

27.02.2024

B+ TREES → Search tree

BST

↳ Ordering

Multistage (M-way) Trees ↳ To reduce height

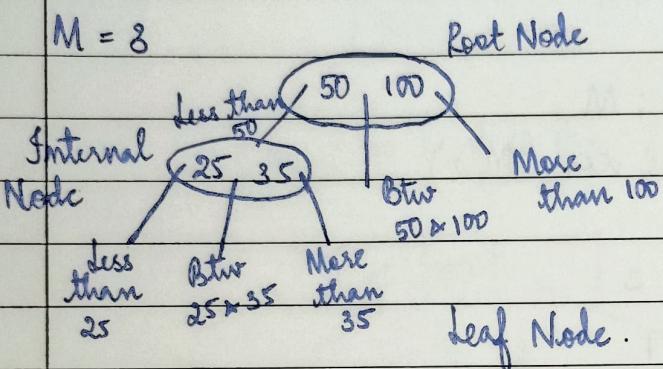
↳ Order of tree ↳ of a tree

Search Tree

For a BST $\Rightarrow M=2$

$$\text{key} = M-1 = 2-1 = 1$$

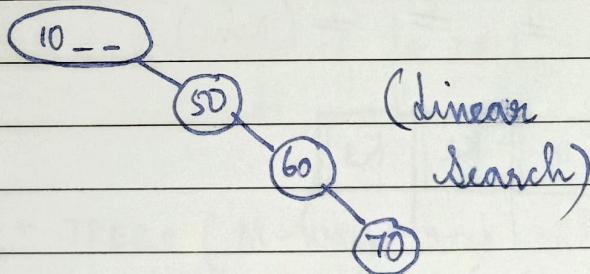
property
Height: $\log_2 n$
No. of nodes (n) is
very high \Rightarrow
Search time will
be very high



Search time becomes better

M can start with 3 and go on to any value.

$M=4$

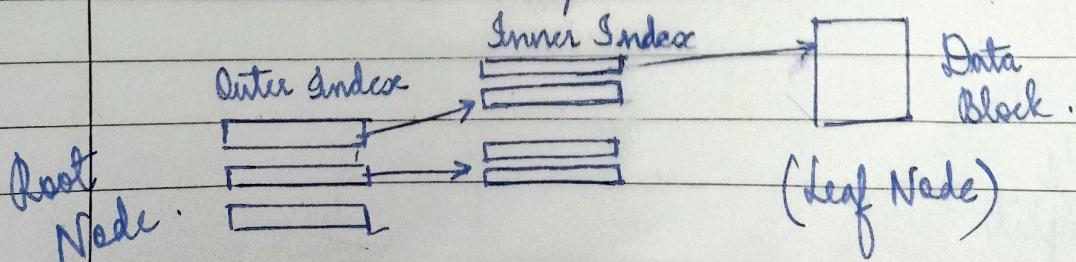


B & B+ tree \rightarrow Balanced M-way tree

↳ MORE
ADVANTAGE

Hddisk \rightarrow Chunk of blocks

MULTILEVEL INDEXING \rightarrow Faster access.



B⁺-tree → Search tree

M-way tree

all leaf node at same level.

Pointer for all block available at leaf node.

NOT APPLICABLE
TO A
LEAF NODE

M-way:

Keys → Max: M-1 keys.

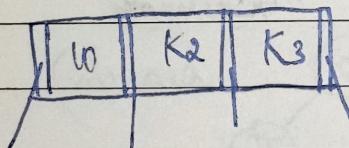
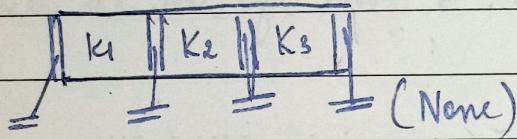
Min: ceil(M/2) - 1

Children → Max: M

Min: ceil(M/2)

M=4 Key = m-1 = 3

Min = 1

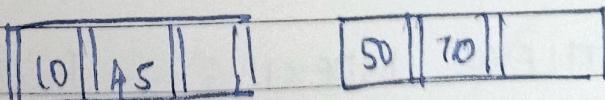


Keys → descending order

[[10 || 50 || 70]] (One level of index)

10 45 50 70 (Not possible (overflow))

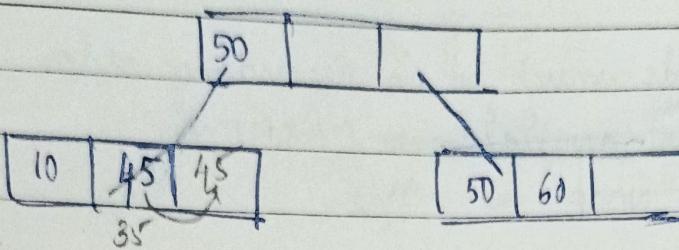
Overflow



Leaf Node

Split

Btree \rightarrow Keys need not be in leaf
has to be only there in internal nodes



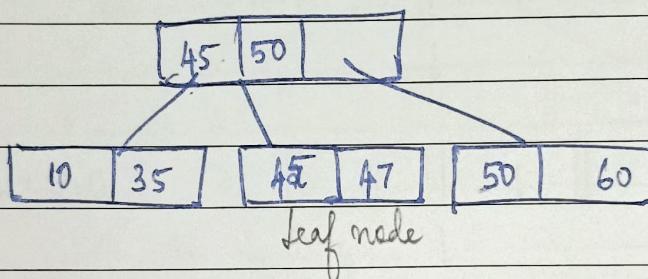
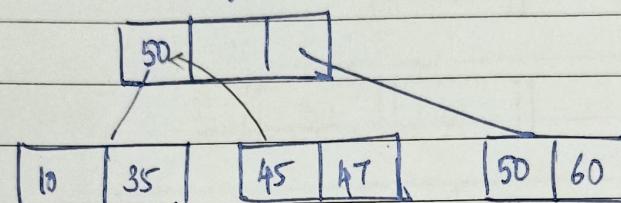
leaf node has exact data,

Insert 35

Insert 47

Overflow

All these data is organised linearly in hard disk
leaf node \rightarrow singly linked list



Internal node \rightarrow Node going to root must not be in internal node.

29.02.2024

PROPS OF B+ TREES (M-way trees with foll. props)

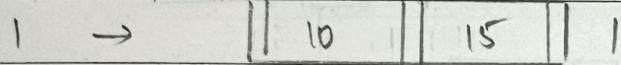
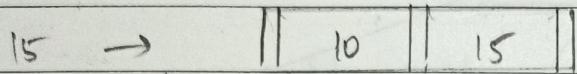
- 1) Order m
- 2) Children max : m min : ceil($m/2$)
- 3) Keys max : m - 1 min : ceil($m/2$) - 1
- 4) All keys should be there in leaf node.
- 5) All leaf same level.

$M = 3 \rightarrow$ Maximum 3 children
Maximum 2 keys

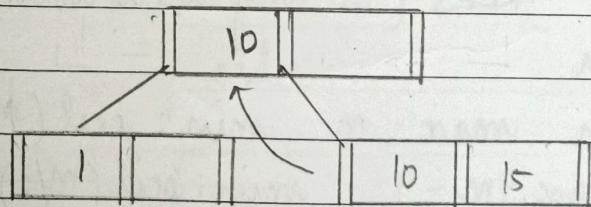
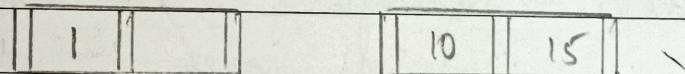


Keys on the node must be in ascending order.
More keys than available \Rightarrow OVERFLOW.
 \Rightarrow SPLIT NODE $\Rightarrow m/2$.

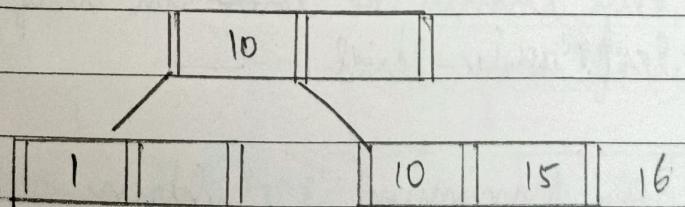
10, 15, 1, 16, 7, 25, 23, 17

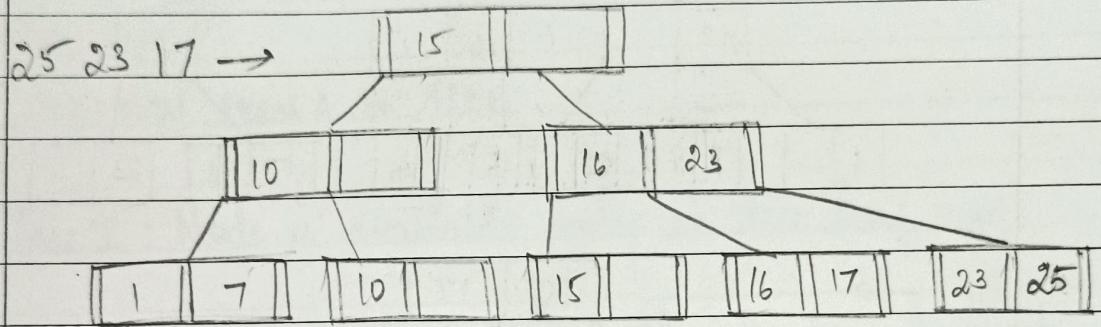
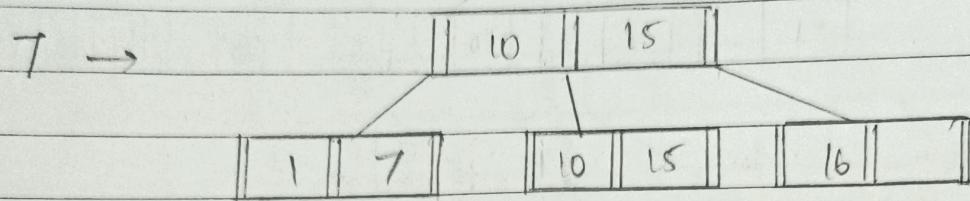
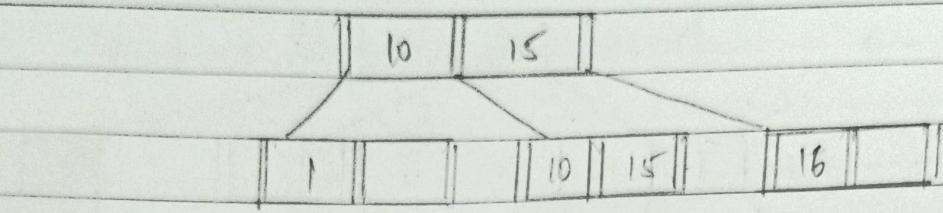


|| 1 || 10 || 15 OVERFLOW

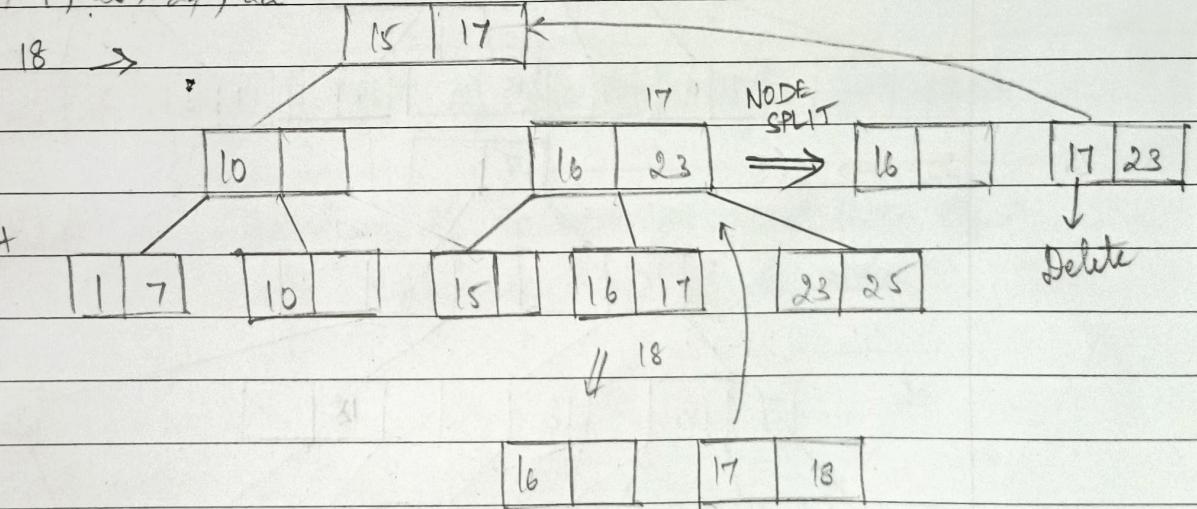


16 \rightarrow

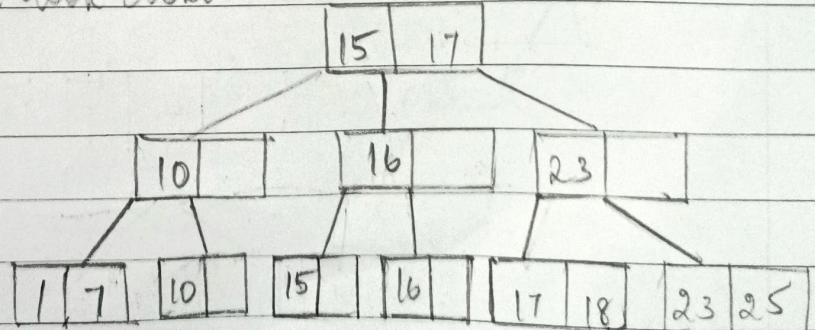




18, 9, 28, 24, 22



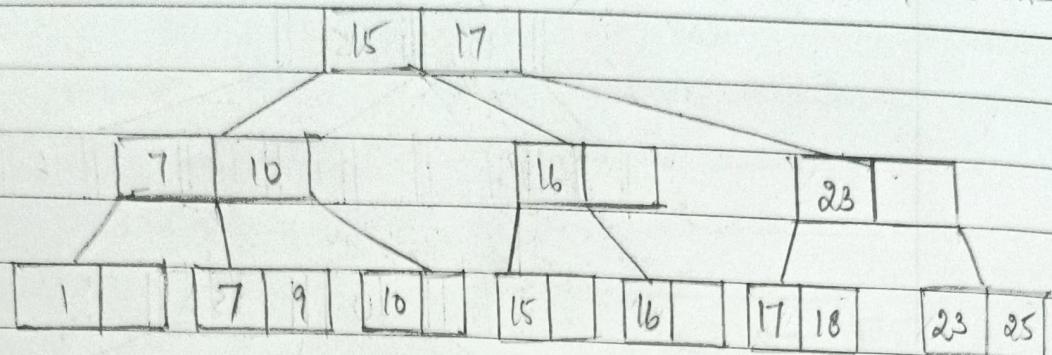
Now this will look like



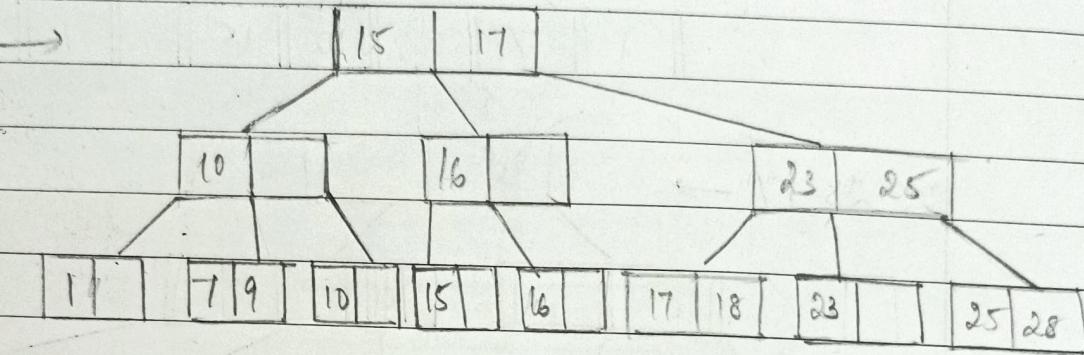
When data increases \Rightarrow

MULTI LEVEL INDEXING / INCREASES

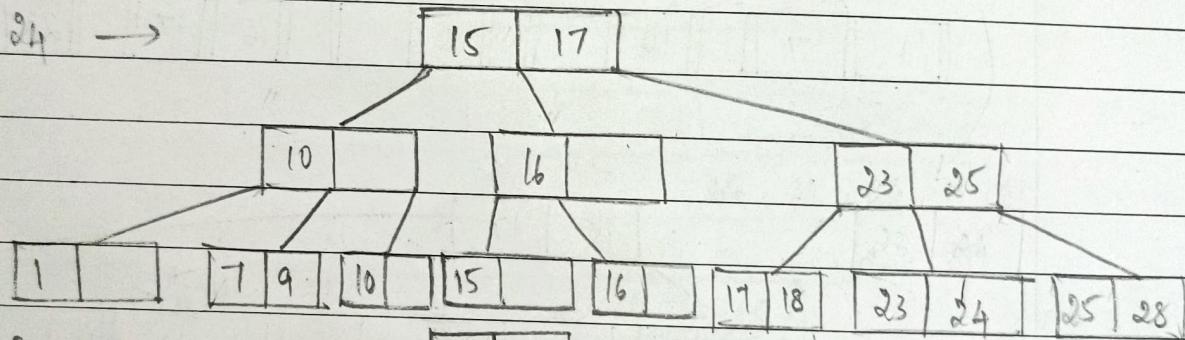
9 \rightarrow



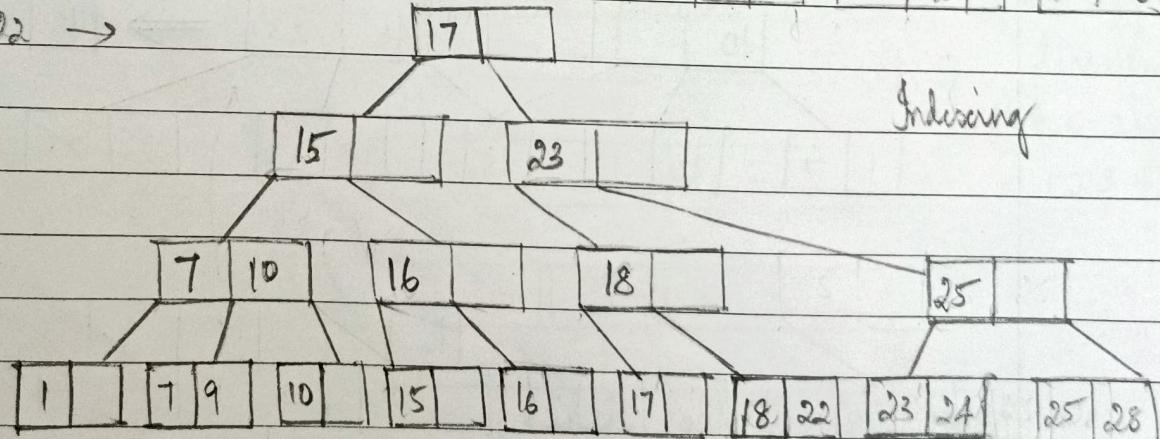
28 \rightarrow



34 \rightarrow



22 \rightarrow

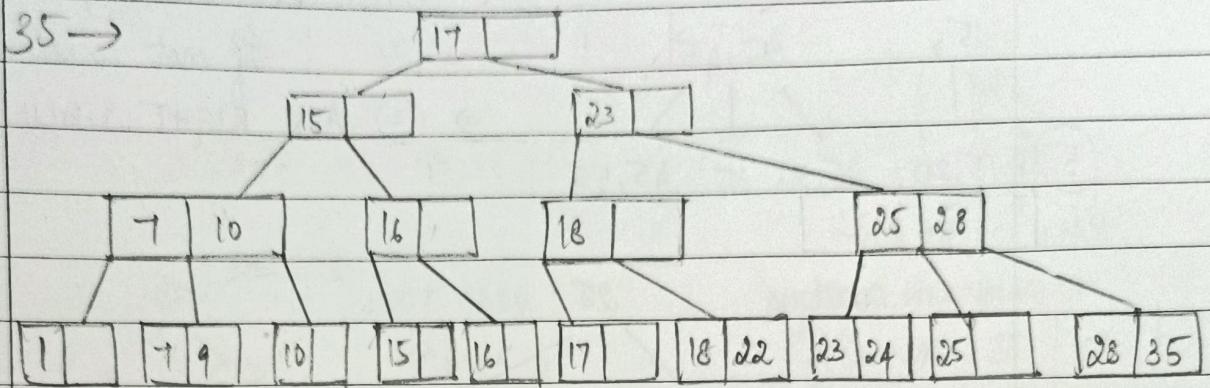


Indexing

AVL → Balanced
BST

B⁺ TREES → Self Balanced _ / _
No rotations

B⁺ TREES → Bottom up approach,



01.03.2024. DELETION FROM A B⁺ TREE

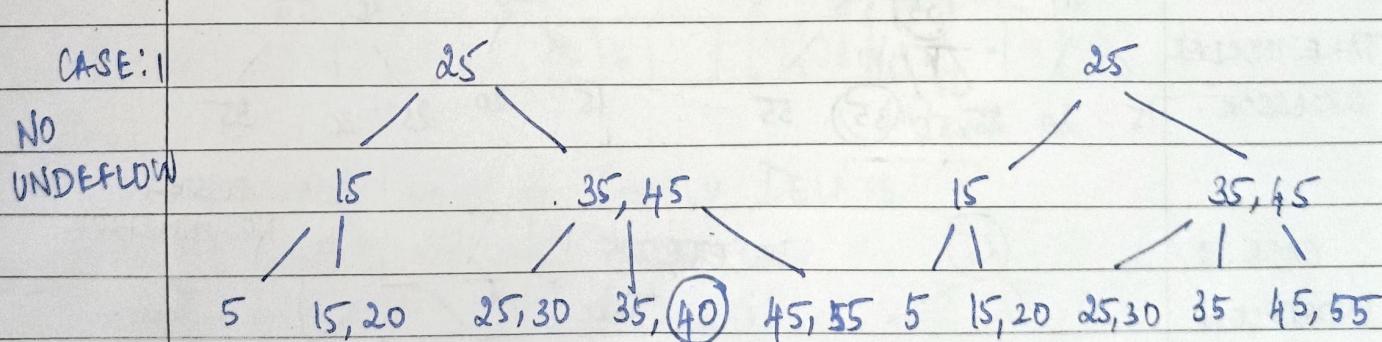
Case I : Node is available only in the leaf node

Case II : Node is available in the leaf node and internal node

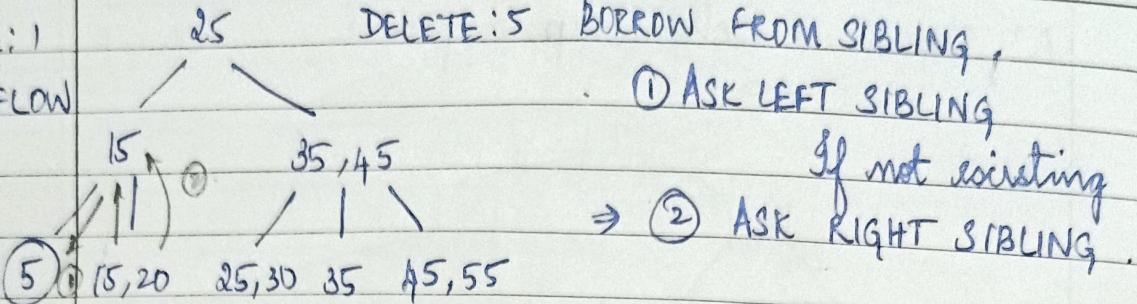
Case III : If the height of the tree gets shrunk

Underflow: Violation of minimum condition of a node (Minimum keys of a node)

CASE: I

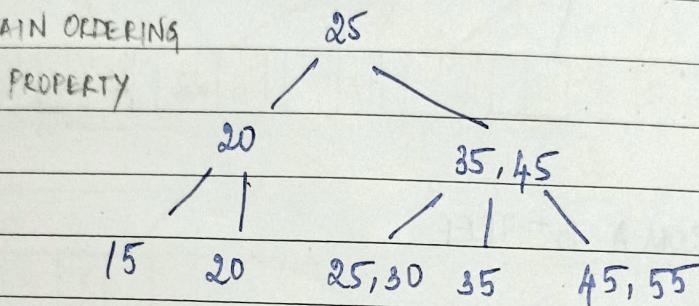


CASE: 1
UNDERFLOW

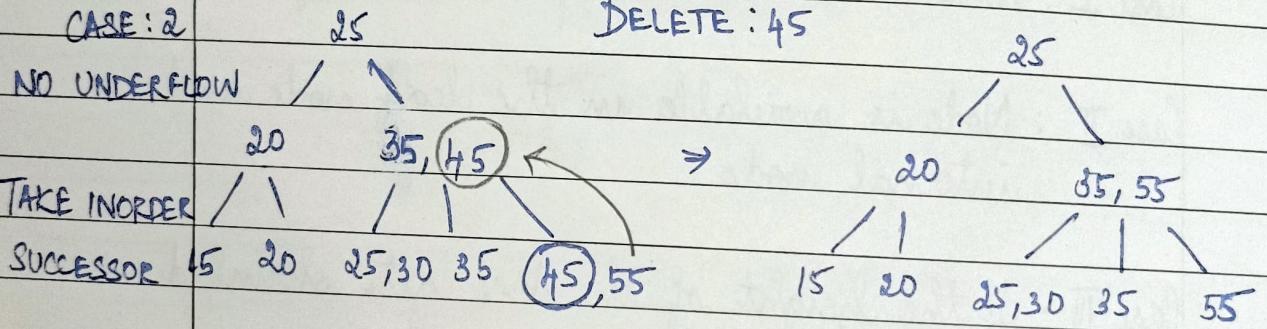


MAKE 20 AS ROOT

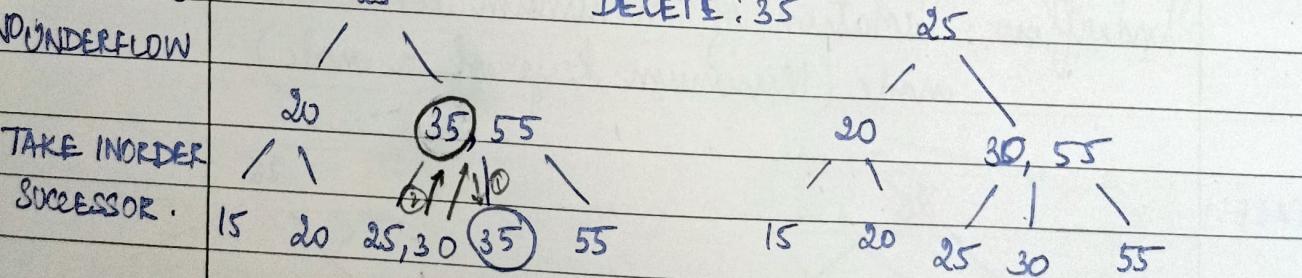
TO MAINTAIN ORDERING



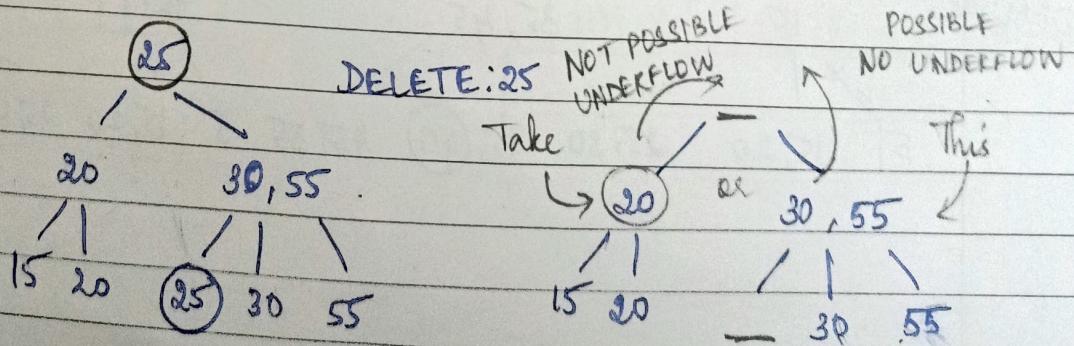
CASE: 2

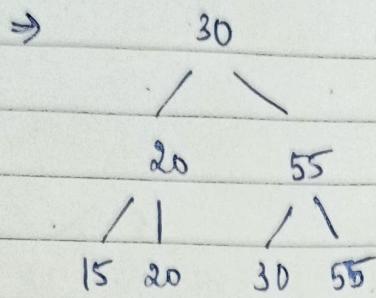


CASE: 2

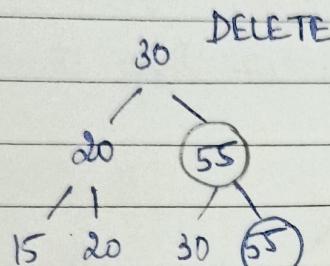


CASE: 2
UNDERFLOW



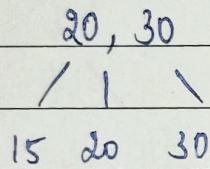
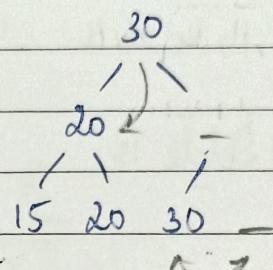


CASE : 3



NOT CASE : 1

CASE : 2
UNDERFLOW

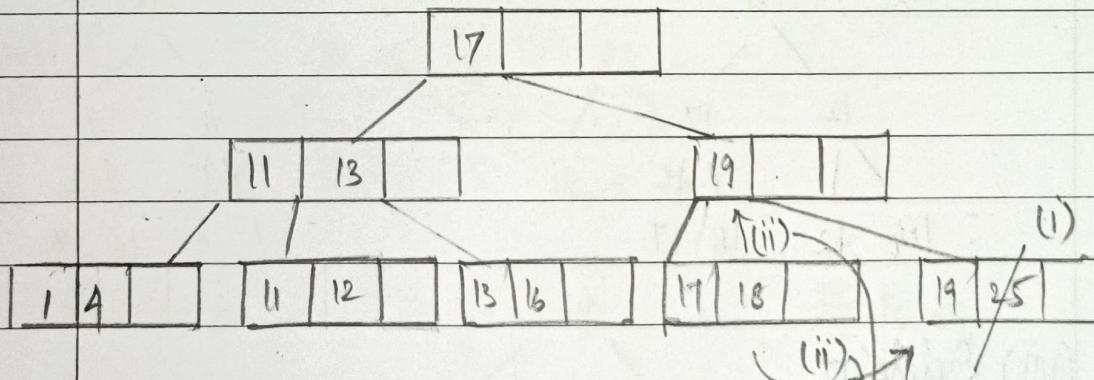


HEIGHT OF TREE

Merge

REDUCES BY 1.

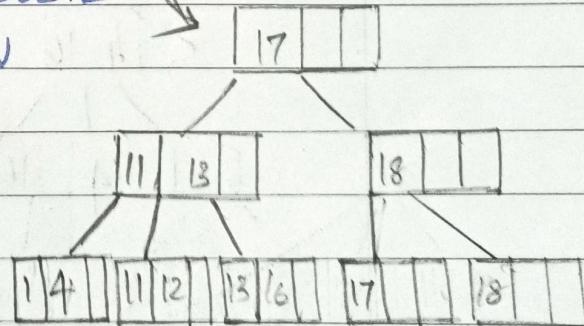
Ques: Order : 4



(i) Delete 25 → SIMPLY DELETE

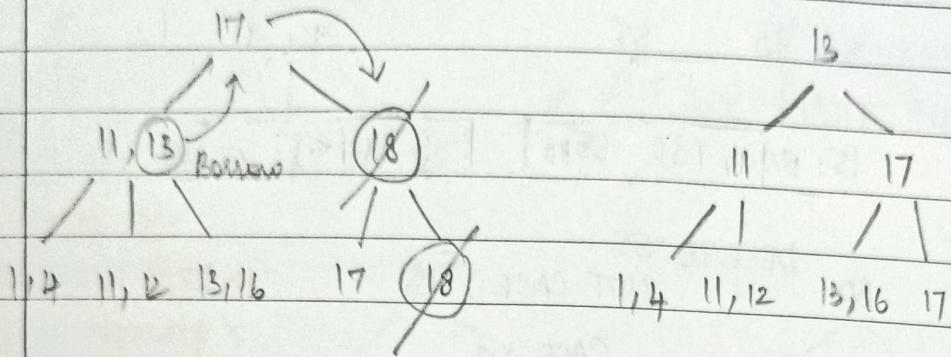
(ii) Delete 19 → UNDERFLOW

CASE (i)

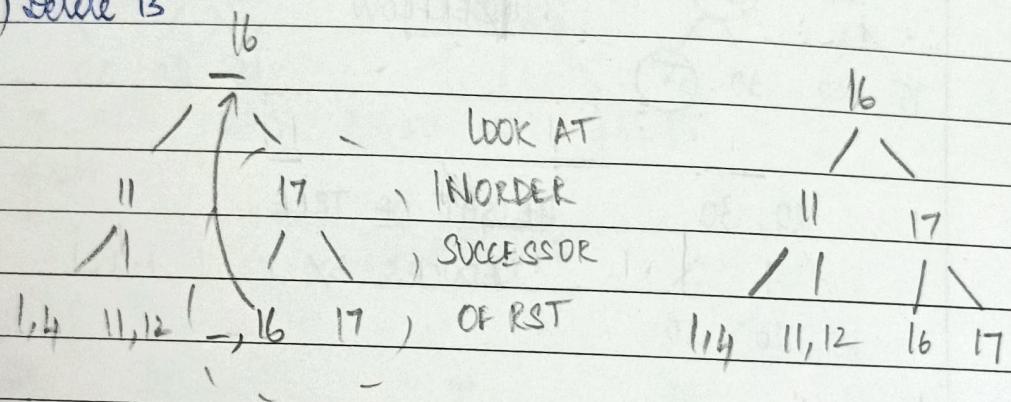


Any node can be in internal node once and leaf node

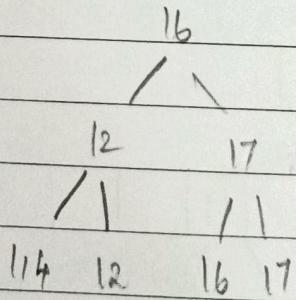
(iii) Delete 18.



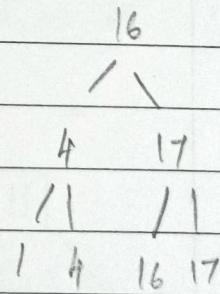
(iv) Delete 13



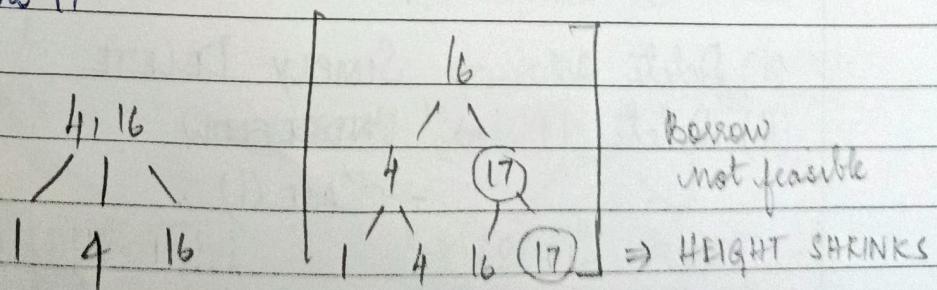
(v) Delete 11

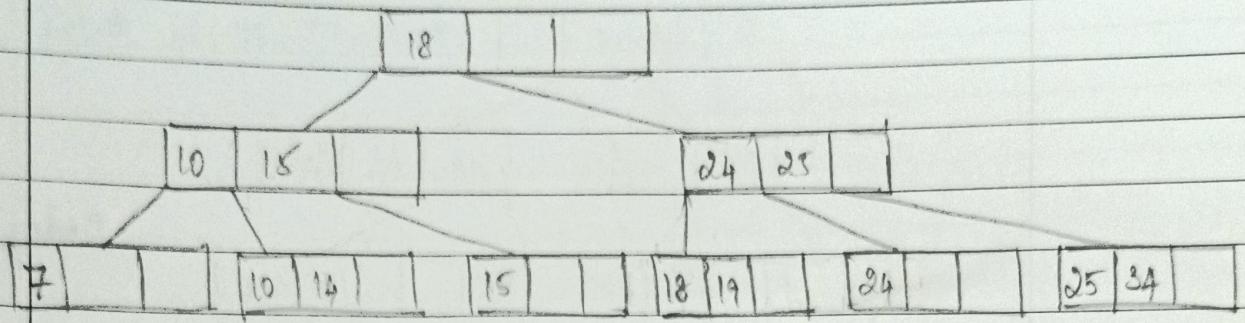


(vi) Delete 12

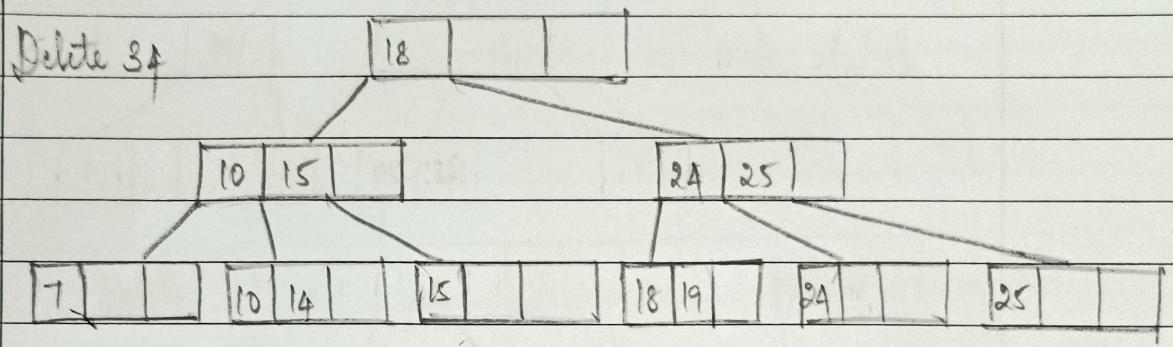


(vii) Delete 17

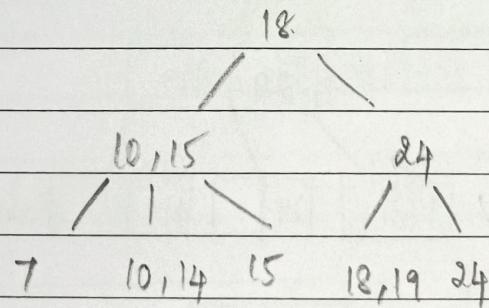




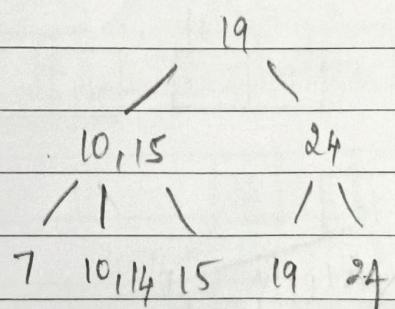
Delete 34, 25, 18, 24, 19, 7



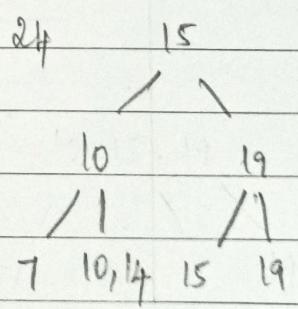
Delete 25



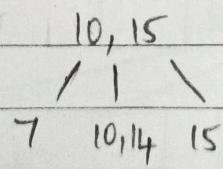
Delete 18



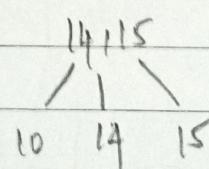
Delete 24

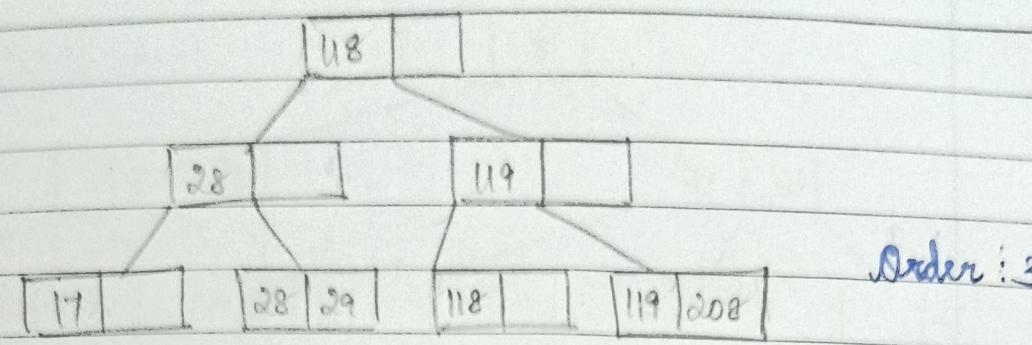


Delete 19

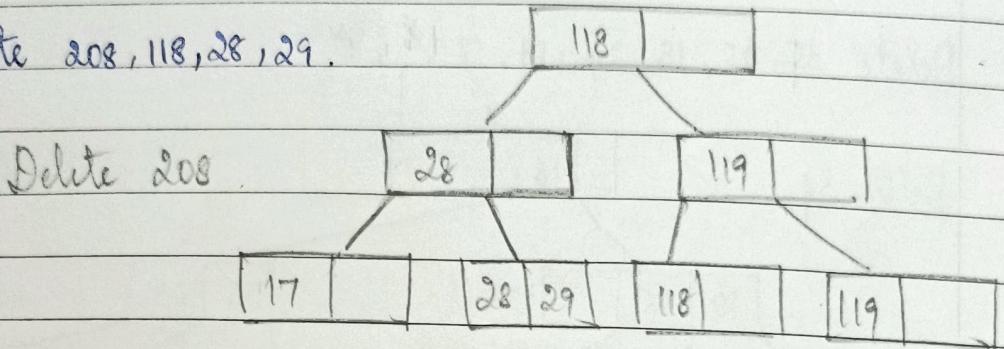


Delete 7

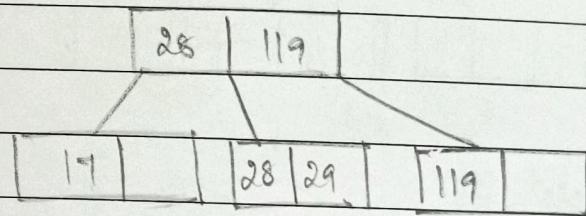




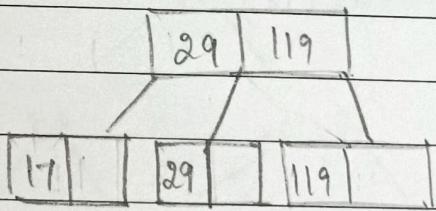
Delete 208, 118, 28, 29.



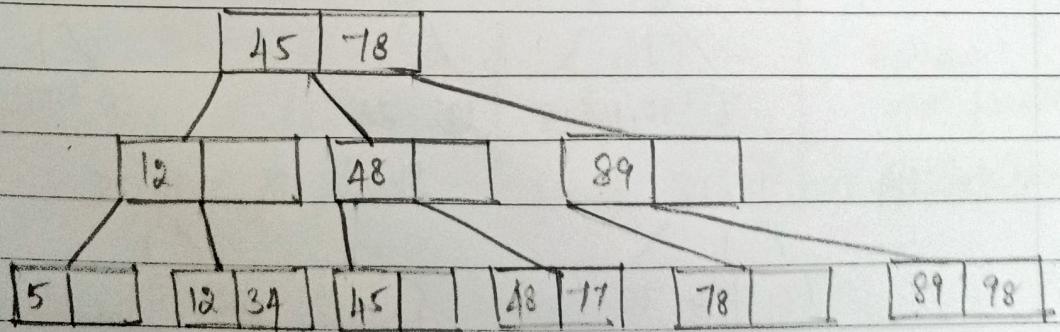
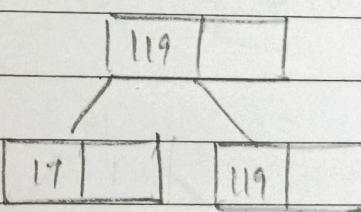
Delete 118



Delete 28



Delete 29



Delete 48, 78, 77, 98, 89

45 72

Delete 48

12 77 89

5 12 34 45 77 78 89 98

Delete 78

45 89

12 77 98

5 12 34 45 77 89 98

Delete 77

89

12 45 98

5 12 34 45 89 98

Reduce one of the key
of the parent
Underflow

Delete 98

45

12 89

5 12 34 45 89

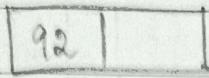
Delete 89.

12 45

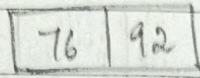
5 12 35 45

Insert 92, 76, 45, 33, 1, 94, 78, 85 Order m=3.

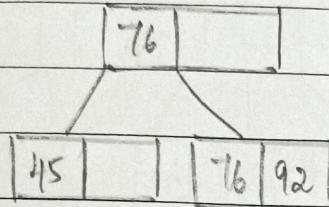
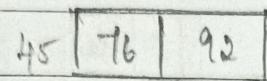
Insert 92.



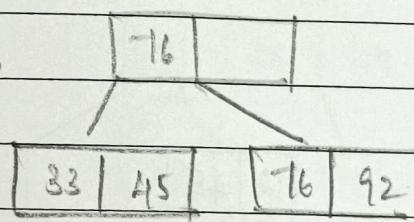
Insert 76



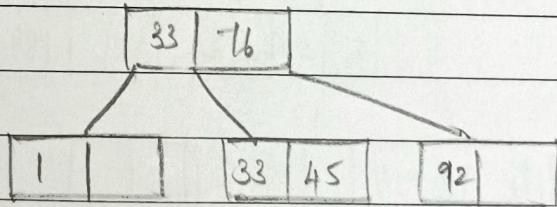
Insert 45 45



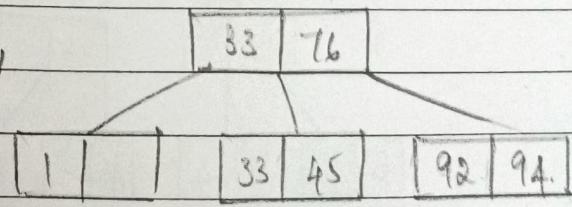
Insert 33



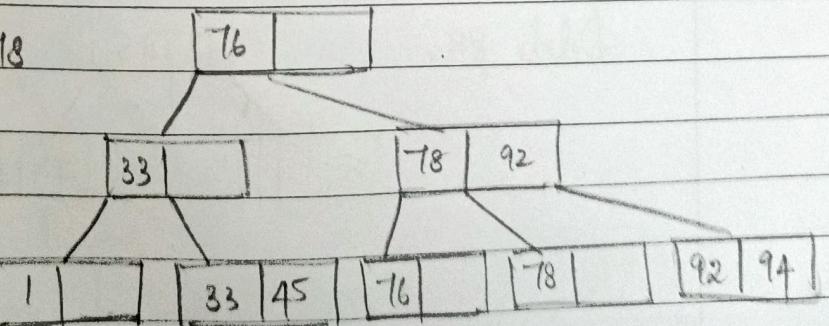
Insert 1



Insert 94



Insert 78



— / / —

Insert 85

