DATA STRUCTURES AND ALGORITHMS

Circular Queue Data Structure

Ву

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Content

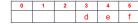
- · Limitation of Linear Queue
- · Introduction to Circular Queue
- · Properties of a Circular Queue
- · Operations of Circular Queue
- · Applications of Circular Queue

Limitation of Linear Queue

The only limitation of a linear queue is that- If the last position of the queue is occupied, it is not possible to enqueue anymore elements even though some positions are vacant.

Operation
Enqueue(g)

Rear	front
5	3



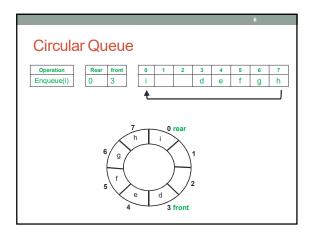
Error: Queue is overflow

- Solution: Circular Queue
- This limitation can be overcome by moving the rear back to index '0', if front is >0

Circular Queue

- Circular Queue is the advanced form of Queue data structure.
- Like a linear Queue, elements are added from an end i.e. rear, and removed from another end that is known as the front.
- It also ensures the first-in-first-out (FIFO) or last-in-last-out (LILO) order of insertion and deletion.
- However, Unlike linear queue, in circular queue rear is reset to index '0', if there are some vacant slots at the beginning.

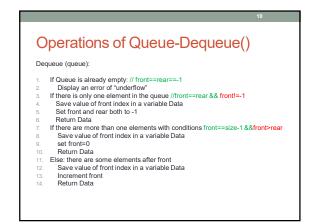
Linear Queue vs. Circular Queue Linear Queue Operation Enqueue(g) Error: Queue is overflow Circular Queue Operation Enqueue(g) Operation Enqueue(g)

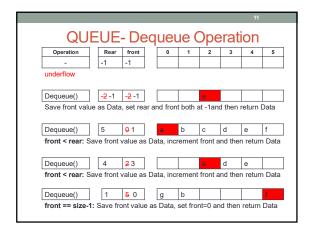


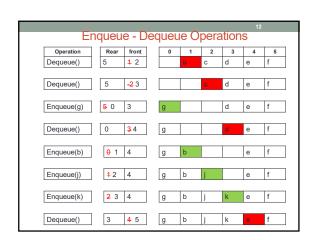
Operations of Queue The common operations of queue are as follow: enqueue() dequeue() isEmpty() isFull() frontValue() rearValue()

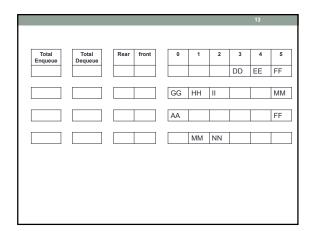
Operations of Queue-Enqueue(item) Enqueue (queue, item) 1. If queue is already full: //front=0 & rear=size-1 or rear+1=front 2. Display an error of "overflow" 3. If queue is empty and this is the first item to be inserted in that queue //rear==front==-1 4. Increment rear and front both 5. Insert item at rear index 6. If queue is not empty and there are some vacant slots in the begining: // front > 0 and rear==size-1 7. Set rear=0 8. Insert item at rear index 9. Else: there are some slots after rear 10. Increment rear 11. Insert item at rear index

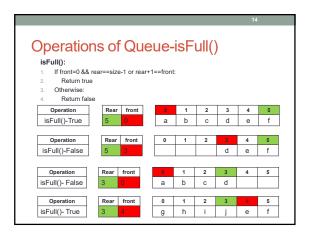
QUEUE- Enqueue Operation Operation Rear front -1 -1 -1 Enqueue(a) Increment rear and front both and then insert item at rear Enqueue(g) 5 0 a b c d e f Overflow because front is at 0 and rear is at size-1 Enqueue(e) 3 4 2 c d e Increment rear and then Insert item at rear index Enqueue(g) 5 0 2 g c d e f								9
Enqueue(a)	QUE	EUE- End	luei	ue (Эре	erat	tion	1
Increment rear and front both and then insert item at rear Enqueue(g) 5 0 a b c d e f Overflow because front is at 0 and rear is at size-1 Enqueue(e) 3 4 2 c d e Increment rear and then Insert item at rear index	Operation -		0	1	2	3	4	5
Overflow because front is at 0 and rear is at size-1 Enqueue(e) 3 4 2 c d e Increment rear and then Insert item at rear index	,		_	sert ite	em at r	ear		
Enqueue(e) 3 4 2 c d e Increment rear and then Insert item at rear index	1 (0)		-	-	_	d	е	f
Increment rear and then Insert item at rear index	Overnow beca	use noncis at o and		at SIZ	e- i			
	Enqueue(e)	3 4 2			С	d	е	
Enqueue(g) 5 0 2 g c d e f	Increment rear	and then Insert ite	m at re	ar inde	Х			
	Enqueue(g)	5 0 2	g		С	d	е	f

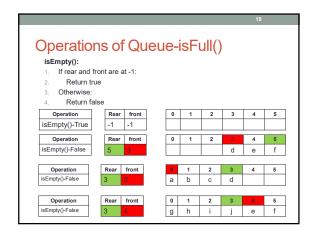


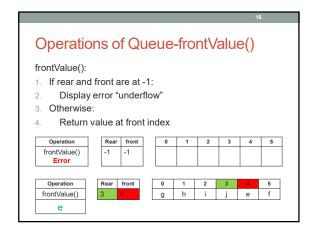


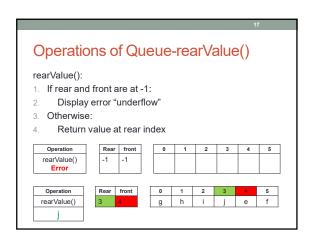


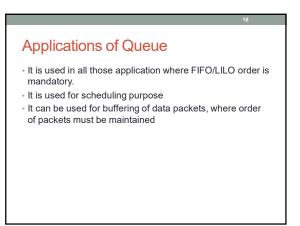












Thank You