Punch holes in your large pages for benefits (performance, flexibility) with less overhead



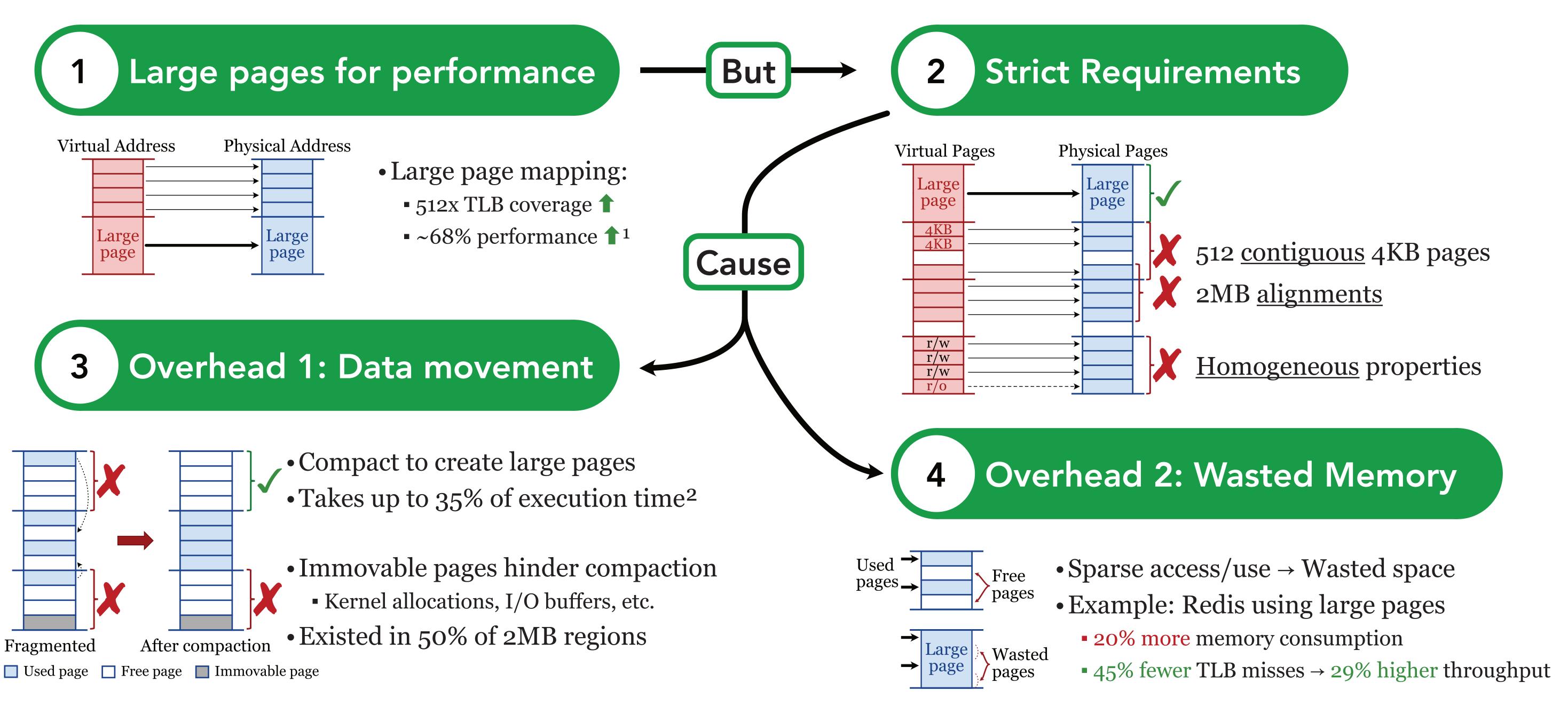


Perforated Pages: Supporting Fragmented Memory Allocation for Large Pages (ISCA 2020)



Chang Hyun Park, Sanghoon Cha, Bokyeong Kim, Youngjin Kwon, David Black-Schaffer, and Jaehyuk Huh



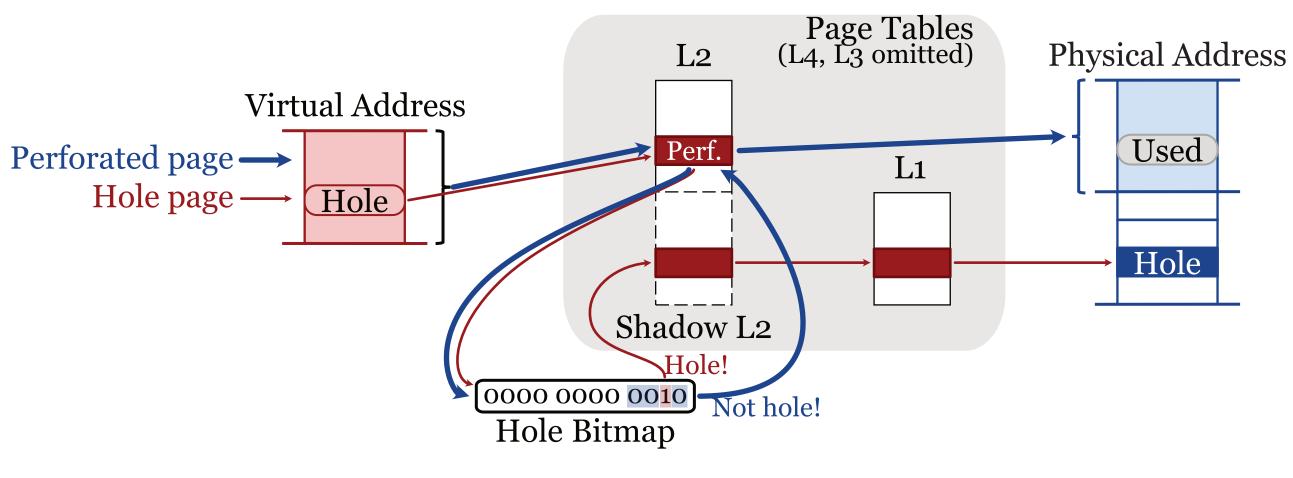


Our Solution: Perforated Pages

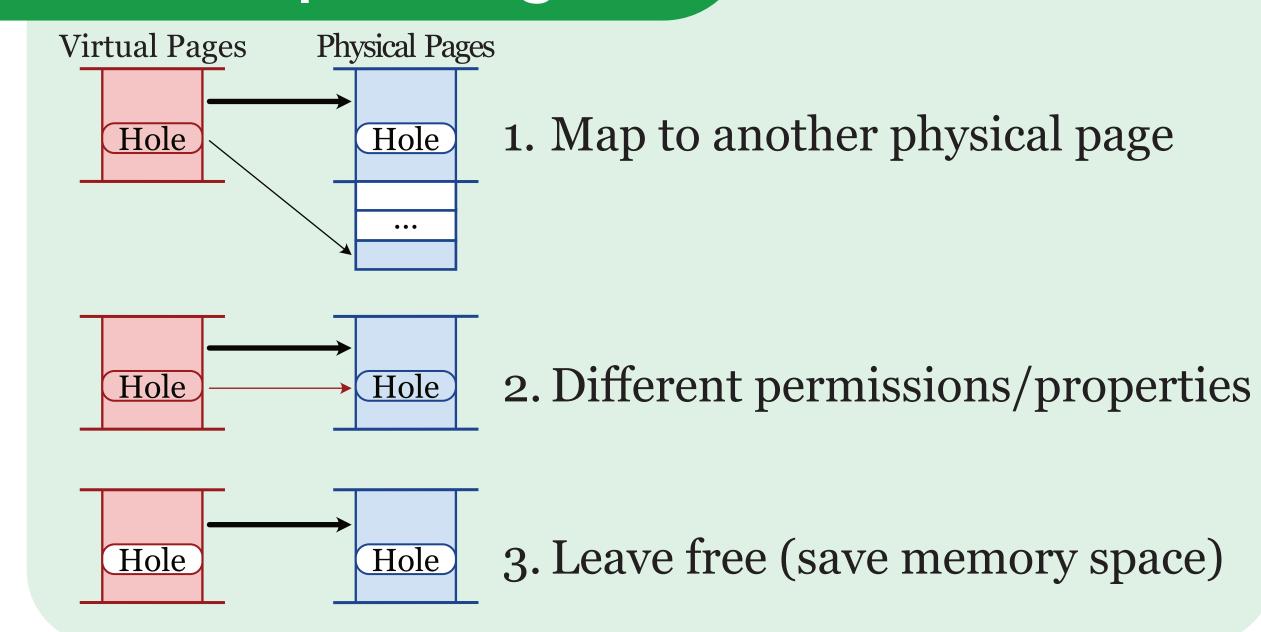
Physical Pages Virtual Pages Hole Hole

- Perforated page (2MB): maps common case
 - Large page-like coverage and mapping
- Hole Page (4KB): maps special case
- Flexibility for specific mappings

How it works - Page table



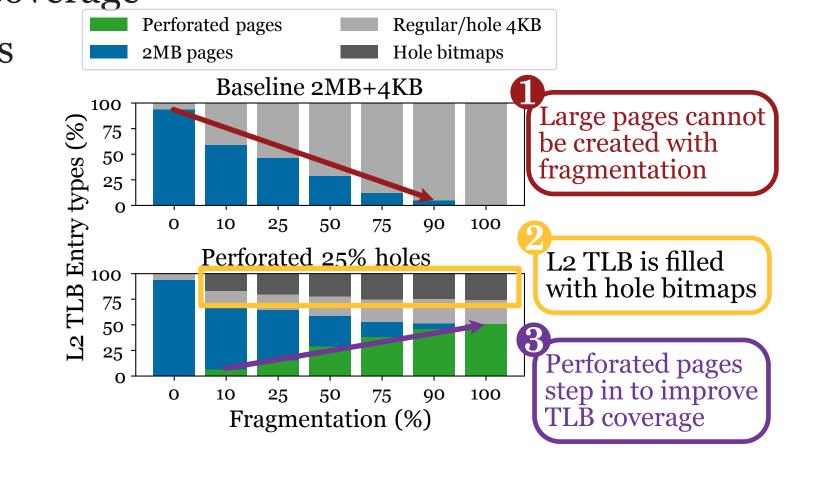
Example Usages +



- Hole Bitmap: Identifies hole pages stored in DRAM
- Shadow L2: provides an additional pointer in the page table
 - Main L2 PTE points to perforated page mapping (2MB)
 - Shadow L2 PTE points to L1 PT that holds hole page mappings (4KB)

How it works - L2 TLB

- Cache and translate perforated/hole pages at L2 TLB
- L2 TLB caches additional entries:
 - Perforated mappings: 2MB page-like coverage
 - Hole mappings: 4KB flexible mappings
 - Hole bitmaps: cached in L2 TLB
- Optimization: bitmap filter
- Identify non-hole regions
- Stored in perforated mappings



Evaluation Summary

