

Sorting

"Sorting is the foundation of organization, and organization is the key to efficiency" - Anonymous

Sorting? arrangement of data, in particular order, on
the basis of some rule

2 3 9 12 17 19

sorted in asc order

19 6 5 2 -1 -9

sorted in dsc order

1 13 9 6 12

Sorted based on how many factors they have in increasing order.

1 2 3 4 6

Sort → Price.

arg :

advanced

How ~~to~~ sort

count of factors

why sorting?

- searching made easy

⇒ inbuilt sort function

sort()

T.C. $O(n \log n)$

S.C. $\underline{O(N)}$

//in advanced classes

✗ T.C. = $O(N^2)$



would sorting help to reduce the time complexity?

Q. Given an array of N integers, minimise the cost to empty the array, where cost of removing an element is equal to sum of all elements left in the array.

ex. $A = [2, 1, \cancel{4}]$, ans = 11

Step I remove 4 $\Rightarrow 4+2+1 = 7$

II. remove 2 $\Rightarrow 2+1 = 3$

III. remove 1 $\Rightarrow 1 = 1$

total cost = 11.

another way

$$A = [\cancel{2}, 1, 4]$$

(I) remove 2 $\Rightarrow 2+1+4 = 7$

(#) remove 4 $\Rightarrow 1+4 = 5$

(##) remove 1 $\Rightarrow 1 = 1$

total cost = 13

Quiz

$$A = [\cancel{4}, \cancel{5}, 1]$$

(I)
(II)
(##)

remove 6 $\Rightarrow 4+6+1 = 11$

remove 4 $\Rightarrow 4+1 = 5$

remove 1 $\Rightarrow 1$

17

another way

$$A = [\cancel{4}, \cancel{5}, 1]$$

(I) remove 4 $\Rightarrow 4+6+1 = 11$

(#) remove 6 $\Rightarrow 6+1 = 7$

(##) remove 1 $\Rightarrow 1 = 1$

19

not minimum

Quiz: $A = [\cancel{3}, \cancel{5}, \cancel{-3}]$

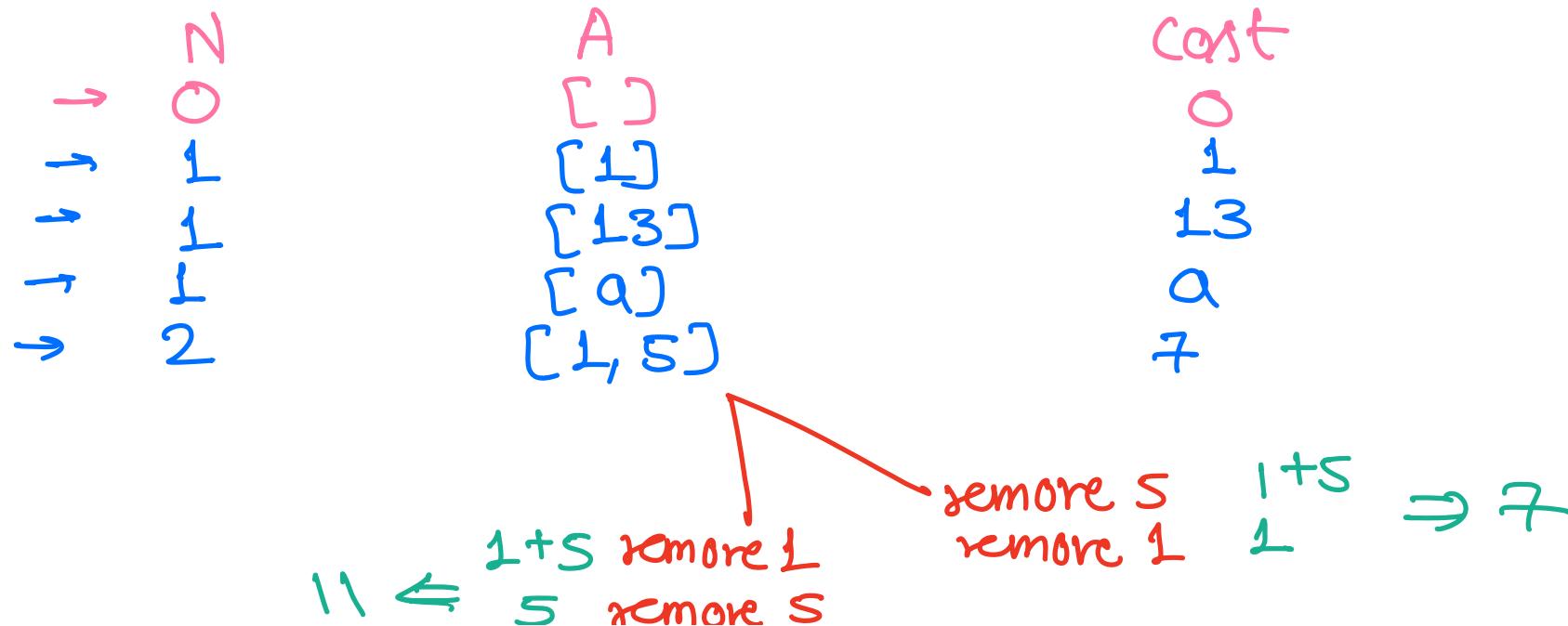
$$\begin{array}{lll}
 \textcircled{I} & \text{remove } 5 \Rightarrow 3+5+1+(-3) = 6 \\
 \textcircled{\#} & \text{remove } 3 \Rightarrow 3+1+(-3) = 1 \\
 \textcircled{III} & \text{remove } 1 \Rightarrow 1+(-3) = -2 \\
 \textcircled{IV} & \text{remove } -3 \Rightarrow -3 = -3
 \end{array}$$

~~2~~

Brute force: try all the possibilities $\Rightarrow n!$ ~~X~~

Possibilities \nwarrow not a good idea

S. optimal solⁿ:



$\Rightarrow 2$

$[10, 30]$

50

$\Rightarrow 2$

$[a, b]$

$$\begin{array}{ccc}
 & \xrightarrow{x a} & \xrightarrow{x b} \\
 [a, b] & \downarrow & \downarrow \\
 a+2b & & 2a+b
 \end{array}$$

min.

if $a < b \Rightarrow 2a+b$ would be smaller than $a+2b$.

$\Rightarrow [a, b, c, d]$

remove a

\Rightarrow

$a+b+c+d$

remove b

\Rightarrow

$b+c+d$

remove c

\Rightarrow

$c+d$

remove d

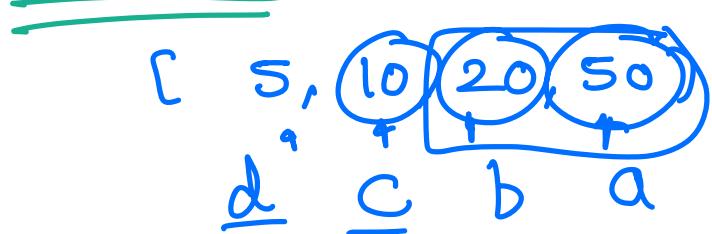
\Rightarrow

d

total cost $\Rightarrow \overbrace{a+2b} + 3c + 4d$

11

4 numbers minimise this



$\Rightarrow [50, 20, 10, 5]$

idea: ① Sort the array (in dsc. order)

$$\textcircled{2} \quad \text{cost} = \text{const} + A[i]^*c[i+1]$$

```
int minCost ( A ) { A = [2, 1, 4]
```

```
int N = A.size(); N = 3 ~\downarrow  
reverse_sort (A); // O(N log N) A = [4, 2, 1]
```

```

int cost = 0;
for (i=0; i< N; i++) {
    cost += A[i] * (i+1)
}
return cost;

```

~~cost = $\sum_{i=0}^{N-1} i \cdot A[i]$~~

~~i = 0, 1, 2, ..., N-1~~

T.C. = $O(N \log N)$
S.C. = $O(N)$

Q. find count of noble integers in the array of
(count of) distinct elements
A[i] is noble, if ^ elements smaller than A[i] is
equal to A[i],

Ex: $A = [1, -5, 3, 5, -10, 4]$

0	1	2	3	4	5
↓	↓	↓	↓	↓	↓
2	1	3	5	0	4
X	X	✓	✗	X	✗

Count = 3

Quiz: $A = [-3, 0, 2, 5]$

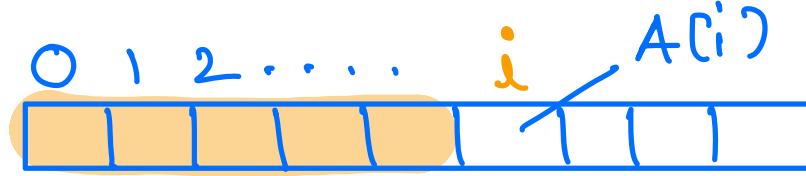


Brute force:

- ① for every element count smaller element
 - ② gf $\text{countSmallerElements} == A[i]$
ans++
- T.C. = $O(N^2)$
S.C. = $O(1)$

Optimal solⁿ:

Sort the data, asc. order



\downarrow smaller than $A[i]$

$$\text{count} = \underbrace{(0 \text{ to } i-1)}_{\downarrow (i-1-0+1)}$$

(L, R)
 $(R-L+1)$

$$= \underline{\underline{i}}$$

Obsⁿ: gf $i == A[i] \Rightarrow A[i]$ is noble,

idea: ① Sort the array
② count for $i == A[i]$ for all $i(s)$

```
// Sort the data
count = 0
for (i=0; i<N; i++) {
    if (A[i] == i) {
        count++
    }
}
```

T.C. = $O(N \log N)$
S.C. = $O(N)$
(merge sort)

* what if there are duplicates.

* understanding
the
problem.

quiz: $A = [-10, 1, 1, 3, 100]$

0 1 2 3 4	
$\downarrow \downarrow \downarrow \downarrow \downarrow$	
0 1 1 3 4	
$\times \checkmark \checkmark \checkmark \times$	

Ans = 3

Quiz: $A = [-10, 1, 1, 2, 4, 4, 4, 8, 10]$

0 1 2 3 4 5 6 7 8	
$\downarrow \downarrow \downarrow \downarrow \downarrow \backslash \backslash \searrow$	
1 1 1 1 1 1 1 1	

Ans = 5

↓ ↗ ↘ ↙ ↖ ↙ ↘ ↙ ↖ ↗
○ 1 1 3 4 4 4 7 8
✗ ✓ ✓ ✗ ✓ ✓ ✓ ✓ ✗ ✗

count = 7

Quiz: $A = \{ -3, 0, 1, 2, 2, 2, 4, 4, 4, 4, 8, 8, 8, 10, 10, 10, 10, 14 \}$

Smaller elements
~~not~~ ?

g_f Current element != previous element

(Count Smaller Element) = i

idea : ① sort the array

② if (current element) != previous element {

gt CA[i] == i)

cout << i;

③ else remember / calculate smaller elements

10:36

Code

- Idea:
- ① Sort the array
 - ② Keep record of countSmallerElements
 - ③ If $A[i] \neq A[i-1]$
the count of smallerElements = i
 - ④ else: don't update count of smallerElements

```
int countnobles(A){  
    // Sort the data → // T.C. = O(N log N)  
    int countofSmallerElements = 0  
    int nobleElements = 0  
    if (A[0] == 0) { nobleElements++; }  
    for (i=1; i<N; i++) {  
        |  
        if (A[i] != A[i-1]) {  
            |  
            T.C. = O(N log N)  
            S.C = O(N)  
            (merge sort)
```

```

    |     countOfsmallerElement) = i
    |
    |     }
    |
    |     if (countOfsmallerElement) == A[i]){
    |         nobleElements++;
    |
    |     }
    |
    |     return nobleElements;
}

```

Q. we have used library sort fn.

Q. Sort the numbers based on count of factors.
if count of factors are equal sort the number
on the basis of magnitude.

$A = [9, 3, 10, 6, 4]$

↓ ↓ ↓ ↓ ↓

3 2 4 4 3



$$\underline{\text{ans}} = [3, 4, 9, 6, 10]$$

Que: How we can use the in-built library function
to sort based on some rule.
↓
Count of factors.

- Sort () < by default 1 parameter → Array
- Sort (--), (array index i, array index j)
 to sort the numbers b/w
 index i to J.
- Sort (--, --,)
 comparator function
 rule for starting

```
int x, int y
```

```
if (x < y)
```

put x first

```
else
```

put y first

by default, it arranges
the integers based on
its magnitude (ASC)

sort \Rightarrow asc order, based on magnitude

```
int x, int y
```

```
if (countfactors(x) < countfactors(y))  
 $\Rightarrow$  put x first
```

```
if (countfactors(y) < countfactors(x))  
 $\Rightarrow$  put y first
```

```
gf (countFactors(x) == countFactors(y)) {  
    gf (x < y)  
        put x first  
    else  
        put y first  
}  
bool comp (int x, int y){  
    datatype datatype  
    + ↓ → gf true put x  
    | first  
    → gf false put y  
    first
```

```
int countX = countOfFactors(x)  
int countY = countOfFactors(y)
```

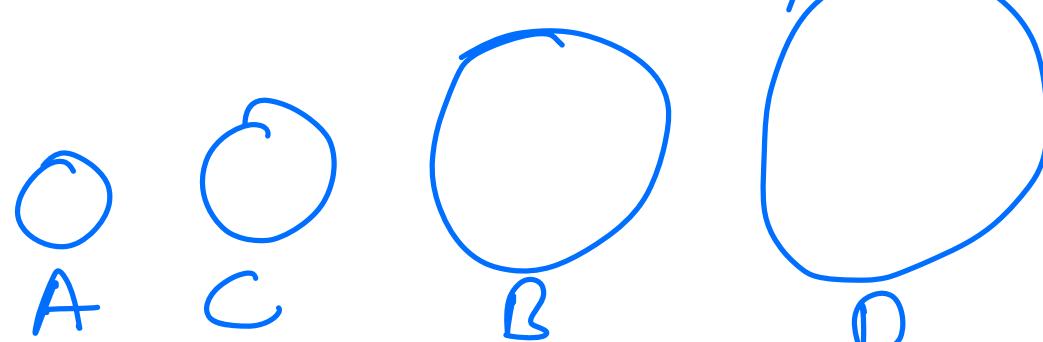
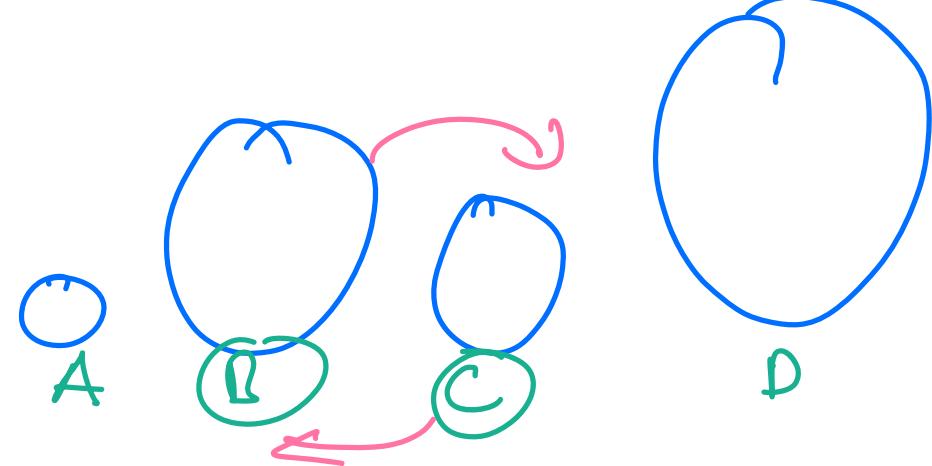
```
gf (countX < countY ) {  
    return true  
}  
else if (countY < countX ) {  
    return false  
}
```

else { // both are equal

if ($x < y$)
| return true
else
| return false

}

e.g.



Comparator fn in languages:

Java, Python, JS, C#, ruby (should come)

→ gt first parameter (X) before second (Y)
-re value should be returned

⇒ gt second parameter (Y) should come before first
+re value should be returned.

⇒ gt both are same, then 0 should be returned

in C++

⇒ in sorted form, if first parameter should
come first return true

⇒ else return false.

ex

[10 -5 2 5 8]

Sort them in $\text{desc} \downarrow$ order.

default library fn sort() returns in $\text{asc} \uparrow$ order.

⇒ we need to write our own comparator fn

code, C++

```
bool Comparator (int x, int y) {  
    if (x < y) {  
        // put y first  
        return false  
    }  
    else {  
        return true;  
    }  
}
```

⇒ Sort (A.begin(), A.end(), comparator);

Q. write the comparator fn to sort N strings.
based on string length

ex. ['aa' , 'b' , 'ccc' , 'aaa']
⇒ ['b' , 'aa' , 'ccc' , 'aaa']

Comparator function ?

```
bool comparator ( string x, string y ) {  
    int xlen = x.size()  
    int ylen = y.size()  
    if (xlen < ylen) {  
        return true;  
    }  
    ... }
```

else
return false;

}

}

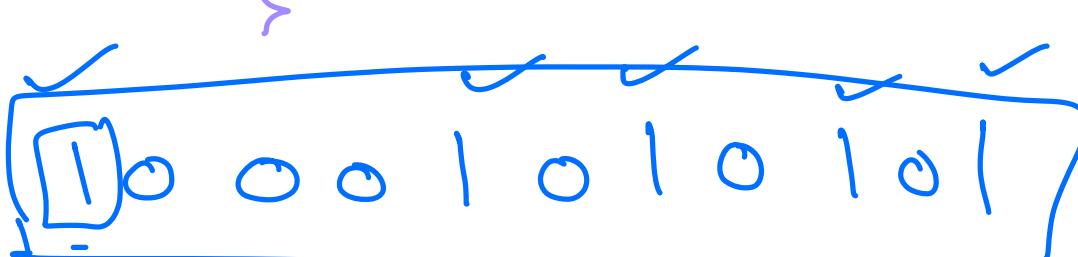
Doubt session

while ($A > 0$)

$A = A \& (A - 1)$ // unset the
last set bit
from the right

$$\begin{array}{r} 111 \\ 100 \\ \hline 000 \end{array}$$

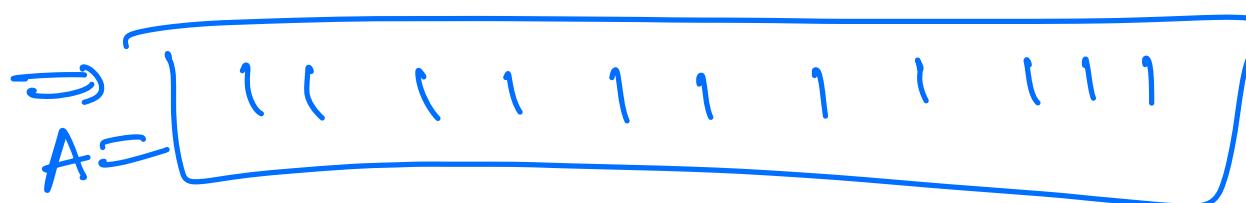
}



③

$$\Rightarrow \log_2^3$$

$\Rightarrow \boxed{1.7}$



$$\Rightarrow 2$$

$\boxed{1.5}$

worst case
count in $\log_2(A)$

Count of set bits

$$T.G = \underline{\underline{\log_2(A)}}$$

int 32 bit

$$T.C. = O(1)$$



$$\rightarrow \begin{array}{l} 11 \\ 5 \quad \log_2 5 \\ 2^n \\ \Rightarrow 101 \end{array}$$

* intuition behind bit-manipulation:

Q

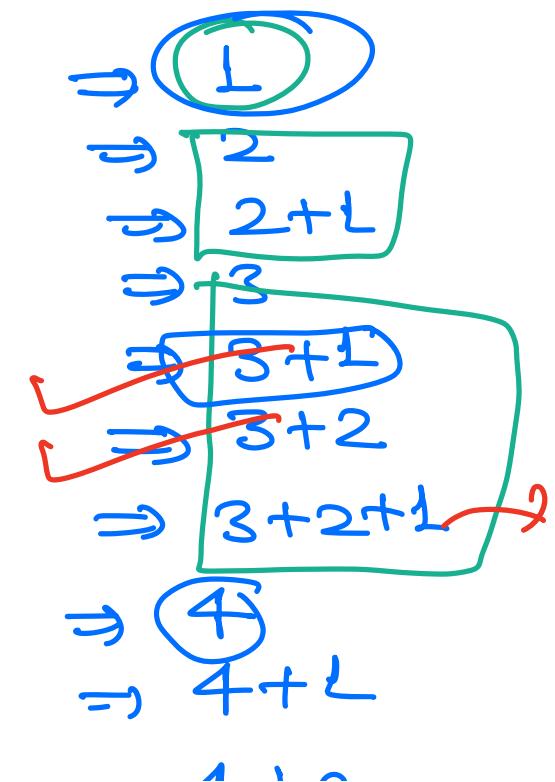
$$\begin{array}{l} 1 \rightarrow 5 \\ 2 \rightarrow 25 \\ 3 \rightarrow 30 \end{array}$$

$$\begin{array}{l} 4 \rightarrow 125 \\ 5 \rightarrow 130 \\ \vdots \end{array}$$

A

?

$$\begin{array}{c} 5^3 \quad 5^2 \quad 5^1 \\ \downarrow \quad \downarrow \quad \downarrow \\ S^3 + S^2 \rightarrow 130 \\ S^4 + S^3 \end{array}$$



$$\begin{array}{ccccc}
 & s^5 & & & \\
 \Rightarrow & 0 & 0 & s^4 & \\
 & & & 0 & s^3 \\
 \Rightarrow & 0 & 0 & & 0 \\
 \Rightarrow & 0 & 0 & & 1 \\
 \Rightarrow & 0 & 1 & & 0 \\
 \Rightarrow & 0 & 1 & 0 & 1 \\
 \Rightarrow & 0 & 1 & 1 & 0 \\
 \Rightarrow & 0 & 1 & 1 & 1
 \end{array}$$

\equiv

$$\begin{aligned}
 & \Rightarrow 4+2 \\
 & \Rightarrow 4+2+1 \\
 & \Rightarrow 4+3 \\
 & \Rightarrow 4+2+1 \\
 & \Rightarrow 4+3+2 \\
 & \Rightarrow 4+3+2+1
 \end{aligned}$$

polyⁿ

$$a_n x^n + a_{n-1} x^{n-1}$$