

Modular Arithimetic Introduction

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Reachable in Scaler Lounge 

"Everything in moderation; including moderation!"

Today's content:

- % operator
- modular arithmetic
- One easy problem
- One hard problem.

int range: $[-2 \times 10^9, 2 \times 10^9]$ $[-2^{31}, 2^{31}-1]$
(4 bytes)

long range: $[-8 \times 10^18, 8 \times 10^{18}]$

% Basics:

$n \% a$ = remainder when n is divided by a

$r = \text{dividend} - (\text{greatest multiple of divisor} \leq \text{dividend})$

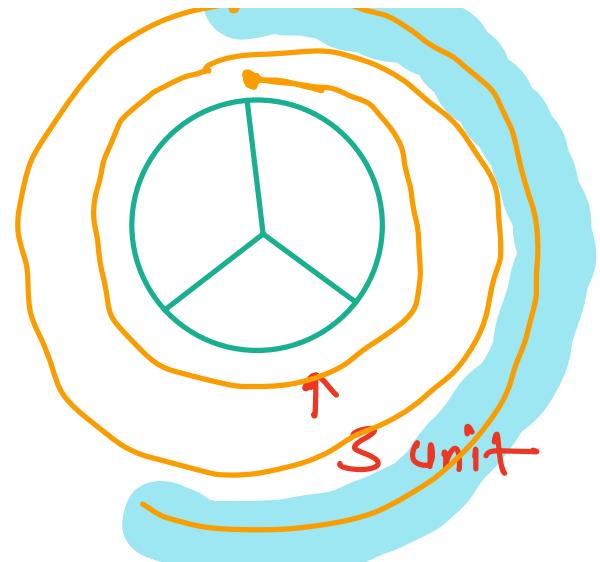
dividend $\frac{A}{B}$, Quotient Q + Remainder R
divisor



$$\frac{8 \% 3}{(0, 5, 10, 15)} = 8 - (3 * 2) = 2$$

$$13 \% 5 = 13 - (10) = 3$$

↑
(0, 5, 10, 15)



Reminder = dividend - divisor * Quotient

↳ (greatest multiple of divisor \leq dividend)

Quiz $150 \% 11 = ?$

$$= 150 - (\text{greatest multiple of } 11 \leq 150)$$

$$= 150 - (143)$$

$$= 7$$

$$\Rightarrow \frac{150}{11} =$$

11	150	13
$\frac{11}{40}$		
$\frac{33}{7}$		

Quiz 2

$$100 \% 7 = ?$$

$$\begin{array}{r} 100 \\ 7) \underline{100} \\ -7 \\ \hline 30 \\ -28 \\ \hline 2 \end{array}$$



$$= 100 - (\text{greatest multiple of } 7 \leq 100)$$

$$= 100 - (98)$$

$$= 2$$

Quiz 3

$$-40 \% 7 = ?$$

$$= -40 - (\text{greatest multiple of } 7 \leq -40)$$

$$\downarrow$$

($-63, -56, -49, \text{ -42, } -35,$
 $-28, -21, -14$
 $-7, 0, 7, 21$)

$$= -40 - (-42)$$

$$= 2.$$

Quiz 4



$$\text{Ques. 5. } -60 \% 9 = ?$$

$= -60 - (\text{greatest multiple of } 9 \leq -60)$

$$\begin{aligned} &= -60 - (-63) \quad \dots \\ &= \underline{\underline{3}}. \end{aligned}$$

$(-90, -81, -72, \textcolor{red}{-63}, -54, -45)$

* 1. operator in programming languages $-40 - (-35) = 5$

ex.

Python

Java/C/C++/C#/JS

$$-40 \% 7$$

2

$$\longleftrightarrow -5$$

$$(-5+7)$$

$$-60 \% 9$$

3

$$\longleftrightarrow -6$$

$$(-6+9)$$

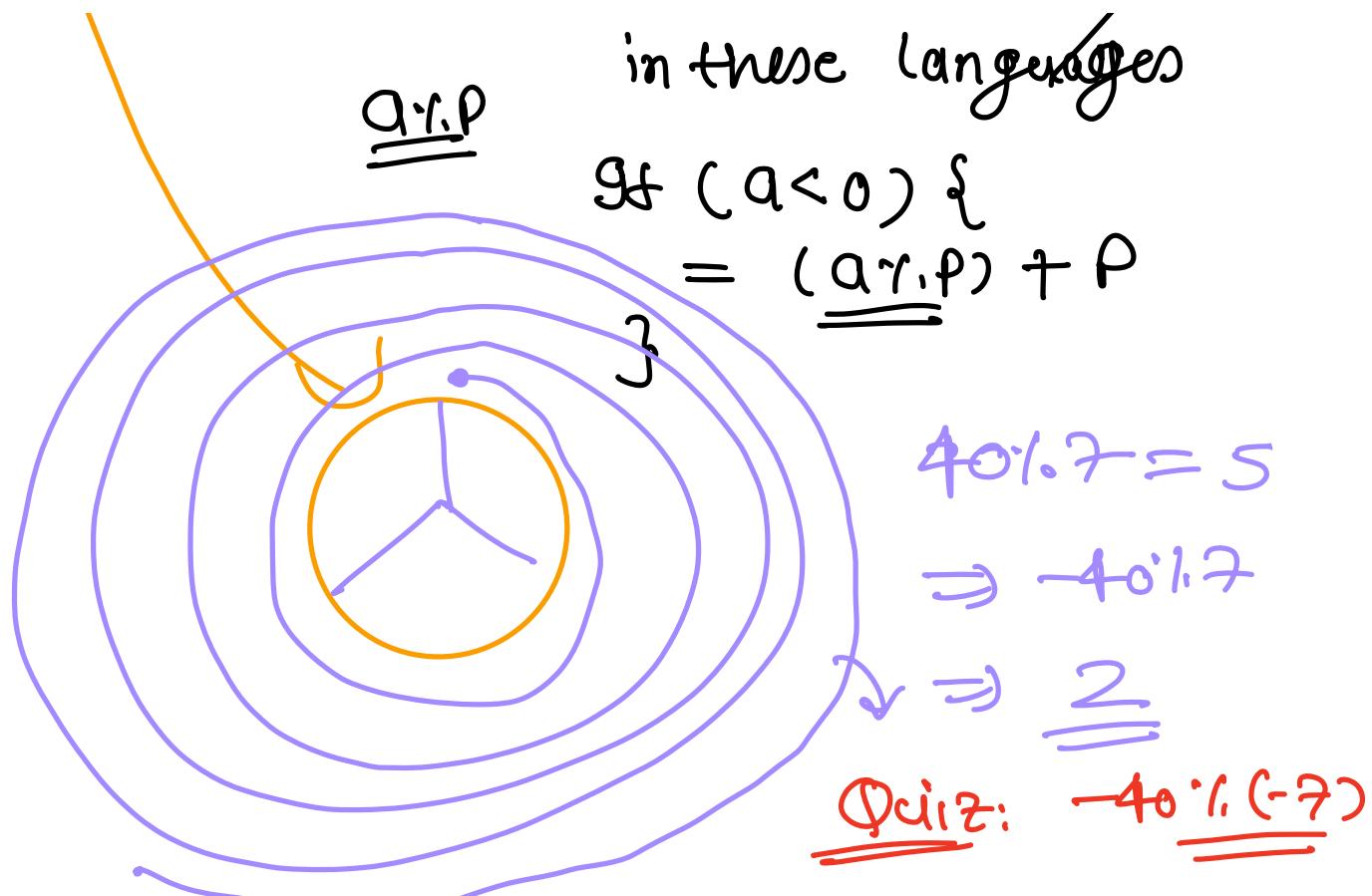
$$-40 \% 9$$

5

$$\longleftrightarrow -4$$

$$(-4+9)$$

Obsⁿ: difference is equal to the divisor



*

Why?

→ limit the scope of our data.

$$-\infty \left[\begin{array}{c} \downarrow \\ \cdot 1.0 \\ \uparrow \\ +\infty \end{array} \right] = \underline{\underline{[0, 9]}}$$

$$-\infty \nwarrow \cdot 1.p = [0, p-1]$$

Club \Rightarrow A_r B_r C
↓
serial number r, 3

$+ \infty$

\equiv



Hashing (DSA)

consistent (HLP, LLD)

Hashing

Properties.



$$\textcircled{2} \quad (a+b) \% M = ((a \% M) + (b \% M)) \% M$$

$$a \quad b \quad M \\ 8 \quad 6 \quad 10$$

$$(a+b) \% M$$

$$4$$

$$((a \% M) + (b \% M)) \% M \\ (8 \% 10) + (6 \% 10) = 4$$

$$18 \quad 11 \quad 10$$

$$9$$

$$(8 + 1) \% 10 = 9$$

$$18 \quad 36 \quad 10$$

$$4$$

$$((8 + 6)) \% 10 = 4 \\ (18 \% 10) + (36 \% 10)$$



$$\textcircled{2} \quad (a * b) \% M = ((a \% M) * (b \% M)) \% M$$

$$a \quad b \quad M \\ 8 \quad 7 \quad 6$$

$$(a * b) \% M \\ 2$$

$$((a \% M) * (b \% M)) \% M \\ ((2) * 1) \% M$$

5 4 6

2

$(5 * 4) \% 6$

$(a - b) \% M$

$(a / b) \% M$

\rightarrow Inverse modulo

advanced module

③

$$(a \% P) \% P = a \% P$$

$$\underline{\underline{[0, P-1] \% P}} = \underline{\underline{[0, P-1]}}$$

$$\begin{aligned} & (16 \% 10) \% 10 \\ & \quad \uparrow \\ & (6 \% 10) \Rightarrow 6 \\ & \rightarrow (16 \% 10) \end{aligned}$$

④

$$(a \% P * b \% P) \% P = (a * b \% P)$$

$$x = a \% P, y = b$$

$$\begin{aligned} (x * y \% P) \% P &= ((x \% P * y \% P) \% P) \% P \\ &= (((a \% P * b \% P) \% P) \% P) \% P \end{aligned}$$

$$\begin{aligned}
 &= ((a \cdot r \cdot p) \star (b \cdot r \cdot p)) \cdot r \cdot p \\
 &= \underline{(a \star b) \cdot r \cdot p}
 \end{aligned}$$

Quiz: $13 = ?$

$\frac{232}{\times}$	$\frac{4560}{\checkmark}$	$\frac{238}{\times}$	$\frac{4333}{\times}$
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④ Divisibility rules:

$12 =$ (if the last digit is $(0, 2, 4, 6, 8)$
is divisible by 2)

$13 =$ sum of digits divisible by 3

$14 =$ last 2 digits / divisible by 4

$18 =$ last 3 digits / divisible by 8

$19 =$ sum of digits divisible by 9

H.W. $\Rightarrow \underline{\underline{f7 = ?}}$

$\Rightarrow 4563 \cdot 1.3 = ?$

$$= (4000 + 500 + 60 + 3) \cdot 1.3$$

$$= ((4000 \cdot 1.3) + (500 \cdot 1.3) + (60 \cdot 1.3) + (3 \cdot 1.3)) \cdot 1.3$$

$$= ((4 * 1000) \cdot 1.3) + \text{---}$$

$$= ((4 \cdot 1.3) * (1000 \cdot 1.3)) \cdot 1.3 + \text{---}$$

$$= (1 * 1 \cdot 1.3) * \text{---}$$

$$= 1 + ((5 \cdot 1.3) * (100 \cdot 1.3)) \cdot 1.3 + \text{---}$$

$$= 1 + (2 * 1) \cdot 1.3 + \text{---}$$

$$= (1 + 2 + ((6 * 3) * (10 \cdot 1.3)) \cdot 1.3 + 3 \cdot 1.3) \cdot 1.3$$

$$\begin{aligned}
 &= (1 + 2 + (0 * 1) \% 3) \% 3 \\
 &= (1 + 2 + 0 \% 3) \% 3 \\
 &= 0
 \end{aligned}$$

Obsⁿ:

$$\begin{array}{ll}
 10 \% 3 = 1 & = ? \text{ divisibility rule for } 3 \\
 100 \% 3 = 1 & 10 \% 4 = 0 \\
 1000 \% 3 = 1 & 100 \% 4 = 0 \\
 \vdots & \vdots \\
 10^n \% 3 = 1 & 10^n \% 4 = 0 \quad n \geq 2
 \end{array}$$

10 : 13

* 2347 / 13 = ?

$$\begin{aligned}
 &= \underbrace{(2 * 1000) \% 13}_{\% 3} + (3 * 100) \% 13 + (4 * 10) \% 13 + (7 \% 13) \% 13
 \end{aligned}$$

$$\begin{aligned}
 &= ((2 \cdot 3) + (3 \cdot 3) + (4 \cdot 3) + (7 \cdot 3)) \cdot 3 \\
 &= (2+3+4+7) \cdot 3 \\
 &= 16 \cdot 3 \\
 &= 1
 \end{aligned}$$

* divisibility by 2

$$2457 \cdot 2 = ?$$

$$\begin{aligned}
 &= [(2 * 1000) \cdot 2 + (4 * 100) \cdot 2 \\
 &\quad + (5 * 10) \cdot 2 + (7 * 1) \cdot 2] \cdot 2 \\
 &= (0 + 0 + 0 + (7 \cdot 2)) \cdot 2 \\
 &= \underline{\underline{7 \cdot 2}}
 \end{aligned}$$

* divisibility by 4.

$$2457 \% 4 = ?$$

$$= ((2 * 1000) \% 4 + (4 * 100) \% 4 + (57 \% 4)) \% 4$$

$$\therefore (57 \% 4)$$

= If the last 2 digit num is divisible by 4, then the given number is divisible by 4.

Q. Given a, n, p , calculate $a^n \% p$ without using 'inbuilt' function.

Ex. $a = 3, n = 4, p = 7$

$$\Rightarrow 3^4 \% 7 \Rightarrow (3 \times 3 \times 3 \times 3) \% 7$$

$$\Rightarrow 81 \% 7 \Rightarrow 4$$

$$\Rightarrow (5 \cdot 1 \cdot 2 * 10 \cdot 1 \cdot 2) \% 2$$

$$\Rightarrow (1 * 0) \% 2 = 0$$

$$\begin{aligned} 1 &\leq a \leq 10^9 \\ 2 &\leq p \leq 10^9 \\ 1 &\leq n \leq 10^5 \end{aligned}$$

~~Q.R.~~ $a=4, n = 5, p = 6$

$$\Rightarrow (4^5) \% 6 \Rightarrow (16 * 16 * 4) \% 6$$
$$\Rightarrow (256 * 4) \% 6$$
$$= (1024) \% 6$$
$$= \underline{\underline{4}}$$

~~Q.R.~~ Brute force:

idea: ① calculate a^n (build pow fn)
② $a^n \% p = \underline{\underline{a^n}}$.

```
int modulo (a, n, p) {  
    int ans = 1  
    for (i=1; i<=n; i++) {
```

// $a * a * a * ..$
- n times
 $\underline{\underline{a^n}}$

```

    O |
    }   ans *= a
}
return ans % p;
}

```

T.C. = O(N)
S.C. = O(1)

NO TLE
partial correct \Rightarrow WRONG

* $a = 2, n = 30, p = 45$

$\Rightarrow a^n \Rightarrow \underline{\underline{2^{30}}}$ (we can store this in
won't overflow $(-2^{31}, \underline{\underline{2^{31}}})$)

* $a = 2, n = 60, p = 45$

$\Rightarrow a^n \Rightarrow \underline{\underline{2^{60}}}$ (we can NOT store this
in int)
Solⁿ \Rightarrow instead of int we use long

* $a = 2, n = 100, p = 45$

$\Rightarrow a^n = \underline{\underline{2^{100}}} \Rightarrow$ (No data structure can hold this value)

Optimal approach:

Idea: take modulo.

$$\underbrace{(a * a) \% P}_{\downarrow} \Rightarrow \underbrace{((a \% P) * (a \% P)) \% P}_{\downarrow}$$

this might overflow

this won't overflow

```
int modulo (a, n, P) {  
    a = a \% P  
    int ans = 1
```

```
    for (i=1; i<=n; i++) {  
        ans = (ans * a) \% P  
    }
```

$$\overbrace{((1 * a \% P) * (a \% P)) \% P}^{\overbrace{* (a \% P)}^{\text{in denominator}} \% P}$$

\downarrow - - in denominator

}

return ans;

(ans is array)
bln (0, P-1)

Dry Run:

a, P, n=5

10^9

10^9

No overflow

ans

i

$i \leq 5$

$ans = (ans * a) \% P$

1

1

✓

$ans = \underline{\underline{a \% P}}$ ^{no overflow}

$a \% P$

2

✓

$$\begin{aligned} ans &= (a \% P * a) \% P \\ &= ((a \% P) \% P) * \\ &\quad a \% P \% P \end{aligned}$$

$$= ((a \% P) * (a \% P)) \% P$$

$$= \underline{\underline{(a * a) \% P}}$$

$$= a^2 \% P \rightarrow \text{no overflow}$$

$a^2 \cdot P$ 3 ✓ $\text{ans} = \overbrace{(a^2 \cdot P * a)}^{\text{---}} \cdot P$

$= \overbrace{a^3 \cdot P}^{\text{---}} \rightarrow \text{no overflow}$

$a^3 \cdot P$ 4 ✓ $\text{ans} = (a^3 \cdot P * a) \cdot P$

$= (a^3 * a) \cdot P$

$= a^4 \cdot P \rightarrow \text{no overflow}$

$a^4 \cdot P$ 5 ✓ $\text{ans} = ((a^4 \cdot P) * a) \cdot P$

$= a^5 \cdot P \rightarrow \text{no overflow}$

$a^5 \cdot P$ 6 ✗ $\Rightarrow \text{return ans}$

Q. Given a number in array format, calculate A%P.

$$1 \leq \text{Array size} \leq 10^5$$

$$0 \leq A[i] \leq 9$$

$$2 \leq P \leq 10^9$$

Ex. $N = 5, A = [6, 2, 3, 4, 3], P = 49$

$$\Rightarrow (62343) \% 49 = 15 = \text{ans}$$

Ex. ② $N = 4, A = [2, 4, 3, 7], P = 16$

$$\Rightarrow (2437) \% 16 = 5 = \text{ans}$$

Brute force soln:

idea: ① construct the number num

② num % P

flow: ① overflow

* Optimised approach:

i) take modulo inside

ex. $N = 5, A = [6, 2, 3, 4, 3]$
 $P = 49$

$$= \underbrace{[6 \times 10^4 + 2 \times 10^3 + 3 \times 10^2 + 4 \times 10^1 + 3 \times 10^0]}_{} \text{ } \% \text{ } P$$

option
 $\text{num} = 62343 \% 49$

$\therefore P$
 $N = 10^5$
 $\Rightarrow 4 \times 10^{10}$

$$\frac{6 \times 10^{10} \text{ \% P}}{8 \times 10^{10} \text{ \% P}}$$

easier

$$6 \times 10^4 \text{ \% P}$$

$$\begin{aligned}
 & (3\% \cdot P * (10^t \cdot P))^n \cdot P \\
 & (4 * 10^t) \% \cdot P \\
 & (4 * (10^t \cdot P))^n \cdot P \\
 & (4 * (t * 10)) \% \cdot P \\
 & t = \underline{\underline{(t * 10)) \% P}} \\
 & 3 * (t * 10) \% \cdot P \\
 & t = \underline{\underline{(t * 10)) \% P}} \\
 & 2 * (t * 10) \% \cdot P
 \end{aligned}$$

$$\Rightarrow ((6 \% \cdot P)^* (10^4 \% \cdot P))^n \cdot P$$

$$\Rightarrow \bar{[(G \% P) * (10^* \underline{\underline{10^3 \% P}})] \% P}$$

$$\Rightarrow ((G \% P) * (10^* t)) \% P$$

$$t = 10^x \% P$$

Brute force:

```
int arrModulo ( A , P ) {
    int power= 1; int answer= 0
    for ( i= N-1; i >= 0; i++ ) {
        answer = ( answer + A [ i ] * power ) \% P
        power = power * 10 \% P
    }
}
```

6234]

\downarrow
4 \% P

$3 \times 10 \% P$

$(2 \times 100) \% P$
 $(6 \times 1000) \% P$

\nwarrow

..s

~~power = Power * 10~~ X
 } power = 10^0
 } → Overflow

Optimal Solution

```

int arrModulo (
  int t = 1;        //  $10^0 \% P = 1 \% P = 1$ 
  int answer = 0;

  for (i=N-1; i ≥ 0; i--) {
    answer = (answer + ((A[i] % P) * t)) % P // No overflow
    t = (t * 10) % P                                // No overflow
  }
  return answer;

```

T.C. = $O(N)$
 S.C. = $O(1)$

3