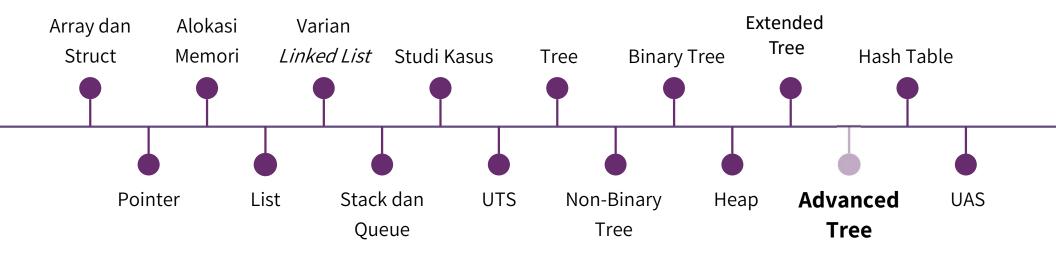
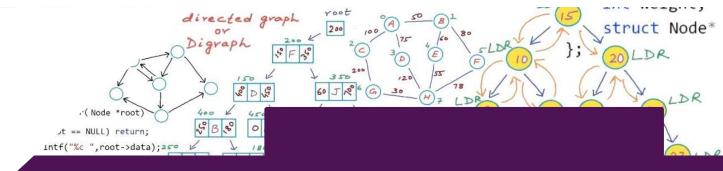


Pekan 14



Tujuan



- 1 Mahasiswa memahami tree yang digunakan untuk aplikasi khusus.
- 2 Mahasiswa memahami cara kerja trie dan balanced tree
- Mahasiswa mampu mengimplementasikan trie dan balanced tree.

Object Space Decomposition

- The decomposition of the key range is driven by the objects (i.e., the key values of the data records) stored in the tree.
- The shape of the tree will depend on the order of key insertion.
- Example: BST

Key Space Decomposition

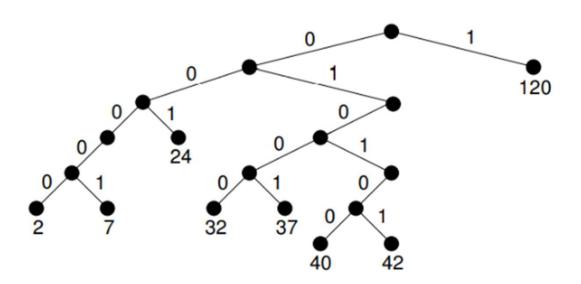
- Splitting based on predetermined subdivisions of the key range (the root could be predefined to split the key range into two equal halves, regardless of the particular values or order of insertion for the data records)
- The shape of the tree will not depend on the order of key insertion.
- The depth of the tree will be limited by the resolution of the key range.
- Example: Trie

Key Space Decomposition

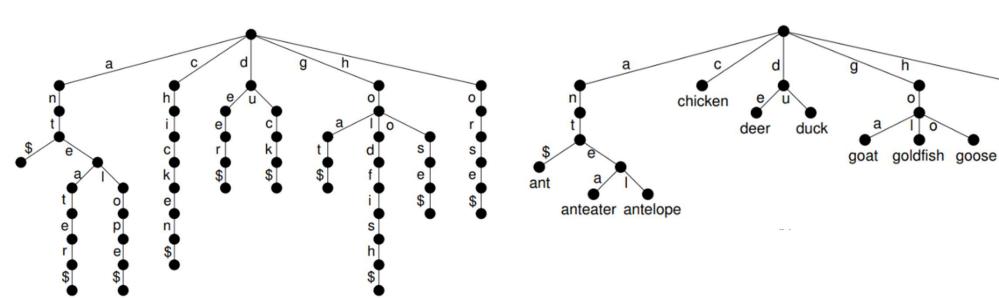
- If the keys are integers in the range 0 to 1023
 - \circ The resolution for the key is log 1023 = 10 bits.
 - As a result, the tree will never be more than ten levels deep.
 - In contrast, a BST containing n records could be as much as n levels deep.

Trie

Binary Trie

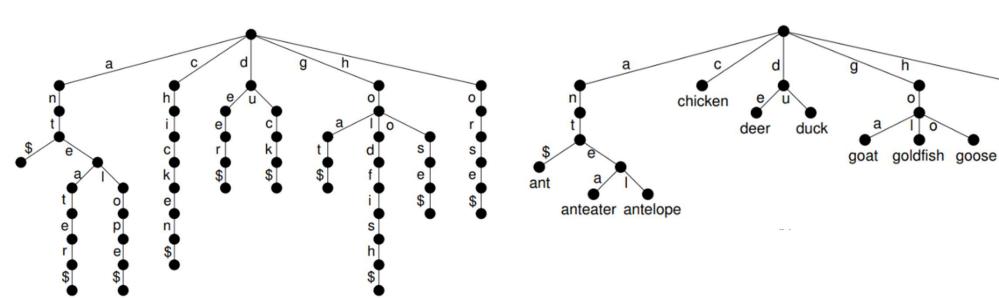


Alphabet Trie



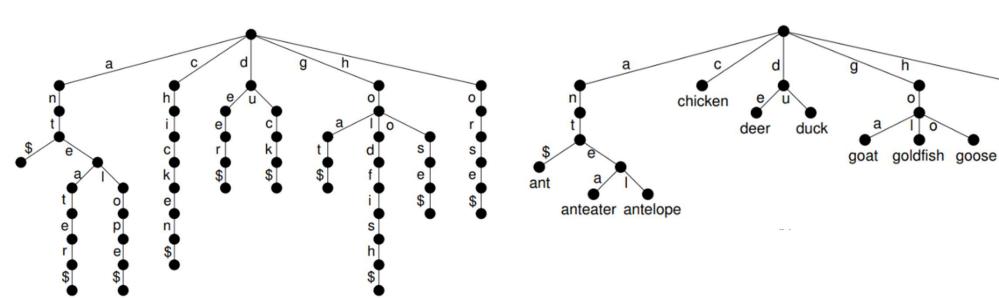
horse

Alphabet Trie

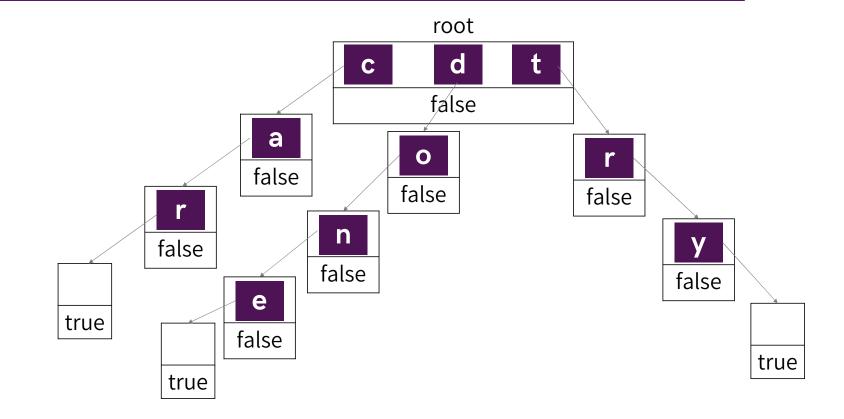


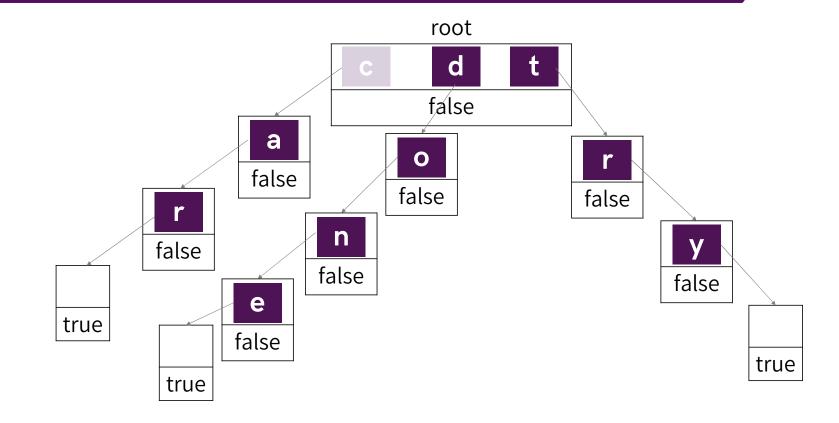
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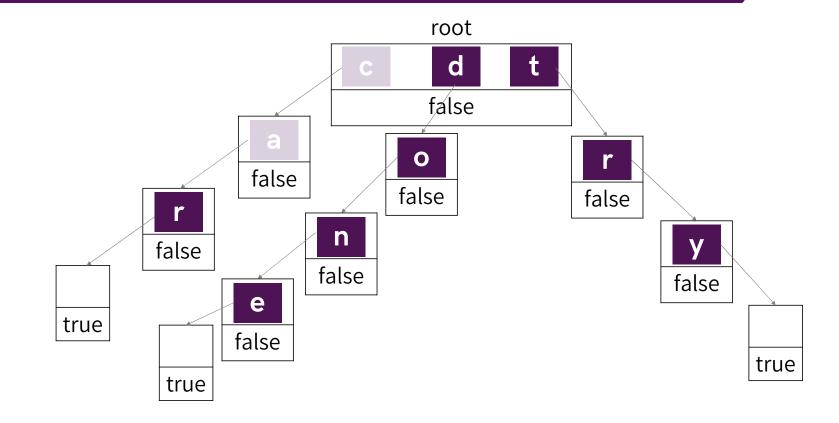
Alphabet Trie

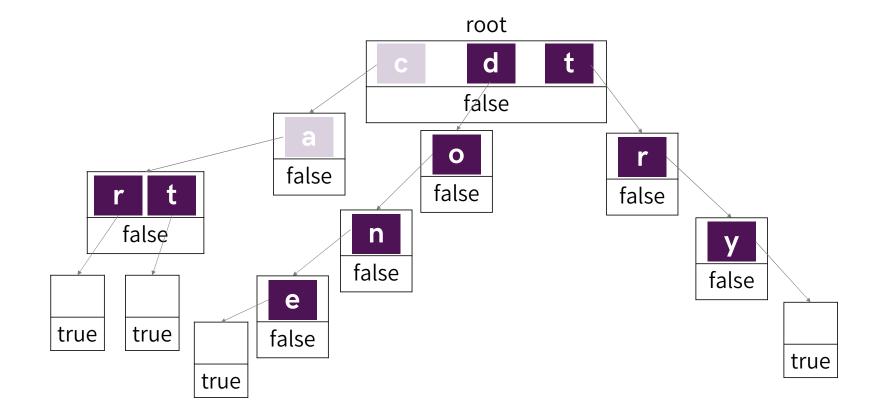


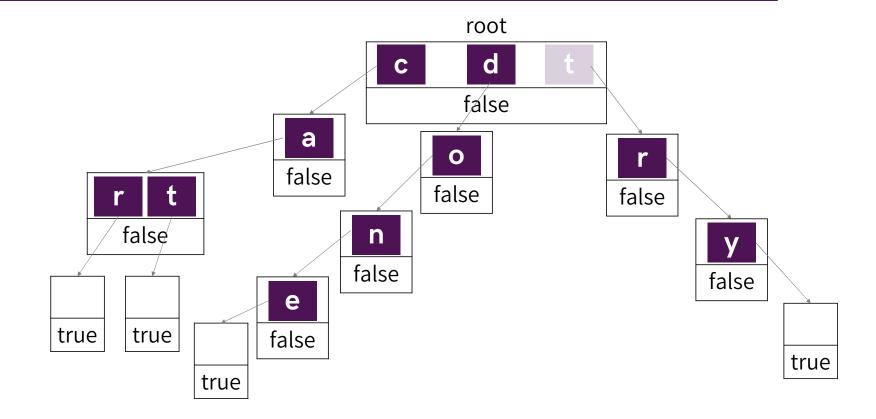
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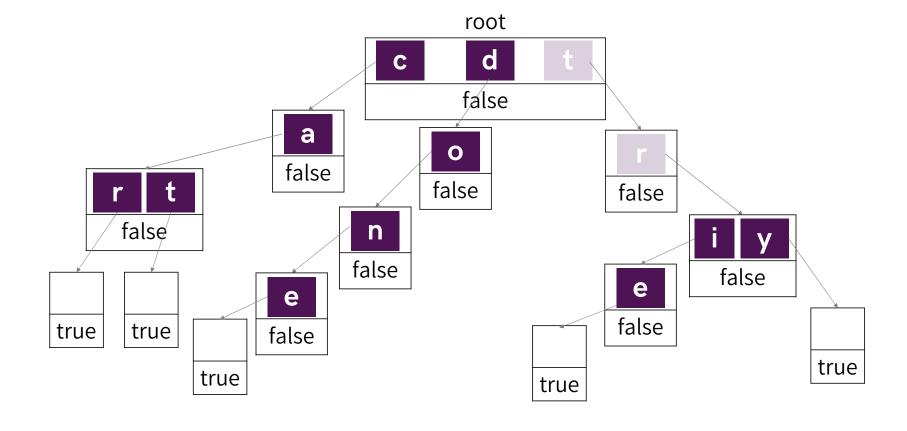


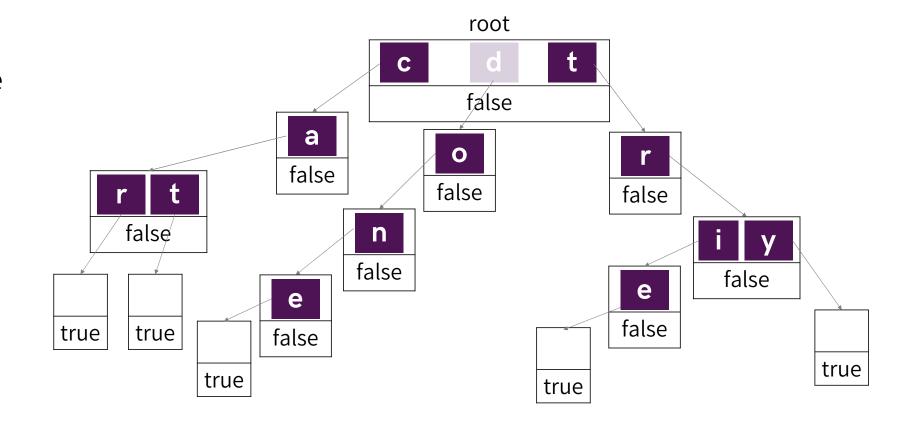


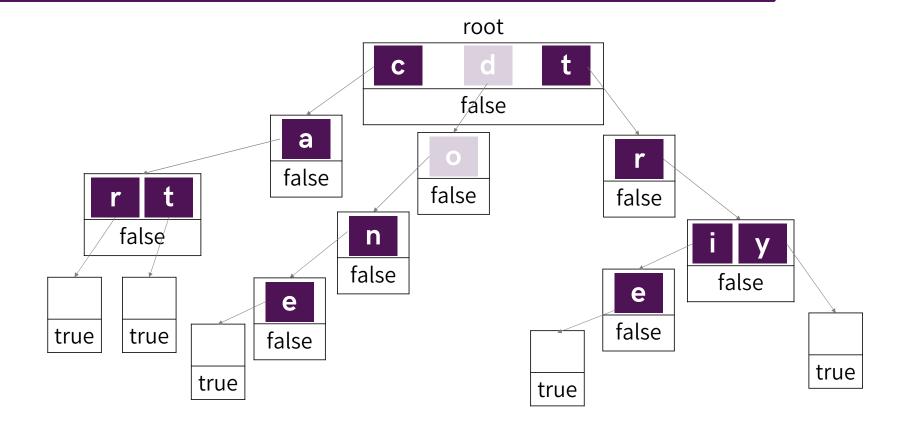


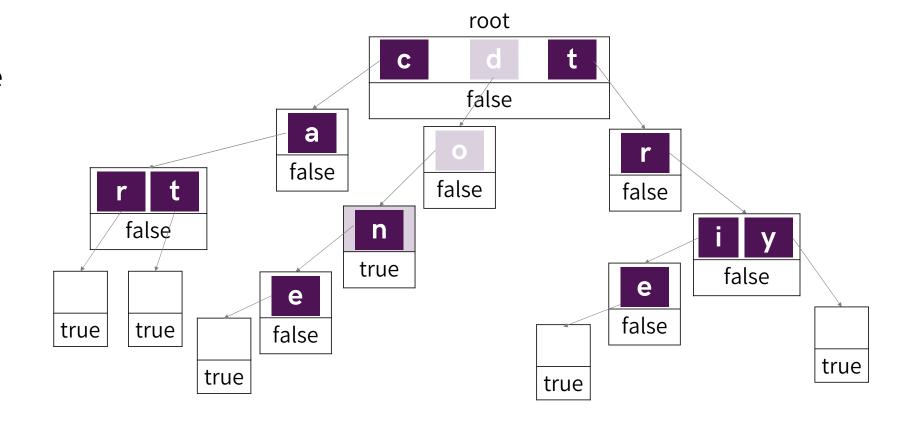




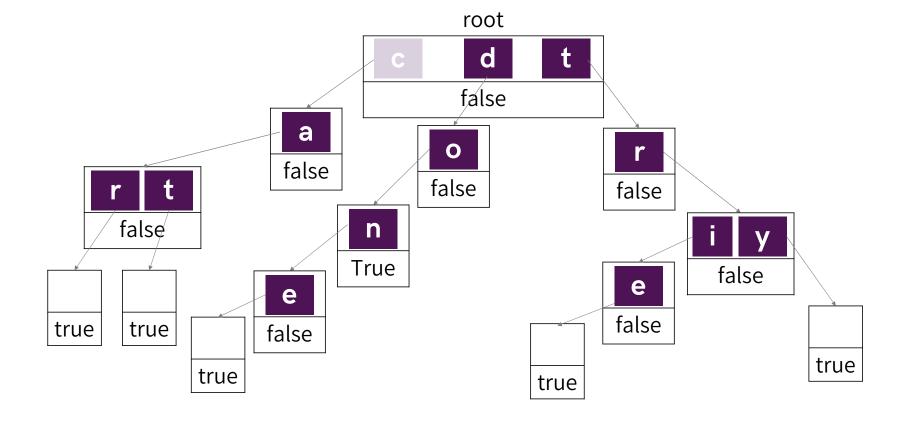




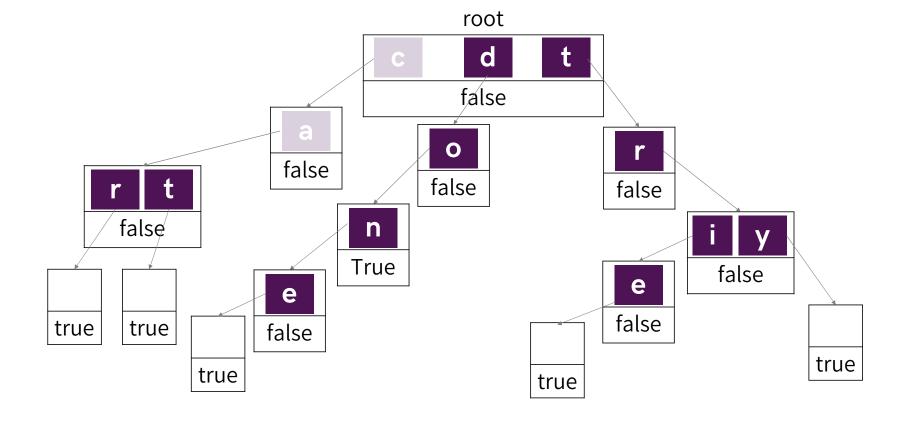




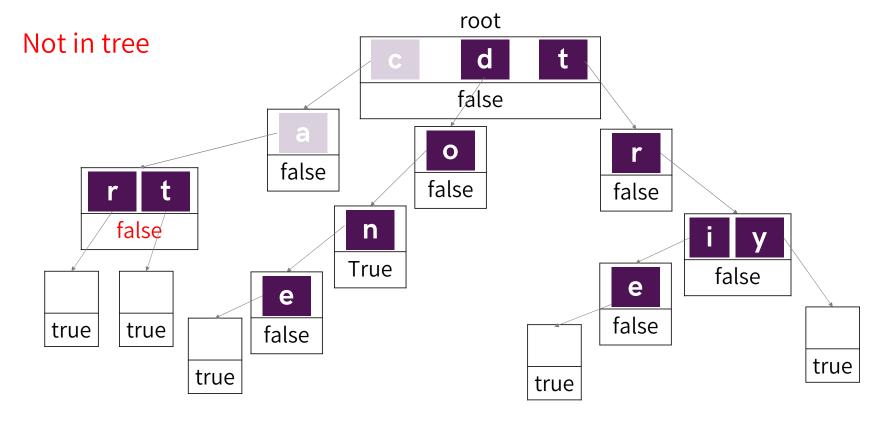
<u>ca</u> zebra try card

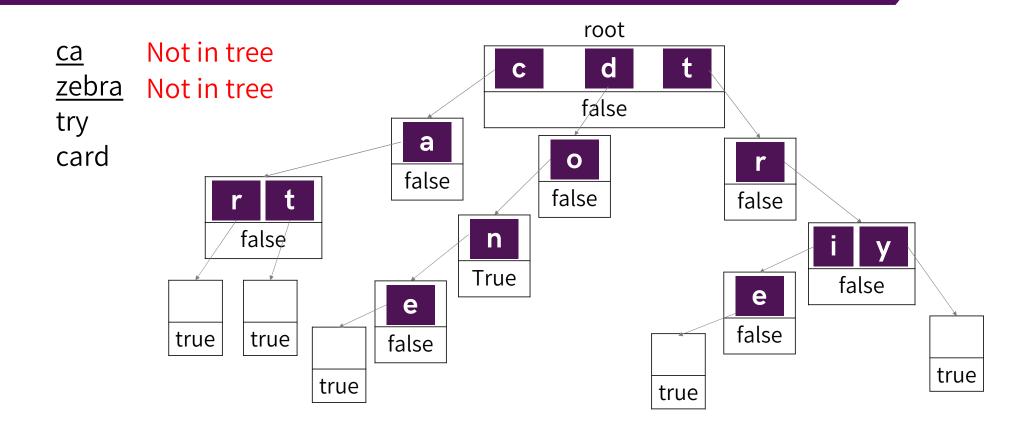


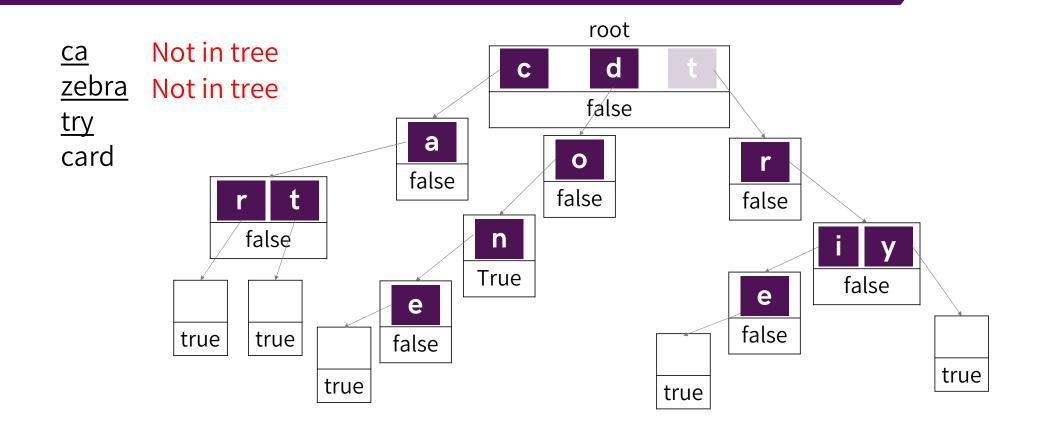
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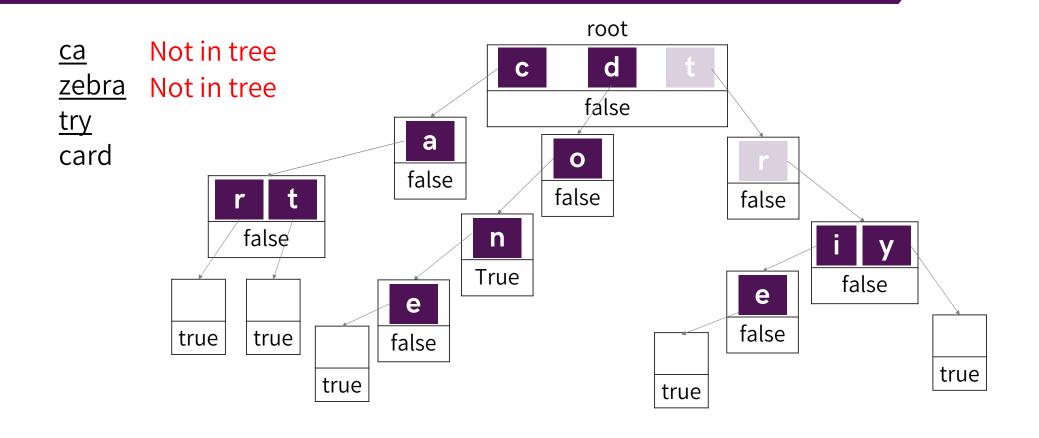


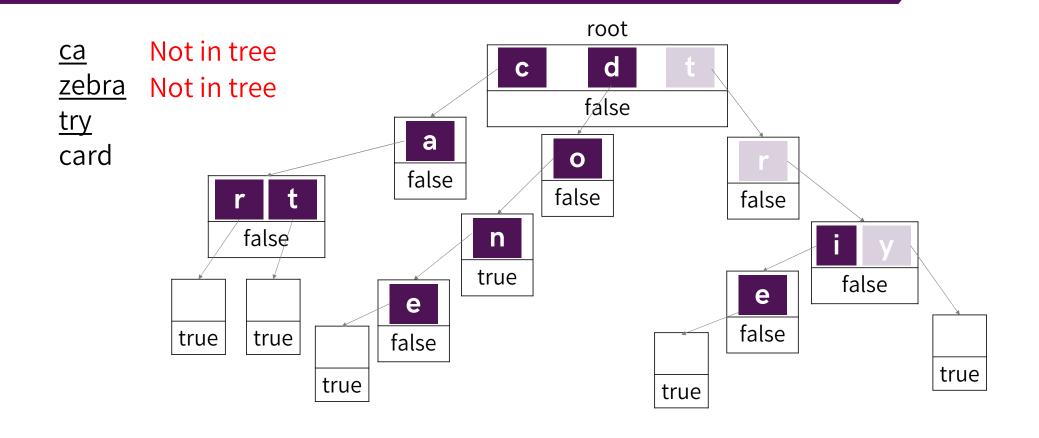
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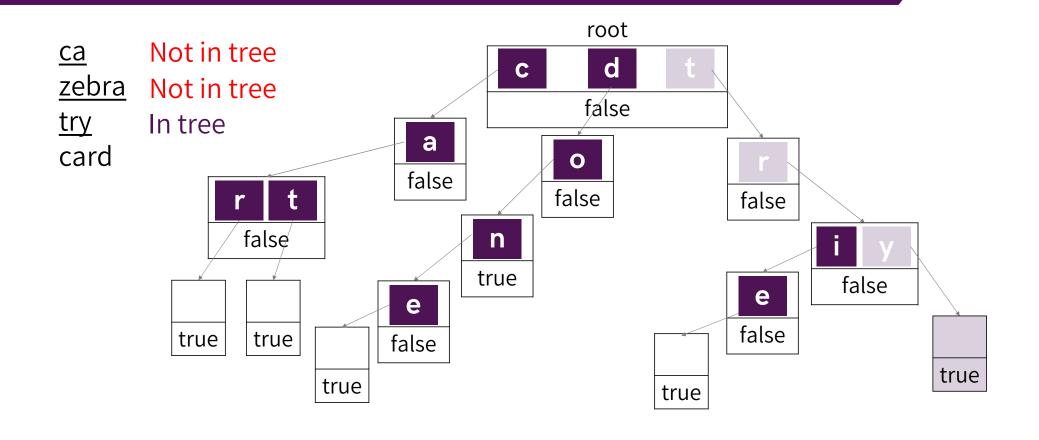


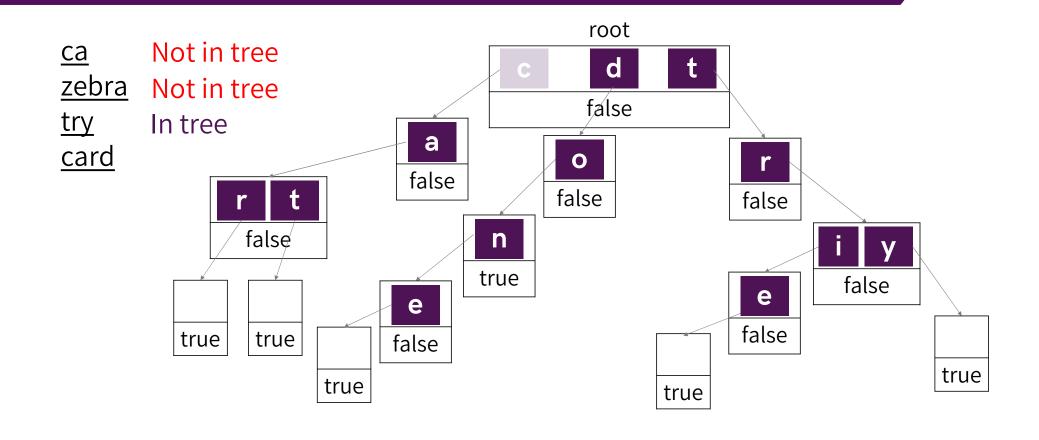


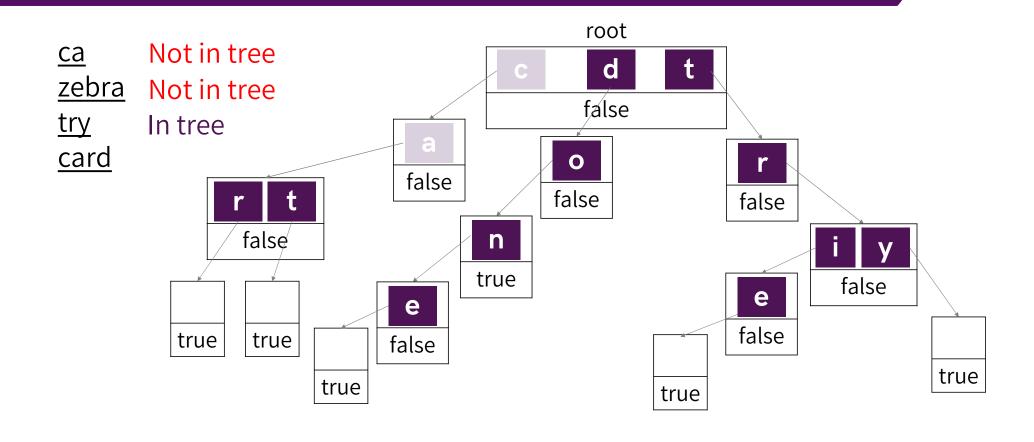


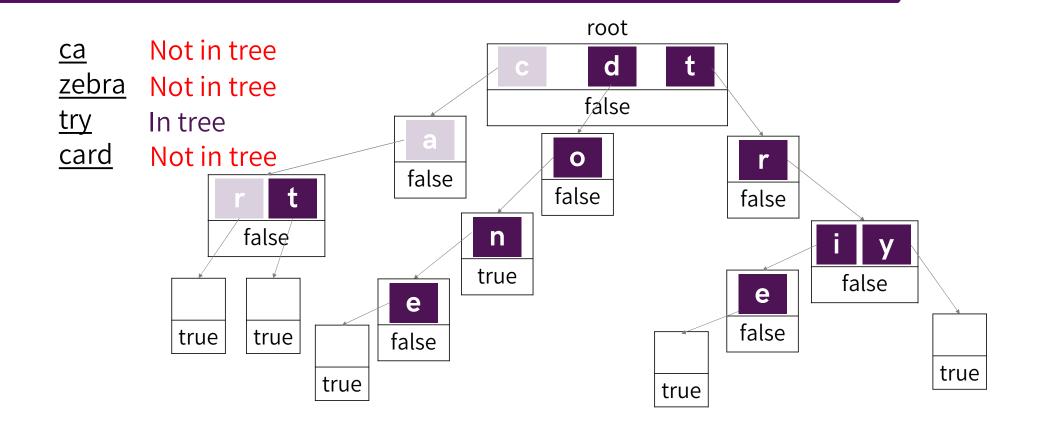




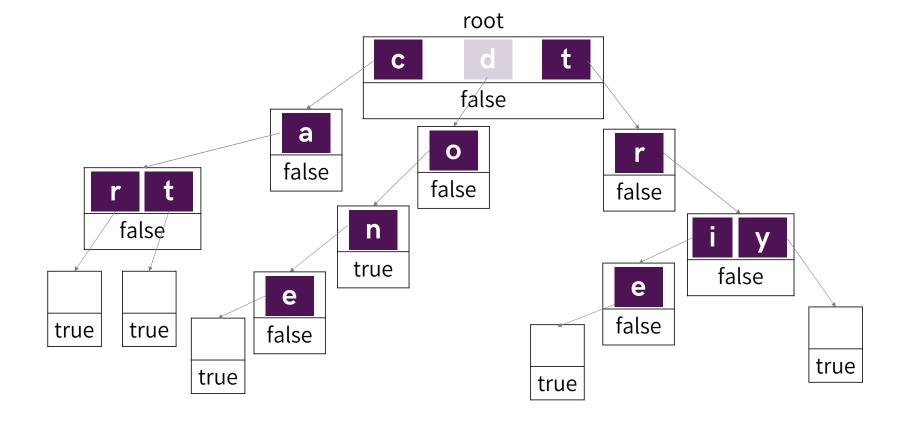




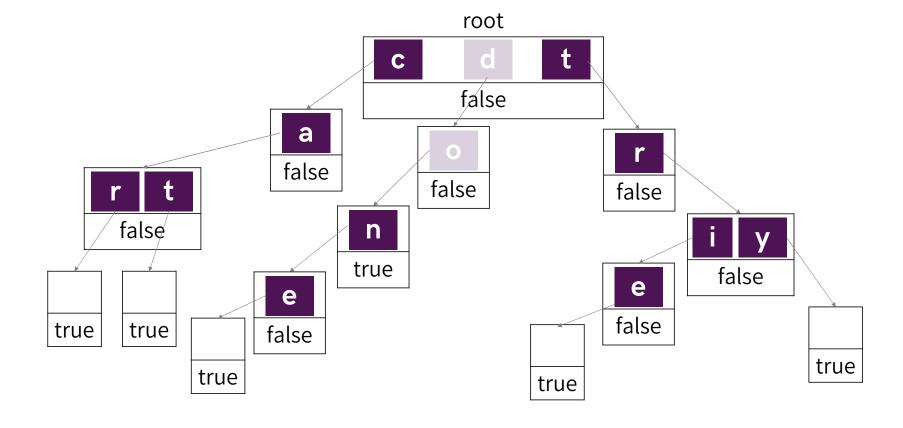




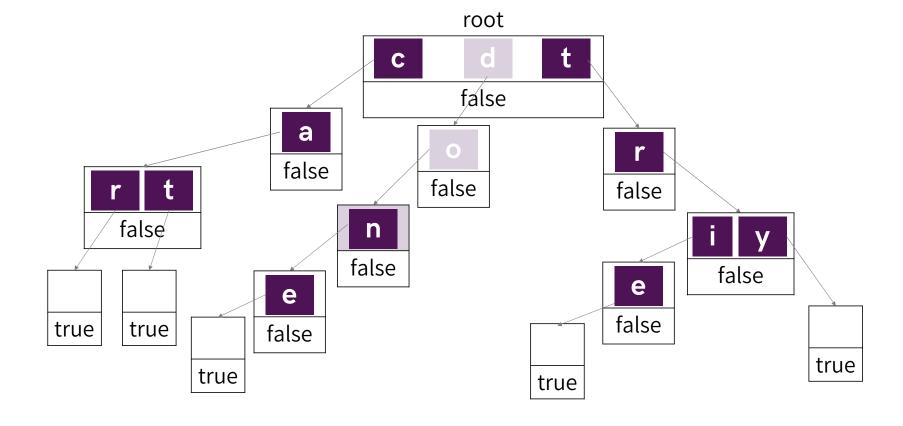
<u>do</u> trie ca

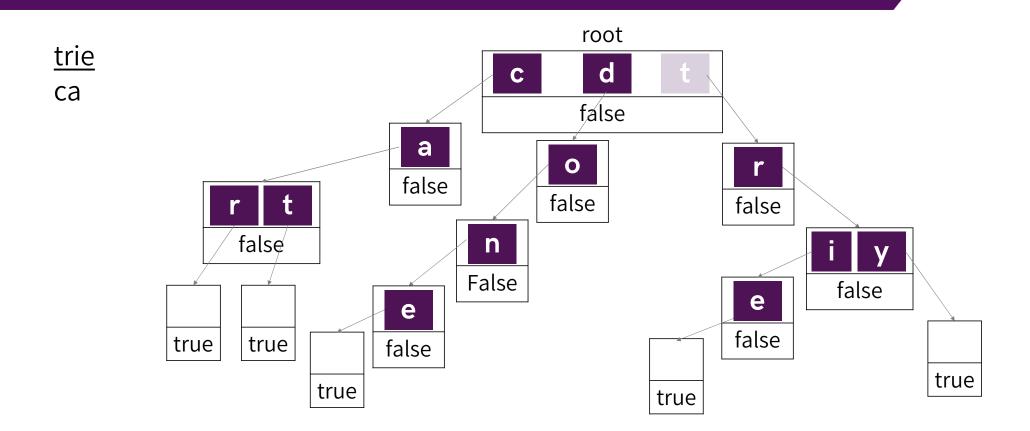


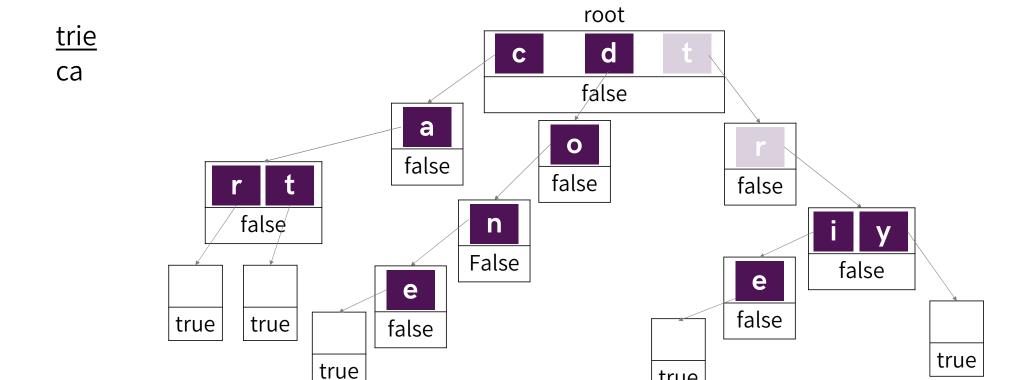
<u>do</u> trie ca



<u>do</u> trie ca

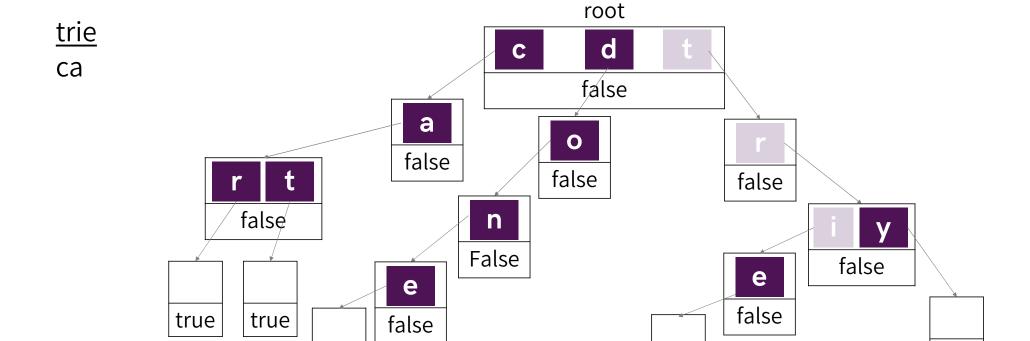






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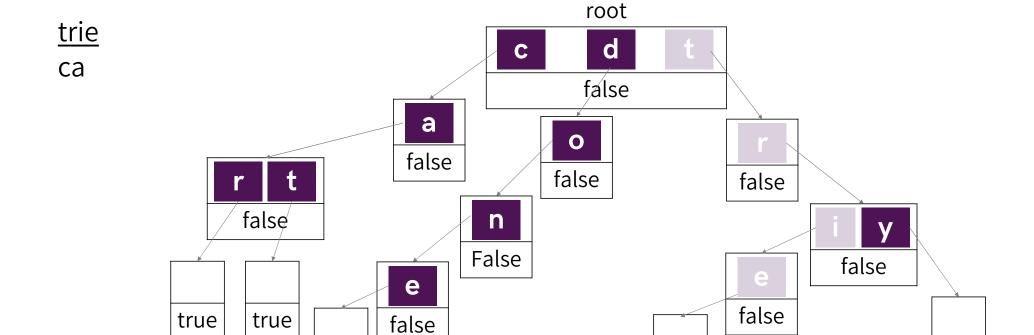
true



true

true

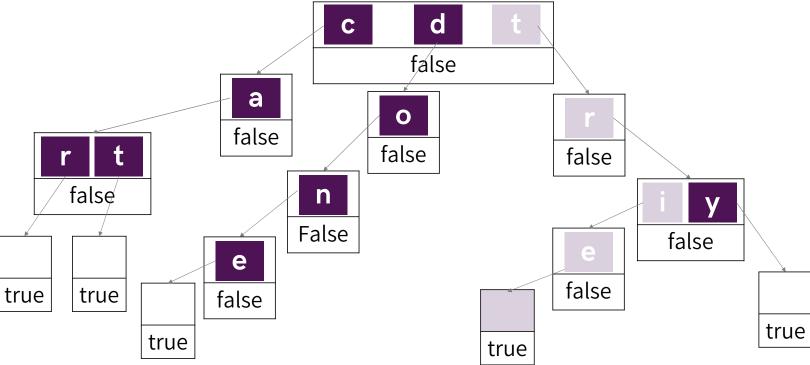
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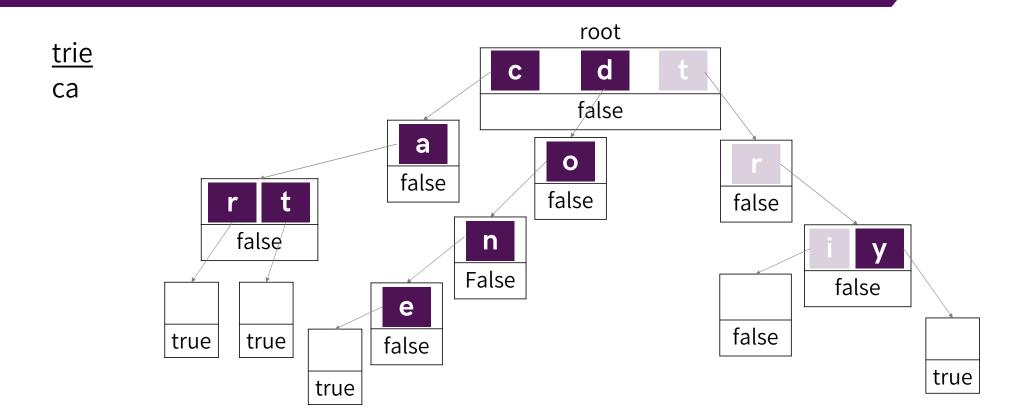
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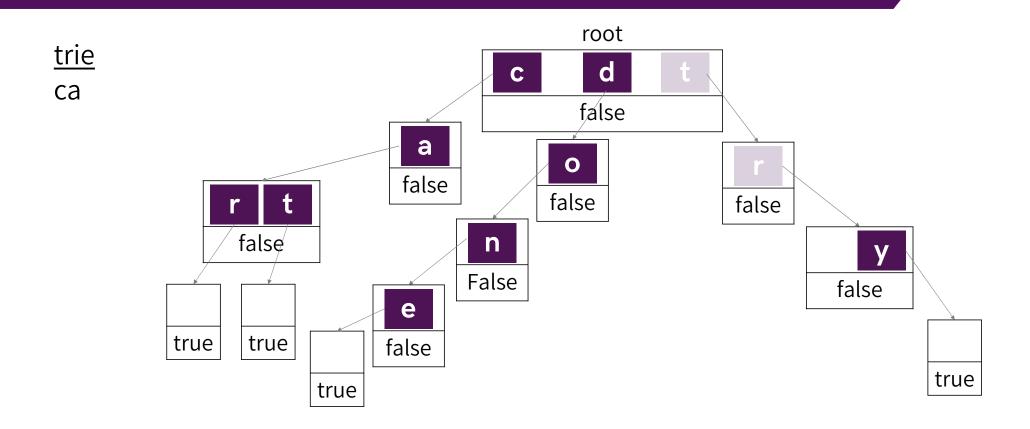
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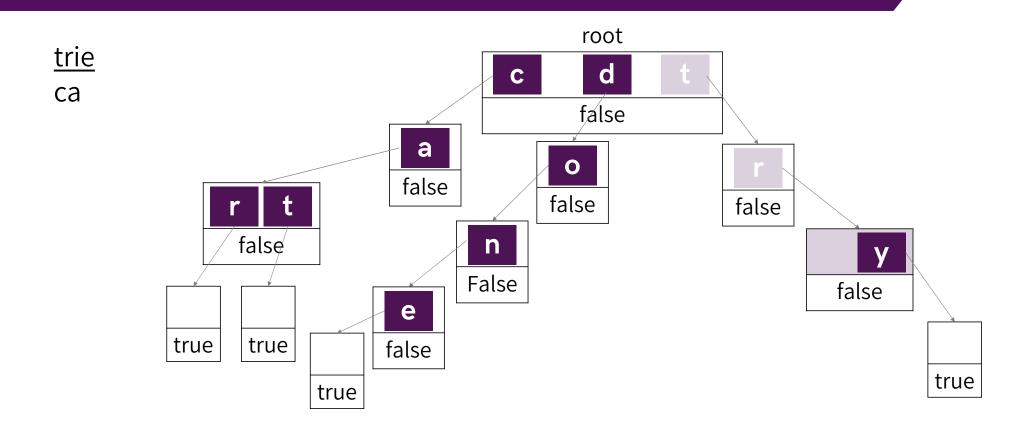


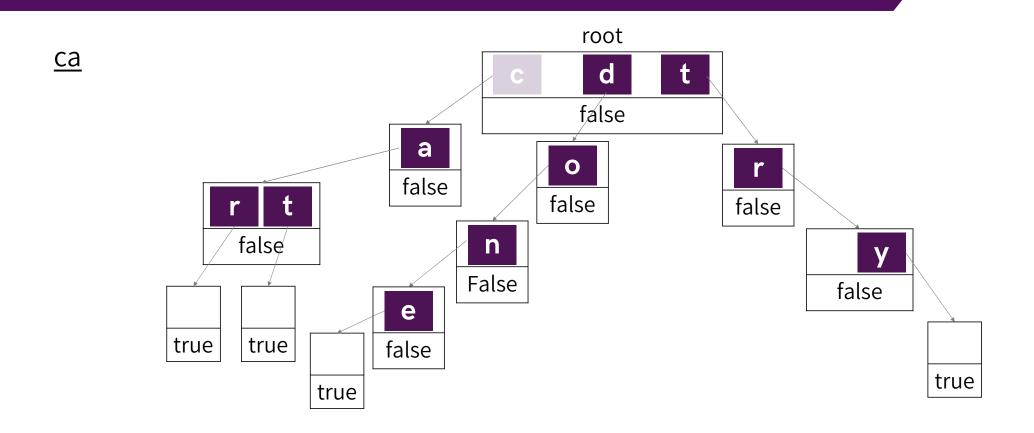


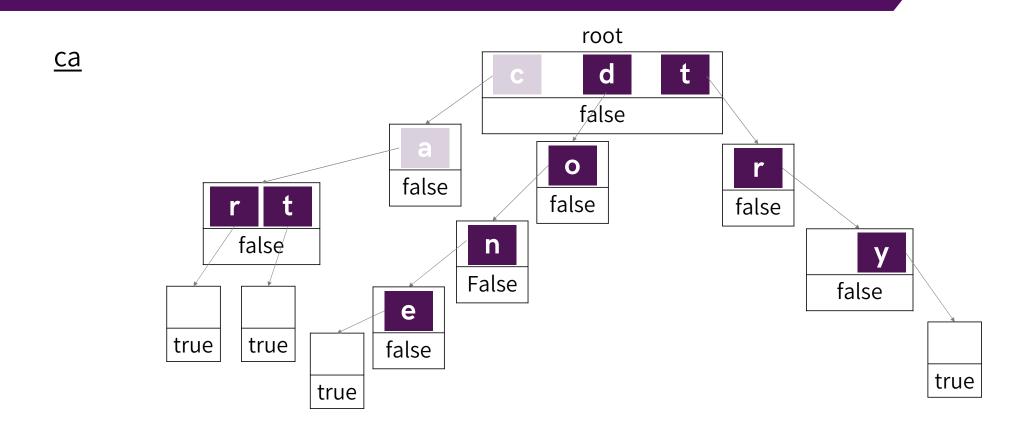
root

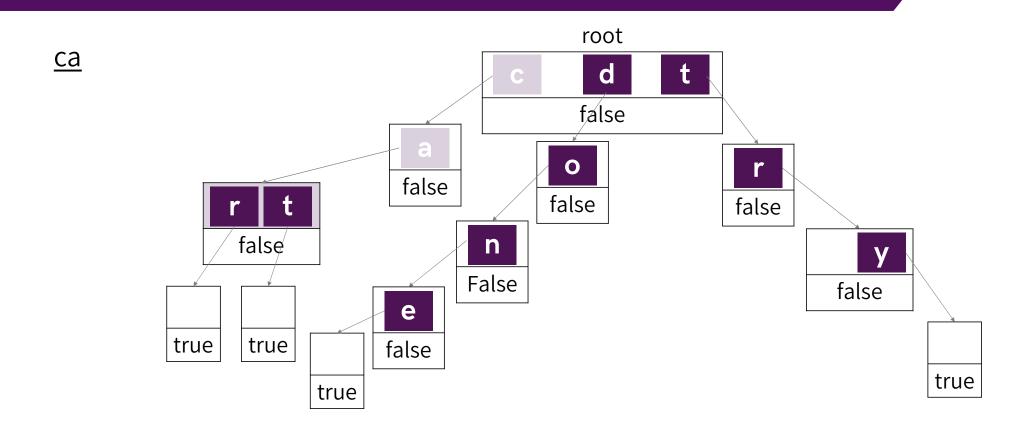












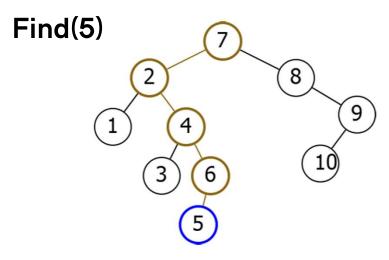
Alphabet Trie

- To determine if a given word is in the dictionary
- Only the letter that leads to a word are shown as branches
- Node's depth depends on the number of character required to distinguish this node from any other
- Implementation:
 - Array of 27 pointers indexed by letter
 - A linked list of pointers to the child nodes

Balanced Tree

BST Runtime

How long do Binary Search Tree operations take?



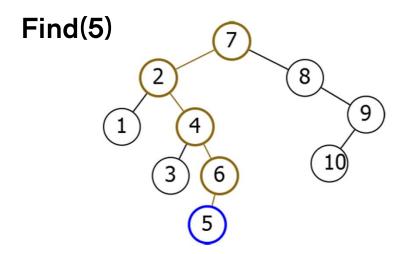
Number of operations=O(Depth)



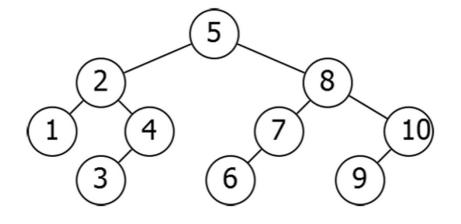
Depth can be as bad as **n**

BST Runtime

How long do Binary Search Tree operations take?



Number of operations=O(Depth)



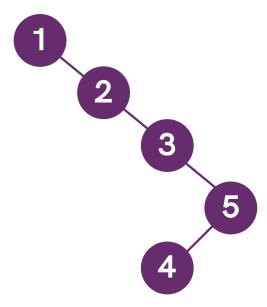
Depth can be much smaller

Balance

- Want left and right subtrees to have approximately the same size.
- Suppose perfectly balanced:
 - Each subtree half the size of its parent.
 - After log2(n) levels, subtree of size 1.
 - Operations run in O(log(n)) time.

Problem

Insertion and Deletion can destroy balance

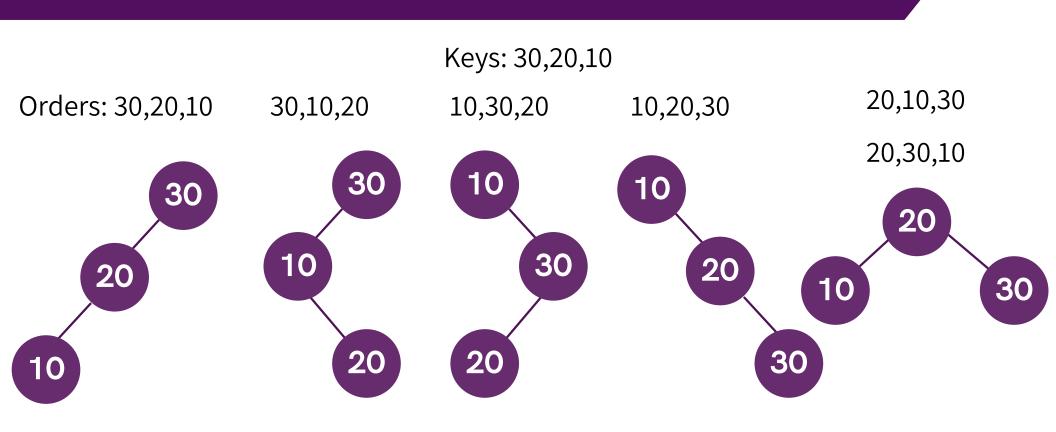


Rebalancing

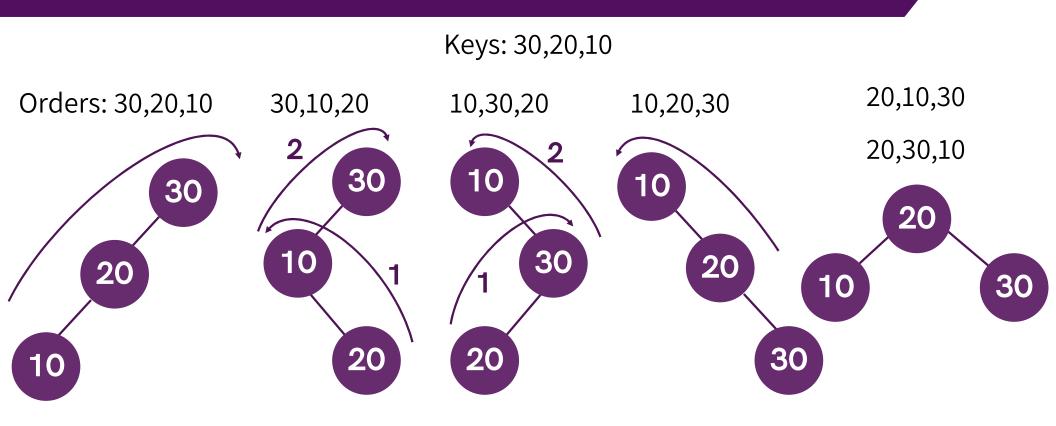
Idea: Rearrange tree to maintain balance.

Problem: How do we arrange tree while maintaining order?

Rotation



Rotation



AVL Tree

Height

The **height** of a node is the maximum depth of its subtree.

Height = 4

Height - Recursive Definition

N.Height equals:

- 1 if N is a leaf,
- 1+max(N.Left.Height, N.Right.Height) otherwise

Balance

- Height is a rough measure of subtree size.
- Want size of subtrees roughly the same.
- Force heights to be roughly the same.

AVL Tree (Adelson-Velskii and Landis)

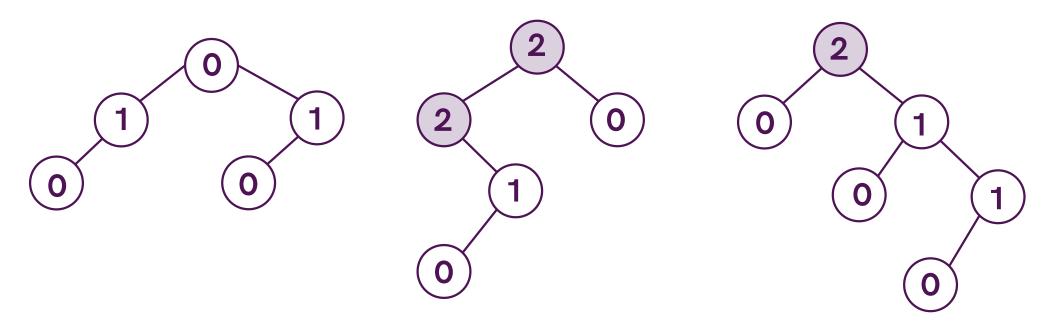
A BST with the following additional property: for every node N, the height of its left and right subtrees differ by at most 1.

 $|N.Left.Height - N.Right.Height| \le 1$

If you can maintain the AVL property, you can perform operations in O(log(n)) time.

Balance

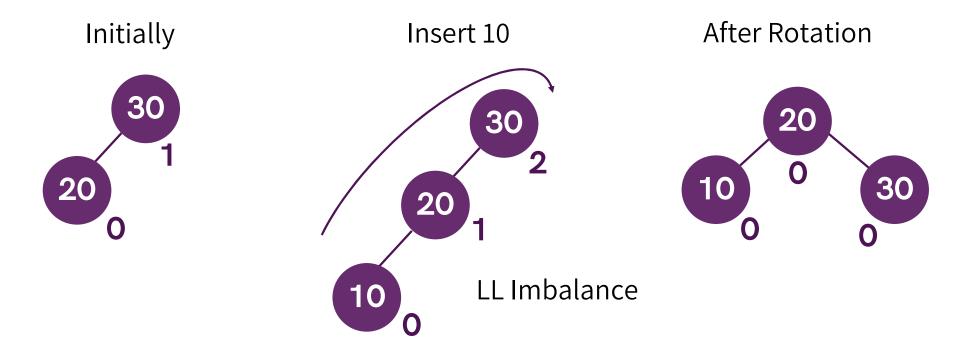
 $|N.Left.Height - N.Right.Height| \le 1$



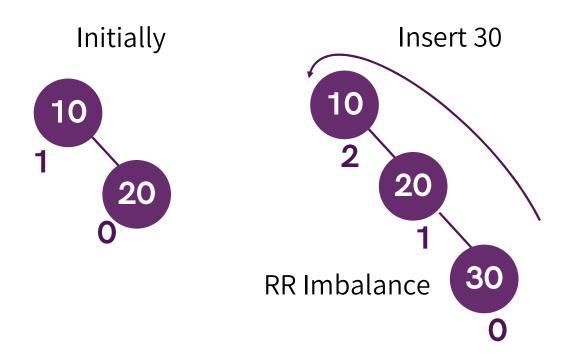
Unbalanced Cost

- 1. The extra node is in the left child of the left child of S (LL) Single Rotation
- 2. The extra node is in the right child of the left child of S (RL) 🛑 Double Rotation
- 3. The extra node is in the left child of the right child of S (LR) Double Rotation
- 4. The extra node is in the right child of the right child of S (RR) Single Rotation

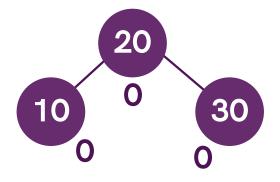
Single Rotation



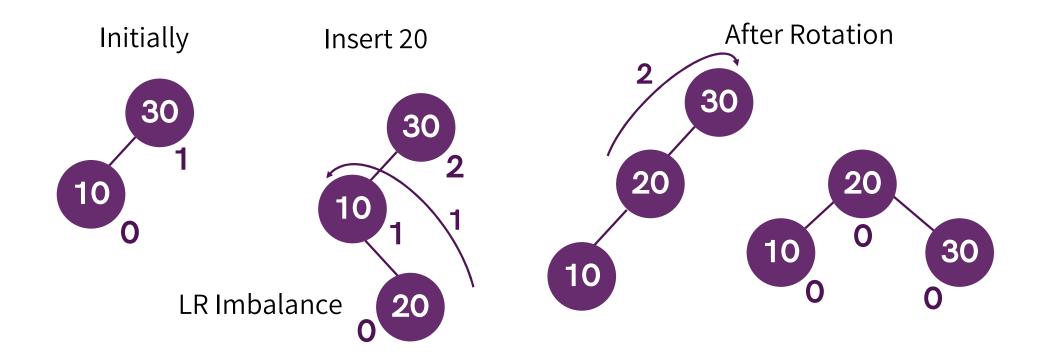
Single Rotation



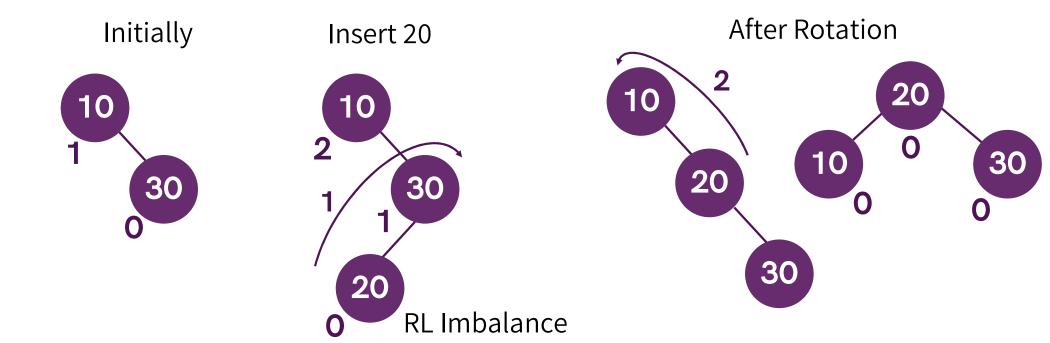




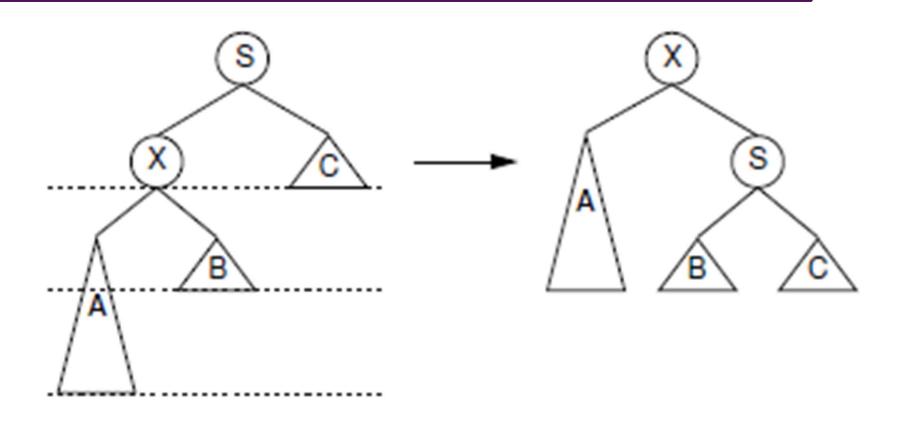
Double Rotation



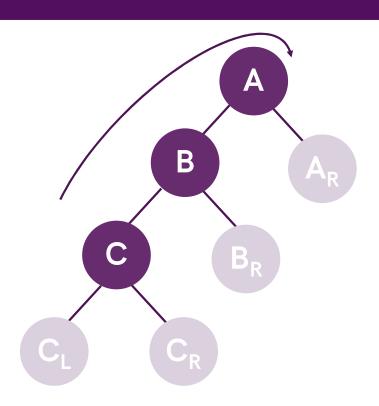
Double Rotation

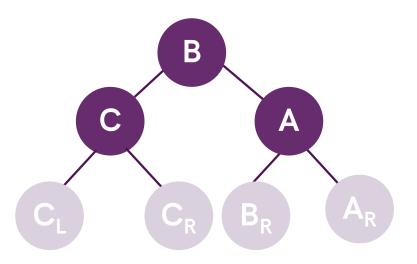


Single Rotation

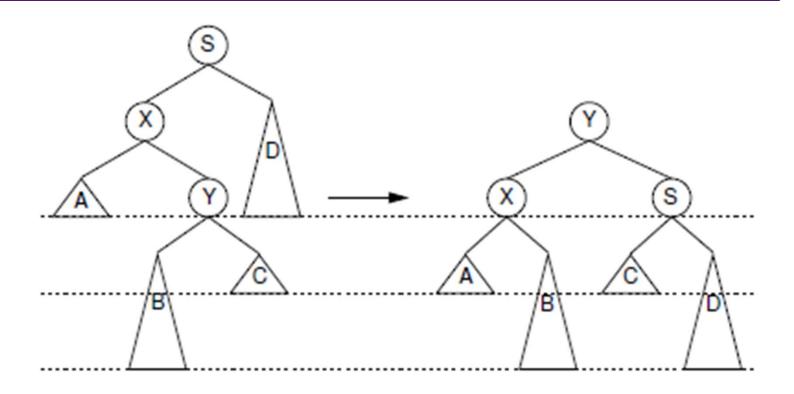


Single Rotation

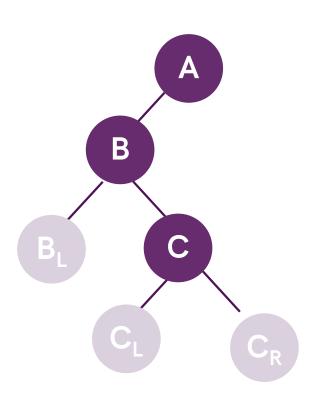


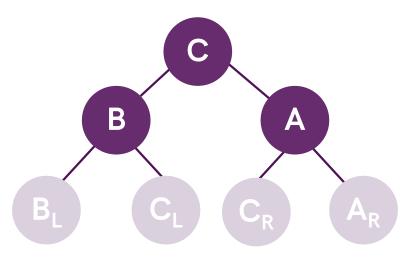


Double Rotation



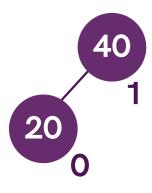
Double Rotation

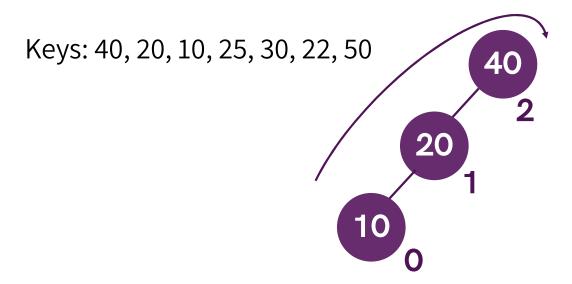


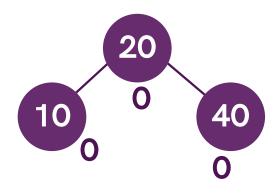


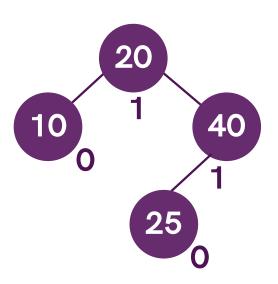
AVL Tree Implementation

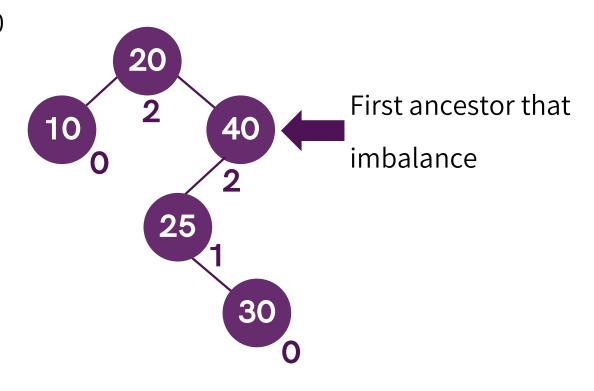
- Normal BST insert/delete
- Perform appropriate rotation on any node that is found to be unbalanced

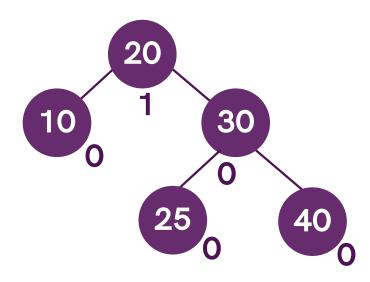


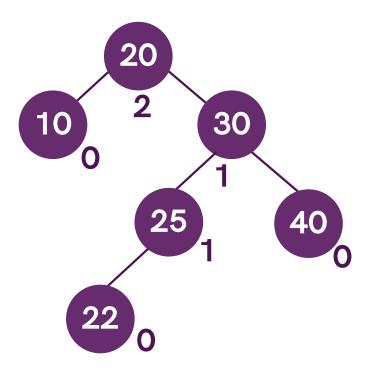


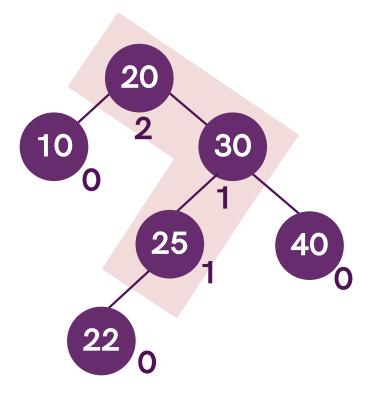


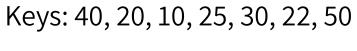


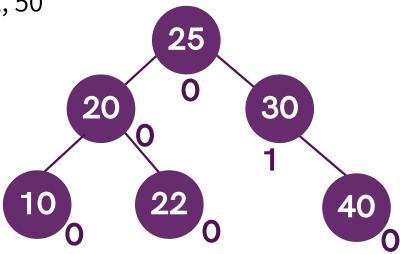


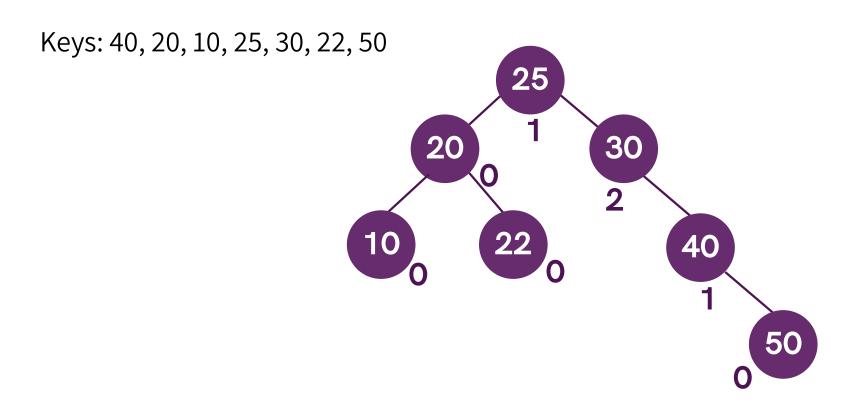


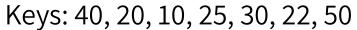


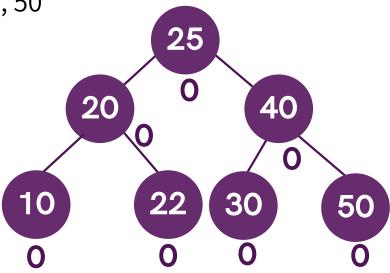




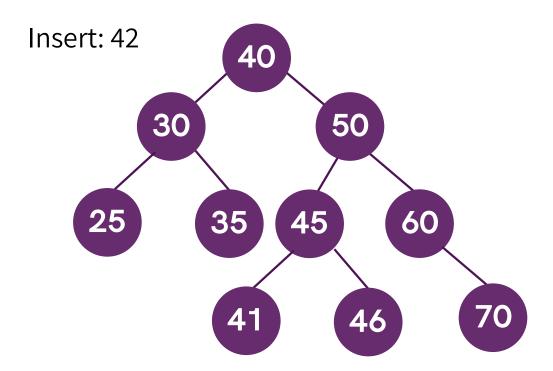








Latihan Soal





Tugas 4

Pemilik kedai kopi menyimpan data menu agar mudah dicari oleh pelanggan. Berikut data menu yang ingin disimpan:

Café Long Black

Café Cappuccino

Café Oat Cappuccino

Café Mocha

Café Latte

Vanilla Latte

Caramel Latte

Hazelnut Latte

Cinnamon Coffee Latte

Café Espresso

Café Macchiato

Café Piccolo

Café Flat White

Babyccino

- 1. Buatlah program untuk menyimpan data menu menggunakan trie!
- 2. Buatlah program untuk mencari apakah suatu item ada dalam menu!
- 3. Buatlah program untuk menghapus menu tertentu dari trie!

Terapkan modular programming dalam mengimplementasikan program