
        T mid = midpoint(l, r);
        T ans = cur->li(x):
        ans = get best(ans, guery(x, cur->left, l, mid));
        ans = get best(ans, guery(x, cur->right, mid + 1, r));
        return ans:
    T query(T x) { return query(x, root, MAXL, MAXR); }
    void add(line li, node *&cur, T l, T r) {
        if (cur == nullptr) {
            cur = new node(li);
            return:
        T mid = midpoint(l, r);
        if (first best(li(mid), cur->li(mid))) swap(li, cur->li);
        if (first best(li(l), cur->li(l))) add(li, cur->left, l, mid);
        if (first best(li(r), cur->li(r))) add(li, cur->right, mid + 1, r)
    void add(T m, T b) { add({m, b}, root, MAXL, MAXR); }
    void addSegment(line li, node *&cur, T l, T r, T lseg, T rseg) {
        if (r < lseg || l > rseg) return;
        if (cur == nullptr) cur = new node;
        if (lseg \leftarrow 1 && r \leftarrow rseg) {
            add(li, cur, l, r);
            return;
        T mid = midpoint(l, r):
        if (l != r) {
            addSegment(li, cur->left, l, mid, lseg, rseg);
```

```
addSegment(li, cur->right, mid + 1, r, lseg, rseg);
}

void addSegment(T a, T b, T l, T r) {
    addSegment({a, b}, root, MAXL, MAXR, l, r);
};
```

4.13 Merge sort tree

Description: Like a segment tree but each node stores the ordered subsegment it represents.

Usage:

• inrange(l, r, a, b): counts the number of positions $i, l \le i \le r$ such that $a \le x_i \le b$. **Time**: Build $O(N \log N^2)$, inrange $O(\log N^2)$

Memory: $O(n \log N)$

```
template <class T>
struct MergeSortTree {
    int n:
    vector<vector<T>> st;
    MergeSortTree(vector<T> &xs) : n(len(xs)), st(n << 1) {</pre>
        rep(i, 0, n) st[i + n] = vector<T>({xs[i]});
        rrep(i, n - 1, 0) {
            st[i].resize(len(st[i << 1]) + len(st[i << 1 | 1]));
            merge(all(st[i << 1]), all(st[i << 1 | 1]), st[i].begin());
    int count(int i, T a, T b) {
        return upper bound(all(st[i]), b) - lower bound(all(st[i]), a);
    int inrange(int l, int r, T a, T b) {
        int ans = 0:
        for (l += n, r += n + 1; l < r; l >>= 1, r >>= 1) {
            if (l & 1) ans += count(l++, a, b);
            if (r \& 1) ans += count(--r, a, b);
        return ans;
};
```

4.14 Mex with update

Description: This DS allows you to mantain an array of elments, insert, and remove, and query the MEX at any time.

Ûsage:

- Mex(mxsz): Initialize the DS, mxsz must be the maximum number of elements that the structure may have.
- add(x): just adds one copy of x.
- rmv(x): just remove a copy of x.
- operator(): returns the MEX.

Time:

• Mex(mxsz): $O(\log mxsz)$

- add(x): $O(\log mxsz)$
- rmv(x): $O(\log mxsz)$
- *operator()*: *O*(1)

```
struct Mex {
    int mx_sz;
    vi hs:
    set<int> st;
    Mex(int mx sz) : mx sz(mx sz), hs(mx sz + 1) {
        auto it = st.begin();
        rep(i, 0, mx sz + 1) it = st.insert(it, i);
    void add(int x) {
        if (x > mx sz) return;
        if (!hs[x]++) st.erase(x);
    void rmv(int x) {
        if (x > mx sz) return;
        if (!--hs[x]) st.emplace(x);
    int operator()() const { return *st.begin(); }
    /*
      Optional, you can just create with size
      len(xs) add N elements :D
    Mex(const vi &xs, int mx sz = -1) : Mex(\sim mx sz ? mx sz : len(xs)) {
        for (auto xi : xs) add(xi):
};
```

4.15 Orderd Set (GNU PBDS)

Usage: If you need an ordered multi set you may add an id to each value. Using greater equal, or less equal is considered undefined behavior.

- order_of_key (k): Number of items strictly smaller/greater than k.
- find_by_order(k): K-th element in a set (counting from zero).

Time: Both $O(\log N)$

Warning: Is 2 or 3 times slower then a regular set/map.

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
template <typename T>
using ordered_set =
    tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
```

4.16 Prefix Sum 2D

Description: Given an 2D array with N lines and M columns, find the sum of the subarray that have the left upper corner at (x1, y1) and right bottom corner at (x2, y2). **Time**: Build $O(N \cdot M)$, Query O(1).

```
template <typename T>
struct psum2d {
   vector<vector<T>> s;
   vector<vector<T>> psum;
   psum2d(vector<vector<T>> &grid, int n, int m)
        s(n + 1, vector < T > (m + 1)), psum(n + 1, vector < T > (m + 1)) {
        for (int i = 1; i \le n; i++)
            for (int j = 1; j <= m; j++) {
                s[i][i] = s[i][i-1] + qrid[i-1][i-1];
                psum[i][j] = psum[i - 1][j] + s[i][j];
   T query(int x1, int y1, int x2, int y2) \{
        T ans = psum[x2 + 1][y2 + 1] + psum[x1][y1];
        ans -= psum[x2 + 1][y1] + psum[x1][y2 + 1];
        return ans:
   }
};
```

4.17 Segment Tree Update Range Query (bottom-up)

```
* @Description:
        merge should be function<T(T,T)>, that
        makes the necessary operation between two
        nodes in the segment tree
 * */
#include "../../Contest/template.cpp"
template <typename T, T identity, auto merge>
struct SegTreeBottomUp {
   int size;
   vector<T> arr;
   SegTreeBottomUp(int n) {
        for (size = 1; size < n; size <<= 1);</pre>
        arr.resize(size << 1);
   void assign(int pos, const T &val) {
        for (arr[pos += size] = val; pos >>= 1;)
            arr[pos] = merge(arr[pos << 1], arr[pos << 1 | 1]);</pre>
   T query(int l, int r) {
       T ans l = identity, ans r = identity;
        for (l += size, r += size + 1; l < r; l >>= 1, r >>= 1) {
            if (l & 1) ans l = merge(ans l, arr[l++]);
            if (r \& 1) ans r = merge(arr[--r], ans r);
        return merge(ans_l, ans_r);
```

```
}
SegTreeBottomUp(const vector<T> &vec) : SegTreeBottomUp(len(vec)) {
    copy(all(vec), begin(arr) + size);
    rrep(i, size - 1, 0) arr[i] = merge(arr[i << 1], arr[i << 1 | 1]);
};
using SegTreeBottomUpSumQuery =
    SegTreeBottomUp<ll, Oll, [](ll a, ll b) { return a + b; }>;

4.18 Sparse table

template <typename T = ll,
</pre>
```

```
vi logs;
vector<vector<T>> st;
public:
   SparseTable(const vector<T> &v) : sz(len(v)), logs(sz + 1) {
      rep(i, 2, sz + 1) logs[i] = logs[i >> 1] + 1;
      st.resize(logs[sz] + 1, vector<T>(sz));
      rep(i, 0, sz) st[0][i] = v[i];
      for (int k = 1; (1 << k) <= sz; k++) {</pre>
```

for (int i = 0; $i + (1 << k) <= sz; i++) {$

cmp(st[k-1][i], st[k-1][i+(1 << (k-1))], st[k][i])

```
}
}
T query(int l, int r) {
    r++;
    const int k = logs[r - l];
    T ret;
    cmp(st[k][l], st[k][r - (1 << k)], ret);
    return ret;
}</pre>
```

4.19 Static range queries

int sz:

};

```
}
T operator()(int l, int r) {
    T lv = (l ? acc[l - 1] : T());
    T ret;
    invop(acc[r], lv, ret);
    return ret;
}
```

4.20 Venice Set

Description: A container that you can insert q copies of element e, increment every element in the container in x, query which is the best element and it's quantity and also remove k copies of the greatest element.

Time:

- add elment O(log N)
 remove O(log N)
 update: O(1)
- query O(1)

```
template <typename T = ll>
struct VeniceSet {
    using T2 = pair<T, ll>;
    priority queue<T2, vector<T2>, greater<T2>> pg;
    VeniceSet() : acc() {}
    void add element(const T &e, const ll q) { pq.emplace(e - acc, q); }
    void update all(const T &x) { acc += x; }
    T2 best() {
        auto ret = pg.top();
        ret.first += acc:
        return ret;
   void pop() { pq.pop(); }
    void pop k(int k) {
        auto [e, q] = pq.top();
        pq.pop();
        q -= k;
        if (q) pq.emplace(e, q);
};
```

4.21 Venice Set (complete)

Description: A container which you can insert elements update all at once and also make a few queries

Usage:

- $add_element(e, q)$: adds q copies of e, if no q is provided adds a single one
- $update_all(x)$: increment every value by x
- erase(e): removes every copy of e, and returns how much was removed.
- count(e): returns the number of e in the container

- high()/low(): returns the hightest/lowest element, and it's quantity
- $pop_low(q)/pop_high(q)$: removes q copies of the lowest/highest elements if no q is provided removes all copies of the lowest/highest element.

You may answer which is the K-th value and it's quantity using an $ordered_set$. Probably works with other operations

Time: Considering N the number of distinct numbers in the container

- $add\ element(e,q):\ O(\log(N))$
- $update \ all(x):O(1)$
- erase(e): $O(\log(N))$
- count(e): $O(\log(N))$
- high()/low(): O(1)
- $pop_low(q)/pop_high(q)$: worst case is $O(N \cdot \log{(N)})$ if you remove all elements and so on...

Warning: There is no error handling if you try to *pop* more elements than exists or related stuff

```
struct VeniceSet {
    set<pll> st;
    ll acc;
    VeniceSet() : acc() {}
    ll add element(ll e, ll q = 1) {
        q = erase(e);
        ė −= acc;
        st.emplace(e, q);
        return q;
    void update all(ll x) { acc += x; }
    ll erase(ll e) {
        e = acc;
        auto it = st.lb({e, LLONG MIN});
        if (it == end(st) | | (*it).first != e) return 0;
        ll ret = (*it).second;
        st.erase(it);
        return ret;
    ll count(ll x) {
        x \rightarrow acc;
        auto it = st.lb({x, LLONG MIN});
        if (it == end(st) || (*it).first != x) return 0;
        return (*it).second;
    pll high() { return *rbegin(st); }
    pll low() { return *begin(st); }
    void pop high(ll q = -1) {
        if (\overline{q} == -1) q = high().second;
        while (a) {
            auto [e, eq] = high();
            st.erase(prev(end(st)));
            if (eq > q) add_element(e, eq - q);
            q = max(0ll, q - eq);
        }
    void pop_low(ll q = -1) {
```

```
if (q == -1) q = low().second;
while (q) {
    auto [e, eq] = low();
    st.erase(st.begin());
    if (eq > q) add_element(e, eq - q);
    q = max(0ll, q - eq);
}
};
```

4.22 Wavelet tree

```
using ll = long long;
template <typename T>
struct WaveletTree {
   struct Node {
       T lo, hi;
       int left child, right child;
       vector<int> pcnt;
       vector<ll> psum;
       Node(int lo_, int hi_)
            : lo(lo), hi(hi), left child(0), right child(0), pcnt(),
   psum() {}
   };
   vector<Node> nodes;
   WaveletTree(vector<T> v) {
       nodes.reserve(2 * v.size());
       auto [mn, mx] = minmax element(v.begin(), v.end());
        auto build = [&](auto &&self, Node &node, auto from, auto to) {
            if (node.lo == node.hi or from >= to) return;
            auto mid = midpoint(node.lo, node.hi);
            auto f = [&mid](T x) { return x <= mid; };</pre>
            node.pcnt.reserve(to - from + 1);
            node.pcnt.push back(0);
            node.psum.reserve(to - from + 1);
            node.psum.push_back(0);
            T left upper = node.lo, right lower = node.hi;
            for (auto it = from; it != to; it++) {
                auto value = f(*it);
                node.pcnt.push_back(node.pcnt.back() + value);
                node.psum.push back(node.psum.back() + *it);
                if (value)
                    left upper = max(left upper, *it);
                    right lower = min(right lower, *it);
            }
            auto pivot = stable partition(from, to, f);
            node.left child = make node(node.lo, left upper);
            self(self, nodes[node.left_child], from, pivot);
            node.right child = make node(right lower, node.hi);
            self(self, nodes[node.right_child], pivot, to);
        build(build, nodes[make_node(*mn, *mx)], v.begin(), v.end());
```

```
T kth element(int L, int R, int K) const {
        auto f = [&](auto &&self, const Node &node, int l, int r, int k)
   -> T {
            if (l > r) return 0;
            if (node.lo == node.hi) return node.lo;
            int lb = node.pcnt[l], rb = node.pcnt[r + 1], left size = rb -
    lb:
            return (left size > k
                         ? self(self, nodes[node.left child], lb, rb - 1, k
                         : self(self, nodes[node.right_child], l - lb, r -
    rb.
                                k - left size));
        return f(f, nodes[0], L, R, K);
    pair<int, ll> count and sum in range(int L, int R, T a, T b) const {
        auto f = [&](auto &&self, const Node &node, int l,
                     int r) -> pair<int, ll> {
            if (l > r \text{ or node.} lo > b \text{ or node.} hi < a) return {0, 0};
            if (a <= node.lo and node.hi <= b)</pre>
                return \{r - l + 1,
                         (node.lo == node.hi ? (r - l + 1ll) * node.lo
                                             : node.psum[r + 1] - node.psum
    [l])};
            int lb = node.pcnt[l], rb = node.pcnt[r + 1];
            auto [left cnt, left sum] =
                self(self, nodes[node.left child], lb, rb - 1);
            auto [right_cnt, right sum] =
                self(self, nodes[node.right_child], l - lb, r - rb);
            return {left cnt + right cnt, left sum + right sum};
        return f(f, nodes[0], L, R);
    inline int count in range(int L, int R, T a, T b) const {
        return count_and_sum_in_range(L, R, a, b).first;
    inline ll sum in range(int L, int R, T a, T b) const {
        return count and sum in range(L, R, a, b).second;
   private:
    int make_node(T lo, T hi) {
        int id = (int)nodes.size();
        nodes.emplace back(lo, hi);
        return id;
};
```

5 Dynamic Programming

5.1 Binary Knapsack (bottom up)

Description: Given the points each element have, and it repespective cost, computes the maximum points we can get if we can ignore/choose an element, in such way that the sum of costs don't exceed the maximum cost allowed.

Time: O(N*W)

Warning: The vectors VS and WS starts at one, so it need an empty value at index 0.

```
const int MAXN(1'000), MAXCOST(1'000 * 20);
ll dp[MAXN + 1][MAXCOST + 1];
bool ps[MAXN + 1][MAXCOST + 1];
pair<ll, vi> knapsack(const vll &points, const vi &costs, int maxCost) {
   int n = len(points) - 1; // ELEMENTS START AT INDEX 1 !
    for (int m = 0; m <= maxCost; m++) {</pre>
        dp[0][m] = 0;
   for (int i = 1; i <= n; i++) {
        dp[i][0] = dp[i - 1][0] + (costs[i] == 0) * points[i];
        ps[i][0] = costs[i] == 0;
   for (int i = 1; i <= n; i++) {
        for (int m = 1; m <= maxCost; m++) {</pre>
            dp[i][m] = dp[i - 1][m], ps[i][m] = 0;
            int w = costs[i]:
            ll v = points[i];
            if (w \le m \text{ and } dp[i - 1][m - w] + v > dp[i][m]) {
                dp[i][m] = dp[i - 1][m - w] + v, ps[i][m] = 1;
        }
   }
   vi is;
   for (int i = n, m = maxCost; i >= 1; --i) {
        if (ps[i][m]) {
            is.emplace back(i);
            m -= costs[i];
    return {dp[n][maxCost], is};
```

5.2 Edit Distance

```
Time: O(N * M)
#include "../Contest/template.cpp"

ll edit_distance(const string &a, const string &b) {
   int n = a.size();
   int m = b.size();
   vll2d dp(n + 1, vi(m + 1, 0));
   const ll ADD = 1, DEL = 1, CHG = 1;
   for (int i = 0; i <= n; ++i) {</pre>
```

```
dp[i][0] = i * DEL;
}
for (int i = 1; i <= m; ++i) {
    dp[0][i] = ADD * i;
}

for (int i = 1; i <= n; ++i) {
    for (int j = 1; j <= m; ++j) {
        int add = dp[i][j - 1] + ADD;
        int del = dp[i - 1][j] + DEL;
        int chg = dp[i - 1][j - 1] + (a[i - 1] != b[j - 1]) * CHG;
        dp[i][j] = min({add, del, chg});
    }
}
return dp[n][m];</pre>
```

5.3 Knapsack

Description: Finds the maximum score you can achieve, given that you have N items, each item has a cost, a point and a quantity, you can spent at most maxcost and buy each item the maximum quantity it has.

Time: $O(n \cdot maxcost \cdot \log maxqtd)$

Memory: O(maxcost).

```
ll knapsack(const vi &weight, const vll &value, const vi &qtd, int maxCost
    vi costs;
    vll values;
    for (int i = 0; i < len(weight); i++) {
         ll a = atd[i];
         for (ll x = 1; x \le q; q = x, x \le 1) {
             costs.eb(x * weight[i]);
             values.eb(x * value[i]);
        if (q) {
             costs.eb(q * weight[i]);
             values.eb(q * value[i]);
    vll dp(maxCost + 1):
    for (int i = 0; i < len(values); i++) {
   for (int j = maxCost; j > 0; j--) {
             if (i \ge costs[i]) dp[i] = max(dp[i], values[i] + dp[i - costs])
    [i]]);
    return dp[maxCost];
```

5.4 Longest Increasing Subsequence

Description: Find the pair (sz, psx) where sz is the size of the longest subsequence and psx is a vector where psx_i tells the size of the longest increase subsequence that ends at

position i. $get_i dx$ just tells which indices could be in the longest increasing subsequence. **Time**: $O(n \log n)$

```
#include "../Contest/template.cpp"
template <tvpename T>
pair<int, vi> lis(const vector<T> &xs, int n) {
   vector<T> dp(n + 1, numeric limits<T>::max());
   dp[0] = numeric limits<T>::min();
   int sz = 0:
   vi psx(n);
   rep(i, 0, n) {
       int pos = lower bound(all(dp), xs[i]) - dp.begin();
       sz = max(sz, pos);
       dp[pos] = xs[i];
       psx[i] = pos;
    return {sz, psx};
template <typename T>
vi get idx(vector<T> xs) {
   int n = xs.size();
   auto [sz1, psx1] = lis(xs, n);
   transform(rall(xs), xs.begin(), [](T x) { return -x; });
   auto [sz2, psx2] = lis(xs, n);
   vi ans;
   rep(i, 0, n) {
       int l = psx1[i];
       int r = psx2[n - i - 1];
       if (l + r - 1 == sz1) ans.eb(i);
    return ans;
```

5.5 Monery sum

Description: Find every possible sum using the given values only once.

```
}
}
for (int i = 0; i < _n; ++i)
    for (int j = 0; j <= _m; ++j)
        if (_dp[i][j]) _ans.insert(j);
return _ans;
}</pre>
```

5.6 Steiner tree

```
template <typename T>
T steinerCost(const vector<vector<T>> &adj, const vi ks,
              T inf = numeric limits<T>::max()) {
    int k = len(ks), n = len(\overline{adj});
    vector<vector<T>> dp(n, vector<T>(1 << k, inf));</pre>
    vi inks(n):
    trav(ki, ks) inks[ki] = 1;
    trav(ki, ks) {
        rep(j, 0, n) {
            if (count(all(ks), j) == 0) {
                dp[j][1 << ki] = adj[ki][j];
        }
    rep(mask, 2, (1 << k)) {
        rep(i, 0, n) {
            if (inks[i]) continue;
            for (int mask2 = (mask - 1) & mask; mask2 >= 1;
                 mask2 = (mask2 - 1) \& mask) {
                int mask3 = mask ^ mask2;
                chmin(dp[i][mask], dp[i][mask2] + dp[i][mask3]);
            rep(j, 0, n) {
                if (inks[j]) continue;
                chmin(dp[i][mask], dp[i][mask] + adi[i][i]);
        }
    T ans = inf:
    rep(i, 0, n) chmin(ans, dp[i][(1 << k) - 1]);
    return ans;
```

5.7 Sum of Subsets

Description: Allows you to find if some mask X is a super mask of any of the given masks **Usage**: Call *build* with the masks then it returns a vector of bool V where V_X says if X is a super mask of any of the initial maks

You can change it to count how many submasks of each mask exsists, by changing the bitwise or by a plus sign...

Time: $O(LOG \cdot 2^{LOG})$ Memory: $O(LOG^2 \cdot 2^{LOG})$ Warning: Remember to set LOG with the highest bit possible

```
const int LOG = 20;
vc build(const vi &masks) {
    vc ret(1 << LOG);
    trav(mi, masks) ret[mi] = 1;
    rep(b, 0, LOG) {
        rep(mask, 0, (1 << LOG)) {
            if (mask & (1 << b)) ret[mask] |= ret[mask ^ (1 << b)];
        }
    }
    return ret;
}</pre>
```

5.8 Travelling Salesman Problem

```
Time: O(N^2 \cdot 2^N)

Memory: O(N^2 \cdot 2^N)

vll2d dist;

vll memo;

int tsp(int i, int mask, int N) {

   if (mask == (1 << N) - 1) return dist[i][0];

   if (memo[i][mask] != -1) return memo[i][mask];

   int ans = INT_MAX << 1;

   for (int j = \overline{0}; j < N; ++j) {

      if (mask & (1 << j)) continue;

      auto t = tsp(j, mask | (1 << j), N) + dist[i][j];

      ans = min(ans, t);
   }

   return memo[i][mask] = ans;
}
```

6 Extras

6.1 Binary to gray

```
string binToGray(string bin) {
    string gray(bin.size(), '0');
    int n = bin.size() - 1;
    gray[0] = bin[0];
    for (int i = 1; i <= n; i++) {
        gray[i] = '0' + (bin[i - 1] == '1') ^ (bin[i] == '1');
    }
    return gray;
}</pre>
```

6.2 Get permutation cycles

Description: Receives a permutation [0, n-1] and return a vector 2D with each cycle.

```
vll2d getPermutationCicles(const vll &ps) {
    ll n = len(ps);
    vector<char> visited(n);
    vector<vll> cicles;
    rep(i, 0, n) {
        if (visited[i]) continue;
        vll cicle;
        ll pos = i;
        while (!visited[pos]) {
            cicle.pb(pos);
            visited[pos] = true;
            pos = ps[pos];
        }
        cicles.push_back(vll(all(cicle)));
    }
    return cicles;
}
```

6.3 Max & Min Check

Description: Returns the min/max value in range [l, r] that satisfies the lambda function check, if there is no such value the 'nullopt' is returned.

Usage: check must be a function that receives an integer and return a boolean.

Time: $O(\log r - l + 1)$

```
template <typename T>
optional<T> maxCheck(T l, T r, auto check) {
    optional<T> ret;
    while (l <= r) {
        T m = midpoint(l, r);
        if (check(m))
            ret ? chmax(ret, m) : ret = m, l = m + 1;
        else
            r = m - 1:
    return ret;
template <typename T>
optional<T> minCheck(T l, T r, auto check) {
    optional<T> ret;
    while (l <= r) {
        T m = midpoint(l, r);
        if (check(m))
            ret ? chmin(ret, m) : ret = m, r = m - 1;
        else
            l = m + 1;
    return ret;
```

6.4 Merge Interals

Time: $(N \log N)$

```
Warning: It destroys the original array
#include "../Contest/template.cpp"
template <typename T>
vector<pair<T, T>> merge intervals(vector<pair<T, T>> &intervals) {
    if (!len(intervals)) return {};
    using Pt = pair<T, T>;
    sort(all(intervals)):
    vector<Pt> ret{intervals.front()};
    rep(i, 1, len(ret)) {
        auto &[pl, pr] = ret.back();
        auto &[l, r] = intervals[i];
        if (l <= pr)
            chmax(pr, r);
        else
            ret.eb(l, r);
    return ret;
```

6.5 Mo's algorithm

```
template <typename T, typename Tans>
struct Mo {
    struct Query {
        int l, r, idx, block;
        Ouery(int l, int _r, int _idx, int _block)
            : l(l), r(r), idx(idx), block(block) {}
        bool operator<(const Query &g) const {</pre>
            if (block != q.block) return block < q.block;</pre>
            return (block & 1 ? (r < q.r) : (r > q.r));
   };
   vector<T> vs:
   vector<Query> qs;
   const int block size;
   Mo(const vector<T> &a) : vs(a), block size((int)ceil(sqrt(a.size())))
   {}
   void add query(int l, int r) {
        qs.emplace_back(l, r, qs.size(), l / block_size);
   auto solve() {
        // get answer return type
        vector<Tans> answers(qs.size());
        sort(all(qs));
        int cur l = 0, cur r = -1;
        for (auto q : qs) {
            while (cur_l > q.l) add(--cur_l);
```

```
while (cur_r < q.r) add(++cur_r);
    while (cur_l < q.l) remove(cur_l++);
    while (cur_r > q.r) remove(cur_r--);
        answers[q.idx] = get_answer();
    }
    return answers;
}

private:
    // add value at idx from data structure
    inline void add(int idx) {}
    // remove value at idx from data structure
    inline void remove(int idx) {}
    // extract current answer of the data structure
    inline Tans get_answer() {}
};
```

6.6 int128t stream

```
void print( int128 x) {
    if (x < 0) {
        cout << '-':
        X = -X;
    if (x > 9) print(x / 10);
    cout << (char)((x % 10) + '0');
int128 read() {
    string s;
    cin >> s:
    int128 x = 0;
    for (auto c : s) {
        if (c != '-') x += c - '0';
        x *= 10;
    x /= 10:
    if (s[0] == '-') x = -x;
    return x;
```

7 Geometry

7.1 All i know about 2D stuff

```
enum PointPosition { IN, ON, OUT };
template <class Point>
vector<Point> segInter(Point a, Point b, Point c, Point d);
template <typename T>
bool equals(T a, T b) {
    if (std::is floating_point<T>::value)
        return \overline{f}abs(a -\overline{b}) < EPS;
    else
        return a == b:
template <class T>
struct Point {
    typedef Point P;
    T x, y;
    explicit Point(T x = 0, T y = 0) : x(x), y(y) {}
    bool operator<(P p) { return tie(x, y) < tie(p.x, p.y); }</pre>
    bool operator>(P& rhs) { return rhs < *this; }</pre>
    bool operator==(P p) { return tie(x, y) == tie(p.x, p.y); }
    P operator+(P p) { return P(x + p.x, y + p.y); }
    P operator-(P p) { return P(x - p.x, y - p.y); }
    P operator*(T d) { return P(x * d, y * d); }
    P operator/(T d) { return P(x / d, y / d); }
    T dot(P p) \{ return x * p.x + y * p.y; \}
    T cross(P p) { return x * p.y - y * p.x; }
    T cross(P a, P b) { return (a - *this).cross(b - *this); }
    T dist2() { return x * x + y * y; }
    double dist() { return sqrt((double)dist2()); }
    // angle to x-axis in interval [-pi, pi]
    double angle() { return atan2(y, x); }
    P unit() { return *this / dist(); } // makes dist()=1
    P perp() \{ return P(-y, x); \}
                                         // rotates +90 degrees
    P normal() { return perp().unit(); }
    // returns point rotated 'a' radians ccw around
    // the origin
    P rotate(double a) {
        return P(x * cos(a) - y * sin(a), x * sin(a) + y * cos(a));
    pair<T, T> slope(Point<T>& o) {
        auto a = o.x - x;
        auto b = o.y - y;
        if (!is floating point<T>::value) {
            auto g = gcd(a, b);
            if (a) a /= a, b /= a:
        return {b, a};
    friend ostream& operator<<(ostream& os, P p) {</pre>
        return os << "(" << p,x << "," << p,v << ")";
    double distanceTo(Point<T>& other) {
        return hypot(other.x - x, other.y - y);
};
```

```
template <typename T>
struct Line {
    T a, b, c;
    Point<T> p1, p2;
    Line(T a = 0, T b = 0, T c = 0) : a(a), b(b), c(c) {
        if (a != 0) {
            double x = 0;
            double y = (-c) / b;
            p1 = Point < T > (x, v):
        if (b != 0) {
            double y = 0;
            double x = (-c) / a;
            p2 = Point < T > (x, y);
    Line(Point<T>& p, Point<T>& q) {
        a = p.y - q.y;
        b = q.x - p.x;
        c = p.cross(q);
        p1 = p, p2 = q;
    bool operator==(Line<T>& other) {
        return tie(a, b, c) == tie(other.a, other.b, other.c);
    // Less-than operator
    bool operator<(Line& rhs) {</pre>
        return tie(a, b, c) < tie(rhs.a, rhs.b, rhs.c);</pre>
    bool operator>(Line& rhs) { return rhs < *this; }</pre>
    Line<T> norm() {
        T d = a == 0 ? b : a;
        return Line(a / d, b / d, c / d);
    bool contains(PointT>& p) { return equals(a * p.x + b * p.y + c, (T)
    0); }
    bool parallel(Line<T>& r) {
        auto det = a * r.b - b * r.a;
        return equals(det, 0) and !(*this == r);
    bool orthogonal(Line<T>& r) { return equals(a * r.a + b * r.b, 0); }
    T direction(Point<T>& p3) { return p1.cross(p2, p3); }
    friend ostream& operator<<(ostream& os. Line 1) {</pre>
        return os << fixed << setprecision(6) << "(" << l.a << "," << l.b</pre>
    << ","
                  << l.c << ")":
    double distance(Point<T>& p) {
        return (a * p.x + b * p.y + c) / hypot(a, b);
    Point<T> closest(Point<T>& p) {
        auto den = (a * a + b * b):
```

```
auto x = (b * (b * p.x - a * p.y) - a * c) / den;
                                                                                    Point<T> c:
        auto v = (a * (-b * p.x + a * p.v) - b * c) / den:
                                                                                    Tr;
                                                                                    Circle(Point<T> c, T r) : c( c), r( r) {}
        return Point<T>{x, y};
                                                                                    Circle(T r) : Circle(\overline{Point} < T > (0, 0), r) {}
                                                                                    ld area() const { return PI * r * r; }
};
                                                                                    ld perimeter() const { return 2.0 * PI * r; }
                                                                                    ld arc(ld theta) const { return theta * r; }
   */
                                                                                    ld chord(ld theta) const { return 2.0 * r * sin(theta / 2.0); }
                                                                                    ld sector(ld theta) const { return (theta * r * r) / 2.0; }
template <typename T>
                                                                                    ld segment(ld theta) const { return ((theta - sin(theta)) * r * r) /
struct LineSeament {
    Point<T> p1, p2:
                                                                                    2.0:  }
                                                                                    PointPosition position(const Point<T>& p) const {
   LineSegment(Point<T> p, Point<T> q) { p1 = p, p2 = q; }
                                                                                        auto d = c.dist(p);
    LineSegment(T a, T b, T c, T d)
                                                                                        return equals(d, r) ? ON : (d < r ? IN : OUT);</pre>
        : LineSegment(Point<T>(a, b), Point<T>(c, d)) {}
    bool operator==(LineSegment<T>& other) {
                                                                                };
        return tie(p1, p2) == tie(other.p1, other.p2);
                                                                                /*
    // Less-than operator
    bool operator<(LineSegment& rhs) {</pre>
                                                                                template <typename T>
        return tie(p1, p2) < tie(rhs.p1, rhs.p2);</pre>
                                                                                struct Rectangle {
                                                                                    Point<T> P, 0;
    bool operator>(LineSegment& rhs) { return rhs < *this; }</pre>
                                                                                    T b, h;
   T direction(Point<T>& p3) { return p1.cross(p2, p3); }
                                                                                    Rectangle(const Point<T>& p. const Point<T>& a) : P(P), O(a) {
    friend ostream& operator<<(ostream& os, LineSegment l) {</pre>
                                                                                        assert(P != 0):
        return os << "(" << l.p1 << "," << l.p2 << ")";
                                                                                        b = \max(P.x, Q.x) - \min(P.x, Q.x);
                                                                                        h = max(P.y, Q.y) - min(P.y, Q.y);
    vector<Point<T>> intersection(LineSegment<T>& other) {
        return seaInter(pl. p2, other.p1, other.p2):
                                                                                    Rectangle(T base, T height)
                                                                                        : P(0, 0), Q(base, height), b(base), h(height) {}
                                                                                    T perimeter() const { return 2 * b + 2 * h; }
    // Verifica se o ponto P da reta r que écontm A e B pertence ao
                                                                                    T area() const { return b * h; }
   seamento
                                                                                    optional<Rectangle> intersection(const Rectangle& r) const {
    bool contains(Point<T>& P) {
                                                                                        using pt = pair<T, T>;
        return equals(p1.x, p2.x)
                   ? \min(p1.y, p2.y) \le P.y \text{ and } P.y \le \max(p1.y, p2.y)
                                                                                        auto i = pt(min(P.x, Q.x), max(P.x, Q.x));
                   : min(p1.x, p2.x) \le P.x \text{ and } P.x \le max(p1.x, p2.x);
                                                                                        auto u = pt(min(r.P.x, r.Q.x), max(r.P.x, r.Q.x));
                                                                                        auto a = max(i.first, u.first);
                                                                                        auto b = min(i.second, u.second);
    // Ponto mais óprximo de P no segmento AB
                                                                                        i = pt(min(P.y, Q.y), max(P.y, Q.y));
    Point<T> closest(Point<T>& P) {
                                                                                        u = pt(min(r.P.y, r.Q.y), max(r.P.y, r.Q.y));
        Line<T> r(p1, p2);
                                                                                        auto c = max(i.first, u.first);
        auto Q = r.closest(P);
                                                                                        auto d = max(i.second, u.first);
        if (this->contains(0)) return 0;
                                                                                        if (d < c or b < a) return nullopt:
        auto distp1 = P.distanceTo(p1);
                                                                                        return {{a, c}, {b, d}};
        auto distp2 = P.distanceTo(p2);
        if (distp1 <= distp2)</pre>
                                                                                };
            return p1:
        else
            return p2;
                                                                                template <typename T>
};
                                                                                struct Trapezium {
                                                                               ___ T B, b, h;
                                                                                    T area() const { return ((b + B) * h) / 2; }
                                                                                };
template <typename T>
                                                                               /*
struct Circle {
```

```
template <tvpename T>
struct Triangle {
   Point<T> A, B, C;
   enum SidesClass { EQUILATERAL, ISOCELES, SCALENE };
   SidesClass classification by sides() const {
        auto a = A.distanceTo(B);
        auto b = B.distanceTo(C);
        auto c = C.distanceTo(A):
        if (equals(a, b) && equals(b, c)) return EQUILATERAL;
        if (equals(a, b) or equals(a, c) or equals(b, c)) return ISOCELES:
        return SCALENE;
   enum AnglesClass { RIGHT, ACUTE, OBTUSE };
   AnglesClass classification by angles() const {
        auto a = dist(A, B);
        auto b = dist(B, C);
        auto c = dist(C, A):
        auto alpha = acos((a * a - b * b - c * c) / (-2 * b * c));
        auto beta = acos((b * b - a * a - c * c) / (-2 * a * c));
        auto gamma = acos((c * c - a * a - b * b) / (-2 * a * b));
        auto right = PI / 2.0:
        if (equals(alpha, right) || equals(beta, right) || equals(gamma,
   right))
            return RIGHT:
        if (alpha > right || beta > right || gamma > right) return OBTUSE;
        return ACUTE:
   }
   double perimeter() const {
        auto a = dist(A, B), b = dist(B, C), c = dist(C, A);
        return a + b + c;
   double area() const {
        Line<T> r(A, B);
        auto b = dist(A, B);
        auto h = r.distance(C);
        return (b * h) / 2;
};
template <typename T>
Point<T> triangleBarycenter(const Point<T>& a, const Point<T>& b,
                            const Point<T>& c) {
    return Point<T>((a.x + b.x + c.x) / 3.0, (a.y + b.y + c.y) / 3.0);
template <typename T>
Point<T> triangleOrthocenter(const Point<T>& a, const Point<T>& b,
                             const Point<T>& c) {
   LineT> r(a, b), s(a, c):
   Line<T> u\{r.b, -r.a, -(c.x * r.b - c.y * r.a)\};
   Line<T> v\{s.b, -s.a, -(b.x * s.b - b.y * s.a)\};
   auto det = u.a * v.b - u.b * v.a:
```

```
=== auto x = (-u.c * v.b + v.c * u.b) / det;
    auto y = (-v.c * u.a + u.c * v.a) / det;
    return {x, y};
template <typename T>
Point<double> triangleIncenter(const Point<T>& a, const Point<T>& b,
                               const Point<T>& c) {
    auto dab = distance(a, b):
    auto dbc = distance(b, c);
    auto dca = distance(c, a);
    auto p = dab + dbc + dca;
    auto \dot{x} = (a.x * dab + b.x * dbc + b.x * dca) / (p);
    auto y = (a.y * dab + b.y * dbc + b.y * dca) / (p);
    return Point<double>(x, v);
template <typename T>
Point<T> triangleCircumcenter(const Point<T>& A, const Point<T>& B,
                              const Point<T>& C) {
    auto D = 2 * (A.x * (B.v - C.v) + B.x * (C.v - A.v) + C.x * (A.v - B.v)
    auto A2 = A.x * A.x + A.v * A.v:
    auto B2 = B.x * B.x + B.y * B.y;
    auto C2 = C.x * C.x + C.y * C.y;
    auto x = (A2 * (B.y - C.y) + B2 * (C.y - A.y) + C2 * (A.y - B.y)) / D;
    auto y = (A2 * (C.x - B.x) + B2 * (A.x - C.x) + C2 * (B.x - A.x)) / D:
    return {x, v};
template <tvpename T>
Point<T> triangleCircumradius(const Point<T>& a, const Point<T>& b,
                              const Point<T>& c) {
    auto dab = distance(a, b);
    auto dbc = distance(b, c);
    auto dca = distance(c, a);
    return (dab + dbc + dca) / triangleArea(a, b, c);
}
template <class Point>
vector<Point> segInter(Point a, Point b, Point c, Point d) {
    auto oa = c.cross(d, a), ob = c.cross(d, b), oc = a.cross(b, c),
         od = a.cross(b, d);
    // Checks if intersection is single non-endpoint
    if (sgn(oa) * sgn(ob) < 0 \&\& sgn(oc) * sgn(od) < 0)
        return \{(a * ob - b * oa) / (ob - oa)\};
    set<Point> s:
    if (onSegment(c, d, a)) s.insert(a);
    if (onSegment(c, d, b)) s.insert(b);
    if (onSegment(a, b, c)) s.insert(c);
    if (onSegment(a, b, d)) s.insert(d);
    return {all(s)};
```

```
*/
                                                                                           st.pop back();
template <tvpename T>
double angle(const Point<T>& P, const Point<T>& Q, const Point<T>& R.
                                                                                       st.push back(pts[i]);
             const Point<T>& S) {
    auto ux = P.x - Q.x;
                                                                                   pts = st;
    auto uy = P.y - Q.y;
    auto vx = R.x - S.x:
    auto vy = R.y - S.y;
    auto num = ux * vx + uv * vv;
                                                                               template <tvpename T>
    auto den = hypot(ux, uy) * hypot(vx, vy);
   // Caso especial: se den == 0, algum dos vetores \acute{e} degenerado: os dois
    // pontos ãso iguais. Neste caso, o ângulo ãno áest definido
    return acos(num / den);
                                                                               template <tvpename T>
                                                                                   Point<T> b = C - A, c = B - A:
struct pt {
    double x, y;
    int id:
                                                                               template <typename T>
};
int orientation(pt a, pt b, pt c) {
    double v = a.x * (b.y - c.y) + b.x * (c.y - a.y) + c.x * (a.y - b.y);
                                                                                   Point<T> o = ps[0];
    if (v < 0) return -1; // clockwise</pre>
                                                                                   double r = 0, EPS = 1 + 1e-8;
   if (v > 0) return +1: // counter-clockwise
    return 0:
                                                                                       o = ps[i], r = 0:
                                                                                           o = (ps[i] + ps[i]) / 2;
bool cw(pt a, pt b, pt c, bool include collinear) {
                                                                                            r = (o - ps[i]).dist();
    int o = orientation(a, b, c);
    return o < 0 || (include collinear && o == 0);
bool collinear(pt a, pt b, pt c) { return orientation(a, b, c) == 0; }
void convex hull(vector<pt>& pts, bool include collinear = false) {
    pt p0 = *min element(all(pts), [](pt a, pt b) {
        return make pair(a.y, a.x) < make pair(b.y, b.x);</pre>
                                                                                   return {o, r};
    sort(all(pts), [&p0](const pt& a, const pt& b) {
        int o = orientation(p0, a, b);
        if (o == 0)
            return (p0.x - a.x) * (p0.x - a.x) + (p0.v - a.v) * (p0.v - a.x)
                                                                               template <tvpename T>
   y) <
                   (p0.x - b.x) * (p0.x - b.x) + (p0.y - b.y) * (p0.y - b.y)
                                                                                   auto a = 2 * (0.x - P.x);
                                                                                   auto b = 2 * (0.v - P.v);
        return o < 0;
                                                                                   return {a, b, c};
   if (include collinear) {
                                                                               }
        int i = len(pts) - 1;
        while (i \ge 0 \& collinear(p0, pts[i], pts.back())) i--;
        reverse(pts.begin() + i + 1, pts.end());
    vector<pt> st;
                                                                               ll polygonArea(vector<pll>& pts) {
    for (int i = 0; i < len(pts); i++) {
                                                                                   ll^{\tilde{a}ts} = 0;
        while (st.size() > 1 \&\&
```

```
!cw(st[len(st) - 2], st.back(), pts[i], include collinear))
double ccRadius(const Point<T>& A, const Point<T>& B, const Point<T>& C) {
    return (B - A).dist() * (C - B).dist() * (A - C).dist() /
          abs((B - A).cross(C - A)) / 2;
Point<T> ccCenter(const Point<T>& A, const Point<T>& B, const Point<T>& C)
    return A + (b * c.dist2() - c * b.dist2()).perp() / b.cross(c) / 2;
pair<Point<T>, double> mec(vector<Point<T>> ps) {
    shuffle(all(ps), mt19937(time(0)));
    rep(i, 0, len(ps)) if ((o - ps[i]).dist() > r * EPS) {
        rep(i, 0, i) if ((o - ps[i]).dist() > r * EPS) {
            rep(k, 0, j) if ((o - ps[k]).dist() > r * EPS) {
                o = ccCenter(ps[i], ps[i], ps[k]);
                r = (o - ps[i]).dist();
Line<T> perpendicular bisector(const Point<T>& P, const Point<T>& Q) {
    auto c = (P.x * P.x + P.y * P.y) - (Q.x * Q.x + Q.y * Q.y);
ll cross(ll x1, ll y1, ll x2, ll y2) { return x1 * y2 - x2 * y1; }
```

```
for (int i = 2; i < len(pts); i++)
        ats += cross(pts[i].first - pts[0].first, pts[i].second - pts[0].
    second,
                     pts[i - 1].first - pts[0].first,
                     pts[i - 1].second - pts[0].second);
    return abs(ats / 211):
ll boundary(vector<pll>& pts) {
    ll ats = pts.size();
    for (int i = 0; i < len(pts); i++) {</pre>
        il deltax = (pts[i].first - pts[(i + 1) % pts.size()].first);
ll deltay = (pts[i].second - pts[(i + 1) % pts.size()].second);
        ats += abs( gcd(deltax, deltay)) - 1;
    return ats;
pll latticePoints(vector<pll>& pts) {
    ll bounds = boundary(pts);
    ll area = polygonArea(pts);
    ll inside = area + 1ll - bounds / 2ll;
    return {inside, bounds};
            _____
template <typename T>
bool contains(const Point<T>& A, const Point<T>& B, const Point<T>& P) {
    // Verifica se P áest na ãregio retangular
    auto xmin = min(A.x, B.x);
    auto xmax = max(A.x, B.x);
    auto ymin = min(A.y, B.y);
    auto ymax = max(A.y, B.y);
    if (P.x < xmin \mid\mid P.x > xmax \mid\mid P.y < ymin \mid\mid P.y > ymax) return false
    // Verifica çãrelao de csemelhana no âtringulo
    return equals((P.y - A.y) * (B.x - A.x), (P.x - A.x) * (B.y - A.y);
}
// the polygon area of a intersection between a circle and a ccw polygon
template <typename T>
#define arg(p, q) atan2(p.cross(q), p.dot(q))
double circlePoly(Point<T> c, double r, vector<Point<T>> ps) {
    auto tri = [&](Point<T> p, Point<T> q) {
        auto r2 = r * r / 2;
        Point<T> d = q - p;
        auto a = d.dot(p) / d.dist2(), b = (p.dist2() - r * r) / d.dist2()
        auto det = a * a - b;
        if (det \le 0) return arg(p, q) * r2;
        auto s = max(0., -a - sqrt(det)), t = min(1., -a + sqrt(det));
        if (t < 0 | | 1 \le s) return arg(p, q) * r2;
        Point<T> u = p + d * s, v = p + d * t;
        return arg(p, u) * r2 + u.cross(v) / 2 + arg(v, q) * r2;
```

7.2 Angle between three points

Description: Computes the angle apb in radians **Warning**: a is equal to b then the angle isn't defined.

```
#include "./template.cpp"
template <typename T>
ld angle(const Point<T>& p, const Point<T>& a, const Point<T>& b) {
    auto ux = p.x - a.x;
    auto uy = p.y - a.y;
    auto vx = p.x - b.x;
    auto vy = p.y - b.y;
    auto num = ux * vx + uy * vy;
    auto den = hypot(ux, uy) * hypot(vx, vy);
    return acos(num / den);
}
```

7.3 Area of union of rectangles

```
using SegT = ll;
const SeqT eSeq = 1e9;
struct QueryT {
    SeqT q, v;
    QueryT(): q(0), v(eSeq) {}
=== QueryT(SegT _v) : q(1), v(_v) {}
inline QueryT combine(QueryT ln, QueryT rn, pii lr1, pii lr2) {
    QueryT ret;
    if (ln.v < rn.v) ret = ln;</pre>
    if (rn.v < ln.v) ret = rn;</pre>
    if (rn.v == ln.v) {
        ret.v = ln.v;
        ret.q = ln.q + rn.q;
    return ret;
using LazvT = SegT:
inline OueryT applyLazyInOuery(QueryT q, LazyT l, pii lr) {
    if (l == LazyT()) return q;
    if (q.v == eSeq) q.v = 0, q.q = 1;
```

```
q.v += l;
   return q;
inline LazyT applyLazyInLazy(LazyT a, LazyT b) { return a + b; }
using UpdateT = SegT;
inline QueryT applyUpdateInQuery(QueryT q, UpdateT u, pii lr) {
    return applyLazyInQuery(q, u, lr);
inline LazyT applyUpdateInLazy(LazyT l, UpdateT u, pii lr) { return l + u;
template <typename Qt = QueryT, typename Lt = LazyT, typename Ut = UpdateT
          auto C = combine, auto ALQ = applyLazyInQuery,
          auto ALL = applyLazyInLazy, auto AUQ = applyUpdateInQuery,
          auto AUL = applyUpdateInLazy>
struct LazySegmentTree {
   int n, h;
   vector<Qt> ts;
   vector<Lt> ds;
   vector<pii> lrs;
   LazySegmentTree(int n)
        : n(_n),
         h(sizeof(int) * 8 - builtin clz(n)),
          ts(n \ll 1),
          ds(n),
          lrs(n << 1) {
        rep(i, 0, n) lrs[i + n] = {i, i};
        rrep(i, n - 1, 0) {
            lrs[i] = {lrs[i << 1].first, lrs[i << 1 | 1].second};</pre>
   LazySegmentTree(const vector<Qt> &xs) : LazySegmentTree(len(xs)) {
       copy(all(xs), ts.begin() + n);
        rep(i, 0, n) lrs[i + n] = \{i, i\};
        rrep(i, n - 1, 0) {
            ts[i] = C(ts[i << 1], ts[i << 1 | 1], lrs[i << 1], lrs[i << 1
   | 11);
   void set(int p, Qt v) {
       ts[p + n] = v;
       build(p + n);
   void upd(int l, int r, Ut v) {
       l += n. r += n + 1:
       int 10 = 1, r0 = r;
        for (; l < r; l >>= 1, r >>= 1) {
            if (l & 1) apply(l++, v);
            if (r & 1) apply(--r, v);
       build(l0), build(r0 - 1);
   Qt qry(int l, int r) {
```

```
l += n, r += n + 1:
        push(l), push(r - 1);
        Qt resl = Qt(), resr = Qt();
        pii lr1 = \{l, l\}, lr2 = \{r, r\};
        for (; l < r; l >>= 1, r >>= 1) {
            if (l & 1) resl = C(resl, ts[l], lr1, lrs[l]), l++;
            if (r \& 1) r--, resr = C(ts[r], resr, lrs[r], lr2);
        return C(resl, resr, lr1, lr2);
    }
    void build(int p) {
        while (p > 1) {
            p >>= 1;
            ts[p] =
                ALQ(C(ts[p << 1], ts[p << 1 | 1], lrs[p << 1], lrs[p << 1]
    | 1]),
                    ds[p], lrs[p]);
    }
    void push(int p) {
        rrep(s, h, 0) {
            int i = p \gg s;
            if (ds[i] != Lt()) {
                apply(i << 1, ds[i]), apply(i << 1 | 1, ds[i]);
                ds[i] = Lt();
        }
    inline void apply(int p, Ut v) {
        ts[p] = AUQ(ts[p], v, lrs[p]);
        if (p < n) ds[p] = AUL(ds[p], v, lrs[p]);
};
ll areaOfRectanglesUnion(
    const vector<pair<Point<int>, Point<int>>> &rectangles) {
    if (!size(rectangles)) return 0;
    int maxy = INT MIN;
    for (auto &[p1, p2] : rectangles) {
        assert(p1.x < p2.x && p1.y < p2.y);
        maxy = max(\{maxy, p1.y, p2.y\});
    vector<array<int, 4>> sl;
    sl.reserve(size(rectangles) * 2);
    for (auto &[p1, p2] : rectangles) {
        sl.push_back({p1.x, p1.y, p2.y - 1, 1});
        sl.push back(\{p2.x, p1.y, p2.y - 1, -1\});
    sort(sl.begin(), sl.end());
    vector<QueryT> aux(maxy, QueryT(0));
    LazySegmentTree seg(aux);
    // memset(seg vec, 0, sizeof(ll) * maxy);
    // seg::build(maxy, seg_vec);
    int prevx = get<0>(sl.front());
    ll ans = 0;
```

```
for (auto &[curx, ys, yf, inc] : sl) {
    auto [q, v] = seg.qry(0, maxy - 1);
    // auto [q, v] = seg::query(0, maxy - 1);
    ans += (ll)(curx - prevx) * (v ? maxy : maxy - q);
    seg.upd(ys, yf, inc);
    prevx = curx;
}
return ans;
```

7.4 Area: polygon

```
#include "./template.cpp"
template <typename T>
ld area(const vector<Point<T>>& pts) {
    ld a = 0.0;
    int n = size(pts);
    for (int i = 0; i < n; i++) {
        a += pts[i].x * pts[(i + 1) % n].y;
        a -= pts[i].y * pts[(i + 1) % n].x;
    }
    return fabs(a) / (ld)2;
}</pre>
```

7.5 Check if point belongs to line

7.6 Check if point belongs to segment

```
#include "./template.cpp"
template <class P>
bool segmentContainsPoint(const P& p, const P& a, const P& b) {
    auto xmin = min(a.x, b.x);
    auto xmax = max(a.x, b.x);
    auto ymin = min(a.y, b.y);
    auto ymax = max(a.y, b.y);
    if (p.x < xmin or p.x > xmax or p.y < ymin or p.y > ymax) return false;
    return equals((p.y - a.y) * (b.x - a.x), (p.x - a.x) * (b.y - a.y));
}
```

7.7 Check if point is inside polygon

Description: checks if the point p is inside the polygon with vertices in pts, works for both convex and concave polygons.

```
#pragma once
#include "./Angle between three points.cpp"
#include "./Check if point belongs to segment.cpp"
#include "./Determinant.cpp"
#include "./template.cpp"
template <tvpename T>
bool contains(const vector<Point<T>>& pts, const Point<T>& p) {
    int n = size(pts):
    if (n < 3) return false; // may treat it appart</pre>
    T sum = 0.0:
    for (int i = 0; i < n; i++) {
        auto d = determinant(p, pts[i], pts[(i + 1) % n]);
        auto a = angle(p, pts[i], pts[(i + 1) % n]);
        sum += d > 0 ? a : (d < 0 ? -a : 0):
    return equals(fabs(sum), 2 * PI);
// 0: outside, 1: inside, 2: boundary
template <class P>
int pointInPolygon(const vector<P>& pts, const P& p) {
    if (contains(pts, p)) return 1;
    int n = size(pts):
    for (int i = 0; i < n; i++) {
        if (segmentContainsPoint(p, pts[i], pts[(i + 1) % n])) {
            return 2:
    return 0;
```

7.8 Convex hull

```
#include "../Contest/template.cpp"
#include "./Determinant.cpp"
#include "./template.cpp"

template <typename T>
vector<Point<T>> convexHull(vector<Point<T>> pts) {
    if (len(pts) <= 1) return pts;
        sort(all(pts));
        vector<Point<T>> h(len(pts) + 1);
    int s = 0, t = 0;
    for (int it = 2; it--; s = --t, reverse(all(pts)))
        for (Point<T> p: pts) {
             while (t >= s + 2 && determinant(h[t - 2], h[t - 1], p) <= 0)
        t--;
             h[t++] = p;
        }
}</pre>
```

```
return \{h.begin(), h.begin() + t - (t == 2 \& h[0] == h[1])\};
template <typename T>
vector<Point<T>> convexHull2(vector<Point<T>> pts) {
    int n = len(pts);
    sort(pts.begin(), pts.end());
   vector<Point<T>> l, u;
   for (int i = 0; i < n; i++) {
        while (len(l) >= 2 \&\&
               determinant(l[len(l) - 1], l[len(l) - 2], pts[i]) < 0) {
            l.pop back();
        l.push back(pts[i]);
   for (int i = n - 1; \simi; --i) {
        while (len(u) \geq 2 &&
               determinant(u[len(u) - 1], u[len(u) - 2], pts[i]) < 0)
            u.pop back();
        u.push back(pts[i]);
   u.pop back(), l.pop back();
   u.reserve(len(u) + len(l));
   u.insert(u.end(), all(l));
    return u:
```

7.9 Cross product between points

```
#pragma once
#include "./template.cpp"
template <typename T>
T cross(const Point<T>& p, const Point<T>& q) {
    return p.x * q.y - p.y * q.x;
}
```

7.10 Define line from two points

```
#pragma once
#include "./template.cpp"
template <typename T>
inline tuple<T, T, T> defineLine(const Point<T>& p, const Point<T>& q) {
    return {p.y - q.y, q.x - p.x, cross(p, q)};
}
```

7.11 Determinant

```
#pragma once
#include "./template.cpp"
template <typename T>
```

7.12 Distance: point to point

```
#include "./template.cpp"
template <typename T>
T distance(const Point<T>& p, const Point<T>& q) {
   return hypot(p.x - q.x, p.y - q.y);
}
```

7.13 Halfplane intersection

```
#pragma once
#include "./Point.cpp"
#include "./template.cpp"
// Basic half-plane struct.
struct Halfplane {
    // 'p' is a passing point of the line and 'pg' is the direction vector
    οf
    // the line.
    Point<ld> p, pq;
    long double angle;
    Halfplane() {}
    Halfplane(const Point<ld>& a, const Point<ld>& b) : p(a), pq(b - a) {
        angle = atan2l(pq.y, pq.x);
    // Check if point 'r' is outside this half-plane.
    // Every half-plane allows the region to the LEFT of its line.
    bool out(const Point<ld>& r) { return cross(pg, r - p) < -EPS; }
    // Intersection point of the lines of two half-planes. It is assumed
   they're
    // never parallel.
    friend Point<ld> inter(const Halfplane& s, const Halfplane& t) {
        long double alpha = cross((t.p - s.p), t.pq) / cross(s.pq, t.pq);
        return s.p + (s.pg * alpha);
};
// Actual algorithm
// receive it by reference if don't care messing with it
vector<Point<ld>> hp intersect(vector<Halfplane> H) {
    const ld inf = 2e6:
    Point<ld> box[4] = \{// Bounding box in CCW order
                        Point<ld>(inf, inf), Point<ld>(-inf, inf),
                        Point<ld>(-inf, -inf), Point<ld>(inf, -inf));
    for (int i = 0; i < 4; i++) { // Add bounding box half-planes.
        Halfplane aux(box[i], box[(i + 1) % 4]);
```

```
H.push back(aux);
}
// Sort by angle and start algorithm
sort(H.begin(), H.end(), [&](const Halfplane& a, const Halfplane& b) {
    return a.angle < b.angle:
deque<Halfplane> dg;
int len = 0:
for (int i = 0; i < int(H.size()); i++) {</pre>
    // Remove from the back of the deque while last half-plane is
redundant
    while (len > 1 && H[i].out(inter(dg[len - 1], dg[len - 2]))) {
        dq.pop back();
        --len:
    // Remove from the front of the deque while first half-plane is
    // redundant
    while (len > 1 \&\& H[i].out(inter(dg[0], dg[1])))  {
        dq.pop front();
        --len:
    }
    // Special case check: Parallel half-planes
    if (len > 0 \& fabsl(cross(H[i].pq, dq[len - 1].pq)) < EPS) {
        // Opposite parallel half-planes that ended up checked against
each
        if (dot(H[i].pq, dq[len - 1].pq) < 0.0) return vector<Point<ld
>>();
        // Same direction half-plane: keep only the leftmost half-
plane.
        if (H[i].out(dq[len - 1].p)) {
            dq.pop_back();
            --len:
        } else
            continue:
    // Add new half-plane
    dq.push back(H[i]);
    ++len:
// Final cleanup: Check half-planes at the front against the back and
while (len > 2 \& dq[0].out(inter(dq[len - 1], dq[len - 2]))) {
    dq.pop back();
    --len:
while (len > 2 && dq[len - 1].out(inter(dq[0], dq[1]))) {
    dq.pop front();
    --len:
// Report empty intersection if necessary
if (len < 3) return vector<Point<ld>>();
// Reconstruct the convex polygon from the remaining half-planes.
```

```
vector<Point<ld>> ret(len);
for (int i = 0; i + 1 < len; i++) {
    ret[i] = inter(dq[i], dq[i + 1]);
}
ret.back() = inter(dq[len - 1], dq[0]);
return ret;
}</pre>
```

7.14 Lattice points

```
#pragma once
#include "../Contest/template.cpp"
#include "./Area: polygon.cpp"
#include "./template.cpp"
template <tvpename T>
pair<ll, ll> latticePoints(const vector<Point<T>> &pts) {
    ll bounds = pts.size();
    int n = pts.size();
    for (int i = 0; i < n; i++) {
        ll deltax = (pts[i].x - pts[(i + 1) % n].x);
        ll deltay = (pts[i].y - pts[(i + 1) % n].y);
        bounds += abs( gcd(deltax, deltay)) - 1;
    ll a = area(pts);
    ll inside = a + 1 - bounds / 2ll;
    return {inside, bounds};
}
```

7.15 Left of polygon cut

Warning: if some vertex lies exactly on the line A B, theese vertex will be included in teh answer

```
#include "./Determinant.cpp"
#include "./template.cpp"
template <tvpename T>
vector<Point<T>> leftOfPolygonCut(const vector<Point<T>>& vs, const Point<</pre>
   T>& A,
                                   const Point<T>& B) {
    // çãInterseo entre a reta AB e o segmento de reta PQ
    auto intersection = [&](const Point<T>& P, const Point<T>& Q,
                             const Point<T>& A, const Point<T>& B) -> Point
    <T> {
        auto a = B.y - A.y;
        auto b = A.x - B.x;
        auto c = B.x * A.y - A.x * B.y;
        auto u = fabs(a * P.x + b * P.y + c);
        auto v = fabs(a * 0.x + b * 0.v + c):
        // éMdia ponderada pelas âdistncias de P e O éat a reta AB
        return \{(P.x * v + Q.x * u) / (u + v), (P.y * v + Q.y * u) / (u + v)\}
   v)};
    };
    vector<Point<T>> points;
```

```
int n = size(vs);
for (int i = 0; i < n; ++i) {
    auto d1 = determinant(A, B, vs[i]);
    auto d2 = determinant(A, B, vs[(i + 1) % n]);

    // éVrtice à esquerda da reta
    if (d1 > -EPS) points.push_back(vs[i]);

    // A aresta cruza a reta
    if (d1 * d2 < -EPS)
        points.push_back(intersection(vs[i], vs[(i + 1) % n], A, B));
}
return points;</pre>
```

7.16 Perimeter: polygon

```
#include "./Distance: point to point.cpp"
#include "./template.cpp"

template <typename T>
T perimeter(const vector<Point<T>>& pts) {
    T p = 0.0;
    int n = size(pts);
    for (int i = 0; i < n; i++) {
        p += distance(pts[i], pts[(i + 1) % n]);
    }
    return p;
}</pre>
```

7.17 Point

```
// Basic point/vector struct.
template <typename T>
struct Point {
   T x, y;
   Point(T x = 0, T y = 0) : x(x), y(y) {}
   // Addition, substraction, multiply by constant, dot product, cross
   product.
   friend Point<T> operator+(const Point<T>& p, const Point<T>& q) {
        return PointT>(p.x + q.x, p.y + q.y);
   }
   friend Point<T> operator-(const Point<T>& p, const Point<T>& q) {
        return Point<T>(p.x - q.x, p.y - q.y);
   template <typename T2>
   friend Point<T> operator*(const Point<T>& p, T2 k) {
        return Point<T>(p.x * k, p.y * k);
   friend T dot(const Point<T>& p, const Point<T>& q) {
        return p.x * q.x + p.y * q.y;
   friend T cross(const Point<T>& p, const Point<T>& q) {
```

```
return p.x * q.y - p.y * q.x; };
```

7.18 Polygon (regular): apothem

```
#include "./Distance: point to point.cpp"
#include "./template.cpp"
template <typename T>
ld apothem(const vector<Point<T>>& pts) {
    auto s = distance(pts[0], pts[1]);
    int n = size(pts);
    return (s / 2.0) * (1.0 / tan(PI / n));
}
```

7.19 Polygon (regular): circumradius

```
#include "./Distance: point to point.cpp"
#include "./template.cpp"

template <typename T>
ld circumradius(const vector<Point<T>>& pts) {
    auto s = distance(pts[0], pts[1]);
    int n = size(pts);
    return (s / 2.0) * (1.0 / sin(PI / (ld)n));
}
```

7.20 Polygon: check if is convex

```
#include "./Determinant.cpp"
#include "./template.cpp"

template <typename T>
bool checkIfPolygonIsConvex(vector<Point<T>>& pts) {
    int n = size(pts);
    if (n < 3) return false;
    int l, g, e;
    l = g = e = 0;
    for (int i = 0; i < n; i++) {
        auto d = determinant(pts[i], pts[(i + 1) % n], pts[(i + 2) % n]);
        d ? (d > 0 ? g++ : l++) : e++;
    }
    return l == n or g == n;
}
```

7.21 Rectangle intersection

```
Assumes that the points P, Q that define
  a rectangle are the bottom-left and top-right
   corner, and also that the sides are parallel to the axis.
#pragma once
#include "../Contest/template.cpp"
#include "./Point.cpp"
template <typename T>
optional<pair<Point<T>, Point<T>>> rectangleIntersection(
    const pair<Point<T>, Point<T>> &r1, const pair<Point<T>, Point<T>> &r2
    assert(r1.first.x < r1.second.x && r1.first.y < r1.second.y);
   assert(r2.first.x < r2.second.x && r2.first.y < r2.second.y);</pre>
   T x1 = max(r1.first.x, r2.first.x);
   T \times 2 = min(r1.second.x, r2.second.x);
   T y1 = max(r1.first.y, r2.first.y);
   T y2 = min(r1.second.y, r2.second.y);
   if (x1 \ge x2 \text{ or } y1 \ge y2) \text{ return nullopt};
    return pair<Point<T>, Point<T>>{{x1, y1}, {x2, y2}};
```

7.22 template

```
#pragma once
#include <bits/stdc++.h>
using namespace std;
using ld = long double;
template <tvpename T>
using Point = pair<T, T>:
#define x first
#define y second
const double EPS{1e-6};
const ld PI = acos(-1);
template <tvpename T>
bool equals(T a, T b) {
    if (std::is_floating_point<T>::value)
        return fabs(a - b) < EPS;</pre>
    else
        return a == b:
template <typename T>
bool equals(Point<T> a, Point<T> b) {
    if (std::is floating point<T>::value)
        return fabs(a,x - b,x) < EPS && fabs(a,v - b,v) < EPS:
   else
        return a == b;
}
```

8 Graphs

8.1 Heavy-Light Decomposition (point update)

8.1.1 Maximum number on path

```
struct Node {
    ll value:
    Node()
        : value(numeric limits<ll>::min()) {}; // Neutral
                                                // element
    Node(ll v) : value(v) {};
Node combine(Node l, Node r) {
    Node m;
    m.value = max(l.value, r.value);
    return m:
template <typename T = Node, auto F = combine>
struct SegTree {
    int n:
    vector<T> st;
    SegTree(int n) : n(n), st(n \ll 1) {}
    void set(int p, const T &k) {
        for (st[p += n] = k; p >>= 1;) st[p] = F(st[p << 1], st[p << 1])
   1]);
   T query(int l, int r) {
        T ansl. ansr:
        for (l += n, r += n + 1; l < r; l >>= 1, r >>= 1) {
            if (l & 1) ansl = F(ansl, st[l++]);
            if (r \& 1) ansr = F(st[--r], ansr);
        return F(ansl, ansr);
};
template <typename SegT = Node, auto SegOp = combine>
struct HeavyLightDecomposition {
    int n:
    vi ps, ds, sz, heavy, head, pos;
    SegTree<SegT, SegOp> seg;
    HeavyLightDecomposition(const vi2d &q, const vector<SegT> &v, int root
    = 0)
        : n(len(q)), seq(n) {
        ps = ds = sz = heavy = head = pos = vi(n, -1);
        auto dfs = [&](auto &&self, int u) -> void {
            sz[u] = 1;
            int mx = 0:
            for (auto x : q[u])
                if (x != ps[u]) {
                    ps[x] = u:
                    ds[x] = ds[u] + 1;
                    self(self, x);
                    sz[u] += sz[x];
```

```
if (sz[x] > mx) mx = sz[x], heavy[u] = x;
    };
    dfs(dfs, root);
    for (int i = 0, cur = 0; i < n; i++) {
        if (ps[i] == -1 \text{ or heavy}[ps[i]] != i)
            for (int j = i; j != -1; j = heavy[j]) {
                head[i] = i;
                pos[j] = cur++;
    }
    rep(i, 0, n) seq.set(pos[i], v[i]);
vector<pii> disjoint ranges(int u, int v) {
    vector<pii> ret:
    for (; head[u] != head[v]; v = ps[head[v]]) {
        if (ds[head[u]] > ds[head[v]]) swap(u, v);
        ret.eb(pos[head[v]], pos[v]);
    if (ds[u] > ds[v]) swap(u, v);
    ret.eb(pos[u], pos[v]);
    return ret;
SegT query path(int u, int v) {
    SegT res;
    for (auto [l, r] : disjoint ranges(u, v)) {
        res = SegOp(res, seg.query(l, r));
    return res;
SegT query subtree(int u) const {
    return seq.query(pos[u], pos[u] + sz[u] - 1);
void set(int u, SeqT x) { seq.set(pos[u], x); }
```

8.2 2-SAT

};

Description: Calculates a valid assignment to boolean variables a, b, c,... to a 2-SAT problem, so that an expression of the type (a||b)&&(!a||c)&&(d||!b)&&... becomes true, or reports that it is unsatisfiable.

Usage: Negated variables are represented by bit-inversions (\tilde{x}) .

Returns true iff it is solvable ts.values [0..N-1] holds the assigned values to the vars.

Time: O(N + E), where N is the number of boolean variables, and E is the number of clauses.

```
/
struct TwoSat {
    int N;
    vector<vi> gr;
    vi values; // 0 = false, 1 = true
    TwoSat(int n = 0) : N(n), gr(2 * n) {}
```

```
int addVar() { // (optional)
        ar.eb():
        gr.eb();
        return N++;
    void either(int f, int j) {
        f = max(2 * f, -1 - 2 * f);
        i = \max(2 * i, -1 - 2 * i);
        gr[f].pb(j ^ 1);
        gr[j].pb(f ^ 1);
    void setValue(int x) { either(x, x); }
    void implies(int f, int j) { either(\simf, j); } // (optional)
    void atMostOne(const vi &li) { // (optional)
        if (len(li) <= 1) return;</pre>
        int cur = \sim li[0];
        rep(i, 2, len(li)) {
            int next = addVar();
            either(cur, ~li[i]);
            either(cur, next);
            either(\simli[i], next);
            cur = \simnext;
        either(cur, \simli[1]);
    vi val, comp, z;
    int time = 0;
    int dfs(int i) {
        int low = val[i] = ++time, x;
        z.pb(i);
        for (int e : gr[i])
            if (!comp[e]) low = min(low, val[e] ?: dfs(e));
        if (low == val[i]) do {
                 x = z.back();
                 z.ppb();
                 comp[x] = low;
                 if (values[x >> 1] == -1) values[x >> 1] = x & 1;
            } while (x != i);
        return val[i] = low:
    bool solve() {
        values.assign(N, -1);
        val.assign(2 * N, 0);
        comp = val;
        rep(i, 0, 2 * N) if (!comp[i]) dfs(i);
        rep(i, 0, N) if (comp[2 * i] == comp[2 * i + 1]) return 0;
        return 1:
};
```

8.3 BFS-01

Description: Similar to a Dijkstra given a weighted graph finds the distance from source s to every other node.

```
Warning: Applicable only when the weight of the edges \in \{0, x\}
vector<pair<ll, int>> adi[maxn];
ll dists[maxn];
int s, n;
void bfs 01() {
    fill(dists, dists + n, oo);
    dist[s] = 0;
    deque<int> q;
    g.emplace back(s);
    while (not q.empty()) {
        auto u = q.front();
        q.pop_front();
        for (auto [v, w] : adi[u]) {
            if (dist[v] <= dist[u] + w) continue;</pre>
            dist[v] = dist[u] + w;
            w ? q.emplace back(v) : q.emplace front(v);
```

8.4 Bellman ford

Description: Find shortest path from a single source to all other nodes. Can detect negative cycles.

Time: $O(V \cdot E)$

Time: O(V+E)

```
bool bellman ford(const vector<vector<pair<int, ll>>> &g, int s,
                  vector<ll> &dist) {
   int n = (int)q.size();
   dist.assign(n, LLONG MAX);
   vector<int> count(n);
   vector<char> in queue(n);
   queue<int> q;
   dist[s] = 0;
   q.push(s);
   in queue[s] = true;
   while (not q.empty()) {
        int cur = q.front();
        q.pop();
        in queue[cur] = false;
        for (auto [to, w] : g[cur]) {
            if (dist[cur] + w < dist[to]) {</pre>
                dist[to] = dist[cur] + w;
                if (not in queue[to]) {
                    q.push(to);
                    in queue[to] = true;
                    count[to]++:
                    if (count[to] > n) return false;
            }
       }
```

```
}
return true;
```

8.5 Bellman-Ford (find negative cycle)

Description: Given a directed graph find a negative cycle by running n iterations, and if the last one produces a relaxation than there is a cycle.

Time: $O(V \cdot E)$

```
const ll oo = 2500 * 1e9;
using graph = vector<vector<pair<int, ll>>>;
vi negative cycle(graph &g, int n) {
    vll d(n, oo);
    vi p(n, -1);
    int x = -1:
    d[0] = 0;
    for (int i = 0; i < n; i++) {
        x = -1;
        for (int u = 0; u < n; u++) {
            for (auto &[v, l] : q[u]) {
                if (d[u] + l < d[v]) {
                    d[v] = d[u] + 1;
                    p[v] = u;
                     \dot{x} = v;
    }
    if (x == -1)
        return {};
    else {
        for (int i = 0; i < n; i++) x = p[x];
        vi cycle;
        for (int v = x; v = p[v]) {
            cvcle.eb(v);
            if (v == x and len(cycle) > 1) break;
        reverse(all(cycle));
        return cycle;
}
```

8.6 Biconnected Components

Description: Build a vector of vectors, where the i-th vector correspond to the nodes of the i-th, biconnected component, a biconnected component is a subset of nodes and edges in which there is no cut point, also exist at least two distinct routes in vertex between any two vertex in the same biconnected component.

Time: O(N+M)

```
const int maxn(5 '00' 000);
int tin[maxn], stck[maxn], bcc_cnt, n, top = 0, timer = 1;
vector<int> q[maxn], nodes[maxn];
```

```
int tarjan(int u, int p = -1) {
    int lowu = tin[u] = timer++;
    int son cnt = 0;
    stck[++top] = u;
    for (auto v : g[u]) {
        if (!tin[v]) {
            son cnt++;
            int lowx = tarjan(v, u);
            lowu = min(lowu, lowx);
            if (lowx >= tin[u]) {
                while (top != -1 \&\& stck[top + 1] != v)
                    nodes[bcc cnt].emplace back(stck[top--]);
                nodes[bcc cnt++] emplace back(u);
        } else {
            lowu = min(lowu, tin[v]);
    if (p == -1 \&\& son cnt == 0) {
        nodes[bcc cnt++].emplace back(u):
    return lowu:
void build bccs() {
    timer = 1:
    top = -1;
    memset(tin, 0, sizeof(int) * n);
    for (int i = 0; i < n; i++) nodes[i] = {};
    bcc cnt = 0;
    for (int u = 0; u < n; u++)
        if (!tin[u]) tarjan(u);
}
```

8.7 Binary Lifting/Jumping

Description: Given a function/successor grpah answers queries of the form which is the node after k moves starting from u.

Time: Build $O(N \cdot MAXLOG2)$, Query O(MAXLOG2).

```
const int MAXN(2e5), MAXLOG2(30);
int bl[MAXN][MAXLOG2 + 1];
int N;
int jump(int u, ll k) {
    for (int i = 0; i <= MAXLOG2; i++) {
        if (k & (1ll << i)) u = bl[u][i];
    }
    return u;
}
void build() {
    for (int i = 1; i <= MAXLOG2; i++) {
        for (int j = 0; j < N; j++) {
            bl[j][i] = bl[bl[j][i - 1]][i - 1];
        }
</pre>
```

```
}
```

8.8 Bipartite Graph

Description: Given a graph, find the 'left' and 'right' side if is a bipartite graph, if is not then a empty vi2d is returned

Time: O(N+M)

```
vi2d bipartite graph(vi2d &adj) {
    int n = len(adj);
    vi side(n, -1);
    vi2d ret(2);
    rep(u, 0, n) {
        if (side[u] == -1) {
            queue<int> q;
            a.emp(u):
            side[u] = 0;
            ret[0].eb(u);
            while (len(q)) {
                int u = q.front();
                q.pop();
                for (auto v : adj[u]) {
                     if (side[v] == -1) {
                         side[v] = side[u] ^ 1;
                         ret[side[v]].eb(v);
                         q.push(v);
                    } else if (side[u] == side[v])
                         return {};
        }
    return ret;
```

8.9 Block-Cut Tree * *

Description: Builds the Block-Cut of a undirected graph. **

Usage: isGraphCutpoint[u] answers how many connected components * are created when the node u is removed from the graph, if * isGraphCutpoint[u] is greater than 1, it means that u is a * cutpoint. * *

Time: $O(N + \hat{M})$ * * Memory: O(N) * *

Warning: Always careful with disconnected graphs ! you may end up having * multiple trees. * *

```
#pragma once
#include "../Contest/template.cpp"
struct BlockCutTree {
   int n;
   vi idOnTree, tin, low, stk, isGraphCutpoint, isTreeCutpoint;
   vi2d comps, treeAdj;
```

```
BlockCutTree(vi2d &g)
    : n(len(q)), idOnTree(n), tin(n), low(n), isGraphCutpoint(n) {
    rep(i, 0, n) {
        if (!tin[i]) {
            int timer = 0;
            dfs(i, -1, timer, g);
    }
    buildTree();
void buildTree() {
    int node id = 0;
    rep(u, 0, n) {
        if (isGraphCutpoint[u]) {
            idOnTree[u] = node id++;
            isTreeCutpoint.eb(true);
            treeAdj.pb({});
        }
    }
    for (auto &comp : comps) {
        int node = node id++;
        treeAdj.pb({});
        isTreeCutpoint.eb(false);
        for (int u : comp) {
            if (!isGraphCutpoint[u]) {
                idOnTree[u] = node:
            } else {
                treeAdi[node].eb(idOnTree[u]),
                    treeAdj[idOnTree[u]].eb(node);
        }
void dfs(int u, int p, int &timer, vi2d &g) {
    tin[u] = low[u] = ++timer:
    stk.eb(u);
    for (auto v : q[u]) {
        if (v == p) continue;
        if (!tin[v]) {
            dfs(v, u, timer, g);
            chmin(low[u], low[v]);
            if (low[v] >= tin[u]) {
                isGraphCutpoint[u] += (tin[u] > 1 or tin[v] > 2);
                comps.pb({u}):
                while (comps.back().back() != v) {
                    comps.back().eb(stk.back());
                    stk.ppb();
        } else
            low[u] = min(low[u], tin[v]);
    }
int countMandatoryNodesOnPath(int startNode, int endNode);
```

};

8.10 Centroid Decomposition

Description: Builds a vector fat where fat_i is who is the father of the node i in the centroid decomposed tree.

```
#pragma once
#include "../Contest/template.cpp"
vi centroidDecomposition(const vi2d &a) {
    int n = len(q);
    vi fat(n, -1), szt(n), tk(n);
    function<int(int, int)> calcsz = [&](int x, int f) {
        szt[x] = 1;
        for (auto y : g[x])
            if (y != f \&\& !tk[y]) szt[x] += calcsz(y, x);
        return szt[x]:
    function<void(int, int, int)> cdfs = [&](int x, int f, int sz) {
        if (sz < 0) sz = calcsz(x, -1);
        for (auto y : g[x])
            if (!tk[y] \&\& szt[y] * 2 >= sz) {
                szt[x] = 0;
                cdfs(y, f, sz);
                return:
        tk[x] = true;
        fat[x] = f;
        for (auto y : g[x])
            if (!tk[y]) cdfs(y, x, -1);
    cdfs(0, -1, -1);
    return fat;
}
```

8.11 Count mandatory nodes on a single path * *

Description: Given a startNode and an endNode, count the mandatory nodes * in the path from startNode to endNode, that is the number of nodes such * that are present in every possible such path. * *

Time: O(N+M) * * Memory: O(N) * *

Warning: The startNode and endNode is always included in the counting, * ajust your final answer depending on the problem. Be careful with a * disconnected graph where the path may not exist, treat it appart !. * *

```
#pragma once
#include "../Contest/template.cpp"
#include "./Block-Cut tree.cpp"

int BlockCutTree::countMandatoryNodesOnPath(int startNode, int endNode) {
    startNode = idOnTree[startNode], endNode = idOnTree[endNode];
    int ans = !isTreeCutpoint[startNode] + !isTreeCutpoint[endNode];
    int artPoints = 0;
```

```
function<void(int, int)> dfsCount = [&](int u, int p) {
    artPoints += isTreeCutpoint[u];
    if (u == endNode) ans += artPoints;
    for (auto v : treeAdj[u]) {
        if (v != p) {
            dfsCount(v, u);
        }
    }
    artPoints -= isTreeCutpoint[u];
};
dfsCount(startNode, -1);
return ans;
}
```

8.12 DSU query

```
struct DSU {
    V<ii>p;
    V<int> s;
    int sum = 0;
    DSU(int n) : p(n, \{-1, -1\}), s(n, 1) \{\}
    int find(int x) {
        if (p[x].ff < 0) return x;
        return find(p[x].ff);
    void join(int x, int y, int w) {
        x = find(x);
        v = find(v);
        if (x == y) return;
        sum += w;
        if (s[x] < s[y]) swap(x, y);
        s[x] += s[y];
        p[y] = mp(x, w);
    int query(int x, int y) {
        int r = 0;
        while (x != y) {
            if (s[x] < s[y])
                r = max(r, p[x].ss), x = p[x].ff;
            else
                r = max(r, p[y].ss), y = p[y].ff;
        return r:
};
```

8.13 D'Escopo-Pape

Description: Is a single source shortest path that works faster than Dijkstra's algorithm and the Bellman-Ford algorithm in most cases, and will also work for negative edges. However not for negative cycles. There exists cases where it runs in exponential time.

Usage: Returns a pair containing two vectors, the first one with the distance from s to every other node, and another one with the ancestor of each node, note that the ancestor of s is -1

```
using Edge = pair<ll, int>;
using Adj = vector<vector<Edge>>;
pair<vll, vi> desopo pape(int s, int n, const Adj &adj) {
    vll ds(n, LLONG \overline{M}AX), ps(n, -1);
    ds[s] = 0;
    vi ms(n, 2);
    deque<int> q;
    q.eb(s);
    while (len(q)) {
        int u = q.front();
        q.pop_front();
        ms[u] = 0;
        for (auto [w, v] : adj[u]) {
            if (chmin(ds[v], w + ds[u])) {
                 ps[v] = u;
                 if (ms[v] == 2)
                     ms[v] = 1, q.pb(v);
                 else if (ms[v] == 0)
                     ms[v] = 1, q.pf(v);
        }
    return {ds, ps};
```

8.14 Dijkstra

```
const int MAXN = 1'00'000:
const ll MAXW = 1'000'0001;
constexpr 11\ 00 = MAXW * MAXN + 1:
using Edge = pair<ll, int>; // { weigth, node}
using Adi = vector<vector<Edge>>;
template <tvpename T>
using min heap = priority queue<T, vector<T>, greater<T>>;
pair<vll, vi> dijkstra(const Adj &g, int s) {
    int n = len(q);
    min heap<Edge> pg;
    vll ds(n, 00);
    vi ps(n, -1);
    pq.emp(0, s);
    ds[s] = 0;
    while (len(pq)) {
        auto [du, u] = pq.top();
        pq.pop();
        if (ds[u] < du) continue;
        for (auto [w, v] : g[u]) {
            ll ndv = du + w;
            if (chmin(ds[v], ndv)) {
                ps[v] = u;
                pq.emp(ndv, v);
```

```
}
}
return {ds, ps};

// optional !
vi recover_path(int source, int ending, const vi &ps) {
    if (ps[ending] == -1) return {};
    int cur = ending;
    vi ans;
    while (cur != -1) {
        ans.eb(cur);
        cur = ps[cur];
    }
    reverse(all(ans));
    return ans;
}
```

8.15 Dijkstra (K-shortest pahts)

```
const ll oo = 1e9 * 1e5 + 1;
using adj = vector<vector<pll>>>;
vector<priority queue<ll>> dijkstra(const vector<vector<pll>>> &g, int n,
   int s,
                                     int k) {
   priority queue<pll, vector<pll>, greater<pll>>> pq;
   vector<priority queue<ll>> dist(n);
   dist[0].emplace(0);
   pq.emplace(0, s);
   while (!pq.empty()) {
        auto [d1, v] = pq.top();
        pq.pop();
        if (not dist[v].empty() and dist[v].top() < d1) continue;</pre>
        for (auto [d2, u] : g[v]) {
            if (len(dist[u]) < k) {</pre>
                pq.emplace(d2 + d1, u);
                dist[u].emplace(d2 + d1);
            } else ·
                if (dist[u].top() > d1 + d2) {
                    dist[u].pop();
                    dist[u].emplace(d1 + d2);
                    pq.emplace(d2 + d1, u);
            }
        }
    return dist;
```

8.16 Extra Edges to Make Digraph Fully Strongly Connected

Description: Given a directed graph G find the necessary edges to add to make the graph a single strongly connected component.

```
Time: O(N+M)
Memory: O(N)
struct SCC {
    int n, num sccs;
    vi2d adj;
    vi scc id;
    SCC(int _n) : n(_n), num_sccs(0), adj(n), scc_id(n, -1) {}
    SCC(const vi2d & adj) : SCC(len( adj)) {
        adj = adj;
        find_sccs();
    void add edge(int u, int v) { adj[u].eb(v); }
    void find sccs() {
        int timer = 1;
        vi tin(n), st;
        st.reserve(n);
        function<int(int)> dfs = [&](int u) -> int {
            int low = tin[u] = timer++, siz = len(st);
            st.eb(u);
            for (int v : adj[u])
                 if (scc id[v] < 0) low = min(low, tin[v] ? tin[v] : dfs(v)
    );
            if (tin[u] == low) {
                 rep(i, siz, len(st)) scc_id[st[i]] = num_sccs;
                 st.resize(siz):
                 num sccs++;
            return low:
        };
        for (int i = 0; i < n; i++)
            if (!tin[i]) dfs(i);
};
vector<array<int, 2>> extra_edges(const vi2d &adj) {
    SCC scc(adi):
    auto scc_id = scc.scc_id;
    auto num sccs = scc.num sccs;
    if (num sccs == 1) return {};
    int n = len(adj);
    vi2d scc adj(num sccs);
    vi zero \overline{in}(num \ s\overline{ccs}, 1);
    rep(u, \overline{0}, n) {
        for (int v : adj[u]) {
            if (scc id[u] == scc id[v]) continue;
             scc adj[scc id[u]].eb(scc id[v]);
            zero in[scc id[v]] = 0;
    int random source = max element(all(zero in)) - zero in.begin();
    vi vis(num sccs);
    function\langle int(int) \rangle dfs = [&](int u) {
```

```
if (empty(scc adj[u])) return u;
    for (int v : scc adj[u])
        if (!vis[v]) -{
            vis[v] = 1;
            int zero_out = dfs(v);
            if (zero out != -1) return zero_out;
    return (int)-1;
};
vector<array<int, 2>> edges;
vi in unused;
rep(i, 0, num_sccs) {
    if (zero \overline{in}[i]) {
        vis[i] = 1;
        int zero out = dfs(i);
        if (zero out != -1)
            edges.push back({zero out, i});
            in unused.push back(i);
    }
}
rep(i, 1, len(edges)) { swap(edges[i][0], edges[i - 1][0]); }
rep(i, 0, num_sccs) {
    if (scc adj[i].empty() && !vis[i]) {
        if (!in unused.empty()) {
            edges.push back({i, in unused.back()});
            in unused.pop_back();
        } else {
            edges.push back({i, random source});
    }
for (int u : in unused) edges.push back({0, u});
vi to node(num sccs);
rep(i, 0, n) to node[scc id[i]] = i;
for (auto &[u, \overline{v}] : edges) u = to_node[u], v = to_node[v];
return edges;
```

8.17 Find Articulation/Cut Points

Description: Given an **undirected** graph find it's articulation points.

Time: O(N+M)

}

Warning: A vertex u can be an articulation point if and only if has at least 2 adjascent vertex

```
const int MAXN(100);
int N;
vi2d G;
int timer;
int tin[MAXN], low[MAXN];
set<int> cpoints;
int dfs(int u, int p = -1) {
```

```
int cnt = 0:
    low[u] = tin[u] = timer++;
    for (auto v : G[u])
        if (not tin[v]) {
            cnt++;
            dfs(v, u);
            if (low[v] >= tin[u]) cpoints.insert(u);
            low[u] = min(low[u], low[v]);
        } else if (v != p)
            low[u] = min(low[u], tin[v]);
    return cnt;
void getCutPoints() {
    memset(low, 0, sizeof(low));
    memset(tin, 0, sizeof(tin));
    cpoints.clear();
    timer = 1:
    for (int i = 0; i < N; i++) {
        if (tin[i]) continue;
        int cnt = dfs(i);
        if (cnt == 1) cpoints.erase(i);
}
```

8.18 Find Bridge-Tree components

Usage: label2CC(u, p) finds the 2-edge connected component of every node.

Time: O(n+m)

```
const int maxn(3 '00' 000);
int tin[maxn], compId[maxn], qtdComps;
vi q[maxn], stck;
int n;
int dfs(int u, int p = -1) {
    int low = tin[u] = len(stck);
    stck.emplace back(u);
    bool multEdge = false;
    for (auto v : q[u]) {
        if (v == p and !multEdge) {
            multEdge = 1;
            continue;
        low = min(low, tin[v] == -1 ? dfs(v, u) : tin[v]);
    if (low == tin[u]) {
        for (int i = tin[u]; i < len(stck); i++) compId[stck[i]] =</pre>
    atdComps:
        stck.resize(tin[u]);
        qtdComps++;
    }
    return low:
```

```
void label2CC() {
    memset(compId, -1, sizeof(int) * n);
    memset(tin, -1, sizeof(int) * n);
    stck.reserve(n);
    for (int i = 0; i < n; i++) {
        if (tin[i] == -1) dfs(i);
    }
}</pre>
```

8.19 Find Bridges

Description: Find every bridge in a **undirected** connected graph.

Warning: Remember to read the graph as pair where the second is the id of the edge! @Time: O(N + M) const int MAXN(10000), MAXM(100000):

```
int N, M, clk, tin[MAXN], low[MAXN], isBridge[MAXM];
vector<pii> G[MAXN];
void dfs(int u, int p = -1) {
    tin[u] = low[u] = clk++;
    for (auto [v, i] : G[u]) {
        if (v == p) continue;
        if (tin[v]) {
            low[u] = min(low[u], tin[v]);
        } else {
            dfs(v, u);
            low[u] = min(low[u], low[v]);
            if (low[v] > tin[u]) {
                isBridge[i] = 1;
        }
   }
void findBridges() {
    fill(tin, tin + N, 0);
    fill(low, low + N, 0);
    fill(isBridge, isBridge + M, 0);
    clk = 1;
    for (int i = 0; i < N; i++) {
        if (!tin[i]) dfs(i);
}
```

8.20 Find Centroid

Description: Given a tree (don't forget to make it 'undirected'), find it's centroids. @Time : O(V)

```
#pragma once
#include "../Contest/template.cpp"
void dfs(int u, int p, int n, vi2d &g, vi &sz, vi &centroid) {
    sz[u] = 1;
    bool iscentroid = true;
    for (auto v : g[u])
```

```
if (v != p) {
          dfs(v, u, n, g, sz, centroid);
          if (sz[v] > n / 2) iscentroid = false;
          sz[u] += sz[v];
     }
     if (n - sz[u] > n / 2) iscentroid = false;
     if (iscentroid) centroid.eb(u);
}
vi getCentroid(vi2d &g, int n) {
     vi centroid;
     vi sz(n);
     dfs(0, -1, n, g, sz, centroid);
     return centroid;
}
```

8.21 Find bridges (online)

```
// O((n+m)*log(n))
struct BridgeFinder {
    // 2ecc = 2 edge conected component
    // cc = conected component
    vector<int> parent, dsu_2ecc, dsu_cc, dsu_cc_size;
    int bridges, lca iteration;
    vector<int> last visit;
    BridgeFinder(int n)
        : parent(n, -1),
          dsu 2ecc(n),
          dsu cc(n),
          dsu cc size(n, 1),
          bridges(0),
          lca iteration(0),
          last visit(n) {
        for (int i = 0; i < n; i++) {
            dsu \ 2ecc[i] = i;
            dsu cc[i] = i;
    int find 2ecc(int v) {
        if (\overline{v} == -1) return -1;
        return dsu 2ecc[v] == v ? v : dsu 2ecc[v] = find 2ecc(dsu 2ecc[v])
    int find cc(int v) {
        v = find 2ecc(v);
        return dsu cc[v] == v ? v : dsu cc[v] = find cc(dsu cc[v]);
    void make root(int v) {
        v = find 2ecc(v);
        int root = v;
        int child = -1;
        while (v != -1) {
            int p = find 2ecc(parent[v]);
```

```
parent[v] = child;
         dsu cc[v] = root;
         child = v;
         v = p;
    dsu cc size[root] = dsu cc size[child];
void merge_path(int a, int b) {
    ++lca iteration;
    vector<int> path a, path b;
    int lca = -1:
    while (lca == -1) {
         if (a != -1) {
             a = find 2ecc(a);
             path a.push back(a);
             if (last visit[a] == lca iteration) {
                 lca = a:
                 break:
             last visit[a] = lca iteration;
             a = parent[a];
         if (b != -1) {
             b = find 2ecc(b);
             path b.push back(b);
             if (last visit[b] == lca iteration) {
                 lca = b:
                 break:
             last visit[b] = lca iteration;
             b = parent[b];
    for (auto v : path a) {
         dsu \ 2ecc[v] = \overline{l}ca;
         if (v == lca) break;
         --bridges;
    for (auto v : path b) {
        dsu \ 2ecc[v] = \overline{l}ca;
         if (v == lca) break;
         --bridges;
    }
}
void add edge(int a, int b) {
    a = \overline{find} \ 2ecc(a);
    b = find 2ecc(b);
    if (a == b) return;
    int ca = find cc(a);
    int cb = find cc(b);
    if (ca != cb) {
         ++bridges;
         if (dsu_cc_size[ca] > dsu_cc_size[cb]) {
             swap(a, b);
```

```
swap(ca, cb);
}
make_root(a);
parent[a] = dsu_cc[a] = b;
dsu_cc_size[cb] += dsu_cc_size[a];
} else {
    merge_path(a, b);
}
};
```

8.22 Floyd Warshall

Description: Simply finds the minimal distance for each node to every other node. $O(V^3)$

8.23 Functional/Successor Graph

Description: Given a functional graph find the vertice after k moves starting at u and also the distance between u and v, if it's impossible to reach v starting at u returns -1.

Time: build $O(N \cdot MAXLOG2)$, kth O(MAXLOG2), dist O(MAXLOG2)

```
const int MAXN(2 '000' 000), MAXLOG2(24);
int N;
vi2d succ(MAXN, vi(MAXLOG2 + 1));
vi dst(MAXN, 0);
int vis[MAXN];
void dfsbuild(int u) {
    if (vis[u]) return;
    vis[u] = 1;
    int v = succ[u][0];
    dfsbuild(v);
    dst[u] = dst[v] + 1;
void build() {
    for (int i = 0; i < N; i++) {
        if (not vis[i]) dfsbuild(i);
    for (int k = 1; k \le MAXLOG2; k++) {
        for (int i = 0; i < N; i++) {
            succ[i][k] = succ[succ[i][k-1]][k-1];
```

```
}
int kth(int u, ll k) {
    if (k <= 0) return u;
    for (int i = 0; i <= MAXL0G2; i++)
        if ((1ll << i) & k) u = succ[u][i];
    return u;
}
int dist(int u, int v) {
    int cu = kth(u, dst[u]);
    if (kth(u, dst[u] - dst[v]) == v)
        return dst[u] - dst[v];
    else if (kth(cu, dst[cu] - dst[v]) == v)
        return dst[u] + (dst[cu] - dst[v]);
    else
        return -1;
}
</pre>
```

8.24 Heavy light decomposition (supreme)

```
struct HLD {
   int V:
   int id:
   int nb_heavy_path;
   std::vector<std::vector<int>> g;
   std::vector<pair<int, int>> edges; // edges of the tree
                                        // par[i] = parent of
   std::vector<int> par;
                                        // vertex i (Default: -1)
   std::vector<int> depth;
                                        // depth[i] = distance between
   root
                                        // and vertex i
                                        // subtree sz[i] = size of
   std::vector<int> subtree sz;
                                        // subtree whose root is i
   std::vector<int> heavy child;
                                        // heavy child[i] = child of
                                        // vertex i on heavy path
                                        // (Default: -1)
   std::vector<int> tree id;
                                        // tree id[i] = id of tree vertex
                                        // i belongs to
   std::vector<int> aligned id,
       aligned id inv;
                                     // aligned id[i] = aligned
                                     // id for vertex i
                                     // (consecutive on heavy
                                     // edges)
   std::vector<int> head;
                                     // head[i] = id of vertex on heavy
                                     // path of vertex i, nearest to root
   std::vector<int> head ids;
                                     // consist of head vertex id's
   std::vector<int> heavy path id; // heavy path id[i] =
                                     // heavy_path_id for vertex
                                     // [i]
   HLD(const std::vector<std::vector<int>> &e, vector<int> roots = {0})
        : HLD((int)e.size()) {
       q = e;
       build(roots);
   HLD(int sz = 0)
        : V(sz),
```

```
id(0),
      nb heavy path(0),
      g(sz),
      par(sz),
      depth(sz),
      subtree_sz(sz),
      heavy child(sz),
      tree id(sz, -1),
      aligned id(sz),
      aligned id inv(sz),
      head(sz),
      heavy path id(sz, -1) {}
void add edge(int u, int v) {
    edges.emplace back(u, v);
    q[u].emplace \overline{b}ack(v);
    g[v].emplace back(u);
void build dfs(int root) {
    std::stack<std::pair<int, int>> st;
    par[root] = -1;
    depth[root] = 0;
    st.emplace(root, 0);
    while (!st.empty()) {
        int now = st.top().first;
        int &i = st.top().second;
        if (i < (int)g[now] size()) {</pre>
             int nxt = q[now][i++];
             if (nxt == par[now]) continue;
             par[nxt] = now;
             depth[nxt] = depth[now] + 1;
             st.emplace(nxt, 0);
        } else {
             st.pop();
            int max_sub_sz = 0;
             subtree sz[now] = 1;
             heavy c\overline{h}ild[now] = -1;
             for (auto nxt : q[now]) {
                 if (nxt == par[now]) continue;
                 subtree sz[now] += subtree sz[nxt];
                 if (max sub sz < subtree sz[nxt])</pre>
                     max sub sz = subtree sz[nxt], heavy child[now] =
nxt;
        }
    }
void build bfs(int root, int tree id now) {
    std::queue<int> q({root});
    while (!q.empty()) {
        int h = q.front();
        q.pop();
        head ids.emplace back(h);
         for (int now = h; now != -1; now = heavy child[now]) {
             tree id[now] = tree id now;
```

```
aligned id[now] = id++;
            aligned id inv[aligned id[now]] = now;
            heavy path id[now] = nb heavy path;
            head[now] = h:
            for (int nxt : q[now])
                if (nxt != par[now] and nxt != heavy child[now])
                     a.push(nxt):
        nb heavy path++;
void build(std::vector<int> roots = {0}) {
    int tree id now = 0;
    for (auto r : roots) build_dfs(r), _build_bfs(r, tree_id_now++);
// data[i] = value of vertex i
template <class T>
std::vector<T> segtree rearrange(const std::vector<T> &data) const {
    assert(int(data.size()) == V);
    std::vector<T> ret;
    ret.reserve(V);
    for (int i = 0; i < V; i++) ret.emplace back(data[aligned id inv[i</pre>
]]);
    return ret;
// data[i] = weight of edge[i]
template <class T>
std::vector<T> segtree rearrange weighted(
    const std::vector<T> &data) const {
    assert(data.size() == edges.size());
    vector<T> ret(V);
    for (int i = 0; i < (int)edges.size(); i++) {</pre>
        auto [u, v] = edges[i];
        if (depth[u] > depth[v]) swap(u, v);
        ret[aligned_id[v]] = data[i];
    return ret;
int segtree edge index(int i) const {
    auto [u, v] = edges[i];
    if (depth[u] > depth[v]) swap(u, v);
    return aligned id[v];
// query for vertices on path [u, v] (INCLUSIVE)
void for each vertex(int u, int v, const auto &f) const {
    static assert(std::is invocable r v<void, decltype(f), int, int>);
    assert(tree id[u] == tree id[v] and tree id[u] >= 0);
    while (true) {
        if (aligned id[u] > aligned id[v]) std::swap(u, v);
        f(std::max(\overline{aligned}_id[head[\overline{v}]), aligned_id[u]), aligned_id[v])
        if (head[u] == head[v]) break;
        v = par[head[v]];
}
```

```
void for each vertex noncommutative(int from, int to, const auto &fup,
                                     const auto &fdown) const {
    static assert(std::is invocable r v<void, decltype(fup), int, int
>);
    static assert(std::is invocable r v<void, decltype(fdown), int,
int>);
    assert(tree id[from] == tree id[to] and tree id[from] >= 0);
    int u = from, v = to:
    const int lca = lowest common ancestor(u, v), dlca = depth[lca];
    while (u >= 0 and depth[u] > dlca) {
        const int p = (depth[head[u]] > dlca ? head[u] : lca);
        fup(aligned id[p] + (p == lca), aligned id[u]), u = par[p];
    static std::vector<std::pair<int, int>> lrs;
    int sz = 0:
    while (v >= 0 \text{ and } depth[v] >= dlca) {
        const int p = (depth[head[v]] >= dlca ? head[v] : lca);
        if (int(lrs.size()) == sz) lrs.emplace back(0, 0);
        lrs.at(sz++) = \{p, v\}, v = par.at(p);
    while (sz--)
        fdown(aligned id[lrs.at(sz).first], aligned id[lrs.at(sz).
second]);
// query for edges on path [u, v]
void for_each_edge(int u, int v, const auto &f) const {
    static assert(std::is invocable r v<void, decltype(f), int, int>);
    assert(tree_id[u] == tree_id[v] and tree_id[u] >= 0);
    while (true) {
        if (aligned_id[u] > aligned_id[v]) std::swap(u, v);
        if (head[u] != head[v]) {
            f(aligned id[head[v]], aligned id[v]);
            v = par[head[v]];
        } else {
            if (u != v) f(aligned id[u] + 1, aligned id[v]);
            break:
    }
// lowest common ancestor: O(log V)
int lowest common ancestor(int u, int v) const {
    assert(tree id[u] == tree id[v] and tree id[u] >= 0);
    while (true) {
        if (aligned_id[u] > aligned_id[v]) std::swap(u, v);
        if (head[u] == head[v]) return u;
        v = par[head[v]];
int distance(int u, int v) const {
    assert(tree id[u] == tree id[v] and tree id[u] >= 0);
    return depth[u] + depth[v] - 2 * depth[lowest common ancestor(u, v
)];
// Level ancestor, O(log V)
```

```
// if k-th parent is out of range, return -1
    int kth parent(int v, int k) const {
        if (k < 0) return -1;
        while (v \ge 0) {
            int h = head.at(v), len = depth.at(v) - depth.at(h);
            if (k <= len) return aligned id inv.at(aligned id.at(v) - k);</pre>
            k = len + 1, v = par.at(h);
        return -1;
   }
    // Jump on tree, O(log V)
    int s to t by k steps(int s, int t, int k) const {
        if (\overline{k} < 0) return -1;
        if (k == 0) return s:
        int lca = lowest common ancestor(s, t);
        if (k <= depth.at(s) - depth.at(lca)) return kth parent(s, k);</pre>
        return kth parent(t, depth.at(s) + depth.at(t) - depth.at(lca) * 2
    - k);
};
```

8.25 Kruskal

Description: Find the minimum spanning tree of a graph. **Time**: $O(E \log E)$

```
#include "./../Data Structures/DSU.cpp"
vector<tuple<ll, int, int>> kruskal(int n, vector<tuple<ll, int, int>> &
    edges) {
    DSU dsu(n);
    vector<tuple<ll, int, int>> ans;
    sort(all(edges));
    for (auto [a, b, c] : edges) {
        if (dsu.same_set(b, c)) continue;
        ans.emplace_back(a, b, c);
        dsu.union_set(b, c);
    }
    return ans;
}
```

8.26 Lowest Common Ancestor

Description: Given two nodes of a tree find their lowest common ancestor, or their distance

```
#pragma once
#include "../Contest/template.cpp"
template <typename T>
struct SparseTable {
   vector<T> v;
   int n;
   static const int b = 30;
   vi mask, t;
```

```
int op(int x, int y) { return v[x] < v[y] ? x : y; }
    int msb(int x) { return builtin clz(1) - builtin clz(x); }
    SparseTable() {}
    SparseTable(const vector<T> &v ) : v(v ), n(v.size()), mask(n), t(n) {
        for (int i = 0, at = 0; i < n; \max \overline{k}[i++] = at |= 1) {
            at = (at << 1) & ((1 << b) - 1);
            while (at and op(i, i - msb(at & -at)) == i) at ^= at & -at;
        for (int i = 0; i < n / b; i++)
            t[i] = b * i + b - 1 - msb(mask[b * i + b - 1]);
        for (int j = 1; (1 << j) <= n / b; j++)
            for (int i = 0; i + (1 << j) <= n / b; i++)
                t[n / b * i + i] = op(t[n / b * (j - 1) + i],
                                       t[n / b * (j - 1) + i + (1 << (j -
   1))]);
    int small(int r, int sz = b) { return r - msb(mask[r] \& ((1 << sz) -
   1)): }
    T query(int l, int r) {
        if (r - l + 1 \le b) return small(r, r - l + 1);
        int ans = op(small(l + b - 1), small(r));
        int x = l / b + 1, y = r / b - 1;
        if (x \le y) {
            int j = msb(y - x + 1);
                op(ans, op(t[n / b * j + x], t[n / b * j + y - (1 << j) +
   1]));
        return ans:
};
struct LCA {
    SparseTable<int> st;
    int n;
    vi v, pos, dep;
    vll wdep;
    LCA(const Graph \&g, int root) : n(len(g)), pos(n), wdep(n) {
        dfs(root, 0, -1, q);
        st = SparseTable<int>(vector<int>(all(dep)));
    void dfs(int i, int d, int p, const Graph &g) {
        v.eb(len(dep)) = i, pos[i] = len(dep), dep.eb(d);
        for (auto [w, j] : g[i])
            if (j != p) {
                wdep[j] = wdep[i] + w;
                dfs(i, d + 1, i, q);
                v.eb(len(dep)) = i, dep.eb(d);
    int lca(int a, int b) {
        int l = min(pos[a], pos[b]);
        int r = max(pos[a], pos[b]);
        return v[st.query(l, r)];
    il dist(int a, int b) { return wdep[a] + wdep[b] - 2ll * wdep[lca(a, b
```

```
)]; }
};
```

8.27 Lowest Common Ancestor (Binary Lifting)

Description: Given a directed tree, finds the LCA between two nodes using binary lifting, and answer a few queries with it.

Usage:

- lca: returns the LCA between the two given nodes
- on path: fids if c is in the path from a to b

Time: build $O(N \cdot MAXLOG_2)$, all queries $O(MAXLOG_2)$

```
struct LCA {
   int n;
    const int maxlog;
   vector<vector<int>> up;
   vector<int> depth:
   LCA(const vector<vector<int>> &tree)
        : n(tree.size()),
          maxlog(ceil(log2(n))),
          up(n, vector<int>(maxlog + 1)),
          depth(n, -1) {
        for (int i = 0; i < n; i++) {
            if (depth[i] == -1) {
                depth[i] = 0;
                dfs(i, -1, tree);
            }
        }
   void dfs(int u, int p, const vector<vector<int>> &tree) {
        if (p != -1) {
            depth[u] = depth[p] + 1;
            up[u][0] = p;
            for (int i = 1; i <= maxlog; i++) {
                up[u][i] = up[up[u][i - 1]][i - 1];
        for (int v : tree[u]) {
            if (v == p) continue;
            dfs(v, u, tree);
   int kth jump(int u, int k) {
        for (int i = maxlog; i \ge 0; i--) {
            if ((1 << i) & k) {
                u = up[u][i];
        return u;
   int lca(int u, int v) {
        if (depth[u] < depth[v]) swap(u, v);</pre>
        int diff = depth[u] - depth[v];
        u = kth jump(u, diff);
```

```
if (u == v) return u;
for (int i = maxlog; i >= 0; i--) {
        if (up[u][i] != up[v][i]) {
            u = up[u][i];
            v = up[v][i];
        }
    }
    return up[u][0];
}
bool on_path(int u, int v, int s) {
        int uv = lca(u, v), us = lca(u, s), vs = lca(v, s);
        return (uv == s or (us == uv and vs == s) or (vs == uv and us == s));
}
int dist(int u, int v) {
        return depth[u] + depth[v] - 2 * depth[lca(u, v)];
}
};
```

8.28 Maximum flow (Dinic)

Description: Finds the **maximum flow** in a graph network, given the **source** s and the **sink** t. Add edge from a to b with capcity c.

Time: In general $O(E \cdot V^2)$, if every capacity is 1, and every vertice has in degree equal 1 or out degree equal 1 then $O(E \cdot \sqrt{V})$,

Warning: Suffle the edges list for every vertice may take you out of the worst case

```
struct Dinic {
    struct Edge {
        int to, rev;
        ll c, oc;
        Il flow() { return max(oc - c, OLL); } // if you need flows
    };
    vi lvl, ptr, q;
    vector<vector<Edge>> adj;
    Dinic(int n) : lvl(n), ptr(n), q(n), adj(n) {}
    void addEdge(int a, int b, ll c, ll rcap = 0) {
        adj[a].pb({b, len(adj[b]), c, c});
        adj[b].pb({a, len(adj[a]) - 1, rcap, rcap});
    ll dfs(int v, int t, ll f) {
        if (v == t || !f) return f;
        for (int &i = ptr[v]; i < len(adj[v]); i++) {</pre>
            Edge &e = adj[v][i];
            if (lvl[e.to] == lvl[v] + 1)
                if (ll p = dfs(e.to, t, min(f, e.c))) {
                    e.c -= p, adj[e.to][e.rev].c += p;
                     return p;
        return 0:
```

```
ll maxFlow(int s, int t) {
        Il flow = 0;
        q[0] = s;
        rep(L, 0, 31) {
            do { // 'int L=30' maybe faster for random
                  // data
                lvl = ptr = vi(len(q));
                int qi = 0, qe = lvl[s] = 1;
                while (qi < qe && !lvl[t]) {</pre>
                    int v = q[qi++];
                    for (Edge e : adi[v])
                        if (!lvl[e.to] && e.c >> (30 - L))
                             q[qe++] = e.to, lvl[e.to] = lvl[v] + 1;
                while (ll p = dfs(s, t, LLONG MAX)) flow += p;
            } while (lvl[t]);
        return flow;
    bool leftOfMinCut(int a) { return lvl[a] != 0; }
};
```

8.29 Minimum Cost Flow

Description: Given a network find the minimum cost to achieve a flow of at most f. Works with **directed** and **undirected** graphs **Usage**:

- add(u, v, c, w): adds an edge from u to v with capacity c and cost w.
- flow(f): return a pair (flow, cost) with the maximum flow until f with source at s and sink at t, with the minimum cost possible.

Time: $O(N \cdot M + f \cdot m \log n)$

```
template <typename T>
struct MinCostFlow {
    struct Edge {
        int to;
        ll c, rc; // capcity, residual capacity
                   // cost
       T w:
    };
   <u>int</u> n, s, t;
   vector<Edge> edges;
   vi2d q;
   vector<T> dist;
   vi pre;
   MinCostFlow() {}
   MinCostFlow(int n_, int _s, int _t) : n(n_), s(_s), t(_t), g(n) \{ \}
   void addEdge(int u, int v, ll c, T w) {
        g[u].pb(len(edges));
        edges.eb(v, c, 0, w);
        g[v].pb(len(edges));
        edges.eb(u, 0, 0, -w);
   // {flow, cost}
   pair<ll, T> flow(ll flow limit = LLONG MAX) {
```

```
II flow = 0:
     T cost = 0:
     while (flow < flow_limit and dijkstra(s, t)) {</pre>
         ll aug = LLONG MAX;
         for (int i = t; i != s; i = edges[pre[i] ^ 1].to) {
             aug = min({flow limit - flow, aug, edges[pre[i]].c});
         for (int i = t; i != s; i = edges[pre[i] ^ 1].to) {
             edges[pre[i]].c -= aug;
             edges[pre[i] ^1].c += aug;
             edges[pre[i]].rc += aug;
             edges[pre[i] ^ 1].rc -= aug;
         flow += aug;
         cost += (T)aug * dist[t];
     return {flow, cost};
 // Needs to be called after flow method
 vi2d paths() {
     vi2d p;
     for (;;) {
         int cur = s;
         auto &res = p.eb();
         res.pb(cur);
         while (cur != t) {
             bool found = false;
             for (auto i : q[cur]) {
                 auto &[v, , c, cost] = edges[i];
                 if (c > 0) {
                     --c:
                     res.pb(cur = v);
                     found = true:
                     break:
                 }
             if (!found) break;
         if (cur != t) {
             p.ppb();
             break;
     }
     return p;
private:
 bool bellman ford(int s, int t) {
     dist.assign(n, numeric limits<T>::max());
     pre.assign(n, -1);
     vc ing(n, false);
     queue<int> q;
     dist[s] = 0;
     q.push(s);
```

```
while (len(q)) {
        int u = q.front();
        q.pop();
        inq[u] = false;
        for (int i : q[u]) {
            auto [v, c, w, _] = edges[i];
             auto new dist = dist[u] + w;
             if (c > \overline{0} \text{ and } dist[v] > \text{new dist})  {
                 dist[v] = new dist;
                 pre[v] = i;
                 if (not ing[v]) {
                     ina[v] = true:
                     q.push(v);
                 }
            }
        }
    return dist[t] != numeric limits<T>::max();
bool dijkstra(int s, int t) {
    dist.assign(n, numeric limits<T>::max());
    pre.assign(n, -1);
    dist[s] = 0;
    using PQ = pair<T, int>;
    pgmn<PQ> pg;
    pq.emp(0, s);
    while (len(pg)) {
        auto [cost, u] = pq.top();
        pq.pop();
        if (cost != dist[u]) continue;
        for (int i : g[u]) {
             auto [v, c, , w] = edges[i];
             auto new dist = dist[u] + w;
             if (c > \overline{0} \text{ and } dist[v] > new_dist) {
                 dist[v] = new dist;
                 pre[v] = i;
                 pq.emp(new dist, v);
        }
    return dist[t] != numeric limits<T>::max();
```

8.30 Minimum Vertex Cover (already divided)

};

Description: Given a bipartite graph g with n vertices at left and m vertices at right, where g[i] are the possible right side matches of vertex i from left side, find a minimum vertex cover. The size is the same as the size of the maximum matching, and the complement is a maximum independent set.

```
vector<int> min_vertex_cover(vector<vector<int>> &g, int n, int m) {
   vector<int> match(m, -1), vis;
```

```
auto find = [&](auto &&self, int j) -> bool {
    if (match[i] == -1) return 1;
    vis[i] = 1;
    int di = match[j];
    for (int e : q[di])
        if (!vis[e] and self(self, e)) {
            match[e] = di;
            return 1;
    return 0;
};
for (int i = 0; i < (int)q.size(); i++) {
    vis.assign(match.size(), 0);
    for (int j : q[i]) {
        if (find(find, j)) {
            match[j] = i;
            break:
        }
    }
int res = (int)match.size() - (int)count(match.begin(), match.end(),
vector<char> lfound(n, true), seen(m);
for (int it : match)
    if (it != -1) lfound[it] = false;
vector<int> q, cover;
for (int i = 0; i < n; i++)
    if (lfound[i]) q.push back(i);
while (!q.empty()) {
    int i = q.back();
    g.pop back();
    lfound[i] = 1;
    for (int e : g[i])
        if (!seen[e] and match[e] !=-1) {
            seen[e] = true;
            q.push back(match[e]);
for (int i = 0; i < n; i++)
    if (!lfound[i]) cover.push back(i);
for (int i = 0; i < m; i++)
    if (seen[i]) cover.push back(n + i);
assert((int)size(cover) == res);
return cover;
```

8.31 Prim (MST)

Description: Given a graph with N vertex finds the minimum spanning tree, if there is no such three returns inf, it starts using the edges that connect with each $s_i \in s$, if none is provided than it starts with the edges of node 0.

Time: $O(V \log E)$

}

```
#include "../Contest/template.cpp"
const int MAXN(1'00'000);
int N:
vector<pair<ll, int>> G[MAXN];
ll prim(vi s = vi(1, 0)) {
   priority queue<pair<ll, int>, vector<pair<ll, int>>, greater<pair<ll,</pre>
   int>>>
        pq;
   vector<char> ingraph(MAXN);
   int ingraphcnt(0);
   for (auto si : s) {
        ingraphcnt++;
        ingraph[si] = true;
        for (auto &[w, v] : G[si]) pg.emplace(w, v);
   ll\ mstcost = 0;
   while (ingraphent < N and !pg.empty()) {</pre>
        ll w:
        int v;
        do {
            tie(w, v) = pq.top();
            pq.pop();
        } while (not pg.empty() and ingraph[v]);
        mstcost += w, ingraph[v] = true, ingraphcnt++;
        for (auto &[w2, v2] : G[v]) {
            pq.emplace(w2, v2);
    return ingraphent == N ? mstcost : oo:
```

8.32 Reachability Tree

```
#include <bits/stdc++.h>
using namespace std;
const int MAXN = 20000 + 100;
int dsu[MAXN];
int n:
const int MAXM = 100000:
int U[MAXM], V[MAXM];
vector<int> adi[MAXN];
int getRoot(int u) {
   if (u == dsu[u]) return u;
   return dsu[u] = getRoot(dsu[u]);
void addEdge(int u, int v) {
   u = aetRoot(u):
   v = getRoot(v);
   dsu[n] = n:
   dsu[u] = dsu[v] = n;
```

```
adj[n].push_back(u);
    if (u != v) adj[n].push_back(v);
    ++n;
}

void build() {
    for (int i = 0; i < n; ++i) dsu[i] = i;
    for (int i = 0; i < m; ++i) addEdge(U[i], V[i]);
}
int32_t main() {
    ios_base::sync_with_stdio(0);
    cin.tie(0);
}</pre>
```

8.33 Shortest Path With K-edges

Description: Given an adjacency matrix of a graph, and a number K computes the shortest path between all nodes that uses exactly K edges, so for $0 \le i, j \le N - 1$ ans[i][j] = "the shortest path between i and j that uses exactly K edges, remember to initialize the adjacency matrix with ∞ .

Time: $O(N^3 \cdot \log K)$

```
template <typename T>
vector<vector<T>> prod(vector<vector<T>> &a, vector<vector<T>> &b) {
    const T oo = numeric limits<T>::max();
    int n = a.size();
    vector<vector<T>> c(n, vector<T>(n, _oo));
    for (int i = 0; i < n; i++)
        for (int j = 0; j < n; j++)
            for (int k = 0; k < n; k++)
                if (a[i][k] != oo and b[k][j] != oo)
                    c[i][i] = m\overline{i}n(c[i][j], a[i][k] + b[k][j]);
    return c;
template <typename T>
vector<vector<T>> shortest with k moves(vector<vector<T>> adj, long long k
   ) {
    if (k == 1) return adj;
    auto ans = adj;
    k--;
    while (k) {
        if (k & 1) ans = prod(ans, adj);
        k >>= 1:
        adj = prod(adj, adj);
    return ans;
```

8.34 Strongly Connected Components (struct)

Description: Find the connected component for each edge (already in a topological order), some additional functions are also provided.

```
Time: Build: O(V+E)
struct SCC {
   int n, num_sccs;
   vi2d adi;
   vi scc_id;
```

```
SCC(int n) : n(n), num sccs(0), adj(n), scc id(n, -1) {}
   void add edge(int u, int v) { adj[u].eb(v); }
   void find sccs() {
        int timer = 1:
        vi tin(n), st:
        st.reserve(n);
        function<int(int)> dfs = [&](int u) -> int {
            int low = tin[u] = timer++, siz = len(st);
            st.eb(u);
            for (int v : adi[u])
                if (scc id[v] < 0) low = min(low, tin[v] ? tin[v] : dfs(v)
   );
            if (tin[u] == low) {
                rep(i, siz, len(st)) scc id[st[i]] = num sccs;
                st.resize(siz):
                num_sccs++;
            return low;
        };
        for (int i = 0; i < n; i++)
            if (!tin[i]) dfs(i);
   vector<set<int>> build qscc() {
        vector<set<int>> qscc;
        for (int i = 0; i < len(adj); ++i)</pre>
            for (auto j : adj[i])
                if (scc_id[i] != scc_id[j]) gscc[scc_id[i]].emplace(scc_id
   [j]);
        return gscc;
   vi2d per comp() {
        vi2d ret(num sccs);
        rep(i, 0, n) ret[scc id[i]].eb(i);
        reverse(all(ret)); // already in topological order ;)
        return ret:
};
      Topological Sorting (Kahn)
```

Description: Finds the topological sorting in a **DAG**, if the given graph is not a **DAG** than an empty vector is returned, need to 'initialize' the **INCNT** as you build the graph. Time: O(V+E)

```
const int MAXN(2 '00' 000);
int INCNT[MAXN];
vi2d GOUT(MAXN);
int N:
```

```
vi toposort() {
    vi order;
    queue<int> q;
    for (int i = 0; i < N; i++)
        if (!INCNT[i]) g.emplace(i);
    while (!q.empty()) {
        auto u = q.front();
        q.pop();
        order emplace back(u);
        for (auto v : GOUT[u]) {
            INCNT[v]--;
            if (INCNT[v] == 0) g.emplace(v);
    }
    return len(order) == N ? order : vi();
```

8.36 Topological Sorting (Tarjan)

Description: Finds a the topological order for the graph, if there is no such order it means the graph is cyclic, then it returns an empty vector

Time: O(V+E)

```
const int maxn(1 '00' 000);
int n, m;
vi q[maxn];
int not found = 0, found = 1, processed = 2;
int state[maxn];
bool dfs(int u, vi &order) {
    if (state[u] == processed) return true;
    if (state[u] == found) return false;
    state[u] = found;
    for (auto v : q[u]) {
        if (not dfs(v, order)) return false;
    state[u] = processed:
    order.emplace back(u);
    return true;
vi topo sort() {
    vi order:
    memset(state, 0, sizeof state);
    for (int u = 0: u < n: u++) {
        if (state[u] == not found and not dfs(u, order)) return {};
    reverse(all(order));
    return order;
```

8.37 Tree Isomorphism (not rooted)

Description: Two trees are considered **isomorphic** if the hash given by thash() is the same. **Time**: $O(V \cdot \log V)$

```
map<vi, int> mphash;
struct Tree {
    int n;
    vi2d q;
    vi sz, cs;
   Tree(int n): n(n), g(n), sz(n) {}
    void add edge(int u, int v) {
        g[u].emplace back(v);
        g[v].emplace back(u);
    void dfs centroid(int v, int p) {
        sz[v] = 1:
        bool cent = true;
        for (int u : q[v])
            if (u != p) {
                dfs centroid(u, v);
                sz[v] += sz[u];
                cent &= not(sz[u] > n / 2);
        if (cent and n - sz[v] <= n / 2) cs.push back(v);</pre>
    int fhash(int v, int p) {
        vi h:
        for (int u : g[v])
            if (u != p) h.push back(fhash(u, v));
        sort(all(h));
        if (!mphash.count(h)) mphash[h] = mphash.size();
        return mphash[h];
    ll thash() {
        cs.clear();
        dfs centroid(0, -1);
        if (cs.size() == 1) return fhash(cs[0], -1);
        ll h1 = fhash(cs[0], cs[1]), h2 = fhash(cs[1], cs[0]);
        return (min(h1, h2) \ll 30ll) + max(h1, h2);
};
```

8.38 Tree Isomorphism (rooted)

Description: Given a rooted tree find the hash of each subtree, if two roots of two distinct trees have the same hash they are considered isomorphic

Time: hash first time in $O(\log N_v \cdot N_v)$ where (N_v) is the of the subtree of v

```
map<vi, int> hasher;
int hs = 0;
struct RootedTreeIso {
   int n;
```

```
vi2d adj;
vi hashes;
RootedTreeIso(int _n) : n(_n), adj(_n), hashes(_n, -1) {};
void add_edge(int u, int v) {
    adj[u].emplace_back(v);
    adj[v].emplace_back(u);
}
int hash(int u, int p = -1) {
    if (hashes[u] != -1) return hashes[u];
    vi children;
    for (auto v : adj[u])
        if (v != p) children.emplace_back(hash(v, u));
    sort(all(children));
    if (!hasher.count(children)) hasher[children] = hs++;
    return hashes[u] = hasher[children];
}
};
```

8.39 Tree diameter (DP)

```
const int MAXN(1 '000' 000);
int N;
vi G[MAXN];
int diameter, toLeaf[MAXN];
void calcDiameter(int u = 0, int p = -1) {
   int d1, d2;
   d1 = d2 = -1;
   for (auto v : G[u]) {
        if (v != p) {
            calcDiameter(v, u);
            d1 = max(d1, toLeaf[v]);
            tie(d1, d2) = minmax({d1, d2});
        }
   }
   toLeaf[u] = d2 + 1;
   diameter = max(diameter, d1 + d2 + 2);
}
```

8.40 Tree edge queries

```
h(n, -1),
      par(n, vector<int>(LOG + 1, root)),
      ed(n, vector < T > (LOG + 1, E)) {
    h[root] = 0, dfs(root, q);
void dfs(int u, const Graph& g) {
    for (auto& [w, v] : g[u]) {
        if (h[v] == -1)
            h[v] = h[u] + 1, par[v][0] = u, ed[v][0] = w;
            for (int k = 0, p; k < LOG; k++) {
                p = par[v][k];
                par[v][k + 1] = par[p][k];
                ed[v][k + 1] = F(ed[v][k], ed[p][k]);
            dfs(v, g);
        }
    }
pair<int, T> up(int u, int dis) {
    T res = E;
    for (int k = 0; k \le LOG; k++) {
        if (dis & (1 << k)) {
            res = F(res, ed[u][k]);
            u = par[u][k];
        }
    return {u, res};
pair<int, T> lca(int u, int v) {
    if (h[u] > h[v]) swap(u, v);
    T res = E:
    tie(v, res) = up(v, h[v] - h[u]);
    if (v == u) return {v, res};
    for (int k = LOG; \sim k; k--) {
        if (par[u][k] != par[v][k]) {
            res = F(res, ed[v][k]):
            res = F(res, ed[u][k]);
            u = par[u][k], v = par[v][k];
        }
    res = F(res, ed[v][0]);
    res = F(res, ed[u][0]);
    return {par[v][0], res};
```

8.41 Virtual Tree

};

```
#pragma once
#include "../Contest/template.cpp"
#include "./Lowest common ancestor (sparse table).cpp"
struct VTree {
   int n;
   LCA lca;
```

```
VTree(const Graph\& q, int root = 0) : n(len(q)), lca(q, root) {}
    pair<vector<tuple<li, int, int>>, int> vtree(vector<int> vs) {
        sort(vs.begin(), vs.end(),
             [&](int u, int v) { return lca.pos[u] < lca.pos[v]; });</pre>
        for (int i = 0, n = size(vs); i + 1 < n; i++) {
            vs.eb(lca.lca(vs[i], vs[i + 1]));
        sort(vs.begin(), vs.end(),
             [&](int u, int v) { return lca.pos[u] < lca.pos[v]; });
        vs.erase(unique(all(vs)), vs.end());
        vi st{vs.front()};
        vector<tuple<ll, int, int>> ret;
        for (int i = 1; i < len(vs); i++) {
            int v = vs[i];
            while (len(st) >= 2 \& lca.lca(v, st.back()) != st.back()) {
                int a = end(st)[-2];
                int b = st.back();
                ll c = lca.dist(a, b);
                ret.eb(c, a, b);
                st.pop back();
            st.pb(v);
        while (len(st) \geq 2) {
            int a = end(st)[-2]:
            int b = st.back();
            ll c = lca.dist(a, b);
            ret.eb(c, a, b);
            st.pop back();
        return {ret, st.back()};
};
```

9 Linear Algebra

9.1 Matrix (primitive)

```
#include "../Contest/template.cpp"
template <typename T>
struct Matrix {
   int n, m;
   valarray<valarray<T>> v;
   Matrix(int _n, int _m, int id = 0) : n(_n), m(_m), v(valarray<T>(m), n
) {
    if (id) {
       rep(i, 0, n) v[i][i] = 1;
    }
}
valarray<T> &operator[](int x) { return v[x]; }
Matrix transpose() {
    Matrix newv(m, n);
    rep(i, 0, n) rep(j, 0, m) newv[j][i] = (*this)[i][j];
```

```
return newv;
Matrix operator+(Matrix &b) {
    Matrix ret(*this);
    return ret.v += b.v;
Matrix & operator += (Matrix &b) { return v += b.v; }
Matrix operator*(Matrix b) {
    Matrix ret(n, b.m);
    rep(i, 0, n) rep(i, 0, m) rep(k, 0, b.m) ret[i][k] +=
        v[i][j] * b.v[j][k];
    return ret:
Matrix & operator *= (Matrix b) { return *this = *this * b; }
Matrix power(ll exp) {
    Matrix in = *this;
    Matrix ret(n, n, 1);
    while (exp) {
        if (exp & 1) ret *= in:
        in *= in:
        \exp >>= 1;
    return ret;
 * Alters current matrix.
 * Does gaussian elimination and puts matrix in
 * upper echelon form (possibly reduced).
 * Returns the determinant of the square matrix
 * with side equal to the number of rows of the
 * original matrix.
T gaussjordanize(int reduced = 0) {
    T \det = T(1);
    int line = 0:
    rep(col, 0, m) {
        int pivot = line:
        while (pivot < n && v[pivot][col] == T(0)) pivot++;
        if (pivot >= n) continue;
        swap(v[line], v[pivot]);
        if (line != pivot) det *= T(-1);
        det *= v[line][line]:
        v[line] /= T(v[line][col]);
        if (reduced) rep(i, 0, line) {
                v[i] = T(v[i][col]) * v[line];
        rep(i, line + 1, n) { v[i] -= T(v[i][col]) * v[line]; }
        line++:
    return det * (line == n);
```

```
* Needs to be called in a square matrix that
* represents a system of linear equations. Returns {possible solution
* number of solutions (2 if infinite solutions)}
pair<vector<T>, int> solve system(vector<T> results) {
    Matrix aux(n, m + 1);
    rep(i, 0, n) {
        rep(i, 0, m) aux[i][i] = v[i][i];
        aux[i][m] = results[i]:
   T det = aux.gaussjordanize(1);
    int ret = 1 + (det == T(0));
    int n = results.size():
    rrep(i, n - 1, 0 - 1) {
        int pivot = 0;
        while (pivot < n && aux[i][pivot] == T(0)) pivot++;</pre>
        if (pivot == n) {
            if (aux[i][m] != T(0)) ret = 0;
        } else
            swap(aux[i], aux[pivot]);
    rrep(i, n - 1, 0 - 1) rep(j, i + 1, n) aux[i][m] =
        aux[i][i] * aux[i][m];
    rep(i, 0, n) results[i] = aux[i][m];
    rep(i, 0, n) rep(j, 0, m) v[i][j] = aux[i][j];
    return {results, ret}:
/* Does not alter current matrix. Returns {inverse matrix, is curent
* matrix invertable} */
pair<Matrix<T>, bool> find inverse() {
    int n = v.size();
    Matrix<T> aug(n, 2 * n);
    rep(i, 0, n) rep(i, 0, n) aug[i][i] = v[i][i];
    rep(i, 0, n) aug[i][n + i] = 1;
   T det = aug.gaussjordanize(1);
   Matrix<T> ret(n, n):
    rep(i, 0, n) ret[i] = valarray < T > (aug[i][slice(n, n, 1)]);
    return {ret, det != T(0)};
/* Returns rank of matrix. does not alter it. */
int get rank() const {
   if (m == 0) return 0;
   Matrix<T> aux(*this);
    aux.gaussjordanize();
    int resp = 0:
    rep(i, 0, n) resp += (aux[i] != valarray<T>(m)).sum();
    return resp;
```

```
};
    Math
10.1 Arithmetic Progression Sum
Usage:
  • s: first term
  • d: common difference
  • n: number of terms
ll arithmeticProgressionSum(ll s, ll d, ll n) {
    return (s + (s + d * (n - 1))) * n / 2ll;
10.2 Binomial
Time: O(N \cdot K)
Memory: O(K)
ll binom(ll n, ll k) {
    if (k > n) return 0;
    vll dp(k + 1, 0);
    dp[0] = 1;
    for (ll i = 1; i <= n; i++)
        for (ll i = k; i > 0; i--) dp[i] = dp[i] + dp[i - 1];
    return dp[k];
}
10.3 Binomial MOD
Description: find \binom{n}{k} (mod MOD)
  • precompute: on first call it takes O(MAXNBIN) to precompute the factorials
  • query: O(1).
Memory: O(MAXNBIN)
Warning: Remember to set MAXNBIN properly!
const ll MOD = 998244353:
inline ll binom(ll n, ll k) {
    static const int BINMAX = 2'000'000;
    static vll FAC(BINMAX + 1), FINV(BINMAX + 1);
    static bool done = false;
    if (!done) {
        vll INV(BINMAX + 1);
        FAC[0] = FAC[1] = INV[1] = FINV[0] = FINV[1] = 1;
        for (int i = 2; i <= BINMAX; i++) {</pre>
            FAC[i] = FAC[i - 1] * i % MOD;
            INV[i] = MOD - MOD / i * INV[MOD % i] % MOD;
```

FINV[i] = FINV[i - 1] * INV[i] % MOD;

done = true:

```
if (n < k \mid | n < 0 \mid | k < 0) return 0;
    return FAC[n] * FINV[k] % MOD * FINV[n - k] % MOD;
10.4 Chinese Remainder Theorem
Description: Find the solution X to the N modular equations.
                                 x \equiv a_1(modm_1)
                                x \equiv a_n(modm_n)
The m_i don't need to be coprime, if there is no solution then it returns -1.
Warning: Make sure all M_i are coprime!
#include "../Contest/template.cpp"
tuple<ll, ll, ll> ext gcd(ll a, ll b) {
    if (!a) return {b, 0, 1};
    auto [g, x, y] = ext_gcd(b % a, a);
    return \{q, y - b / a * x, x\};
template <typename T = ll>
struct crt {
    T a, m;
    crt() : a(0), m(1) {}
    crt(T a_, T m_) : a(a_), m(m_) {}
    crt operator*(crt C) {
        auto [g, x, y] = ext gcd(m, C.m);
        if ((a - C.a) \% q != 0) a = -1;
        if (a == -1 \text{ or } C.a == -1) \text{ return } crt(-1, 0);
        T lcm = m / g * C.m;
        T ans = a + (x * (C.a - a) / g % (C.m / g)) * m;
        return crt((ans % lcm + lcm) % lcm, lcm);
};
template <typename T = ll>
struct Congruence {
    T a, m;
};
template <typename T = ll>
T chinese remainder theorem(const vector<Congruence<T>> &equations) {
    crt<T> ans:
```

for (auto &[a , m] : equations) {

return ans.a;

ans = ans * crt<T>(a , m);

(1)

}

10.5 Derangement / Matching Problem

```
D_N = N! \left(1 - \frac{1}{1!} + \frac{1}{2!} - \frac{1}{3!} + \dots + (-1)^N \frac{1}{N!}\right)
Time: O(N)
#warning Remember to call precompute !
const l\bar{l} MOD = le9 + 7:
const int MAXN(1 '000' 000);
ll fats[MAXN + 1];
void precompute() {
    fats[0] = 1;
    for (ll i = 1; i <= MAXN; i++) {
         fats[i] = (fats[i - 1] * i) % MOD;
}
ll fastpow(ll a, ll p, ll m) {
    ll ret = 1:
    while (p) {
         if (p & 1) ret = (ret * a) % MOD;
         p >>= 1;
         a = (a * a) % MOD;
    return ret;
ll divmod(ll a, ll b) { return (a * fastpow(b, MOD - 2, MOD)) % MOD; }
ll derangement(const ll n) {
    ll \ ans = fats[n]:
    for (ll i = 1; i <= n; i++) {
         ll k = divmod(fats[n], fats[i]);
         if (i & 1) {
             ans = (ans - k + MOD) % MOD;
         } else {
             ans = (ans + k) % MOD;
    return ans;
```

Description: Computes the derangement of N, which is given by the formula:

10.6 Euler Phi

Description: Computes the number of positive integers less than N that are coprimes with N, in $O(\sqrt{N})$.

```
int phi(int n) {
    if (n == 1) return 1;
    auto fs = factorization(n);  // a vctor of pair or a map
    auto res = n;
    for (auto [p, k] : fs) {
        res /= p;
        res *= (p - 1);
    }
    return res;
}
```

10.7 Euler phi (in range)

Description: Computes the number of positive integers less than n that are coprimes with N, in the range [1, N], in $O(N \log N)$.

10.8 Extended Euclidian algorithm

Description: Finds the gcd between a and b and x and y such that ax + by = g**Time:** $O(\log min(a,b))$

Warning: If a = b = 0 then there is infity solutions, but 0 is returned. Be careful about overflow.

```
#pragma once
#include "../Contest/template.cpp"
template <typename T>
tuple<T, T, T> extGcd(T a, T b) {
    if (!b) return {a, 1, 0};
    auto [d, x1, y1] = extGcd(b, a % b);
    T x = y1, y = x1 - y1 * (a / b);
    return {d, x, y};
}
```

10.9 FFT convolution and exponentiation

```
const ld PI = acos(-1);
/* change the ld to doulbe may increase
 * performance =D */
struct num {
    ld a{0.0}, b{0.0};
    num() {}
    num(ld na) : a{na} {}
```

```
num(ld na, ld nb) : a{na}, b{nb} {}
   const num operator+(const num &c) const { return num(a + c.a, b + c.b)
   const num operator-(const num &c) const { return num(a - c.a. b - c.b)
    const num operator*(const num &c) const {
        return num(a * c.a - b * c.b, a * c.b + b * c.a);
    const num operator/(const ll &c) const { return num(a / c, b / c); }
};
void fft(vector<num> &a, bool invert) {
   int n = len(a);
   for (int i = 1, j = 0; i < n; i++) {
        int bit = n \gg 1;
        for (; j & bit; bit >>= 1) j ^= bit;
        i ^= bit;
        if (i < j) swap(a[i], a[j]);</pre>
   for (int sz = 2; sz <= n; sz <<= 1) {
        ld ang = 2 * PI / sz * (invert ? -1 : 1);
        num wsz(cos(ang), sin(ang));
        for (int i = 0; i < n; i += sz) {
            num w(1);
            rep(j, 0, sz / 2)  {
                num u = a[i + j], v = a[i + j + sz / 2] * w;
                a[i + i] = u + v;
                a[i + j + sz / 2] = u - v;
                W = W * WSZ:
   if (invert)
        for (num \&x : a) x = x / n:
vi conv(vi const a, vi const b) {
   vector<num> fa(all(a));
   vector<num> fb(all(b));
   int n = 1;
    while (n < len(a) + len(b)) n <<= 1;
   fa.resize(n);
   fb.resize(n):
   fft(fa, false);
   fft(fb, false);
    rep(i, 0, n) fa[i] = fa[i] * fb[i];
   fft(fa, true):
   vi result(n);
   rep(i, 0, n) result[i] = round(fa[i].a);
   while (len(result) and result.back() == 0) result.pop back();
   /* Unconment this line if you want a boolean
    * convolution*/
   // for (auto &xi : result) xi = min(xi, 1ll);
   return result:
vll poly exp(vll &ps, int k) {
   vll ret(len(ps));
```

```
auto base = ps;
ret[0] = 1;
while (k) {
    if (k & 1) ret = conv(ret, base);
    k >>= 1;
    base = conv(base, base);
}
return ret;
}
```

10.10 Factorial Factorization

Description: Computes the factorization of N! in $\varphi(N) * \log N$ **Time**: $O(\varphi(N) \cdot \log N)$

```
ll E(ll n, ll p) {
    ll k = 0, b = p;
    while (b <= n) {
        k += n / b;
        b *= p;
    }
    return k;
}
map<ll, ll> factorial_factorization(ll n, const vll &primes) {
        map<ll, ll> fs;
        for (const auto &p : primes) {
             if (p > n) break;
             fs[p] = E(n, p);
        }
        return fs;
}
```

10.11 Factorization

Description: Computes the factorization of N.

Time: $O(\sqrt{n})$.

```
map<ll, ll> factorization(ll n) {
    map<ll, ll> ans;
    for (ll i = 2; i * i <= n; i++) {
        ll count = 0;
        for (; n % i == 0; count++, n /= i);
        if (count) ans[i] = count;
    }
    if (n > 1) ans[n]++;
    return ans;
}
```

10.12 Factorization (Pollard's Rho)

Description: Factorizes a number into its prime factors.

Time: $O(N^{(\frac{1}{4})} * \log(N))$.

```
ll mul(ll a, ll b, ll m) {
    ll ret = a * b - (ll)((ld)1 / m * a * b + 0.5) * m;
    return ret < 0 ? ret + m : ret;
ll pow(ll a, ll b, ll m) {
   ll\ ans = 1:
    for (; b > 0; b /= 2ll, a = mul(a, a, m)) {
        if (b % 2ll == 1) ans = mul(ans, a, m);
    return ans;
bool prime(ll n) {
   if (n < 2) return 0:
    if (n <= 3) return 1;
    if (n % 2 == 0) return 0;
    ll r = builtin ctzll(n - 1), d = n >> r;
    for (int a: {2, 325, 9375, 28178, 450775, 9780504, 795265022}) {
        ll x = pow(a, d, n);
        if (x == 1 \text{ or } x == n - 1 \text{ or a } % n == 0) continue;
        for (int j = 0; j < r - 1; j++) {
            x = mul(x, x, n);
            if (x == n - 1) break;
        if (x != n - 1) return 0;
    return 1;
ll rho(ll n) {
    if (n == 1 or prime(n)) return n;
    auto f = [n](ll x) \{ return mul(x, x, n) + 1; \};
    ll x = 0, y = 0, t = 30, prd = 2, x0 = 1, q;
    while (t % 40 != 0 or qcd(prd, n) == 1) {
        if (x == y) x = ++x0, y = f(x);
        q = mul(prd, abs(x - y), n);
        if (q != 0) prd = q;
        x = f(x), y = f(f(y)), t++;
    return gcd(prd, n);
vector<ll> fact(ll n) {
    if (n == 1) return {};
    if (prime(n)) return {n};
    ll d = rho(n);
    vector<ll> l = fact(d), r = fact(n / d);
    l.insert(l.end(), r.begin(), r.end());
    return l:
10.13 Fast Pow
```

Description: Computes $a^b \pmod{m}$ **Time:** $O(\log B)$.

```
ll fpow(ll a, ll b, ll m) {
```

```
ll ret = 1;
while (b) {
    if (b & 1) ret = (ret * a) % m;
    b >>= 1;
    a = (a * a) % m;
}
return ret;
}
```

10.14 Find linear recurrence (Berlekamp-Massey)

Description: Given the first N terms of a linear recurrence finds the smallest recurrence that matches the sequence.

Time: $O(N^2)$

Warning: Works faster if the *mod* is const but can be also be a parameter.

Absolute magic!

```
const ll mod = 998244353;
ll modpow(ll b, ll e) {
    ll ans = 1:
    for (; e; b = b * b % mod, e /= 2)
        if (e & 1) ans = ans * b % mod;
    return ans:
vl berlekampMassey(vll s) {
    int n = len(s), L = 0, m = 0;
    if (!n) return {};
    vll C(n), B(n), T;
    C[0] = B[0] = 1;
    ll b = 1;
    rep(i, 0, n) {
        ++m:
        ll d = s[i] \% mod;
        rep(j, 1, L + 1) d = (d + C[j] * s[i - j]) % mod;
        if (!d) continue;
        T = C;
        ll coef = d * modpow(b, mod - 2) % mod;
        rep(j, m, n) C[j] = (C[j] - coef * B[j - m]) % mod;
        if (2 * L > i) continue;
        L = i + 1 - L;
        B = T;
        b = d;
        m = 0:
    C.resize(L + 1):
    C.erase(C.begin());
    for (ll \&x : C) x = (mod - x) % mod;
    return C:
```

10.15 Find multiplicatinve inverse

```
ll inv(ll a, ll m) { return a > 1ll ? m - inv(m % a, a) * m / a : 1ll; }
```

10.16 Floor division

```
template <typename T1, typename T2>
constexpr typename std::common_type<T1, T2>::type floor_div(T1 x, T2 y) {
    assert(y != 0);
    if (y < 0) x = -x, y = -y;
    return x < 0 ? (x - y + 1) / y : x / y;
}</pre>
```

10.17 GCD

```
template <typename T>
T gcd(T a, T b) {
    return b ? gcd(b, a % b) : a;
}
```

10.18 Gauss XOR elimination / XOR-SAT

Description: Execute gaussian elimination with xor over the system Ax = b in. The add method must receive a bitset indicating which variables are present in the equation, and the solution of the equation.

Time: $O(\frac{nm^2}{64})$

```
const int MAXXI = 2009;
using Equation = bitset<MAXXI>;
struct GaussXor {
   vector<char> B;
   vector<Equation> A:
   void add(const Equation &ai, bool bi) {
        A.push back(ai):
        B.push back(bi):
   pair<bool, Equation> solution() {
        int cnt = 0. n = A.size():
        Equation vis:
        vis.set();
        Equation x;
        for (int j = MAXXI - 1, i; j >= 0; j--) {
            for (i = cnt; i < n; i++) {
                if (A[i][j]) break;
            if (i == n) continue;
            swap(A[i], A[cnt]), swap(B[i], B[cnt]);
            i = cnt++;
            vis[j] = 0;
            for (int k = 0; k < n; k++) {
                if (i == k || !A[k][j]) continue;
                A[k] ^= A[i]:
                B[k] ^= B[i];
            }
        x = vis:
```

```
for (int i = 0; i < n; i++) {
    int acum = 0;
    for (int j = 0; j < MAXXI; j++) {
        if (!A[i][j]) continue;
        if (!vis[j]) {
            vis[j] = 1;
            x[j] = acum ^ B[i];
        }
        acum ^= x[j];
    }
    if (acum != B[i]) return {false, Equation()};
}
return {true, x};
}</pre>
```

10.19 Guess K-th (Berlekamp-Massey)

```
/* Berlekamp-Massey algorithm
 * Given the first n terms of a linear recurrence relation, this algorithm
 * finds the shortest linear recurrence relation that generates the given
 * sequence.
  Note: mod needs to have inverse
 * Time complexity: O(n^2)
template <typename T>
vector<T> berlekamp massey(const vector<T> &s) {
    vector<T> cur, best;
    int lf, ld;
    for (int i = 0; i < (int)s.size(); i++) {
        T delta = 0;
        for (int j = 0; j < (int)cur.size(); j++)</pre>
            delta += s[i - j - 1] * cur[j];
        if (delta == s[i]) continue;
        if (cur.empty()) {
            cur.resize(i + 1);
            lf = i;
            ld = (int)(delta - s[i]).value();
            continue:
        T coef = -(s[i] - delta) / ld;
        vector<T> c(i - lf - 1);
        c.push back(coef);
        for (auto &x : best) c.push back(-x * coef);
        if (c.size() < cur.size()) c.resize(cur.size());</pre>
        for (int j = 0; j < (int)cur.size(); j++) c[j] += cur[j];</pre>
        if (i - lf + (int)best.size() >= (int)cur.size())
            best = cur, lf = i, ld = (int)(delta - s[i]).value();
        cur = c;
    return cur;
template <typename T>
T get kth(const vector<T> &rec, const vector<T> &dp, ll k) {
```

```
int n = (int)rec.size():
    assert(rec.size() <= dp.size());
    // use fft to speed up
    auto mul = [&](const vector<T> &a, const vector<T> &b) {
        vector<T> res(2 * n);
        for (int i = 0; i < n; i++)
            for (int j = 0; j < n; j++) res[i + j] += a[i] * b[j];
        for (int i = 2 * n - 1; i >= n; i--)
            for (int j = 1; j \le n; j++) res[i - j] += res[i] * rec[j - j]
    11;
        res.resize(n);
        return res:
    vector<T> a(n), x(n);
    x[0] = 1;
    if (n != 1)
        a[1] = 1;
        a[0] = rec[0]:
    while (k) {
        if (k \& 1) x = mul(x, a);
        a = mul(a, a):
        k >>= 1:
    T res = 0:
    for (int i = 0; i < n; i++) res += x[i] * dp[i];
    return res;
template <typename T>
T guess kth term(const vector<T> &s, ll k) {
    if (k < (int)s.size()) return s[k];</pre>
    auto coef = berlekamp massey(s);
    if (coef.empty()) return 0;
    return get kth(coef, s, k);
```

10.20 Integer partition

Description: Find the total of ways to partition a given number N in such way that none of the parts is greater than K.

```
Time: O(N \cdot min(N, K))
```

Memory: O(N)

Warning: Remember to memset everything to -1 before using it

```
const ll MOD = 10000000007;
const int MAXN(100);
ll memo[MAXN + 1];
ll dp(ll n, ll k = oo) {
    if (n == 0) return 1;
    ll &ans = memo[n];
    if (ans != -1) return ans;
    ans = 0;
    for (int i = 1; i <= min(n, k); i++) {
        ans = (ans + dp(n - i, k)) % MOD;
    }
    return ans;</pre>
```

10.21 LCM

}

```
ll gcd(ll a, ll b) { return b ? gcd(b, a % b) : a; }
ll lcm(ll a, ll b) { return a / gcd(a, b) * b; }
```

10.22 Linear Recurrence

Description: Find the n-th term of a linear recurrence, given the recurrence rec and the first K values of the recurrence, remember that first_k[i] is the value of f(i), considering 0-indexing.

Usage: Suppose you want the N-th term of Fibonacci the first k should be 1,1, and the rec should be 0.1,1.1.

Time: $O(K^3 \log N)$

```
template <typename T>
vector<vector<T>> prod(vector<vector<T>> &a, vector<vector<T>> &b,
                       const ll mod) -
    assert(a.back().size() == b.size());
    int n = a.size();
    int m = a.back().size();
    vector<vector<T>> c(n, vector<T>(m));
    for (int i = 0: i < n: i++) {
        for (int i = 0; i < m; i++) {
            for (int k = 0; k < m; k++) {
                c[i][j] = (c[i][j] + ((a[i][k] * b[k][j]) % mod)) % mod;
    return c;
template <tvpename T>
vector<vector<T>> fpow(vector<vector<T>> &xs, ll p, ll mod) {
    vector<vector<T>> ans(xs.size(), vector<T>(xs.size()));
    for (int i = 0; i < (int)xs.size(); i++) ans[i][i] = 1;</pre>
    for (auto b = xs; p; p >>= 1, b = prod(b, b, mod))
        if (p \& 1) ans = prod(ans, b, mod);
    return ans:
ll linear reg(vector<vector<ll>> rec, vector<ll> first k, ll n, const ll
   mod) {
    int k = first k.size();
    if (n <= k) return first k[n - 1];</pre>
    ll n2 = n - k:
    rec = fpow(rec, n2, mod);
    ll ret = 0:
    for (int i = 0; i < k; i++)
        ret = (ret + (rec.back()[i] * first_k[i]) % mod) % mod;
    return ret;
```

10.23 Linear diophantine equation (count)

Description:

```
Time: O(\log min(a,b))
#pragma once
#include "../Contest/template.cpp"
#include "./Extended Euclidian algorithm.cpp"
#include "./Linear diophantine equation (solve).cpp"
template <typename T>
T countSolutionsInRange(T a, T b, T c, T minX, T maxX, T minY, T maxY) {
   auto ss = [\&](T \&x, T \&y, T a, T b, T cnt) { x += cnt * b, y -= cnt *}
   assert(a and b):
   auto sol = diophantineEquationSolution(a, b, c);
   if (!sol) return 0;
   auto [x, y] = *sol;
   auto g = get<0>(extGcd(a, b));
   a /= q:
   b /= a:
   int signA = a > 0 ? +1 : -1;
   int signB = b > 0 ? +1 : -1;
   ss(x, y, a, b, (minX - x) / b);
   if (x < minX) ss(x, y, a, b, signB);
   if (x > maxX) return 0;
   int lx1 = x;
    ss(x, y, a, b, (maxX - x) / b);
   if (x > maxX) ss(x, y, a, b, -signB);
   int rx1 = x:
    ss(x, y, a, b, -(minY - y) / a);
   if (y < minY) ss(x, y, a, b, -signA);
   if (y > maxY) return 0;
   int lx2 = x:
   ss(x, y, a, b, -(maxY - y) / a);
   if (y > maxY) ss(x, y, a, b, signA);
   int rx2 = x;
   if (lx2 > rx2) swap(lx2, rx2);
   int lx = max(lx1, lx2);
   int rx = min(rx1, rx2);
   if (lx > rx) return 0;
   return (rx - lx) / abs(b) + 1;
```

10.24 Linear diophantine equation (solve)

Description: Finds a solution for ax + by = c, where a, b, c, are given and x and y unknown.

Time: $O(\log min(a, b))$

```
#pragma once
#include "../Contest/template.cpp"
#include "./Extended Euclidian algorithm.cpp"
```

```
template <typename T>
optional<pair<T, T>> diophantineEquationSolution(T a, T b, T c) {
    if (a == 0 \text{ and } b == 0) {
        if (c)
            return nullopt;
        else
            return pair<T, T>{(T)0, (T)0};
    }
    auto [g, x0, y0] = extGcd(a < 0 ? a * -1 : a, b < 0 ? b * -1 : b);
    if (c % g) return nullopt;
    x0 *= c / q, y0 *= c / q;
    if (a < 0) x0 = -x0:
    if (b < 0) v0 = -v0;
    pair<T, T> ret;
    ret.first = x0, ret.second = y0;
    return ret;
```

10.25 List N elements choose K

Description: Process every possible combination of K elements from N elements, thoose index marked as 1 in the index vector says which elements are choosed at that moment.

Time: $O(\binom{N}{K} \cdot O(process))$

```
void process(vi &index) {
    for (int i = 0; i < len(index); i++) {
        if (index[i]) cout << i << " \n"[i == len(index) - 1];
    }
}
void n_choose_k(int n, in k) {
    vi index(n);
    fill(index.end() - k, index.end(), 1);
    do {
        process(index);
    } while (next_permutation(all(index)));
}</pre>
```

10.26 List primes (Sieve of Eratosthenes)

```
const ll MAXN = 2e5;
vll list_primes(ll n = MAXN) {
    vll ps;
    bitset<MAXN + 1> sieve;
    sieve.set();
    sieve.reset(1);
    for (ll i = 2; i <= n; ++i) {
        if (sieve[i]) ps.push_back(i);
        for (ll j = i * 2; j <= n; j += i) {
            sieve.reset(j);
        }
    }
}</pre>
```

```
return ps;
}
```

10.27 Matrix exponentiation

```
const ll MOD = 1 '000' 000'007;
template <typename T>
vector<vector<T>> prod(vector<vector<T>> &a, vector<vector<T>> &b) {
   int n = len(a):
   vector<vector<T>> c(n, vector<T>(n));
   for (int i = 0; i < n; i++) {
        for (int j = 0; j < n; j++) {
            for (int k = 0; k < n; k++) {
                c[i][j] = (c[i][j] + ((a[i][k] * b[k][j]) % MOD)) % MOD;
       }
   }
    return c;
template <typename T>
vector<vector<T>> fpow(vector<vector<T>> &xs, ll p) {
   vector<vector<T>> ans(len(xs), vector<T>(len(xs)));
   for (int i = 0; i < len(xs); i++) ans[i][i] = 1;
   auto b = xs;
   while (p) {
        if (p \& 1) ans = prod(ans, b);
        p >>= 1;
        b = prod(b, b);
    return ans;
}
```

10.28 NTT integer convolution and exponentiation

Time:

- Convolution $O(N \cdot \log N)$,
- Exponentiation: $O(\log K \cdot N \cdot \log N)$

```
mint(ll v ) {
        if (v_ >= _mod or v_ <= -_mod) v_ %= _mod;
        if (v_{-} < 0) v_{-} += \underline{mod}; v_{-} += \underline{mod};
    m &operator+=(const m &a) {
        v += a.v;
        if (v \ge mod) v = mod;
        return *this:
    m &operator-=(const m &a) {
        \dot{v} -= a.v:
        if (v < 0) v += mod;
        return *this;
    m &operator*=(const m &a) {
        v = v * ll(a.v) % _mod;
        return *this;
    m &operator/=(const m &a) {
        v = v * inv(a.v) % mod;
        return *this;
    m operator-() { return m(-v); }
    m & operator^=(ll e) {
        if (e < 0) {
            v = inv(v):
             e = -e;
        v = expo(v, e);
        // possivel otimizacao:
        // cuidado com 0^0
        // v = \exp(v, e^{(p-1)});
        return *this;
    bool operator==(const m &a) { return v == a.v; }
    bool operator!=(const m &a) { return v != a.v; }
    friend istream &operator>>(istream &in, m &a) {
        ll val:
        in >> val;
        a = m(val);
        return in;
    friend ostream &operator<<(ostream &out, m a) { return out << a.v; }</pre>
    friend m operator+(m a, m b) { return a += b; }
    friend m operator-(m a, m b) { return a -= b;
    friend m operator*(m a, m b) { return a *= b; }
    friend m operator/(m a, m b) { return a /= b; }
    friend m operator^(m a, ll e) { return a ^= e; }
};
const ll MOD1 = 998244353;
const ll MOD2 = 754974721;
const ll MOD3 = 167772161;
template <int mod>
void ntt(vector<mint<_mod>> &a, bool rev) {
    int n = len(a);
```

```
auto b = a:
   assert(!(n \& (n - 1)));
   mint < mod > g = 1;
   while ((g \land (mod / 2)) == 1) g += 1;
   if (rev) q = 1 / q;
   for (int step = n / 2; step; step /= 2) {
        mint < mod > w = q ^ (mod / (n / step)), wn = 1;
        for (int i = 0; i < n / 2; i += step) {
            for (int j = 0; j < step; j++) {
                auto u = a[2 * i + j], v = wn * a[2 * i + j + step];
                b[i + j] = u + v;
                b[i + n / 2 + i] = u - v:
            \dot{w}n = wn * w;
        swap(a, b);
   if (rev) {
        auto n1 = mint< mod>(1) / n;
        for (auto &x : \overline{a}) x *= n1;
template <ll mod>
vector<mint< mod>> convolution(const vector<mint< mod>> &a,
                                const vector<mint< mod>> &b) {
   vector<mint< mod>> l(all(a)), r(all(b));
   int N = len(\overline{l}) + len(r) - 1, n = 1;
   while (n \le N) n *= 2;
   l.resize(n), r.resize(n);
   ntt(l, false), ntt(r, false);
   for (int i = 0; i < n; i++) l[i] *= r[i];
   ntt(l, true);
   l.resize(N);
   // Uncommnent for a boolean convolution :)
    for (auto& li : l) {
     li.v = min(li.v, 1ll);
    return l;
template <ll mod>
vector<mint< mod>> poly exp(vector<mint< mod>> &ps, int k) {
   vector<mint< mod>> ret(len(ps));
   auto base = \overline{ps}:
   ret[0] = 1;
   while (k) {
        if (k & 1) ret = convolution(ret, base);
        k >>= 1:
        base = convolution(base, base);
    return ret:
```

10.29 NTT integer convolution and exponentiation (2 mods) modules)

Description: Computes the convolution between the two polynomials and.

Time: $O(N \log N)$

Warning: This is pure magic!

```
template <int mod>
struct mint {
    ll expo(ll b, ll e) {
        ll ret = 1;
        while (e) {
            if (e % 2) ret = ret * b % mod;
            e /= 2, b = b * b % mod;
        return ret;
    ll inv(ll b) { return expo(b, mod - 2); }
    usina m = mint:
    ll v:
    mint() : v(0) {}
    mint(ll v ) {
        if (v_ >= _mod or v_ <= -_mod) v_ %= _mod;
        if (v_{-} < 0) v_{-} += _{mod};
    m &operator+=(const m &a) {
        \dot{v} += a.v:
        if (v >= mod) v -= mod;
        return *this;
    m &operator-=(const m &a) {
        \dot{v} -= a.v;
        if (v < 0) v += mod;
        return *this;
    m &operator*=(const m &a) {
        v = v * ll(a.v) % mod;
        return *this;
    m &operator/=(const m &a) {
        v = v * inv(a.v) % mod;
        return *this:
    m operator-() { return m(-v); }
    m &operator^=(ll e) {
        if (e < 0) {
            v = inv(v);
            e = -e;
        v = expo(v, e);
        // possivel otimizacao:
        // cuidado com 0^0
        // v = \exp(v, e^{(p-1)});
        return *this;
```

```
bool operator==(const m &a) { return v == a.v; }
    bool operator!=(const m &a) { return v != a.v; }
    friend istream &operator>>(istream &in, m &a) {
        ll val:
        in >> val;
        a = m(val);
        return in;
    friend ostream &operator<<(ostream &out, m a) { return out << a.v; }</pre>
    friend m operator+(m a, m b) { return a += b; }
    friend m operator-(m a, m b) { return a -= b; }
    friend m operator*(m a, m b) { return a *= b; }
    friend m operator/(m a, m b) { return a /= b; }
    friend m operator^(m a, ll e) { return a ^= e; }
};
const ll MOD1 = 998244353;
const ll MOD2 = 754974721;
const ll MOD3 = 167772161;
template <int mod>
void ntt(vector<mint<_mod>> &a, bool rev) {
    int n = len(a):
    auto b = a;
    assert(!(n \& (n - 1)));
    mint < mod > q = 1;
    while ((q ^ (mod / 2)) == 1) q += 1;
   if (rev) q = 1 / q;
    for (int step = n / 2; step; step /= 2) {
        mint < mod > w = g ^ (mod / (n / step)), wn = 1;
        for (int i = 0; i < n / 2; i += step) {
            for (int j = 0; j < step; j++) {
                auto u = a[2 * i + j], v = wn * a[2 * i + j + step];
                b[i + j] = u + v;
                b[i + n / 2 + i] = u - v;
            \dot{w}n = wn * w;
        swap(a, b);
    if (rev) {
        auto n1 = mint<_mod>(1) / n;
        for (auto \&x : \overline{a}) x *= n1;
tuple<ll, ll, ll> ext gcd(ll a, ll b) {
   if (!a) return {b, 0, 1};
    auto [g, x, y] = ext gcd(b % a, a);
    return \{g, y - b / a * x, x\};
template <typename T = ll>
struct crt {
   T a, m;
    crt() : a(0), m(1) {}
    crt(T a_, T m_) : a(a_), m(m_) {}
    crt operator*(crt C) {
```

```
auto [q, x, y] = ext qcd(m, C.m);
        if ((a - C.a) % q != 0) a = -1;
        if (a == -1 \text{ or } C.a == -1) \text{ return } crt(-1, 0);
        T lcm = m / g * C.m;
        T ans = a + (x * (C.a - a) / g % (C.m / g)) * m;
        return crt((ans % lcm + lcm) % lcm, lcm);
};
template <typename T = ll>
struct Congruence {
    T a, m;
};
template <typename T = ll>
T chinese remainder_theorem(const vector<Congruence<T>> &equations) {
    crt<T> ans:
    for (auto &[a_, m_] : equations) {
        ans = ans * crt<T>(a , m );
    return ans.a;
#define int long long
template <ll m1, ll m2>
vll merge_two_mods(const vector<mint<ml>> &a, const vector<mint<m2>> &b) {
    int n = len(a):
    vll ans(n);
    for (int i = 0; i < n; i++) {
        auto cur = crt<ll>():
        auto ai = a[i].v;
        auto bi = b[i].v;
        cur = cur * crt<ll>(ai, m1);
        cur = cur * crt < ll > (bi, m2);
        ans[i] = cur.a:
    }
    return ans;
vll convolution 2mods(const vll &a, const vll &b) {
    vector<mint<MOD1>> l(all(a)), r(all(b));
    int N = len(l) + len(r) - 1, n = 1;
    while (n \le N) n *= 2;
    l.resize(n), r.resize(n);
    ntt(l, false), ntt(r, false);
    for (int i = 0; i < n; i++) l[i] *= r[i];</pre>
    ntt(l, true);
    l.resize(N);
    vector<mint<MOD2>> l2(all(a)), r2(all(b));
    l2.resize(n), r2.resize(n);
    ntt(l2, false), ntt(r2, false);
    rep(i, 0, n) l2[i] *= r2[i];
    ntt(l2, true);
    l2.resize(N);
    return merge_two_mods(l, l2);
vll poly exp(const vll &xs, ll k) {
    vll ret(len(xs));
```

```
ret[0] = 1;
auto base = xs;
while (k) {
    if (k & 1) ret = convolution_2mods(ret, base);
    k >>= 1;
    base = convolution_2mods(base, base);
}
return ret;
```

10.30 Polynomial Taylor Shift

```
using C = complex<double>;
const ll mod = 998244353;
void fft(vector<C> &a) {
   int n = len(a), L = 31 - \_builtin\_clz(n);
   static vector<complex<long double>> R(2, 1);
    static vector<C> rt(2, 1);
   for (static int k = 2; k < n; k *= 2) {
        R.resize(n);
        rt.resize(n);
        auto x = polar(1.0L, acos(-1.0L) / k);
        for (int i = k; i < 2 * k; i++)
            rt[i] = R[i] = i \& 1 ? R[i / 2] * x : R[i / 2];
   vector<int> rev(n):
   for (int i = 0; i < n; i++) rev[i] = (rev[i / 2] | (i \& 1) << L) / 2;
   for (int i = 0; i < n; i++)
        if (i < rev[i]) swap(a[i], a[rev[i]]);</pre>
   for (int k = 1; k < n; k *= 2) {
        for (int i = 0; i < n; i += 2 * k)
            for (int j = 0; j < k; j++) {
                auto x = (double *)&rt[j + k], y = (double *)&a[i + j + k]
   ];
                C z(x[0] * y[0] - x[1] * y[1], x[0] * y[1] + x[1] * y[0]);
                a[i + j + k] = a[i + j] - z;
                a[i + j] += z;
            }
   }
vector<double> conv(const vector<double> &a, const vector<double> &b) {
   if (a.empty() || b.empty()) return {};
   vector<double> res(len(a) + len(b) - 1);
   int L = 32 - builtin clz(len(res)), n = 1 \ll L;
   vector<C> in(\overline{n}), out(\overline{n});
    copy(a.begin(), a.end(), begin(in));
   for (int i = 0; i < len(b); i++) in[i].imag(b[i]);
   fft(in):
    for (C &x : in) x *= x;
    for (int i = 0; i < n; i++) {
        out[i] = in[-i \& (n - 1)] - conj(in[i]);
   fft(out);
   for (int i = 0; i < len(res); i++) {
```

```
res[i] = imag(out[i]) / (4 * n);
    return res;
template <ll M>
vector<ll> convMod(const vector<ll> &a, const vector<ll> &b) {
    if (a.empty() || b.empty()) return {};
    vector<ll> res(len(a) + len(b) + 1);
    int B = 32 - builtin clz(len(res)), n = 1 \ll B, cut = int(sqrt(M));
    vector<C> L(n), R(n), outs(n), outl(n);
    for (int i = 0; i < len(a); i++) {
        L[i] = C((int)a[i] / cut, (int)a[i] % cut);
    for (int i = 0; i < len(b); i++) {
        R[i] = C((int)b[i] / cut, (int)b[i] % cut);
    fft(L), fft(R);
    for (int i = 0; i < n; i++) {
        int j = -i \& (n - 1);
        outl[j] = (L[i] + conj(L[j])) * R[i] / (2.0 * n);
        outs[j] = (L[i] - conj(L[j])) * R[i] / (2.0 * n) / 1i;
    fft(outl), fft(outs);
    for (int i = 0; i < len(res); i++) {
        ll av = ll(real(outl[i]) + .5), cv = ll(imag(outs[i]) + .5);
        ll bv = ll(imag(outl[i]) + .5) + ll(real(outs[i]) + .5);
        res[i] = ((av % M * cut + bv) % M * cut + cv) % M;
    }
    return res;
ll fexp(ll b, ll e) {
    ll res = 1;
    while (e > 0) {
        if (e & 1) res = res * b % mod;
        b = b * b % mod;
        e >>= 1:
    return res;
ll inv(ll n) { return fexp(n, mod - 2); }
vector<ll> shift(vector<ll> &a, ll v) {
    int n = len(a) - 1;
    vector<ll> f(n + 1), g(n + 1);
    vector<ll> i fact(n + 1);
    f[0] = a[0];
    q[n] = 1;
    i fact[0] = 1;
    l\bar{l} fact = 1, potk = 1;
    for (int i = 1; i < n + 1; i++) {
        fact = fact * i % mod;
        f[i] = fact * a[i] % mod;
        potk = (potk * v % mod + mod) % mod;
        q[n-i] = ((potk * inv(fact)) % mod + mod) % mod;
        i fact[i] = inv(fact);
    auto p = convMod<mod>(f, g);
```

```
vector<ll> res(n + 1);
for (int i = 0; i < n + 1; i++) {
    res[i] = (p[i + n] * i_fact[i] % mod + mod) % mod;
}
return res;
}</pre>
```

10.31 Polyominoes

Usage: buildPolyominoes(x) creates every polyomino until size x, and put it in polyominoes[x], access polyomino.v to find the vector of pairs representing the coordinates of each piece, considering that the polyomino was 'rooted' in coordinate (0,0).

Warning: note that when accessing polyominoes[x] only the first x coordinates are valid.

```
#include "../Contest/template.cpp"
const int MAXP = 10;
using pii = pair<int, int>;
// This implementation considers the rotations as
// distinct
                 0, 10, 10+9, 10+9+8...
int pos[11] = \{0, 10, 19, 27, 34, 40, 45, 49, 52, 54, 55\};
struct Polyominoes {
    pii v[MAXP];
   ll id:
   int n;
   Polyominoes() {
        n = 1;
        v[0] = \{0, 0\};
        normalize();
   pii &operator[](int i) { return v[i]; }
   bool add(int a, int b) {
        for (int i = 0; i < n; i++)
            if (v[i].first == a and v[i].second == b) return false;
        v[n++] = pii(a, b);
        normalize();
        return true:
   void normalize() {
        int mnx = 100, mny = 100;
        for (int i = 0; i < n; i++)
            mnx = min(mnx, v[i].first), mny = min(mny, v[i].second);
        id = 0:
        for (int i = 0; i < n; i++) {
            v[i].first -= mnx, v[i].second -= mny;
            id |= (1LL << (pos[v[i].first] + v[i].second));</pre>
   }
vector<Polyominoes> polyominoes[MAXP + 1];
void buildPolyominoes(int mxN = 10) {
   vector<pair<int, int>> dt({{1, 0}, {-1, 0}, {0, -1}, {0, 1}});
    for (int i = 0; i <= mxN; i++) polyominoes[i].clear();</pre>
   Polyominoes init;
   queue<Polyominoes> q;
```

```
unordered set<int64 t> used;
g.push(init);
used.insert(init.id);
while (!q.empty()) {
    Polyominoes u = q.front();
    q.pop();
    polyominoes[u.n].push back(u);
    if (u n == mxN) continue;
    for (int i = 0; i < u.n; i++) {
        for (auto [dx, dy] : dt) {
            Polyominoes to = u;
            bool ok = to.add(to[i].first + dx, to[i].second + dy);
            if (ok and !used.count(to.id)) {
                q.push(to);
                used.insert(to.id);
   }
```

11 Primitives

11.1 Bigint

```
const int maxn = 1e2 + 14, lg = 15;
const int base = 10000000000:
const int base digits = 9;
struct bigint ₹
    vi a;
    int sign;
    int size() {
        if (a.empty()) return 0;
        int ans = (a.size() - 1) * base digits;
        int ca = a.back();
        while (ca) ans++, ca /= 10;
        return ans;
    bigint operator^(const bigint &v) {
        bigint ans = 1, a = *this, b = v;
        while (!b.isZero()) {
            if (b \% 2) ans *= a;
            a *= a, b /= 2;
        return ans;
    string to_string() {
        stringstream ss;
        ss << *this;
        string s;
        ss >> s;
        return s;
    int sumof() {
        string s = to string();
```

```
int ans = 0:
    for (auto c : s) ans += c - '0';
    return ans;
/*</arpa>*/
bigint() : sign(1) {}
bigint(long long v) { *this = v; }
bigint(const string &s) { read(s); }
void operator=(const bigint &v) {
    sign = v.sign;
    a = v.a:
void operator=(long long v) {
    sign = 1;
    a.clear();
    if (v < 0) sign = -1, v = -v;
    for (; v > 0; v = v / base) a.push back(v % base);
bigint operator+(const bigint &v) const {
    if (sign == v.sign) {
        bigint res = v;
        for (int i = 0, carry = 0;
             i < (int)max(a.size(), v.a.size()) || carry; ++i) {</pre>
            if (i == (int)res.a.size()) res.a.push back(0);
            res.a[i] += carry + (i < (int)a.size() ? a[i] : 0);
            carry = res.a[i] >= base;
            if (carry) res.a[i] -= base;
        return res;
    return *this - (-v);
bigint operator-(const bigint &v) const {
    if (sign == v.sign) {
        if (abs() >= v.abs()) {
            bigint res = *this;
            for (int i = 0, carry = 0; i < (int)v.a.size() || carry;
++i) {
                res.a[i] \rightarrow carry + (i < (int)v.a.size() ? v.a[i] : 0)
                carrv = res.a[i] < 0:
                if (carry) res.a[i] += base;
            res.trim():
            return res;
        return -(v - *this);
    return *this + (-v);
void operator*=(int v) {
    if (v < 0) sign = -sign, v = -v;
    for (int i = 0, carry = 0; i < (int)a.size() || carry; ++i) {
        if (i == (int)a.size()) a.push back(0);
```

```
long long cur = a[i] * (long long)v + carry;
        carry = (int)(cur / base);
        a[i] = (int)(cur % base);
        // asm("divl %%ecx" : "=a"(carry),
        // "=d"(a[i]) : "A"(cur), "c"(base));
    trim();
bigint operator*(int v) const {
    bigint res = *this:
    res *= v;
    return res;
void operator*=(long long v) {
    if (v < 0) sign = -sign, v = -v;
    if (v > base) {
        *this = *this * (v / base) * base + *this * (v % base);
    for (int i = 0, carry = 0; i < (int)a.size() || carry; ++i) {</pre>
        if (i == (int)a.size()) a.push back(0);
        long long cur = a[i] * (long long)v + carry;
        carry = (int)(cur / base);
        a[i] = (int)(cur % base);
        // asm("divl %%ecx" : "=a"(carry),
        // "=d"(a[i]) : "A"(cur), "c"(base));
    trim():
bigint operator*(long long v) const {
    bigint res = *this;
    res *= v;
    return res:
friend pair<br/>
bigint, bigint> divmod(const bigint &a1, const bigint &b1)
    int norm = base / (b1.a.back() + 1);
    bigint a = a1.abs() * norm;
    bigint b = b1.abs() * norm;
    biaint a, r:
    q.a.resize(a.a.size());
    for (int i = a.a.size() - 1; i \ge 0; i--) {
        r *= base;
        r += a.a[i];
        int s1 = r.a.size() <= b.a.size() ? 0 : r.a[b.a.size()];</pre>
        int s2 = r.a.size() \le b.a.size() - 1 ? 0 : r.a[b.a.size() -
1];
        int d = ((long long)base * s1 + s2) / b.a.back();
        r -= b * d:
        while (r < 0) r += b, --d;
        q.a[i] = d;
    q.sign = a1.sign * b1.sign;
    r.sign = a1.sign;
    q.trim();
```

```
r.trim();
    return make pair(q, r / norm);
bigint operator/(const bigint &v) const { return divmod(*this, v).
bigint operator%(const bigint &v) const { return divmod(*this, v).
second; }
void operator/=(int v) {
    if (v < 0) sign = -sign, v = -v;
    for (int i = (int)a.size() - 1, rem = 0; i >= 0; --i) {
        long long cur = a[i] + rem * (long long)base;
        a[i] = (int)(cur / v);
        rem = (int)(cur % v);
    trim();
bigint operator/(int v) const {
    bigint res = *this;
    res /= v;
    return res;
int operator%(int v) const {
    if (v < 0) v = -v;
    int m = 0:
    for (int i = a.size() - 1; i >= 0; --i)
        m = (a[i] + m * (long long)base) % v;
    return m * sign;
void operator+=(const bigint &v) { *this = *this + v; }
void operator==(const bigint &v) { *this = *this - v; }
void operator*=(const bigint &v) { *this = *this * v; }
void operator/=(const bigint &v) { *this = *this / v; }
bool operator<(const bigint &v) const {</pre>
    if (sign != v.sign) return sign < v.sign;</pre>
    if (a.size() != v.a.size())
        return a.size() * sign < v.a.size() * v.sign;</pre>
    for (int i = a.size() - 1; i >= 0; i--)
        if (a[i] != v.a[i]) return a[i] * sign < v.a[i] * sign;</pre>
    return false:
bool operator>(const bigint &v) const { return v < *this; }</pre>
bool operator<=(const bigint &v) const { return !(v < *this); }</pre>
bool operator>=(const bigint &v) const { return !(*this < v); }</pre>
bool operator==(const bigint &v) const {
    return !(*this < v) && !(v < *this);
bool operator!=(const bigint &v) const { return *this < v || v < *this
; }
void trim() {
    while (!a.empty() && !a.back()) a.pop back();
    if (a.empty()) sign = 1;
bool isZero() const { return a.empty() || (a.size() == 1 \&\& !a[0]); }
```

```
bigint operator-() const {
    bigint res = *this:
    res.sign = -sign;
    return res;
bigint abs() const {
    bigint res = *this;
    res.sian *= res.sian:
    return res;
long longValue() const {
    long long res = 0;
    for (int i = a.size() - 1; i \ge 0; i--) res = res * base + a[i];
    return res * sign:
friend bigint gcd(const bigint &a, const bigint &b) {
    return b.isZero() ? a : gcd(b, a % b);
friend bigint lcm(const bigint &a, const bigint &b) {
    return a / \gcd(a, b) * b;
void read(const string &s) {
    sign = 1;
    a.clear();
    int pos = 0;
    while (pos < (int)s.size() && (s[pos] == '-' || s[pos] == '+')) {
        if (s[pos] == '-') sign = -sign:
        ++pos;
    for (int i = s.size() - 1; i \ge pos; i -= base digits) {
        int x = 0:
        for (int j = max(pos, i - base digits + 1); j <= i; j++)
            x = x * 10 + s[i] - '0';
        a.push back(x);
    trim();
friend istream &operator>>(istream &stream, bigint &v) {
    string s:
    stream >> s:
    v.read(s);
    return stream;
friend ostream &operator<<(ostream &stream, const bigint &v) {</pre>
    if (v.sign == -1) stream << '-';
    stream << (v.a.empty() ? 0 : v.a.back());</pre>
    for (int i = (int)v.a.size() - 2; i >= 0; --i)
        stream << setw(base digits) << setfill('0') << v.a[i];</pre>
    return stream:
static vector<int> convert_base(const vector<int> &a, int old digits,
                                 int new digits) {
    vector<long long> p(max(old digits, new digits) + 1);
    p[0] = 1;
```

```
for (int i = 1; i < (int)p.size(); i++)p[i] = p[i-1] * 10;
    vector<int> res:
    long long cur = 0;
    int cur digits = 0;
    for (int i = 0; i < (int)a.size(); i++) {
        cur += a[i] * p[cur digits];
        cur digits += old digits;
        while (cur digits >= new digits) {
            res.push back(int(cur % p[new digits]));
            cur /= p[new digits]:
            cur digits -= new_digits;
        }
    res.push back((int)cur);
    while (!res.empty() && !res.back()) res.pop back();
    return res;
typedef vector<long long> vll;
static vll karatsubaMultiply(const vll &a, const vll &b) {
    int n = a.size();
    vll res(n + n):
    if (n <= 32) {
        for (int i = 0; i < n; i++)
            for (int i = 0; i < n; i++) res[i + i] += a[i] * b[i];
        return res:
    }
    int k = n \gg 1;
    vll a1(a.begin(), a.begin() + k);
    vll a2(a.begin() + k, a.end());
    vll b1(b.begin(), b.begin() + k);
    vll b2(b.beain() + k.b.end()):
    vll alb1 = karatsubaMultiply(a1, b1);
    vll a2b2 = karatsubaMultiply(a2, b2);
    for (int i = 0; i < k; i++) a2[i] += a1[i];
    for (int i = 0; i < k; i++) b2[i] += b1[i];
    vll r = karatsubaMultiply(a2, b2);
    for (int i = 0; i < (int)alb1.size(); i++) r[i] -= alb1[i];
    for (int i = 0; i < (int)a2b2.size(); i++) r[i] -= a2b2[i];
    for (int i = 0; i < (int)r.size(); i++) res[i + k] += r[i];
    for (int i = 0; i < (int)alb1.size(); i++) res[i] += alb1[i];
    for (int i = 0; i < (int)a2b2.size(); i++) res[i + n] += a2b2[i];
    return res;
bigint operator*(const bigint &v) const {
    vector<int> a6 = convert base(this->a, base digits, 6);
    vector<int> b6 = convert base(v.a, base digits, 6);
    vll a(a6.begin(), a6.end\overline{()});
    vll b(b6.begin(), b6.end());
    while (a.size() < b.size()) a.push_back(0);</pre>
    while (b.size() < a.size()) b.push back(0);</pre>
    while (a.size() \& (a.size() - 1)) \overline{a.push back(0)}, b.push back(0);
    vll c = karatsubaMultiply(a, b);
    bigint res;
```

```
res.sign = sign * v.sign;
for (int i = 0, carry = 0; i < (int)c.size(); i++) {
    long long cur = c[i] + carry;
    res.a.push_back((int)(cur % 1000000));
    carry = (int)(cur / 1000000);
}
res.a = convert_base(res.a, 6, base_digits);
res.trim();
return res;
}
};</pre>
```

11.2 Integer Mod

```
#include "../Contest/template.cpp"
template <ll m>
struct mod int {
    ll x;
    mod int(ll v = 0)  {
        x = v % m;
        if (x < 0) v += m;
    mod int &operator+=(mod int const &b) {
        x += b.x;
        if (x \ge m) x = m:
        return *this;
    mod int &operator=(mod int const &b) {
        x \rightarrow b.x:
        if (x < 0) x += m:
        return *this;
    mod int &operator*=(mod int const &b) {
        x = (ll)x * b.x % m;
        return *this;
    friend mod int mpow(mod int a, ll e) {
        mod int res = 1:
        while (e) {
            if (e & 1) res *= a;
            a ∗= a;
            e >>= 1:
        return res;
    friend mod int inverse(mod int a) { return mpow(a, m - 2); }
    mod int &operator/=(mod int const &b) { return *this *= inverse(b); }
    friend mod int operator+(mod int a, mod int const b) { return a += b;
    mod_int operator++(int) { return this->x = (this->x + 1) % m; }
    mod int operator++() { return this->x = (this->x + 1) % m; }
    friend mod int operator-(mod int a, mod int const b) { return a -= b;
```

```
friend mod_int operator-(mod_int const a) { return 0 - a; }
mod_int operator--(int) { return this->x = (this->x - 1 + m) % m; }
mod_int operator--() { return this->x = (this->x - 1 + m) % m; }
friend mod_int operator*(mod_int a, mod_int const b) { return a *= b; }
friend mod_int operator/(mod_int a, mod_int const b) { return a /= b; }
friend ostream & operator<<(ostream & os, mod_int const & a) {
    return os << a.x; }
friend bool operator==(mod_int const & a, mod_int const & b) {
    return a.x == b.x; }
friend bool operator!=(mod_int const & a, mod_int const & b) {
    return a.x != b.x; }
};</pre>
```

11.3 Integer Mod (complete)

```
#include "../Contest/template.cpp"
template <ll Mod>
struct modint {
    static constexpr ll mod = Mod;
   ll v;
   modint() : v(0) \{ \}
   template <ll Mod2>
   modint(const modint<Mod2> &x) : v(x.value()) {}
   modint(ll x) : v(safe_mod(x)) {}
   ll &value() { return \overline{v}; }
   const ll &value() const { return v; }
   static ll safe mod(ll x) {
        return x \ge 0 ? x \% \mod : ((x \% \mod) + \mod) \% \mod;
   template <typename T>
   explicit operator T() const {
        return (T)v;
    bool operator==(const modint rhs) const noexcept { return v == rhs.v;
    bool operator!=(const modint rhs) const noexcept { return v != rhs.v;
   bool operator<(const modint rhs) const noexcept { return v < rhs.v; }</pre>
   bool operator<=(const modint rhs) const noexcept { return v <= rhs.v;</pre>
   bool operator>(const modint rhs) const noexcept { return v > rhs.v; }
   bool operator>=(const modint rhs) const noexcept { return v >= rhs.v;
   modint operator++(int) {
        modint res = *this:
        *this += 1;
        return res;
   modint operator--(int) {
```

```
modint res = *this:
    *this -= 1:
    return res;
modint &operator++() { return *this += 1; }
modint &operator--() { return *this -= 1; }
modint operator+() const { return modint(*this); }
modint operator-() const { return mod - modint(*this): }
friend modint operator+(const modint lhs, const modint rhs) noexcept {
    return modint(lhs) += rhs;
friend modint operator-(const modint lhs, const modint rhs) noexcept {
    return modint(lhs) -= rhs:
friend modint operator*(const modint lhs, const modint rhs) noexcept {
    return modint(lhs) *= rhs;
friend modint operator/(const modint lhs, const modint rhs) noexcept {
    return modint(lhs) /= rhs;
modint &operator+=(const modint rhs) {
    v += rhs.v:
    if (v \ge mod) v = mod;
    return *this;
modint & operator = (const modint rhs) {
    if (v < rhs.v) v += mod:
    v = rhs.v:
    return *this;
modint & operator *= (const modint rhs) {
    v = v * rhs.v % mod;
    return *this;
modint &operator/=(modint rhs) { return *this *= rhs.inv(); }
modint pow(ll p) const {
    static assert(mod < static cast<ll>(1) << 32,
                  "Modulus must be less than 2**32");
    modint res = 1, a = *this;
    while (p) {
        if (p \& 1) res *= a;
        a *= a;
        p >>= 1;
    return res;
modint inv() const { return pow(mod - 2); }
modint sqrt() const {
    modint b = 1;
    while (b.pow((mod - 1) >> 1) == 1) b += 1;
    ll m = mod - 1. e = 0:
    while (\sim m \& 1) m >>= 1. e++:
    auto x = pow((m - 1) >> 1);
    auto v = *this * x * x:
    x *= *this;
    auto z = b.pow(m);
    while (y != 1) {
```

```
ll i = 0:
             for (modint t = y; t != 1; t *= t, ++j);
            z.pow(111 << (e - j - 1));
            X *= Z;
Z *= Z;
            V *= Z;
            e = i;
        return x;
    friend ostream &operator<<(ostream &s, const modint &x) {</pre>
        s << x.value():
        return s:
    friend istream &operator>>(istream &s, modint &x) {
        ll value:
        s >> value:
        x = \{value\};
        return s;
};
```

11.4 Matrix

```
template <tvpename T>
struct Matrix {
   vector<vector<T>> d;
   Matrix() : Matrix(0) {}
   Matrix(int n) : Matrix(n, n) {}
   Matrix(int n, int m) : Matrix(vector<vector<T>>(n, vector<T>(m))) {}
   Matrix(const vector<vector<T>> &v) : d(v) {}
   constexpr int n() const { return (int)d.size(); }
   constexpr int m() const { return n() ? (int)d[0].size() : 0; }
   void rotate() { *this = rotated(); }
   Matrix<T> rotated() const {
       Matrix<T> res(m(), n());
       for (int i = 0; i < m(); i++) {
            for (int j = 0; j < n(); j++) {
                res[i][i] = d[n() - i - 1][i];
        return res;
   Matrix<T> pow(int power) const {
       assert(n() == m()):
       auto res = Matrix<T>::identity(n());
       auto b = *this;
       while (power) {
            if (power & 1) res *= b;
            b *= b:
            power >>= 1;
        return res:
```

```
Matrix<T> submatrix(int start i, int start j, int rows = INT MAX,
                    int cols = INT MAX) const {
    rows = min(rows, n() - start i);
    cols = min(cols, m() - start i):
    if (rows \leq 0 or cols \leq 0) return \{\}:
    Matrix<T> res(rows, cols);
    for (int i = 0: i < rows: i++)
        for (int j = 0; j < cols; j++)
            res[i][i] = d[i + start i][j + start_j];
    return res;
}
Matrix<T> translated(int x, int y) const {
    Matrix<T> res(n(), m());
    for (int i = 0; i < n(); i++) {
        for (int j = 0; j < m(); j++) {
            if (i + x < 0 \text{ or } i + x >= n() \text{ or } j + y < 0 \text{ or } j + y >= m()
                 continue:
            res[i + x][j + y] = d[i][j];
    return res;
static Matrix<T> identity(int n) {
    Matrix<T> res(n):
    for (int i = 0; i < n; i++) res[i][i] = 1;
    return res;
vector<T> &operator[](int i) { return d[i]; }
const vector<T> &operator[](int i) const { return d[i]; }
Matrix<T> & operator += (T value) {
    for (auto &row : d) {
        for (auto &x : row) x += value;
    return *this:
Matrix<T> operator+(T value) const {
    auto res = *this;
    for (auto &row : res) {
        for (auto &x : row) x = x + value;
    return res;
Matrix<T> &operator==(T value) {
    for (auto &row : d) {
        for (auto &x : row) x -= value;
    return *this:
Matrix<T> operator-(T value) const {
    auto res = *this;
    for (auto &row : res) {
        for (auto &x : row) x = x - value:
    return res:
```

```
Matrix<T> &operator*=(T value) {
    for (auto &row : d) {
        for (auto &x : row) x *= value;
    return *this:
Matrix<T> operator*(T value) const {
    auto res = *this:
    for (auto &row : res) {
        for (auto &x : row) x = x * value;
    return res;
Matrix<T> &operator/=(T value) {
    for (auto &row : d) {
        for (auto &x : row) x /= value;
    return *this;
Matrix<T> operator/(T value) const {
    auto res = *this;
    for (auto &row : res) {
        for (auto \&x : row) x = x / value;
    return res;
Matrix<T> & operator += (const Matrix<T> &o) {
    assert(n() == o.n() and m() == o.m());
    for (int i = 0; i < n(); i++) {
        for (int i = 0; i < m(); i++) {
            d[i][i] += o[i][i];
    return *this:
Matrix<T> operator+(const Matrix<T> &o) const {
    assert(n() == o.n() and m() == o.m());
    auto res = *this;
    for (int i = 0; i < n(); i++) {
        for (int j = 0; j < m(); j++) {
            res[i][j] = res[i][j] + o[i][j];
    return res;
Matrix<T> & Operator = (const Matrix<T> & O) {
    assert(n() == o.n() and m() == o.m()):
    for (int i = 0; i < n(); i++) {
        for (int j = 0; j < m(); j++) {
            d[i][i] -= o[i][i];
    return *this;
Matrix<T> operator-(const Matrix<T> &o) const {
    assert(n() == o.n() and m() == o.m());
    auto res = *this;
```

```
for (int i = 0; i < n(); i++) {
            for (int i = 0; i < m(); i++) {
                res[i][j] = res[i][j] - o[i][j];
        return res;
    Matrix<T> & operator*=(const Matrix<T> &o) {
        *this = *this * 0:
        return *this:
    Matrix<T> operator*(const Matrix<T> &o) const {
        assert(m() == o.n()):
        Matrix<T> res(n(), o.m());
        for (int i = 0; i < res.n(); i++) {
            for (int j = 0; j < res.m(); j++) {
                auto &x = res[i][i];
                for (int k = 0; k < m(); k++) {
                    x += (d[i][k] * o[k][i]):
            }
        return res:
    friend istream &operator>>(istream &is. Matrix<T> &mat) {
        for (auto &row : mat)
            for (auto &x : row) is >> x;
        return is:
    friend ostream &operator<<(ostream &os, const Matrix<T> &mat) {
        bool frow = 1:
        for (auto &row : mat) {
            if (not frow) os << '\n';</pre>
            bool first = 1:
            for (auto &x : row) {
                if (not first) os << ' ':
                os << x:
                first = 0:
            frow = 0:
        return os;
    auto begin() { return d.begin(); }
    auto end() { return d.end(); }
    auto rbegin() { return d.rbegin(); }
    auto rend() { return d.rend(); }
    auto begin() const { return d.begin(); }
    auto end() const { return d.end(); }
    auto rbegin() const { return d.rbegin(); }
    auto rend() const { return d.rend(); }
};
```

12 Problems

12.1 2081 - Fixed-Lenght Paths II

```
#include <bits/stdc++.h>
using namespace std;
const int MAXN = 2'00'000;
int N, K1, K2;
vector<int> ADJ[MAXN];
int64 t ans = 0;
int sz[MAXN], removed[MAXN];
void calcSize(int u, int p = -1) {
    sz[u] = 1:
    for (int v : ADJ[u]) {
        if (v != p and !removed[v]) {
            calcSize(v, u);
            sz[u] += sz[v];
   }
}
int findCentroid(int u, int mxSz, int p = -1) {
    for (int v : ADJ[u]) {
        if (!removed[v] \text{ and } v != p \text{ and } sz[v] * 2 >= mxSz)
            return findCentroid(v, mxSz, u);
    return u;
int64 t cnt[MAXN], totCnt[MAXN], initialSum;
int mxD;
void dfs(int u, int p, int d) {
   if (d > K2) return;
    cnt[d]++:
    mxD = max(mxD, d);
    if (K1 - 1 \le d \text{ and } d \le K2 - 1) initialSum++;
    for (int v : ADJ[u]) {
        if (v != p and !removed[v]) {
            dfs(v, u, d + 1);
void solve(int curRoot) {
    calcSize(curRoot);
    int centroid = findCentroid(curRoot, sz[curRoot]);
    removed[centroid] = true;
    int totMxD = 0:
    initialSum = (K1 == 1);
    // cerr << "centroid: " << centroid << '\n';</pre>
    for (int v : ADJ[centroid]) {
        if (!removed[v]) {
            // cerr << "v: " << v << '\n';
            m \times D = 0:
            int64 t curSum = initialSum;
            dfs(v. centroid, 1):
            totMxD = max(totMxD, mxD);
```

```
for (int d = 1; d <= mxD; d++) {</pre>
                 // cerr << "d : " << d << " curSum: " << curSum << '\n';
                 ans += (curSum * cnt[d]);
                 int pl = max(0, K1 - d) - 1;
                 if (pl >= 0) curSum += totCnt[pl];
                 int pr = K2 - d:
                 curSum -= totCnt[pr];
            for (int d = 1; d <= mxD; d++) totCnt[d] += cnt[d];</pre>
            fill(\&cnt[1], \&cnt[1] + mxD + 1, 0);
    // cerr << "centroid: " << centroid
    //<< " ans: " << ans << '\n';
    for (int d = 1; d \le totMxD; d++) totCnt[d] = 0;
    for (int v : ADJ[centroid])
        if (!removed[v]) solve(v);
int32_t main() {
    ios base::sync with stdio(!cin.tie(0));
    tot\overline{C}nt[0] = 1;
    cin >> N >> K1 >> K2;
    for (int i = 0; i < N - 1; i++) {
        int u. v:
        cin >> u >> v:
        u--, v--;
        ADJ[u] emplace back(v);
        ADJ[v].emplace_back(u):
    solve(0):
    cout << ans << '\n':
// AC, centroid decomposition
```

12.2 Fixed lenght pahts I

```
#include <bits/stdc++.h>
using namespace std;

const int MAXN = 2'00'000;
int N, K;
vector<int> ADJ[MAXN];
int64_t ans;
bool removed[MAXN];
int cnt[MAXN];
int sz[MAXN];
void calcSize(int u, int p = -1) {
    sz[u] = 1;
    for (int v : ADJ[u]) {
        if (v != p and !removed[v]) {
            calcSize(v, u);
            sz[u] += sz[v];
```

```
}
}
int getCentroid(int mxSz, int u, int p = -1) {
    for (int v : ADJ[u]) {
        if (v != p and !removed[v] and sz[v] >= mxSz)
            return getCentroid(mxSz, v, u);
    return u;
}
int mxd:
void dfs(int u, int p, bool upd, int d = 1) {
    if (d > K) return:
    mxd = max(mxd, d);
    upd ? cnt[d]++ : ans += cnt[K - d];
    for (int v : ADJ[u]) {
        if (v != p \text{ and } !removed[v]) dfs(v, u, upd, d + 1);
}
void solve(int u) {
    calcSize(u);
    int c = getCentroid(sz[u] >> 1, u);
    removed[c] = true;
    mxd = 0:
    cnt[0] = 1;
    for (int v : ADJ[c]) {
        if (!removed[v]) {
            dfs(v, c, false);
            dfs(v, c, true);
    for (int i = 0; i \le mxd; i++) cnt[i] = 0;
    for (int v : ADJ[c]) {
        if (!removed[v]) solve(v);
int32 t main() {
    ios base::sync with stdio(0);
    cin.tie(0);
    cin >> N >> K;
    for (int i = 0; i < N - 1; i++) {
        int u, v;
        cin >> u >> v;
        u--, v--;
        ADJ[u].emplace back(v);
        ADJ[v].emplace back(u);
    }
    solve(0);
    cout << ans << '\n';
    return 0;
```

12.3 Fixed length paths II

```
#include <bits/stdc++.h>
using namespace std;
const int MAXN = 2'00'000;
int N, K1, K2;
vector<int> ADJ[MAXN];
int64 t ans = 0;
int sz[MAXN], removed[MAXN];
void calcSize(int u, int p = -1) {
    sz[u] = 1:
    for (int v : ADJ[u]) {
        if (v != p and !removed[v]) {
            calcSize(v, u);
            sz[u] += sz[v];
int findCentroid(int u, int mxSz, int p = -1) {
    for (int v : ADJ[u]) {
        if (!removed[v] \text{ and } v != p \text{ and } sz[v] * 2 >= mxSz)
            return findCentroid(v, mxSz, u);
    return u;
int64 t cnt[MAXN], totCnt[MAXN], initialSum;
int mxD:
void dfs(int u, int p, int d) {
    if (d > K2) return;
    cnt[d]++:
    mxD = max(mxD, d):
    if (K1 - 1 \le d \text{ and } d \le K2 - 1) initialSum++;
    for (int v : ADJ[u]) {
        if (v != p and !removed[v]) {
            dfs(v, u, d + 1);
void solve(int curRoot) {
    calcSize(curRoot);
    int centroid = findCentroid(curRoot, sz[curRoot]);
    removed[centroid] = true;
    int totMxD = 0:
    initialSum = (K1 == 1);
    // cerr << "centroid: " << centroid << '\n';</pre>
    for (int v : ADJ[centroid]) {
        if (!removed[v]) {
            // cerr << "v: " << v << '\n':
            mxD = 0:
            int64_t curSum = initialSum;
            dfs(v, centroid, 1);
            totMxD = max(totMxD, mxD);
            for (int d = 1; d <= mxD; d++) {</pre>
                 // cerr << "d : " << d << " curSum: " << curSum << '\n';
```

```
ans += (curSum * cnt[d]);
                int pl = \max(0, K1 - d) - 1;
                if (pl >= 0) curSum += totCnt[pl];
                int pr = K2 - d:
                curSum -= totCnt[pr];
            for (int d = 1: d \le mxD: d++) totCnt[d] += cnt[d]:
            fill(\&cnt[1], \&cnt[1] + mxD + 1, 0):
    // cerr << "centroid: " << centroid
    //<< " ans: " << ans << '\n';
    for (int d = 1; d \le totMxD; d++) totCnt[d] = 0;
    for (int v : ADJ[centroid])
        if (!removed[v]) solve(v);
int32 t main() {
    ios base::sync with stdio(!cin.tie(0));
    totCnt[0] = 1:
    cin >> N >> K1 >> K2:
    for (int i = 0; i < N - 1; i++) {
        int u, v;
        cin >> u >> v;
        u--, v--;
        ADJ[u].emplace back(v);
        ADJ[v].emplace back(u);
    solve(0);
    cout << ans << '\n';
// AC, centroid decomposition
```

13 Strings

13.1 Z-Function

13.1.1 Pattern Matching

Description: The Pattern Matching algorithm uses the Z-function to efficiently find all occurrences of a pattern p in a text S. The Z-function is first computed for the concatenation of the pattern and text (p + \$ + S), where \$ is a unique separator character that does not appear in either p or S. The Z-array allows us to quickly determine where in the text the pattern occurs by comparing segments of the text with the pattern.

Usage: The function $pattern_matching(s, p)$ takes two arguments: s, the text, and p, the pattern to search for. It returns a vector of integers, each of which represents the starting index of a match of the pattern p in the text S.

```
string text = "abacabadabacaba";
string pattern = "abac";
vector<int> result = pattern_matching(text, pattern);
// result = [0, 8]
string text_v = "ababab";
```

```
string pattern_v = "ab";
vector<int> result_v = pattern_matching(text_v, pattern_v);
// result v = [0, 2, 4]
```

Time: O(n+m), where n is the length of the text S and m is the length of the pattern p. **Memory:** O(n+m)

Warning: The pattern must be non-empty, and the text must also be non-empty. Ensure that the separator character \$ is not present in the pattern or text.

```
#pragma once
#include "../../Contest/template.cpp"
#include "../z-function-build.cpp"

template <typename Seq>
vector<int> z_function_pattern_matching(const Seq& s, const Seq& p) {
    if (s.empty() or p.empty()) return {};
    Seq concatenated = p + "$" + s;
    vector<int> z = z_function_build(concatenated);
    vector<int> occurrences;
    int m = p.size();
    for (int i = m + 1; i < z.size(); ++i) {
        if (z[i] == m) occurrences.push_back(i - m - 1);
    }
    return occurrences;
}</pre>
```

13.1.2 Z-function building

Description: The Z-function is an algorithm used to compute the Z-array of a given string. For a string s, Z[i] represents the length of the longest common prefix between the string s and the suffix of s starting from the index i.

Usage: The function z_function_build(s) takes a single argument s, which is a string (or any container-like structure), and returns a vector of integers representing the Z-function of the input.

```
string s = "abacaba";
vector<int> result = z_function_build(s);
// result = [0, 0, 1, 0, 3, 0, 1]
vector<int> v = {1, 2, 3, 1, 2, 3};
vector<int> result_v = z_function_build(v);
// result_v = [0, 1, 2, 3, 0, 1]
```

Time: O(n)Memory: O(n)

Warning: By definition Z[0] = 0, remember to treat it appart.

```
#pragma once
#include "../../Contest/template.cpp"
template <typename Seq>
vector<int> z_function_build(const Seq& s) {
   int n = int(s.size());
   vector<int> z(n);
   for (int i = 1, l = 0, r = 0; i < n; ++i) {</pre>
```

```
if (i <= r) z[i] = min(r - i + 1, z[i - l]);
while (i + z[i] < n && s[z[i]] == s[i + z[i]]) ++z[i];
if (i + z[i] - 1 > r) l = i, r = i + z[i] - 1;
}
return z;
}
```

13.2 Count distinct anagrams

```
const ll\ MOD = 1e9 + 7:
const int maxn = 1e6;
vll fs(maxn + 1);
void precompute() {
    fs[0] = 1;
    for (ll i = 1; i <= maxn; i++) {
        fs[i] = (fs[i - 1] * i) % MOD;
}
ll fpow(ll a, int n, ll mod = LLONG MAX) {
    if (n == 0) return 1:
    if (n == 1) return a;
    ll x = fpow(a, n / 2, mod) % mod;
    return ((x * x) % mod * (n & 1 ? a : 111)) % mod;
ll distinctAnagrams(const string &s) {
    precompute();
    vi hist('z' - 'a' + 1, 0);
    for (auto &c : s) hist[c - 'a']++;
    ll ans = fs[len(s)];
    for (auto &q : hist) {
        ans = (ans * fpow(fs[q], MOD - 2, MOD)) % MOD;
    return ans;
}
```

13.3 Double hash range query

```
int n;
    vll h, h2, hi, hi2;
    Hash() {}
    Hash(const string \&s) : n(s.size()), h(n), h2(n), hi(n), hi2(n) {
        h[0] = h2[0] = s[0];
        for (int i = 1; i < n; i++)
            h[i] = (s[i] + h[i - 1] * P) % MOD
            h2[i] = (s[i] + h2[i - 1] * P) % MOD2;
        hi[n - 1] = hi2[n - 1] = s[n - 1];
        for (int i = n - 2; i >= 0; i--)
            hi[i] = (s[i] + hi[i + 1] * P) % MOD
            hi2[i] = (s[i] + hi2[i + 1] * P) % MOD2;
    pll query(int l, int r) {
        ll hash = (h[r] - (l ? h[l - 1] * p[r - l + 1] % MOD : 0));
        ll hash2 = (h2[r] - (l? h2[l - 1] * p2[r - l + 1] % MOD2 : 0));
        return {(hash < 0 ? hash + MOD : hash),</pre>
                \{hash2 < 0 ? hash2 + MOD2 : hash2\};
    pll query inv(int l, int r) {
        ll hash = (hi[l] - (r + 1 < n ? hi[r + 1] * p[r - l + 1] % MOD :
   0));
        ll\ hash2 =
            (hi2[l] - (r + 1 < n ? hi2[r + 1] * p2[r - l + 1] % MOD2 : 0))
        return {(hash < 0 ? hash + MOD : hash),</pre>
                \{hash2 < 0 ? hash2 + MOD2 : hash2\};
};
```

13.4 Hash range query

```
#include "../Contest/template.cpp"
const ll P = 31;
const ll\ MOD = 1e9 + 9;
const int MAXN(1e6):
ll ppow[MAXN + 1];
void pre calc() {
    ppow[0] = 1;
    for (int i = 1; i \le MAXN; i++) ppow[i] = (ppow[i - 1] * P) % MOD;
struct Hash {
    int n:
    vll h, hi;
    Hash(const string &s) : n(s.size()), h(n), hi(n) {
        h[0] = s[0];
        hi[n - 1] = s[n - 1];
        for (int i = 1; i < n; i++) {
            h[i] = (s[i] + h[i - 1] * P) % MOD;
            hi[n - i - 1] = (s[n - i - 1] + hi[n - i - 1] * P) % MOD;
    ll gry(int l, int r) {
```

```
ll hash = (h[r] - (l ? h[l - 1] * ppow[r - l + 1] % MOD : 0));
    return hash < 0 ? hash + MOD : hash;
}
ll qry_inv(int l, int r) {
    ll hash = (hi[l] - (r + 1 < n ? hi[r + 1] * ppow[r - l + 1] % MOD : 0));
    return hash < 0 ? hash + MOD : hash;
}
};</pre>
```

13.5 Hash unsigned long long $2^{64} - 1$

Description: Arithmetic mod $2^{64} - 1$. 2x slower than mod 2^{64} and more code, but works on evil test data (e.g. Thue-Morse, where ABBA... and BAAB... of length 2^{10} hash the same mod 2^{64}).

"typedef ull H;" instead if you think test data is random.

```
#include "../Contest/template.cpp"
struct H {
    ull x;
    H(ull x = 0) : x(x) {}
    H operator+(H o) { return x + o.x + (x + o.x < x); }
    H operator-(H o) { return *this + \sim0.x; }
    H operator*(H o) {
        auto m = (uint128 t)x * o.x;
        return H((\overline{ul})m) + (\overline{ul})(m >> 64):
    ull get() const { return x + !\sim x; }
    bool operator==(H o) const { return get() == o.get(); }
    bool operator<(H o) const { return get() < o.get(); }</pre>
static const H C = (long long)lel1 + 3; // (order \sim 3e9; random also ok)
struct Hash {
    int n;
    vector<H> ha, pw;
    Hash(string \&str) : n(str.size()), ha((int)str.size() + 1), pw(ha) {
        for (int i = 0; i < (int)str.size(); i++)</pre>
             ha[i + 1] = ha[i] * C + str[i], pw[i + 1] = pw[i] * C;
    H query(int a, int b) { // hash [a, b]
        return ha[b] - ha[a] * pw[b - a];
    }
};
vector<H> getHashes(string &str, int length) {
    if ((int)str.size() < length) return {};</pre>
    H h = 0. pw = 1:
    for (int i = 0; i < length: i++) h = h * C + str[i]. pw = pw * C:
    vector<H> ret = {h};
    for (int i = length; i < (int)str.size(); i++)</pre>
        ret.push back(h = h * C + str[i] - pw * str[i - length]);
    return ret:
```

```
}
H hashString(string &s) {
    H h{};
    for (char c : s) h = h * C + c;
    return h;
}
```

13.6 K-th digit in digit string

Description: Find the k-th digit in a *digit string*, only works for $1 \le k \le 10^{18}$! **Time**: precompute O(1), query O(1)

```
using vull = vector<ull>:
vull pow10:
vector<array<ull, 4>> memo;
void precompute(int maxpow = 18) {
    ull atd = 1:
    ull start = 1:
    ull end = 9:
    ull curlenght = 9;
    ull startstr = 1:
    ull endstr = 9;
    for (ull i = 0, j = 111; (int) i < maxpow; i++, j *= 1011) pow10.eb(j);
    for (ull i = 0; i < maxpow - 1ull; i++) {
        memo.push back({start, end, startstr, endstr});
        start = end + 1ll:
        end = end + (9ll * pow10[atd]);
        curlenght = end - start + 1ull;
        atd++:
        startstr = endstr + 1ull;
        endstr = (endstr + 1ull) + (curlenght)*qtd - 1ull;
char kthDigit(ull k) {
    int atd = 1:
    for (auto [s, e, ss, es] : memo) {
        if (k \ge ss and k \le ss) {
            ull pos = k - ss;
            ull index = pos / qtd;
            ull nmr = s + index:
            int i = k - ss - atd * index;
            return ((nmr / pow10[atd - i - 1]) \% 10) + '0':
        qtd++;
    }
    return 'X';
```

13.7 KMP

```
* Author: Johan Sannemo
 * Date: 2016-12-15
 * License: CC0
 * Description: pi[x] computes the length of the longest prefix of s that
 * at x, other than s[0...x] itself (abacaba \rightarrow 0010123). Can be used to
 * all occurrences of a string. Time: O(n) Status: Tested on
 * kattis:stringmatching
 * @Title: Prefix function - Knuth-Morris-Pratt
 * @Description: Given a string $S$ builds an array $A$ such that
 * $A i$ is the longest suffix that ends in $i$ and is also a prefix
 * of $S$.
#pragma once
vi pi(const string& s) {
    vi p(sz(s));
    rep(i, 1, sz(s)) {
        int g = p[i - 1];
        while (g \&\& s[i] != s[g]) g = p[g - 1];
        p[i] = g + (s[i] == s[g]);
    return p;
vi match(const string& s, const string& pat) {
    vi p = pi(pat + '\0' + s), res;
    rep(i, sz(p) - sz(s), sz(p)) if (p[i] == sz(pat))
        res.push back(i - 2 * sz(pat));
    return res:
}
```

13.8 Longest Palindrome Substring (Manacher)

Description: Finds the longest palindrome substring, manacher returns a vector where the i-th position is how much is possible to grow the string to the left and the right of i and keep it a palindrome.

Time: O(N)

```
vi manacher(const string &s) {
    int n = len(s) - 2;
    vi p(n + 2);
    int l = 1, r = 1;
    for (int i = 1; i <= n; i++) {
        p[i] = max(0, min(r - i, p[l + (r - i)]));
        while (s[i - p[i]] == s[i + p[i]]) p[i]++;
        if (i + p[i] > r) l = i - p[i], r = i + p[i];
        p[i]--;
    }
    return p;
}
string longest_palindrome(const string &s) {
```

```
string t("$#");
for (auto c : s) t.push_back(c), t.push_back('#');
t.push_back('^');
vi xs = manacher(t);
int mpos = max_element(all(xs)) - xs.begin();
string p;
for (int k = xs[mpos], i = mpos - k; i <= mpos + k; i++)
    if (t[i] != '#') p.push_back(t[i]);
return p;
}</pre>
```

13.9 Longest palindrome

```
string longest palindrome(const string &s) {
    int n = (int)s.size();
    vector<array<int, 2>> dp(n);
    pii odd(0, -1), even(0, -1);
    pii ans;
    for (int i = 0; i < n; i++) {
        int k = 0;
        if (i > odd.second)
             k = 1;
        else
             k = min(dp[odd.first + odd.second - i][0], odd.second - i + 1)
        while (i - k >= 0 \text{ and } i + k < n \text{ and } s[i - k] == s[i + k]) k++;
        dp[i][0] = k--;
        if (i + k > odd.second) odd = \{i - k, i + k\};
        if (2 * dp[i][0] - 1 > ans.second) ans = \{i - k, 2 * dp[i][0] - 1\}
   1};
        k = 0:
        if (i <= even.second)</pre>
             k = min(dp[even.first + even.second - i + 1][1],
                      even.second -i + 1;
        while (i - k - 1 \ge 0 \text{ and } i + k < n \text{ and } s[i - k - 1] == s[i + k])
    k++;
        dp[i][1] = k--;
        if (i + k > even.second) even = \{i - k - 1, i + k\};
        if (2 * dp[i][1] > ans.second) ans = \{i - k - 1, 2 * dp[i][1]\};
    return s.substr(ans.first, ans.second);
}
```

13.10 Lyndon factorization

```
vi lyndon_factorization(string S) {
   auto sa = suffix_array(S);
   vi ans;
   vi mex(len(S) + 1, 0);
   int p = 0;
   rtrav(si, sa) {
      if (si == p) {
        ans.eb(si);
   }
}
```

```
}
    mex[si] = 1;
    while (mex[p]) p++;
}
ans.eb(len(S));
return ans;
```

13.11 Rabin-Karp

```
size_t rabin_karp(const string &s, const string &p) {
   \overline{i}f (s.si\overline{z}e() < p.size()) return 0;
    auto n = s.size(), m = p.size();
    const ll p1 = 31, p2 = 29, q1 = 1e9 + 7, q2 = 1e9 + 9;
    const ll p1 1 = fpow(p1, q1 - 2, q1), p1 2 = fpow(p1, m - 1, q1);
    const ll p2_1 = fpow(p2, q2 - 2, q2), p2_2 = fpow(p2, m - 1, q2);
    pair<ll, ll> hs, hp;
    for (int i = (int)m - 1; \sim i; --i) {
        hs.first = (hs.first * p1) % q1;
        hs.first = (hs.first + (s[i] - 'a' + 1)) % q1;
        hs.second = (hs.second * p2) % q2;
        hs.second = (hs.second + (s[i] - 'a' + 1)) % q2;
        hp.first = (hp.first * p1) % q1;
        hp.first = (hp.first + (p[i] - 'a' + 1)) % q1;
        hp.second = (hp.second * p2) % q2;
        hp.second = (hp.second + (p[i] - 'a' + 1)) % q2;
   size t occ = 0;
   for (size t i = 0; i < n - m; i++) {
        occ += (hs == hp);
        int fi = s[i] - 'a' + 1;
        int fm = s[i + m] - 'a' + 1;
        hs.first = (hs.first - fi + al) % al:
        hs.first = (hs.first * p1 1) % q1;
        hs.first = (hs.first + fm * p1 2) % q1;
        hs.second = (hs.second - fi + \overline{q}2) % q2;
        hs.second = (hs.second * p2 1) % q2;
        hs.second = (hs.second + fm * p2 2) % q2;
   occ += hs == hp;
    return occ;
```

13.12 Suffix Automaton

Description: A suffix automaton A for a string s is a minimal finite automaton that recognizes the suffixes of s.

```
#pragma once
#include "../Contest/template.cpp"
struct SuffixAutomaton {
```

13.13 Suffix array

```
#include <bits/stdc++.h>
using namespace std;
#ifdef LOCAL
#include "debug.cpp"
#else
#define dbg(...)
#endif
#define endl '\n'
#define fastio
    ios base::sync with stdio(0); \
    cin.tie(0);
#define int long long
#define all(j) j.begin(), j.end()
#define rall(j) j.rbegin(), j.rend()
#define len(j) (int)j.size()
#define rep(i, a, b) \
    for (common type t < decltype(a), decltype(b) > i = (a); i < (b); i++)
#define rrep(i, a, b) \
    for (common_type_t<decltype(a), decltype(b)> i = (a); i > (b); i--)
#define trav(xi, xs) for (auto &xi : xs)
#define rtrav(xi, xs) for (auto &xi : ranges::views::reverse(xs))
#define pb push back
#define pf push front
#define ppb pop back
#define ppf pop front
#define eb emplace back
#define lb lower bound
#define ub upper bound
#define fi first
#define se second
#define emp emplace
#define ins insert
#define divc(a, b) ((a) + (b) - 111) / (b)
using str = string;
using ll = long long;
using ull = unsigned long long;
using ld = long double;
using vll = vector<ll>:
```

```
using pll = pair<ll, ll>;
using vll2d = vector<vll>;
using vi = vector<int>;
using vi2d = vector<vi>;
using pii = pair<int, int>;
using vpii = vector<pii>;
using vc = vector<char>;
using vs = vector<str>;
template <typename T, typename T2>
using umap = unordered_map<T, T2>;
template <typename T>
using pgmn = priority queue<T, vector<T>, greater<T>>;
template <typename T>
using pqmx = priority queue<T, vector<T>>;
template <typename T, typename U>
inline bool chmax(T &a, U const &b) {
    return (a < b ? a = b, 1 : 0);
template <typename T, typename U>
inline bool chmin(T &a, U const &b) {
    return (a > b ? a = b, 1 : 0);
vector<int> sort cyclic shifts(string const &s) {
   int n = s.size();
    const int alphabet = 128;
   vector<int> p(n), c(n), cnt(max(alphabet, n), 0);
   for (int i = 0; i < n; i++) cnt[s[i]]++;
   for (int i = 1; i < alphabet; i++) cnt[i] += cnt[i - 1];</pre>
   for (int i = 0; i < n; i++) p[--cnt[s[i]]] = i;
   c[p[0]] = 0;
   int classes = 1;
   for (int i = 1; i < n; i++) {
        if (s[p[i]] != s[p[i - 1]]) classes++;
        c[p[i]] = classes - 1;
   vector<int> pn(n), cn(n);
   for (int h = 0; (1 << h) < n; ++h) {
        for (int i = 0; i < n; i++) {
            pn[i] = p[i] - (1 << h);
            if (pn[i] < 0) pn[i] += n;
        fill(cnt.begin(), cnt.begin() + classes, 0);
        for (int i = 0; i < n; i++) cnt[c[pn[i]]]++;
        for (int i = 1; i < classes; i++) cnt[i] += cnt[i - 1];</pre>
        for (int i = n - 1; i \ge 0; i--) p[--cnt[c[pn[i]]]] = pn[i];
        cn[p[0]] = 0;
        classes = 1:
        for (int i = 1; i < n; i++) {
            pair < int, int > cur = {c[p[i]], c[(p[i] + (1 << h)) % n]};
            pair<int, int> prev = \{c[p[i-1]], c[(p[i-1] + (1 << h)) \%
   n]};
            if (cur != prev) ++classes;
            cn[p[i]] = classes - 1;
        c.swap(cn);
```

```
}
    return p;
vector<int> suffix array(string s) {
    s += "$";
    vector<int> p = sort cyclic shifts(s);
    p.erase(p.begin());
    return p;
vector<int> longestCommonPrefix(const string &s, const vector<int> &suf) {
    int n = s.size():
    vector<int> isuf(n), res(n - 1);
    for (int i = 0; i < n; ++i) isuf[suf[i]] = i;
    int k = 0;
    for (; isuf[k] != n - 1; ++k) {
        int cmp i = suf[isuf[k] + 1];
        int r = k == 0 ? 0 : max(res[isuf[k - 1]] - 1, (int)0);
        while (k + r < n \& \& cmp i + r < n \& \& s[k + r] == s[cmp i + r]) ++r
        res[isuf[k]] = r;
    ++k;
    for (int i = k; i < n; ++i) {
        int cmp i = suf[isuf[i] + 1];
        int r = i == k ? 0 : max(res[isuf[i - 1]] - 1, (int)0);
        while (i + r < n \&\& cmp_i + r < n \&\& s[i + r] == s[cmp_i + r]) ++r
        res[isuf[i]] = r;
    return res;
ll distinct substrings(const string &s, const vi &sa) {
    int n = len(s);
    vi lcp = longestCommonPrefix(s, sa);
    ll ans = n - sa[0];
    rep(i, 1, n) \{ ans += n - sa[i] - lcp[i - 1]; \}
    return ans;
void run();
int32 t main() {
#ifndef LOCAL
    fastio:
#endif
    int T = 1;
    /*cin >> T:*/
    rep(t, 0, T) {
        dbg(t);
        run();
void run() {
    string S;
    cin >> S;
    auto sa = suffix array(S);
    cout << distinct_substrings(S, sa) << endl;</pre>
```

```
}
```

13.14 Suffix array (supreme)

```
template <typename T = ll,
          auto cmp = [](T \& src1, T \& src2, T \& dst) \{ dst = min(src1, src2); \}
class SparseTable {
   private:
   int sz;
    vi logs;
    vector<vector<T>> st;
   public:
    SparseTable() {}
    SparseTable(const vector<T> \&v) : sz(len(v)), logs(sz + 1) {
        rep(i, 2, sz + 1) logs[i] = logs[i >> 1] + 1;
        st.resize(logs[sz] + 1, vector<T>(sz));
        rep(i, 0, sz) st[0][i] = v[i];
        for (int k = 1; (1 << k) <= sz; k++) {
            for (int i = 0; i + (1 << k) <= sz; i++) {
                cmp(st[k-1][i], st[k-1][i+(1 << (k-1))], st[k][i])
        }
    T query(int l, int r) {
        const int k = logs[r - l]:
        T ret;
        cmp(st[k][l], st[k][r - (1 << k)], ret);
        return ret:
};
template <typename T>
using RMQ = SparseTable<T, [](T \&a, T \&b, T \&c) \{ c = min(a, b); \}>;
// éCrditos: ShahjalalShohag
// O(N)
struct SA {
    string s;
    int n:
    vector<int> sa, lcp, pos;
    RMQ<int> rmg;
    void induced sort(vector<int> &vec, int val, vector<int> &sa,
                      vector<bool> &sl, vector<int> &lms) {
        vector<int> l(val), r(val);
        for (int c : vec) {
            if (c + 1 < val) | (c + 1) + +:
            r[c]++;
        partial sum(l.begin(), l.end(), l.begin());
        partial sum(r.begin(), r.end(), r.begin());
        fill(sa.begin(), sa.end(), -1);
        for (int i = lms.size() - 1; i >= 0; i--) sa[--r[vec[lms[i]]]] =
   lms[i];
```

```
for (int i : sa) {
                       if (i >= 1 \&\& sl[i - 1]) sa[l[vec[i - 1]]++] = i - 1;
            fill(r.begin(), r.end(), 0);
            for (int c : vec) r[c]++:
            partial_sum(r.begin(), r.end(), r.begin());
           for (int k = sa.size() - 1, i = sa[k]; k >= 1; --k, i = sa[k]) {
                       if (i \ge 1 \&\& !sl[i-1]) sa[--r[vec[i-1]]] = i-1;
 }
 vector<int> build sa(vector<int> &vec, int val) {
            int n = vec.s\overline{i}ze():
            vector<int> sa(n), lms;
            vector<bool> sl(n);
            sl[n - 1] = false;
            for (int i = n - 2; i \ge 0; i--) {
                       sl[i] =
                                  (\text{vec}[i] > \text{vec}[i + 1] \mid | (\text{vec}[i] == \text{vec}[i + 1] \& sl[i + 1])
11));
                      if (sl[i] \&\& !sl[i+1]) lms.push back(i+1);
            reverse(lms.begin(), lms.end());
            induced sort(vec, val, sa, sl, lms);
            vector<int> new_lms(lms.size()), lms_vec(lms.size());
            for (int i = 0, k = 0; i < n; i++) {
                      if (!sl[sa[i]] \&\& sa[i] >= 1 \&\& sl[sa[i] - 1]) new_lms[k++] =
sa[i];
            int cur = 0;
            sa[n - 1] = cur:
            for (int k = 1; k < (int)new_lms.size(); k++) {</pre>
                       int i = new lms[k - 1], \overline{j} = new lms[k];
                       if (vec[i] != vec[i]) {
                                 sa[j] = ++cur;
                                 continue;
                       bool flag = false:
                       for (int a = i + 1, b = j + 1; ++a, ++b) {
                                 if (vec[a] != vec[b]) {
                                            flag = true;
                                            break:
                                 if ((!sl[a] && sl[a - 1]) || (!sl[b] && sl[b - 1])) {
                                            flag = !((!sl[a] \&\& sl[a - 1]) \&\& (!sl[b] \&\& sl[b - 1]) \&\& (!sl[b] \&\&
1]));
                                            break;
                       sa[j] = (flag ? ++cur : cur);
           for (int i = 0; i < (int)lms.size(); i++) lms vec[i] = sa[lms[i]];</pre>
            if (cur + 1 < (int)lms.size()) {</pre>
                       auto lms sa = build sa(lms vec, cur + 1);
                       for (int i = 0; i < (int)lms.size(); i++)</pre>
                                 new lms[i] = lms[lms_sa[i]];
            }
```

```
induced sort(vec, val, sa, sl, new lms);
    return sa;
vector<int> suffix_array() {
    vector < int > vec(n + 1);
    copy(begin(s), end(s), begin(vec));
    vec.back() = '$';
    auto sa = build sa(vec, 256);
    sa.erase(sa.begin());
    return sa;
vector<int> build lcp() {
    int n = (int)\overline{s}.size(), k = 0;
    vector<int> rank(n), lcp(n);
    for (int i = 0; i < n; i++) rank[sa[i]] = i;
    for (int i = 0; i < n; i++, k -= !!k) {
        if (rank[i] == n - 1) {
            k = 0;
            continue;
        int j = sa[rank[i] + 1];
        while (i + k < n \&\& j + k < n \&\& s[i + k] == s[j + k]) k++;
        lcp[rank[i]] = k;
    return lcp;
}
SA() {}
SA(string s) : s(s), n(len(s)), pos(n) {
    sa = suffix array();
    lcp = build lcp();
    rmq = RMQ<int>(lcp);
    for (int i = 0; i < n; i++) pos[sa[i]] = i;</pre>
int get lcp(int i,
            int j) { // lcp na çãposio i, indica o lcp
                       // das çõposies i e i+1 do sa
    if (i == j) return n - i;
    int l = pos[i], r = pos[j];
    if (l > r) swap(l, r);
    return rmg.query(l, r);
// string s = a + '+' + b;
tuple<int, int, int> lcs(int n) { // m é o tamanho da string a
    int m = len(s) - n - 1;
    int best len = 0;
    int index s = 0;
    int index t = 0:
    for (int \overline{i} = 0; i < n + m; ++i) {
        if ((sa[i] < n \&\& sa[i + 1] >= n + 1) ||
            (sa[i] >= n + 1 \&\& sa[i + 1] < n)) {
            if (lcp[i] > best len) {
                 best len = lcp[i];
                index s = min(sa[i], sa[i + 1]);
```

13.15 Suffix automaton

```
#include <bits/stdc++.h>
using namespace std;
#ifdef LOCAL
#include "debug.cpp"
#else
#define dbg(...)
#endif
#define endl '\n'
#define fastio
    ios_base::sync_with_stdio(0); \
    cin.tie(0);
#define int long long
#define all(j) j.begin(), j.end()
#define rall(j) j.rbegin(), j.rend()
#define len(j) (int)j.size()
#define rep(i, a, b) \
    for (common type t < decltype(a), decltype(b) > i = (a); i < (b); i++)
#define rrep(i, a, b) \
    for (common type t<decltype(a), decltype(b) > i = (a); i > (b); i--)
#define trav(xi, xs) for (auto &xi : xs)
#define rtrav(xi, xs) for (auto &xi : ranges::views::reverse(xs))
#define pb push back
#define pf push front
#define ppb pop back
#define ppf pop front
#define eb emplace back
#define lb lower bound
#define ub upper bound
#define fi first
#define se second
#define emp emplace
#define ins insert
#define divc(a, b) ((a) + (b) - 111) / (b)
```

```
using str = string;
using ll = long long:
using ull = unsigned long long;
using ld = long double;
using vll = vector<ll>;
using pll = pair<ll, ll>;
using vll2d = vector<vll>;
using vi = vector<int>;
using vi2d = vector<vi>;
using pii = pair<int, int>;
using vpii = vector<pii>;
using vc = vector<char>;
using vs = vector<str>;
template <typename T, typename T2>
using umap = unordered map<T, T2>;
template <typename T>
using pgmn = priority queue<T, vector<T>, greater<T>>;
template <tvpename T>
using pqmx = priority queue<T, vector<T>>;
template <typename T, typename U>
inline bool chmax(T &a, U const &b) {
   return (a < b ? a = b, 1 : 0);
template <typename T, typename U>
inline bool chmin(T &a, U const &b) {
   return (a > b ? a = b, 1 : 0);
struct SuffixAutomaton {
   struct state {
       int len, link, cnt, firstpos;
       // this can be optimized using a vector with
       // the alphabet size
       map<char, int> next;
       vi inv link;
   };
   vector<state> st;
   int sz = 0;
   int last;
   vc cloned:
   SuffixAutomaton(const string &s, int maxlen)
        : st(maxlen * 2), cloned(maxlen * 2) {
       st[0].len = 0;
       st[0].link = -1;
       SZ++;
       last = 0;
       for (auto &c : s) add char(c);
       // precompute for count occurences
       for (int i = 1; i < sz; i++) {
            st[i].cnt = !cloned[i];
       vector<pair<state, int>> aux;
       for (int i = 0; i < sz; i++) {
            aux.push back({st[i], i});
       sort(all(aux),
```

```
[](const pair<state, int> &a, const pair<state, int> &b) {
             return a.fi.len > b.fi.len:
         });
    for (auto &[stt, id] : aux) {
        if (stt.link !=-1) {
            st[stt.link].cnt += st[id].cnt;
    }
    // for find every occurende position
    for (int v = 1; v < sz; v++) {
        st[st[v].link].inv link.push back(v);
void add char(char c) {
    int cur = sz++;
    st[cur].len = st[last].len + 1;
    st[cur].firstpos = st[cur].len - 1;
    int p = last:
    // follow the suffix link until find a
    // transition to c
    while (p != -1 \text{ and } !st[p].next.count(c))  {
        st[pl.next[c] = cur;
        p = st[p].link;
    // there was no transition to c so create and
    // leave
    if (p == -1) {
        st[curl.link = 0:
        last = cur:
        return;
    }
    int q = st[p].next[c];
    if (st[p].len + 1 == st[q].len) {
        st[cur].link = q;
    } else {
        int clone = sz++:
        cloned[clone] = true;
        st[clone].len = st[p].len + 1;
        st[clone].next = st[q].next;
        st[clone].link = st[q].link;
        st[clone].firstpos = st[q].firstpos;
        while (p != -1 \text{ and } st[p].next[c] == q) {
            st[p].next[c] = clone;
            p = st[p].link;
        st[q].link = st[cur].link = clone;
    last = cur;
bool checkOccurrence(const string &t) { // O(len(t))
    int cur = 0;
    for (auto &c : t) {
        if (!st[cur].next.count(c)) return false;
        cur = st[cur].next[c];
    }
```

```
return true:
   ll totalSubstrings() { // distinct, O(len(s))
       ll tot = 0:
       for (int i = 1: i < sz: i++) {
            tot += st[i].len - st[st[i].link].len;
        return tot;
   }
   // count occurences of a given string t
   int countOccurences(const string &t) {
       int cur = 0:
       for (auto &c : t) {
            if (!st[cur].next.count(c)) return 0;
            cur = st[curl.next[c];
        return st[cur].cnt;
   // find the first index where t appears a
   // substring O(len(t))
   int firstOccurence(const string &t) {
       int cur = 0;
        for (auto c : t) {
            if (!st[cur].next.count(c)) return -1;
            cur = st[curl.next[c]:
        return st[cur].firstpos - len(t) + 1;
   vi everyOccurence(const string &t) {
       int cur = 0:
       for (auto c : t) {
            if (!st[cur].next.count(c)) return {};
            cur = st[cur].next[c];
       vi ans:
       getEveryOccurence(cur, len(t), ans);
        return ans;
   void getEveryOccurence(int v, int P length, vi &ans) {
       if (!cloned[v]) ans.pb(st[v].firstpos - P length + 1);
       for (int u : st[v].inv link) getEveryOccurence(u, P length, ans);
   }
void run();
int32 t main() {
#ifndef LOCAL
   fastio;
#endif
   int T = 1:
   /*cin >> T;*/
   rep(t, 0, T) {
        dbg(t);
       run();
```

```
void run() {
   string S;
   cin >> S;
   SuffixAutomaton sa(S, len(S));
   cout << sa.totalSubstrings() << endl;
}</pre>
```

13.16 Suffix-Tree (Ukkonen's Algorithm)

```
/**
* Author: Unknown
* Date: 2017-05-15
* Source: https://e-maxx.ru/algo/ukkonen
* Description: Ukkonen's algorithm for online suffix tree construction.
* Each node contains indices [l, r) into the string, and a list of child
* nodes. Suffixes are given by traversals of this tree, joining [l, r)
* substrings. The root is 0 (has l = -1, r = 0), non-existent children
* To get a complete tree, append a dummy symbol -- otherwise it may
   an incomplete path (still useful for substring matching, though).
* Time: $0(26N)$
* Status: stress-tested a bit
#pragma once
#include "../Contest/template.cpp"
struct SuffixTree {
    enum { N = 200010, ALPHA = 26 }; // N \sim 2*maxlen+10
    int toi(char c) { return c - 'a': }
    string a; // v = cur node, g = cur position
    int t[N][ALPHA], l[N], r[N], p[N], s[N], v = 0, q = 0, m = 2;
    void ukkadd(int i. int c) {
    suff:
        if (r[v] <= q) {</pre>
            if (t[v][c] == -1) {
                t[v][c] = m;
                l[m] = i:
                p[m++] = v;
                v = s[v]:
                a = r[v]:
                goto suff:
            v = t[v][c];
            a = l[v];
        if (q == -1 || c == toi(a[q]))
            a++;
        else {
            l[m + 1] = i;
            p[m + 1] = m:
            l(m) = l(v):
            r[m] = q;
            [v]q = [m]q
            t[m][c] = m + 1:
```

```
t[m][toi(a[q])] = v;
        l[v] = a:
        p[v] = m;
        t[p[m]][toi(a[l[m]])] = m;
        v = s[p[m]]:
        q = l[m];
        while (q < r[m]) {
            v = t[v][toi(a[q])];
            q += r[v] - l[v];
        if (q == r[m])
            s[m] = v:
        else
            s[m] = m + 2;
        q = r[v] - (q - r[m]);
        m += 2;
        qoto suff;
SuffixTree(string a) : a(a) {
    fill(r, r + N, len(a));
    memset(s, 0, sizeof s);
memset(t, -1, sizeof t);
    fill(t[1], t[1] + ALPHA, 0);
    s[0] = 1:
    \bar{1}[0] = \bar{1}[1] = -1;
    r[0] = r[1] = p[0] = p[1] = 0;
    rep(i, 0, len(a)) ukkadd(i, toi(a[i]));
// example: find longest common substring (uses ALPHA = 28)
pii best;
int lcs(int node, int i1, int i2, int olen) {
    if (l[node] <= i1 && i1 < r[node]) return 1;</pre>
    if (l[node] <= i2 && i2 < r[node]) return 2;</pre>
    int mask = 0, len = node ? olen + (r[node] - l[node]) : 0;
    rep(c, 0, ALPHA) if (t[node][c] != -1) mask |=
        lcs(t[node][c], i1, i2, len);
    if (mask == 3) best = max(best, {len, r[node] - len});
    return mask;
static pii LCS(string s, string t) {
    SuffixTree st(s + (char)('z' + 1) + t + (char)('z' + 2));
    st.lcs(0, len(s), len(s) + 1 + len(t), 0);
    return st.best;
```

13.17 Trie

};

Description:

- build with the size of the alphabet (sigma) and the first char (norm)
- insert(s) insert the string in the trie O(|s| * sigma)
- erase(s) remove the string from the trie O(|s|)
- find(s) return the last node from the string s, 0 if not found O(|s|)

```
#include "../Contest/template.cpp"
struct Trie {
    vi2d to;
    vi end, pref;
    int sigma;
    char norm;
    Trie(int sigma = 'z' - 'a' + 1, char norm = 'a')
        : sigma(sigma ), norm(norm ) {
        to = {vector<int>(sigma)};
        end = \{0\}, pref = \{0\};
    int next(int node, char key) { return to[node][key - norm]; }
    void insert(const string &s) {
        int x = 0:
        for (auto c : s) {
            int &nxt = to[x][c - norm];
            if (!nxt) {
                nxt = len(to):
                to.push back(vi(sigma));
                end.emplace back(0), pref.emplace back(0);
            x = nxt, pref[x]++;
        end[x]++, pref[0]++;
    }
    void erase(const string &s) {
        int x = 0:
        for (char c : s) {
            int &nxt = to[x][c - norm];
            x = nxt, pref[x]--;
            if (!pref[x]) nxt = 0;
        end[x]--, pref[0]--;
    int find(const string &s) {
        int x = 0:
        for (auto c : s) {
            x = to[x][c - norm];
            if (!x) return 0;
        return x;
};
```

13.18 Z-function get occurrence positions

Time: O(len(s) + len(p))
#include "../Contest/template.cpp"
vi get0ccPos(string& s, string& p) {
 // Z-function
 char delim = '#';
 string t{p + delim + s};
 vi zs(len(t));

```
for (int i = 1, l = 0, r = 0; i < len(t); i++) {
        if (i \le r) zs[i] = min(zs[i - l], r - i + 1);
        while (zs[i] + i < len(t) \text{ and } t[zs[i]] == t[i + zs[i]]) zs[i]++;
        if (r < i + zs[i] - 1) l = i, r = i + zs[i] - 1;
   // Iterate over the results of Z-function to get
   // ranges
   vi ans;
   int start = len(p) + 1 + 1 - 1;
    for (int i = start; i < len(zs); i++) {</pre>
        if (zs[i] == len(p)) {
            int l = i - start;
            ans.emplace_back(l);
    return ans;
template <class T>
std::vector<int> z algorithm(const std::vector<T>& s) {
    int n = int(s.size());
   if (n == 0) return {};
    std::vector<int> z(n);
    for (int i = 1, j = 0; i < n; i++) {
        int & k = z[i];
        k = (j + z[j] \le i) ? 0 : std::min(j + z[j] - i, z[i - j]);
        while (i + k < n \&\& s[k] == s[i + k]) k++;
        if (j + z[j] < i + z[i]) j = i;
```

```
z[0] = n;
    return z;
std::vector<int> z algorithm(const std::string& s) {
    int n = int(s.size());
    std::vector<int> s2(n);
    for (int i = 0; i < n; i++) {
        s2[i] = s[i];
    return z algorithm(s2);
// Pattern matching com um erro ??
// prefixo do apadro dando certo
// depois o sufixo continua dando certo...
// Pattern matching com x erros
// únmero de substrings distintas...
// únmero de bustrings distintas
// oline ??
// achar o íperodo
//
// kmp é áintercambivel com z
```