

# Source Coding

Joe Louis Paul. I



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# Coding Theory

- The purpose is to design an efficient and reliable data transmission methods.
- This typically involves the removal of redundancy and the correction (or detection) of errors in the transmitted data.
- It also includes the study of the properties of codes and their fitness for a specific application.

# Two Aspects of Coding Theory

- **Data compression** (or, *source coding*)
- **Error correction** (or, *channel coding*)

# Source Encoding Vs Channel Encoding

- Source encoding, attempts to compress the data from a source in order to transmit it more efficiently.
  - This practice is found every day on the Internet where the common "Zip" data compression is used to reduce the network load and make files smaller.
- The second, channel encoding, adds extra data bits to make the transmission of data more robust to disturbances present on the transmission channel.

# Need of Source Coding

- The important problem in communication is that **how the efficient representation of data produced by discrete source can be achieved?**
- The process by which this representation is achieved called as **source coding** (i.e. **representation of symbols in binary**).
- The aim of source coding is to take the source data and make it smaller.

# Source Codes

- Basically source codes try to reduce the redundancy present in the source, and represent the source with fewer bits that carry more information.

# Source Encoder

- The **device or an element which involves with source coding in a communication system.**



# Types of Source Coding

- **Fixed Length Coding (FLC)**
- **Variable Length Coding (VLC)**

# Fixed Length Coding (FLC)

- Modern codes
  - used for the interchange and processing of information
  - encode each character by one or more octets, an octet being a sequence of 8 bits.
- *A code is fixed length* if each character of the code is represented by the same number of octets.

# Example: Fixed Length Coding

- **ASCII**

- One of the most well known coded character sets is ASCII, the *American Standard Code for Information Interchange*.
- This represents **128** characters and uses all possible combinations of **7** bits.

# Table I: 8-bit Extended ASCII

65	01000001	A	Uppercase A
66	01000010	B	Uppercase B
67	01000011	C	Uppercase C
68	01000100	D	Uppercase D
69	01000101	E	Uppercase E
70	01000110	F	Uppercase F
71	01000111	G	Uppercase G
72	01001000	H	Uppercase H
73	01001001	I	Uppercase I
74	01001010	J	Uppercase J
75	01001011	K	Uppercase K
76	01001100	L	Uppercase L
77	01001101	M	Uppercase M
78	01001110	N	Uppercase N
79	01001111	O	Uppercase O
80	01010000	P	Uppercase P
81	01010001	Q	Uppercase Q
82	01010010	R	Uppercase R
83	01010011	S	Uppercase S
84	01010100	T	Uppercase T
85	01010101	U	Uppercase U
86	01010110	V	Uppercase V
87	01010111	W	Uppercase W
88	01011000	X	Uppercase X
89	01011001	Y	Uppercase Y
90	01011010	Z	Uppercase Z

# Variable Length Coding (VLC)

- In *coding theory* a **variable-length code** is a code which maps source symbols to a *variable* number of bits.

# Example: Variable Length Coding

- Morse Code

- Letters of the alphabet and numbers are encoded into sequence of marks and spaces denoted as “.” and “-” respectively.
- Example
  - Letter E denoted as “.”
  - Letter Q denoted as “--.-”

# Table II: Morse Code-A Variable Length Code

A	.-	N	-. .	1	.----	.	.-.-.-
B	-... .	O	--- .	2	..---	,	--.-.-
C	-.-. .	P	.-- .	3	...--	?	..-.-.
D	-.. .	Q	--.-	4	....-	(	-.--.
E	. .	R	.-. .	5	.....	)	-.-.-.
F	..-. .	S	... .	6	-....	-	-....-
G	--. .	T	- .	7	--...-	"	.-.-.-.
H	.... .	U	..- .	8	---..	_	.-.-.-.
I	.. .	V	...- .	9	-----	+	..-.-.
J	.----	W	.-.-	0	-----	:	---... .
K	-.-. .	X	-.-.-	/	-.-.-.	;	-.-.-.
L	.-... .	Y	-.--	+	.-.-.-.	\$	...-.-.-
M	-- .	Z	--.. .	=	-...-		

# Fixed Length and Variable Length Codes (1/3)

- A **fixed length code**, or **block code** is a code whose codewords all have the same length  $n$ .
- In this case, the number  $n$  is also called as the length of the code.
- If a code  $C$  contains codewords of varying lengths then it is called as a variable length code.



# Fixed Length and Variable Length Codes (2/3)

- Fixed length codes have advantages and disadvantages over variable length codes.
- One advantage is that **they never require a special symbol to separate the characters in the message being coded.**
- For example, consider the ***encoded ASCII message,***  
01000011010011110100010001000101

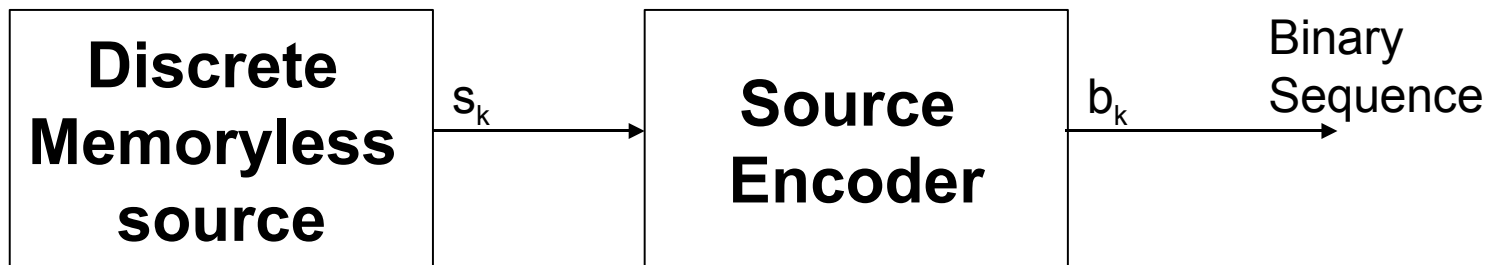
# Fixed Length and Variable Length Codes (3/3)

- Binary (ASCII) code is a fixed length code whose codewords have length 8, we know that the first 8 bit represents the first character of the original message. Similarly the second set of 8 bits represents the second character in that message.
- Perhaps, the main disadvantage of fixed length codes, such as the ASCII code, is that characters that are used frequently such as the letter e, have codes as long as characters that are used infrequently, such as the space character.
- On the other hand, variable length codes, which can encode frequently used characters using shorter codewords can save a great deal of space and time.

# Development of an Efficient Source Encoder

- The statistics of the source is required.
- **Two Fundamental Requirements**
  - The codewords produced by the source encoder are in binary form.
  - The source coding scheme should be *uniquely decodable* so that the original source sequence can be reconstructed exactly from the encoded binary sequence.

# A Typical Source Coding Scheme



- Let us assume that
  - Source sequence  $s_k$
  - Binary sequence  $b_k$
  - Length of the binary sequence  $l_k$

# Source Coding

- Consider a source that has an alphabet of  $K$  different symbols such that  $\mathbf{s}_k$  is the  $k^{\text{th}}$  symbol in the alphabet.
- Also let  $\mathbf{s}_k$  occurs with probability  $\mathbf{p}_k$  and let the binary codeword assigned to  $\mathbf{s}_k$  by the encoder have length  $l_k$  measured in bits.

# Average Codeword Length $\bar{L}$

- The average codeword length of the source encoder is defined as

$$\bar{L} = \sum_{k=0}^K p_k l_k$$

- The unit of average codeword length is **bits per symbol**.

# Coding Efficiency $\eta$

- $\bar{L}$  represents the average number of bits per source symbol used in the source encoding process.
- If  $L_{\min}$  denotes the minimum possible codeword length then, the coding efficiency of the source encoder is defined as

$$\eta = \frac{L_{\min}}{\bar{L}}$$

- The coding efficiency will be equal to one if  $\bar{L} = L_{\min}$  i.e. efficient source input.

# Shannon's Source Coding Theorem

- With  $\bar{L} \geq L_{\min}$ , then  $\eta \leq 1$ .
- The source encoder is said to be efficient when  $\eta$  approaches unity.
- **Source coding Theorem**
  - Given a discrete memoryless source of entropy  $H(S)$ , the average code-word  $\bar{L}$  length for any source encoding is bounded as

$$\bar{L} \geq H(S)$$



# Determining the Value of $L_{\min}$

- From Coding efficiency formula,

$$\bar{L} \geq L_{\min}$$

- From Shannon's Source Coding Theorem,

$$\bar{L} \geq H(S)$$

- From these above equations,

$$L_{\min} = H(S)$$

# Coding Efficiency using $H(S)$

- If  $H(S)$  denotes the minimum possible codeword length then, the coding efficiency of the source encoder is defined as

$$\eta = \frac{H(S)}{\bar{L}}$$