

# Advanced Algorithm Analysis and Design

## Assignment 2

Due: Nov. 11, 2024

1. Write down the main steps of proving the NP-Completeness of a problem.
2. Given a graph, a dominating set is a subset of vertices such that any vertex not in this set is adjacent to at least one vertex in this set. The dominating set problem is to check whether a given graph has a dominating set of size at most  $k$ .
  - 2.1 Prove that the dominating set problem is in NP.
  - 2.2 Prove that the dominating set problem is NP-hard.
3. Prove that: if we can check whether a graph has a clique (a complete graph) of size  $k$  in polynomial time then we can also find a clique of size  $k$  in polynomial time.
4. A graph is called a 2-plex if each vertex in the graph is not adjacent to at most one other vertex. Prove that it is NP-complete to check whether an input graph has a subgraph of at least  $k$  vertices that is a 2-plex.
5. In the multiway cut problem, we are given a undirected graph  $G=(V,E)$  and some special vertices in  $V$  (called terminals). The problem asks us to delete the minimum number of edges from the graph such that no pair of terminals is connected. Please give a 2-approximation algorithm for this problem.