

DATA STRUCTURES AND ALGORITHMS

Arrays

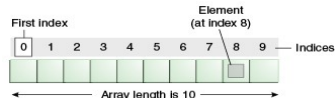
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Content

- Concept of Arrays
- Array Representation
- Operations performed on arrays
- Limitation of arrays
- Application of linear Arrays

Arrays

- A linear array is a list of finite number of homogeneous data elements such that
 - The elements of the array are referenced by an index set consisting of n consecutive integer numbers
 - The elements of the array are stored in consecutive memory locations



Array Representation in Memory

- Memory of a computer system is simply a sequence of addressed location.

Computer	Content	Name	Type
00000000	00		
00000001	00		
00000002	00		
00000003	FF		
00000004	FF		
00000005	FF		
00000006	FF		
00000007	FF		
00000008	FF		
00000009	FF		
0000000A	FF		
0000000B	FF		
0000000C	FF		
0000000D	FF		
0000000E	00		
0000000F	00		
00000010	00		
00000011	00		
00000012	00		
00000013	00		
00000014	00		
00000015	00		
00000016	00		
00000017	00		
00000018	00		
00000019	00		
0000001A	00		
0000001B	00		
0000001C	00		
0000001D	00		
0000001E	00		
0000001F	00		

- As arrays are stored in consecutive memory locations, the system need not to keep track of the address of every element of that array, but needs to keep the address of first element only (base address)

Array Representation in Memory

Address	Slot
1197	
1198	15
1199	
1200	14
1201	
1202	23
1203	
1204	11
1205	
1206	42
1207	
1208	
1209	

0	1	2	3	4
15	14	23	11	42

array

Loc (array[k])= base(array) + w * k
w is # of bytes per element
K is index Number

Loc(array[3])= 1197+2*3=1203

Operations performed on arrays

- Traversal Operation
- Searching Operation
 - Linear Search
 - Binary Search
- Insertion Operation
- Deletion Operation
- Sorting Operation
 - Selection Sort
 - Bubble Sort
 - Insertion Sort
 - Merge Sort
 - Quick Sort etc.

Traversal Operation

TraverseLinearArray(a, n)

Here **a** is a linear array of size **n**. This algorithm traverses the array and applies certain operation to each element of the array.

1. Set $i=0$ //Initialize counter
2. Repeat steps 3 and 4 while $i \leq (n-1)$
3. Apply process ($a[i]$) //visit element
4. Set $i=i+1$ //increment counter
5. Endwhile
6. Exit

0	1	2	3	4	5
15	14	23	11	42	03

Searching

Searching Operation: Linear Search

LinearSearch(a, n, item, loc)

Here **a** is a linear array of size **n**. This algorithm finds the location of the **item** in linear array **a**. If search end in success it sets **loc** to the index of the element; otherwise it sets **loc** to -1.

1. Set $i=0$ //Initialize counter
2. Repeat steps 3 and 4 while $i \leq (n-1)$
3. If ($a[i]=item$) then
Set $loc=i$ //item found at location **i**
Exit
Endif
4. Set $i=i+1$ //increment counter
5. Endwhile
6. Set $loc=-1$ //item not found
7. Exit

0	1	2	3	4	5	6
3	10	15	20	35	40	60

Searching Operation: Binary Search

BinarySearch(a, n, item, loc)

Here **a** is a linear array of size **n**. This algorithm finds the location of the **item** in sorted (ascending) linear array **a**. If search end in success it sets **loc** to the index of the element; otherwise it sets **loc** to -1. Here variables **beg** and **end** are used to keep track of the first and last element of the array and variable **mid** is used as index of the middle element of the array under consideration.

1. Set $beg=0$
2. Set $end=n-1$
3. Set $mid=(beg+end)/2$
4. Repeat steps 5 and 6 while ($beg \leq end$) && ($a[mid] \neq item$)
5. If ($item < a[mid]$) then
Set $end=mid-1$
else
Set $beg=mid+1$
Endif
6. Set $mid=(beg+end)/2$ // shift mid
7. Endwhile
8. If ($beg > end$) then
Set $loc=-1$ //item not found
Else
Set $loc=mid$ //item found at mid
Endif
9. Exit

Searching Operation: Binary Search

Solution:

Given Array **a**

a[0]	a[1]	a[2]	a[3]	a[4]	a[5]	a[6]
3	10	15	20	35	40	60

To start with, we take $beg = 0$, $end = 6$, and compute location of middle element as

$$mid = (beg + end) / 2 = (0 + 6) / 2 = 3$$

Since $a[mid]$, i.e., $a[3] \neq 15$, and $beg \leq end$. We start next iteration.

As $a[mid] = 20 > 15$, therefore, we take $end = mid - 1 = 3 - 1 = 2$, where as beg remains unchanged. Thus

$$mid = (beg + end) / 2 = (0 + 2) / 2 = 1$$

Since $a[mid]$, i.e., $a[1] \neq 15$, and $beg \leq end$. We start next iteration.

As $a[mid] = 10 < 15$, therefore, we take $beg = mid + 1 = 1 + 1 = 2$, where as end remains unchanged.

Since $beg = end$, again compute location of the middle element as

$$mid = (beg + end) / 2 = (2 + 2) / 2 = 2$$

Since $a[mid]$, i.e., $a[2] = 15$, the search terminates on success.

The element is found at index 2.

0	1	2	3	4	5	6
3	10	15	20	35	40	60

0	1	2	3	4	5	6
3	10	15	20	35	40	60

0	1	2	3	4	5	6
3	10	15	20	35	40	60

Class Task

Write a function to find the location of largest element in an array.

Signature of function must as follow
int Largest(int arr[], int n)

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Bubble Sort

```
- void bubbleSort(int arr[], int size)
- {
-   for (int i = size-1; i > 0; i--)
-   {
-     for (int k = 0; k < i; k++)
-     {
-       if (arr[k] > arr[k+1])
-       {
-         int temp=arr[k];
-         arr[k]=arr[k+1];
-         arr[k+1]=temp;
-       } //if
-     } //for with k
-   } //for with i
- } //function
-
```

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Thank You