# Solution of Nonlinear Equations

Roots of non-linear Equations:

$$f(x) = 0$$

f may be a function belonging to any class: algebraic, trigonometric, hyperbolic, polynomials, logarithmic, exponential, etc.

We will learn the following groups of methods:

- ✓ **Bracketing methods:** Bisection, Regula-Falsi
- ✓ Open methods: Fixed point, Newton-Raphson, Secant, Muller
- ✓ **Special methods for polynomials:** Bairstow

Background assumed (MTH 101): intermediate value theorem; nested interval theorem; Cauchy sequence and convergence; Taylor's and Maclaurin's series; etc.

### **Bracketing Methods**

- ✓ Intermediate value theorem: Let f be continuous on [a, b] and let f(a) < s < f(b), then there exists at least one x such that a < x < b and f(x) = s.
  - $\checkmark$  Bracketing methods are application of this theorem with s=0
- Nested interval theorem: For each n, let  $I_n = [a_n, b_n]$  be a sequence of (non-empty) bounded intervals of real numbers such that  $I_1 \supset I_2 \supset \cdots I_n \supset I_{n+1} \supset \cdots$  and  $\lim_{n \to \infty} (a_n b_n) = 0$ , then  $\bigcap_{n=1}^{\infty} I_n$  contains only one point.
  - ✓ This guarantees the convergence of the bracketing methods to the root.
- ✓ In bracketing methods, a sequence of nested interval is generated such that each interval follows the *intermediate value theorem* with *s* = 0. Then the method converges to the root by the one point specified by the *nested interval theorem*. Methods only differ in ways to generate the nested intervals.

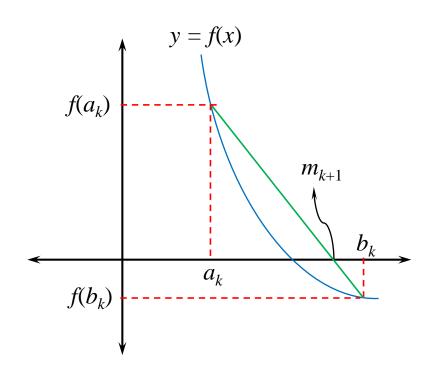
### **Bisection Method**

- ✓ Principle: Choose an initial interval based on intermediate value theorem and halve the interval at each iteration step to generate the nested intervals.
- ✓ Initialize: Choose  $a_0$  and  $b_0$  such that,  $f(a_0)f(b_0) < 0$ . This is done by trial and error.
- ✓ Iteration step k:
  - ✓ Compute mid-point  $m_{k+1} = (a_k + b_k)/2$  and functional value  $f(m_{k+1})$
  - ✓ If  $f(m_{k+1}) = 0$ ,  $m_{k+1}$  is the root. (It's your lucky day!)
  - ✓ If  $f(a_k)f(m_{k+1}) < 0$ :  $a_{k+1} = a_k$  and  $b_{k+1} = m_{k+1}$ ; else,  $a_{k+1} = m_{k+1}$  and  $b_{k+1} = b_k$
  - ✓ After *n* iterations: size of the interval  $d_n = (b_n a_n) = 2^{-n} (b_0 a_0)$ , stop if  $d_n \le \varepsilon$
  - $\checkmark$  Estimate the root  $(x = \alpha \text{ say!})$  as:  $\alpha = m_{n+1} \pm 2^{-(n+1)} (b_0 a_0)$

### Regula-Falsi or Method of False Position

- ✓ Principle: In place of the mid point, the function is assumed to be linear within the interval and the root of the linear function is chosen.
- ✓ Initialize: Choose  $a_0$  and  $b_0$  such that,  $f(a_0)f(b_0) < 0$ . This is done by trial and error.
- ✓ Iteration step k:
  - A straight line passing through two points  $(a_k, f(a_k))$  and  $(b_k, f(b_k))$  is given by:  $\frac{x a_k}{f(x) f(a_k)} = \frac{b_k a_k}{f(b_k) f(a_k)}$
  - $\checkmark$  Root of this equation at f(x) = 0 is:  $x = m_{k+1} = a_k \frac{b_k a_k}{f(b_k) f(a_k)} f(a_k)$
  - ✓ If  $f(m_{k+1}) = 0$ ,  $m_{k+1}$  is the root. (It's your lucky day!)
  - ✓ If  $f(a_k)f(m_{k+1}) < 0$ :  $a_{k+1} = a_k$  and  $b_{k+1} = m_{k+1}$ ; else,  $a_{k+1} = m_{k+1}$  and  $b_{k+1} = b_k$
  - ✓ After *n* iterations: size of the interval  $d_n = (b_n a_n)$ , stop if  $d_n \le \varepsilon$
  - $\checkmark$  Estimate the root  $(x = \alpha \text{ say!})$  as:  $\alpha = a_n \frac{b_n a_n}{f(b_n) f(a_n)} f(a_n)$

### Regula-Falsi or Method of False Position



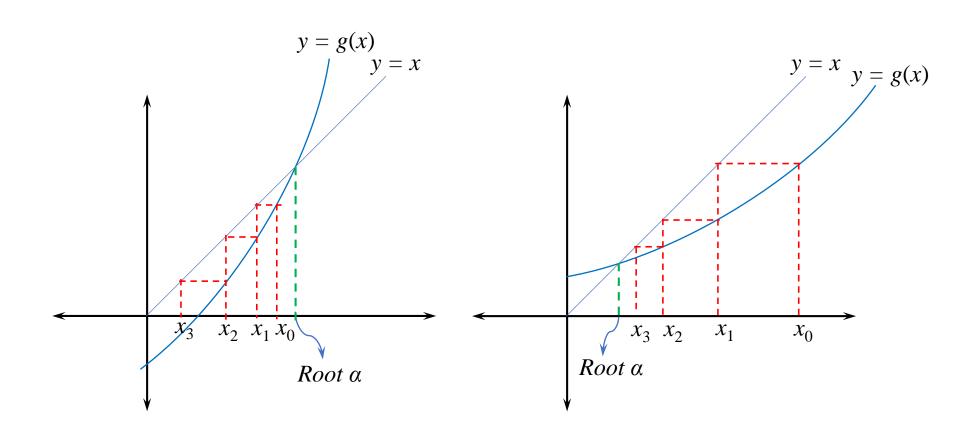
### Open Methods: Fixed Point

- ✓ Problem: f(x) = 0, find a root  $x = \alpha$  such that  $f(\alpha) = 0$ 
  - ✓ Re-arrange the function: f(x) = 0 to x = g(x)
  - ✓ Iteration:  $x_{k+1} = g(x_k)$
  - ✓ Stopping criteria:  $\left| \frac{x_{k+1} x_k}{x_k} \right| \le \varepsilon$
- $\checkmark$  Convergence: after n iterations,
  - $\checkmark$  At the root:  $\alpha = g(\alpha)$  or  $\alpha x_{n+1} = g(\alpha) g(x_n)$
  - ✓ Mean Value Theorem:  $\frac{g(\alpha)-g(x_n)}{\alpha-x_n} = g'(\xi)$  for some  $\xi \in (\alpha, x_n)$

$$(\alpha - x_{n+1}) = g'(\xi)(\alpha - x_n) \text{ or } e_{n+1} = g'(\xi) e_n \text{ or } \frac{|e_{n+1}|}{|e_n|} = |g'(\xi)|$$

- ✓ Condition for convergence:  $|g'(\xi)| < 1$
- $\checkmark$  As  $x_n \to \alpha$ ,  $\frac{|e_{n+1}|}{|e_n|} = |g'(\alpha)| = \text{constant}$

### Open Methods: Fixed Point



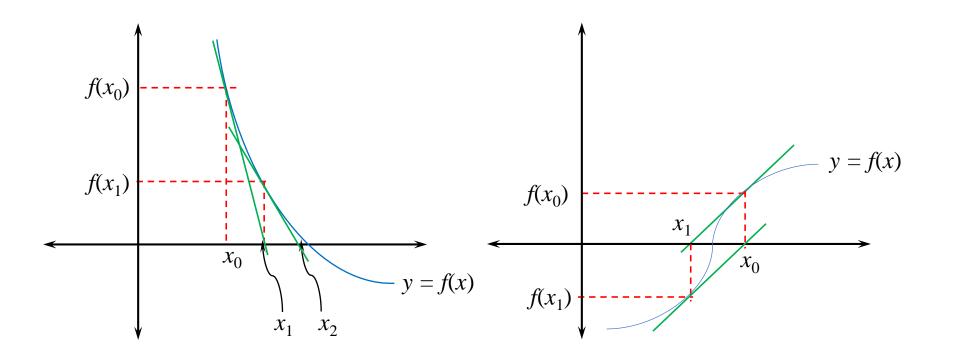
# Open Methods: Newton-Raphson

- ✓ Principle: Approximate the function as a straight line having same slope as the original function at the point of iteration.
- ✓ Problem: f(x) = 0, find a root  $x = \alpha$  such that  $f(\alpha) = 0$ 
  - ✓ Iteration Step *k*: Taylor's Theorem

$$f(x_{k+1}) = f(x_k) + (x_{k+1} - x_k) f'(x_k) + \frac{1}{2!} (x_{k+1} - x_k)^2 f''(x_k) + \dots + \frac{1}{n!} (x_{k+1} - x_k)^n f^n(x_k) + \frac{1}{(n+1)!} (x_{k+1} - x_k)^{n+1} f^{n+1}(\xi)$$
for some  $\xi \in (x_k, x_{k+1})$ 

- ✓ Assumptions: Neglect 2<sup>nd</sup> and higher order terms and assume that the root is arrived at the (k+1)<sup>th</sup> iteration, i.e.,  $f(x_{k+1}) = 0$
- ✓ Iteration Formula:  $x_{k+1} = x_k \frac{f(x_k)}{f'(x_k)}$
- ✓ Stopping criteria:  $\left| \frac{x_{k+1} x_k}{x_k} \right| \le \varepsilon$

### Open Methods: Newton-Raphson



Newton-Raphson method may get stuck!

# Open Methods: Newton-Raphson

 $\checkmark$  Convergence: Taylor's Theorem after n iterations,

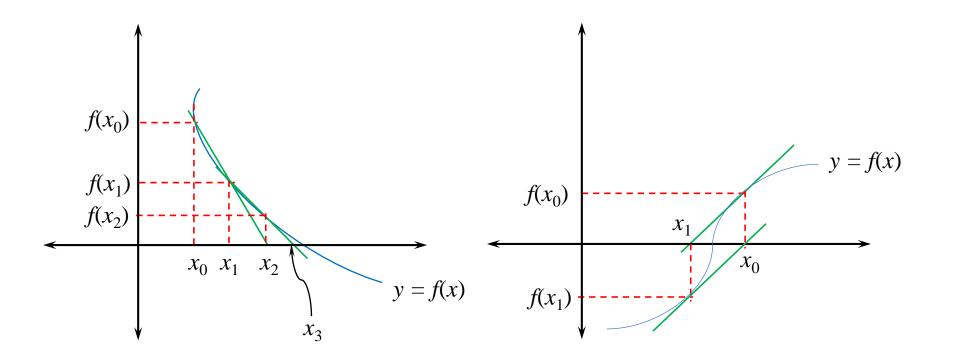
$$0 = f(\alpha) = f(x_n) + (\alpha - x_n) f'(x_n) + \frac{1}{2!} (\alpha - x_n)^2 f''(\xi)$$
  
for some  $\xi \in (x_n, \alpha)$ 

✓ Re-arrange: 
$$\frac{f(x_n)}{f'(x_n)} + (\alpha - x_n) = -\frac{1}{2!}(\alpha - x_n)^2 \frac{f''(\xi)}{f'(x_n)}$$

or 
$$(\alpha - x_{n+1}) = -\frac{1}{2!}(\alpha - x_n)^2 \frac{f''(\xi)}{f'(x_n)}$$
 or  $e_{n+1} = -\frac{1}{2!}e_n^2 \frac{f''(\xi)}{f'(x_n)}$ 

$$\checkmark$$
 As  $x_n \to \alpha$ ,  $\frac{|e_{n+1}|}{|e_n|^2} = \left| -\frac{1}{2} \frac{f''(\alpha)}{f'(\alpha)} \right| = \text{constant}$ 

- ✓ Principle: Use a difference approximation for the slope or derivative in the Newton-Raphson method. This is equivalent to approximating the tangent with a secant.
- ✓ Problem: f(x) = 0, find a root  $x = \alpha$  such that  $f(\alpha) = 0$ 
  - ✓ Initialize: choose two points  $x_0$  and  $x_1$  and evaluate  $f(x_0)$  and  $f(x_1)$
  - ✓ Approximation:  $f'(x_k) \approx \frac{f(x_k) f(x_{k-1})}{x_k x_{k-1}}$ , replace in Newton-Raphson
  - ✓ Iteration Formula:  $x_{k+1} = x_k f(x_k) \frac{x_k x_{k-1}}{f(x_k) f(x_{k-1})}$
  - ✓ Stopping criteria:  $\left| \frac{x_{k+1} x_k}{x_k} \right| \le \varepsilon$



✓ Convergence: Newton's polynomial,

$$0 = f(\alpha) = f(x_n) + (\alpha - x_n)f[x_n, x_{n-1}] + \frac{1}{2}(\alpha - x_n) (\alpha - x_{n-1}) f''(\xi)$$
for some  $\xi \in (x_{n-1}, x_n, \alpha)$ 

Divided difference: 
$$f[x_n, x_{n-1}] = \frac{f(x_n) - f(x_{n-1})}{x_n - x_{n-1}}$$

✓ Re-arrange:

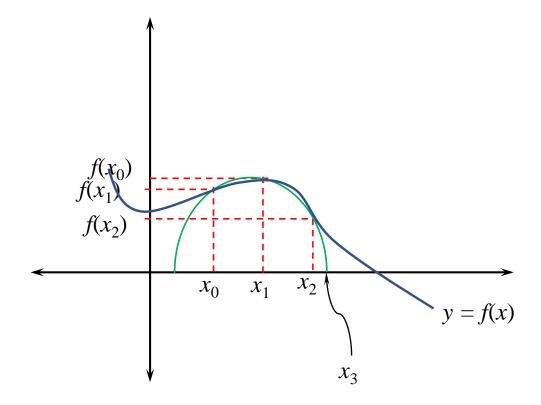
$$\frac{f(x_n)}{f[x_{n-1},x_n]} + (\alpha - x_n) = (\alpha - x_{n+1}) = -\frac{1}{2}(\alpha - x_n) (\alpha - x_{n-1}) \frac{f''(\xi)}{f[x_{n-1},x_n]}$$

$$e_{n+1} = -\frac{1}{2} e_n e_{n-1} \frac{f''(\xi)}{f'(\xi')}, \quad \xi' \in (x_{n-1}, x_n)$$
 [using mean value theorem]

$$\checkmark$$
 As  $x_n \to \alpha$ ,  $\frac{|e_{n+1}|}{|e_n||e_{n-1}|} = \left|-\frac{1}{2} \frac{f''(\alpha)}{f'(\alpha)}\right| = \text{constant}$ 

- ✓ Derivative of the Newton-Raphson method is evaluated numerically using difference approximation.
- ✓ Numerical methods for estimation of derivative of a function will be covered in detail later.
- ✓ Rest of the method is same.

- ✓ Principle: fit a quadratic polynomial through three points to approximate the function.
- ✓ Problem: f(x) = 0, find a root  $x = \alpha$  such that  $f(\alpha) = 0$



- ✓ Initialize: choose three points  $x_0$ ,  $x_1$ ,  $x_2$  and evaluate  $f(x_0)$ ,  $f(x_1)$ ,  $f(x_2)$ . Denote  $f(x_k) = f_k$ 
  - ✓ Fit a quadratic polynomial through three points  $x_k$ ,  $x_{k-1}$ ,  $x_{k-2}$  as:

$$p(x) = a(x - x_k)^2 + b(x - x_k) + c$$

#### ✓ Equations:

$$p(x_k) = f_k = c$$

$$p(x_{k-1}) = f_{k-1} = a(x_{k-1} - x_k)^2 + b(x_{k-1} - x_k) + f_k$$

$$p(x_{k-2}) = f_{k-2} = a(x_{k-2} - x_k)^2 + b(x_{k-2} - x_k) + f_k$$

✓ Define Divided Differences:

- ✓ 1st Divided Difference:  $f[x_k, x_{k-1}] = \frac{f(x_k) f(x_{k-1})}{x_k x_{k-1}}$
- ✓ 2<sup>nd</sup> Divided Difference:  $f[x_k, x_{k-1}, x_{k-2}] = \frac{f[x_k, x_{k-1}] f[x_{k-1}, x_{k-2}]}{x_k x_{k-2}}$

$$p(x_k) = f_k = c$$

$$p(x_{k-1}) = f_{k-1} = a(x_{k-1} - x_k)^2 + b(x_{k-1} - x_k) + f_k$$

$$p(x_{k-2}) = f_{k-2} = a(x_{k-2} - x_k)^2 + b(x_{k-2} - x_k) + f_k$$

#### ✓ Express as:

$$\frac{f_k - f_{k-1}}{x_k - x_{k-1}} = f[x_k, x_{k-1}] = a(x_{k-1} - x_k) + b$$

$$\frac{f_k - f_{k-2}}{x_k - x_{k-2}} = f[x_k, x_{k-2}] = a(x_{k-2} - x_k) + b$$

#### ✓ Solutions:

$$a = \frac{f[x_k, x_{k-1}] - f[x_k, x_{k-2}]}{x_{k-1} - x_{k-2}}$$
$$b = f[x_k, x_{k-2}] + a(x_k - x_{k-2})$$

### Properties of Divided Differences

$$\begin{aligned} & 2^{\text{nd}} \text{ Divided Difference: } f[x_k, x_{k-1}, x_{k-2}] = \frac{f[x_k, x_{k-1}] - f[x_{k-1}, x_{k-2}]}{x_k - x_{k-2}} \\ & a = \frac{f[x_k, x_{k-1}] - f[x_k, x_{k-2}]}{x_{k-1} - x_{k-2}} \\ & = \frac{\frac{f_k - f_{k-1}}{x_k - x_{k-1}} - \frac{f_k - f_{k-2}}{x_k - x_{k-2}}}{x_{k-1} - x_{k-2}} \\ & = \frac{x_k f_k - x_k f_{k-1} - x_{k-2} f_k + x_{k-2} f_{k-1} - \frac{x_k f_k}{x_k} + x_k f_{k-2} + x_{k-1} f_k - x_{k-1} f_{k-2}}{(x_k - x_{k-1})(x_{k-1} - x_{k-2})(x_k - x_{k-2})} \\ & = \frac{(x_{k-1} - x_{k-2}) f_k - (x_{k-1} - x_{k-2}) f_{k-1} - (x_k - x_{k-1}) f_{k-1} + (x_k - x_{k-1}) f_{k-2}}{(x_k - x_{k-1})(x_{k-1} - x_{k-2})(x_k - x_{k-2})} \\ & = \frac{\frac{f_k - f_{k-1}}{x_k - x_{k-1}} - \frac{f_{k-1} - f_{k-2}}{x_{k-1} - x_{k-2}}}{x_{k-1} - x_{k-2}} = \frac{f[x_k, x_{k-1}] - f[x_{k-1}, x_{k-2}]}{x_k - x_{k-2}} = f[x_k, x_{k-1}, x_{k-2}] \end{aligned}$$

### Note: Properties of Divided Differences

✓ 1<sup>st</sup> Divided Difference:

$$f[x_k, x_{k-1}] = \frac{f_k - f_{k-1}}{x_k - x_{k-1}} = \frac{f_{k-1} - f_k}{x_{k-1} - x_k} = f[x_{k-1}, x_k]$$

✓ 2nd Divided Difference:

$$f[x_k, x_{k-1}, x_{k-2}] = \frac{f[x_k, x_{k-1}] - f[x_{k-1}, x_{k-2}]}{x_k - x_{k-2}} = \frac{f[x_k, x_{k-1}] - f[x_k, x_{k-2}]}{x_{k-1} - x_{k-2}} = \frac{f[x_k, x_{k-2}] - f[x_{k-1}, x_{k-2}]}{x_k - x_{k-1}}$$

$$f[x_k, x_{k-1}, x_{k-2}] = f[x_{k-1}, x_k, x_{k-2}] = f[x_{k-2}, x_{k-1}, x_k] = f[x_k, x_{k-2}, x_{k-1}]$$

We shall use these properties for the *Theory of Approximation*!

- ✓ Principle: fit a quadratic polynomial through three points to approximate the function.
- ✓ Problem: f(x) = 0, find a root  $x = \alpha$  such that  $f(\alpha) = 0$ 
  - ✓ Initialize: choose three points  $x_0$ ,  $x_1$ ,  $x_2$  and evaluate  $f(x_0)$ ,  $f(x_1)$ ,  $f(x_2)$
  - ✓ Fit a quadratic polynomial through three points  $x_k$ ,  $x_{k-1}$ ,  $x_{k-2}$  as:

$$p(x) = a(x - x_k)^2 + b(x - x_k) + c$$

- Constants:  $c = f(x_k)$ ;  $b = f[x_k, x_{k-1}] + (x_k x_{k-1})f[x_k, x_{k-1}, x_{k-2}]$ ;  $a = f[x_k, x_{k-1}, x_{k-2}]$
- ✓ Iteration step k:  $x_{k+1} = x_k \frac{b \pm \sqrt{b^2 4ac}}{2a}$  or  $x_{k+1} = x_k \frac{2c}{b \pm \sqrt{b^2 4ac}}$
- ✓ Stopping criteria:  $\left|\frac{x_{k+1}-x_k}{x_k}\right| \le \varepsilon$

✓ Convergence: Analysis is similar to Secant method, using Newton's Polynomial.

$$e_{n+1} = -\frac{1}{6}e_n e_{n-1} e_{n-2} \frac{f'''(\xi)}{f'(\xi')},$$

$$\xi \in (x_{n-2}, x_{n-1}, x_n, \alpha), \qquad \xi' \in (x_{n-1}, x_n)$$
as  $x_n \to \alpha$ ,  $\frac{|e_{n+1}|}{|e_n||e_{n-1}||e_{n-2}|} = \left| -\frac{1}{6} \frac{f'''(\alpha)}{f'(\alpha)} \right| = \text{constant}$ 

### Order of Convergence

✓ Definition: Let  $\{x_0, x_1, x_2, x_3 ...\}$  be a sequence which converges to  $\alpha$ . Define  $e_n = x_n - \alpha$ . If there exists a number p and a constant  $C \neq 0$  such that,

$$\lim_{n \to \infty} \frac{|e_{n+1}|}{|e_n|^p} = C$$

Then, p is called the order of convergence of the sequence and C is the asymptotic error constant.

- ✓ Fixed Point:  $\frac{|e_{n+1}|}{|e_n|} = |g'(\alpha)| = C$ , 1st Order.
- ✓ Newton Raphson:  $\frac{|e_{n+1}|}{|e_n|^2} = \left| -\frac{1}{2} \frac{f''(\alpha)}{f'(\alpha)} \right| = C$ , 2<sup>nd</sup> Order
- ✓ Secant:  $\frac{|e_{n+1}|}{|e_n||e_{n-1}|} = \left|-\frac{1}{2} \frac{f''(\alpha)}{f'(\alpha)}\right| = C$ , mixed order, ≈ 1.6
- ✓ Muller:  $\frac{|e_{n+1}|}{|e_n||e_{n-1}||e_{n-2}|} = \left| -\frac{1}{6} \frac{f'''(\alpha)}{f'(\alpha)} \right| = C$ , mixed order, ≈ 1.84

# Polynomial Methods: Single Root

$$p_n(x) = \sum_{k=0}^n a_k x^k = a_0 + a_1 x + a_2 x^2 + \dots + a_n x^n$$

If we divide by a factor (x - r) such that,  $r = \alpha$  is a root of the polynomial, we will get an exact polynomial of order (n - 1).

$$q_{n-1}(x) = \sum_{k=0}^{n-1} b_{k+1} x^k = b_1 + b_2 x + b_3 x^2 + \dots + b_n x^{n-1}$$

If  $r \neq \alpha$ , dividing by a factor (x - r) will have a remainder  $b_0$ .

# Polynomial Methods: Single Root

$$p_n(x) = \sum_{k=0}^n a_k x^k = a_0 + a_1 x + a_2 x^2 + \dots + a_n x^n$$

$$= (x - r)q_{n-1}(x) + b_0 = (x - r)\sum_{k=0}^{n-1} b_{k+1}x^k + b_0$$

$$= b_0 + b_1(x-r) + b_2x(x-r) + b_3x^2(x-r) + \dots + b_{n-2}x^{n-3}(x-r) + b_{n-1}x^{n-2}(x-r) + b_nx^{n-1}(x-r)$$

$$= (b_0 - rb_1) + x(b_1 - rb_2) + x^2(b_2 - rb_3) + \dots + x^{n-2}(b_{n-2} - rb_{n-1}) + x^{n-1}(b_{n-1} - rb_n) + b_n x^n$$

$$b_n = a_n$$
;  $b_i = a_i + rb_{i+1}$ ;  $i = (n-1), (n-2), \dots 2, 1, 0$ 

# Polynomial Methods: Single Root

 $b_0$  is a function of  $r \to b_0(r)$ , at  $r = \alpha$ ,  $b_0(r) = 0$ 

Problem: f(x) = 0, find a root  $x = \alpha$  such that  $f(\alpha) = 0$ 

Problem:  $b_0(r) = 0$ , find a root  $r = \alpha$  such that  $b_0(\alpha) = 0$ 

Apply Newton-Raphson:

Iteration Formula for Step *k*:

$$x_{k+1} = x_k - \frac{f(x_k)}{f'(x_k)}$$
 or  $r_{k+1} = r_k - \frac{b_0(r_k)}{b'_0(r_k)}$ 

$$b_0 = a_0 + rb_1 \rightarrow b_0'(r) = b_1 \rightarrow r_{k+1} = r_k - \frac{b_0(r_k)}{b_1(r_k)}$$

Assume a value of r, estimate  $b_0$  and  $b_1$ , compute new r.

Continue until  $b_0$  becomes zero. (with acceptable relative error)

$$p_n(x) = \sum_{k=0}^n a_k x^k = a_0 + a_1 x + a_2 x^2 + \dots + a_n x^n$$

Let us divide by a factor  $(x^2 - rx - s)$ . If the factor is exact, the resulting polynomial will be of order (n - 2). Two roots of the polynomial can be estimated simultaneously as the roots of the quadratic factor. For the complex roots, they will be the complex conjugates.

$$q_{n-2}(x) = \sum_{k=0}^{n-2} b_{k+2} x^k = b_2 + b_3 x + b_4 x^2 + \dots + b_n x^{n-2}$$

If the factor  $(x^2 - rx - s)$  is not exact, there will be two remainder terms, one function of x and another constant.

Let us express the remainder term as  $b_1(x - r) + b_0$ . This form instead of the standard  $b_1x + b_0$  is chosen to device a convenient iteration formula!

$$\begin{split} p_n(x) &= \sum_{k=0}^n a_k x^k = a_0 + a_1 x + a_2 x^2 + \dots + a_n x^n \\ &= (x^2 - rx - s) q_{n-2}(x) + b_1 (x - r) + b_0 \\ &= (x^2 - rx - s) \sum_{k=0}^{n-2} b_{k+2} x^k + b_1 (x - r) + b_0 \\ &= b_0 + b_1 (x - r) + b_2 (x^2 - rx - s) + b_3 x (x^2 - rx - s) + \dots \\ &+ b_{n-2} x^{n-4} (x^2 - rx - s) + b_{n-1} x^{n-3} (x^2 - rx - s) + b_n x^{n-2} (x^2 - rx - s) \\ &= (b_0 - rb_1 - sb_2) + x (b_1 - rb_2 - sb_3) + x^2 (b_2 - rb_3 - sb_4) + \dots \\ &+ x^{n-2} (b_{n-2} - rb_{n-1} - sb_n) + x^{n-1} (b_{n-1} - rb_n) + b_n x^n \end{split}$$

 $b_n = a_n$ ;  $b_{n-1} = a_{n-1} + rb_n$ ;  $b_i = a_i + rb_{i+1} + sb_{i+2}$ ;  $i = (n-2), \dots 2, 1, 0$ 

 $b_0$  and  $b_1$  are functions of r and  $s \rightarrow b_0(r, s)$  and  $b_1(r, s)$ 

Expand in Taylor's series: Apply 2-d Newton-Raphson

$$0 = b_0(r + \Delta r, s + \Delta s) = b_0 + \frac{\partial b_0}{\partial r} \Delta r + \frac{\partial b_0}{\partial s} \Delta s + HOT$$
$$0 = b_1(r + \Delta r, s + \Delta s) = b_1 + \frac{\partial b_1}{\partial r} \Delta r + \frac{\partial b_1}{\partial s} \Delta s + HOT$$

$$\begin{bmatrix} \frac{\partial b_0}{\partial r} & \frac{\partial b_0}{\partial s} \\ \frac{\partial b_1}{\partial r} & \frac{\partial b_1}{\partial s} \end{bmatrix} \begin{bmatrix} \Delta r \\ \Delta s \end{bmatrix} = \begin{bmatrix} -b_0 \\ -b_1 \end{bmatrix}$$

Need to evaluate:  $\frac{\partial b_0}{\partial r}$ ,  $\frac{\partial b_0}{\partial s}$ ,  $\frac{\partial b_1}{\partial r}$  and  $\frac{\partial b_1}{\partial s}$ 

Partial differentials with respect to *r*:

$$b_{n} = a_{n} \rightarrow \frac{\partial b_{n}}{\partial r} = 0; \qquad b_{n-1} = a_{n-1} + rb_{n} \rightarrow \frac{\partial b_{n-1}}{\partial r} = b_{n} = c_{n}$$

$$b_{n-2} = a_{n-2} + rb_{n-1} + sb_{n} \rightarrow \frac{\partial b_{n-2}}{\partial r} = b_{n-1} + r\frac{\partial b_{n-1}}{\partial r} + s\frac{\partial b_{n}}{\partial r}$$

$$= b_{n-1} + rc_{n} = c_{n-1}$$

$$b_{n-3} = a_{n-3} + rb_{n-2} + sb_{n-1} \rightarrow \frac{\partial b_{n-3}}{\partial r} = b_{n-2} + r\frac{\partial b_{n-2}}{\partial r} + s\frac{\partial b_{n-1}}{\partial r}$$

$$= b_{n-2} + rc_{n-1} + sc_{n} = c_{n-2}$$

$$c_n = b_n$$
;  $c_{n-1} = b_{n-1} + rc_n$ ;  $c_i = b_i + rc_{i+1} + sc_{i+2}$ ;  $i = (n-2), \dots 2, 1, 0$   
$$\frac{\partial b_i}{\partial r} = c_{i+1}$$
;  $i = (n-1), \dots 2, 1, 0$ 

Partial differentials with respect to *s*:

$$\begin{split} b_{n} &= a_{n} \quad \rightarrow \frac{\partial b_{n}}{\partial s} = 0; \ b_{n-1} = a_{n-1} + rb_{n} \quad \rightarrow \frac{\partial b_{n-1}}{\partial s} = 0 \\ b_{n-2} &= a_{n-2} + rb_{n-1} + sb_{n} \rightarrow \frac{\partial b_{n-2}}{\partial s} = b_{n} + r\frac{\partial b_{n-1}}{\partial s} + s\frac{\partial b_{n}}{\partial s} = b_{n} = c_{n} \\ b_{n-3} &= a_{n-3} + rb_{n-2} + sb_{n-1} \rightarrow \frac{\partial b_{n-3}}{\partial s} = b_{n-1} + r\frac{\partial b_{n-2}}{\partial s} + s\frac{\partial b_{n-1}}{\partial s} \\ &= b_{n-1} + rc_{n} = c_{n-1} \\ b_{n-4} &= a_{n-4} + rb_{n-3} + sb_{n-2} \rightarrow \frac{\partial b_{n-4}}{\partial s} = b_{n-2} + r\frac{\partial b_{n-3}}{\partial s} + s\frac{\partial b_{n-2}}{\partial s} \\ &= b_{n-2} + rc_{n-1} + sb_{n} = c_{n-2} \end{split}$$

$$c_n = b_n$$
;  $c_{n-1} = b_{n-1} + rc_n$ ;  $c_i = b_i + rc_{i+1} + sc_{i+2}$ ;  $i = (n-2), \dots 2, 1, 0$   
$$\frac{\partial b_i}{\partial s} = c_{i+2}$$
;  $i = (n-2), \dots 2, 1, 0$ 

$$\frac{\partial b_i}{\partial r} = c_{i+1}; \ i = (n-1), \dots 2, 1, 0 \text{ and } \frac{\partial b_i}{\partial s} = c_{i+2}; \ i = (n-2), \dots 2, 1, 0$$

$$\frac{\partial b_0}{\partial r} = c_1; \ \frac{\partial b_1}{\partial r} = c_2; \ \frac{\partial b_0}{\partial s} = c_2 \text{ and } \frac{\partial b_1}{\partial s} = c_3$$

$$\begin{bmatrix} \frac{\partial b_0}{\partial r} & \frac{\partial b_0}{\partial s} \\ \frac{\partial b_1}{\partial r} & \frac{\partial b_1}{\partial s} \end{bmatrix} \begin{bmatrix} \Delta r \\ \Delta s \end{bmatrix} = \begin{bmatrix} c_1 & c_2 \\ c_2 & c_3 \end{bmatrix} \begin{bmatrix} \Delta r \\ \Delta s \end{bmatrix} = \begin{bmatrix} -b_0 \\ -b_1 \end{bmatrix}$$

For any given polynomial, we know  $\{a_0, a_1, \dots a_n\}$ . Assume r and s. Compute  $\{b_0, b_1, \dots b_n\}$  and  $\{c_0, c_1, \dots c_n\}$ . Compute  $\Delta r$  and  $\Delta s$ .

### Polynomial Methods: Bairstow Algorithm

- ✓ Step 1: input  $a_0, a_1, \dots a_n$  and initialize r and s.
- $\checkmark$  Step 2: compute  $b_0, b_1, \dots b_n$

$$b_n = a_n; b_{n-1} = a_{n-1} + rb_n; b_i = a_i + rb_{i+1} + sb_{i+2}; i = (n-2), \dots 2, 1, 0$$

 $\checkmark$  Step 3: compute  $c_0, c_1, \dots c_n$ 

$$c_n = b_n$$
;  $c_{n-1} = b_{n-1} + rc_n$ ;  $c_i = b_i + rc_{i+1} + sc_{i+2}$ ;  $i = (n-2), \dots 2, 1, 0$ 

- ✓ Step 4: compute  $\Delta r$  and  $\Delta s$  from  $\begin{bmatrix} c_1 & c_2 \\ c_2 & c_3 \end{bmatrix} \begin{bmatrix} \Delta r \\ \Delta s \end{bmatrix} = \begin{bmatrix} -b_0 \\ -b_1 \end{bmatrix}$
- ✓ Step 5: compute  $r_{new} = r + \Delta r$ ,  $s_{new} = s + \Delta s$
- ✓ Step 6: check for convergence,  $\left|\frac{r_{new}-r}{r}\right|$ ,  $\left|\frac{s_{new}-s}{s}\right| \le \varepsilon$  and  $b_0, b_1 \le \varepsilon'$
- ✓ Step 7: Stop if all convergence checks are satisfied. Else, set  $r = r_{new}$ ,  $s = s_{new}$  and go to step 2.

### Multiple Roots

✓ Definition: A root  $\alpha$  of the equation f(x) = 0 is said to have a multiplicity of q if,

$$g(x) = \frac{f(x)}{(x-\alpha)^q}$$
  $0 \neq g(\alpha) < \infty$ 

when, q > 1, the order of convergence are no longer valid.

✓ Solution: Suppose a function f(x) is q-times continuously differentiable in the neighbourhood of a root  $\alpha$  of multiplicity q,

$$f(x) = \frac{1}{q!}(x - \alpha)^q f^q(\xi) \text{ and } f'(x) = \frac{1}{(q-1)!}(x - \alpha)^{q-1} f^q(\xi')$$
where  $\xi, \xi' \in (x, \alpha)$ 

$$\text{Define } u(x) = \frac{f(x)}{f'(x)} = \frac{1}{q}(x - \alpha) \frac{f^q(\xi)}{f^q(\xi')}$$

$$\lim_{x \to \alpha} \frac{u(x)}{(x - \alpha)} = \frac{1}{q}$$

Therefore,  $\alpha$  is a root of f(x) of multiplicity q but is a simple root of u(x)!