

Bit Manipulation - 2

Jul 19, 2023

AGENDA

- Left shift and right shift operator
- Bit masking and related problems
- Negative nos.
- Ranges and overflows

Revision

$$11 \mid 1 = 11$$

$$\begin{array}{l} a \& 1 \\ \text{(for some no. } a) \end{array} = \begin{array}{l} 1 \text{ if } a \text{ is odd} \\ 0 \text{ if } a \text{ is even.} \end{array}$$

$$a/a = a$$

$$a^{\wedge} a = 0$$

Left shift

\ll

It left-shifts all the bits.

00101011 $\ll 1$: left shift all bits by 1.

01010110

[0 is put in LSB, all other bits are shifted to left by one].

00001010 $\ll 3$: left shift all bits by 3.
 \Rightarrow 01010000

Consider 8 bits

$$5 \ll 1 \\ = 10$$

$$13 \ll 1 \\ = 26$$

$$15 \ll 1 \\ = 30$$

00000101
 \hookrightarrow 00001010

00001101
 \hookrightarrow 00011010

00001111
 \hookrightarrow 00011110

(MSB) This is always 0.

6 5 4 3 2 1 0
 0001011

6 5 4 3 2 1 0
 0010110

$$= 0 \times 2^6 + 0 \times 2^5 + 0 \times 2^4 + 1 \times 2^3 + 0 \times 2^2 + 1 \times 2^1 + 1 \times 2^0$$

$$= 0 \times 2^6 + 0 \times 2^5 + 1 \times 2^4 + 0 \times 2^3 + 1 \times 2^2 + 1 \times 2^1 + 0 \times 2^0$$

$$* \quad a \ll 1 = 2 * a$$

$$* \quad \begin{aligned} a \ll 2 &= (a * 2) * 2 \\ a \ll 3 &= ((a * 2) * 2) * 2 \\ a \ll 5 &= a * 2^5 \end{aligned}$$

$$* \quad a \ll m = a * 2^m \rightarrow O(1) \text{ operation}$$

Right shift

\gg

$$\begin{aligned} &0000\overset{\text{pink}}{1}\overset{\text{pink}}{0}\overset{\text{pink}}{1}\overset{\text{pink}}{1} \gg 1 \\ = &00000101 \end{aligned}$$

(All bits right shift by 1, LSB disappears).

$$\begin{aligned} 5 \gg 1 \\ = 2 \end{aligned}$$

$$\begin{aligned} 13 \gg 1 \\ = 6 \end{aligned}$$

$$\begin{aligned} 12 \gg 1 \\ = 6 \end{aligned}$$

$$\begin{aligned} 14 \gg 1 \\ = 7 \end{aligned}$$

$$\begin{aligned} 101 \gg 1 \\ = 10 \end{aligned}$$

$$\begin{aligned} 1101 \gg 1 \\ = 110 \end{aligned}$$

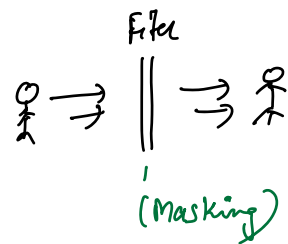
$$\begin{aligned} 1100 \gg 1 \\ = 110 \end{aligned}$$

$$\begin{aligned} 1110 \gg 1 \\ = 111 \end{aligned}$$

* $a \gg 1 = a/2$ [Integer division]

* $a \gg m = a/2^m$

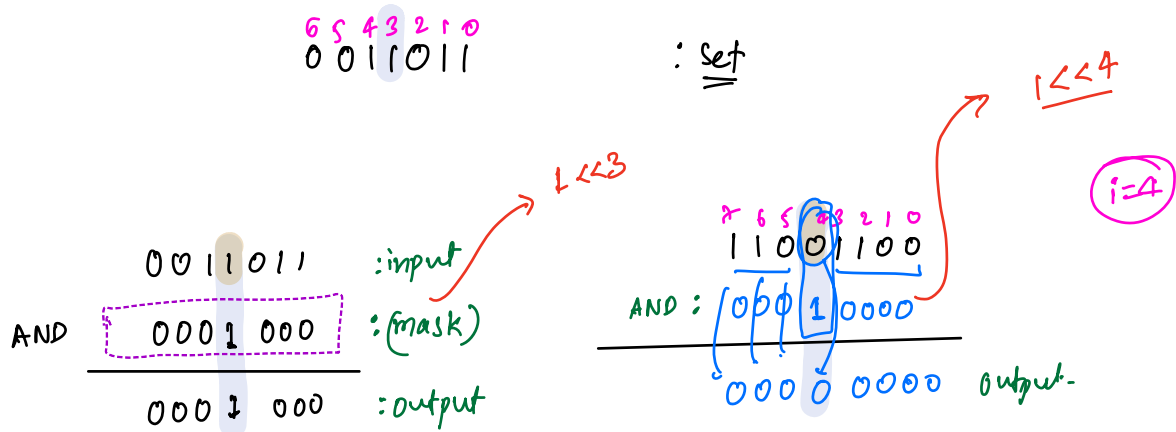
Bit masking



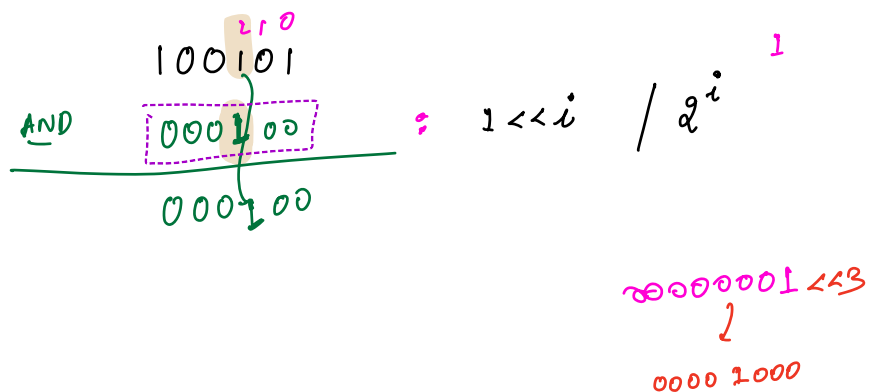
Q. Given a no. N , check if i^{th} bit is set or not.
 ↳ 0-indexed from the right.

Set bit : 1
 Unset bit : 0

e.g.. $N = 27$, $i = 3$



$N = 37$, $i = 2$, check whether i^{th} bit is set.



Code

```
checkIthbit(int N, int i)
{
    mask = 1 << i
    res = N & mask
    if(res == 0)
        return false    // Bit was unset.
    else
        return true     // Bit was set.
}
```

Q.

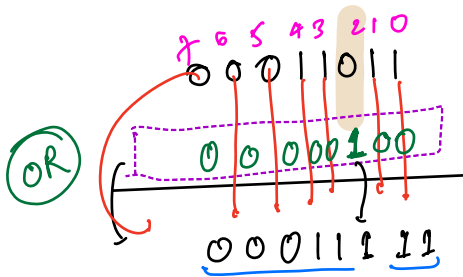
Given a no. N , set the i th bit.

; Make i th bit 1.

(no change if it is already one),

$N = 27$,

$i = 2$



setIthbit (int N , int i)

{

mask = $1 \ll i$;

return $N | \text{mask}$;

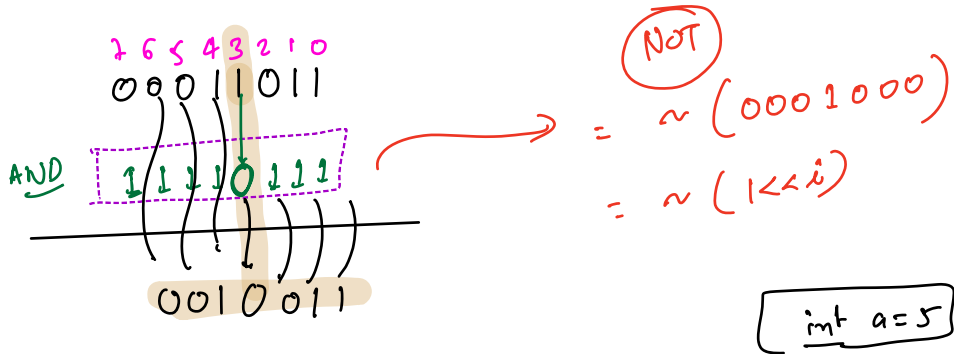
}

T.C. $\rightarrow O(1)$

Q.

Given a no. N , clear the i th bit.
(Make the i th bit 0.)

$N = 27$, $i = 3$



```
clearIthbit(int N, int i)
{
    mask = ~ (1 << i);
    return N & mask;
}
```

Q. Given a no. N , toggle i th bit.

\rightarrow $1 \rightarrow 0$
 $0 \rightarrow 1$

$N = 27$, $i = 2$

$\begin{array}{r} \text{7 6 5 4 3 2 1 0} \\ 0 \ 0 \ 0 \ 1 \ 1 \ 0 \ 1 \ 1 \\ \text{xor} \quad \boxed{0 \ 0 \ 0 \ 0 \ 0 \ 1 \ 0 \ 0} \\ \hline 0 \ 0 \ 0 \ 1 \ 1 \ 1 \ 1 \end{array}$

`toggleIthbit(int N, int i)`
{

`mask = 1 << i`

`return N ^ mask ;`

}

$N = 11011$
 $\wedge \ 00100$
 $\hline 11111 \checkmark$

Break till 8:24 AM

Q. Count no. of set bits in N.

$N=27$

c 000000000000000011011

ans = 4

if int \rightarrow 32 bits
long \rightarrow 64 bits.

```
cnt = 0
for (int i = 0; i < 32; i++)
{
    if (checkBit(N, i))
        cnt++;
}
return cnt.
```

$N=27$

c 000000000000000011011

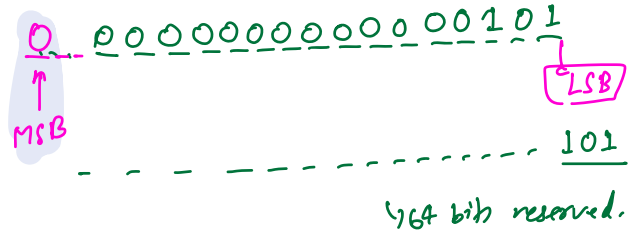
ans = 4

```
cnt = 0
while (N > 0)
{
    if (N & 1 == 1)
        cnt++;
    N = N >> 1;
}
return cnt.
```

Negative nos.

int $a = 5$;
↓
32 bits.

```
long a = 5;
```



(MSB is reserved as a sign bit.)

↳ MSB is 1 if it is a negative no.
MSB is 0 if it is a positive no.

8 bit nos.

$$0000\ 0101 = 5$$

(+) $1000\ 0101$ $\neq -5$
 $\neq 0$ $= 0$

How to represent negative no.?

"2's complement"

To get binary representⁿ of -N

- 1. Get Binary representⁿ of N.
- 2. Invert the no. (flip all bits)
- 3. Add 1.

Binary rep. of -5

0000 0101 = 5

0000 0101

flip ↙

1111 1010

+ 1

1111 1011 : -5

(Add) 0000 0101 + 5

0000 0000

$$\begin{array}{r}
 3: \quad 00000011 \\
 \quad 1111100 \\
 \quad + 1 \\
 \hline
 \quad 1111101 = -3
 \end{array}$$

$$\begin{array}{r}
 10: \quad 00001010 \\
 \quad 11110101 \\
 \quad + 1 \\
 \hline
 \quad 11110110 = -10
 \end{array}$$

Binary to Decimal

$$\begin{array}{c}
 7 \ 6 \ 5 \ 4 \ 3 \ 2 \ 1 \ 0 \\
 00101010
 \end{array}
 : 2^5 + 2^3 + 2^1$$

$$\begin{array}{c}
 7 \ 6 \ 5 \ 4 \ 3 \ 2 \ 1 \ 0 \\
 10101111
 \end{array}
 : 2^7 + 2^5 + 2^3 + 2^2 + 2^1 \quad (\text{up till now}).$$

$$\begin{array}{c}
 7 \ 6 \ 5 \ 4 \ 3 \ 2 \ 1 \ 0 \\
 10101111
 \end{array}
 = -2^7 + 2^5 + 2^3 + 2^2 + 2^1$$

place value = -2⁷

sign bits

`int a = 5`

$$\begin{aligned}
 & -128 + 32 + 8 + 4 + 2 \\
 & = -128 + 32 + 14 \\
 & = -128 + 46 \\
 & = -82
 \end{aligned}$$

4 bit nos.

If 4 bits,
how many nos. can
you represent?

— — — —
 $2^4 \text{ nos.} = 16 \text{ nos.}$

If sign is not considered,

0000 → 0

1111 → 15

0000	:	0
0001	:	1
0010	:	2
0011	:	3
0100	:	4
0101	:	5
0110	:	6
0111	:	7
1000	:	-8
1001	:	-7
1010	:	-6
1011	:	-5
1100	:	-4
1101	:	-3
1110	:	-2
<u>1111</u>	:	<u>-1</u>

[Min no. = -8
Max no. = 7]

→ $-2^3 + 2^2 + 2^1 + 2^0 = -1$

sign

— — —

→

0 to 15?

or -8 to 7?

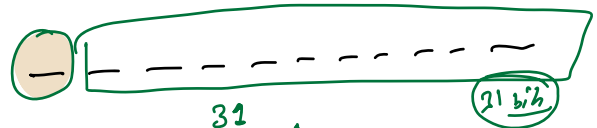
----- : 0-15 if unsigned
1

: -8 to 7 if signed.

int a=5;

a = -3 ✓

[int no. → 32 bits]
compiler assumes that it is signed.



Max: $2^{31} - 1$

Min: -2^{31}

unsigned int a=5;

a = -3; ? ~~XX~~
error

[32 bits
No sign bit]

Max: $2^{32} - 1$

Min: 0

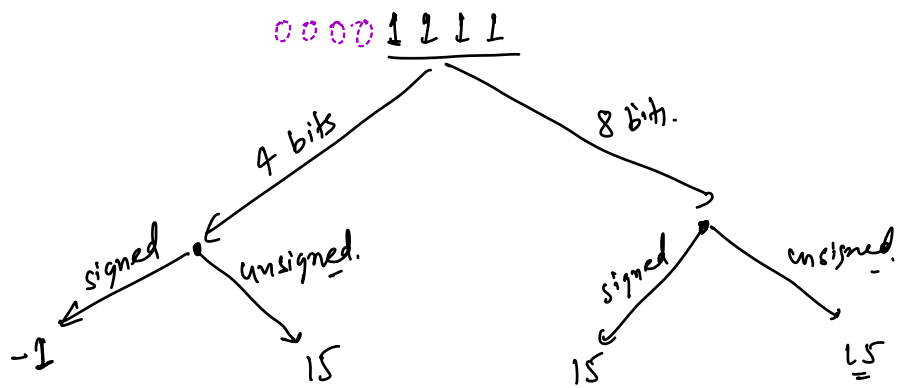
1 0 0 0 0 0 0 0

3 2 1 0
1 1 1 1

$$= 2^4 - 1 = \underline{15}$$

$$2^3 + 2^2 + 2^1 + 2^0 = \frac{1 * (2^n - 1)}{2 - 1}$$

$$= \underline{2^n - 1}$$



Range of Data types

		Min	Max
int (32 bits)	signed.	-2^{31}	$2^{31} - 1$
	unsigned	0	$2^{32} - 1$
long (64 bit)	<u>Signed.</u>	-2^{63}	$2^{63} - 1$
	unsigned.	0	$2^{64} - 1$

Constraint

int $a = 10^5$
int $b = 10^6$

int $c = a * b$

! Overflow

long $c = a * b$ ~~xxx~~

long $c = (\text{long}) a * b$

$1 \leq N \leq 10^5$

$1 \leq \text{arr}[i] \leq 10^6$

Sum of nos. in the array.

```
int long sum = 0  
for (int i = 0; i < n; i++)  
{  
    sum += arr[i]  
}  
return sum
```

32bit

$10^6 \quad 10^6 \quad 10^6 \quad 10^6 \quad \dots$
 $10^6 * 10^5$

Max possible sum value

$= 10^{11}$

10^{11} cannot be stored in int

Range of data types & Constraints.

Bit masking:

$$10 = 2^{\log_2 10}$$

$$= \left(2^{\log_2 10} \right)^{10}$$

$$= 2^{\log_2 10 \times 10}$$

$$= 2^{3.1 \times 10}$$

$$= 2^{31}$$

3.

5 & 2

→ Bitwise operator

32 bit

84 bit

int a = 5;

0 : 0 0 0 0 0

→ 0th value

1 : 0 0 0 0 1

→ 1st

2 : 0 0 0 1 0

→ 2nd

3 : 0 0 0 1 1

→ 3rd value

4 :

$c_1 \in \phi$ $mai = 1$
 $c_2 \in \phi$ $mai = 2$
 ϕ
 0

1 1 1 2 3 4 7