

## Definitions

**Tree:** connected graph with no cycles.

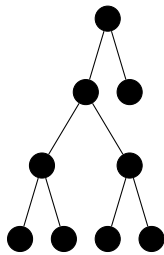
**Rooted:** if one node identified as root.

**$m$ -ary:** if every parent has  $m$  children.

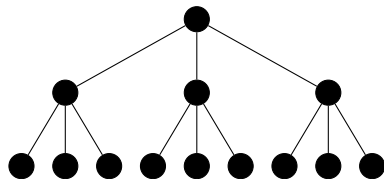
**Height:** maximum length of path to a leaf.

**Leaf:** is a node with no children.

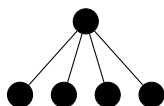
## Examples



Binary (2-ary) tree of height 3.



Ternary (3-ary) tree of height 2.



4-ary tree of height 1.

## Theorem

$m$ -ary tree with  $l$  leaves has height at least  $\log_m l$ .

## Log

$$b^a = c \Leftrightarrow \log_b c = a$$

$$10^2 = 100 \Leftrightarrow \log_{10} 100 = 2$$

## Rationale

Maximum leaves is  $l \leq m^h$ .

## Proof of theorem

$$h \geq \log_m l \Leftrightarrow m^h \geq m^{\log_m l} \Leftrightarrow m^h \geq l$$

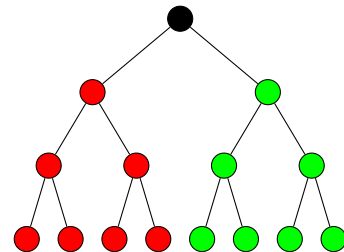
By induction on  $h$ :

$h = 0$ :  $m^0 = 1$ ,  $l = 1$ ,  $1 \geq 1$ .

$h = n$ : Assume true for  $h = n - 1$ . Removing root gives  $m$  trees with maximum  $m^{n-1}$  leaves each. In total, we get maximum of  $m(m^{n-1}) = m^n$  leaves.

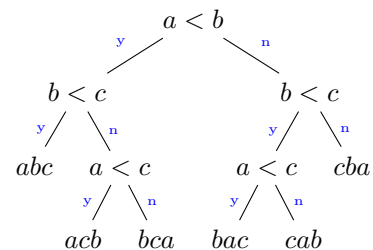
True for  $h = 0$ . Therefore true for  $h = 0 + 1 = 1$ . Therefore true for  $h = 1 + 1 = 2$ . Therefore true for  $h = 2 + 1 = 3$ . And so on.

## Example



## Application

To sort a list of  $n$  items we have to choose between  $n!$  permutations. So, in the worst case, using binary decisions we make at least  $\log_2 n!$  comparisons.



$$\log_2 3! = \log_2 6 > 2$$