

# Seeing the Mountain

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*Seeing the mountain as the mountain;  
Seeing the mountain as not a mountain;  
Seeing the mountain as still a mountain.*

QINGYUAN XINGSI

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# Chapter 0

## Introduction

### 0.1 What is Algebraic Geometry?

Algebraic geometry is the study of geometric spaces that locally arise as solutions to polynomial equations. This study can be approached at two levels:

- **Local:** This involves studying the geometric properties of solution sets to polynomial equations. Since we work not only over  $\mathbb{R}$  or  $\mathbb{C}$  but potentially over arbitrary fields or rings, we must construct a more intrinsic notion of geometry associated with algebra—namely, the spectrum of a ring. This leads to the study of affine varieties and affine schemes, allowing a translation between algebra and geometry.
- **Synthetic:** Analogous to how manifolds generalize open subsets of  $\mathbb{R}^n$ , we study spaces that are locally isomorphic to affine varieties. This broader perspective leads to the study of varieties and schemes through various formal frameworks.

This lecture is part of an algebraic geometry course that emphasizes the second, synthetic level. However, the underlying question—“How can we construct and classify generalized spaces from certain building blocks?”—is relevant across many branches of geometry. Thus, this part can stand alone as a form of *post-modern*<sup>1</sup> synthetic geometry.

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<sup>1</sup>A commonly accepted definition is “after World War II.”

## 0.2 Essentialism vs. Structuralism

Synthetic geometry formalizes geometric concepts through axioms that directly address fundamental entities—such as points and lines—rather than relying on a background space like Cartesian coordinates. In contrast to the analytic viewpoint, where every geometric object is composed of points, synthetic geometry treats lines, curves, and other entities as primitive. The focus shifts to the **relationships** between these objects, such as “a point lies on a line” or “two lines intersect.”

This structuralism perspective emphasizes understanding objects through their interactions. To formalize these relationships, we use **category theory** and **sheaf theory**. Moreover, by viewing categories themselves as spaces, we uncover deeper geometric structures. This formalism, developed through algebraic geometry, now plays a central role in various fields: differential geometry, topology, quantum field theory, and beyond.

## 0.3 What is Geometry in Physics?

Quantum physics connects to both the local and synthetic levels of algebraic geometry:

- **Local:** The concept of the spectrum of a commutative ring originates from  $C^*$ -algebras and operator theory, both fundamental in quantum physics. The duality between algebra and geometry is already present in the Heisenberg and Schrödinger formulations.
- **Synthetic:** In quantum field theory, we study the space of field configurations (histories), which is infinite-dimensional and behaves irregularly. Nevertheless, we aim to treat it as a “manifold” to define metrics, path integrals, and differential forms. Synthetic geometry provides a rigorous framework for these constructions (e.g., via smooth sets). Gauge theory, supergeometry, and related topics can also be unified within this framework [2].

A concrete example is mirror symmetry, which links Gromov–Witten invariants with Hodge theory.

## 0.4 Course Plan

We will focus on three types of geometry and study them comparatively using synthetic methods:

- **Smooth Sets / Manifolds:** Built from open subsets of  $\mathbb{R}^n$ , these provide intuitive and geometric examples, serving as a bridge to classical geometry.
- **Simplicial Sets / Kan Complexes:** Constructed from simplices  $\Delta$ , these offer the simplest abstract examples.
- **Algebraic Sets / Schemes:** Built from spectra of commutative rings, these are the central objects of study in this course.





# Chapter 1

## Category

There are many references available for category theory; in this course, we follow [3]

### 1.1 Definition and Examples

**Definition 1.1.1** (Category). A **category**  $\mathcal{C}$  consists of:

- A class of objects  $\text{Ob}(\mathcal{C})$ .
- For each pair  $A, B$ , a set of morphisms  $\text{Hom}_{\mathcal{C}}(A, B)$ .
- A composition operation  $\circ$  of morphisms and identity morphisms  $\text{id}_A$  for each object  $A$ .

Subject to:

- **Associativity:**  $h \circ (g \circ f) = (h \circ g) \circ f$
- **Identity:**  $\text{id}_B \circ f = f = f \circ \text{id}_A$

For simplicity, we denote objects in a category  $\mathcal{C}$  by  $A \in \mathcal{C}$  and morphisms by  $f : A \rightarrow B$ . There are various ways to understand what a category is; let us explore this through examples:

**Example 1.1.2** (Concrete Category). A concrete category can be viewed as a collection of mathematical structures, where morphisms are maps that preserve those structures.

- **Set:** Objects are sets<sup>1</sup>; morphisms are functions between sets.
- **Grp:** Objects are groups; morphisms are group homomorphisms.
- **Vect:** Objects are vector spaces; morphisms are linear maps.
- **Top:** Objects are topological spaces; morphisms are continuous maps.

You can construct infinitely many examples from the mathematical structures you are familiar with.

At first glance, this abstraction may seem unnecessary. However, we can generalize familiar notions from set theory. For instance, consider a morphism  $f : A \rightarrow B$  in a category  $\mathcal{C}$ :

- **Isomorphism:**  $f$  is an *isomorphism* if there exists a morphism  $g : B \rightarrow A$  such that:

$$g \circ f = \text{id}_A \quad \text{and} \quad f \circ g = \text{id}_B \text{ (i.e. } g = f^{-1} \text{)}$$

In this case,  $A$  and  $B$  are said to be *isomorphic*.

- **Monomorphism:**  $f$  is a *monomorphism* (or *mono*) if for all morphisms  $g_1, g_2 : X \rightarrow A$ , we have:

$$f \circ g_1 = f \circ g_2 \Rightarrow g_1 = g_2$$

That is,  $f$  is left-cancellable.

- **Epimorphism:**  $f$  is an *epimorphism* (or *epi*) if for all morphisms  $h_1, h_2 : B \rightarrow Y$ , we have:

$$h_1 \circ f = h_2 \circ f \Rightarrow h_1 = h_2$$

That is,  $f$  is right-cancellable.

**Exercise 1.1.1.** Verify that in the category **Set**, these definitions correspond to bijections, injections, and surjections, respectively.

Besides the concrete categories, in which objects have their own meaning, we can have abstract categories where objects are meaningless without the context of the category.

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<sup>1</sup>Since the collection of all sets does not itself form a set, we refer instead to a *class* of objects. However, if we restrict our attention to *small* sets, then the collection of objects can be treated as a set. In this course, we will ignore set-theoretic subtleties and proceed informally.

**Example 1.1.3** (Classifying Category). A category can be viewed as an algebraic structure generalizing a monoid<sup>2</sup>. Indeed, for any object  $A \in \mathcal{C}$ , the set of endomorphisms  $\text{Hom}_{\mathcal{C}}(A, A)$  forms a monoid. Conversely, any monoid  $M$  can be associated with a **classifying category**  $BM$ , which has a single object  $\bullet$  and morphisms  $\text{Hom}_{BM}(\bullet, \bullet) = M$ . In particular, for any group  $G$ , we obtain a **classifying space**  $BG$ , where all morphisms are isomorphisms.

We refer to  $BG$  as a space because we can interpret its isomorphisms as paths in a certain topological space:

**Example 1.1.4** (Groupoid). A **groupoid** is a category in which every morphism is an isomorphism. Given a groupoid  $\mathcal{X}$ , we can construct its **geometric realization**  $|\mathcal{X}|$  as follows:

- Take the objects  $x \in \mathcal{X}$  as points.
- For each morphism  $f : x \rightarrow y$ , attach a segment from point  $x$  to point  $y$ .
- For each relation  $f \circ g = h$ , attach a triangle with edges labeled by  $f$ ,  $g$ , and  $h$ .
- Continue this process for higher-dimensional cells ...

We will formalize this construction when we introduce simplicial sets.

**Exercise 1.1.2.** Show that the set of connected components  $\pi_0(|\mathcal{X}|, x)$  corresponds bijectively to the isomorphism class of  $x$  in  $\mathcal{X}$ . Furthermore, observe that  $\text{Hom}_{\mathcal{X}}(x, x)$  is a group, and prove that the fundamental group  $\pi_1(|\mathcal{X}|, x) \cong \text{Hom}_{\mathcal{X}}(x, x)$ .

Conversely, for a topological space  $S$ , we can define the **fundamental groupoid**  $\Pi_1(S)$ , where:

- Objects are points of  $S$ .
- Morphisms  $\text{Hom}_{\Pi_1(S)}(a, b)$  are homotopy classes of paths from  $a$  to  $b$ .

**Exercise 1.1.3.** Describe the composition law in  $\Pi_1(S)$  and verify that it satisfies the axioms of a groupoid.

The examples above illustrate that morphisms in a category can carry rich structure. On the other hand, if morphisms are trivial (i.e., at most one between any two objects), we obtain a partially ordered set:

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<sup>2</sup>An algebraic structure similar to a group, but without requiring inverses

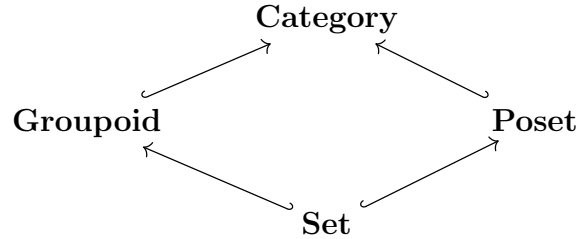
**Example 1.1.5** (Poset). Let  $(P, \leq)$  be a partially ordered set. We can regard  $P$  as a category as follows:

- **Objects:** Elements of  $P$ .
- **Morphisms:** For  $x, y \in P$ , there exists a unique morphism  $f : x \rightarrow y$  if and only if  $x \leq y$ .
- **Composition:** If  $x \leq y$  and  $y \leq z$ , then  $x \leq z$ , so the morphism  $x \rightarrow z$  is the composition of  $x \rightarrow y$  and  $y \rightarrow z$ .
- **Identity:** For each  $x \in P$ , the identity morphism  $\text{id}_x : x \rightarrow x$  corresponds to the reflexivity  $x \leq x$ .

This category is called *thin*, meaning there is at most one morphism between any two objects.

A frequently used example is the poset of open subsets of a topological space  $S$ , denoted  $(\text{Op}(S), \subseteq)$ .

Note that a set  $A$  can be viewed as a category in two distinct ways: either as a trivial groupoid or as a trivial poset. And a category can be viewed as a combination of this two case<sup>3</sup>.



*Remark 1.1.6.* In the definition of a category, if we allow the morphism set  $\text{Hom}_{\mathcal{C}}(A, B)$  to carry additional structure—such as an Abelian group, vector space, groupoid, topological space, or even another category—we obtain an **enriched category**. In fact, it is often more natural to think of  $\text{Hom}_{\mathcal{C}}(A, B)$  as the set of connected components of a space:

$$\text{Hom}_{\mathcal{C}}(A, B) \cong \pi_0 \text{Map}_{\mathcal{C}}(A, B).$$

This reflects a general principle of the **Univalence Foundation**: mathemat-

<sup>3</sup>It is useful to think category as an oriented graph

ical structures should be treated as spaces (or types) from the outset, and the classical set-theoretic version can be recovered by taking the set of connected components. We will explore how to identify and work with these underlying geometric structures later in the course.

## 1.2 Functor

To define a map between two categories, it is natural to require that such a map respects the structure of morphisms.

**Definition 1.2.1** (Functor). Let  $\mathcal{C}$  and  $\mathcal{D}$  be categories. A **functor**  $F : \mathcal{C} \rightarrow \mathcal{D}$  consists of:

- A function that assigns to each object  $A \in \mathcal{C}$  an object  $F(A) \in \mathcal{D}$ .
- A function that assigns to each morphism  $f : A \rightarrow B$  in  $\mathcal{C}$  a morphism  $F(f) : F(A) \rightarrow F(B)$  in  $\mathcal{D}$ .

such that:

- **Identity preservation:**  $F(\text{id}_A) = \text{id}_{F(A)}$ .
- **Composition preservation:** For all  $f : A \rightarrow B$  and  $g : B \rightarrow C$  in  $\mathcal{C}$ ,

$$F(g \circ f) = F(g) \circ F(f).$$

We denote the set (later, the category) of functors from  $\mathcal{C}$  to  $\mathcal{D}$  by  $\text{Fun}(\mathcal{C}, \mathcal{D})$ .

**Example 1.2.2** (Functors Between Concrete Categories). Functors between concrete categories respect the underlying structures:

- **Forgetful Functor:**

$$U : \mathbf{Grp} \rightarrow \mathbf{Set}$$

assigns to each group its underlying set, and to each group homomorphism the same function viewed as a map of sets. Similar forgetful functors exist for **Vect**, **Top**, etc.

- **Free Functor:**

$$\text{Free} : \mathbf{Set} \rightarrow \mathbf{Grp}$$

assigns to each set the free group it generates, and to each function between sets the induced group homomorphism.

**Exercise 1.2.1.** Verify that  $\text{Free}$  is indeed a functor. Then show:

$$\text{Hom}_{\mathbf{Grp}}(\text{Free}(A), G) \cong \text{Hom}_{\mathbf{Set}}(A, U(G)).$$

How would you define  $\text{Free}$  for  $\mathbf{Vect}$ ?

- **Discrete Functor:**

$$\text{Disc} : \mathbf{Set} \rightarrow \mathbf{Top}$$

assigns to each set the discrete topological space, and to each function the same map viewed as continuous.

- **Connected Components:**

$$\pi_0 : \mathbf{Top} \rightarrow \mathbf{Set}$$

assigns to each topological space its set of connected components.

**Exercise 1.2.2.** Show:

$$\text{Hom}_{\mathbf{Top}}(\text{Disc}(A), S) \cong \text{Hom}_{\mathbf{Set}}(A, U(S)), \quad \text{Hom}_{\mathbf{Top}}(S, \text{Disc}(A)) \cong \text{Hom}_{\mathbf{Set}}(\pi_0(S), A).$$

**Exercise 1.2.3.** 1. Let  $G, H$  be groups (or more generally, monoids). Show that functors  $F : BG \rightarrow BH$  correspond bijectively to group (monoid) homomorphisms  $f : G \rightarrow H$ .

2. Show that a functor  $F : BG \rightarrow \mathbf{Vect}$  correspond bijectively to representation of  $G$ .

**Exercise 1.2.4.** 1. Let  $\mathcal{X}, \mathcal{Y}$  be groupoids. Show that a functor  $F : \mathcal{X} \rightarrow \mathcal{Y}$  induces a continuous map between their geometric realizations:

$$|F| : |\mathcal{X}| \rightarrow |\mathcal{Y}|.$$

Let  $\mathbf{Grpd}$  be the category of groupoids with functors as morphisms. Show that geometric realization defines a functor:

$$|\cdot| : \mathbf{Grpd} \rightarrow \mathbf{Top}.$$

2. Similarly, show that the fundamental groupoid construction defines a functor:

$$\Pi_1 : \mathbf{Top} \rightarrow \mathbf{Grpd}.$$

3. \* Show the adjunction:

$$\mathrm{Hom}_{\mathbf{Top}}(S, |\mathcal{X}|) \cong \mathrm{Hom}_{\mathbf{Grpd}}(\Pi_1(S), \mathcal{X}).$$

**Exercise 1.2.5.** Let  $P, Q$  be posets. Show that a functor  $F : P \rightarrow Q$  is equivalent to an order-preserving function.

**Example 1.2.3** (Ring–Space Correspondence). Given a topological space  $X$ , its ring of real continuous functions  $C(X)$  defines a contravariant functor. A map  $p : X \rightarrow Y$  induces a pullback:

$$p^* : C(Y) \rightarrow C(X), \quad f \mapsto f \circ p.$$

To formalize this, we introduce the **opposite category**  $\mathcal{C}^{\mathrm{op}}$ , which has the same objects as  $\mathcal{C}$  but reverses the direction of morphisms:

$$\mathrm{Hom}_{\mathcal{C}^{\mathrm{op}}}(A, B) = \mathrm{Hom}_{\mathcal{C}}(B, A).$$

Then,

$$C(-) : \mathbf{Top}^{\mathrm{op}} \rightarrow \mathbf{Ring}$$

is a functor.

*Remark 1.2.4.* A central theme in algebraic geometry is reversing this functor: constructing a space from a commutative ring. That is, defining a functor:

$$\mathrm{Spec} : \mathbf{Ring}^{\mathrm{op}} \rightarrow \mathbf{Top},$$

such that:

$$\mathrm{Hom}_{\mathbf{Ring}}(R, C(X)) \cong \mathrm{Hom}_{\mathbf{Top}}(X, \mathrm{Spec}(R)).$$

*Exercise 1.2.6.* \* Show that the underlying set of  $\mathrm{Spec}(R)$  is:

$$U(\mathrm{Spec}(R)) = \mathrm{Hom}_{\mathbf{Ring}}(R, \mathbb{R}).$$

What topology should be given to  $\mathrm{Spec}(R)$ ?

In practice, we impose additional structure to make this correspondence well-behaved:

- Between  $C^*$ -algebras and locally compact Hausdorff spaces, via the Gelfand representation theorem—where the term “spectrum” originates.
- Between commutative rings and locally ringed spaces, which is foundational in algebraic geometry.

**Example 1.2.5** (Hom Functor). For a fixed object  $A$  in a category  $\mathcal{C}$ , the **Hom functor** is defined as:

$$h^A := \text{Hom}_{\mathcal{C}}(A, -) : \mathcal{C} \rightarrow \mathbf{Set},$$

which assigns to each object  $B$  the set  $\text{Hom}_{\mathcal{C}}(A, B)$ , and to each morphism  $f : B \rightarrow C$  the function:

$$h^A(f) : \text{Hom}_{\mathcal{C}}(A, B) \rightarrow \text{Hom}_{\mathcal{C}}(A, C), \quad g \mapsto f \circ g.$$

Similarly, we define the contravariant version:

$$h_A := \text{Hom}_{\mathcal{C}}(-, A) : \mathcal{C}^{\text{op}} \rightarrow \mathbf{Set}.$$

**Exercise 1.2.7.** Show that:

- $f$  is a monomorphism if and only if  $h^A(f)$  is injective for all  $A$ .
- $f$  is an epimorphism if and only if  $h_A(f)$  is injective for all  $A$ .
- $f$  is an isomorphism if and only if both  $h^A(f)$  and  $h_A(f)$  are bijective for all  $A$ .

*Remark 1.2.6.* We can think of objects in  $\mathcal{C}$  as “test objects,” and a functor as a way of encoding data about how these tests behave. For example, in a sigma model with target space  $M$ , let  $\mathcal{C}$  be the category of space-times. For each space-time  $\Sigma$ , the collection of fields  $\text{Map}(\Sigma, M)$  defines such a functor. The fundamental question is: given such a functor, can we reconstruct the underlying “space”?

## 1.3 Presheaf

We now introduce a central concept for understanding “generalized spaces” in this course.



**Definition 1.3.1** (Presheaf). Let  $\mathcal{C}$  be a category. A **presheaf** on  $\mathcal{C}$  is a functor:

$$F : \mathcal{C}^{\text{op}} \rightarrow \mathbf{Set}$$

That is,  $F$  assigns:

- To each object  $U$  in  $\mathcal{C}$ , a set  $F(U)$ .
- To each morphism  $f : V \rightarrow U$  in  $\mathcal{C}$ , a function:

$$F(f) : F(U) \rightarrow F(V)$$

such that:

- $F(\text{id}_U) = \text{id}_{F(U)}$
- $F(g \circ f) = F(f) \circ F(g)$  for composable morphisms  $f$  and  $g$  in  $\mathcal{C}$

Let  $\text{PSh}(\mathcal{C}) = \text{Fun}(\mathcal{C}^{\text{op}}, \mathbf{Set})$  denote the category of presheaves on  $\mathcal{C}$ .

*Remark 1.3.2.* To understand why presheaves represent generalized spaces, consider the analogy with distributions as generalized functions:

- We begin with a space of test functions. Let  $\mathcal{D}(\Omega)$  denote the space of smooth functions with compact support in  $\Omega$ . There is a natural pairing:

$$\langle f, g \rangle := \int_{\Omega} fg$$

This allows us to associate to each test function  $f \in \mathcal{D}(\Omega)$  a continuous functional  $T_f \in \mathcal{D}'(\Omega)$  defined by:

$$T_f(\phi) := \int_{\Omega} f\phi.$$

The space  $\mathcal{D}'(\Omega)$  of distributions is thus a space of generalized functions.

- Similarly, starting from a category of test spaces  $\mathcal{C}$ , we consider the Hom functor:

$$\mathrm{Hom}_{\mathcal{C}}(-, -) : \mathcal{C}^{\mathrm{op}} \times \mathcal{C} \rightarrow \mathbf{Set}.$$

For each object  $A \in \mathcal{C}$ , the presheaf  $h_A := \mathrm{Hom}_{\mathcal{C}}(-, A)$  is *representable*. Presheaves generalize these representable ones, just as distributions generalize smooth functions.

- In solving differential equations, we often first obtain a distributional solution, and then study its *regularity*—how closely it resembles a smooth function locally—allowing us to conclude that the distribution is an actual function.
- In geometry, for example in moduli problems, we may first define a (pre)sheaf<sup>a</sup> and then study its *representability*—whether it locally resembles a test space—allowing us to conclude that we have constructed an actual space.

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<sup>a</sup>sheaf to presheaf, is like continuous functional to functional, we ask some continuous condition on presheaf

As we've seen, for any  $A \in \mathcal{C}$ , the presheaf  $h_A$  is called representable. Just as there are distributions that are not smooth functions, there are presheaves that are not representable. Here are some concrete examples:

**Example 1.3.3** (Presheaf of Functions). Let  $X$  be a topological space. For each open set  $U \in \mathrm{Op}(X)$ , the set of continuous functions  $C(U)$  defines a presheaf  $C \in \mathrm{PSh}(\mathrm{Op}(X))$ . For  $U \subseteq V$ , we have a restriction map  $C(V) \rightarrow C(U)$ . More generally, for any topological space  $Y$ , we can define a presheaf  $C(-, Y) \in \mathrm{PSh}(\mathrm{Op}(X))$ . Similarly, one can define presheaves of smooth, analytic, or locally constant functions.

**Example 1.3.4** (Presheaf of Sections). Let  $E \rightarrow X$  be a vector bundle over a topological space  $X$ . Define a presheaf  $\mathcal{F}$  by assigning to each open set  $U \subseteq X$  the set of continuous (or smooth) sections of  $E$  over  $U$ :

$$\mathcal{F}(U) = \{s : U \rightarrow E \mid s \text{ is a section of } E \text{ over } U\}.$$

**Example 1.3.5** (Smooth Set). Let **Cart** be the category of Cartesian spaces, with objects  $\mathbb{R}^n$  and morphisms given by smooth maps. A presheaf on **Cart** is called a **smooth set**.

- **Manifolds:** For any smooth manifold  $M$ , define a presheaf  $M(\mathbb{R}^n) := \{f : \mathbb{R}^n \rightarrow M \mid f \text{ is smooth}\}$ . Thus, every manifold defines a smooth set.

- **Differential Forms:** Consider the presheaf  $\Omega^k$  assigning to each  $\mathbb{R}^n$  the space of  $k$ -forms  $\Omega^k(\mathbb{R}^n)$ , with pullbacks along smooth maps. Show that  $\Omega^k$  is not representable by a manifold.

**Example 1.3.6** (Simplicial Set). example

**Example 1.3.7** (Algebraic Set). In algebraic geometry, we treat  $\mathbf{Ring}^{\text{op}}$  as the category of test spaces. A presheaf  $\mathcal{F} \in \text{PSh}(\mathbf{Ring}^{\text{op}}) = \text{Fun}(\mathbf{Ring}, \mathbf{Set})$  is then a functor from rings to sets, which we call an **algebraic set**.

- **Affine Variety:** We begin with the geometry of zero set of polynomials. Let  $P_1, \dots, P_m \in \mathbb{Z}[x_1, \dots, x_n]$  be polynomials. Define a functor  $V_P : \mathbf{Ring} \rightarrow \mathbf{Set}$  by:

$$V_P(R) = \{(r_1, \dots, r_n) \in R^n \mid P_i(r_1, \dots, r_n) = 0 \text{ for all } i\}.$$

**Exercise 1.3.1.** Let  $R_P = \mathbb{Z}[x_1, \dots, x_n]/(P_1, \dots, P_m)$ . Show that  $V_P = h^{R_P}$ , i.e.,  $V_P(R) = \text{Hom}_{\mathbf{Ring}}(R_P, R)$ .

- **Projective Space:** Then we consider a non-representable example. Define a functor  $\mathbb{P}^n : \mathbf{Ring} \rightarrow \mathbf{Set}$  by:

$$\mathbb{P}^n(R) = \{(r_0, \dots, r_n) \in R^{n+1} \mid \exists (u_0, \dots, u_n) \in R^{n+1} \text{ s.t. } \sum u_i r_i = 1\} / \sim,$$

where  $\sim$  identifies tuples under scalar multiplication by units in  $R$ .

**Exercise 1.3.2.** Show that this defines a functor. Compare this with the definition of projective space in differential geometry when  $R = \mathbb{R}$ .

As we indicated earlier,  $\text{PSh}(\mathcal{C})$  should itself form a category. Let us now determine what the morphisms between presheaves ought to be.

Since representable presheaves  $h_A$  and  $h_B$  are thought of as generalized spaces, it is natural to expect that morphisms between them should reflect the morphisms in the original category  $\mathcal{C}$ . Indeed, given a morphism  $f : A \rightarrow B$  in  $\mathcal{C}$ , we can define a map of sets:

$$f \circ - : h_A(C) \rightarrow h_B(C), \quad \text{for each } C \in \mathcal{C},$$

However, for this to define a morphism of presheaves, these maps must be compatible with the structure of the presheaves—i.e., they must commute with the restriction maps for all choices of  $C$  and morphisms between them.

This leads us naturally to the definition of morphisms between presheaves as *natural transformations*, which ensure such compatibility across the entire category.

**Definition 1.3.8** (Morphisms of Presheaves). A *morphism of presheaves*  $\varphi : \mathcal{F} \rightarrow \mathcal{G}$  is a natural transformation between the functors  $\mathcal{F}$  and  $\mathcal{G}$ . That is, for each object  $U$  in  $\mathcal{C}$ , there is a function  $\varphi_U : \mathcal{F}(U) \rightarrow \mathcal{G}(U)$  such that for every morphism  $f : V \rightarrow U$  in  $\mathcal{C}$ , the following diagram commutes:

$$\begin{array}{ccc} \mathcal{F}(U) & \xrightarrow{\varphi_U} & \mathcal{G}(U) \\ \mathcal{F}(f) \downarrow & & \downarrow \mathcal{G}(f) \\ \mathcal{F}(V) & \xrightarrow{\varphi_V} & \mathcal{G}(V) \end{array}$$

**Exercise 1.3.3.** 1. Show that a morphism  $f : A \rightarrow B$  in  $\mathcal{C}$  induces a morphism of presheaves  $\mathfrak{z}(f) : h_A \rightarrow h_B$ . This defines a functor:

$$\mathfrak{z} : \mathcal{C} \rightarrow \text{PSh}(\mathcal{C}), \quad A \mapsto h_A,$$

called the *Yoneda embedding*.

2. \* (Yoneda Lemma) Show that, for  $A \in \mathcal{C}$  and  $F \in \text{PSh}(\mathcal{C})$  there is a bijection:

$$F(A) \cong \text{Hom}_{\text{PSh}(\mathcal{C})}(h_A, F).$$

As a corollary, for any  $A, B \in \mathcal{C}$  there is a bijection of morphisms:

$$\mathfrak{z}(-) : \text{Hom}_{\mathcal{C}}(A, B) \cong \text{Hom}_{\text{PSh}(\mathcal{C})}(h_A, h_B).$$

This lemma is extremely important. It justifies embedding the category  $\mathcal{C}$  into the category of presheaves  $\text{PSh}(\mathcal{C})$  in a way that fully preserves its morphism structure. In other words,  $\mathcal{C}$  can be faithfully represented inside  $\text{PSh}(\mathcal{C})$  via the Yoneda embedding.

**Exercise 1.3.4.** 1. Let  $M$  be a smooth manifold viewed as a smooth set. Show that every differential form  $\omega \in \Omega^k(M)$  induces a morphism of smooth sets:

$$\omega : M \rightarrow \Omega^k.$$

2. \* Then prove that this defines a bijection:

$$\Omega^k(M) \cong \text{Hom}_{\text{sm}}(M, \Omega^k).$$

# Chapter 2

## Sheaf

As we mentioned early, presheaf is analogue of "linear functional", to get a category of generalized space, we need to impose the "continuous" condition, And sheaf is such "continuous" presheaf.

In analysis, continuous means preserve the limit, i.e.  $f(\lim x_i) = \lim f(x_i)$ . So we should also define limit in category, in some sense it describes how to approximate an object by others.

Some reference can be found in [1].

### 2.1 Limit and Colimit

Let's begin with the easiest example of analysis: the limit of an increasing sequence.

If we view the real number poset  $(\mathbb{R}, \leq)$  as a category. Then an increasing sequence is an order preserving map from  $(\mathbb{N}, \leq)$  to  $\mathbb{R}$ . i.e a functor  $a_{(-)} : \mathbb{N} \rightarrow \mathbb{R}$ . Let us unwind the definition of the limit  $\lim a_i$ : it is the supremum of  $\{a_i\}$ , i.e

$$\forall b \in \mathbb{R}, \forall i \in \mathbb{N}, a_i \leq b \Leftrightarrow \lim a_i \leq b$$

Recall for the poset category, morphism  $\text{Hom}_{\mathbb{R}}(a, b)$  can be seen as the proofs of proposition  $a \leq b$ : if there a morphism,  $a \leq b$  is true, otherwise the proofs is empty, it is false. So we can rewrite it as

$$\forall b \in \mathbb{R}, \prod_{i \in \mathbb{N}} \text{Hom}_{\mathbb{R}}(a_i, b) \cong \text{Hom}_{\mathbb{R}}(\lim a_i, b)$$

So we can think the limit is formally defined as a pre(co)sheaf  $b \mapsto \lim h^{a_i}(b)$ , and then if we can find an object who represent this pre(co)sheaf as  $h^{\lim a_i}$ , the limit exists as this object.

Then next example we consider a functor  $X_{(-)} : \mathbb{N} \rightarrow \mathbf{Set}$ , we should intuitively think its limit is  $\bigcup_{i \in \mathbb{N}} X_i$ , in this case it is called **Colimit**. But if we compare to the previous formula, we just get an inclusion:

$$\forall A \in \mathbf{Set}, \prod_{i \in \mathbb{N}} \mathrm{Hom}_{\mathbf{Set}}(X_i, A) \supseteq \mathrm{Hom}_{\mathbf{Set}}\left(\bigcup_{i \in \mathbb{N}} X_i, A\right)$$

The reason for this is we also need to ask the morphisms  $f_i \in \mathrm{Hom}_{\mathbf{Set}}(X_i, A)$  compatible which the morphism from functor  $X_{i \leq j} : X_i \rightarrow X_j$ , that is to say  $f_i = f_j \circ X_{i \leq j}$ .

$$\begin{array}{ccccc} \cdots & \xrightarrow{X_{i-1 \leq i}} & X_i & \xrightarrow{X_{i \leq i+1}} & X_{i+1} & \xrightarrow{X_{i+1 \leq i+2}} & \cdots \\ & & \downarrow f_i & \swarrow f_{i+1} & & & \\ & & A & & & & \end{array}$$

This motivates us to give the definition of Limit and Colimit:

**Definition 2.1.1** (Limit and Colimit of Set). Let  $D : J \rightarrow \mathbf{Set}$  be a diagram (functor) of sets indexed by a small category  $J$ . The *limit* of  $D$ , denoted  $\mathrm{Lim}_J D$ , is the subset of the product

$$\prod_{j \in J} D(j)$$

consisting of all families  $(x_j)_{j \in J}$  such that for every morphism  $f : i \rightarrow j$  in  $J$ , we have  $D(f)(x_i) = x_j$ .

The *colimit* of  $D$ , denoted  $\mathrm{Colim}_J D$ , is the quotient of the disjoint union

$$\bigsqcup_{j \in J} D(j)$$

by the equivalence relation generated by  $x \sim D(f)(x)$  for every morphism  $f : i \rightarrow j$  in  $J$  and every  $x \in D(i)$ .

We will omit  $J$  sometimes.

*Remark 2.1.2.* The intuition of limit is gluing functions, of colimit is gluing space. Image there is a covering of spaces  $\bigcup X_i \rightarrow X$ , to gluing space we start from  $\bigsqcup X_i$  then we identify the intersections  $X_i \cap X_j \hookrightarrow X_i, X_j$ ; To gluing function on  $C(X_i)$  we start with  $\prod C(X_i)$  then we impose compactibility condition when restrict to

$$C(X_i \cap X_j).$$

**Exercise 2.1.1.** If we view  $i \mapsto \text{Hom}_{\mathbf{Set}}(X_i, A)$  as functor  $h_A(X_{(-)}) : \mathbb{N}^{\text{op}} \rightarrow \mathbf{Set}$ , Then we have the relation between limit and colimit:

$$\forall A \in \mathbf{Set}, \lim_{i \in \mathbb{N}^{\text{op}}} \text{Hom}_{\mathbf{Set}}(X_i, A) \cong \text{Hom}_{\mathbf{Set}}(\text{Colim}_{i \in \mathbb{N}} X_i, A)$$

Show that this is hold in general for all category  $J$

$$\forall A \in \mathbf{Set}, \lim_{j \in J^{\text{op}}} \text{Hom}_{\mathbf{Set}}(X_j, A) \cong \text{Hom}_{\mathbf{Set}}(\text{Colim}_{j \in J} X_j, A)$$

$$\forall A \in \mathbf{Set}, \lim_{j \in J} \text{Hom}_{\mathbf{Set}}(A, X_j) \cong \text{Hom}_{\mathbf{Set}}(A, \lim_{j \in J} X_j)$$

**Exercise 2.1.2.** Consider the category of functor  $\text{Fun}(J, \mathbf{Set})$  where the morphism is natural transformation. For  $A \in \mathbf{Set}$  let  $c(A) \in \text{Fun}(J, \mathbf{Set})$  be the const functor. Show that

$$\text{Hom}_{\text{Fun}(J, \mathbf{Set})}(X_{(-)}, c(A)) \cong \lim_{j \in J^{\text{op}}} \text{Hom}_{\mathbf{Set}}(X_j, A)$$

$$\text{Hom}_{\text{Fun}(J, \mathbf{Set})}(c(A), X_{(-)}) \cong \lim_{j \in J} \text{Hom}_{\mathbf{Set}}(A, X_j)$$

**Example 2.1.3. • Product and Coproduct:**

Consider the set  $J$  viewed as a discrete category and no non-identity morphisms. A functor  $D : J \rightarrow \mathbf{Set}$  is simply a family of sets  $\{A_j\}_{j \in J}$ . The limit of  $D$  is the product set  $\prod_{j \in J} A_j$  and the colimit is the coproduct (disjoint union)  $\bigsqcup_{j \in J} A_j$ .

**• Equalizer and Coequalizer:**

Let  $J$  be the category with two objects  $\alpha, \beta$  and two parallel morphisms  $f, g : \alpha \rightarrow \beta$ . A functor  $D : J \rightarrow \mathbf{Set}$  consists of sets  $A, B$  and functions  $f, g : A \rightarrow B$ . The limit (equalizer) is:

$$\lim D = \text{Eq}(f, g) := \{a \in A \mid f(a) = g(a)\}.$$

And the colimit (coequalizer) is the quotient set:

$$\text{Colim } D = \text{Coeq}(f, g) := B / \sim$$

where  $b \sim b'$  if there exists  $a \in A$  such that  $f(a) = b$  and  $g(a) = b'$ .

- **Pullback and Pushout:**

Let  $J$  be the diagram  $\alpha \xrightarrow{f} \gamma \xleftarrow{g} \beta$ . A functor  $D : J \rightarrow \mathbf{Set}$  assigns sets  $X, Y, Z$  and functions  $f : X \rightarrow Z, g : Y \rightarrow Z$ . The limit (pullback) is:

$$\mathrm{Lim} D = X \times_Z Y := \{(x, y) \in X \times Y \mid f(x) = g(y)\}.$$

And the colimit (pushout) is:

$$\mathrm{Colim} D = X \sqcup_Z Y := (X \sqcup Y) / \sim$$

where  $\sim$  is the equivalence relation generated by  $f(z) \sim g(z)$  for all  $z \in Z$ .

*Remark 2.1.4.* Actually these are all essential limit and colimit: limit can be seen as equalizer of product and colimit can be seen as coequalizer of coproduct:

Limit:

- Consider the product  $\prod_{j \in J} D(j)$ .
- For each morphism  $f : i \rightarrow j$  in  $J$ , define two morphisms:

$$\alpha, \beta : \prod_{j \in J} D(j) \rightarrow \prod_{f: i \rightarrow j} D(j)$$

where  $\alpha$  sends  $(x_j)_{j \in J}$  to  $(D(f)(x_i))_{f: i \rightarrow j}$  and  $\beta$  sends  $(x_j)_{j \in J}$  to  $(x_j)_{f: i \rightarrow j}$ .

- The limit  $\mathrm{Lim} D$  is the equalizer of  $\alpha$  and  $\beta$ :

$$\mathrm{Lim} D = \mathrm{Eq}(\alpha, \beta).$$

Colimit:

- Consider the coproduct  $\coprod_{j \in J} D(j)$ .
- For each morphism  $f : i \rightarrow j$  in  $J$ , define two morphisms:

$$\alpha', \beta' : \bigsqcup_{f: i \rightarrow j} D(i) \rightarrow \bigsqcup_{j \in J} D(j)$$

where  $\alpha'$  sends  $x \in D(i)$  (in the  $f : i \rightarrow j$  summand) to  $D(f)(x) \in D(j)$ , and  $\beta'$  sends  $x \in D(i)$  to  $x$  viewed in  $D(i)$ .



- The colimit  $\text{Colim } D$  is the coequalizer of  $\alpha'$  and  $\beta'$ :

$$\text{Colim } D = \text{Coeq}(\alpha', \beta').$$

*Remark 2.1.5* (Homotopy (Co)limit).

To definite limit and colimit for general category, we can make use of morphism:

**Definition 2.1.6** (Limit and Colimit). Let  $\mathcal{C}$  be a category,  $J$  a small category, and  $D : J \rightarrow \mathcal{C}$  a functor (called a diagram in  $\mathcal{C}$ ). A *limit* of  $D$  is an object  $\text{Lim}_J D$  of  $\mathcal{C}$  such that

$$\forall A \in \mathcal{C}, \text{Lim}_{j \in J} \text{Hom}_{\mathcal{C}}(A, D(j)) \cong \text{Hom}_{\mathcal{C}}(A, \text{Lim}_J D)$$

A *colimit* of  $D$  is an object  $\text{Colim}_J D$  of  $\mathcal{C}$  such that

$$\forall A \in \mathcal{C}, \text{Lim}_{j \in J^{\text{op}}} \text{Hom}_{\mathcal{C}}(D(j), A) \cong \text{Hom}_{\mathcal{C}}(\text{Colim}_J D, A)$$

**Example 2.1.7** (Poset and Lattice). In a category come from poset  $P$  we have the  $\text{Lim } p$  is the meet  $\bigwedge p$ , and the colimit  $\text{Colim } p$  is the joint  $\bigvee p$

Thanks to Yoneda lemma, to understand stand (co)limit of general category, we only need to understand its presheaf, and thus limit of **Set**.

Then we introduce an important kind of colimit, *filtered colimit*.

**Definition 2.1.8** (Filtered Category). A small category  $J$  is called *filtered* if:

1. It is non-empty.
2. For every pair of objects  $j_1, j_2 \in \text{Ob}(J)$ , there exists an object  $k \in \text{Ob}(J)$  and morphisms  $j_1 \rightarrow k$  and  $j_2 \rightarrow k$ .
3. For every pair of parallel morphisms  $f, g : j \rightarrow j'$  in  $J$ , there exists an object  $k$  and a morphism  $h : j' \rightarrow k$  such that  $h \circ f = h \circ g$ .

**Example 2.1.9.** For the category of open set  $\text{Op}(X)$ , both  $\text{Op}(X)$  and  $\text{Op}(X)^{\text{op}}$  are filtered

**Exercise 2.1.3.** Let  $I$  be any finite category, i.e with finite objects and morphisms, then for any functor  $F : I \rightarrow J$ , Show that  $J$  is filtered iff there exists a object  $j_F \in J$  and a natural transformation  $F \rightarrow c(j_F)$ .

Intuitively, a filtered category allows us to "coherently glue" data indexed by  $J$ . The colimit definite from  $D : J \rightarrow \mathcal{C}$  is called *filtered colimit*. Filtered colimit can be defined by more explicit quotient instead of that of general colimit.

**Exercise 2.1.4.** Show that for  $J$  filtered  $D : J \rightarrow \mathbf{Set}$ , the filtered colimit  $\text{Colim}_J D = \left( \coprod_{j \in J} F(j) \right) / \sim$ , where

$(x, j) \sim (y, k)$  if there exist morphisms  $f : j \rightarrow l, g : k \rightarrow l$  in  $J$  such that  $D(f)(x) = D(g)(y)$ .

**Exercise 2.1.5.** 1. Let  $I, J$  be any small categories, consider  $D : I \times J \rightarrow \mathcal{C}$ . Then show that:  $\text{Lim}_I \text{Lim}_J D \cong \text{Lim}_{I \times J} D \cong \text{Lim}_J \text{Lim}_I D$ , same for the colimit. (You just need prove it for  $\mathbf{Set}$ .)

2. Let  $I$  be finite, and  $J$  be filtered category, consider  $D : I \times J \rightarrow \mathbf{Set}$ . Then show that  $\text{Lim}_I \text{Colim}_J D \cong \text{Colim}_J \text{Lim}_I D$ . That is to say filtered colimit preserve finite limit.

Limit and Colimit are not necessarily always existing for all  $D : J \rightarrow \mathcal{C}$ , but just like for a space  $X$  we can define its completion  $\hat{X}$  to make limit exist, we can define a (co)completion of a category  $\mathcal{C}$  to make (co)limit exists. An easy observation is limit and colimit are interchanged in  $\mathcal{C}$  and  $\mathcal{C}^{\text{op}}$ , so let us be focus on the case for colimit. Notice that limit and colimit are admitted for  $\mathbf{Set}$ , then so do presheaf category  $\mathbf{PSh}(\mathcal{C})$

We first introduce the category of elements for a presheaf, which will be the diagram for a colim to assembly spaces.

**Definition 2.1.10** (Category of Elements). Let  $\mathcal{C}$  be a category and let  $\mathcal{F} : \mathcal{C}^{\text{op}} \rightarrow \mathbf{Set}$  be a presheaf. The *category of elements* of  $\mathcal{F}$ , denoted  $\int_{\mathcal{C}} \mathcal{F}$ , is defined as follows:

- **Objects:** Pairs  $(C, x)$  where  $C$  is an object of  $\mathcal{C}$  and  $x \in \mathcal{F}(C)$ .
- **Morphisms:** A morphism  $(C, x) \rightarrow (D, y)$  is a morphism  $f : C \rightarrow D$  in  $\mathcal{C}$  such that

$$\mathcal{F}(f)(y) = x.$$

(Note: since  $\mathcal{F}$  is contravariant, the direction of  $f$  is  $C \rightarrow D$ , but the induced map goes  $\mathcal{F}(D) \rightarrow \mathcal{F}(C)$ .)

- **Composition and identities:** Inherited from the category  $\mathcal{C}$ .

**Example 2.1.11** (Category of Elements of Simplicial Set). Let  $X : \Delta^{\text{op}} \rightarrow \mathbf{Set}$  be a simplicial set, Then the category of elements  $\int_{\Delta} X$  will be

- **Objects:**  $X([0])$  collection of points;  $X([1])$  collection of segments;  $X([3])$  collection of triangles, ...
- **Morphisms:** A morphism  $x \in X([n]) \rightarrow y \in X([n+1])$  is a morphism  $\delta_i : [n] \rightarrow [n+1]$  such that  $x$  is the  $i$ -th boundary of  $y$ , etc.

Intuitively we should think  $\int_{\mathcal{C}} \mathcal{F}$  give us the blueprint to reassembly. For any functor  $R : \mathcal{C} \rightarrow \mathcal{D}$ , we should think

$$\int_{C \in \mathcal{C}} \mathcal{F}(C) \times R(C) \left( \text{or } \int_{\mathcal{C}} \mathcal{F} \times R \right) := \text{Colim}_{(C,x) \in \int_{\mathcal{C}} \mathcal{F}} R(C) \in \mathcal{D}$$

Is the assembly of  $\mathcal{F}$  inside  $\mathcal{D}$ .

**Exercise 2.1.6.** Take the Yoneda embedding  $\mathfrak{y} : \mathcal{C} \rightarrow \mathbf{PSh}(\mathcal{C})$ , we have the  $\int_{\mathcal{C}} \mathcal{F} \times \mathfrak{y} \in \mathbf{PSh}(\mathcal{C})$ .

1. Show that for  $B \in \mathcal{C}$ , we have a map

$$\mathcal{F}(B) \rightarrow \int_{C \in \mathcal{C}} \mathcal{F}(C) \times \mathfrak{y}(C)(B) = \text{Colim}_{(C,x) \in \int_{\mathcal{C}} \mathcal{F}} \text{Hom}_{\mathcal{C}}(B, C)$$

defined by

$$b \in \mathcal{F}(B) \mapsto ((B, b), \text{id}_B) \in \text{Colim}_{(C,x) \in \int_{\mathcal{C}} \mathcal{F}} \text{Hom}_{\mathcal{C}}(B, C)$$

which is a bijection. (This is analogue to  $f(b) = \int_{c \in \mathbb{R}} f(c) \delta(c - b)$ )

2. This extended to a natural transformation  $\mathcal{F} \rightarrow \int_{\mathcal{C}} \mathcal{F} \times \mathfrak{y}$  which is an isomorphism.
3. Let  $\mathcal{D}$  be a cocomplete category (i.e. admitted all colimit), then there is a map between

$$\text{Fun}(\mathcal{C}, \mathcal{D}) \rightarrow \text{Fun}^{\text{cocont}}(\mathbf{PSh}(\mathcal{C}), \mathcal{D})$$

where  $\text{Fun}^{\text{cocont}}$  means cocontinuous functor i.e. preserve colimit, the map is defined by

$$R \in \text{Fun}(\mathcal{C}, \mathcal{D}) \mapsto \left( \mathcal{F} \in \text{PSh}(\mathcal{C}) \mapsto \int_{\mathcal{C}} \mathcal{F} \times R \in \mathcal{D} \right)$$

Show that this is a bijection. And in other words,  $\mathfrak{y} : \mathcal{C} \rightarrow \text{PSh}(\mathcal{C})$  is the *free cocompletion* of  $\mathcal{C}$ .

Notice that even if  $\mathcal{C}$  is cocompletion, which doesn't mean  $\mathcal{C} \cong \text{PSh}(\mathcal{C})$ , because  $\mathfrak{y} : \mathcal{C} \rightarrow \text{PSh}(\mathcal{C})$  is not cocontinuous, i.e.  $\mathfrak{y}(\text{Colim } C) \neq \text{Colim } \mathfrak{y}(C)$  in general. This can also be seen from  $\mathcal{F} \in \text{PSh}$  is not continuous in general:

$$\mathcal{F}(\text{Colim } C) = \text{Hom}_{\text{PSh}(\mathcal{C})}(\mathfrak{y}(\text{Colim } C), \mathcal{F}) \neq \text{Hom}_{\text{PSh}(\mathcal{C})}(\text{Colim } \mathfrak{y}(C), \mathcal{F}) = \text{Lim } \mathcal{F}(C)$$

So if we impose some continuous condition, we can get subcategory of continuous presheaf, which are called sheaf, behaves more close to the test category  $\mathcal{C}$ .

**Exercise 2.1.7.** Show that for  $B \in \mathcal{C}$ ,  $\mathfrak{y}(B)$  are continuous, i.e.  $\mathfrak{y}(B)(\text{Colim}_{j \in J} C_j) \cong \text{Lim}_{j \in J} \mathfrak{y}(B)(C_j)$ . In others words  $\mathfrak{y}(B)$  is a sheaf.

## 2.2 Sheaf and Coverage

If we ask to presheaf be continuous for all colimit we will get nothing more than representable ones, instead we can restrict which kinds of colimit that it preserve.

**Example 2.2.1** (Ind-Object). Let  $\mathcal{C}$  have all finite colimit (i.e for  $C : J \rightarrow \mathcal{C}$  with  $J$  finite), Then the ind-object is the presheaf that preserve finite colimit, i.e.  $\mathcal{F}(\text{Colim}_J C) \cong \text{Lim}_J \mathcal{F}(C)$  for  $J$  finite, this is also called *left exact*. Let  $\text{ind}(\mathcal{C}) \subset \text{PSh}(\mathcal{C})$  be the subcategory of ind-object.

**Exercise 2.2.1.** 1. For any  $J$  filtered and  $D : J \rightarrow \mathcal{C}$ , show that  $\text{Colim}_J \mathfrak{y}(D) \in \text{PSh}(\mathcal{C})$  is actually an ind-object.

2. \* For any  $\mathcal{F} \in \text{ind}(\mathcal{C})$ , show that the category of elements  $\int_{\mathcal{C}} \mathcal{F}$  is filtered. This shows that  $\mathcal{F} \cong \int_{\mathcal{C}} \mathcal{F} \mathfrak{y}$  is a filter colimit, and  $\text{ind}(\mathcal{C})$  is generated by representable presheaf with filtered colimit.

In geometry setting, the colimit we want to preserve are come from coverage.

# Bibliography

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