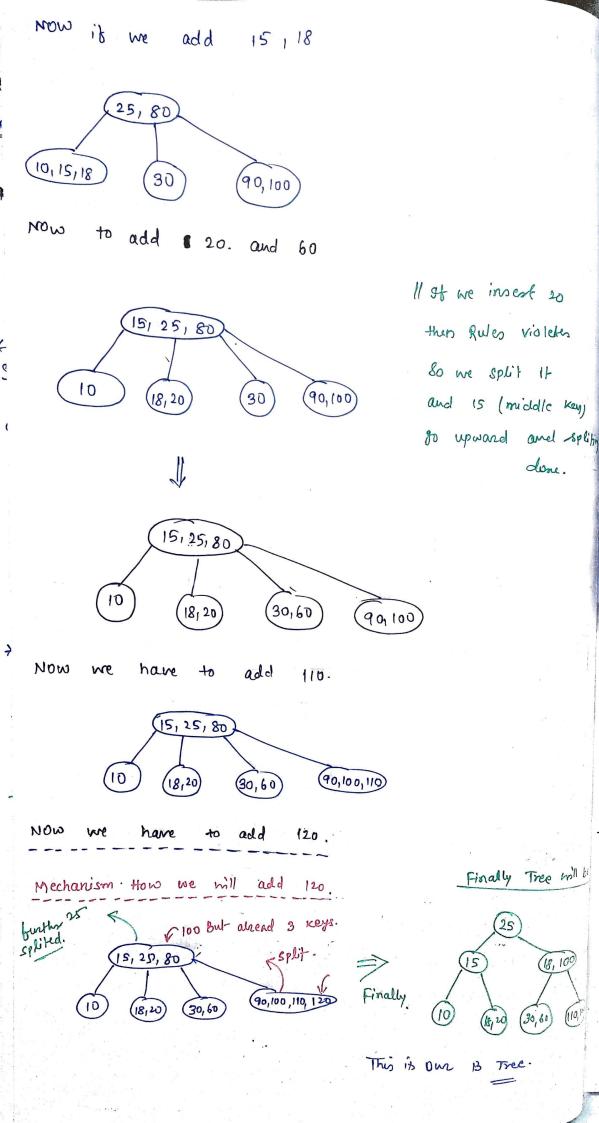
Date: 05/09/2021 19. B and Bt Tree DBMS Notes Written by Abhishek Sharma (GFG counce Q) What are B and 3+ trees are ?? Dr: They are n ways search Trees. * Every Node of a Tree has some no. of values. and the no. of values that a node can contain is defined by Branch factor. . If Branch factor is b, then maximum b children are in the node and minimum [b/2] children. EX. if Branch factor = 4. then minimum no of children = $\left[\frac{4}{2}\right] = 2$ children. maximum no. of children = 4 children. * No. of Keys present in accorde in a ways search Tree node = (no. of children -1) Example: 1. 1. 1 it a Trees has 4 children then No. of keys = (4-1)=3 --> Keys 10, 20, 30 Ex. - No. of 2,5,8 11,19 25,28 35,38 logger less than 6/w 6/10 greater 10-20 20-30 10 than 30

* B and B+ Trees are self Balanced Search Trees. Ex If There are millions of Record and millions of hight then they manage them by themselves. * They make sure that hight never goes beyond o (log n) if these are n records to store Q> HOW B and B+ Trees maintain their Height ?? de we will see via this example. 10, 30, 25, 80, 90, 100 Given: 15, 18, 20, 60, 110, 120. Branch factor = 4. we have to insest these values in the Tree. Let's See. b/2 minimum children bormula is not applicable for root. 10 Now we will issert 90 Now if we try to innest Here 80, then it violents the (30,8999 10 maximum children property. it ? it we add 100 same problem has b=4 children now. so we will split it. (25,80) 90,100 socidate key will go up and split.



Abhishek Shanma Notro. B Trees Height grows upward unlike the normal Binary Search Trees Where Height grows downward. in BST We add always from the leaf and that's insert height always grows downwards. we always insest from in B and B+ Trees the root not and Height growing from In the previous examples we saw that if the Height no of keys increases Height grows upward. * Similarly Deletion Operation, it (B and B+ Trees) srinks or deleted. the Height when no of keys removed to represent index files B Trees are used > so they have pointers of actual records. actual Helphotos of the disk when the records sources and blocks They also have pointers to children rodes. 2 Pointers. i) child pointer and ii) path pointer So B Trees has Bor (Disk Block) How the B Trees store data. ?? Q. 20 10 Poinky 15 Data Pointes for all Attribute Data pointer for all attrible value 20 (Disk Block) 10 (PISK Block) -vache + 11 same

Bti Tree.

It is the improved version of B Tree.

It is the most used Tree in patabase for indexing

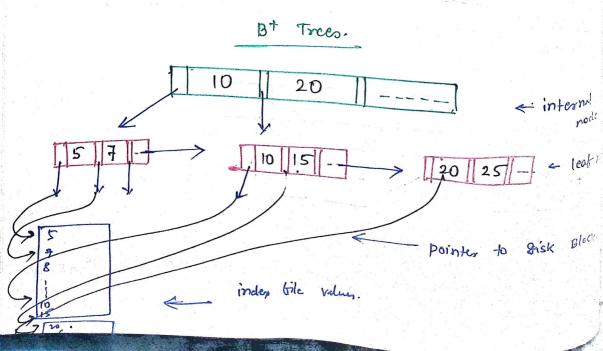
like oracle, Microsoft saul Server. uses this (B+ Tree)

Idea of B+ Tree:

All the leaf nodes are connected with each other and Leaf nodes allows you to sequencially access the data.

Since leaf nodes allows us to seawnially access the data so, every data item has to be present in the leaf nodes. Whatever present in the internal node is also to be present in the leaf nodes. So that we can access all the table entries in a sequencial manner.

* Bt Trees are index blocks.



you do not need 2 pointers. You only need one pointer with every key attributer also we have sequencial Access.

* Drawbacks of B+ Trees.

This might happen that key value might be present.

Two times. (i.e. all the key value that is present in the internal nodes will be also be present in the leat nodes also).

* 8t Tree 1's the fantastic Data Structure that works upon multilevel indexing. They grow and shrinks automatically. They have only one pointer bcz your internal nodes have block pointers pointing to Disk blocks. Which have index file values (as we saw in the Diagram).

* Both B and B+ Trees insure that every leaf node present at the same level.

* Height of these (B and B+) Trees is actually log n.

20. ACID Properties.

Abhishek sharme

Transactions and concurrency control (ACIP property).

Fransaction: A set of logical instruction for the Database.

that Should happen together. Either all of them Happen

or non of them will happen.

Example. In a Bank data of Two people

of and y

Our tank is to send 100 Rs. from x to y.

So the instructions are.

Before Transfaction:

X: 500

y : 200

v ,

Read (x)

2 = x -100

Write (x): 400

Read (Y)

A = A +100

write (Y) : 300

Abter Transaction.

2: 400

y: 300

This is the whole process

for Sending 100 Ro. from

x to y.

So Transaction says est

nothing will happen.

Bich A stop -rel ADI

Step-

Abhishek Sharom Notes * Most of the databases have the syntam available for Fransachor You have to give prostruction to DBMS that this is the Transaction, either it should happen completely or should not

(DBMS)

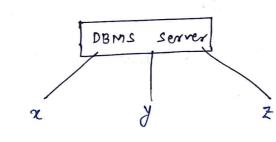
Transaction is one problem that your DBMs should solve Another problem that comes with Transaction is concurrence

Concurrency 1 concurrency says that it a DBMS server is doing multiple transactions. Then they all should in concurrent way.

same as operating system when a proceess that is doing IO, you bring another process that can you use Cpv.

Example.

happen at all.



E DBMS Server

multiple. 4 Transactions

Transaction 1. x = 500 Rs. y = 200 Rs.

Transferry 100 Rs.

Transfering 200 Rm to a account.

7 = 200 Ro.

Transaction 2.

from x to y.

Let's Discuss if a Tranaction is happen in Translation 1. and it stops at beduction on 100 from a account and at the same time I's account Transfer the money to se account

Abhishek sharme notes. we might not get our Desireed Resut. So, the There are the problems happen in Transaction and concurrency. To get Rid of these types of problem we have ACID Properties. in our DBMS. These are 4 certain things that will insure Transaction and Concurency in the DBMS. These are. ACID properties. Atomicity 2) consistency 3) Isolation 4) Durability. Atomicity: Either the whole Transaction happen or nothing will happen. Consistency; your Database should be in consistent stop State After the Transaction has happened. In a Bank., two People x and y have morey Example: 400 Rispectively. After any Transaction happen. 500 and Sum of them -then the remain constart. It should not as 2 = 500 = sum = 900 it any Transaction Sum will be Y = ,400 Happen 900. Bebor Transaction. This is consistency says.

Q

Money should not go anywhere. It should be constant before and after the Transaction. This is consistency says.

Your BBMS should be in consistent state.

Isolation: There are multiple transaction that is doing by

Your DBMS server. Isolation says that every Transaction

feel like that it is running alone. Any Transaction should

not be impacted by another Transaction. All the

Transaction should be running in alone seperately and

in a Sequencial manner.

Ex: It your are doing Internet Banking, and at the same time anyone Sending you money. So there are Two Transaction is happing at the same time. But your internet Banking Transaction will be alone and sate It will not an effect of anyone Sending you money. at the same time.

Duanability: Duanability Says that once your transaction is done, whatever will be the changes it should be in stay your DBMs. It should not be lost. Changes should be durable. It should not be lost.