

# Congratulations for maintaining >80% PSP 🎆🥳

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Today's Content →

Introduction to Graphs

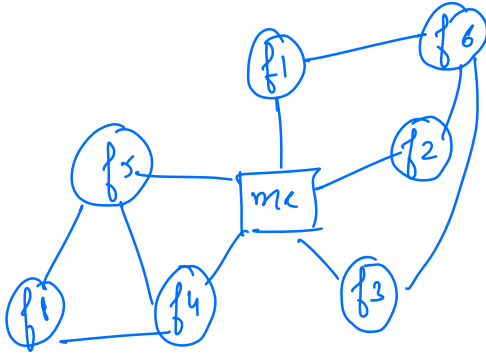
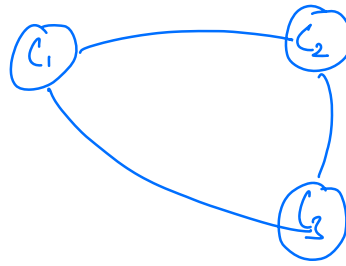
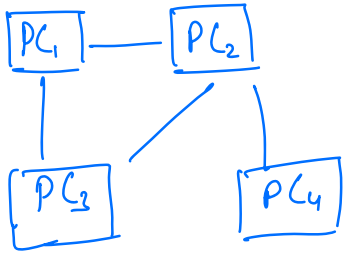
Types of Graphs

DFS

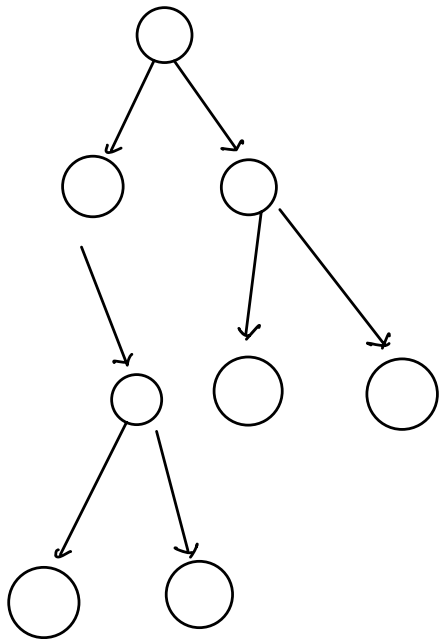
BFS

Detect cycle in directed graph.

Graph  $\rightarrow$  network.



graph  $\rightarrow$  collection of nodes or vertices and edges.



Every tree is a graph. ✓

Every graph is a tree. ✗

$\Downarrow$

[Tree is a subset  
of Graph]

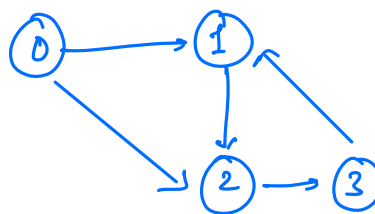
① Tree always has root node.

②  $N$  nodes  $\rightarrow N-1$  edges.

③ Cycle can't be there in tree.

# How to store graphs?

## ① Adjacency Matrix



4 Nodes (N)  
5 edges (E)

	0	1	2	3
0	0	1	1	0
1	0	0	1	0
2	0	0	0	1
3	0	1	0	0

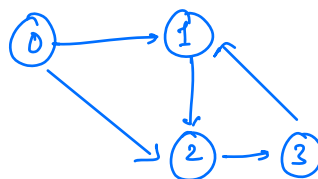
$mat[i][j] = 0 \Rightarrow$  There is no edge from  $i$  to  $j$ .

$mat[i][j] = 1 \Rightarrow$  There is an edge from  $i$  to  $j$ .

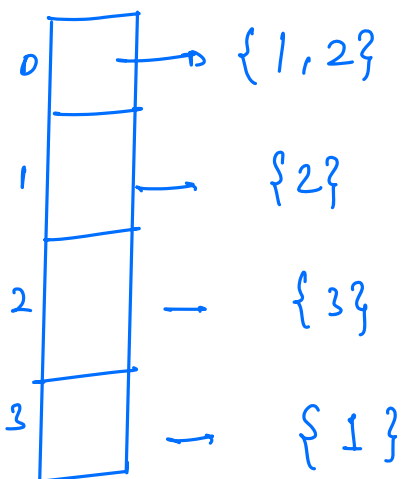
$$SC \rightarrow O(N^2)$$

## ② Adjacency List

(Array of lists)



$N = 4$   
 $E = 5$



$A[i] \rightarrow$  adjacent nodes in the list  
↳ neighbours of  $i$  in the list.

$$\begin{bmatrix} 1 \leq N \leq 10^5 \\ 1 \leq E \leq 10^5 \end{bmatrix}$$

$$SC \rightarrow O(N + E)$$

# Properties / Types of Graph

①

Directed

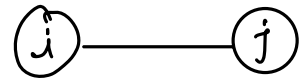


i to j ✓

j to i ✗

(Bi-directional)

Un-directed Graph

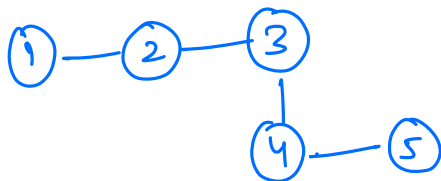


i to j ✓

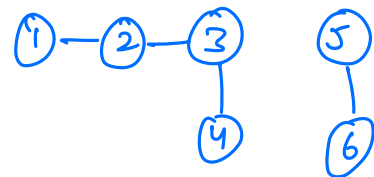
j to i ✓

②

Connected



Disconnected



③

Weighted



Un-weighted



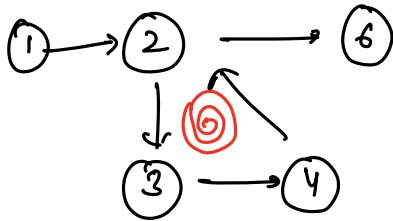
$mat[i][j] = 0$  if there is no edge from i to j  
 $= w_{ij}$  if there is edge from i to j.

$graph[i] = \{ \{1, 5\}, \{2, 3\}, \{4, 7\} \dots \}$   
list of pairs

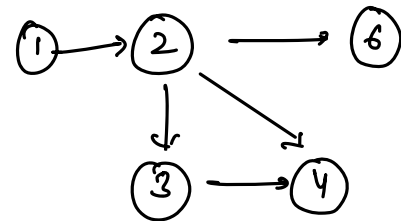
pair {  
vtx :  
edge wt;  
}

(4)

Cyclic.

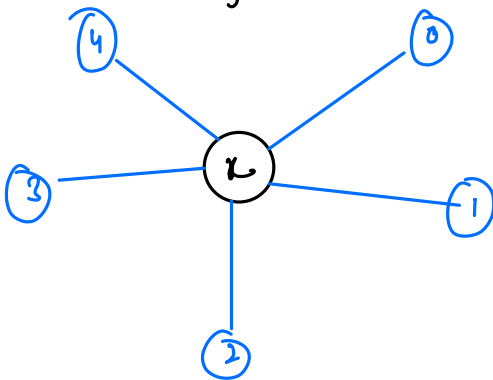


Acyclic



(5)

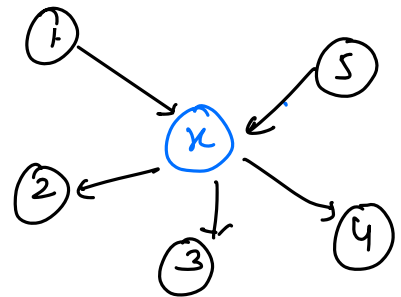
Degree.



$$\text{degree}(x) = 5$$

No. of edges connected with a node.

In-degree / out-degree



in-degree(x)  $\rightarrow$  incoming edges [2]

out-degree(x)  $\rightarrow$  outgoing edges [3]

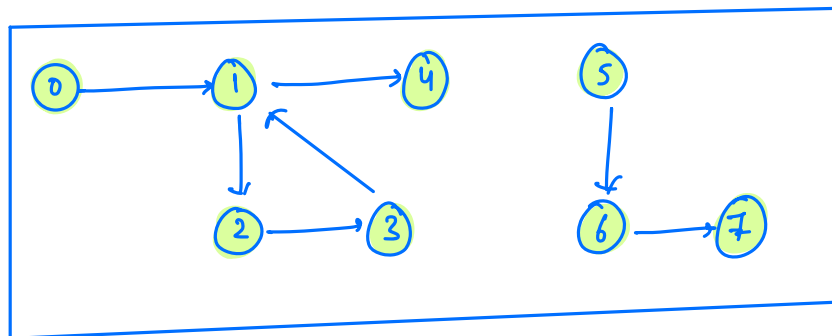
---

N nodes.  $\rightarrow$  numbered from 0 to  $N-1$

if they are not numbered from 0 to  $N-1$ , then we will do the mapping and get the order from 0 to  $N-1$ .

# Traversals

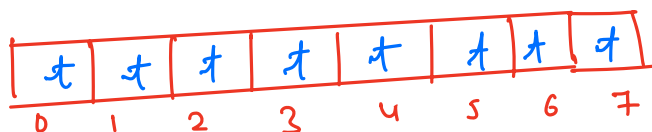
## ① Depth First Traversal (Pre-order traversal of tree)



$N=8$ ,  $E=7$

→ Keep track of visited nodes.

visited →



# code →

// Graph → given

boolean visited[N]; //  $\forall i$ ,  $visited[i] = false$

for (  $i = 0$ ;  $i < N$ ;  $i++$  ) {

if (  $visited[i] == false$  ) {

dfs ( graph, i, visited );

}

}

```
void dfs ( Graph, src, visited[N]) {
```

```
    print (src);
```

```
    visited[src] = true;
```

```
    for ( int nbr : graph[src]) {
```

```
        if ( visited[nbr] == false ) {
```

```
            {
                dfs ( graph, nbr, visited );
            }
        }
    }
```

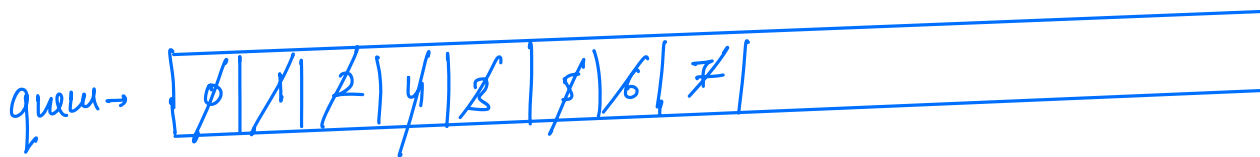
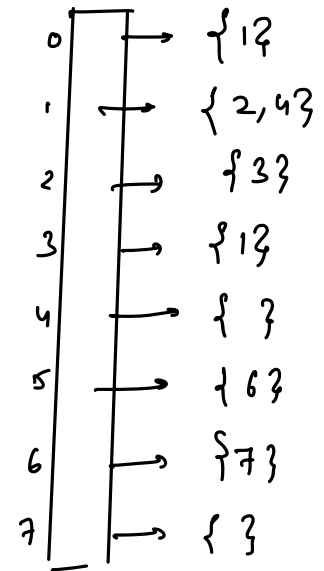
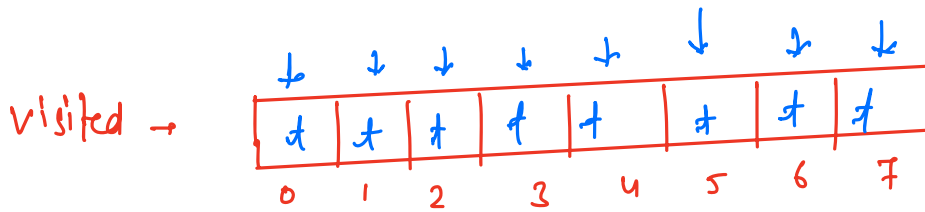
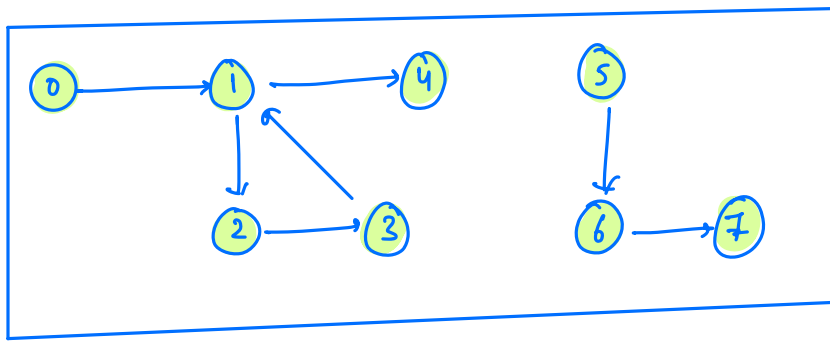
```
}
```

$\left( \begin{array}{l} T.C \rightarrow O(N+E) \\ S.C \rightarrow O(N) \end{array} \right)$

↓

visited[N] +  
max size of stack.

## Breadth first Traversal → level order traversal.



print → 0, 1, 2, 4, 3, 5, 6, 7

# code →

// Graph → given

bool can visited[N]; //  $\forall i, \text{visited}(i) = \text{false}$

for ( i = 0; i < N; i++ ) {

    if ( visited[i] == false ) {

        bfs ( graph, i, visited );

    }

}



```
void bfs ( graph, src, visited) {
```

```
    Queue < int > q;
```

```
    q.enqueue(src) , visited[src] = true;
```

```
    print(src);
```

```
    while ( q.isEmpty() == false) {
```

```
        rv = q.dequeue();
```

```
        for ( int nbr : graph[rv] ) {
```

```
            if ( visited[nbr] == false ) {
```

```
                visited[nbr] = true; print(nbr);
```

```
                q.enqueue(nbr);
```

```
            }
```

```
        }
```

```
    }
```

```
}
```

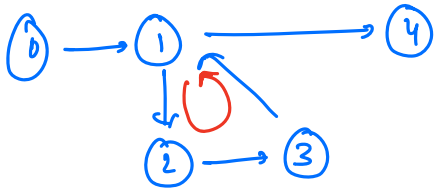
$T.C \rightarrow O(N+E)$   
 $S.C \rightarrow O(N)$

↓

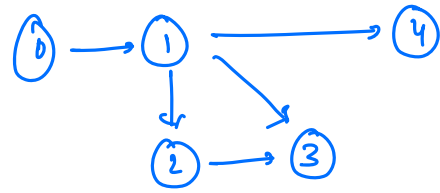
visited[N],

Queue

Q. Check if given directed graph has a cycle or not.



ans  $\rightarrow$  true.



ans  $\rightarrow$  false

$\rightarrow$  if a visited node is encountered again  $\rightarrow$  cycle. X

$\rightarrow$  if a visited node in current path is encountered again  $\rightarrow$  cycle ✓

#code  $\rightarrow$

boolean visited[N], //  $\forall i$ , visited[i] = false;

boolean path[N], //  $\forall i$ , path[i] = false;

for( i = 0; i < N; i++) {

if( visited[i] == false) {

if ( dfs( graph, i, visited, path) == true) {

return true;

}

}

return false;

```
boolean dfs ( Graph , src , visited[N] , path[N] ) {
```

```
    visited[src] = true;
```

```
    path[src] = true;
```

```
    for ( int nbr : graph[src] ) {
```

```
        if ( path[nbr] == true ) { //if nbr is already present  
            //in current path  
            return true;
```

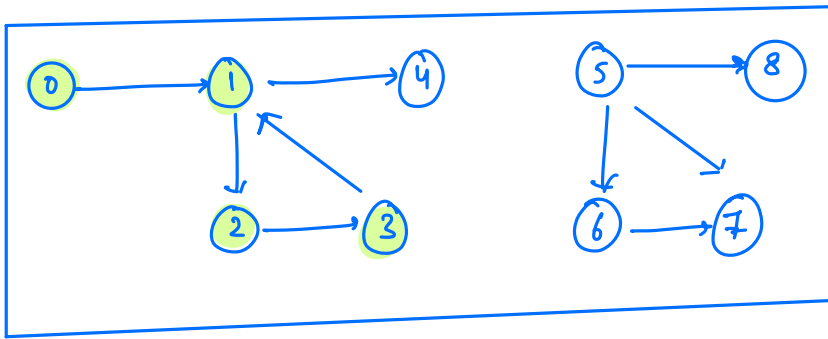
```
        if ( visited[nbr] == false && dfs ( graph , nbr , visited , path ) ) {  
            return true;
```

```
    }
```

```
    path[src] = false; //removing src from curr path before  
    return false;      //returning from src.
```

```
}
```

T.C  $\rightarrow O(N+E)$   
S.C  $\rightarrow O(N)$

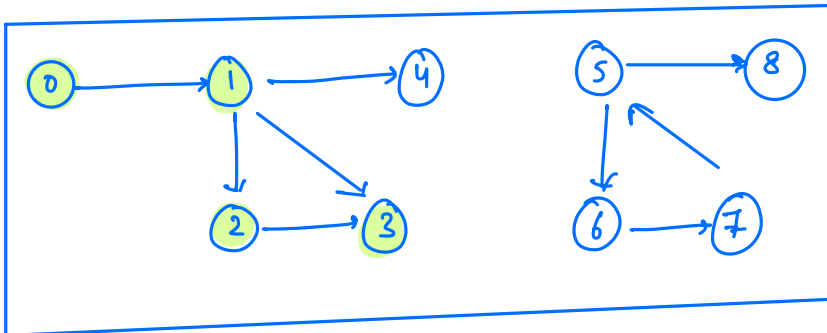
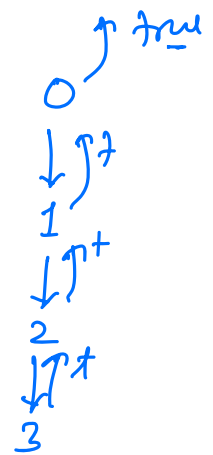


visited [ ] →

↓										
	t	t	t	t	f	f	f	f	f	f
	0	1	2	3	4	5	6	7	8	9

path [ ] →

	t	t	t	t	f	f	f	f	f	f
--	---	---	---	---	---	---	---	---	---	---



60-70%  $\rightarrow$  BFS, DFS