Gauss Elimination Method: Pivoting

- First Step: $R'_{i}=R_{i}-(a_{i1}/a_{11})xR_{1}$ i=2 to n
- Second step: $R''_{i}=R'_{i}-(a'_{i2}/a'_{22})xR'_{2}$ i= 3 to n; and so on..
- The multiplying factor at each step is equal to the corresponding element of the row divided by the pivot element
- If we want round-off errors to be attenuated, this factor should be small in magnitude.
- Since order of equations is immaterial, we can interchange rows without affecting the solution
- Partial Pivoting: Scan the elements below the diagonal in that particular column, find the largest magnitude, and interchange rows so that the pivot element becomes the largest.

Pivoting

Complete Pivoting: Scan the elements in the entire submatrix, find the largest magnitude, and interchange rows and columns, so that the pivot element becomes the largest.

$$\begin{bmatrix} a_{11} & a_{12} & a_{13} & a_{14} & \cdots & a_{1k} & a_{1k+1} & \cdots & a_{1n} \\ 0 & a_{22} & a_{23} & a_{24} & \cdots & a_{2k} & a_{2k+1} & \cdots & a_{2n} \\ 0 & 0 & a_{33} & a_{34} & \cdots & a_{3k} & a_{3k+1} & \cdots & a_{3n} \\ 0 & 0 & 0 & a_{44} & \cdots & a_{4k} & a_{4k+1} & \cdots & a_{4n} \\ \cdots & \cdots \\ 0 & 0 & 0 & 0 & \cdots & a_{kk} & a_{k+1k+1} & \cdots & a_{kn} \\ 0 & 0 & 0 & 0 & \cdots & a_{k+1k} & a_{k+1k+1} & \cdots & a_{k+1n} \\ \cdots & \cdots & \cdots & \cdots & \cdots & \cdots \\ 0 & 0 & 0 & \cdots & a_{nk} & a_{nk+1} & \cdots & a_{nn} \end{bmatrix}$$

Partial Complete

Pivoting

- Parial Pivoting does not need any additional bookkeeping.
 Complete pivoting, due to an exchange of columns, needs to keep track of this exchange.
- If we exchange columns k and l: after the solution, the values of x_k and x_l have to be interchanged.

Examples: Partial Pivoting:

$$\begin{bmatrix} 2 & 4 & 1 \\ 1 & 2 & 3 \\ 3 & 1 & 5 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} 13 \\ 14 \\ 20 \end{bmatrix}; \begin{bmatrix} 3 & 1 & 5 \\ 1 & 2 & 3 \\ 2 & 4 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} 20 \\ 14 \\ 13 \end{bmatrix}; \begin{bmatrix} 3 & 1 & 5 \\ 0 & 5/3 & 4/3 \\ 0 & 10/3 & -7/3 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} 20 \\ 22/3 \\ -1/3 \end{bmatrix}$$

$$\begin{bmatrix} 3 & 1 & 5 \\ 0 & 10/3 & -7/3 \\ 0 & 5/3 & 4/3 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} 20 \\ -1/3 \\ 22/3 \end{bmatrix}; \begin{bmatrix} 3 & 1 & 5 \\ 0 & 10/3 & -7/3 \\ 0 & 0 & 5/2 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} 20 \\ -1/3 \\ 15/2 \end{bmatrix};$$
 Soln:

Example: Complete (or Full) Pivoting

$$\begin{bmatrix} 2 & 4 & 1 \\ 1 & 2 & 3 \\ 3 & 1 & 5 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} 13 \\ 14 \\ 20 \end{bmatrix}; \quad \begin{bmatrix} 3 & 1 & 5 \\ 1 & 2 & 3 \\ 2 & 4 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} 20 \\ 14 \\ 13 \end{bmatrix}; \quad \begin{bmatrix} 5 & 1 & 3 \\ 3 & 2 & 1 \\ 1 & 4 & 2 \end{bmatrix} \begin{bmatrix} x_3 \\ x_2 \\ x_1 \end{bmatrix} = \begin{bmatrix} 20 \\ 14 \\ 13 \end{bmatrix};$$

• Note the interchange, of the variables x_1 and x_3 . While computing, we deal with only [A] and {b}. So we make a note of this exchange and do it after the solution is obtained.

$$\begin{bmatrix} 5 & 1 & 3 \\ 0 & 1.4 & -0.8 \\ 0 & 3.8 & 1.4 \end{bmatrix} \{x\} = \begin{cases} 20 \\ 2 \\ 9 \end{cases}; \qquad \begin{bmatrix} 5 & 1 & 3 \\ 0 & 3.8 & 1.4 \\ 0 & 1.4 & -0.8 \end{bmatrix} \{x\} = \begin{cases} 20 \\ 9 \\ 2 \end{cases};$$

$$\begin{bmatrix} 5 & 1 & 3 \\ 0 & 3.8 & 1.4 \\ 0 & 0 & -25/19 \end{bmatrix} \{x\} = \begin{cases} 20 \\ 9 \\ -25/19 \end{cases}; Solution: 3,2,1. Interchange: 1,2,3$$

Other considerations

- Round-off errors are typically large if the elements of the [A] matrix are of very different magnitudes.
- Scaling of variable and equilibration of equation are used to make the coefficient matrix elements roughly of the same order
- For example, $\begin{bmatrix} 0.02 & 4 & 1 \\ 0.01 & 2 & 3 \\ 0.03 & 1 & 5 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} 13 \\ 14 \\ 20 \end{bmatrix}$ may be because x_1 is in cm and others in m
- Similarly, $\begin{bmatrix} 2 & 4 & 1 \\ 1 & 2 & 3 \\ 1/20 & 1/60 & 1/12 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} 13 \\ 14 \\ 1/3 \end{bmatrix}$; may be because the and others are in minutes
- Equilibrating by multiplying the third equation by 60.

Gauss Elimination Method, LU decomposition

- GE reduces the matrix [A] to upper triangular form
- If there are several RHS vectors, an "augmented" [A]
 matrix could be used to obtain the solution for all of
 these through the same algorithm
- Sometimes, an RHS vector depends on the solution for the previous RHS vector.
- The LU decomposition method is then more efficient, since solution of two triangular systems involves much less computational effort

$$[A] = [L][U] \Rightarrow [L]\{y\} = \{b\}$$
 where, $\{y\} = [U]\{x\}$

 The effort in decomposition is needed only once (The assumption, of course, is that [A] is not changing!)

LU decomposition: Alternative ways

$$\begin{bmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ a_{21} & a_{22} & \dots & a_{2n} \\ \dots & \dots & \dots & \dots \end{bmatrix} = \begin{bmatrix} l_{11} & 0 & \dots & 0 \\ l_{21} & l_{22} & \dots & 0 \\ \dots & \dots & \dots & \dots \\ \vdots & \dots & \dots & \dots \\ l_{n1} & l_{n2} & \dots & l_{nn} \end{bmatrix} \begin{bmatrix} u_{11} & u_{12} & \dots & u_{1n} \\ 0 & u_{22} & \dots & u_{2n} \\ \dots & \dots & \dots & \dots \\ \vdots & \dots & \dots & \dots \\ 0 & 0 & \dots & u_{nn} \end{bmatrix}$$

- n^2 equations, n^2+n "unknowns". The way the "floating" n values are fixed: different algorithms
- Doolittle: All diagonal elements of [L] are 1
- Crout: All diagonal elements of [U] are 1
- Cholesky (some call it Choleski): Works only for symmetric positive definite [A]: n(n+1)/2 eqns. [U]=[L]^T

LU decomposition: Doolittle

$$\begin{bmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ a_{21} & a_{22} & \dots & a_{2n} \\ \dots & \dots & \dots & \dots \\ a_{n1} & a_{n2} & \dots & a_{nn} \end{bmatrix} = \begin{bmatrix} 1 & 0 & \dots & 0 \\ l_{21} & 1 & \dots & 0 \\ \dots & \dots & \dots & \dots \\ \vdots & \dots & \dots & \dots \\ l_{n1} & l_{n2} & \dots & 1 \end{bmatrix} \begin{bmatrix} u_{11} & u_{12} & \dots & u_{1n} \\ 0 & u_{22} & \dots & u_{2n} \\ \dots & \dots & \dots & \dots \\ \vdots & \dots & \dots & \dots \\ 0 & 0 & \dots & u_{nn} \end{bmatrix}$$

$$= \begin{bmatrix} u_{11} & u_{12} & \dots & u_{1n} \\ l_{21}u_{11} & l_{21}u_{12} + u_{22} & \dots & u_{2n} \\ \dots & \dots & \dots & \dots \\ \vdots & \dots & \dots & \dots \\ \vdots & \dots & \dots & \dots \\ \vdots & \dots & \dots & \dots \\ l_{n1}u_{11} & l_{n1}u_{12} + l_{n2}u_{22} & \dots & l_{n1}u_{1n} + l_{n2}u_{2n} + \dots + u_{nn} \end{bmatrix}$$

- n² equations, n² "unknowns": (n²-n)/2 in [L] and (n²+n)/2 in [U]
- Sequential computations: $u_{11}=a_{11}$, $u_{12}=a_{12}$,..., $u_{1n}=a_{1n}$; $l_{21}=a_{21}/u_{11}$, $l_{31}=a_{31}/u_{11}$,..., $l_{n1}=a_{n1}/u_{11}$; $u_{22}=(a_{22}-l_{21}xu_{12})$, $u_{23}=(a_{23}-l_{21}xu_{13})$, ..., $u_{2n}=(a_{2n}-l_{21}xu_{1n})$; $l_{32}=(a_{32}-l_{31}xu_{12})/u_{22}$,...

LU decomposition: Crout

$$\begin{bmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ a_{21} & a_{22} & \dots & a_{2n} \\ \dots & \dots & \dots & \dots \end{bmatrix} = \begin{bmatrix} l_{11} & 0 & \dots & 0 \\ l_{21} & l_{22} & \dots & 0 \\ \dots & \dots & \dots & \dots \\ l_{n1} & l_{n2} & \dots & l_{nn} \end{bmatrix} \begin{bmatrix} 1 & u_{12} & \dots & u_{1n} \\ 0 & 1 & \dots & u_{2n} \\ \dots & \dots & \dots & \dots \\ 0 & 0 & \dots & 1 \end{bmatrix}$$

$$= \begin{bmatrix} l_{11} & l_{11}u_{12} & \dots & l_{1n}u_{1n} \\ l_{21} & l_{21}u_{12} + l_{22} & \dots & l_{21}u_{1n} + l_{22}u_{2n} \\ \dots & \dots & \dots & \dots \\ \vdots & \dots & \dots & \dots \\ l_{n1} & l_{n1}u_{12} + l_{n2} & \dots & l_{n1}u_{1n} + l_{n2}u_{2n} + \dots + l_{nn} \end{bmatrix}$$

- n² equations, n² "unknowns": (n²+n)/2 in [L] and (n²-n)/2 in [U]
- Sequential computations: $I_{11}=a_{11}$, $I_{21}=a_{21}$,..., $I_{n1}=a_{n1}$; $u_{12}=a_{12}/I_{11}$, $u_{13}=a_{13}/I_{11}$,..., $u_{1n}=a_{1n}/I_{11}$; $I_{22}=(a_{22}-I_{21}xu_{12})$, $I_{32}=(a_{32}-I_{31}xu_{12})$, ..., $I_{n2}=(a_{n2}-I_{n1}xu_{12})$; $u_{23}=(a_{23}-I_{21}xu_{13})/I_{22}$,...

LU decomposition: Cholesky

- n(n+1)/2 equations, n(n+1)/2 "unknowns" in [L]
- Sequential computations: $I_{11}=Va_{11}$, $I_{21}=a_{21}/I_{11}$,..., $I_{n1}=a_{n1}/I_{11}$; $I_{22}=V(a_{22}-I_{21}^2)$, $I_{32}=(a_{32}-I_{31}xI_{21})/I_{22}$, ..., $I_{n2}=(a_{n2}-I_{n1}xI_{21})/I_{22}$;

LU decomposition: Examples

$$\begin{bmatrix} 9 & 3 & -2 \\ 3 & 6 & 1 \\ -2 & 1 & 9 \end{bmatrix} \begin{Bmatrix} x_1 \\ x_2 \\ x_3 \end{Bmatrix} = \begin{Bmatrix} 10 \\ 10 \\ 8 \end{Bmatrix}$$

Doolittle:

$$\begin{bmatrix} 9 & 3 & -2 \\ 3 & 6 & 1 \\ -2 & 1 & 9 \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 \\ l_{21} & 1 & 0 \\ l_{31} & l_{32} & 1 \end{bmatrix} \begin{bmatrix} u_{11} & u_{12} & u_{13} \\ 0 & u_{22} & u_{23} \\ 0 & 0 & u_{33} \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 \\ 1/3 & 1 & 0 \\ -2/9 & 1/3 & 1 \end{bmatrix} \begin{bmatrix} 9 & 3 & -2 \\ 0 & 5 & 5/3 \\ 0 & 0 & 8 \end{bmatrix}$$

$$\begin{bmatrix} 1 & 0 & 0 \\ 1/3 & 1 & 0 \\ -2/9 & 1/3 & 1 \end{bmatrix} \begin{bmatrix} y_1 \\ y_2 \\ y_3 \end{bmatrix} = \begin{bmatrix} 10 \\ 10 \\ 8 \end{bmatrix} \qquad \begin{bmatrix} 9 & 3 & -2 \\ 0 & 5 & 5/3 \\ 0 & 0 & 8 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} 10 \\ 20/3 \\ 8 \end{bmatrix}$$

LU decomposition: Examples

$$\begin{bmatrix} 9 & 3 & -2 \\ 3 & 6 & 1 \\ -2 & 1 & 9 \end{bmatrix} \begin{Bmatrix} x_1 \\ x_2 \\ x_3 \end{Bmatrix} = \begin{Bmatrix} 10 \\ 10 \\ 8 \end{Bmatrix}$$

Crout:

$$\begin{bmatrix} 9 & 3 & -2 \\ 3 & 6 & 1 \\ -2 & 1 & 9 \end{bmatrix} = \begin{bmatrix} l_{11} & 0 & 0 \\ l_{21} & l_{22} & 0 \\ l_{31} & l_{32} & l_{33} \end{bmatrix} \begin{bmatrix} 1 & u_{12} & u_{13} \\ 0 & 1 & u_{23} \\ 0 & 0 & 1 \end{bmatrix} = \begin{bmatrix} 9 & 0 & 0 \\ 3 & 5 & 0 \\ -2 & 5/3 & 8 \end{bmatrix} \begin{bmatrix} 1 & 1/3 & -2/9 \\ 0 & 1 & 1/3 \\ 0 & 0 & 1 \end{bmatrix}$$

$$\begin{bmatrix} 9 & 0 & 0 \\ 3 & 5 & 0 \\ -2 & 5/3 & 8 \end{bmatrix} \begin{bmatrix} y_1 \\ y_2 \\ y_3 \end{bmatrix} = \begin{bmatrix} 10 \\ 10 \\ 8 \end{bmatrix} \qquad \begin{bmatrix} 1 & 1/3 & -2/9 \\ 0 & 1 & 1/3 \\ 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} 10/9 \\ 4/3 \\ 1 \end{bmatrix}$$

LU decomposition: Examples

$$\begin{bmatrix} 9 & 3 & -2 \\ 3 & 6 & 1 \\ -2 & 1 & 9 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} 10 \\ 10 \\ 8 \end{bmatrix}$$

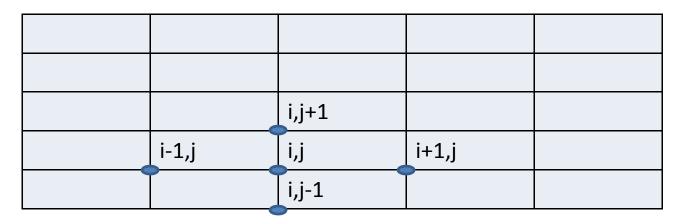
Cholesky:

$$\begin{bmatrix} 9 & 3 & -2 \\ 3 & 6 & 1 \\ -2 & 1 & 9 \end{bmatrix} = \begin{bmatrix} l_{11} & 0 & 0 \\ l_{21} & l_{22} & 0 \\ l_{31} & l_{32} & l_{33} \end{bmatrix} \begin{bmatrix} l_{11} & l_{21} & l_{31} \\ 0 & l_{22} & l_{32} \\ 0 & 0 & l_{33} \end{bmatrix} = \begin{bmatrix} 3 & 0 & 0 \\ 1 & \sqrt{5} & 0 \\ -2/3 & \sqrt{5}/3 & \sqrt{8} \end{bmatrix} \begin{bmatrix} 3 & 1 & -2/3 \\ 0 & \sqrt{5} & \sqrt{5}/3 \\ 0 & 0 & \sqrt{8} \end{bmatrix}$$

$$\begin{bmatrix} 3 & 0 & 0 \\ 1 & \sqrt{5} & 0 \\ -2/3 & \sqrt{5}/3 & \sqrt{8} \end{bmatrix} \begin{bmatrix} y_1 \\ y_2 \\ y_3 \end{bmatrix} = \begin{bmatrix} 10 \\ 10 \\ 8 \end{bmatrix} \begin{bmatrix} 3 & 1 & -2/3 \\ 0 & \sqrt{5} & \sqrt{5}/3 \\ 0 & 0 & \sqrt{8} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} 10/3 \\ 4\sqrt{5}/3 \\ \sqrt{8} \end{bmatrix}$$

Banded and Sparse matrices

Several of the matrices in engineering applications are **sparse**, i.e., have very few non-zero elements. For example, solution of differential equations:



$$\frac{\partial T}{\partial t} = \frac{T_{i,j}^{t+\Delta t} - T_{i,j}^{t}}{\Delta t}; \frac{\partial^{2} T}{\partial x^{2}} = \frac{\frac{T_{i+1,j} - T_{i,j}}{\Delta x} - \frac{T_{i,j} - T_{i-1,j}}{\Delta x}}{\Delta x} = \frac{T_{i+1,j} - 2T_{i,j} + T_{i-1,j}}{\Delta x^{2}}$$

If the non-zero elements occur within a "narrow" band around the diagonal, it is called **banded** matrix. The fact that most of the elements are zero can be used to save computational effort.

Banded Matrix

$$\begin{bmatrix}
\times & \times & . & 0 & 0 \\
\times & \times & . & . & 0 \\
. & . & . & . & . \\
0 & . & . & \times & \times \\
0 & 0 & . & \times & \times
\end{bmatrix}$$

- Band width: If main diagonal, adjacent r row elements, and adjacent c column elements **maybe** non-zero, the band-width is r+c+1. In GE method, since we know that elements beyond the band are already zero, we do not need to perform the computations to make them zero.
- Simplest case: Tridiagonal, r=c=1 (e.g., in 1-d finite diff.)
- Thomas Algorithm