

# CS 731: Blockchain Technology And Applications

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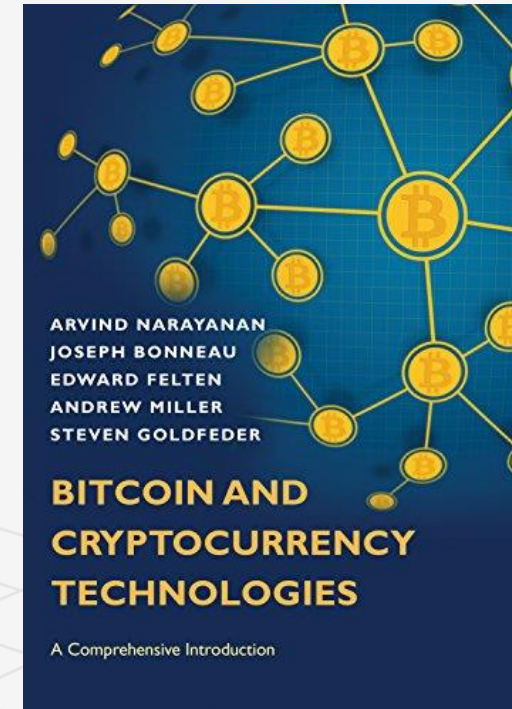
C3I Center



# Acknowledgement

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- The material of this lecture material is mostly due to Prof. Arvind Narayanan's Lecture at Princeton and his book on Bitcoin (Chapter 8 mostly)



# Alternative Mining Puzzles



# Puzzles are the core of Bitcoin

- Incentive system steers participants
- Basic features of Bitcoin's puzzle (recap)
  - The puzzle is difficult to solve, so attacks are costly
  - ... but not too hard, so honest miners are compensated
- What other features could a puzzle have?

# This lecture

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- Alternative puzzle designs  
Used in practice, and speculative
- Variety of possible goals  
ASIC resistance, pool resistance, intrinsic benefits...
- Essential security requirements

# Essential Puzzle Requirements

# Puzzle requirements

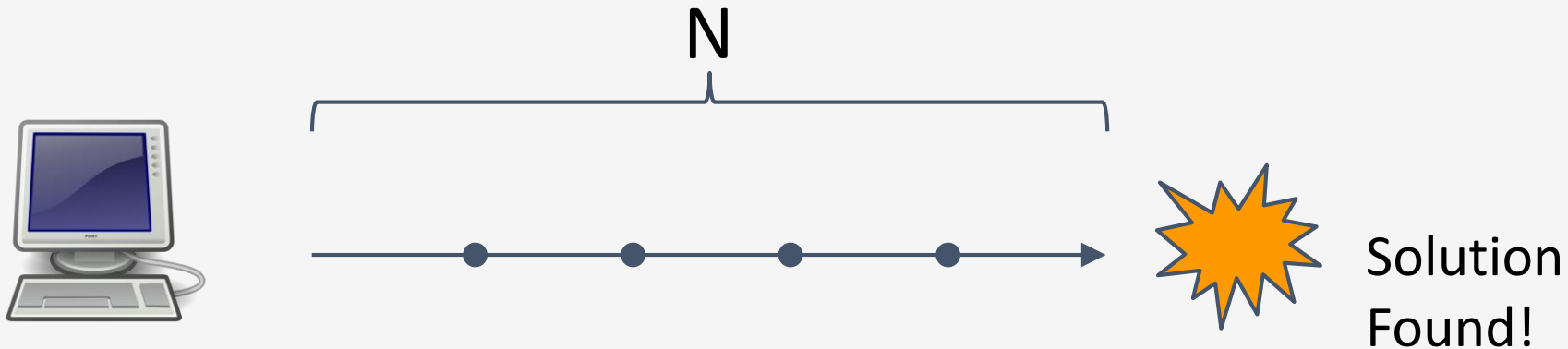
- Cheap to Verify
- Adjustable difficulty

...

- Chance of winning is proportional to hashpower
  - Large players get only proportional advantage
  - Even small players get proportional compensation

# Bad puzzle: a sequential puzzle

Consider a puzzle that takes  $N$  steps to solve  
a “Sequential” Proof of Work



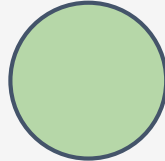
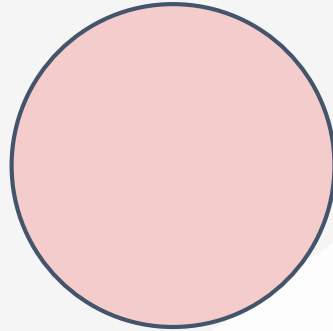
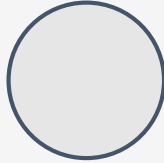


# Bad puzzle: a sequential puzzle

Problem: fastest miner **always** wins the race!



# Good puzzle → Weighted sample



This property is sometimes called “progress-free”

# ASIC Resistant Puzzles

# ASIC resistance - Why? (1 of 2)

Goal: Ordinary people with idle laptops, PCs, or even mobile phones can mine!

Lower barrier to entry

Approach: reduce the gap between custom hardware and general purpose equipment

# ASIC resistance - Why? (2 of 2)

Goal: Prevent large manufacturers from dominating the game

“Burn-in” advantage

In-house designs

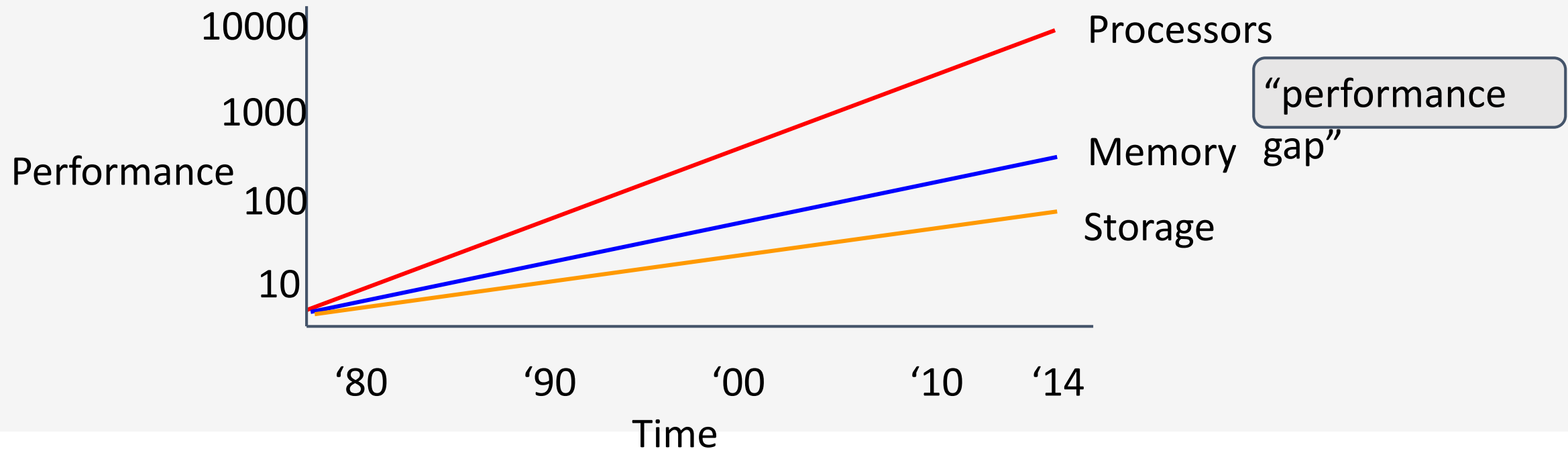
Approach: reduce the “gap” between future hardware and the custom ASICs we already have



# Memory hard puzzles

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Premise: the cost and performance of memory is more stable than for processors



# scrypt

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Colin Percival, 2009

- Memory hard hash function
  - Constant time/memory tradeoff***
- Most widely used alternative Bitcoin puzzle
- Also used elsewhere in security (Password-hashing)

1. Fill memory with random values
2. Read from the memory in random order





## script - step 2 of 2 (read)

Input:  $X$

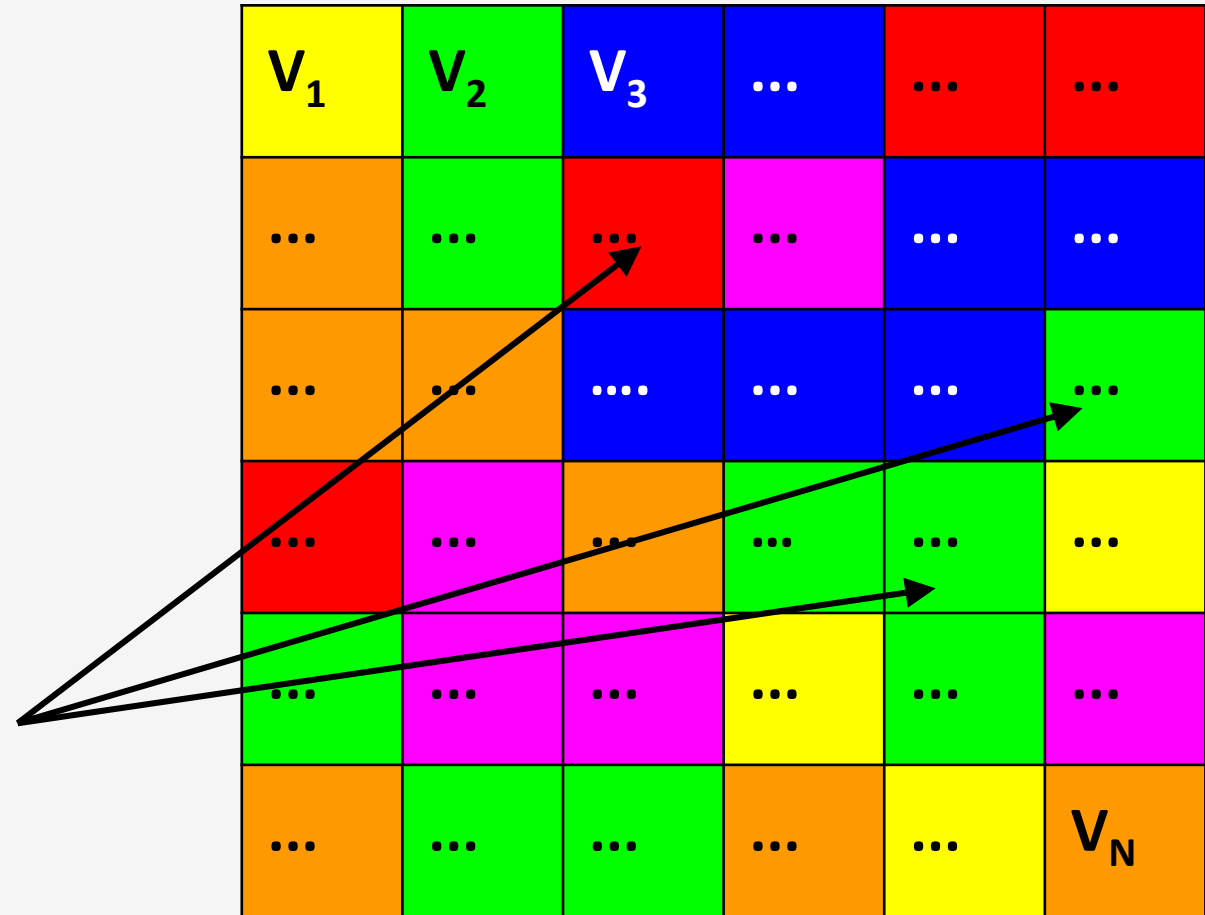
$\mathbf{A} := \mathbf{H}^{N+1}(\mathbf{X})$

For  $N$  iterations:

$\mathbf{i} := \mathbf{A} \bmod N$

$\mathbf{A} := \mathbf{H}(\mathbf{A} \text{ xor } V_{\mathbf{i}})$

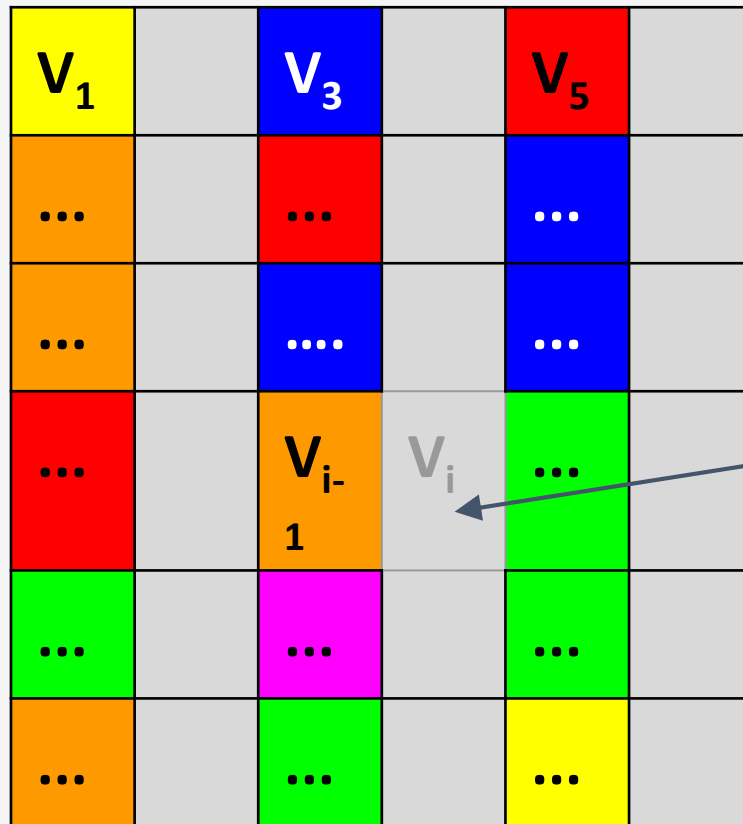
Output:  $\mathbf{A}$



# script - time/memory tradeoff

Why is this memory-hard?

Reduce memory by half, 1.5x the # steps



Need to access  $V_i$  where  $i$  is even?

Access  $V_{i-1}$

Compute  $V_i = H(V_{i-1})$

# script

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Disadvantages:

Also requires  $N$  steps,  $N$  memory to check

Is it actually ASIC resistant?

script ASICs *are* already available



<http://zeusminer.com/>

# Cuckoo hash cycles

John Tromp, 2014

Memory hard puzzle that's cheap to verify

Input:  $X$

For  $i = 1$  to  $E$ :

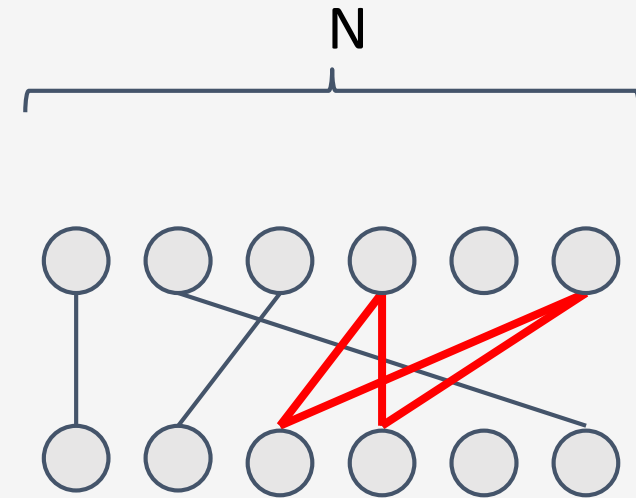
$a := H_0(X + i)$

$b := N + H_1(X + i)$

**edge** ( $a \bmod N$ ,  $b \bmod N$ )

Is there a cycle of size  $K$ ?

If so, Output:  $X$ ,  $K$  edges



# Even more approaches

- More complicated hash functions  
X11: 11 different hash functions combined
- Moving target  
Change the puzzle periodically

# Counter argument: SHA2 is fine

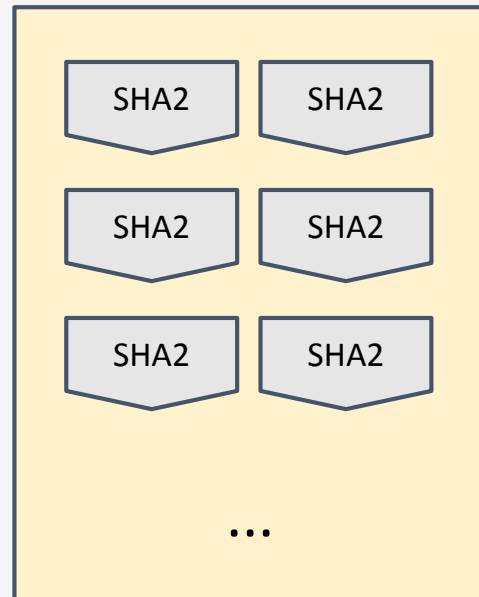
Bitcoin Mining ASICs aren't changing much

Big ASICs only marginally more performant than small ones

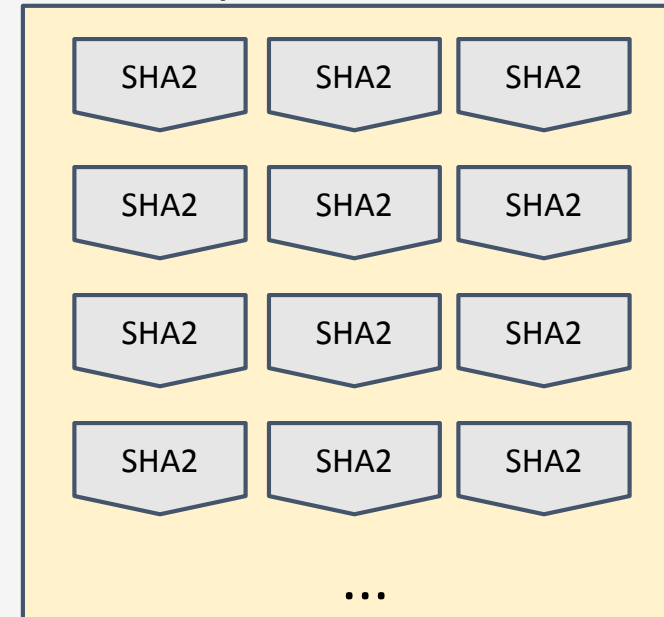
Ordinary SHA2 Circuit



Affordable ASIC



Expensive ASIC



# Proof-of-useful-work

# Recovering wasted work

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Recall:

between 150 MW - 900 MW power consumed (as of mid-2014)

Natural question:

Can we recycle this and do something useful?



# Candidates - needle in a haystack

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- Natural choices:
  - Protein folding (find a low energy configuration)
  - Search for aliens (find an anomalous region of a signal)

These have been successful @Home problems
- Challenges:
  - Randomly chosen instances must be hard

Who chooses the problem?

# Primecoin

Sunny King, 2013



Puzzle based on finding large prime numbers

Cunningham chain:

$p_1, p_2, \dots, p_n$  where  $p_{i+1} = 2 \cdot p_i + 1$

Each  $p_i$  is a large (probable) prime

$p_1$  is divisible by  $H(\text{prev} || \text{mrkl\_root} || \text{nonce})$

# Primecoin

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- Many of the largest known Cunningham chains have come from Primecoin miners
- Hard problem? Studied by others (e.g., PrimeGrid)
- Usefulness? Maybe - at least one known use

# Recovering wasted hardware

Estimate: more than \$100M spent on customized Bitcoin mining hardware

This hardware investment is otherwise useless

Idea: a puzzle where hardware investment is useful, even if the work is wasted?

# Permacoin - Mining with storage

Miller et al.,  
2014

## Bitcoin



## Permacoin



Side effect:  
Massively distributed, replicated storage system

# Permacoin

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Assume we have a large file **F** to store

For simplicity: **F** is chosen globally, at the beginning, by a trusted dealer

Each user stores a random subset of the file

# Storage-based puzzle

1. Build a Merkle tree, where each leaf is a segment of the file

2. Generate a public signing key  $pk$ , which determines a random subset of file segments

$F_1$   $F_2$   $F_4$   $F_5$

3. Each mining attempt:

$F_2$   $F_4$

a) Select a random nonce

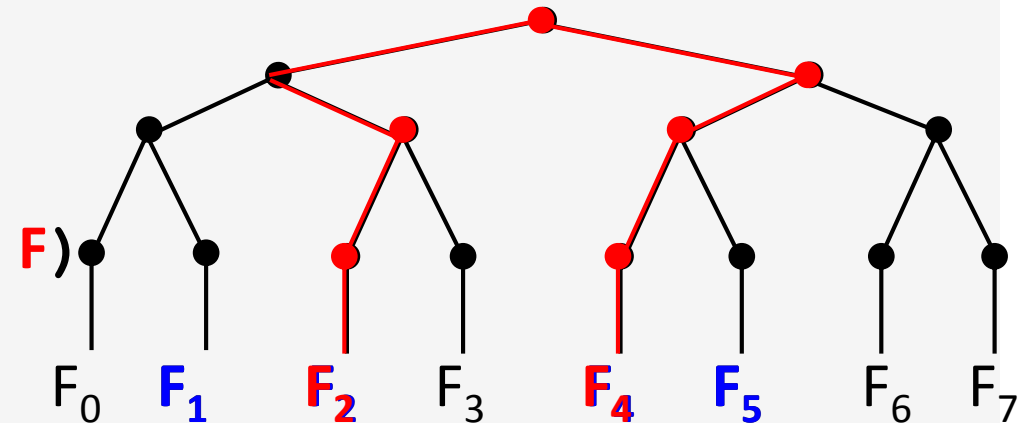
b)  $h1 := H(\text{prev} || \text{mrkl\_root} || PK || \text{nonce})$

c)  $h1$  selects  $k$  segments from subset

d)  $h2 :=$

$H(\text{prev} || \text{mrkl\_root} || PK || \text{nonce} || F)$

e) Winner if  $h2 < \text{TARGET}$



# Reducing Bitcoin's "honesty" cost

"Honest" miners validate every transaction

Validation requires the UTXO database ~200MB

Maintaining the UTXO database doesn't pay

Idea: use Permacoin to reward UTXO storage



# Summary

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- Useful proof-of-work is a natural goal  
(while maintaining security requirements)
- The benefit must be a pure public good
- Viable approaches include storage, prime-finding, others may be possible
- Realized benefit so far has been limited

# Nonoutsourceable Puzzles

# Large mining pools are a threat

- Bitcoin's core value is decentralization
- If power is consolidated in a few large pools, the operators are targets for coercion/hacking
- Position: large pools should be discouraged!  
Analogy to voting: It's illegal (in US) to sell your vote

# Hacking, Distributed

## It's Time For a Hard Bitcoin Fork

Ittay Eyal, and Emin Gün Sirer

Friday June 13, 2014 at 02:05 PM

A Bitcoin mining pool, called GHash and operated by an anonymous entity called CEX.io, just reached 51% of total network mining power today. Bitcoin is no longer decentralized. GHash can control Bitcoin transactions.

### Is This Really Armageddon?

Yes, it is. GHash is in a position to exercise complete control over which



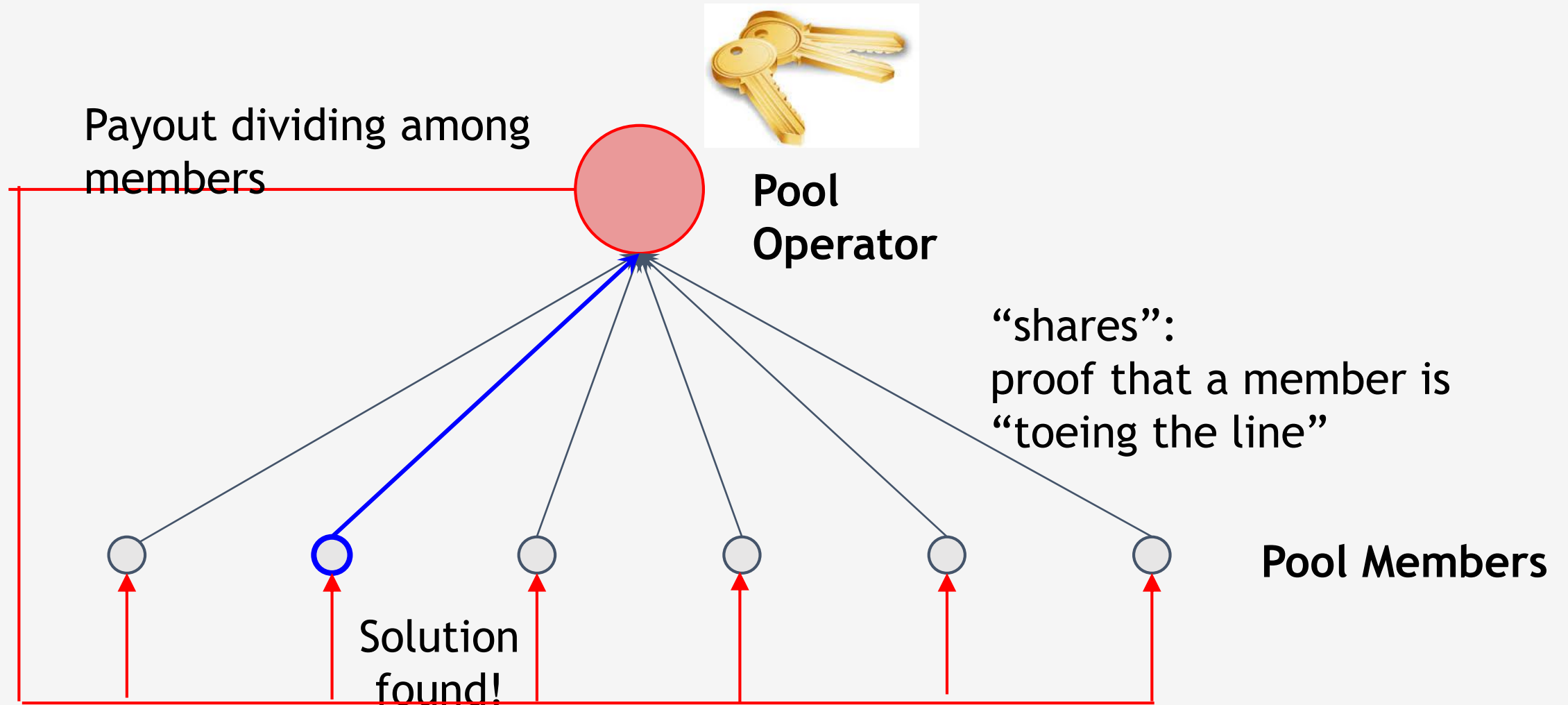
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Observation:

Pool participants don't trust each other

Pools only work because the “shares” protocol lets members ***prove*** cooperation

# Standard Bitcoin mining pool



# The Vigilante Attack

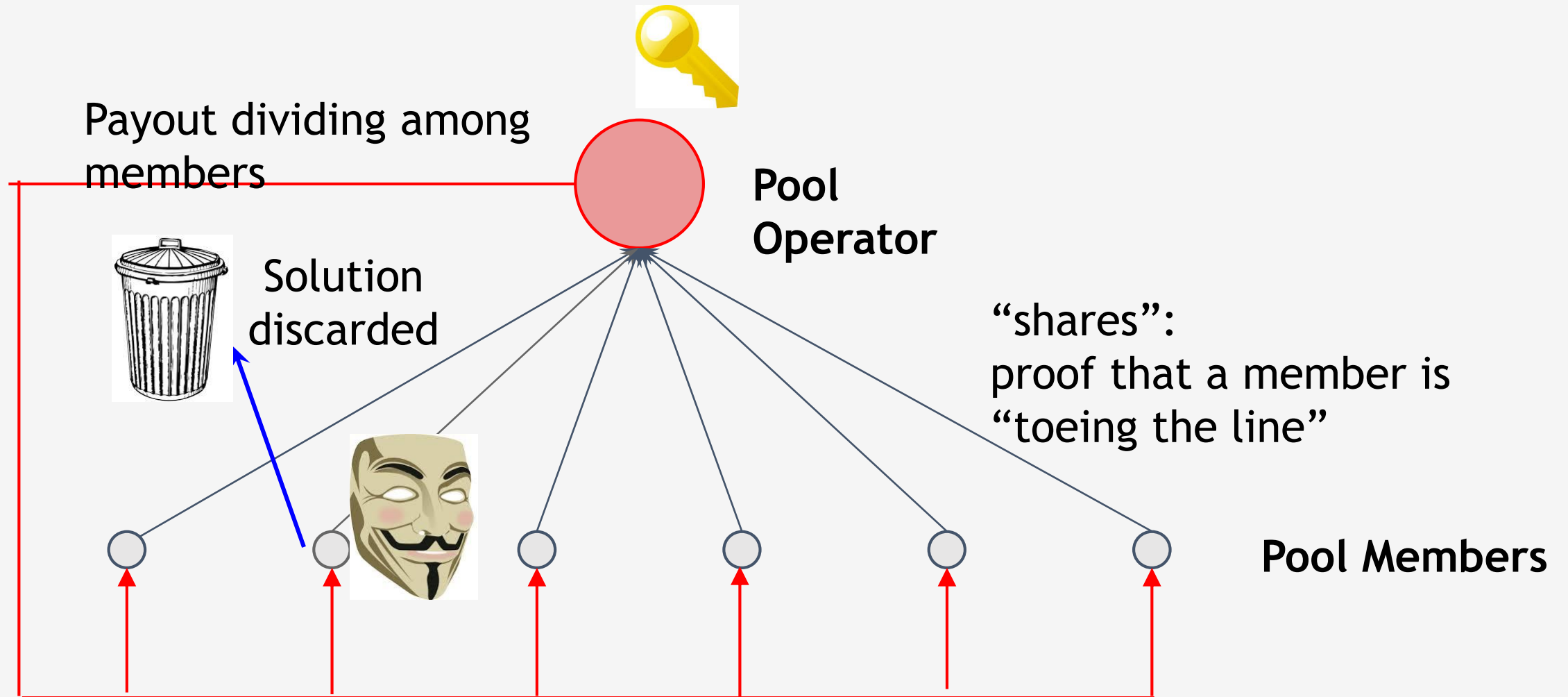
Suppose a Vigilante is angry with a large pool

He submits “shares” like normal....

... but if he finds a real solution, discards it

Pool output is reduced, Vigilante loses a little

# The Vigilante Attack





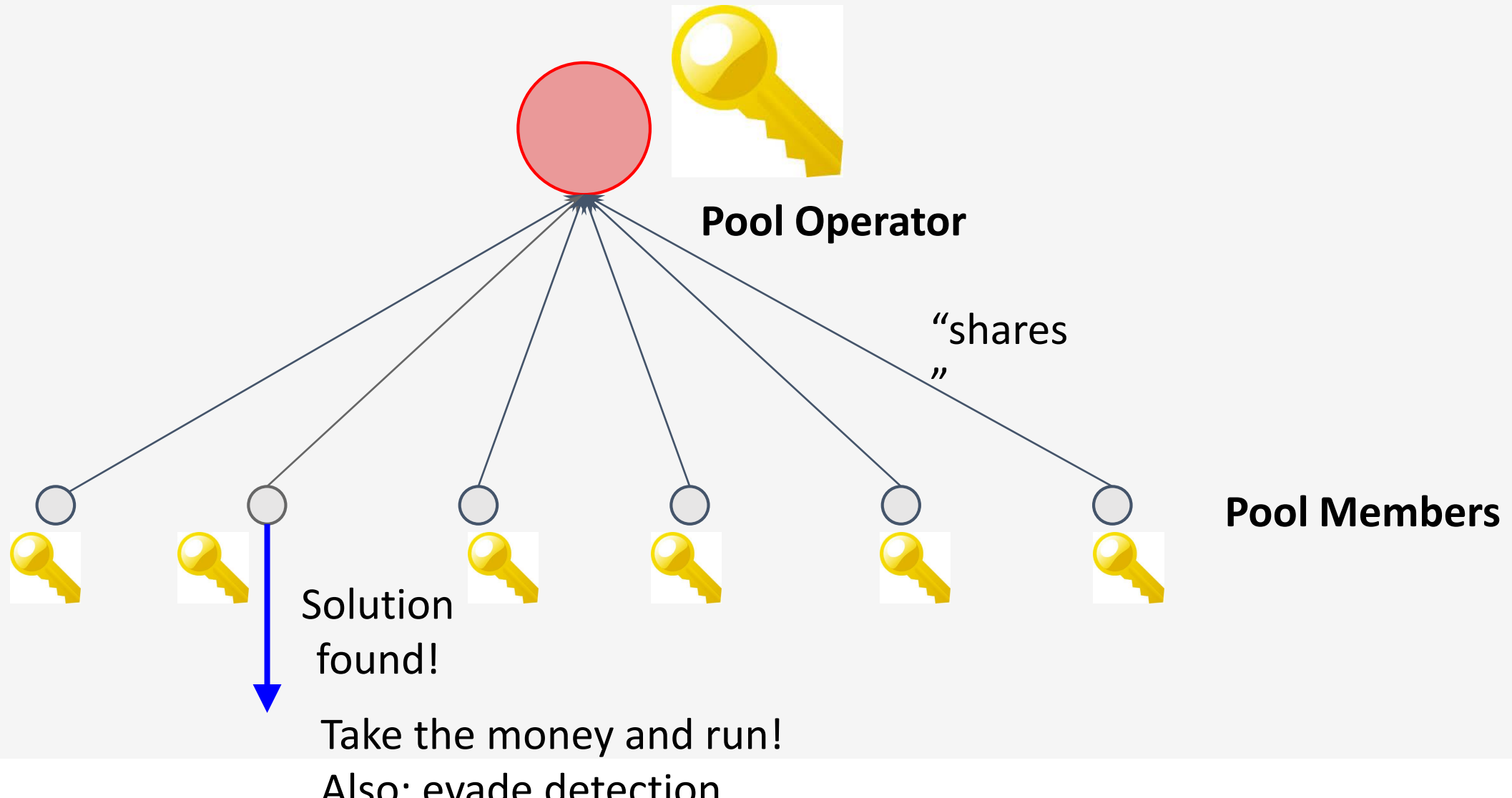
# Encouraging the Vigilante

Whoever **FINDS** a solution spends the reward

Approach:

- searching for a solution requires **SIGNING**, not just hashing. (Knowledge of a private key)
- Private key can be used to spend the reward

# Encouraging the Vigilante



# Nonoutsourceable puzzle

Solution:

**(prev, mrkl\_root, nonce, PK, s1, s2)**

such that:

**$H(\text{prev} || \text{PK} || \text{nonce} || \text{s1}) < \text{TARGET}$**   
**VerifySig(PK, s1, prev || nonce)**  
**VerifySig(PK, s2, prev || mrkl\_root)**

Signature needed to find  
solution

Public Key

s1, s2

Second signature spends  
reward

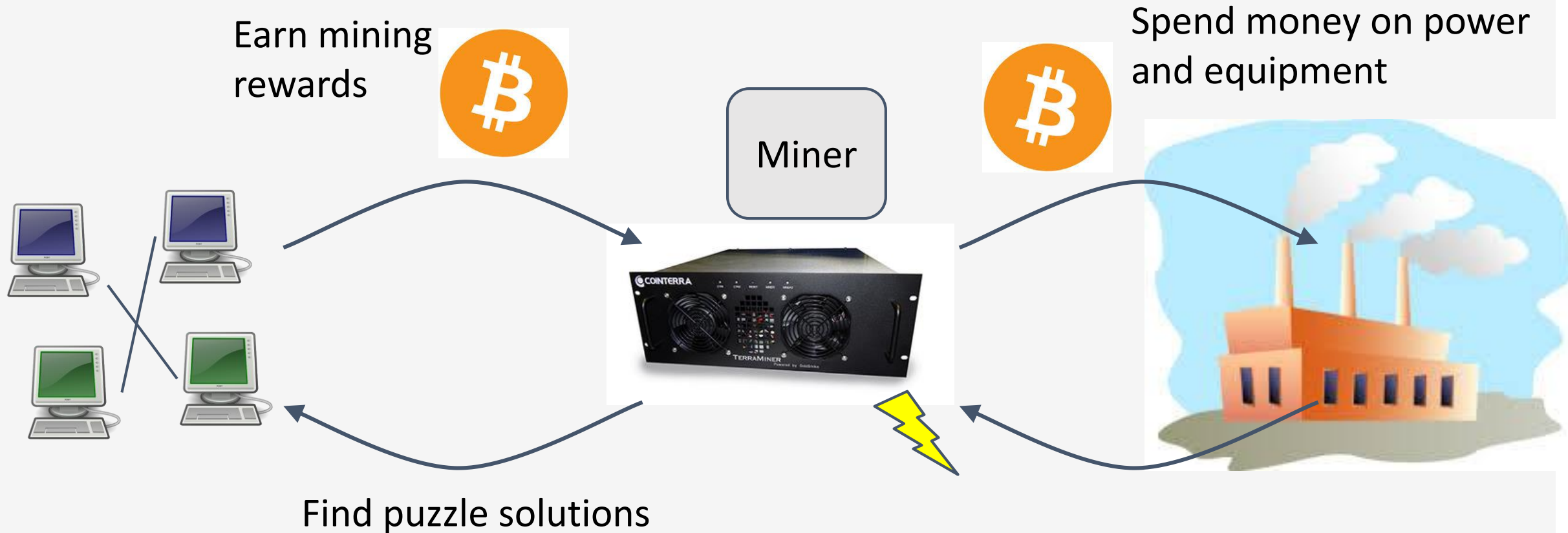
# Nonoutsourceable puzzle concerns

- This puzzle discourages ALL pools including harmless decentralized P2Pools
- Other forms of outsourcing might drive pool members to hosted mining

# Proof-of-Stake “Virtual Mining”

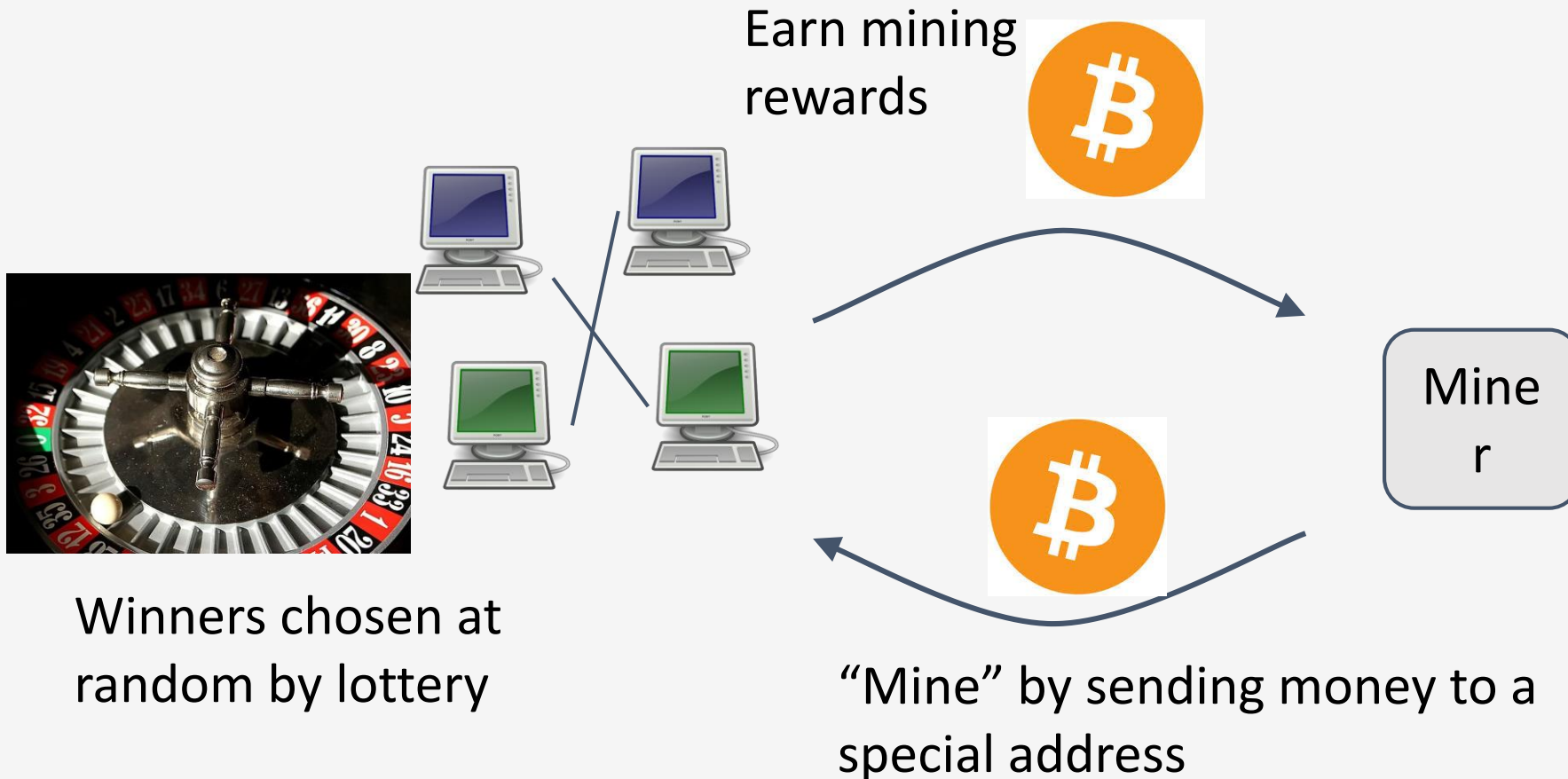
# Mining has an unnecessary step

Proof-of-Work Mining:



# Mining has an unnecessary step

Virtual Mining:



# Potential benefits

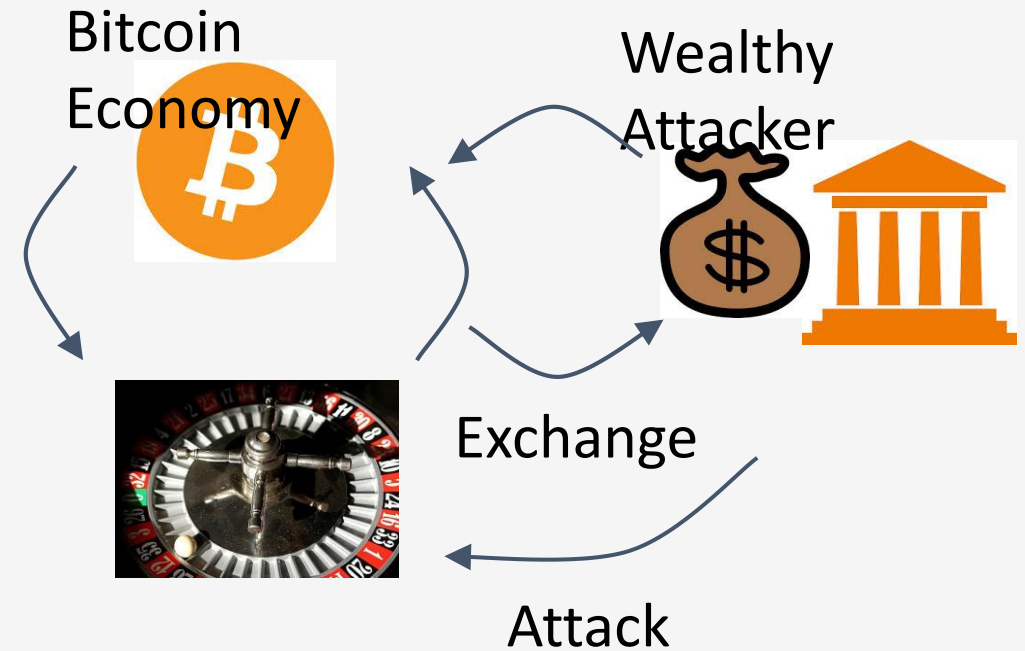
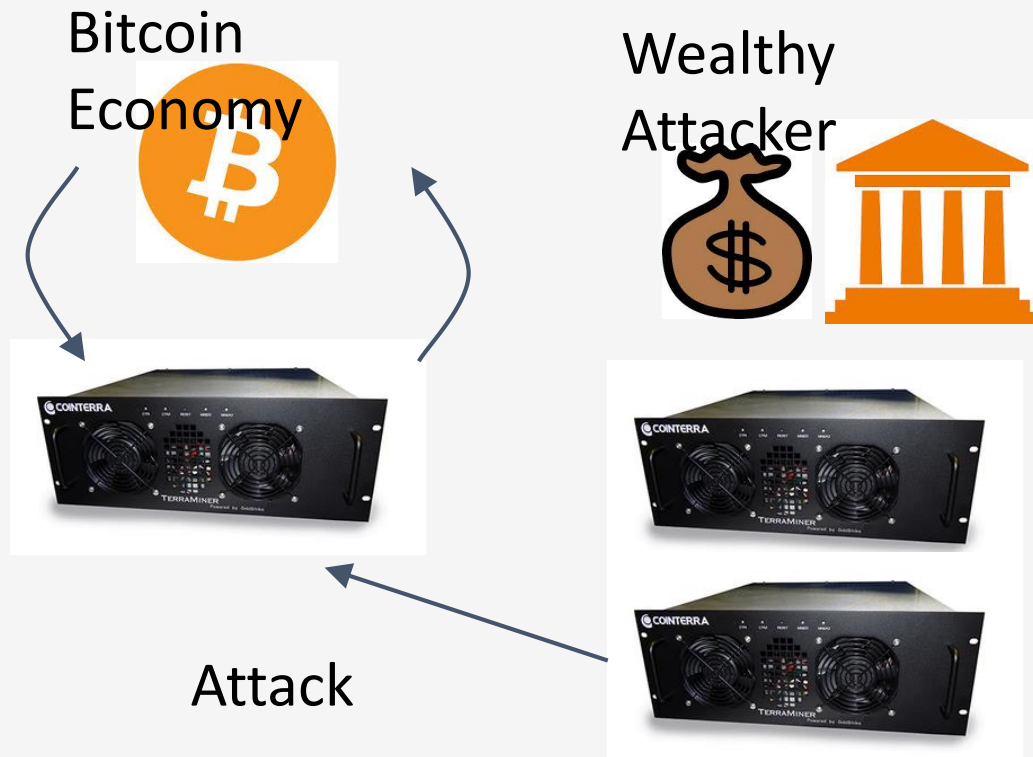
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- Lower overall costs
  - No harm to the environment
  - Savings distributed to all coin holders
- Stakeholder incentives - good stewards?
- No ASIC advantage
- 51% attack is even harder



# 51% attack prevention

The Bitcoin economy is smaller than the world  
Wealth *outside* Bitcoin has to move *inside*



# Variations of Virtual Mining

- Proof-of-Stake: “Stake” of a coin grows over time as long as the coin is unused
- Proof-of-Burn: mining with a coin destroys it
- Proof-of-Deposit: can reclaim a coin after some time
- Proof-of-Activity: any coin might be win (if online)

# Open Questions with Virtual Mining

Is there any security that can only be gained by consuming “real” resources?

- If so, then “waste” is the cost of security
- If not, then PoW mining may go extinct

# Conclusion

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- Many possible design goals
  - Prevent ASIC miners from dominating
  - Prevent large pools from dominating
  - Intrinsic usefulness
    - Eliminate the need for mining hardware at all
- Best tradeoff is unclear for now
- Outlook: alternatives will coexist for the near future