

Lezione 13

venerdì 10 novembre 2023 11:11

PSEUDO RANDOM PERMUTATION

Recap: PRF implies EFFICIENTLY everything in symmetric crypto.
Sometimes need PRP (i.e. modes of operations).

From theory: OWF \Rightarrow PRF, (\Rightarrow PRP) well see Today.

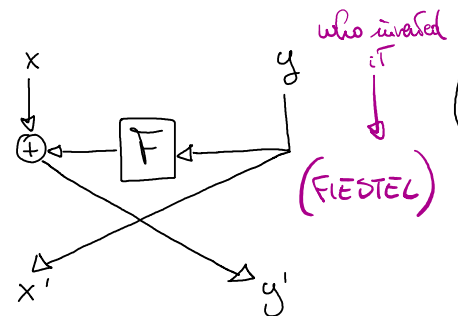
From practice: AES is a PRP. Make some ad-hoc PRF (only heuristic) and then use construction PRF \Rightarrow PRP.
there is a famous one so it's pseudorandom and invertible.

we can't prove much

LUBY-RACKOFF construction

$$F: \{0,1\}^m \rightarrow \{0,1\}^m \text{ (not invertible)}$$

$$\text{construction: } \Psi_F(x, y) = (y, x \oplus F(y)) = (x', y')$$



$$\text{now it's } \{0,1\}^{2m} \rightarrow \{0,1\}^{2m}$$

and we claim it's invertible even if F it's not:

$$\Psi_F^{-1}(x', y') = (F(x') \oplus y', x') = (x, y)$$

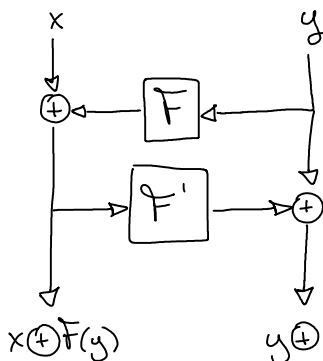
• Question: is it pseudorandom? Assuming F it is.

NO

2° part of the output exactly the 1° part of the input

We can use this construction multiple times try 2 ROUNDS

If this happens we can say for sure that it's FIESTEL and not R



⚠ we use different functions

$$\Psi_{F, F'}(x, y) = \Psi_{F'}(\Psi_F(x, y)) \text{ still invertible}$$

• Question: Is it a PRP? **NO**

We can break it with 2 queries: (x, y) and (x', y') $x \neq x'$
so we get $\Psi_{F, F'}(x, y)$ and $\Psi_{F, F'}(x', y')$

$$\text{If we XOR them: } \Psi_{F, F'}(x, y) \oplus \Psi_{F, F'}(x', y') = (x \oplus x', \text{---})$$

so we can check -

LUCKY - RACKOFF: 3-ROUNDS WORKS

then

Let F be a PRF.

Then $\Psi_F[3] = \Psi_{F_{k_3}}(\Psi_{F_{k_2}}(\Psi_{F_{k_1}}(x, y)))$ it's a PRP
for $k_1, k_2, k_3 \leftarrow \{0, 1\}^L$

Practice: F replaced by heuristic construct (S-BOXES) # of rounds: 18 times
 k_1, \dots, k_{18} all derived from single k

Two steps:

1. If the queries have all y unique then 2-ROUND FIESTEL is secure.
2. the first round makes all queries "y-unique" (all y unique)

LEMMA: For every UNBOUNDED distinguisher asking $q(L) = \text{poly}(L)$ queries, the following are close: STATISTICALLY

→ $S: F, F' \leftarrow \mathcal{R}(L, m, m)$
and answer (x, y) with $\Psi_{F'}(\Psi_F(x, y))$

→ $R: R \leftarrow \mathcal{R}(L, 2m, 2m)$
and answer (x, y) with $R(x, y)$

so long as $(x_1, y_1), \dots, (x_q, y_q)$ are s.t. y-unique
 $y_i \neq y_j \quad \forall i \neq j$.

proof: Hybrid argument: $H_i(L)$ that answers the first i queries using S and all other queries using R . We already know: $H_0 \equiv R$ and $H_q \equiv S$
(we need to prove $H_i \approx S$).

The first i outputs of H_i : 2 ROUND FIESTEL

$(x_j \oplus F(y_j), y_j \oplus F'(x_j \oplus F(y_j)))_{j=1}^{i-1}$
 $x_i \oplus F(y_i), y_i \oplus F'(x_i \oplus F(y_i)), \dots$
at a certain point he will answer R

By the y-UNIQUENESS, $F(y_i)$ is random and independent of the rest:

By the γ -UNIQUENESS, $F(y_i)$ is random and independent of the rest:
so we can write

$$(x_j \oplus F(y_j), y_j \oplus F'(x_j \oplus F(y_j)))_{j=1}^{i-1},$$

$$(x_i \oplus r, y_i \oplus F'(x_i \oplus r)) \dots, \dots, \dots$$

What's now the probability that $x_i \oplus r = x_j \oplus F(y_j)$?

For some $j < i$ the probability $\leq (i-1) \cdot 2^{-m}$. Assuming it doesn't happen:
this becomes

$((x_i \oplus r), (y_i \oplus r'))$ which is $\equiv_R (x_i, y_i)$, uniform

By the lemma of game playing the distance between

$$H_i \text{ and } H_{i-1} \leq \underbrace{q(1)}_{\text{to be precise it should be } (i-1)} \cdot 2^{-m} = \text{negl}(1).$$

to be precise it should be $(i-1)$

$$\Rightarrow \Delta(S, R) \leq \sum (i-1) \cdot 2^{-m} \leq q^2 \cdot 2^{-m} = \text{negl}(1)$$

\square lemma

proof (Thm)

We consider a bunch of experiments.

$$\begin{array}{l} \text{2 exp } \left\{ \begin{array}{l} \cdot T: (x, y) \mapsto \Upsilon_{F_{K_3}}(\Upsilon_{F_{K_2}}(\Upsilon_{F_{K_1}}(x, y))) \quad \text{real world.} \\ \cdot S: (x, y) \mapsto \Upsilon_{F, F'}(\Upsilon_F(x, y)) \quad \begin{array}{l} \text{Truly} \\ \text{random} \\ \text{function} \end{array} : F, F', F'' \leftarrow \$R(1, m, m) \end{array} \right. \\ \\ \cdot R: (x, y) \mapsto R(x, y) \quad R \leftarrow \$R(1, 2m, 2m) \quad \text{RANDOM FUNCTION} \\ \\ \cdot P: (x, y) \mapsto P(x, y) \quad P \leftarrow \$P(1, 2m, 2m) \quad \text{RANDOM PERTURBATION} \end{array}$$

ideal world
RANDOM PERTURBATION

We want to demonstrate that they are all indistinguishable.