o-Minimality and its Variations

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o-Minimality

Assumptions

- $\mathcal{M} = (M, <, ...)$
- < is dense, linear, without endpoints
- definability with parameters

Definition

The structure \mathcal{M} is called *o-minimal* if every definable subset of M is a finite union of singletons and open intervals with endpoints in $M_{\infty} := M \cup \{-\infty, +\infty\}.$

A theory T is called o-minimal if every model \mathcal{M} of T is o-minimal.

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o-Minimality

The class of o-minimal structures is,

- closed under reducts (if < still remains in the language)
- closed under expansions by constants

Some o-minimal structures

- \bullet $(\mathbb{Q},<)$
- \bullet ($\mathbb{Q}, <, +$)
- $\mathcal{R} = (\mathbb{R}, <, +, -, \cdot, 0, 1)$

$(\mathbb{Q}, <, +, \cdot, 0, 1)$ is not o-minimal

The infinite discrete set of perfect squares is definable.



Monotonicity and Finiteness

Monotonicity Theorem

Let \mathcal{M} be an o-minimal structure and $f:(a,b) \to M$ be a definable function with domain (a,b) (possibly $a=-\infty$ or $b=+\infty$). Then, there are points $a=a_0 < a_1 < \cdots < a_{k+1}$ s.t. for each $j=0,\ldots,k,$ $f|_{(a_j,a_{j+1})}$ is either,

- constant, or
- a strictly monotonic and continuous bijection to an interval.

Finiteness Lemma

Let $A\subseteq M^2$ be definable and suppose that for each $x\in M$ the fiber $A_x:=\{y\in M|(x,y)\in A\}$ is finite. Then there is $N<\omega$ s.t. $|A_x|\leq N$ for all $x\in M$.





Applications

Corollary

Let $f:(a,b)\to M$ be definale and continuous. Then f takes a maximum and minimum value on [a,b].

Exchange Principle

Let \mathcal{M} be o-minimal. Let $b, c, a_1, \ldots, a_n \in \mathcal{M}$. If b is definable over c, a_1, \ldots, a_n , and b is not definable over a_1, \ldots, a_n , then c is definable over b, a_1, \ldots, a_n .

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Motivation

Question

What happen with the definable subsets of M^n ?

Notation

Given definable $X \subseteq M^n$

- $C(X) := \{f : X \to M | f \text{ is definable and continuous} \}$
- $C_{\infty}(X) := C(X) \cup \{-\infty, +\infty\}$
- for $f, g \in C_{\infty}(X)$ s.t. $(\forall \bar{x} \in X)(f(\bar{x}) < g(\bar{x}))$ then

$$(f,g)_X := \{(\bar{x},y) \in X \times M | f(\bar{x}) < y < g(\bar{x})\}$$





Cells

Let (i_1, \ldots, i_n) be a binary sequence. Then a (i_1, \ldots, i_n) -cell is a definable subset of M^n defined as follows,

- a (0)-cell is a singleton
- a (1)-cell is an open interval

Suppose that we have defined (i_1, \ldots, i_{n-1}) -cells, then

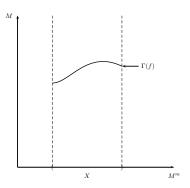
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Suppose that we have defined (i_1, \ldots, i_{n-1}) -cells, then

• a $(i_1, \ldots, i_{n-1}, 0)$ -cell is the graph of some $f \in C(X)$, where X is a (i_1, \ldots, i_{n-1}) -cell



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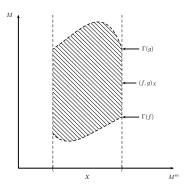
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- a $(i_1, \ldots, i_{n-1}, 0)$ -cell is the graph of some $f \in C(X)$, where X is a (i_1, \ldots, i_{n-1}) -cell
- a $(i_1, \ldots, i_{n-1}, 1)$ -cell is the set $(f, g)_X$ where X is a (i_1, \ldots, i_{n-1}) -cell and $f, g \in C_{\infty}(X)$





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Decompositions

A decomposition of M^n is a partition of M^n into finitely many cells. It is defined recursively.

 \bullet A decomposition of M is a collection,

$$\{(-\infty, a_1), (a_1, a_2), \dots, (a_k, +\infty), \{a_1\}, \dots, \{a_k\}\}$$

where $a_1 < \ldots < a_k$ are points in M;

• a decomposition of M^{n+1} is a finite partition of M^{n+1} into cells C s.t. the set of projections $\pi(C)$ is a decomposition of M^n .

A decomposition \mathcal{D} of M^n is said to partition a set $S \subseteq M^n$ if S is a union of cells in \mathcal{D} .

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Cell Decomposition Theorem

Uniform Finiteness Property

Let $Y \subseteq M^{n+1}$ be definable and also for every $\bar{x} \in M^n$,

$$Y_{\bar{x}} := \{ r \in M | (\bar{x}, r) \in Y \}$$

is finite. Then there exists $N < \omega$ s.t. $Y_{\bar{x}} \leq N$ for all $\bar{x} \in M^n$.

Cell Decomposition Theorem

- Given any definable sets $A_1, \ldots, A_k \subseteq M^n$ there is a decomposition of M^n partitioning each of A_1, \ldots, A_k .
- ② For each definable function $f: A \to M$, $A \subseteq M^n$, there is a decomposition \mathcal{D} of M^n partitioning A s.t. the restriction $f|B: B \to M$ to each cell $B \in \mathcal{D}$ with $B \subseteq A$ is continuous.





Applications

Theorem

If \mathcal{M} is an o-minimal structure, the Th(\mathcal{M}) is an o-minimal theory.

We will show that if $\mathcal{N} \equiv \mathcal{M}$ then \mathcal{N} is o-minimal.



Applications

Theorem

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We will show that if $\mathcal{N} \equiv \mathcal{M}$ then \mathcal{N} is o-minimal.

Let $\phi(x_1,\ldots,x_m,y)$ be a formula.

By o-minimality for all $\bar{a} \in M^m$ the set,

$$\phi(\bar{a}, y)^M := y \in M | \mathcal{M} \vDash \phi(\bar{a}, y)$$

is finite union of singletons and intervals. Consider now $\psi(\bar{x},z)$ s.t.

$$\psi(\bar{x},z)^M := (\bar{x},z)|z \in \mathrm{bd}(y|\mathcal{M} \vDash \phi(\bar{x},y))$$

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- ★ boundary is definable
- \star boundary is finite for definable subsets of M
- $\star \psi$ is finite over M^m





Applications

Theorem

If \mathcal{M} is an o-minimal structure, the $\mathrm{Th}(\mathcal{M})$ is an o-minimal theory.

We will show that if $\mathcal{N} \equiv \mathcal{M}$ then \mathcal{N} is o-minimal.

By UFP, $\psi(\bar{x}, z)$ will be uniformly finite over M^m .

Thus for some $\kappa < \omega$ we have that,

$$\mathcal{M} \vDash (\forall \bar{x})(\{y|\phi(\bar{x},y)\} \text{ has at most } \kappa \text{ boundary points})$$

Therefore for any $\mathcal{N} \equiv \mathcal{M}$ and \bar{a} , $\phi(\bar{a}, y)^N$ has finite boundary points and so it will be a finite union of points and intervals.



Prime models

Definition

Let $\mathcal{M} \models T$ and $A \subseteq \mathcal{M}$. \mathcal{M} is said to be *prime over* A if for any $\mathcal{M}' \models T$ with $A \subseteq \mathcal{M}'$, there is an elementary mapping $f : \mathcal{M} \to \mathcal{M}'$ that is identity on A.

Theorem

Let $A \subseteq \mathcal{M} \models T$, where T is an o-minimal theory. Then there is an a model $\mathcal{M}' \models T$, $A \subseteq \mathcal{M}'$ that is prime over A, and is unique up to isomorphism over A.



Ressayre's Theorem

Definition

Let \mathcal{M} be an L-structure and let $A \subseteq M$. Let δ be an ordinal and $(a_{\alpha} : \alpha < \delta)$ be a sequence of elements from M. Let $A_{\alpha} = A \cup \{a_{\beta} : \beta < \alpha\}$. We call $(a_{\alpha} : \alpha < \delta)$ a construction over A if $\operatorname{tp}(a_{\alpha}/A_{\alpha})$ is isolated for all $\alpha < \delta$.

Theorem [Ressayre]

Suppose that $A \subseteq \mathbb{M}$, $\mathcal{M} \prec \mathbb{M}$, $\mathcal{N} \prec \mathbb{M}$, and \mathcal{M} and \mathcal{N} are constructible over A. Then the identity map on A extends to an isomorphism between \mathcal{M} and \mathcal{N} .



Cuts and Noncuts

Let $\mathcal{M} \models T$. Let $p(x) \in S_1(M)$. We say that p(x) is a *cut* over \mathcal{M} iff there are non empty disjoint subsets C_0 and C_1 of M, s.t. $C_0 \models C_1 = M$. C_2 has no greatest element. C_3 has no least element for

 $C_0 \cup C_1 = M$, C_0 has no greatest element, C_1 has no least element, for all $c \in C_0$, "c < x" $\in p(x)$, and for all $c \in C_1$, "x < c" $\in p(x)$.

If p(x) is not isolated and is not a cut, we call p(x) a noncut.



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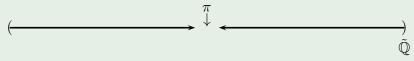
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Example

Let $\tilde{\mathbb{Q}}$ be the field of real algebraic numbers. We know that $\pi \in \mathbb{R} \setminus \tilde{\mathbb{Q}}$, so $\operatorname{tp}(\pi, \tilde{\mathbb{Q}})$ is a cut.

If t is an infinite hyperreal or t is infinitesimally close to an element of $\tilde{\mathbb{Q}}$, then $\operatorname{tp}(t,\tilde{\mathbb{Q}})$ is a noncut.



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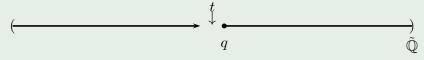
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Cuts and Noncuts

Let $\mathcal{M} \models T$, where T is an o-minimal theory.

Theorem

Let $\sigma(x) \in S_1(M)$ be a cut. Let $\tau(x) \in S_1(M)$ be a noncut, and b be a realization of $\tau(x)$. Then $\sigma(x)$ is omitted in $P(\mathcal{M} \cup \{b\})$.

Theorem

Let $\tau(x) \in S_1(M)$ be a noncut. Let $\sigma(x) \in S_1(M)$ be a cut, and b be a realization of $\sigma(x)$. Then $\tau(x)$ is omitted in $P(\mathcal{M} \cup \{b\})$.

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Ordered Groups and Fields

Theorem

Let $\mathcal{G} = (G, +, 0, <)$ be an o-minimal ordered group. Then \mathcal{G} is a divisible oredered abelian group.

Reminder

An abelian group G is divisible if, for every positive integer n and every $g \in G$, there exists $y \in G$ such that ny = g.

Theorem

Let $\mathcal{R} = (R, +, \cdot, 1, <)$ be an o-minimal ordered ring. Then \mathcal{R} is a real closed field.

Reminder

A field F is real closed if it is elementarily equivalent to \mathbb{R} .



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Trichotomy Theorem

Question

Is there a sense in which an o-minimal structure is either "trivial", or group-like, or field-like?

Any answer to this has to be "local".



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Trichotomy Theorem

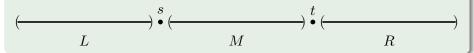
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Example

- \bullet L is carrying the structure of a pure dense linear order
- ullet M is carrying the structure of a divisible abelian ordered group
- R is carrying the structure of a real closed field



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Trichotomy Theorem

Trichotomy Theorem

Let \mathcal{M} be a sufficiently saturated o-minimal structure. Given an $a \in M$ one and only one of the following holds,

- a is trivial,
- the structure that \mathcal{M} induces on some convex neighborhood of a is an ordered vector space over an ordered division ring.
- the structure that \mathcal{M} induces over some open interval around a is an o-minimal expansion of a real closed field.



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Weakly o-minimal structures

Let $\mathcal{M} = (M, <, ...)$ be a linearly ordered structure.

Definition

A set $C \subseteq M$ is called convex, if for any $a,b \in C$ with a < b, and $c \in M$ such that a < c < b, then $c \in C$.

Definition

A structure \mathcal{M} will be called weakly o-minimal, if the definable subsets of \mathcal{M} are finite unions of convex sets in (M,<).

We say that a complete theory T is weakly o-minimal if every model of T is weakly o-minimal.

Theorem

Expanding an o-minimal structure with unary predicates for convex subsets yields a structure with weakly o-minimal theory.



Monotonicity

Let $\mathcal{M} = (M, <, P, Q, f)$ such that.

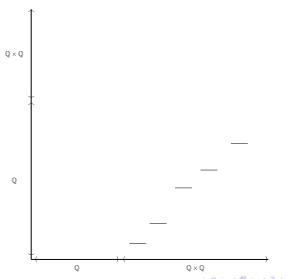
- M is the disjoint union of the interpretations of the unary relations P and Q
- P is the interpretation of \mathbb{Q} with the usual order
- Q is the interpretation of $\mathbb{Q} \times \mathbb{Q}$, lexicographically ordered
- P proceeds Q in < on M
- $f: Q \to P, f((n,m)) = n \text{ for all } n, m \in \mathbb{Q}$

M is weakly o-minimal and also $Th(\mathcal{M})$ is weakly o-minimal.



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Monotonicity



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Weakly o-minimal structures

- \star We have a local monotonicity theorem.
- ★ Weakly o-minimal structures do not neccesserily have weakly o-minimal theory.
- ★ Weakly o-minimal structures do not neccesserily have prime models.

Theorem

Every weakly o-minimal ordered group is divisible and abelian.

Theorem

Every weakly o-minimal ordered field is real closed.



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A definition for minimality

Let $L \subset L^+$ be languages, and \mathcal{K} be an elementary class of L-structures.

Definition

An L^+ -structure \mathcal{M} is \mathcal{K} -minimal if the the reduct $\mathcal{M}|_L$ is in \mathcal{K} and every L^+ -definable subset of M is definable by a quantifier-free L-formula. A complete L^+ -theory is \mathcal{K} -minimal if all its models are \mathcal{K} -minimal.

- ★ o-minimality is a special case of the above definition but not weak o-minimality.
- \star K-minimality is closed under reducts to languages containg L, and under expansion by constants.



C-minimality

Let C(x; y, z) be a ternary realation, $L = \{C\}$, and \mathcal{K}_C be the class of L-structures satisfying the following axioms.

- $(\forall xyz)[C(x;y,z) \to C(x;z,y)]$
- $(\forall xyz)[C(x;y,z) \to C(y;x,z)]$
- $(\forall xyzw)[C(x;y,z) \rightarrow (C(w;y,z) \lor C(x;w,z))]$
- $(\forall xy)[x \neq y \rightarrow (\exists z \neq y)C(x;y,z)]$
- \bullet $(\exists xy)(x \neq y)$

Definition

A structure $\mathcal{M} = (M, C, ...)$ is C-minimal if its theory is \mathcal{K}_C -minimal.

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P-minimality

Definition

Let $L = (+, -, \cdot, 0, 1, (P_n)_{n>1})$, where P_n are unary predicates. Regard \mathbb{Q}_p as an L-structure, letting P_n picking the n^{th} powers in $mathbold Q_p$. Let \mathcal{K}_P be the class of L-structures elementarily equivalent to \mathbb{Q}_p . Then if $L^+ \supseteq L$, an L^+ -structure is P-minimal if all models of its theory are \mathcal{K}_P -minimal

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Summary

	o-minimal	weakly	C-minimal	P-minimal
Monotonicity	✓	local	✓	✓
CDT	✓	✓	✓	iff it has Skolem functions
Prime Model	✓	Х	Х	Х
Groups	DAG	DAG		
Fields	RCF	RCF	ACVF	
Exchange	✓	X	X	✓
IP	X	X	Х	Х

