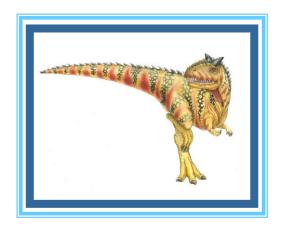
Chapter 4: Threads & Concurrency





Outline

- Overview
- Multicore Programming
- Multithreading Models
- Thread Libraries
- Operating System Examples





Objectives

- Identify the basic components of a thread, and contrast threads and processes
- Describe the benefits and challenges of designing multithreaded applications
- Designing multithreaded applications using the Pthreads





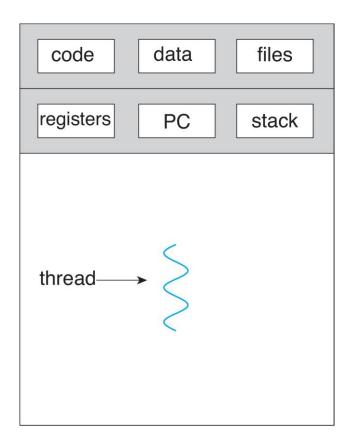
Motivation

- Most modern applications are multithreaded
- Threads run within application
- Multiple tasks with the application can be implemented by separate threads
 - Update display
 - Fetch data
 - Spell checking
 - Answer a network request
- Process creation is heavy-weight while thread creation is light-weight
- Can simplify code, increase efficiency
- Kernels are generally multithreaded

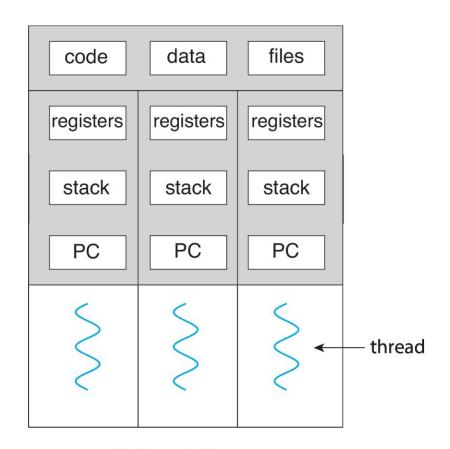




Single and Multithreaded Processes



single-threaded process

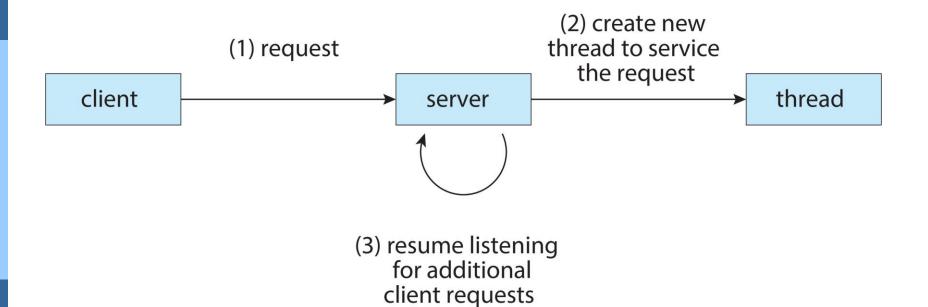


multithreaded process





Multithreaded Server Architecture







Benefits

- Responsiveness may allow continued execution if part of process is blocked, especially important for user interfaces
- Resource Sharing threads share resources of process, easier than shared memory or message passing
- Economy cheaper than process creation, thread switching lower overhead than context switching
- Scalability process can take advantage of multicore architectures





Multicore Programming

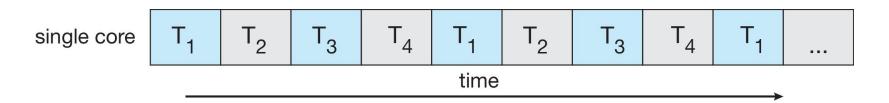
- In response to the need for more computing performance, single-CPU systems evolved into multi-CPU systems.
 - Current trend in system design is to place multiple computing cores on a single processing chip
- Parallelism implies a system can perform more than one task simultaneously
- Concurrency supports more than one task making progress
 - Single processor / core, scheduler providing concurrency



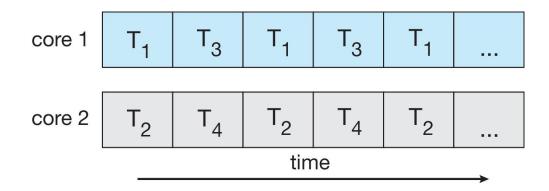


Concurrency vs. Parallelism

Concurrent execution on single-core system:



Parallelism on a multi-core system:







Multicore Programming

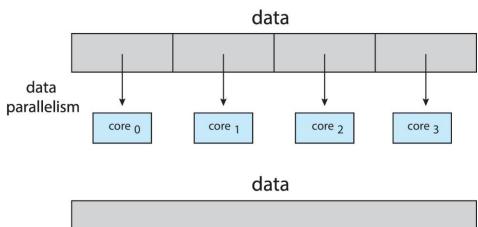
- Multicore or multiprocessor systems puts pressure on programmers, challenges include:
 - Dividing activities
 - Balance
 - Data splitting
 - Data dependency
 - Testing and debugging

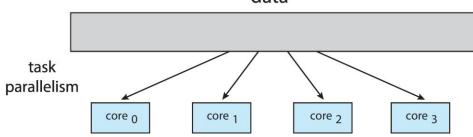




Multicore Programming

- Types of parallelism
 - Data parallelism distributes subsets of the same data across multiple cores, same operation on each
 - Task parallelism distributing threads across cores, each thread performing unique operation









Amdahl's Law

- Identifies performance gains from adding additional cores to an application that has both serial and parallel components
- S is serial portion
- N processing cores

$$speedup \le \frac{1}{S + \frac{(1-S)}{N}}$$

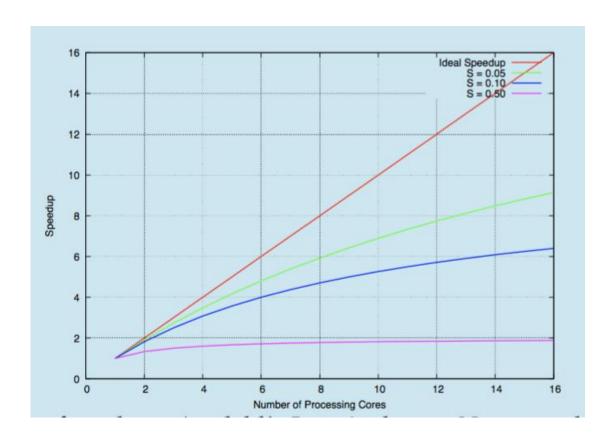
- That is, if application is 75% parallel / 25% serial, moving from 1 to 2 cores results in speedup of 1.6 times
- As N approaches infinity, speedup approaches 1 / S

Serial portion of an application has disproportionate effect on performance gained by adding additional cores

But does the law take into account contemporary multicore systems?



Amdahl's Law







Thread Libraries

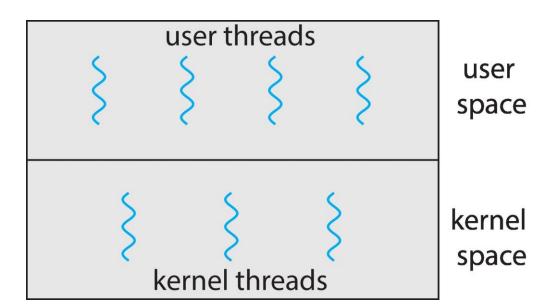
- Thread libraries provide programmers with an API for creating and managing threads.
- Thread libraries may be implemented either in user or in kernel space.
 - User space; API functions implemented solely within user space, with no kernel support.
 - Kernal space; involves system calls, and requires a kernel with thread library support.
 - A few well established primary thread libraries
 - POSIX Pthreads may be provided as either a user or kernel library
 - Win32 threads provided as a kernel-level library on Windows systems.
 - Java threads May be Pthreads or Win32 depending on the OS and hardware the JVM is running.





User Threads and Kernel Threads

- User threads management done by user-level threads library
- Kernel threads Supported by the Kernel
 - Exists virtually in all general purpose OS:
 - Windows, Linux, Mac OS X, iOS, Android
- Even user threads will ultimately need kernel thread support (Why??)







Multithreading Models

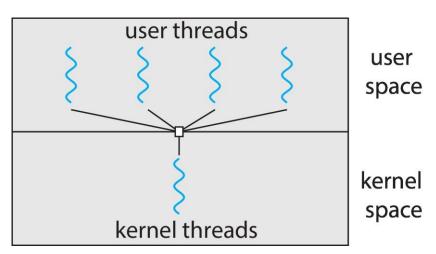
- Many-to-One
- One-to-One
- Many-to-Many





Many-to-One

- Many user-level threads mapped to single kernel thread
- Blocking one thread causes all to block
- Multiple threads may not run in parallel on multicore system because only one may be in kernel at a time
- Old approach: Few systems currently use this model
- Examples:
 - Solaris Green Threads
 - GNU Portable Threads

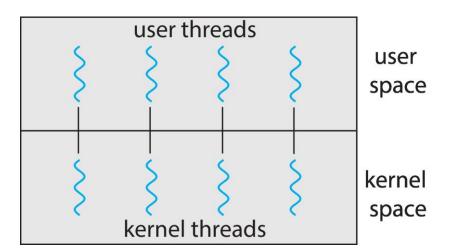






One-to-One

- Each user-level thread maps to kernel thread
- Creating a user-level thread creates a kernel thread
- More concurrency than many-to-one
- Number of threads per process sometimes restricted due to overhead
- Examples
 - Windows
 - Linux

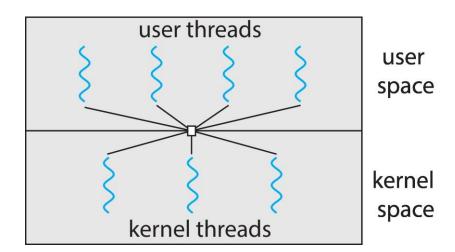






Many-to-Many Model

- Allows many user level threads to be mapped to many kernel threads
- Allows the operating system to create a sufficient number of kernel threads
- Windows with the *ThreadFiber* package
- Otherwise not very common

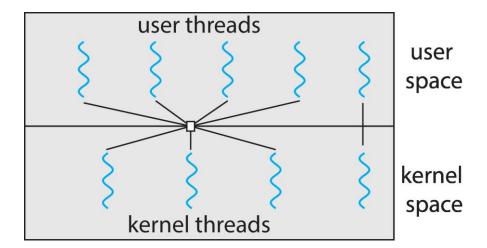






Two-level Model

 Similar to M:M, except that it allows a user thread to be bound to kernel thread







Pthreads

- May be provided either as user-level or kernel-level
- A POSIX standard API for thread creation and synchronization
- Specification, not implementation
 - API specifies behavior of the thread library, implementation is up to development of the library
- Example: Sum of N natural numbers
- Global data: Any variable declared globally are shared among all threads of the same process
- Local data: Data local to a function (running in a thread) are stored in thread stack





#include<stdio.h>
#include<pthread.h>





#include<stdio.h>
#include<pthread.h>

int sum; // global variable shared over threads

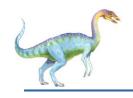




#include<stdio.h>
#include<pthread.h>

int sum; // global variable shared over threads void *runner (void *param); // threads begin execution in a specified function





```
#include<stdio.h>
#include<pthread.h>
```

```
int sum; // global variable shared over threads void *runner (void *param); // threads begin execution in a specified function int main(int argc, char *argv[]){
```





```
#include<stdio.h>
#include<pthread.h>

int sum; // global variable shared over threads

void *runner (void *param); // threads begin execution in a specified function
int main(int argc, char *argv[]){
    pthread_t tid; \\ declares the identifier for the thread
```

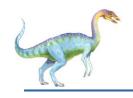




```
#include<stdio.h>
#include<pthread.h>

int sum; // global variable shared over threads
void *runner (void *param); // threads begin execution in a specified function
int main(int argc, char *argv[]){
    pthread_t tid; \\ declares the identifier for the thread
    pthread attr t attr; \\ set of thread attributes
```





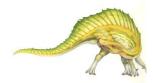
```
#include<stdio.h>
#include<pthread.h>

int sum; // global variable shared over threads
void *runner (void *param); // threads begin execution in a specified function
int main(int argc, char *argv[]){
    pthread_t tid; \\ declares the identifier for the thread
    pthread attr t attr; \\ set of thread attributes
    pthread attr init(&attr); \\ set the default attributes of the thread
```





```
#include<stdio.h>
#include<pthread.h>
int sum; // global variable shared over threads
void *runner (void *param); // threads begin execution in a specified function
int main(int argc, char *argv[]){
          pthread_t tid; \\ declares the identifier for the thread
          pthread attr t attr; \\ set of thread attributes
          pthread attr init(&attr); \\ set the default attributes of the thread
          pthread create(&tid, &attr, runner, argv[1]); \\ create the thread
```

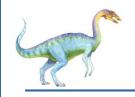




#include<stdio.h>

```
#include<pthread.h>
int sum; // global variable shared over threads
void *runner (void *param); // threads begin execution in a specified function
int main(int argc, char *argv[]){
          pthread_t tid; \\ declares the identifier for the thread
          pthread attr t attr; \\ set of thread attributes
          pthread attr init(&attr); \\ set the default attributes of the thread
          pthread create(&tid, &attr, runner, argv[1]); \\ create the thread
          pthread join(tid,NULL); \\ wait for the thread to exit
```





```
#include<stdio.h>
#include<pthread.h>
int sum; // global variable shared over threads
void *runner (void *param); // threads begin execution in a specified function
int main(int argc, char *argv[]){
         pthread_t tid; \\ declares the identifier for the thread
          pthread attr t attr; \\ set of thread attributes
         pthread attr init(&attr); \\ set the default attributes of the thread
          pthread create(&tid, &attr, runner, argv[1]); \\ create the thread
          pthread join(tid,NULL); \\ wait for the thread to exit
          printf("sum = %d\n",sum);
}
```



```
#include<stdio.h>
#include<pthread.h>
int sum; // global variable shared over threads
void *runner (void *param); // threads begin execution in a specified function
/* The thread will execute in this function */
void *runner(void *param) {
         int i, upper = atoi(param);
         sum = 0;
         for (i = 1; i \le upper; i++)
                  sum += i;
                   pthread exit(0); \\ thread terminates
}
```



- Growing dominance of multicore systems, writing programs containing several threads has become common.
- Simple method for waiting on several threads using the pthread_join()
 function is to enclose the operation within a simple for loop





OpenMP

- Collection of compiler directives and an API for C, C++, FORTRAN
- Provides support for parallel programming in sharedmemory environments
- Identifies parallel regions blocks of code that can run in parallel

#pragma omp parallel

 Create as many threads as there are cores

```
#include <omp.h>
#include <stdio.h>
int main(int argc, char *argv[])
  /* sequential code */
  #pragma omp parallel
    printf("I am a parallel region.");
  /* sequential code */
  return 0;
```



End of Chapter 4

