

"Computer Science is not really about computers and it's not about computers in the same sense that astronomy is not about telescope"

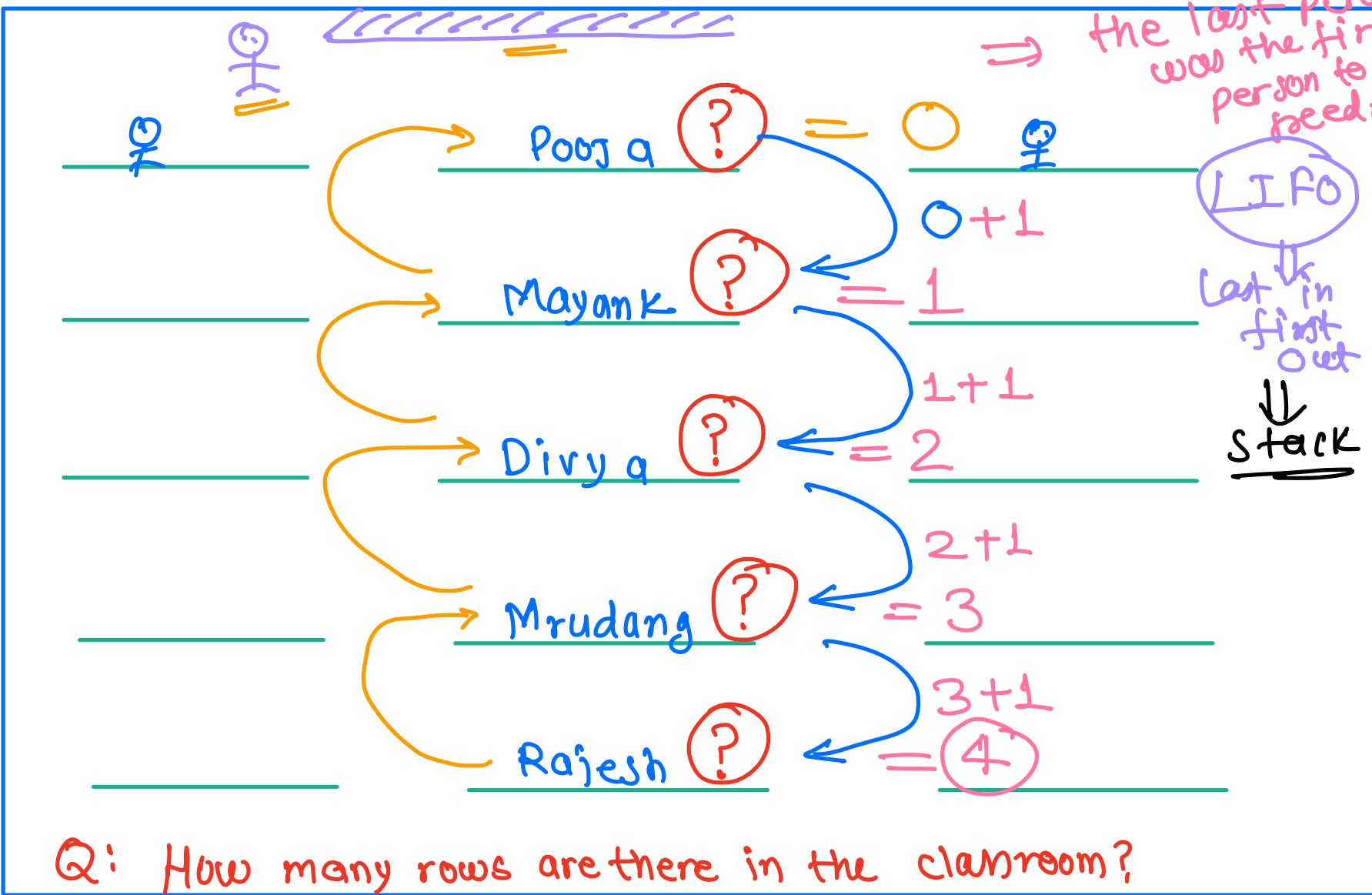
music: "In Motion" from The Social Network movie

### §. why we should recursion?

- ① Recursion makes it easy to solve certain types of problem
- ② A core foundation of problem solving (linked list, trees, graphs, DP)
- ③ Important for interviews.

→ function calling itself: recursive function

Recursion: Recursion is a problem solving method in which we solve a bigger problem by solving smaller instances of the exact same problem. (Reduce & Conquer)



Q: How many rows are there in the classroom?

= How many rows are in front of me  $\Rightarrow ? = 4$   
 $+ \frac{\text{my own row (he is at the last row)}}{1}$

$$= \underline{5}$$

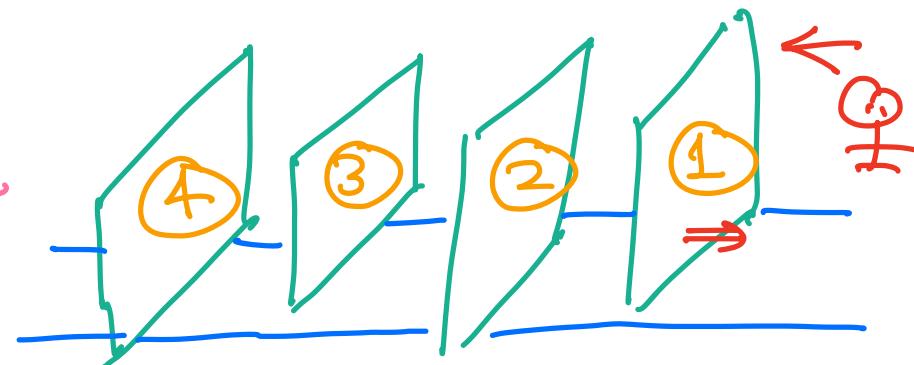
How many rows are in front of the person sitting in front of me + 1

### Russian doll video

Ex:

Mathematical induction:

Obs<sup>n</sup>: ① If we push plate #1, it will fall



② If a plate is falling then the plate next to it will also fall

$\Rightarrow$  If we push 1<sup>st</sup> plate, all the plates will fall.

Obs" I The more cheese you have,  
the more holes you have.

~~X~~ II More holes you have,  
less cheese you have



⇒ the more cheese you have ⇒ the less cheese you have

### S. Abstract thinking:

↳ ignore lower level details.

eg when we book a cab:

- ① Driver license
- ② fuel

hop in-hop out

Recursion is easy: ① If I have a sol<sup>n</sup> for  $f(n-1)$   
how do I get a sol<sup>n</sup> for  $f(n)$

② What is the base condition?

↓ Where we have trivial sol<sup>n</sup>.

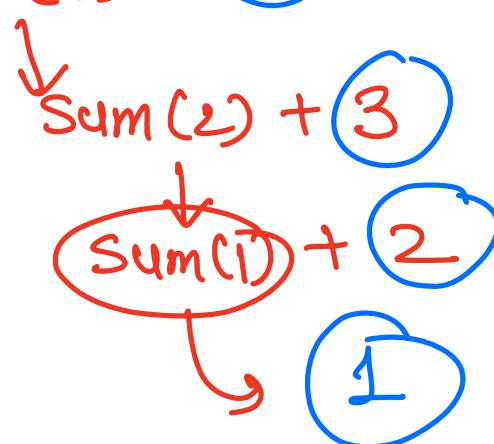
Q: Given a number  $N$ , find the sum of first  $N$  natural numbers.

$$\begin{aligned} \Rightarrow N = 1 &\Rightarrow 1 = 1 \\ \Rightarrow N = 2 &\Rightarrow 1, 2 = 3 \\ \Rightarrow N = 3 &\Rightarrow 1, 2, 3 = 6 \\ N = 4 & \qquad \qquad \qquad 6 + 4 = 10 \end{aligned}$$

$$\text{sum}(N+1) = \underline{\text{sum}(N)} + \underline{\underline{(N+1)}}$$

base condition :

$$\text{sum}(4) = \text{sum}(3) + 4$$



```

int sum (N) {
    if (N == 1) {
        return 1;
    }
    return sum(N-1) + N;
}

```

Q. Given a number  $N$ , find factorial of  $N$ .

Ex:

$$1! = 1$$

$$2! = 2 \times 1$$

$$3! = 3 \times 2 \times 1$$

$$4! = 4 \times 3 \times 2 \times 1$$

$$5! = 5 \times 4 \times 3 \times 2 \times 1 = 120$$

$$!0 = 1 = ?$$

$$(1 = 1 \times 0!)$$

A B C (ABC, ACB, BAC, BCA, CAB, CBA)

$$= 13$$

$$\underline{3} \times \underline{2} \times \underline{1}$$

$n!$  = ways to arrange N items.

⇒ 0 items in how many ways g can arrange them  
= 1

⇒  $\{1, 2, 3\} \Rightarrow [\cancel{\{1\}}, \{1\}, \{2\}, \{3\}, \dots]$

size 3 → 1

size 2 → 2

size 1 → 3

size 0 → 1

$$\text{fact}(N) = \text{fact}(N-1) * N$$

```
int fact(N){  
    if (N==0)  
        return 1;  
    }  
    return fact(N-1)*N;
```

}

N	ans
0	1
1	1
·	:

Q **function call tracing:** involves tracing the sequence of function calls that are made when a program is executed.

ex.

```
int add (x,y) {  
    return x+y;  
}  
  
int mul (x,y) {  
    return x*y;  
}  
  
int sub (x,y) {  
    return x-y;  
}  
  
void print (x) {  
    print (x);  
}
```

```

int main () {
    int x = 10
    int y = 20
    print (sub (mul (add (x,y), 30), 75));
    return 0; // successful
}

```

print (sub (mul (add (x,y), 30), 75)) ; 825 // on the console

$$\text{sub} (\underbrace{\text{mul} (\text{add} (\text{x}, \text{y}), 30)}_{\downarrow}, 75) = 825$$

$$\text{mul} (\underbrace{\text{add} (\text{x}, \text{y}), 30}_{\downarrow}, 30) = 900$$

$$\text{add} (\text{x}, \text{y}) = 30$$

$\Rightarrow$  Recursion uses stack to keep track of

smaller function calls value

## 5. function calls

```
int fact (N=5) {  
    if (N==0)  
        return 1;  
    }  
    return fact(N-1)*N; = 120  
}  
↓
```

```
int fact (N=4) {  
    if (N==0)  
        return 1;  
    }  
    return fact(N-1)*N; = 24  
}  
↓
```

```
int fact (N=3) {
```

```
    if (N==0)
```

$$\text{return } 1; \quad 2 * 3 = 6$$

```
}
```

```
return fact(N-1) * N;
```

```
}
```

```
int fact (N=2) {
```

```
    if (N==0)
```

$$\text{return } 1; \quad 1 * 2 = 2$$

```
}
```

```
return fact(N-1) * N;
```

```
}
```

```
int fact (N=1) {
```

```
    if (N==0)
```

$\dots$   
 $\vdots$   
 $\text{return } \cancel{\frac{\text{fact}(N-1) * N}{0}}$   
 $\}$

int fact (N=0){  
 $\quad \checkmark \text{if } (N==0)$        $\Rightarrow$   
 $\quad \text{return } 1;$   
 $\quad \text{return } \text{fact}(N-1) * N;$   
 $\}$

Obs<sup>n</sup>: In recursive functions, if base condition is not there or not possible to reach , then we will get MLE

why: infinite copies of function call.  
 $\downarrow$   
Stack overflow

## NCERT class XI mathematical induction:

(I)  $f(0)$  is true

(II) If  $f(n)$  is true then  $f(n+1)$  is also true;

⇒ we conclude  $f(n)$  is true for all

numbers  $(0, \infty)$

10:37

Q. Fibonacci numbers: The fibonacci sequence is a series of numbers in which each number is the sum of two previous numbers.

$$\begin{matrix} N = & 0, 1, 1, 2, 3, 5, 8, 13. \\ \text{index} & 0 \ 1 \ 2 \ 3 \ 4 \ 5 \ 6, 7 \end{matrix}$$

$$\underline{\underline{f(7)}} =$$

Sol<sup>n</sup>: ①  $f(n) = f(n-1) + f(n-2)$

$$\text{fib}(0) = 0$$

$$\text{fib}(1) = 1$$

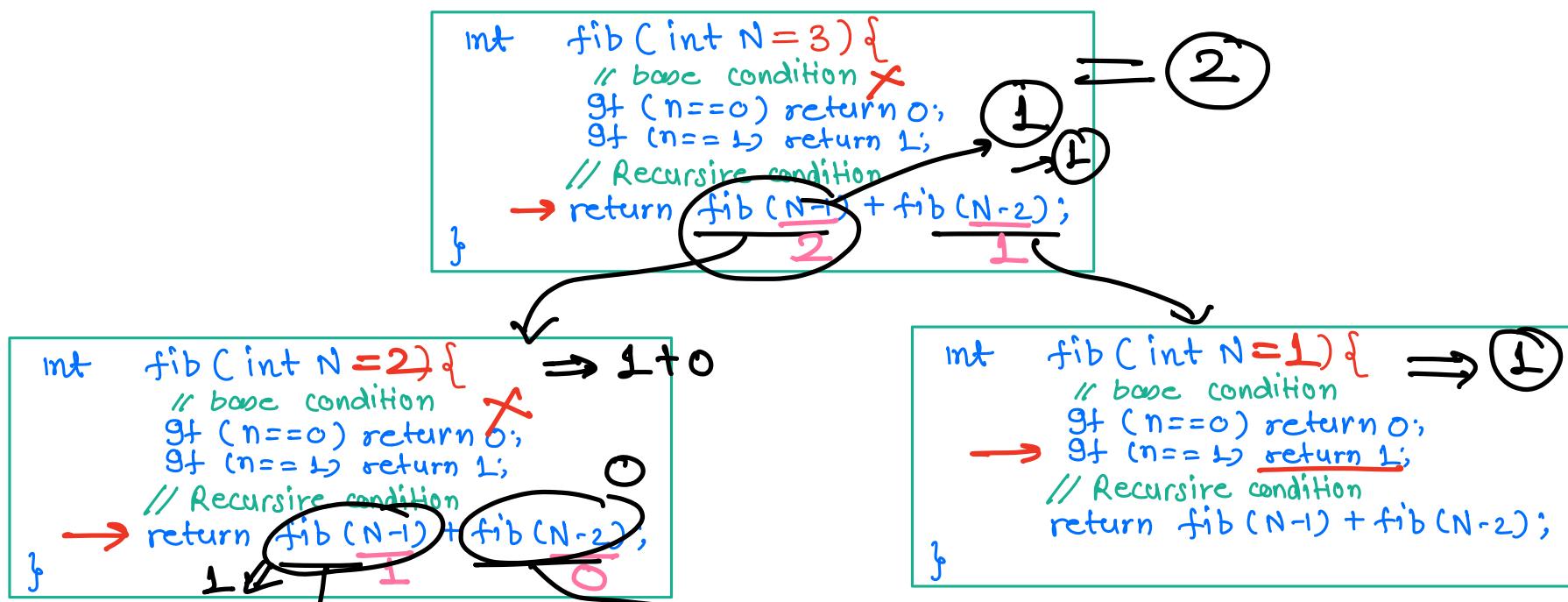
```

int fib(int N) {
    // base condition
    if (n==0) return 0;
    if (n==1) return 1;
    // Recursive condition
    return fib(N-1) + fib(N-2);
}

```

$$\begin{aligned}
f(2) &= f(0)+f(1) \\
f(3) &= f(2)+f(1) \\
f(4) &= f(3)+f(2)
\end{aligned}$$

function call tracking for  $\times \text{N}=3$



```

int fib(int N=1) {
    // base condition
    if (n==0) return 0;
    if (n==1) return 1;
    // Recursive condition
    return fib(N-1) + fib(N-2);
}

```

= 1

```

int fib(int N=0) {
    // base condition
    if (n==0) return 0;
    if (n==1) return 1;
    // Recursive condition
    return fib(N-1) + fib(N-2);
}

```

= 0

function call in the machine:



$$\begin{aligned}
 f(3) &= ? \\
 &= 1 + 1 \\
 &= f(2) + f(1); \\
 &\quad \text{---} \\
 &= f(1) + f(0) \\
 &= 1 + f(0) \\
 &= 1 + 0
 \end{aligned}$$

$f(2) = ?$   
 $f(1) = ?$   
 $f(0) = ?$

Q: Given a positive number  $N$ , write a recursive function to print all numbers from 1 to  $N$  in increasing order.

$N = 1 \Rightarrow 1$   
 $N = 2 \Rightarrow 1 \ 2$   
 $N = 3 \Rightarrow 1 \ 2 \ 3$   
 $N = 4 \Rightarrow 1 \ 2 \ 3 \ 4$   
 $N = 5 \Rightarrow 1 \ 2 \ 3 \ 4 \ \cancel{5}$

ex:  $N = 6 \Rightarrow 1 \ 2 \ 3 \ 4 \ 5 \ 6$

idea:

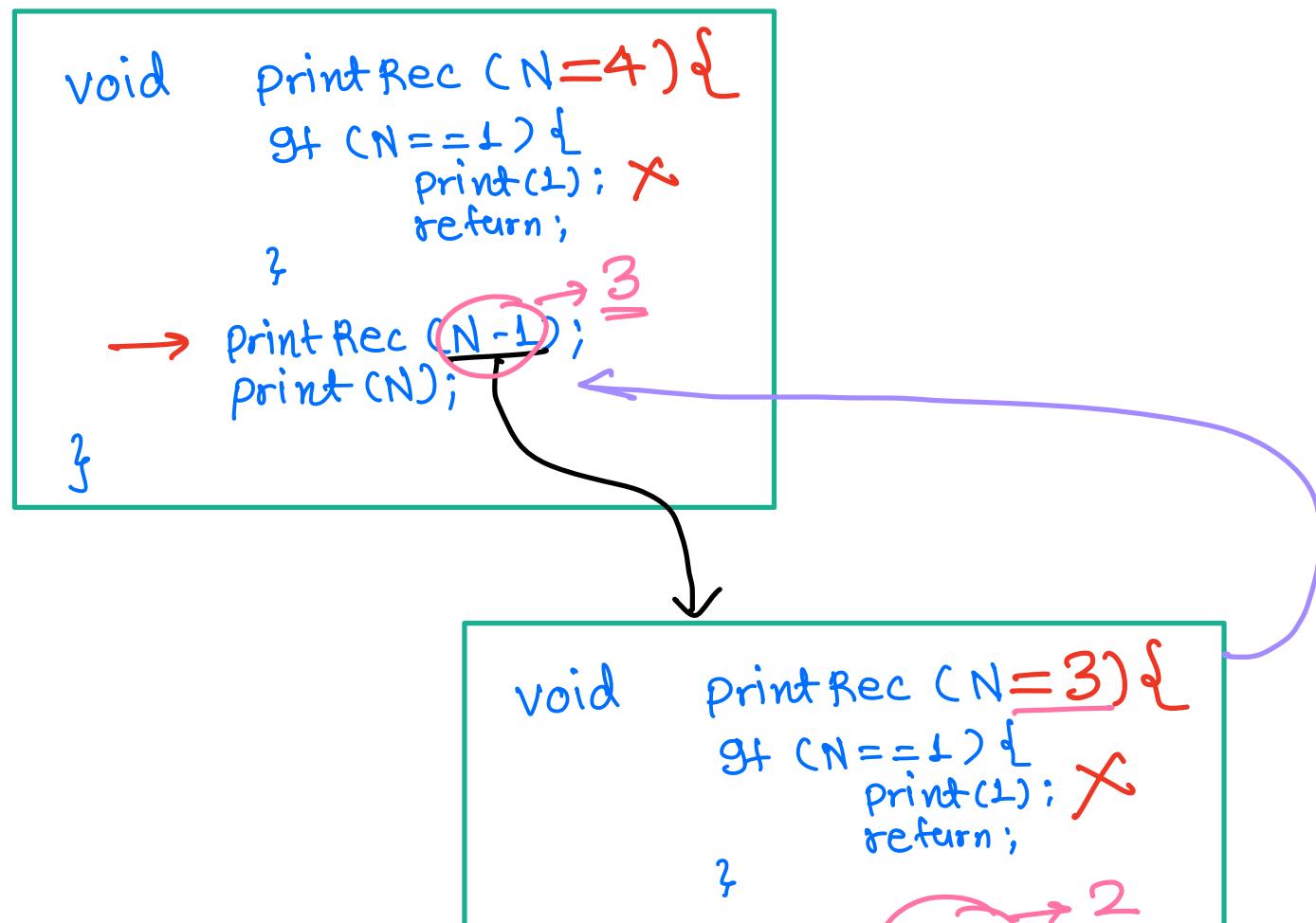
- (1) Start from  $N$
- (2) assume  $N-1$  numbers are already printed
- (3) print  $N$

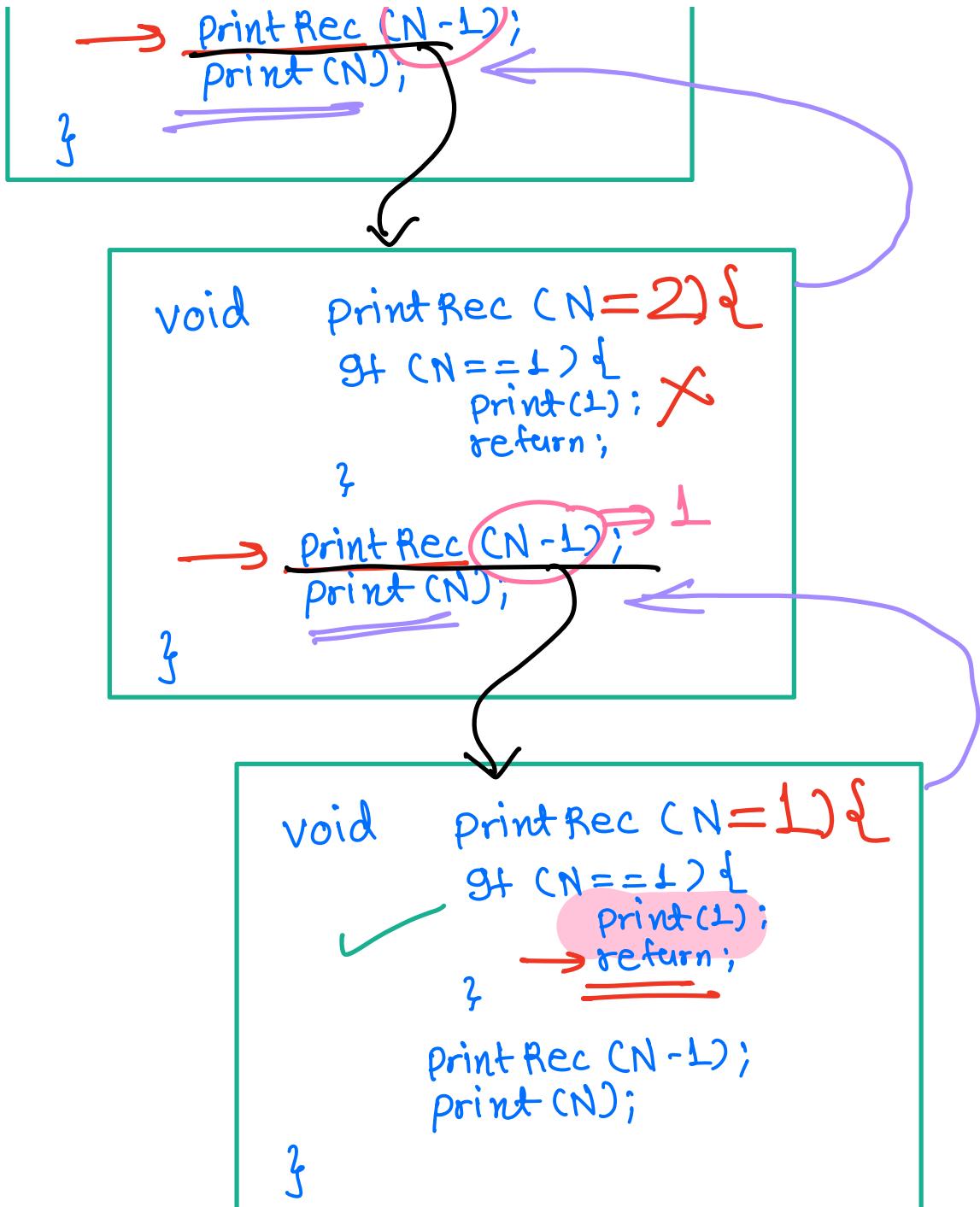
obsn: If we have the sol<sup>n</sup> for  $f(n-1)$   
 $\Rightarrow$  we have the sol<sup>n</sup> for  $f(n)$

void printRec(N) {  
 // Base condition  
 if ( $N == 1$ ) {  
 print(1);  
 return;  
 }  
 // Recursive condition  
 printRec(N-1); // By magic, it prints first
}

print(N);  
}' N-1 number as expected

Dry Run for  $N=4$





Console  $\gg \underline{1} \underline{2} \underline{3} \underline{4}$

H.W.  $\Rightarrow$  Given a positive number N, write a recursive function to print all numbers from 1 to N in decreasing order.

Q: Given a string , check if it is a palindrome or not.

ex:  $s = \text{'madam'}, ans = \text{true}$   
 $s = \text{'abcdx'}, ans = \text{false}$

quiz: A raor B deed C maddam D moon B&C  
X ✓ ✓ X

quiz: A radar B level C melon D hello  
✓ ✓ X X

Sol<sup>n</sup> using recursion:

Obj: How do you reduce the problem size?

ex:  $st = \text{"racecar"}$   $s=0, e=6, st+, e--$

Substring  
↓  
racecar

$st[s]$   
r

$st[e]$   
r

if  $st[s] == st[e]$   
✓

aceca

a

a

✓

cec

c

c

✓

e

e

e

✓

-  
 $st$  is a palindrome.

idea! ①  $st.$  size is  $N$

② check the start & end point  
if they are not equal return false

initially will be called  
with  $(\underline{st}, \underline{0}, \underline{N-1})$

else check smaller subbing is  
palindrome

```
bool isPalindrome (string st, int start, int end){  
    // Base condition  
    // If st is empty return true;    st = "ABC"  
    if (start > end) {  
        return true;  
    }  
  
    if (st[start] != st[end]) {  
        return false;  
    }  
  
    return isPalindrome ( st , start+1, end-1 )  
}
```

st =  $\underline{\underline{0}} \underline{\underline{1}} \underline{\underline{2}}$   
end =  $\underline{\underline{2}} \underline{\underline{1}} \underline{\underline{0}}$

Dry run ①  $st = " = "$      $N=0$      $(0, N-1)$     ✓

②

$st = "P"$

$N = 1$

Palindrome (0,0)

```
bool isPalindrome (string st, int start, int end){  
    if (start > end) {  
        return true;  
    }  
  
    if (st[start] != st[end]) {  
        return false;  
    }  
    return isPalindrome (st, start+1, end-1);  
}
```

1 -1

```
bool isPalindrome (string st, int start, int end){  
    if (start > end) {  
        return true;  
    }  
  
    if (st[start] != st[end]) {  
        return false;  
    }  
    return isPalindrome (st, start+1, end-1);  
}
```



st = "QQ"  $\Rightarrow$  isPalin(st, 0, 1) ✓

```
bool isPalindrome(string st, int start, int end){  
    if (start > end) {  
        return true; ✗  
    }  
  
    if (st[start] != st[end]) {  
        return false; ✗  
    }  
    }  
    return isPalindrome(st, start+1, end-1);  
}
```

Diagram showing the recursive call `isPalindrome(st, start+1, end-1)` with `start=1` and `end=0`. The result is labeled true.

```
bool isPalindrome(string st, int start, int end){  
    if (start > end) { ✓  
        return true; =====  
    }  
  
    if (st[start] != st[end]) {  
        return false;  
    }  
    }  
    return isPalindrome(st, start+1, end-1);  
}
```

Diagram showing the recursive call `isPalindrome(st, start+1, end-1)` with `start=1` and `end=0`. The result is labeled true.



st = "radar", N = 5, isPalind..(st, 0, 4)

```

bool isPalindrome (string st, int start, int end) {
    if (start > end) {
        return true;
    }

    if (st[start] != st[end]) {
        return false;
    }
    return isPalindrome (st, start+1, end-1);
}

```

4

false

```

bool isPalindrome (string st, int start, int end) {
    if (start > end) {
        return true;
    }

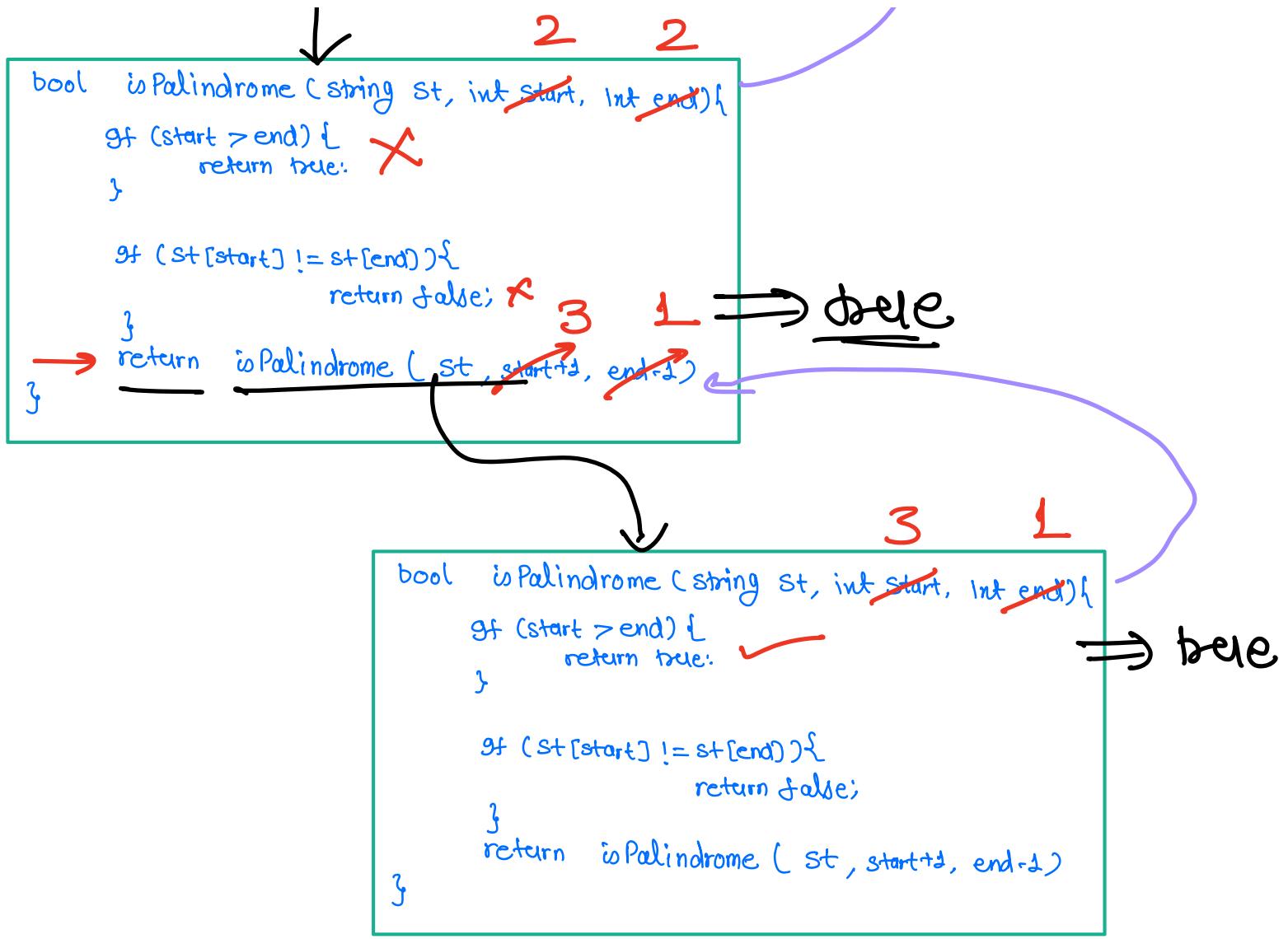
    if (st[start] != st[end]) {
        return false;
    }
    return isPalindrome (st, start+1, end-1);
}

```

1 3

1 3

false



① Office hours. 9pm tomorrow

② Mathematical induction in the class notes

Q:  $\frac{f(n-1)}{\cancel{f(n-1)}} \xrightarrow{\text{now?}} \frac{f(n)}{\cancel{f(n)}}$

$O(1)$   $\Rightarrow$  analysis is required

$\uparrow$   
most cases