

# JavaScript Data Structures & Algorithms

Essential Cheat Sheet

```
class Node {  
  constructor(value) {  
    this.value = value;  
    this.next = null;  
  }  
}
```

Swipe for more >>

# Big O Notation

## Time Complexity Analysis

$O(1)$

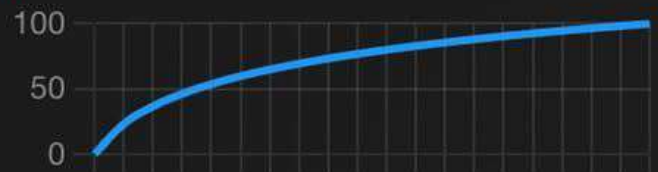
Constant Time



Always takes same time regardless of input size. Like accessing array index.

$O(\log n)$

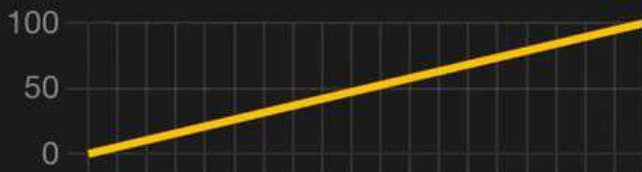
Logarithmic Time



Cuts problem in half each time. Like binary search.

$O(n)$

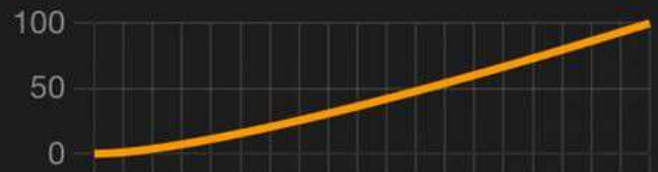
Linear Time



Time grows linearly with input. Like linear search.

$O(n \log n)$

Linearithmic Time



Common in efficient sorting algorithms like merge sort.

$O(n^2)$

Quadratic Time



Nested iterations. Like bubble sort.

$O(2^n)$

Exponential Time



Solutions double with each addition. Like recursive fibonacci.

# Arrays

## Common Operations & Time Complexity

### Access

$O(1)$

Get or set element directly using index position. Fastest array operation.

```
const value = array[5];
```

### Push/Pop

$O(1)$

Add (push) or remove (pop) element at end. Efficient as no reindexing needed.

```
array.push(1); // add  
array.pop();   // remove
```

### Unshift/Shift

$O(n)$

Add (unshift) or remove (shift) at start. Requires reindexing all elements.

```
array.unshift(1); // add  
array.shift();    // remove
```

### Search

$O(n)$

Find element by value. Must check each element until match found.

```
array.indexOf(5);  
array.find(x => x > 3);
```

### Insert

$O(n)$

Insert at position. Requires shifting all subsequent elements right.

```
array.splice(2, 0, 5);
```

### Delete

$O(n)$

Remove at position. Requires shifting all subsequent elements left.

```
array.splice(2, 1);
```

# Linked List

## Linear Data Structure

### Insert Head

$O(1)$

Add new node at beginning. Just update head pointer.

```
node.next = head;
head = node;
```

### Insert Tail

$O(1)*$

Add at end with tail pointer.  $O(n)$  without tail.

```
tail.next = node;
tail = node;
```

### Delete Head

$O(1)$

Remove first node. Update head pointer.

```
head = head.next;
```

### Insert/Delete Middle

$O(n)$

Need to traverse to find position first.

```
prev.next = node;
node.next = current;
```

### Search

$O(n)$

Must traverse nodes sequentially.

```
while(current) {
  if(current.val === x) return true;
  current = current.next;
}
```

### Access by Index

$O(n)$

No direct access. Must traverse from start.

```
for(let i = 0; i < index; i++) {
  current = current.next;
}
```



# Binary Trees

## Common Operations

### Insert

$O(\log n)^*$

Add new node in BST.  $O(n)$  worst case for unbalanced.

```
if (value < node.val)
  node.left = insert(node.left, value);
else
  node.right = insert(node.right, value);
```

### Search

$O(\log n)^*$

Find node in BST.  $O(n)$  worst case for unbalanced.

```
if (!node || node.val === value) return node;
return value < node.val
  ? search(node.left, value)
  : search(node.right, value);
```

### DFS Traversal

$O(n)$

Depth-first: preorder, inorder, postorder.

```
// Inorder (sorted in BST)
inorder(node) {
  inorder(node.left);
  visit(node);
  inorder(node.right);
}
```

### BFS Traversal

$O(n)$

Level by level using queue.

```
const queue = [root];
while (queue.length) {
  const node = queue.shift();
  queue.push(node.left, node.right);
}
```

### Delete

$O(\log n)^*$

Remove node and maintain BST property.

```
// Find min in right subtree
// Replace node with min
// Delete min from right
```

### Height/Depth

$O(n)$

Find max depth of tree.

```
return !node ? 0 : 1 + Math.max(
  height(node.left),
  height(node.right)
);
```

# Graphs

## Common Operations

### Add Vertex

$O(1)$

Add new vertex to adjacency list.

```
graph[vertex] = [];
```

### Add Edge

$O(1)$

Connect two vertices (nodes).

```
graph[v1].push(v2);  
// For undirected:  
graph[v2].push(v1);
```

### BFS

$O(V + E)$

Breadth-first traversal using queue.

```
const q = [start];  
while (q.length) {  
  const v = q.shift();  
  q.push(...graph[v]);  
}
```

### DFS

$O(V + E)$

Depth-first traversal using recursion/stack.

```
function dfs(vertex) {  
  visited.add(vertex);  
  graph[vertex].forEach(v => {  
    if (!visited.has(v)) dfs(v);  
  });  
}
```

### Find Path

$O(V + E)$

Find path between two vertices.

```
// Using BFS/DFS  
while (queue.length) {  
  const v = queue.shift();  
  if (v === target) return path;  
}
```

### Check Cycle

$O(V + E)$

Detect cycle in directed/undirected graph.

```
function hasCycle(v, parent) {  
  if (visited.has(v)) return true;  
  return graph[v].some(next =>  
    next !== parent && hasCycle(next, v)  
  );  
}
```

# Sorting

## Common Algorithms

### Bubble Sort

$O(n^2)$

Compare adjacent elements and swap if needed. Simple but inefficient.

```
for (let i = 0; i < n; i++) {
  for (let j = 0; j < n - i - 1; j++) {
    if (arr[j] > arr[j + 1]) {
      [arr[j], arr[j + 1]] = [arr[j + 1], arr[j]];
    }
  }
}
```

### Quick Sort

$O(n \log n)$ \*

Pick pivot, partition array, recursively sort.  $O(n^2)$  worst case.

```
function quickSort(arr, low, high) {
  if (low < high) {
    let pi = partition(arr, low, high);
    quickSort(arr, low, pi - 1);
    quickSort(arr, pi + 1, high);
  }
}
```

### Merge Sort

$O(n \log n)$

Divide array, sort halves, merge results. Stable sort.

```
function mergeSort(arr) {
  if (arr.length <= 1) return arr;
  let mid = Math.floor(arr.length / 2);
  return merge(
    mergeSort(arr.slice(0, mid)),
    mergeSort(arr.slice(mid))
  );
}
```

### Insertion Sort

$O(n^2)$

Build sorted array one item at a time. Good for small  $n$ .

```
for (let i = 1; i < n; i++) {
  let key = arr[i];
  let j = i - 1;
  while (j >= 0 && arr[j] > key) {
    arr[j + 1] = arr[j];
    j--;
  }
  arr[j + 1] = key;
}
```

### Heap Sort

$O(n \log n)$

Build max heap, repeatedly extract max. In-place sort.

```
function heapSort(arr) {
  buildMaxHeap(arr);
  for (let i = arr.length - 1; i > 0; i--) {
    [arr[0], arr[i]] = [arr[i], arr[0]];
    heapify(arr, 0, i);
  }
}
```

### Selection Sort

$O(n^2)$

Find min element, place at start. Simple but inefficient.

```
for (let i = 0; i < n - 1; i++) {
  let min = i;
  for (let j = i + 1; j < n; j++) {
    if (arr[j] < arr[min]) min = j;
  }
  if (min !== i) [arr[i], arr[min]] = [arr[min], arr[i]];
}
```

# Search

## Common Algorithms

### Linear Search

$O(n)$

Check each element sequentially. Works on unsorted arrays.

```
function linearSearch(arr, target) {
  for (let i = 0; i < arr.length; i++) {
    if (arr[i] === target) return i;
  }
  return -1;
}
```

### Binary Search

$O(\log n)$

Divide and conquer. Requires sorted array.

```
function binarySearch(arr, target) {
  let left = 0, right = arr.length - 1;
  while (left <= right) {
    let mid = Math.floor((left + right) / 2);
    if (arr[mid] === target) return mid;
    if (arr[mid] < target) left = mid + 1;
    else right = mid - 1;
  }
  return -1;
}
```

### Jump Search

$O(\sqrt{n})$

Jump fixed steps, then linear search. For sorted arrays.

```
function jumpSearch(arr, target) {
  const step = Math.floor(Math.sqrt(arr.length));
  let prev = 0;
  while (arr[Math.min(step, arr.length) - 1] < target) {
    prev = step;
    step += Math.floor(Math.sqrt(arr.length));
    if (prev >= arr.length) return -1;
  }
  // Linear search in block
  while (arr[prev] < target) {
    prev++;
    if (prev === Math.min(step, arr.length)) return -1;
  }
  return arr[prev] === target ? prev : -1;
}
```

### Interpolation Search

$O(\log \log n)$

Like binary search but with better position guessing.

```
function interpolationSearch(arr, target) {
  let low = 0, high = arr.length - 1;
  while (low <= high && target >= arr[low] && target <= arr[high]) {
    let pos = low + Math.floor(
      (target - arr[low]) * (high - low) /
      (arr[high] - arr[low])
    );
    if (arr[pos] === target) return pos;
    if (arr[pos] < target) low = pos + 1;
    else high = pos - 1;
  }
  return -1;
}
```



# Thank You!

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