### o-Minimality and its Variations

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### o-Minimality

#### Assumptions

- $\mathcal{M} = (M, <, ...)$
- < is dense, linear, without endpoints
- definability with parameters

#### Definition

The structure  $\mathcal{M}$  is called *o-minimal* if every definable subset of M is a finite union of singletons and open intervals with endpoints in  $M_{\infty} := M \cup \{-\infty, +\infty\}.$ 

A theory T is called o-minimal if every model  $\mathcal{M}$  of T is o-minimal.



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### o-Minimality

The class of o-minimal structures is,

- closed under reducts (if < still remains in the language)
- closed under expansions by constants

#### Some o-minimal structures

- $\bullet$  ( $\mathbb{Q},<$ )
- $\bullet$  ( $\mathbb{Q}, <, +$ )
- $\mathcal{R} = (\mathbb{R}, <, +, -, \cdot, 0, 1)$

#### $(\mathbb{Q}, <, +, \cdot, 0, 1)$ is not o-minimal

The infinite discrete set of perfect squares is definable.



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### Monotonicity and Finiteness

#### Monotonicity Theorem [Pillay & Steinhorn, 1986]

Let  $\mathcal{M}$  be an o-minimal structure and  $f:(a,b) \to M$  be a definable function with domain (a,b) (possibly  $a=-\infty$  or  $b=+\infty$ ). Then, there are points  $a=a_0 < a_1 < \cdots < a_{k+1}$  s.t. for each  $j=0,\ldots,k,$   $f|_{(a_j,a_{j+1})}$  is either,

- constant, or
- a strictly monotonic and continuous bijection to an interval.

#### Finiteness Lemma [Knight et al., 1986]

Let  $A \subseteq M^2$  be definable and suppose that for each  $x \in M$  the fiber  $A_x := \{y \in M | (x,y) \in A\}$  is finite. Then there is  $N < \omega$  s.t.  $|A_x| \leq N$  for all  $x \in M$ .





#### Corollary

Let  $f:(a,b)\to M$  be definale and continuous. Then f takes a maximum and minimum value on [a,b].

#### Exchange Principle [Pillay & Steinhorn, 1986]

Let  $\mathcal{M}$  be o-minimal. Let  $b, c, a_1, \ldots, a_n \in \mathcal{M}$ . If b is definable over  $c, a_1, \ldots, a_n$ , and b is not definable over  $a_1, \ldots, a_n$ , then c is definable over  $b, a_1, \ldots, a_n$ .

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#### Motivation

#### Question

What happen with the definable subsets of  $M^n$ ?

#### Notation

Given definable  $X \subseteq M^n$ 

- $C(X) := \{f : X \to M | f \text{ is definable and continuous} \}$
- $C_{\infty}(X) := C(X) \cup \{-\infty, +\infty\}$
- for  $f, g \in C_{\infty}(X)$  s.t.  $(\forall \bar{x} \in X)(f(\bar{x}) < g(\bar{x}))$  then

$$(f,g)_X := \{(\bar{x},y) \in X \times M | f(\bar{x}) < y < g(\bar{x})\}$$





#### Cells

Let  $(i_1, \ldots, i_n)$  be a binary sequence. Then a  $(i_1, \ldots, i_n)$ -cell is a definable subset of  $M^n$  defined as follows,

- a (0)-cell is a singleton
- a (1)-cell is an open interval

Suppose that we have defined  $(i_1, \ldots, i_{n-1})$ -cells, then

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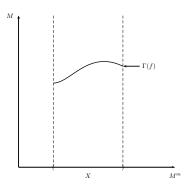
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Suppose that we have defined  $(i_1, \ldots, i_{n-1})$ -cells, then

• a  $(i_1, \ldots, i_{n-1}, 0)$ -cell is the graph of some  $f \in C(X)$ , where X is a  $(i_1, \ldots, i_{n-1})$ -cell



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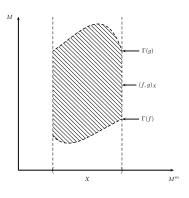
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- a  $(i_1, \ldots, i_{n-1}, 0)$ -cell is the graph of some  $f \in C(X)$ , where X is a  $(i_1, \ldots, i_{n-1})$ -cell
- a  $(i_1, \ldots, i_{n-1}, 1)$ -cell is the set  $(f, g)_X$  where X is a  $(i_1, \ldots, i_{n-1})$ -cell and  $f, g \in C_{\infty}(X)$





### Decompositions

A decomposition of  $M^n$  is a partition of  $M^n$  into finitely many cells. It is defined recursively.

 $\bullet$  A decomposition of M is a collection,

$$\{(-\infty, a_1), (a_1, a_2), \dots, (a_k, +\infty), \{a_1\}, \dots, \{a_k\}\}$$

where  $a_1 < \ldots < a_k$  are points in M;

• a decomposition of  $M^{n+1}$  is a finite partition of  $M^{n+1}$  into cells C s.t. the set of projections  $\pi(C)$  is a decomposition of  $M^n$ .

A decomposition  $\mathcal{D}$  of  $M^n$  is said to partition a set  $S \subseteq M^n$  if S is a union of cells in  $\mathcal{D}$ .

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### Cell Decomposition Theorem

#### Uniform Finiteness Property [Knight et al., 1986]

Let  $Y \subseteq M^{n+1}$  be definable and also for every  $\bar{x} \in M^n$ ,

$$Y_{\bar{x}} := \{ r \in M | (\bar{x}, r) \in Y \}$$

is finite. Then there exists  $N < \omega$  s.t.  $Y_{\bar{x}} \leq N$  for all  $\bar{x} \in M^n$ .

#### Cell Decomposition Theorem [Knight et al., 1986]

- Given any definable sets  $A_1, \ldots, A_k \subseteq M^n$  there is a decomposition of  $M^n$  partitioning each of  $A_1, \ldots, A_k$ .
- ② For each definable function  $f: A \to M$ ,  $A \subseteq M^n$ , there is a decomposition  $\mathcal{D}$  of  $M^n$  partitioning A s.t. the restriction  $f|B: B \to M$  to each cell  $B \in \mathcal{D}$  with  $B \subseteq A$  is continuous.





#### Theorem [Knight et al., 1986]

If  $\mathcal{M}$  is an o-minimal structure, the  $\mathrm{Th}(\mathcal{M})$  is an o-minimal theory.

We will show that if  $\mathcal{N} \equiv \mathcal{M}$  then  $\mathcal{N}$  is o-minimal.



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Let  $\phi(x_1,\ldots,x_m,y)$  be a formula.

By o-minimality for all  $\bar{a} \in M^m$  the set,

$$\phi(\bar{a}, y)^M := y \in M | \mathcal{M} \vDash \phi(\bar{a}, y)$$

is finite union of singletons and intervals. Consider now  $\psi(\bar{x},z)$  s.t.

$$\psi(\bar{x},z)^M := (\bar{x},z)|z \in \mathrm{bd}(y|\mathcal{M} \vDash \phi(\bar{x},y))$$



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- ★ boundary is definable
- $\star$  boundary is finite for definable subsets of M
- $\star \psi$  is finite over  $M^m$



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### Applications

#### Theorem [Knight et al., 1986]

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We will show that if  $\mathcal{N} \equiv \mathcal{M}$  then  $\mathcal{N}$  is o-minimal.

By UFP,  $\psi(\bar{x}, z)$  will be uniformly finite over  $M^m$ .

Thus for some  $\kappa < \omega$  we have that,

$$\mathcal{M} \vDash (\forall \bar{x})(\{y|\phi(\bar{x},y)\} \text{ has at most } \kappa \text{ boundary points})$$

Therefore for any  $\mathcal{N} \equiv \mathcal{M}$  and  $\bar{a}$ ,  $\phi(\bar{a}, y)^N$  has finite boundary points and so it will be a finite union of points and intervals.



#### Prime models

#### Definition

Let  $\mathcal{M} \models T$  and  $A \subseteq \mathcal{M}$ .  $\mathcal{M}$  is said to be *prime over* A if for any  $\mathcal{M}' \models T$  with  $A \subseteq \mathcal{M}'$ , there is an elementary mapping  $f : \mathcal{M} \to \mathcal{M}'$  that is identity on A.

#### Theorem [Pillay & Steinhorn, 1986]

Let  $A \subseteq \mathcal{M} \models T$ , where T is an o-minimal theory. Then there is an a model  $\mathcal{M}' \models T$ ,  $A \subseteq \mathcal{M}'$  that is prime over A, and is unique up to isomorphism over A.

### Ressayre's Theorem

#### Definition

Let  $\mathcal{M}$  be an L-structure and let  $A \subseteq M$ . Let  $\delta$  be an ordinal and  $(a_{\alpha} : \alpha < \delta)$  be a sequence of elements from M. Let  $A_{\alpha} = A \cup \{a_{\beta} : \beta < \alpha\}$ . We call  $(a_{\alpha} : \alpha < \delta)$  a construction over A if  $\operatorname{tp}(a_{\alpha}/A_{\alpha})$  is isolated for all  $\alpha < \delta$ .

#### Theorem [Ressayre]

Suppose that  $A \subseteq \mathbb{M}$ ,  $\mathcal{M} \prec \mathbb{M}$ ,  $\mathcal{N} \prec \mathbb{M}$ , and  $\mathcal{M}$  and  $\mathcal{N}$  are constructible over A. Then the identity map on A extends to an isomorphism between  $\mathcal{M}$  and  $\mathcal{N}$ .



#### Cuts and Noncuts

Let  $\mathcal{M} \models T$ . Let  $p(x) \in S_1(M)$ . We say that p(x) is a *cut* over  $\mathcal{M}$  iff there are non empty disjoint subsets  $C_0$  and  $C_1$  of M, s.t.  $C_0 \cup C_1 = M$ ,  $C_0$  has no greatest element,  $C_1$  has no least element, fo

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If p(x) is not isolated and is not a cut, we call p(x) a noncut.



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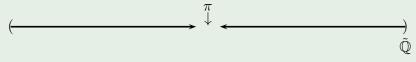
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#### Example

Let  $\tilde{\mathbb{Q}}$  be the field of real algebraic numbers. We know that  $\pi \in \mathbb{R} \setminus \tilde{\mathbb{Q}}$ , so  $\operatorname{tp}(\pi, \tilde{\mathbb{Q}})$  is a cut.

If t is an infinite hyperreal or t is infinitesimally close to an element of  $\tilde{\mathbb{Q}}$ , then  $\operatorname{tp}(t,\tilde{\mathbb{Q}})$  is a noncut.



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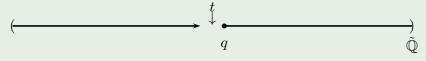
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#### Cuts and Noncuts

Let  $\mathcal{M} \models T$ , where T is an o-minimal theory.

#### Theorem [Marker, 1986]

Let  $\sigma(x) \in S_1(M)$  be a cut. Let  $\tau(x) \in S_1(M)$  be a noncut, and b be a realization of  $\tau(x)$ . Then  $\sigma(x)$  is omitted in  $P(\mathcal{M} \cup \{b\})$ .

#### Theorem [Marker, 1986]

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## Ordered Groups and Fields

#### Theorem [Pillay & Steinhorn, 1986]

Let  $\mathcal{G} = (G, +, 0, <)$  be an o-minimal ordered group. Then  $\mathcal{G}$  is a divisible oredered abelian group.

#### Reminder

An abelian group G is divisible if, for every positive integer n and every  $g \in G$ , there exists  $y \in G$  such that ny = g.

#### Theorem [Pillay & Steinhorn, 1986]

Let  $\mathcal{R} = (R, +, \cdot, 1, <)$  be an o-minimal ordered ring. Then  $\mathcal{R}$  is a real closed field.

#### Reminder

A field F is real closed if it is elementarily equivalent to  $\mathbb{R}$ .



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### Trichotomy Theorem

#### Question

Is there a sense in which an o-minimal structure is either "trivial", or group-like, or field-like?

Any answer to this has to be "local".



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### Trichotomy Theorem

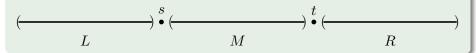
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Any answer to this has to be "local".

#### Example

- ullet L is carrying the structure of a pure dense linear order
- $\bullet$  M is carrying the structure of a divisible abelian ordered group
- R is carrying the structure of a real closed field





### Trichotomy Theorem

#### Trichotomy Theorem [Peterzil & Starchenko, 1998]

Let  $\mathcal{M}$  be a sufficiently saturated o-minimal structure. Given an  $a \in M$  one and only one of the following holds,

- a is trivial,
- the structure that  $\mathcal{M}$  induces on some convex neighborhood of a is an ordered vector space over an ordered division ring.
- the structure that  $\mathcal{M}$  induces over some open interval around a is an o-minimal expansion of a real closed field.

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## Weakly o-minimal structures

Let  $\mathcal{M} = (M, <, ...)$  be a linearly ordered structure.

#### Definition

A set  $C \subseteq M$  is called convex, if for any  $a,b \in C$  with a < b, and  $c \in M$  such that a < c < b, then  $c \in C$ .

#### Definition

A structure  $\mathcal{M}$  will be called weakly o-minimal, if the definable subsets of  $\mathcal{M}$  are finite unions of convex sets in (M,<).

We say that a complete theory T is weakly o-minimal if every model of T is weakly o-minimal.

#### Theorem

Expanding an o-minimal structure with unary predicates for convex subsets yields a structure with weakly o-minimal theory.



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## Monotonicity

Let  $\mathcal{M} = (M, <, P, Q, f)$  such that.

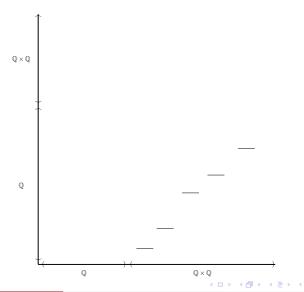
- $\bullet$  M is the disjoint union of the interpretations of the unary relations P and Q
- P is the interpretation of  $\mathbb{Q}$  with the usual order
- Q is the interpretation of  $\mathbb{Q} \times \mathbb{Q}$ , lexicographically ordered
- P proceeds Q in < on M
- $f: Q \to P, f((n,m)) = n \text{ for all } n, m \in \mathbb{Q}$

M is weakly o-minimal and also  $Th(\mathcal{M})$  is weakly o-minimal.





## Monotonicity



## 

### Weakly o-minimal structures

- ★ We have a local monotonicity theorem. [Arefiev, 1997]
- ★ Weakly o-minimal structures do not neccesserily have weakly o-minimal theory.
- ★ Weakly o-minimal structures do not neccesserily have prime models.

#### Theorem [Macpherson $et\ al.$ , 2000]

Every weakly o-minimal ordered group is divisible and abelian.

#### Theorem [Macpherson $et\ al.\ ,\ 2000]$

Every weakly o-minimal ordered field is real closed.



## 

## A definition for minimality

Let  $L \subset L^+$  be languages, and  $\mathcal{K}$  be an elementary class of L-structures.

#### Definition [Macpherson & Steinhorn, 1996]

An  $L^+$ -structure  $\mathcal{M}$  is  $\mathcal{K}$ -minimal if the the reduct  $\mathcal{M}|_L$  is in  $\mathcal{K}$  and every  $L^+$ -definable subset of M is definable by a quantifier-free L-formula. A complete  $L^+$ -theory is  $\mathcal{K}$ -minimal if all its models are  $\mathcal{K}$ -minimal.

- ★ o-minimality is a special case of the above definition but not weak o-minimality.
- $\star$   $\mathcal{K}$ -minimality is closed under reducts to languages containg L, and under expansion by constants.





## C-minimality

Let C(x; y, z) be a ternary realation,  $L = \{C\}$ , and  $\mathcal{K}_C$  be the class of L-structures satisfying the following axioms.

- $(\forall xyz)[C(x;y,z) \to C(x;z,y)]$
- $(\forall xyz)[C(x;y,z) \to C(y;x,z)]$
- $\bullet \ (\forall xyzw)[C(x;y,z) \to (C(w;y,z) \vee C(x;w,z))]$
- $(\forall xy)[x \neq y \rightarrow (\exists z \neq y)C(x; y, z)]$
- $\bullet \ (\exists xy)(x \neq y)$

#### Definition

A structure  $\mathcal{M} = (M, C, ...)$  is C-minimal if its theory is  $\mathcal{K}_C$ -minimal.



## P-minimality

#### Definition

Let  $L = (+, -, \cdot, 0, 1, (P_n)_{n>1})$ , where  $P_n$  are unary predicates. Regard  $\mathbb{Q}_p$  as an L-structure, letting  $P_n$  picking the  $n^{\text{th}}$  powers in  $\mathbf{Q}_p$ . Let  $\mathcal{K}_P$  be the class of L-structures elementarily equivalent to  $\mathbb{Q}_p$ . Then if  $L^+ \supseteq L$ , an  $L^+$ -structure is P-minimal if all models of its theory are  $\mathcal{K}_P$ -minimal

## 

### Summary

	o-minimal	weakly	C-minimal	P-minimal
Monotonicity	1	local	<b>√</b>	✓
CDT	✓	✓	✓	iff it has Skolem functions <sup>1</sup>
Prime Model	✓	Х	Х	Х
Groups	DAG	DAG		
Fields	RCF	RCF	ACVF	
Exchange	✓	X	×	✓
IP	$\mathbf{X}^2$	X3	$\mathbf{X}^4$	<b>X</b> <sup>5</sup>

<sup>&</sup>lt;sup>1</sup>Mourgues [2009]

<sup>&</sup>lt;sup>2</sup>Pillay & Steinhorn [1986]

<sup>&</sup>lt;sup>3</sup>Macpherson et al. [2000]

<sup>&</sup>lt;sup>4</sup>Macpherson & Steinhorn [1996]

<sup>&</sup>lt;sup>5</sup>Haskell & Macpherson [1997]

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