### o-Minimality and its Variations

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### o-Minimality

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### Weakly o-minimal structures

Let  $\mathcal{M} = (M, <, \dots)$  be a linearly ordered structure.

#### Definition

A set  $C \subseteq M$  is called convex, if for any  $a,b \in C$  with a < b, and  $c \in M$  such that a < c < b, then  $c \in C$ .

#### Definition

A structure  $\mathcal{M}$  will be called weakly o-minimal, if the definable subsets of  $\mathcal{M}$  are finite unions of convex sets in (M,<).

We say that a complete theory T is weakly o-minimal if every model of T is weakly o-minimal.

#### Theorem

Expanding an o-minimal structure with unary predicates for convex subsets yields a structure with weakly o-minimal theory.



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## Monotonicity

Let  $\mathcal{M} = (M, <, P, Q, f)$  such that.

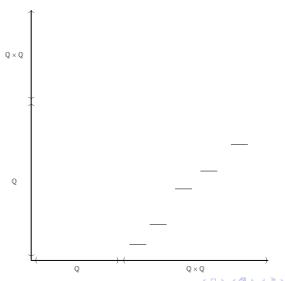
- $\bullet$  M is the disjoint union of the interpretations of the unary relations P and Q
- P is the interpretation of  $\mathbb{Q}$  with the usual order
- Q is the interpretation of  $\mathbb{Q} \times \mathbb{Q}$ , lexicographically ordered
- P proceeds Q in < on M
- $f: Q \to P, f((n,m)) = n \text{ for all } n, m \in \mathbb{Q}$

M is weakly o-minimal and also  $Th(\mathcal{M})$  is weakly o-minimal.





## Monotonicity



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### Weakly o-minimal structures

- $\star$  We have a local monotonicity theorem.
- ★ Weakly o-minimal structures do not neccesserily have weakly o-minimal theory.
- ★ Weakly o-minimal structures do not neccesserily have prime models.

#### Theorem

Every weakly o-minimal ordered group is divisible and abelian.

#### $\overline{\text{Theorem}}$

Every weakly o-minimal ordered field is real closed.



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## A definition for minimality

Let  $L \subset L^+$  be languages, and  $\mathcal{K}$  be an elementary class of L-structures.

#### Definition

An  $L^+$ -structure  $\mathcal{M}$  is  $\mathcal{K}$ -minimal if the the reduct  $\mathcal{M}|_L$  is in  $\mathcal{K}$  and every  $L^+$ -definable subset of M is definable by a quantifier-free L-formula. A complete  $L^+$ -theory is  $\mathcal{K}$ -minimal if all its models are  $\mathcal{K}$ -minimal.

- ★ o-minimality is a special case of the above definition but not weak o-minimality.
- $\star$  K-minimality is closed under reducts to languages containg L, and under expansion by constants.



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## C-minimality

Let C(x; y, z) be a ternary realation,  $L = \{C\}$ , and  $\mathcal{K}_C$  be the class of L-structures satisfying the following axioms.

- $(\forall xyz)[C(x;y,z) \to C(x;z,y)]$
- $(\forall xyz)[C(x;y,z) \to C(y;x,z)]$
- $\bullet \ (\forall xyzw)[C(x;y,z) \to (C(w;y,z) \lor C(x;w,z))]$
- $(\forall xy)[x \neq y \rightarrow (\exists z \neq y)C(x; y, z)]$
- $(\exists xy)(x \neq y)$

#### Definition

A structure  $\mathcal{M} = (M, C, ...)$  is C-minimal if its theory is  $\mathcal{K}_C$ -minimal.

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## P-minimality

#### Definition

Let  $L = (+, -, \cdot, 0, 1, (P_n)_{n>1})$ , where  $P_n$  are unary predicates. Regard  $\mathbb{Q}_p$  as an L-structure, letting  $P_n$  picking the  $n^{\text{th}}$  powers in  $mathbold Q_p$ . Let  $\mathcal{K}_P$  be the class of L-structures elementarily equivalent to  $\mathbb{Q}_p$ . Then if  $L^+ \supseteq L$ , an  $L^+$ -structure is P-minimal if all models of its theory are  $\mathcal{K}_P$ -minimal

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### Summary

	o-minimal	weakly	C-minimal	P-minimal
Monotonicity	✓	local	1	✓
CDT	✓	✓	✓	iff it has Skolem functions
Prime Model	✓	X	Х	Х
Groups	DAG	DAG		
Fields	RCF	RCF	ACVF	
Exchange	✓	Х	X	✓
IP	Х	X	X	Х