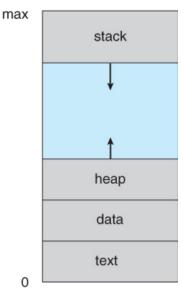
#### **Processes**

#### **Processes**

- A process is a program loaded into main memory, ready for execution.
- The OS assigns **isolated memory** to each process, ensuring no process sneaks into another's memory like a nosy neighbor. This memory is divided into four segments:

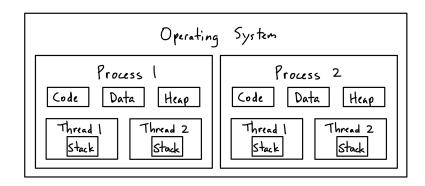


- 1. **Stack**: Contains local variables, parameters, and return addresses.
- 2. **Heap:** Used for dynamic memory allocation. (The program's storage unit.)
- 3. **Data**: Stores global variables, arrays, and structures. It's split into:
- BSS (Block Started by Symbol): For uninitialized variables.
  - **Data**: Contains initialized variables.
- 4. **Text**: The program's compiled machine code.

# • Memory Growth:

- o The stack and heap grow toward each other.
  - Stack Overflow: When the stack crosses the line into the heap.
  - Heap Overflow: When the heap pushes into the stack.

#### **Threads**



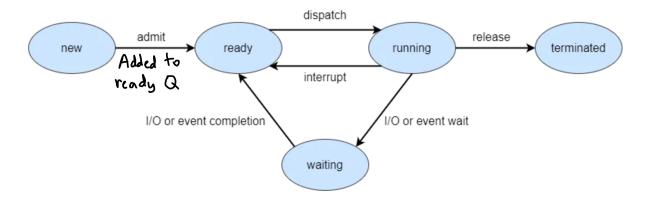
- A process contains at least one thread.
- A thread is the smallest unit of execution within a process and shares the process's memory space.
- Each thread has its own stack within the allocated memory of the process.
- Thread Levels:
  - Threads operate at both user-level and kernel-level.
  - In a multi-threaded environment, thread management can follow these models:
    - Many-to-One
    - One-to-One
    - Many-to-Many
       (Honestly, I skipped reviewing these because I know I won't remember during an interview.)

## **Process Control Block (PCB)**

- The **Process Control Block (PCB)** is a data structure used by the OS to manage process metadata. It's like the process's personal file folder, filled with everything the OS needs to know.
- Key elements of the PCB include:
  - o PID (Process ID): A unique identifier for the process.
  - Parent PID: The ID of the parent process.
  - Child PID: The ID of any child processes.
  - PC (Program Counter): The location of the next instruction to execute.
  - Memory Limits: Defines the boundaries of the process's memory.
  - Other Metadata: Includes process state, scheduling info, and I/O status.

## **Process Creation**

- A new process (**child**) is created by **forking** (using the fork() system call) from an existing process (**parent**).
  - o fork() Return Values:
    - The parent process receives the child's PID.
    - The child process receives **0**.
- All processes are created by the CPU and can exist in one of the following states:
  - 1. **New:** Just created, waiting to be admitted.
  - 2. **Ready:** Prepared to run when the CPU is available.
  - 3. Running: Actively being executed by the CPU.
  - 4. Waiting: Paused, waiting for an event (e.g., I/O operation).
  - 5. **Terminated:** Finished execution, now ready to rest in peace.



### State Transitions:

- A process in the ready state moves to running if it has the highest priority.
- o If it times out, it returns to the ready state.
  - (This timeout involves an interrupt, essentially a polite "hey CPU, your attention please!")

# **Multiprocessing and Multithreading**

## Concurrency vs. Parallelism

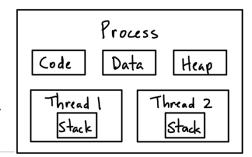
- Concurrency: A single core CPU switches between tasks (context switching).
- Parallelism: A multicore CPU executes multiple tasks simultaneously.

### Multiprocessing

- Multiprocessing involves splitting a task into multiple processes.
  - It's more stable because processes operate independently, but comes with overhead.
  - The overhead arises from context switching, where the CPU pauses the current job, stores its state, and loads the new job's state. This costs both time and memory.
- Each process is allocated its own memory space, ensuring isolation.
- Since processes cannot directly access each other's memory, they communicate via IPC (Inter-Process Communication).

## Multithreading

- Multiple threads exist within a process, each performing different tasks.
- Threads share the **heap**, **data**, **and text** segments but have separate stacks.
- Pros:
  - Less overhead in context switching.
  - No need for IPC.
- Cons:
  - Shared stack requires synchronization.
  - A problem in one thread may affect others.



# **Context Switching**

- Interrupts: Requests made to the CPU for attention due to events like:
  - I/O operations.
  - CPU usage time expiration.

o Creation of child processes.

# Mechanism:

- $\circ$  The CPU can only process one task at a time.
- An interrupt, often caused by the CPU scheduler, switches the CPU to another process.
- Overhead occurs during context switching because the CPU idles while loading a process's state into the registers.

# • Program Counter (PC) and Stack Pointer:

- o **PC:** Holds the address of the next instruction to execute.
- o Stack Pointer: Points to the largest address in the stack.