Stack

(LIFO ordering structure)

Stack ADT

```
template <class T>
class Stack {
public:
  Stack();
  bool IsEmpty() const;
  // T& Top() const;
  void Push(const T& data);
  T Pop();
private:
   Not yet!!
```

```
int main() {
   List<int> mystack;
   mystack.Push(9);
   mystack.Push(3);
   mystack.Push(2);
   mystack.Pop();
   mystack.Push(4);
```

Linked list based implementation

```
template <class T>
class Stack {
public:
  Stack();
  bool IsEmpty() const;
  void Push(const T& d);
  T Pop();
private:
   struct stackNode {
       T data;
       stackNode *next;
       stackNode();
   stackNode *top;
   int size;
};
```

```
template < class T>
     void Stack<T>::Push(const T & d)
        stackNode *newNode =
                  new stackNode(d);
4
         newNode->next = top;
         top = newNode;
         size++;
3
     // Assume stack is not empty
     template < class T>
     T Stack<T>::Pop() {
9
```

Array based implementation

size = 3

```
template <class T>
class Stack {
public:
  Stack();
  bool IsEmpty() const;
  void Push(const T& d);
  T Pop();
private:
   int capacity;
   int size;
   T *stack;
};
```

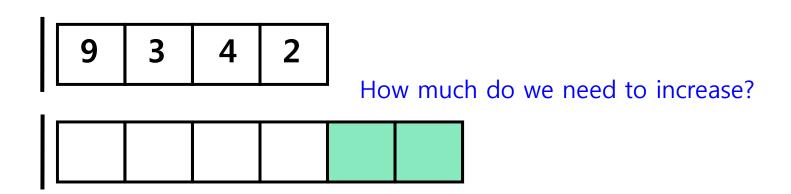
```
template <class T>
void Stack<T>Push(const T &d) {
    if (size == capacity) {
        // resize stack array
    }
    stack[size] = d;
    size++;
}
```

```
3
4
1
3
9
```

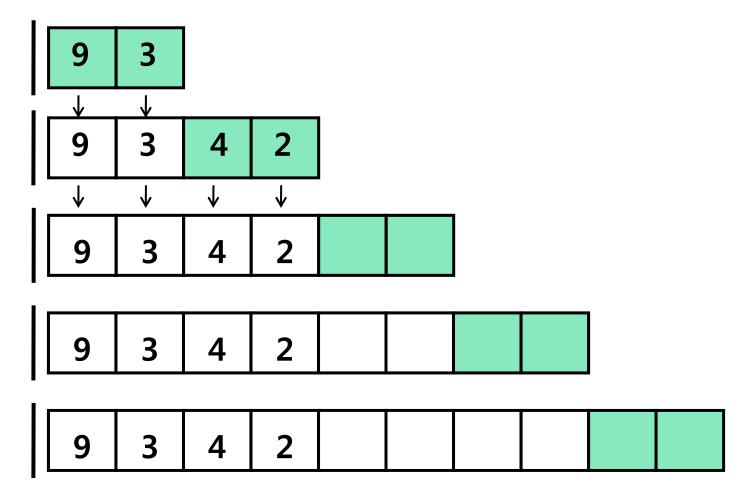
```
template <class T>
Stack<T>::Stack() {
   capacity = 4;
   size = 0;
   stack = new T[capacity];
}
```

Array resizing when inserting (pushing), but the array is full

Steps.

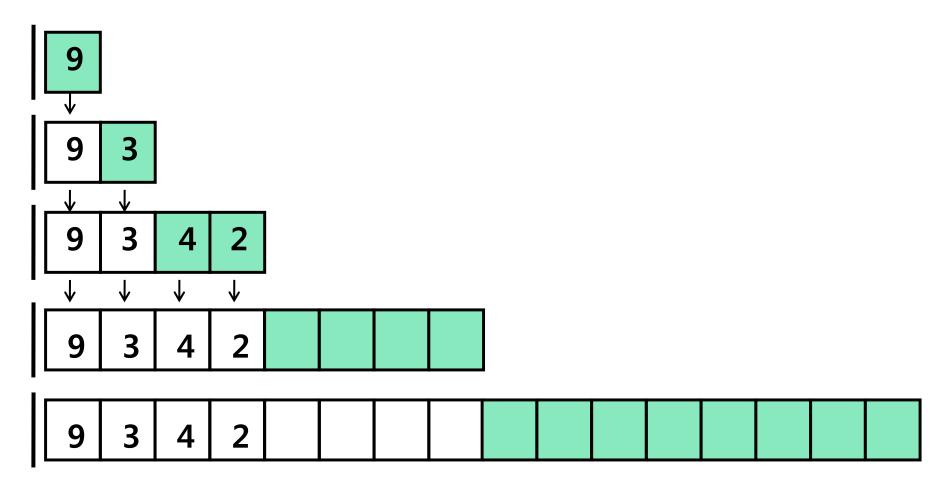


Strategy 1: Increase array size by a constant amount



How much is the cost by this strategy on a sequence of n pushes?

Strategy 2: Double the array size



How much is the cost by this strategy on a sequence of n pushes?

Run time summary

- Linked list based implementation
 - O() per push, pop
- Array based implementation
 - O() per pop
 - O() per push if capacity exists.
 - Cost over n pushes is O() for an average of O() per push.
- Which implementation we choose?

Queue

(FOFO ordering structure)

Queue ADT

```
template <class T>
class Queue
public:
  Queue ();
  Boolean IsEmpty() const;
  // T& Front() const;
  // T& Rear() const;
  void Push(const T& item);
  T Pop();
Private:
  // not yet
```

```
pop()
dequeue()

exit
front

push()
enqueue()
```

Linked list based implementation

```
template <class T>
class Queue {
public:
  Queue();
  bool IsEmpty() const;
  void Push(const T& d);
  T Pop();
private:
```

```
4 9 2
```

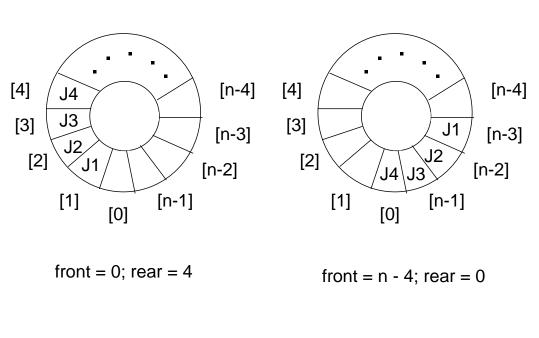
Which pointer is "entry" ("rear") and which is "exit" ("front")?

Running time of Push?

Running time of Pop?

Array based implementation

```
template <class T>
class Queue {
public:
  Queue();
  bool IsEmpty() const;
  void Push(const T& d);
  T Pop();
private:
   T *queue;
   int capacity;
   int front
   int rear;
};
```



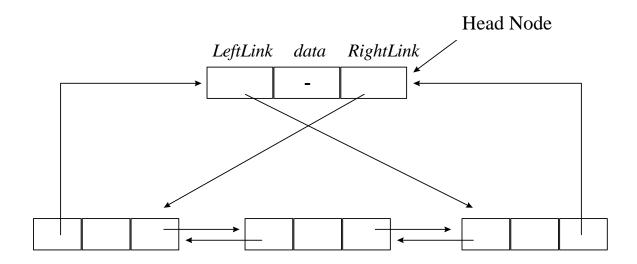
Modeling as a circular list

DLL

(Doubly Linked Lists)

- Problems with singly linked lists
 - can move only in the direction of the links
 - deletion requires knowing the preceding node
- Node
 - at least three fields :
 data, llink (left link), rlink (right link)

DLL with sentinel node



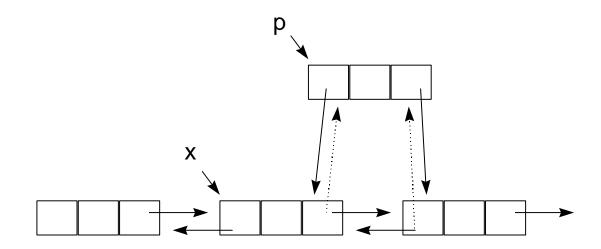
- Head node convenient for algorithms data field usually contains no information
- Essential virtue
 one can go back and forth with equal ease
 p==p→llink→rlink==p→rlink→llink

Implementation

```
class DblList;
class DblListNode {
// Dblist functions are allowed to access
friend class DblList;
private:
  int data;
  DblListNode *left, *right;
class DblList {
public:
 // List manipulation operations
private:
  DblListNode *first; // points to head node
```

Insert a node

```
void DblList::Insert(DblListNode *p, DblListNode *x)
// insert node p to the right of node x
{
    p->left = x; p->right = x->right;
    x->right->left = p; x->right = p;
}
```



Delete a node

```
void DblList::Delete(DblListNode *x)
{
  if ( x == first ) throw "Deletion of headnode not permitted";
  else {
     x->left->right = x->right;
     x->right->left = x->left;
     delete x;
  }
}
```