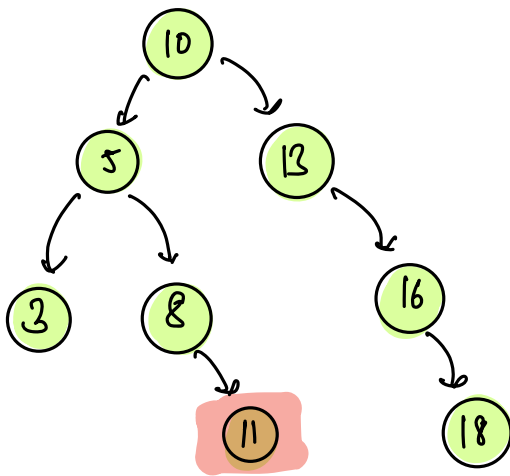
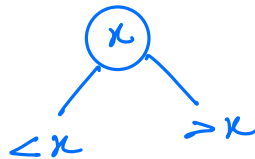


Binary Search Trees (B.S.T)

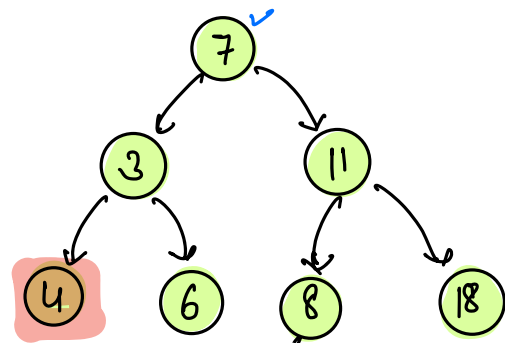
⇓
Binary Tree

*nodes

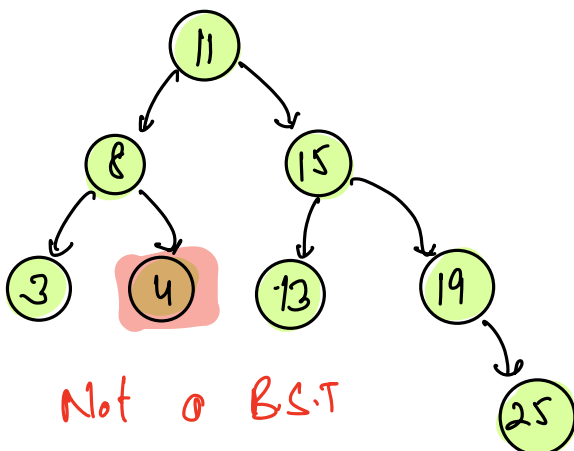
[all values in the l.s.t < node.data < all values in the r.s.t]



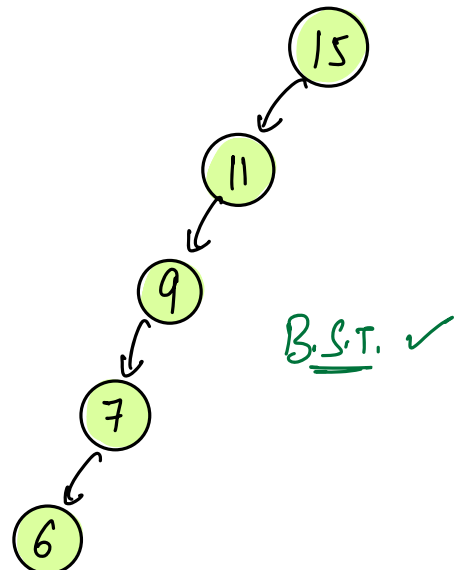
Not a B.S.T



Not a B.S.T



Not a B.S.T



B.S.T. ✓


```
boolean search (Node root, int k){
```

```
    Node temp = root;
```

```
    while (temp != Null){
```

```
        if (temp.val == k){
```

```
            return true;
```

```
        } else if (temp.val < k){
```

```
            temp = temp.right;
```

```
        } else if
```

```
            temp = temp.left;
```

```
    }
```

```
    return false;
```

```
}
```

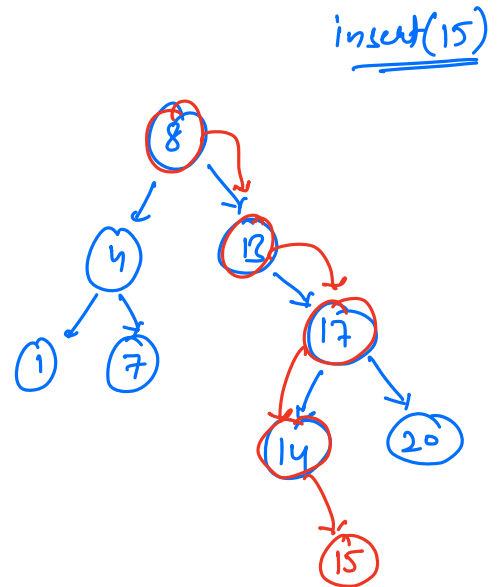
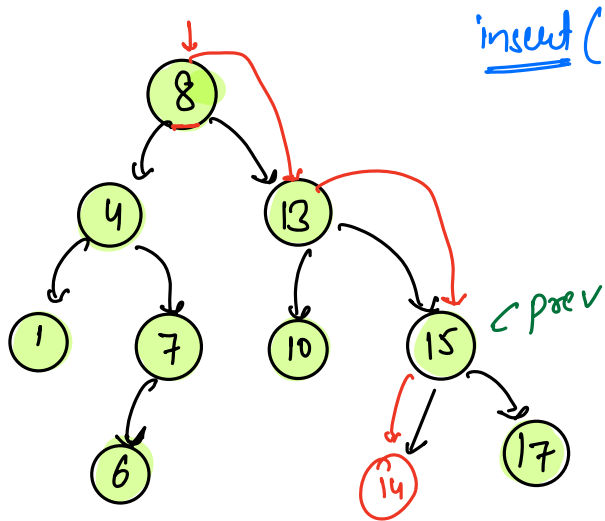
balanced.

$\log_2 N$

skewed
 N

$\left[\begin{array}{l} \text{T.C} \rightarrow O(H) \\ \text{S.C} \rightarrow O(1) \end{array} \right]$

Q1 Insert an element x in B.S.T.



#code:-

Node temp = root, prev = NULL;

while(temp != NULL){

prev = temp;

if (temp.val == x){

{ return;

else if (temp.val < x){

{ temp = temp.right;

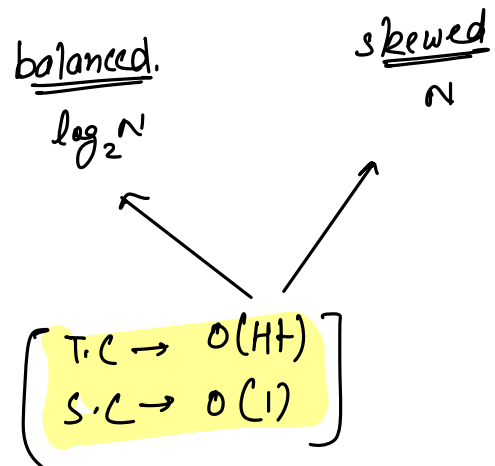
else { temp = temp.left;

}

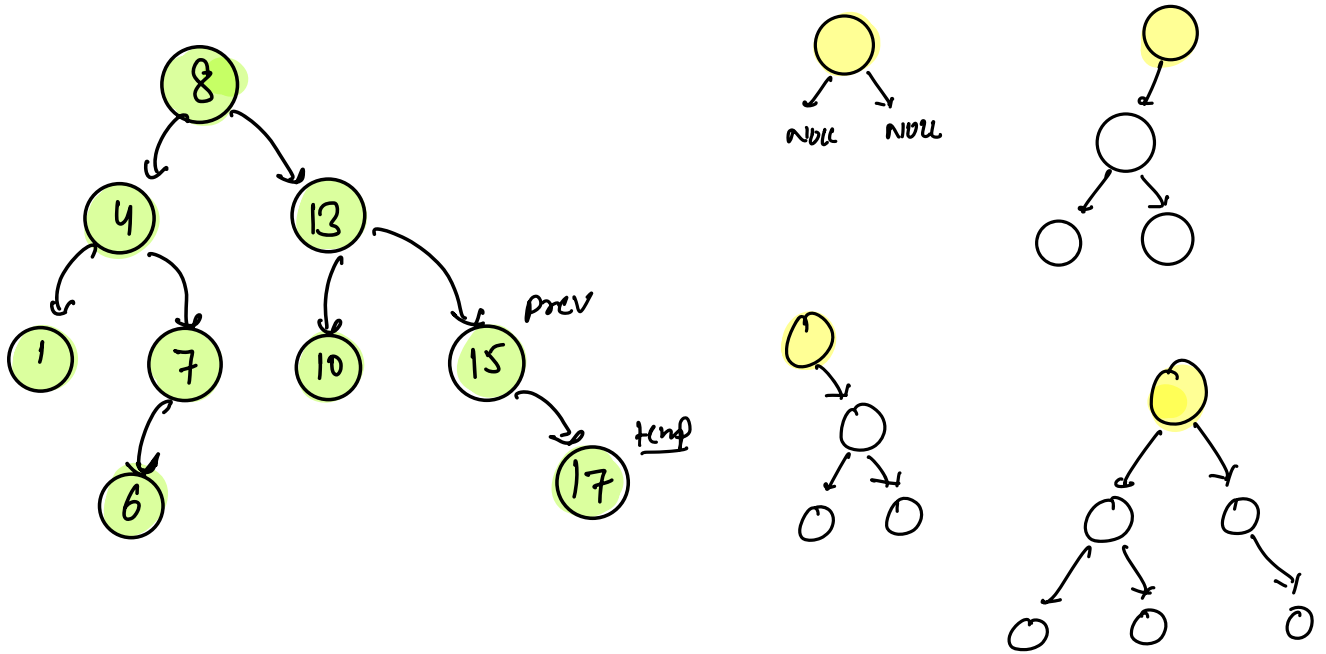
if(prev == NULL){ root = new Node(x), return }

if (x < prev.val) { prev.left = new Node(x) }

else { prev.right = new Node(x) }



Deletion of a Node from B&T.



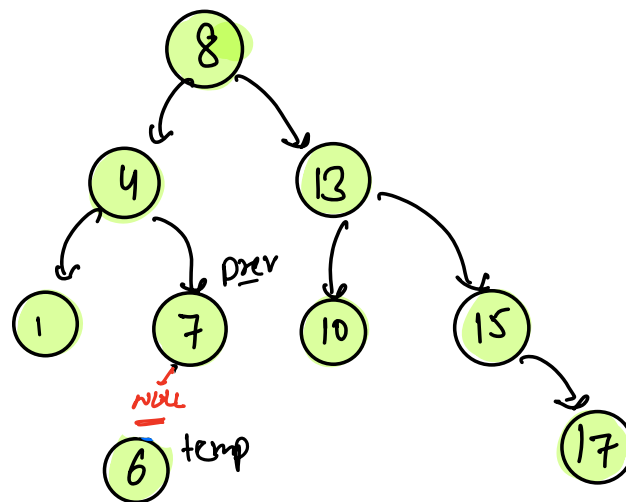
→ search(x) by keeping track of prev node.

- ① if (temp.left == NULL && temp.right == NULL){
 // temp is a leaf node

- ② else if (temp.left == NULL || temp.right == NULL){
 // temp with single child

- ③ else if
 // temp with both the children

Case-1 When temp is leaf node

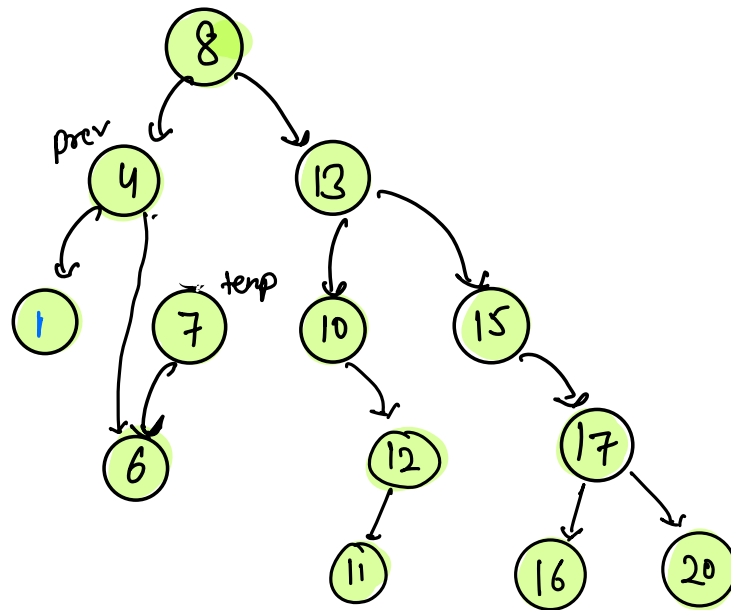


remove(6)

```
if (prev.left == temp) {  
    {  
        prev.left = NULL;  
    }  
else {  
    {  
        prev.right = NULL;  
    }  
}
```

Case-2:

temp with single child



remove(7)

if (temp.left != null) {

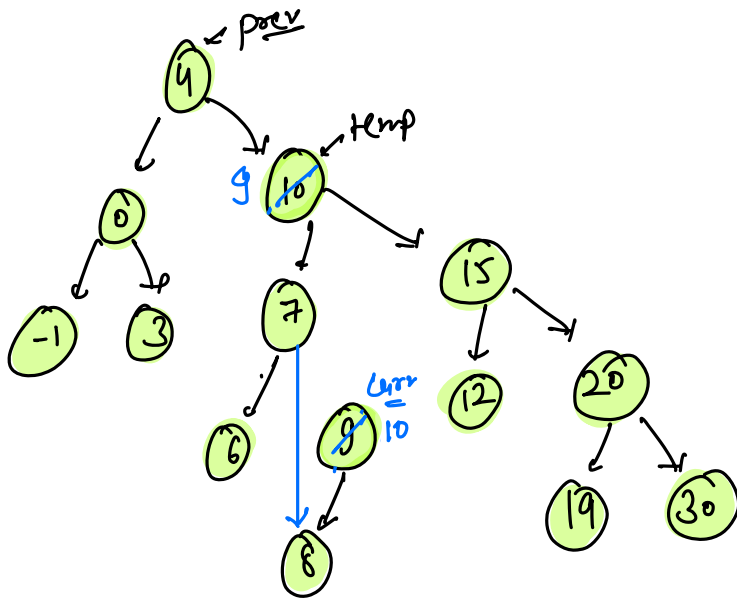
```
    if ( prev.left == temp ) {  
        {  
            prev.left = temp.left;  
        }  
    }  
    else {  
        {  
            prev.right = temp.left;  
        }  
    }
```

}
else if (temp.right != null) {

```
    if ( prev.left == temp ) {  
        {  
            prev.left = temp.right;  
        }  
    }  
    else {  
        {  
            prev.right = temp.right;  
        }  
    }
```

}

Case-2. temp with both the children



remove(10)

$\left[\text{All nodes in l.s.t} < \text{node.data} < \text{nodes in r.s.t} \right]$

\Downarrow

$\left[\text{Max in l.s.t} < \text{node.data} < \text{Min in r.s.t} \right]$

→ Max of l.s.t.

Node curr = temp.left;

if (curr.right != NULL) {
 $\left[\right.$ curr = curr.right;
 $\left. \right]$

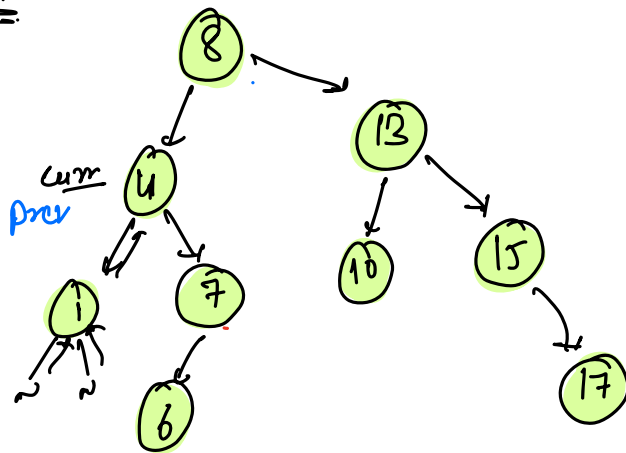
swap(temp.value with curr.value);

// Make a recursive call to remove 10 from temp.left

remove(temp.left, K);

$\left[\begin{array}{l} \text{T.C} \rightarrow O(H) \\ \text{s.c} \rightarrow O(1) \end{array} \right]$

Validate Bst.



[1, 4, 6, 7, 8, 10, 12, 15, 17]

idea-1 Do inorder traversal \rightarrow Array is sorted \rightarrow Yes.

\downarrow
[while doing traversal]
prev, curr

$\left[\begin{array}{l} \text{T.C} \rightarrow O(N) \\ \text{S.C} \rightarrow O(N) \end{array} \right]$

Node prev = NULL, boolean isBST = true;

```
void traversal( Node curr){
```

```
    if( curr == NULL ) { return; }
```

```
    traversal( curr->left );
```

```
    if( prev != NULL && prev->val >= curr->val ){
```

```
        { isBST = false; }
```

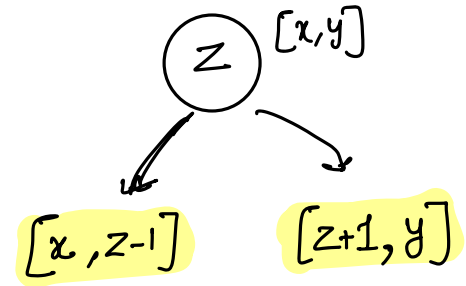
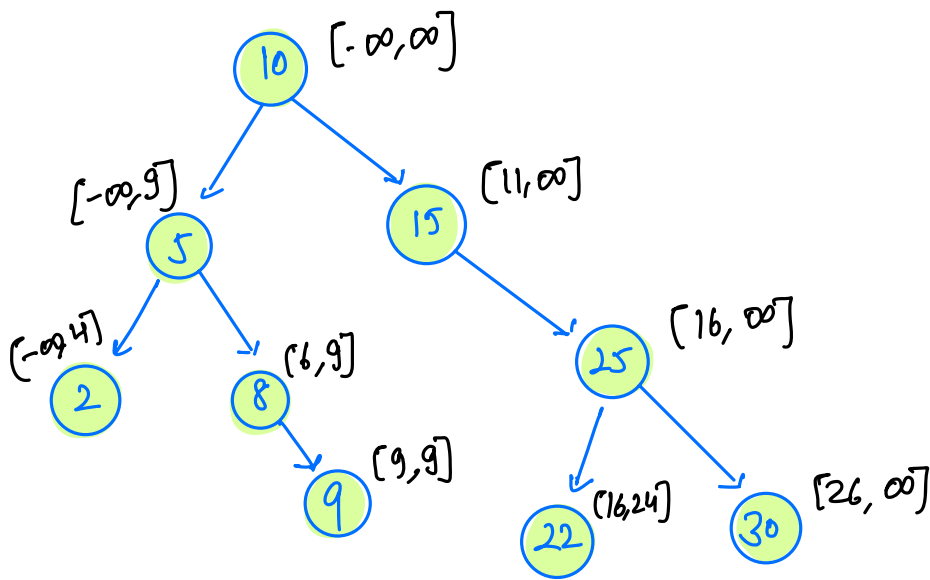
```
        prev = curr;
```

```
    traversal( curr->right );
```

```
}
```

$\left[\begin{array}{l} \text{T.C} \rightarrow O(N) \\ \text{S.C} \rightarrow O(h) \end{array} \right]$

② Using pre-order →



#code →

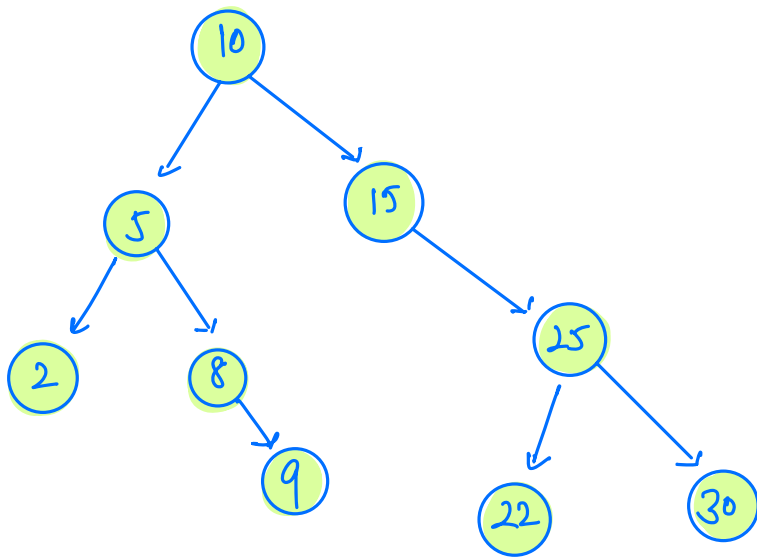
LONG.MIN-VALUE LONG.MAX-VALUE.

```

boolean isBst ( root,  $\downarrow$  l,  $\downarrow$  r ) {
    if ( root == null ) { return true }
    if ( root.val ≥ l && root.val ≤ r ) {
        boolean f1 = isBst ( root.left, l, root.val - 1 );
        boolean f2 = isBst ( root.right, root.val + 1, r );
        return ( f1 && f2 );
    }
    return false;
}

```

$T.C \rightarrow O(N)$
 $S.C \rightarrow O(H+)$



Context → [L.L, Stack, Queue, Trees]