Lists

(ordered set of data objects)

Sequential representation

- successive nodes of the data object are stored a fixed distance apart → array 구조
- (physical) order of elements is the same as in ordered list
- adequate for functions such as accessing an arbitrary node in a table
- operations such as insertion and deletion of arbitrary elements from ordered lists become expensive

Linked representation

- successive items of a list may be placed anywhere in memory (효율성)
- (physical) order of elements need not be the same as order in list
- each data item is associated with a pointer (link) to the next item
- We will study <u>singly linked lists (SLL)</u>

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List ADT (description of functionality of list data structure)

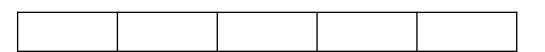
```
template < class T>
class List {
public:
   List();
   ~List();
   void insert(int loc, T d);
   void remove(int loc);
   int getSize();
   T & getData(int loc);
private:
   // not yet !!
};
```

```
int main() {
   List<int> mylist;
   mylist.insert(1, 10);
   mylist.insert(2, 20);
   mylist.insert(3, 15);
   mylist.remove(2);
   mylist.insert(1, 35);
   mylist.insert(3, mylist.getData(2));
   cout << mylist.getSize() << endl;</pre>
   return 0;
```

Implementation by Array

```
template < class T>
class List {
public:
   List() { size = 0; }
   //~List();
   void insert(int loc, T d);
   void remove(int loc);
   int getSize();
   T & getData(int loc);
private:
   ⊤ data[5];
   int size;
};
```

insert a new data object?
remove an existing data object?



Implementation by Links

```
template < class T>
class List {
private:
   class listNode {
   public:
     T data;
     listNode *next;
     listNode(T newItem);
public:
   List(): size(0), head(NULL) { }
   ~List();
   listNode * Find(listNode *nptr, int k);
   void insert(int loc, T d);
   void remove(int loc);
   int getSize() const;
   T & getData(int loc);
private:
   listNode *head;
   int size;
};
```

- head
 - Point to an object of listNode in List.
- **-~**List();
- ✓ Need to delete dynamically allocated objects of listNode.
- Find(nptr, k);
- ✓ Take k steps forward in the list from the object pointed by 'nptr'.
- ✓ Array implementation 의 index 역할.

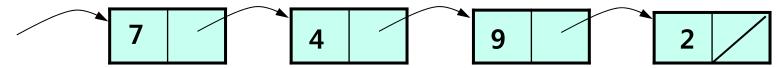
Implementation by Links (계속)

```
template < class T>
class List {
private:
   class listNode {
   public:
     T data;
     listNode *next;
     listNode(T newItem);
public:
   List() { size = 0; head = NULL; }
   ~List();
   listNode * Find(listNode *nptr, int k);
   void insert(int loc, T d);
   void remove(int loc);
   int getSize() const;
   T & getData(int loc);
private:
   listNode *head;
   int size;
};
```

- class listNode { ... }
 - Nested class of List class.
 - Only List functions can create objects of the private class listNode.
 - Only List functions can access listNode objects.
 - Implementation of constructor:

Insert a new node at the front

```
List<int> mylist;
.....
listNode *h = mylist.getHead();
mylist.insertAtFront(h, 10);
```



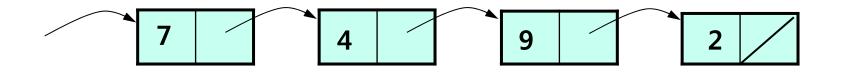
```
// Insert a new node with data d at the front of the list node
// pointed by curr. In addition, curr should be updated to point
// the new node.
void List<T>::insertAtFront(listNode * curr, T d) {
```

}

Time complexity:

For array case ?

Print data in list in reverse order



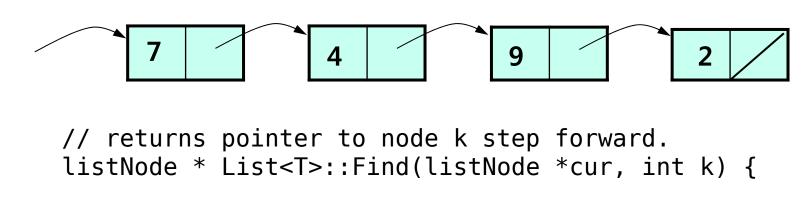
```
// 2 9 4 7 should be printed for the input of above list.
void List<T>::printReverse(listNode *curr) {
```

}

Time complexity:

For array case ?

Find a pointer to node k steps forward from *current



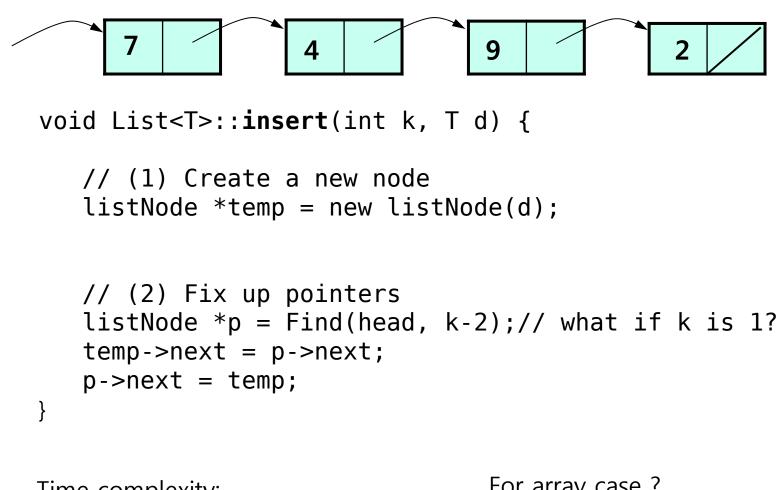
}

Time complexity:

For array case?

Insert a new node in k-th position

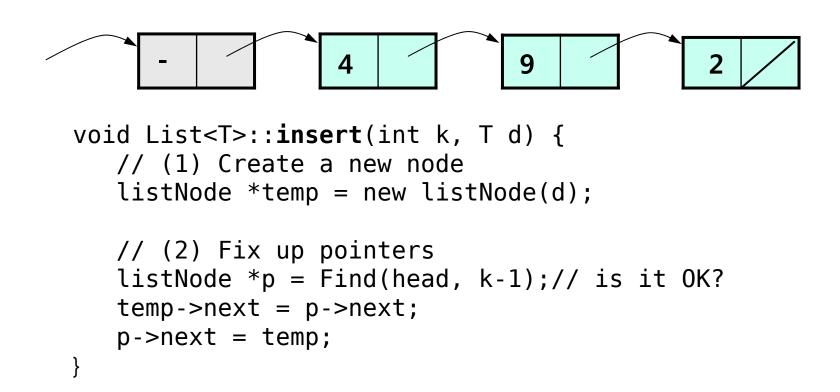
List<int> mylist; mylist.insert(3, 100);



Time complexity:

For array case?

Insert a new node in a list with a sentinel node



Time complexity:

Remove a node in fixed position (given by a node pointer)

```
void List<T>::remove(listNode * curr) {
```

}

Time complexity:

Remove a node in fixed position (given by a node pointer)

```
- 4 6 9 2

// Constant time (trick!!)
void List<T>::remove(listNode * curr) { curr
```

Run times for List functions

	Singly linked List	Array
Insert/Remove at front	O(1)	O(1)
Insert at given location	O(1): inserting a node after the given location	O(n) shift : inserting a node at the given index
Remove at given location	O(1) trick	O(n) shift
Insert at arbitrary location	O(n) find	O(n) shift
Remove at arbitrary location	O(n) find	O(n) shift