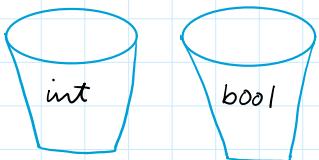


Jo store karega unn salika datatype same rahega.



Creation

Variable →

↳ int a;
↳ int a=5;

garbage value
a

When we talk about array, we can store multiple items

If we want to store 1 Lakh no's then

we make 1 Lakh variables ✗

we make 1 Lakh integer ✗

we have to write 1 Lakh line of code ✗

Let's make array which can store 1 Lakh numbers.

Array Creation

Syntax

int arr [10000];

datatype ↓
Array name Size

int a[5];

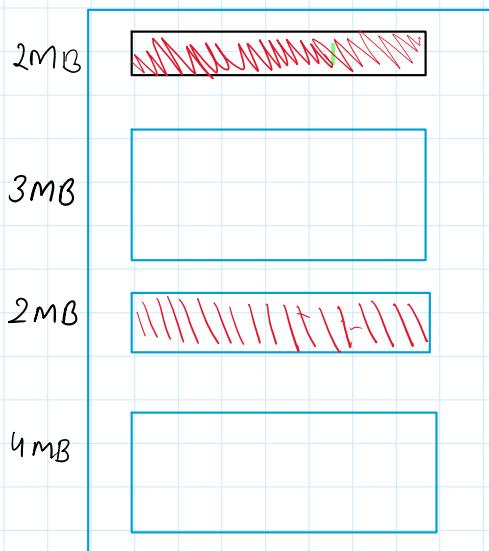


Similar type of element store

some type of
element store

↓

Contiguous → Contiguous Memory Lab.



Q) Create an array of 5MB Space

$$\text{Empty/Free Space} = 3 + 4 = 7 \text{ MB}$$

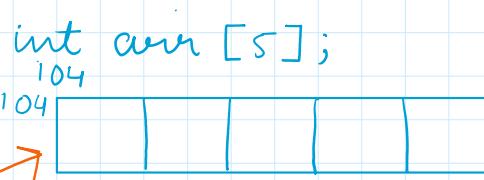
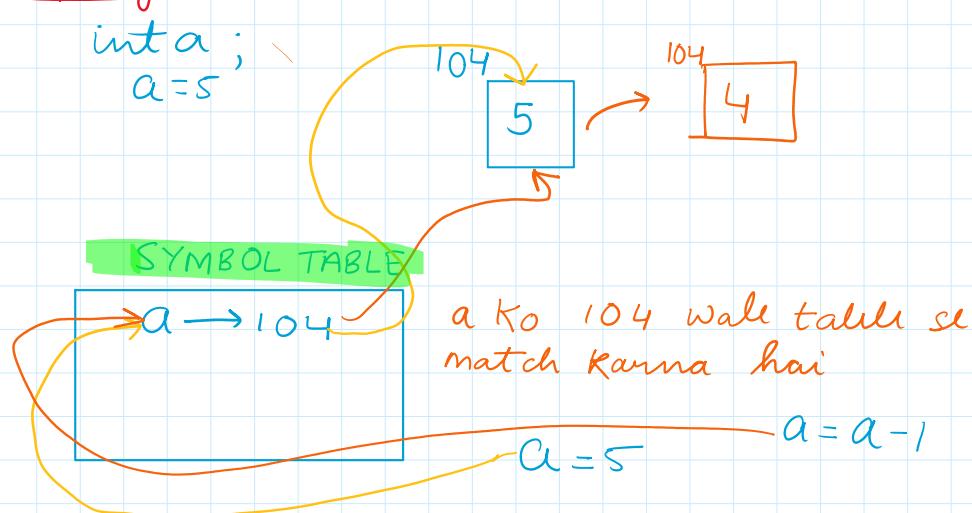
But the array can't be stored because the available required space is not contiguous.

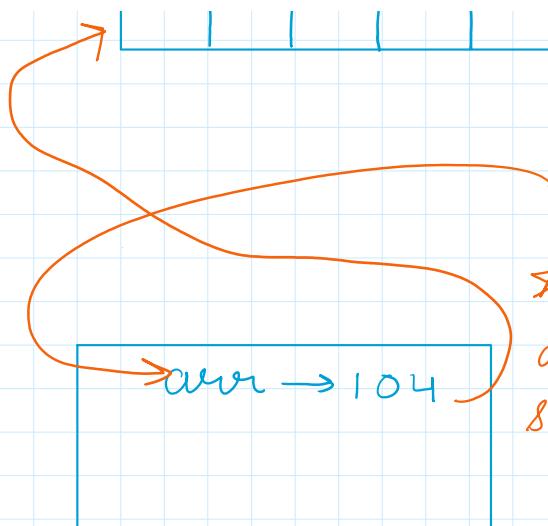
* GALAT

int a; 
memory block name 'a' created

Only one thing is attached with memory
i.e address.

Reality





★ Memory block se koi name attached nahi hota ush sirf saamne ke liye bolte hai

Create

- 10 bool → bool arr [10];
- 15 char → Char arr [15];
- 1 int → int arr [5];
- String

int arr [5]; → Garbage value

Initialisation

int arr [5] = {10, 20, 30, 40, 50};

10	20	30	40	50
----	----	----	----	----

int arr [5] = {10, 20};

10	20	0	0	0
----	----	---	---	---

Imp:- Agar size se kam value de rakhi hai toh Rest bachdi hui space par 0 aayayega.

int arr [4] = {10};

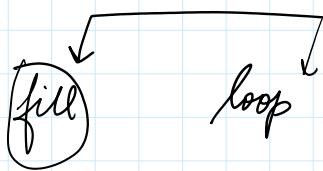
0	0	0	0
---	---	---	---

int arr [3] = {10}

10	0	0
----	---	---

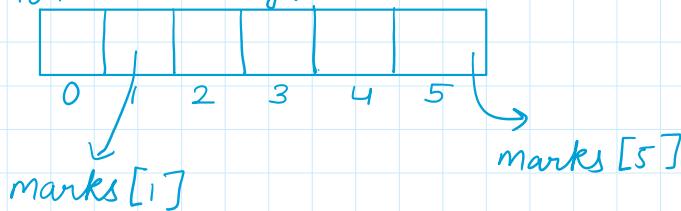
arr → Each member ko → 10 initialise

`arr` → Each member $k_0 \rightarrow 10$ Initialise



Indexing

`int marks[6];`
104 → Starting / Base address



arr size → 6
Index
↓
0-5

Important

arr size → n
index → {0 - (n-1)}

`int arr[3] = {10, 20, 30};`

`Cout << arr[4];`

→ Garbage Value
→ Executes memory depends on compiler to compiler

Fill Function

Segment table

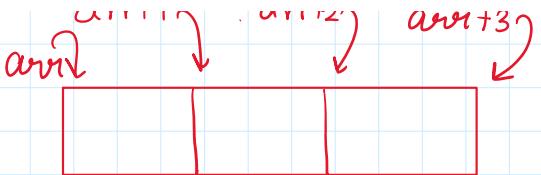
`arr → 104`

104
`arr`

`int arr[3];`
fill (`arr, arr+3, 23`)

Starting address
Ending address
Value of Initialisation

`arr` $arr+1$ $arr+2$ $arr+3$



After running all value should be 23

```

int main() {
    // int arr[10];
    // //error
    // //int brr[]

    int crr[] = {10,20,40,70};           ||
    int drr[4] = {0};

    // //cout << drr[0] << endl;
    // cout << crr[4] << endl;

    int arr[4];
    fill(arr,arr+4,101);
    cout << arr[0] << " " << arr[1] << " " << arr[2]
    << " " << arr[3] << endl;
}
  
```

Input

int a → input → cin >> a;

int arr [5]

$\xrightarrow{5 \text{ Blocks}}$ $arr[0] \rightarrow cin >> arr[0];$ }
 :
 :
 $arr[4] \rightarrow cin >> arr[4];$

Loop

$n=5$

```

for (int i=0 ; i < 5 ; i++) {
    cin >> arr[i];
}
  
```

Why array?

↳ Jab bahut saare same type of data store karna hai toh array.

Because 100 No. ke liye 100 variables

... our own array.

Because 100 No. ke liye 100 Variable
name hum create nahi karenge.

* int size;

cin \rightarrow size;

int arr [size];

Bad Practice

why \rightarrow Study in
Dynamic Memo. allocation

Doubt

\hookrightarrow Size

int a = 5;

Size of (a) \rightarrow 0 \downarrow \rightarrow 4

arr \rightarrow

10	20	30	40
----	----	----	----

$$\frac{\text{Size of (arr)}}{\text{Size of (int)}} = \frac{16}{4} = 4$$

Number of Blocks of array

~~Size~~ It doesn't tell no. of elements
present in an array

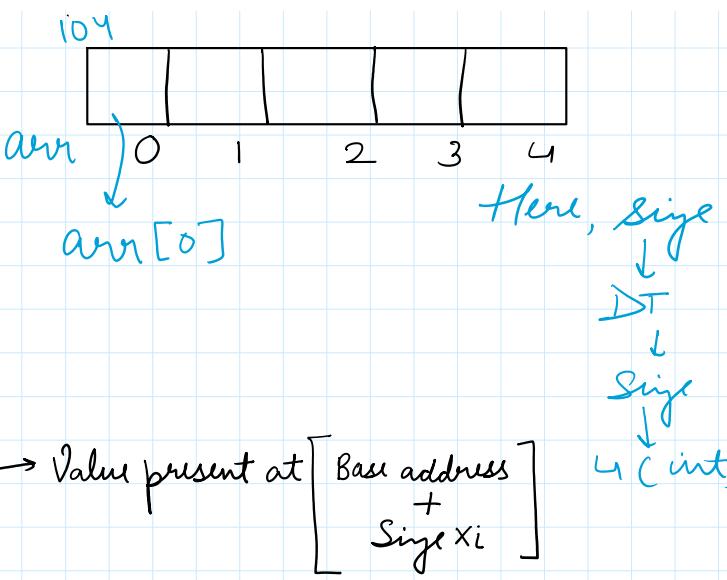
We don't have any formula

Only find by using another variable
separately.

Q Why indexing start from 0?

104





$$arr[0] = \dots [104 + 4 \times 0]$$

Gsp.
 $\text{arr}[i] \rightarrow$ value at $(\text{arr} + (\text{size} \times i))$
 $i[\text{arr}] \rightarrow$ sun le run hoga
 dona chalte hai
 $\text{arr}[i] \rightarrow i[\text{arr}]$ why Equal.
 Behind some yeh chalte raha hota
 $*(\text{arr} + i)$
~~at arr~~
 $*(\text{arr})$
 Value Present at address

Case-1
 $\text{arr}[i] \xrightarrow{\text{Compiler understands it as}} \star(\text{arr} + i)$
 Value present at $(\text{arr} + i)$

$i = 3$
 $\text{val present at } (\text{arr} + 3)$
 $(\text{arr} + 3) \xrightarrow{\downarrow} (\text{int size} \times 3)$
 (114)

Case-2
 $i[\text{arr}] \xrightarrow{\text{Value present at }} \star(i[\text{arr}])$
 $(\text{arr} + 104)$

The diagram illustrates the equivalence between two ways of accessing array elements. On the left, $\text{arr}[i]$ is shown above a double-headed vertical arrow pointing down to $i[\text{arr}]$. To the right of the arrow, the expression $*(\&\text{arr} + i)$ is shown above the text "value at (1112)". Above this, the value $l = 2$ is written next to a circled i . At the bottom right, the value $\text{ARR} = 1010$ is written.

#functions in an array

```
int main() {  
    int a = 5;  
    int c = solve(a);  
}
```

int solve (int a) {

```
int main ()  
{  
    int arr[5] = {0};  
    solve (arr, n);  
}  
  
int  
solve (arr, n)  
{  
    singl  
    solve (arr, singl);  
}
```

```
int solve(int arr[], int size) {  
    // Four parallel lines  
    // Three parallel lines  
    // One line  
    // One line  
}
```

```

int main() {
    // int arr[10];
    // //error
    // //int brr[]

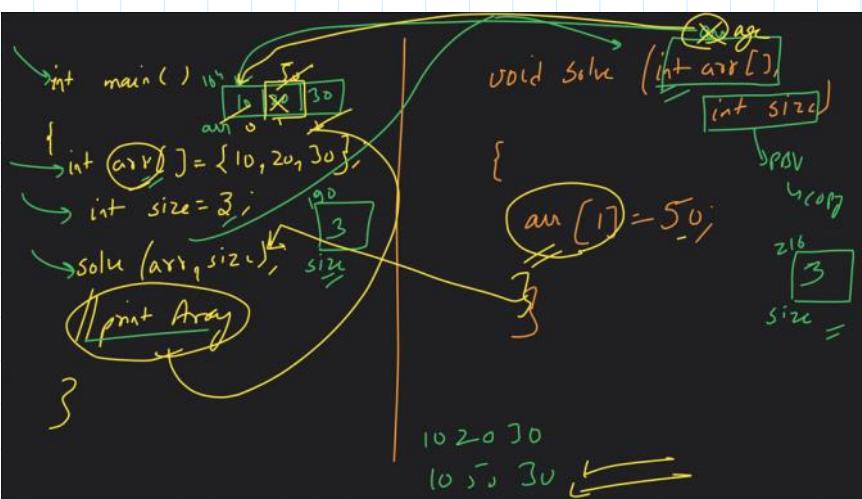
    // int crr[] = {10,20,40,70};
    // int drr[4] = {0};

    // cout << drr[0] << endl;
    // cout << crr[4] << endl;

    int arr[4];
    fill(arr,arr+4,101);
    cout << arr[0] << " " << arr[1] << " " << arr[2]
    << " " << arr[3] << endl;
}

```

Speed



```

int main ()
{
    int arr[4]={0};
    int size=4;
    solve(arr, size);

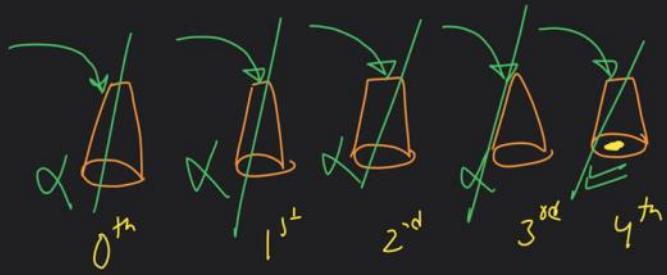
    #for(i=0→i<n)
    {
        cout << arr[i];
    }
    return 0;
}

void solve (int num[], int n)
{
    for (i=0→ i<n)
    {
        num[i]=num[i]+10;
    }
    return;
}

```

Array → Pass By Reference

① Algorithm → Linear Search



ek ki baad ek par jana issi linear search kehti hai

array ↦ int arr[5] → {10, 20, 30, 40, 50}

arr	0	1	2	3	4
	10	20	30	40	50

Find target element present or not
if found return true → element found

& if not present return → element not found.

```

int main() {

    int arr[] = {10, 20, 30, 40, 50};
    int size = 5;
    int target = 55;

    bool ans = findTarget(arr, size,
    target);
    cout << "ans: " << ans << endl;
    // solve1(arr, size);
    // print(arr, size);
}

```

```

bool findTarget(int arr[], int size, int target) {
    //traverse the entire array
    for(int i=0; i<size; i++) {

        int currentElement = arr[i];

        if(currentElement == target) {
            //found
            return true;
        }
    }
    //agar aap yaha tk pohoch gye toh
    //iska mtlb poora loop chal chuka hai
    //iska matlab poore loop me kahin bhi target nahi mila
    //iska mtlb element not found
    //iskat matlab return false;
    return false;
}

```

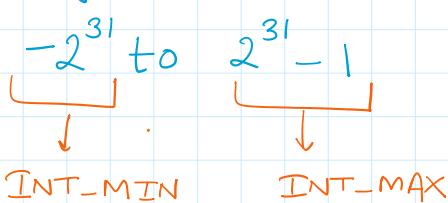
```

bool checkTarget (arr, size, target)
{
    for (i=0; i<size; i++)
    {
        if (arr [i] == target)
        {
            return true;
        }
    }
    return false;
}

```

Integer
 ↓
 Range

Range



arr	13	21	14	17
	0	1	2	3

To Find

$\rightarrow \text{Max}^M \text{ No.} \rightarrow \text{ans} = \text{INT_MIN}$

$\rightarrow \text{Min}^M \text{ No.} \rightarrow \text{ans} = \text{INT_MAX}$

$\text{arr}[i] > \text{ans} \rightarrow \text{ans} = \text{arr}[i]$
 $\text{arr}[i] < \text{ans} \rightarrow \text{ignore}$

```
int to findMax(int arr[], int size) {
    int maxAns = INT_MIN;
    for(int i=0; i<size; i++) {
        maxAns = max(maxAns, arr[i]);
        // if(arr[i] > maxAns) {
        //     maxAns = arr[i];
        // }
    }
}
```

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If Rank
bhi compare
Kar Sakti
hai

To use INT_MIN & INT_MAX → Use ~~#include~~ `<limits.h>`

```

int findMax(arr n)
{
    int maxAns = INT_MIN;
    for (int i=0; i<n; i++)
    {
        maxAns = max(maxAns, arr[i]);
    }
    return maxAns;
}

```

\rightarrow	13	42	55	70	88
arr	0	1	2	3	4

$\maxAns = -2^{31} \Rightarrow 5$

```

int findMax (arr , n)
{
    int maxAns = INT_MIN;
    for (i=0; i<n; i++)
    {
        if (arr [i] > maxAns)
        {
            maxAns = arr [i];
        }
    }
    return maxAns;
}

```

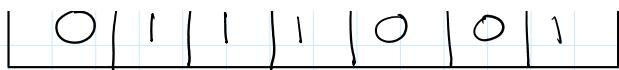
```

int findMin (arr , n)
{
    int minAns = INT_MAX;
    for (i=0; i<n; i++)
    {
        minAns = min (minAns, arr [i]);
    }
    return minAns;
}

```

Q. I/p mai array given hai
 \hookrightarrow Count \rightarrow 0's / 1's

0		1		1		0		0		1
---	--	---	--	---	--	---	--	---	--	---



Zero_Count = 0;

One_Count = 0;

```
void printZeroesAndOnes(int arr[], int n) {
    int zeroCount = 0;
    int oneCount = 0;

    //traverse Array
    for(int i=0; i<n; i++) {
        int currElement = arr[i];

        if(currElement == 0) {
            zeroCount++;
        }
        if(currElement == 1) {
            oneCount++;
        }
    }
    cout << "Total Zeroes: " << zeroCount << endl;
    cout << "Total Ones: " << oneCount << endl;
}
```

Q Extreme Printing:

i/p

10	20	30	40	50	60
----	----	----	----	----	----

o/p

10	60	20	50	40
----	----	----	----	----

2 Pointer Method

Approach

$i = 0$
$j = n + 1$

$\text{cout} \ll i;$

$i++$

$\text{cout} \ll j;$

$j--$

Stop

$j < i$

```

10
void extremePrint(int arr[], int n) {
    int i = 0;
    int j = n-1;

    while(i < j) {
        cout << arr[i] << " ";
        i++;
        cout << arr[j] << " ";
        j--;
    }
}

```

Not Work For Odd No's



Change Condition ($i \leq j$)



It will print middle element 2 times

Method-1

```

Use if (j==i) {
    cout << i << " ";
}

else {
    cout << i << " ";
    i++;
    cout << j << " ";
    j--;
}

```

Method -2

```

cout << i << " ";
i++;

if (j>i) {
    cout << arr[j] << " ";
    j--;
}

```

HOMEWORK QUESTIONS

Important

Q1) Swap int a → 5

int b → 4

↳ Various Methods

($+,-$, temp, \wedge , \vee , Swap)

0/p → a → 4

b → 5

Q2) Reverse an Array ?