

# CS330 - Computer Organization and Assembly Language Programming

## Lecture 15

-Arrays, Heterogeneous Data Structures,  
Combining Control and Data-

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# Agenda

- Array Allocation and Access
  - ❖ One-dimensional
  - ❖ Multi-dimensional (nested)
  - ❖ Multi-level
- Heterogeneous Data Structures
  - ❖ Allocation
  - ❖ Access
  - ❖ Alignment

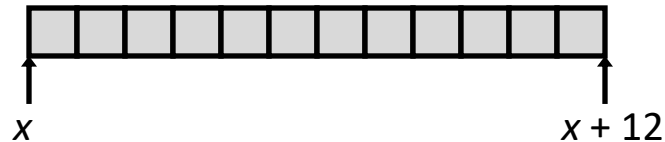
# Array Allocation

- Basic Principle

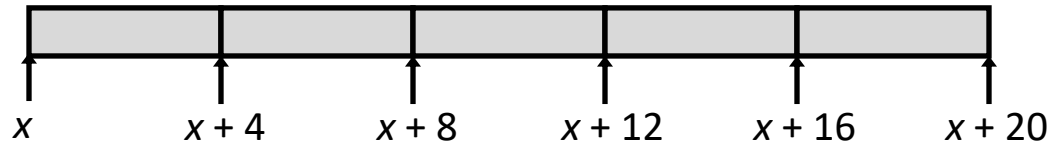
$T$  **A**[ $L$ ];

- Array of data type  $T$  and length  $L$
- Contiguously allocated region of  $L * \text{sizeof}(T)$  bytes in memory

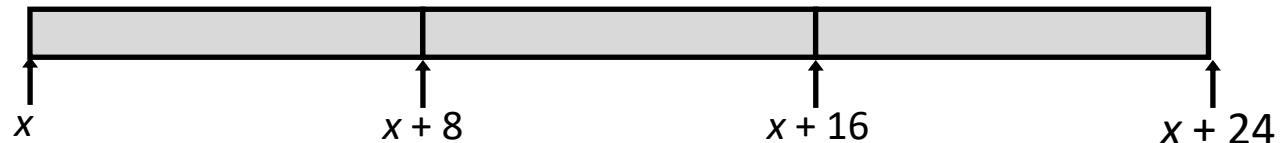
`char string[12];`



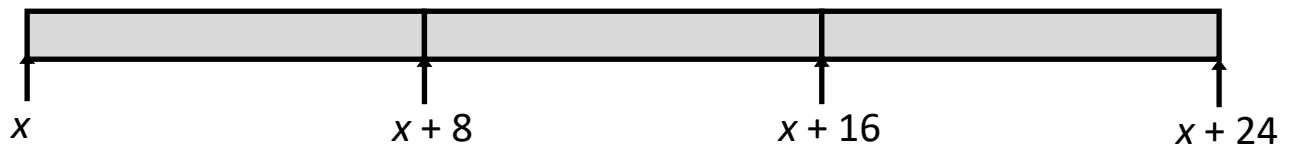
`int val[5];`



`double a[3];`



`char *p[3];`



# recall pointers

- The unary operators & and \* allow generation and dereferencing of pointers
- An expression Expr  $\rightarrow$  denoting some object;
  - &Expr is a pointer giving the address of the object
- An expression AExpr denoting an address;
  - \*AExpr gives the value at that address

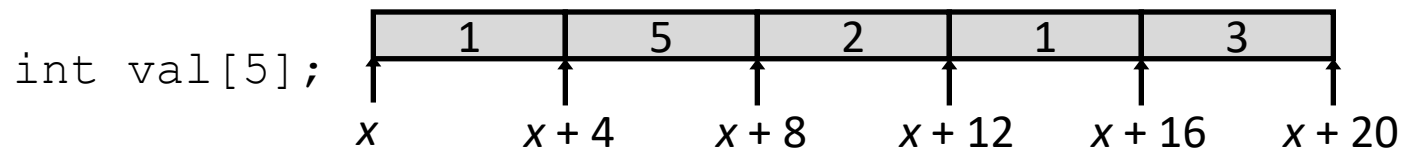
# Array Access

- Basic Principle

$T$  **A**[ $L$ ] ;

- Array of data type  $T$  and length  $L$

- Identifier **A** can be used as a pointer to array element 0: Type  $T^*$



- Reference Type Value

<code>val[4]</code>	<code>int</code>	3
<code>val</code>	<code>int *</code>	$x$
<code>val+1</code>	<code>int *</code>	$x+4$
<code>&amp;val[2]</code>	<code>int *</code>	$x+8$
<code>val[5]</code>	<code>int</code>	??
<code>*(val+1)</code>	<code>int</code>	5
<code>val + i</code>	<code>int *</code>	$x+4i$

# Pointer Arithmetic

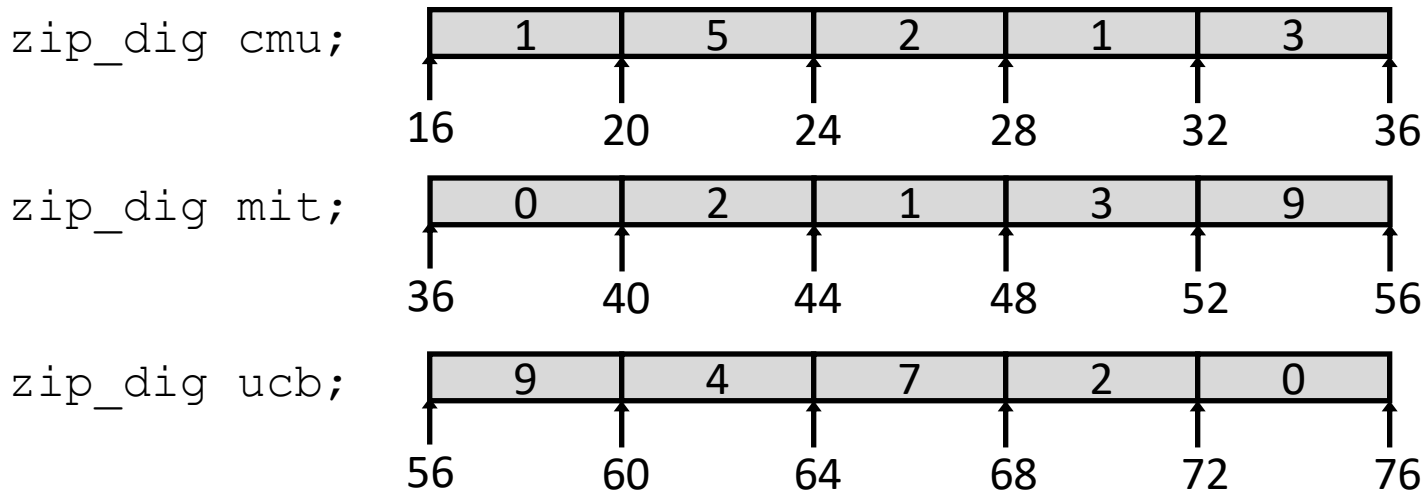
- You can refer to these resources if you want to learn more about pointer arithmetic
- [https://www.learn-c.org/en/Pointer\\_Arithmetics](https://www.learn-c.org/en/Pointer_Arithmetics)
- [https://www.tutorialspoint.com/cprogramming/c\\_pointer\\_arithmetic.htm](https://www.tutorialspoint.com/cprogramming/c_pointer_arithmetic.htm)
- <https://overiq.com/c-programming-101/pointer-arithmetic-in-c/>

- The memory referencing instruction of x86\_64 are designed to simplify array access.
- For ex: suppose E is an array of values of type int and we wish to evaluate E[i], where the **address of E** is stored in register **%rdx** and **i** is stored in register **%rcx**. Then the instruction is:
  - `movl (%rdx, %rcx,4), %eax`
- which performs the address computation  $x_E + 4i$ , read that memory location, and copy the result to register `%eax`

# Array Example

```
#define ZLEN 5
typedef int zip_dig[ZLEN];

zip_dig cmu = { 1, 5, 2, 1, 3 };
zip_dig mit = { 0, 2, 1, 3, 9 };
zip_dig ucb = { 9, 4, 7, 2, 0 };
```

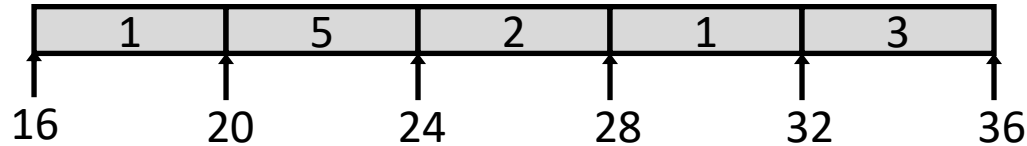


- Declaration “zip\_dig cmu” equivalent to “int cmu[5]”
- Example arrays were allocated in successive 20 byte blocks
  - Not guaranteed to happen in general



# Array Accessing Example

zip\_dig cmu;



```
int get_digit
(zip_dig z, int digit)
{
    return z[digit];
}
```

IA32

```
# %rdi = z
# %rsi = digit
movl (%rdi,%rsi,4), %eax # z[digit]
```

- Register `%rdi` contains starting address of array
- Register `%rsi` contains array index
- Desired digit at  $\text{\%rdi} + 4 * \text{\%rsi}$
- Use memory reference  $(\text{\%rdi}, \text{\%rsi}, 4)$

# Array Loop Example

```
void zincr(zip_dig z) {  
    size_t i;  
    for (i = 0; i < ZLEN; i++)  
        z[i]++;  
}
```

```
# %rdi = z  
movl    $0, %eax          # i = 0  
jmp     .L3               # goto middle  
.L4:                      # loop:  
    addl    $1, (%rdi,%rax,4) # z[i]++  
    addq    $1, %rax        # i++  
.L3:                      # middle  
    cmpq    $4, %rax        # i:4  
    jbe     .L4             # if <=, goto loop  
rep; ret
```

# Multidimensional Arrays


```
int A[5] [3];
```

is equivalent to the declaration;

```
typedef int row3_t;  
row3_t A[5];
```

# Storing in the memory

- The array elements are ordered in memory in ***row-major*** order;
  - all elements of row 0 ( $A[0][i]$ )
  - followed by all elements of row 1 ( $A[1][i]$ )
  - ....
  - ...
- To access elements of multidimensional arrays, the compiler generates code to compute the offset of the desired element and then uses one of the mov instructions with the start of the array as the base address.

# Multidimensional (Nested) Arrays

- Declaration

$T \text{ } \mathbf{A}[R][C];$

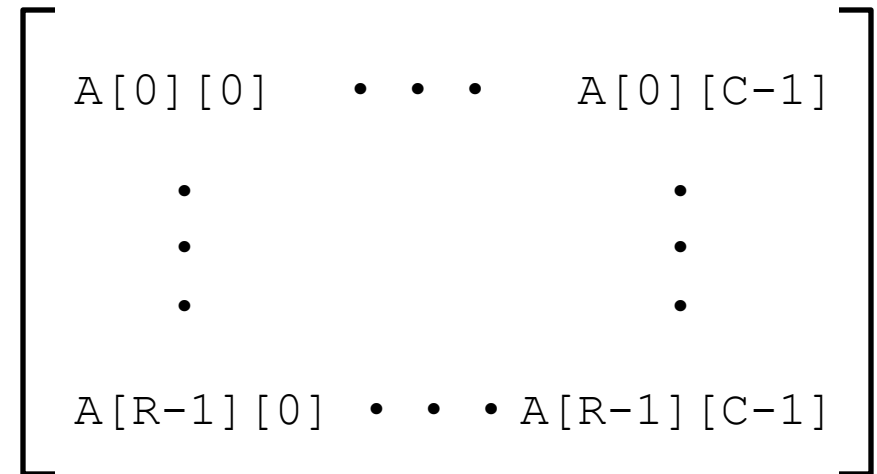
- 2D array of data type  $T$
- $R$  rows,  $C$  columns
- Type  $T$  element requires  $K$  bytes

- Array Size

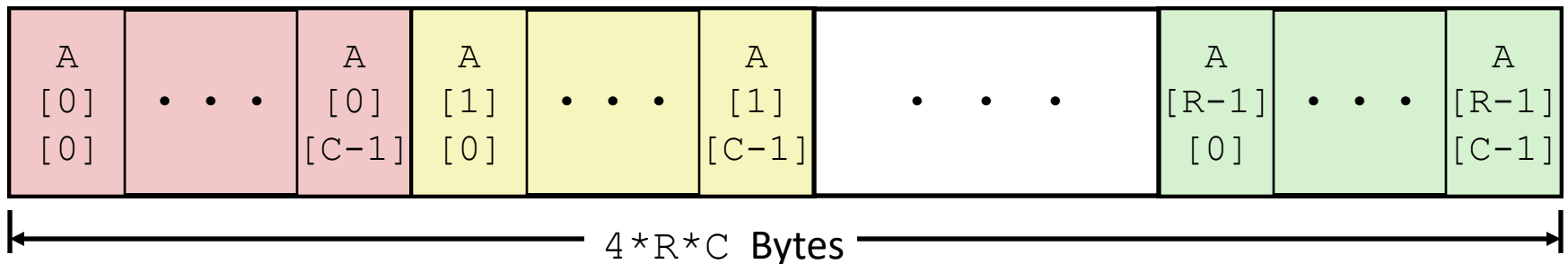
- $R * C * K$  bytes

- Arrangement

- Row-Major Ordering



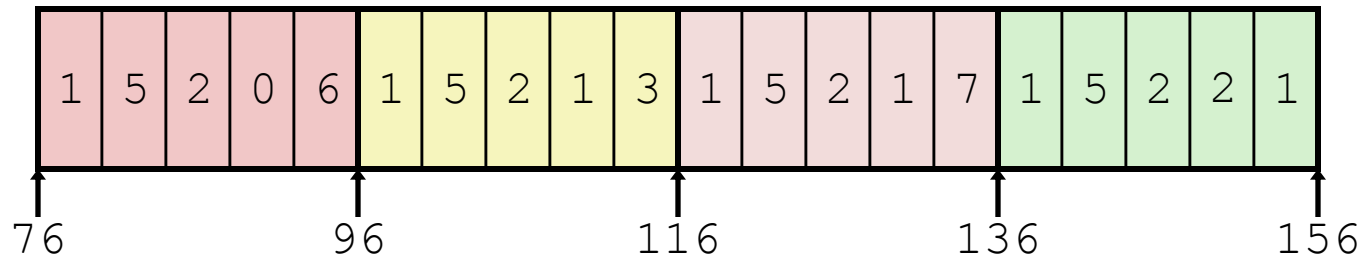
`int A[R][C];`



# Nested Array Example

```
#define PCOUNT 4
zip_dig pgh[PCOUNT] =
    {{1, 5, 2, 0, 6},
     {1, 5, 2, 1, 3 },
     {1, 5, 2, 1, 7 },
     {1, 5, 2, 2, 1 }};
```

zip\_dig  
pgh[4];

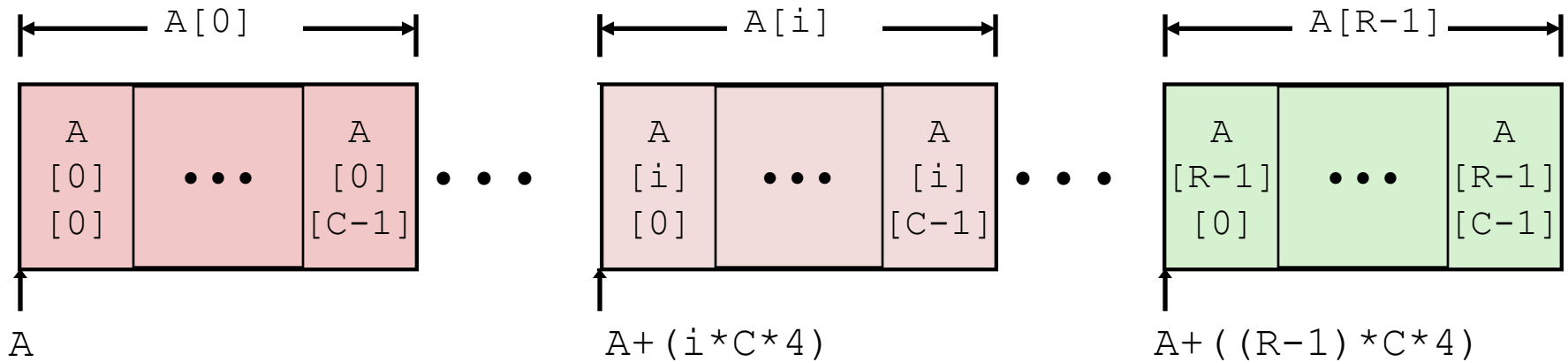


- “zip\_dig pgh[4]” equivalent to “int pgh[4][5]”
  - Variable **pgh**: array of 4 elements, allocated contiguously
  - Each element is an array of 5 **int**’s, allocated contiguously
- “Row-Major” ordering of all elements in memory

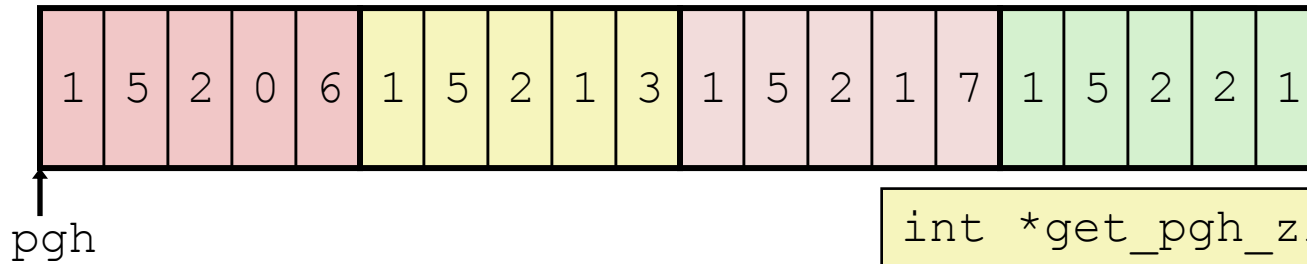
# Nested Array Row Access

- Row Vectors
  - $\mathbf{A}[i]$  is array of  $C$  elements
  - Each element of type  $T$  requires  $K$  bytes
  - Starting address  $\mathbf{A} + i * (C * K)$

```
int A[R][C];
```



# Nested Array Row Access Code



```
int *get_pgh_zip(int index)
{
    return pgh[index];
}
```

```
# %rdi = index
leaq (%rdi,%rdi,4),%rax # 5 * index
leaq pgh(,%rax,4),%rax  # pgh + (20 * index)
```

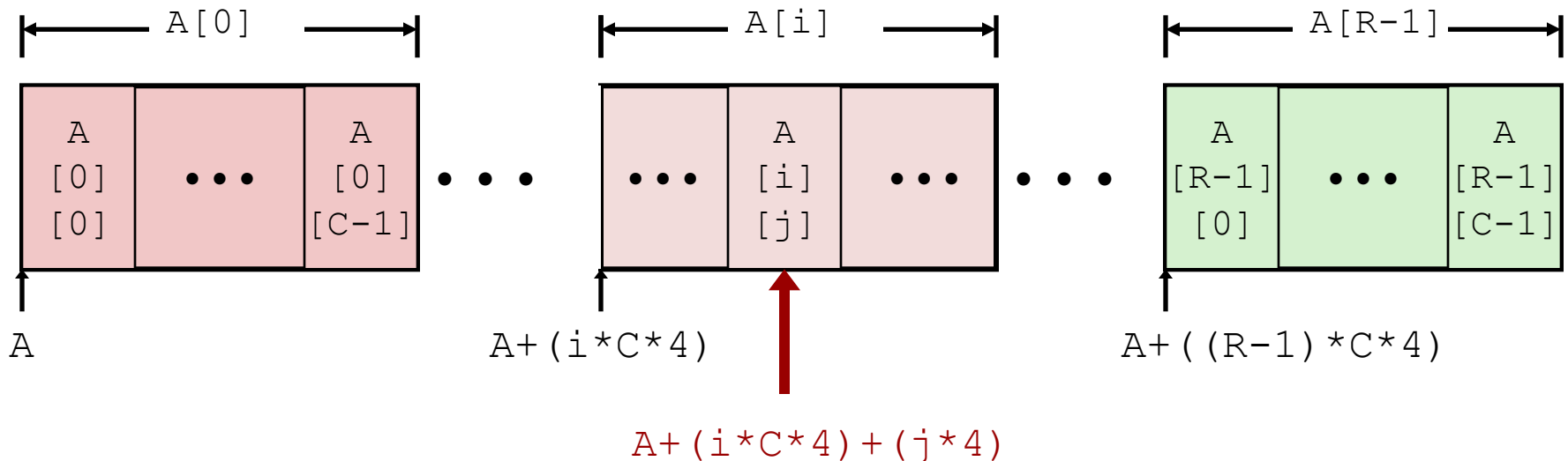
- Row Vector
  - **pgh[index]** is array of 5 **int**'s
  - Starting address **pgh+20\*index**
- Machine Code
  - Computes and returns address
  - Compute as **pgh + 4\*(index+4\*index)**



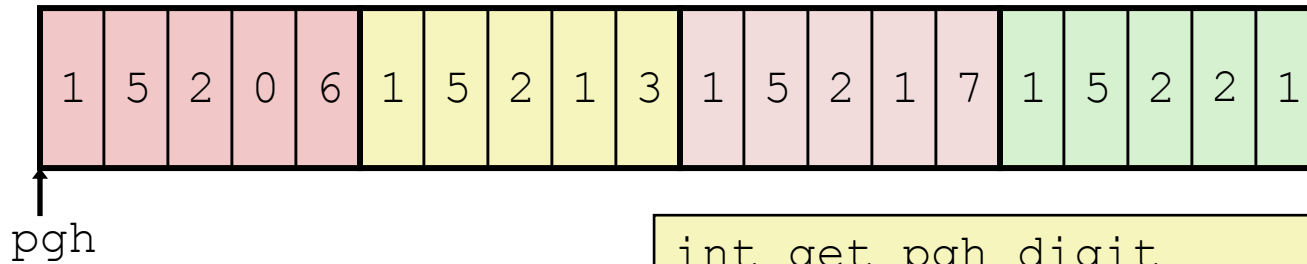
# Nested Array Element Access

- Array Elements
  - $A[i][j]$  is element of type  $T$ , which requires  $K$  bytes
  - Address  $A + i * (C * K) + j * K = A + (i * C + j) * K$

```
int A[R][C];
```



# Nested Array Element Access Code



```
int get_pgh_digit
(int index, int dig)
{
    return pgh[index][dig];
}
```

```
leaq    (%rdi,%rdi,4), %rax    # 5*index
addl    %rax, %rsi            # 5*index+dig
movl    pgh(,%rsi,4), %eax     # M[pgh + 4*(5*index+dig)]
```

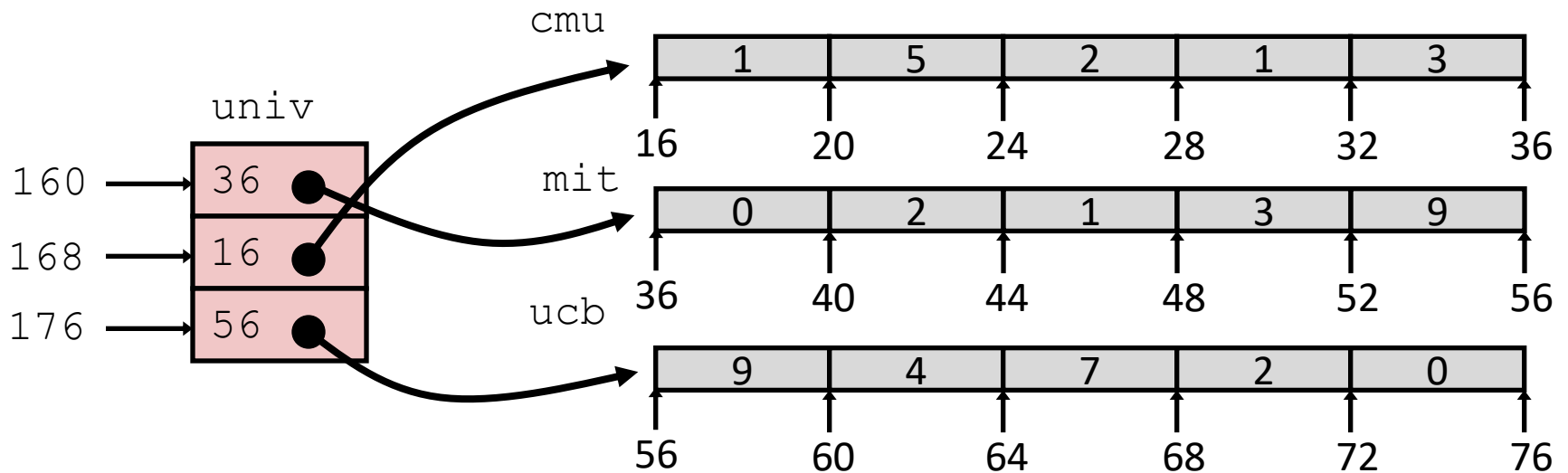
- Array Elements
  - `pgh[index][dig]` is `int`
  - Address: `pgh + 20*index + 4*dig`
    - = `pgh + 4*(5*index + dig)`

# Multi-Level Array Example

```
zip_dig cmu = { 1, 5, 2, 1, 3 };  
zip_dig mit = { 0, 2, 1, 3, 9 };  
zip_dig ucb = { 9, 4, 7, 2, 0 };
```

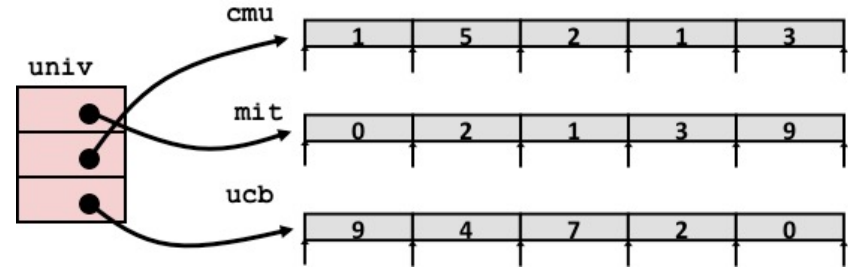
```
#define UCOUNT 3  
int *univ[UCOUNT] = {mit, cmu, ucb};
```

- Variable `univ` denotes array of 3 elements
- Each element is a pointer – 8 bytes
- Each pointer points to array of `int`'s



# Element Access in Multi-Level Array

```
int get_univ_digit
(size_t index, size_t digit)
{
    return univ[index][digit];
}
```



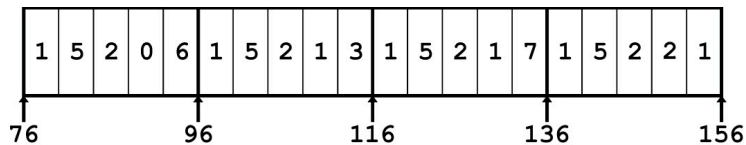
```
salq    $2, %rsi          # 4*digit
addq    univ(,%rdi,8), %rsi # p = univ[index] + 4*digit
movl    (%rsi), %eax       # return *p
ret
```

- Computation
  - Element access **Mem[Mem[univ+8\*index]+4\*digit]**
  - Must do two memory reads
    - First get pointer to row array
    - Then access element within array

# Array Element Accesses

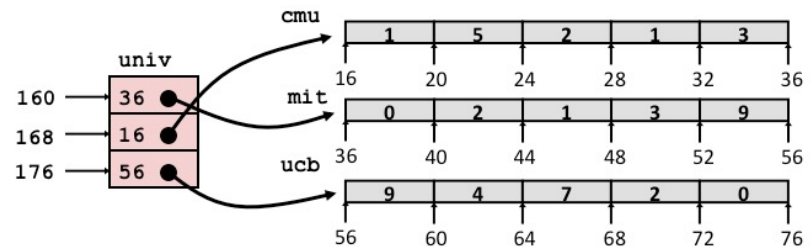
## Nested array

```
int get_pgh_digit
(size_t index, size_t digit)
{
    return pgh[index][digit];
}
```



## Multi-level array

```
int get_univ_digit
(size_t index, size_t digit)
{
    return univ[index][digit];
}
```



Accesses looks similar in C, but address computations very different:

$\text{Mem}[\text{pgh} + 20 * \text{index} + 4 * \text{digit}]$        $\text{Mem}[\text{Mem}[\text{univ} + 8 * \text{index}] + 4 * \text{digit}]$

# N X N Matrix

## Code

- Fixed dimensions
  - Know value of N at compile time
- Variable dimensions, explicit indexing
  - Traditional way to implement dynamic arrays
- Variable dimensions, implicit indexing
  - Now supported by gcc

```
#define N 16
typedef int fix_matrix[N][N];
/* Get element a[i][j] */
int fix_ele(fix_matrix a,
            size_t i, size_t j)
{
    return a[i][j];
}
```

```
#define IDX(n, i, j) ((i)*(n)+(j))
/* Get element a[i][j] */
int vec_ele(size_t n, int *a,
            size_t i, size_t j)
{
    return a[IDX(n,i,j)];
}
```

```
/* Get element a[i][j] */
int var_ele(size_t n, int a[n][n],
            size_t i, size_t j) {
    return a[i][j];
}
```

# 16 X 16 Matrix Access

## ■ Array Elements

- Address  $\mathbf{A} + i * (\mathbf{C} * \mathbf{K}) + j * \mathbf{K}$
- $\mathbf{C} = 16, \mathbf{K} = 4$

```
/* Get element a[i][j] */  
int fix_ele(fix_matrix a, size_t i, size_t j) {  
    return a[i][j];  
}
```

```
# a in %rdi, i in %rsi, j in %rdx  
salq    $6, %rsi          # 64*i  
addq    %rsi, %rdi         # a + 64*i  
movl    (%rdi,%rdx,4), %eax # M[a + 64*i + 4*j]  
ret
```

# n X n Matrix Access

## ■ Array Elements

- Address  $\mathbf{A} + i * (\mathbf{C} * \mathbf{K}) + j * \mathbf{K}$
- $\mathbf{C} = \mathbf{n}, \mathbf{K} = 4$
- Must perform integer multiplication

```
/* Get element a[i][j] */  
int var_ele(size_t n, int a[n][n], size_t i, size_t j)  
{  
    return a[i][j];  
}
```

```
# n in %rdi, a in %rsi, i in %rdx, j in %rcx  
imulq    %rdx, %rdi          # n*i  
leaq     (%rsi,%rdi,4), %rax  # a + 4*n*i  
movl     (%rax,%rcx,4), %eax  # a + 4*n*i + 4*j  
ret
```



# C - Structures

Arrays allow to define type of variables that can hold several data items of the same kind. Similarly **structure** is another user defined data type available in C that allows to combine data items of different kinds.

Structures are used to represent a record. Suppose you want to keep track of your books in a library. You might want to track the following attributes about each book –

- Title
- Author
- Subject
- Book ID

# Defining a Structure

To define a structure, you must use the **struct** statement. The struct statement defines a new data type, with more than one member. The format of the struct statement is as follows –

```
struct [structure tag] {  
    member definition;  
    member definition;  
    ...  
    member definition;  
} [one or more structure variables];
```

The **structure tag** is optional and each member definition is a normal variable definition, such as `int i;` or `float f;` or any other valid variable definition. At the end of the structure's definition, before the final semicolon, you can specify one or more structure variables but it is optional. Here is the way you would declare the Book structure –

```
struct Books {  
    char  title[50];  
    char  author[50];  
    char  subject[100];  
    int   book_id;  
} book;
```

# Accessing Structure Members

To access any member of a structure, we use the **member access operator (.)**. The member access operator is coded as a period between the structure variable name and the structure member that we wish to access. You would use the keyword **struct** to define variables of structure type. The following example shows how to use a structure in a program;

<https://www.tutorialspoint.com/tpcg.php?p=7DL5Jk>

# Structures

To learn more about structures

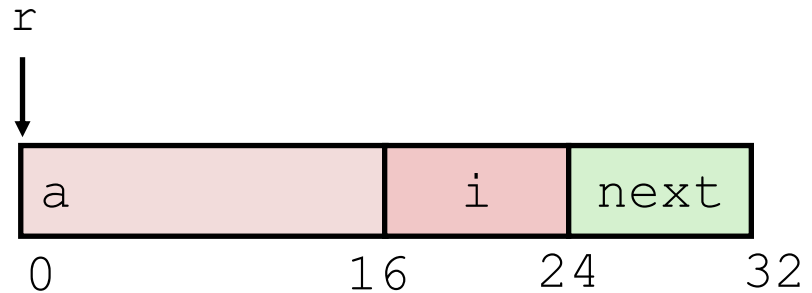
[https://www.tutorialspoint.com/cprogramming/c\\_structures.htm](https://www.tutorialspoint.com/cprogramming/c_structures.htm)

<https://www.learn-c.org/en/Structures>

<https://www.programiz.com/c-programming/c-structures>

# Structure Representation

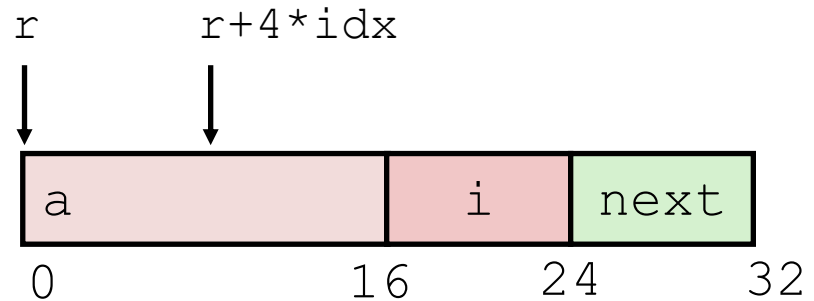
```
struct rec {  
    int a[4];  
    size_t i;  
    struct rec *next;  
};
```



- Structure represented as block of memory
  - **Big enough to hold all of the fields**
- Fields ordered according to declaration
  - **Even if another ordering could yield a more compact representation**
- Compiler determines overall size + positions of fields
  - **Machine-level program has no understanding of the structures in the source code**

# Generating Pointer to Structure Member

```
struct rec {  
    int a[4];  
    size_t i;  
    struct rec *next;  
};
```



- Generating Pointer to Array Element

- Offset of each structure member determined at compile time
- Compute as  **$r + 4 \cdot idx$**

```
int *get_ap  
(struct rec *r, size_t idx)  
{  
    return &r->a[idx];  
}
```

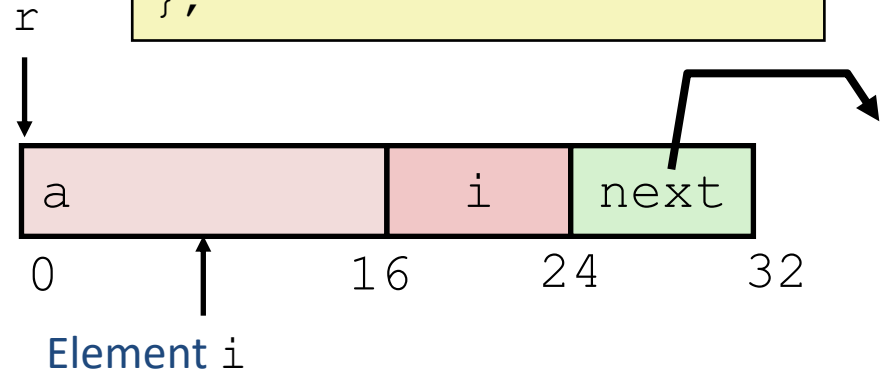
```
# r in %rdi, idx in %rsi  
leaq (%rdi,%rsi,4), %rax  
ret
```

# Following Linked List

- C Code

```
void set_val
(struct rec *r, int val)
{
    while (r) {
        int i = r->i;
        r->a[i] = val;
        r = r->next;
    }
}
```

```
struct rec {
    int a[4];
    int i;
    struct rec *next;
};
```



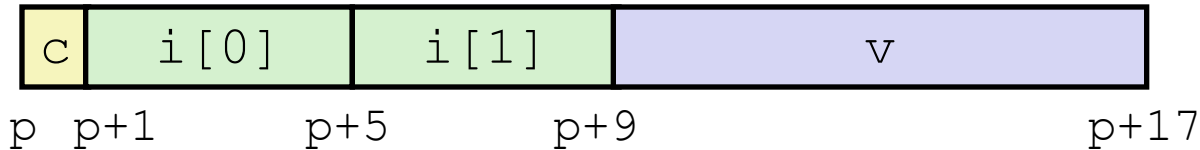
Register	Value
%rdi	r
%rsi	val

```
.L11:                                # loop:
    movslq    16(%rdi), %rax          # i = M[r+16]
    movl      %esi, (%rdi,%rax,4)     # M[r+4*i] = val
    movq      24(%rdi), %rdi          # r = M[r+24]
    testq     %rdi, %rdi              # Test r
    jne       .L11                    # if !=0 goto loop
```



# Structures & Alignment

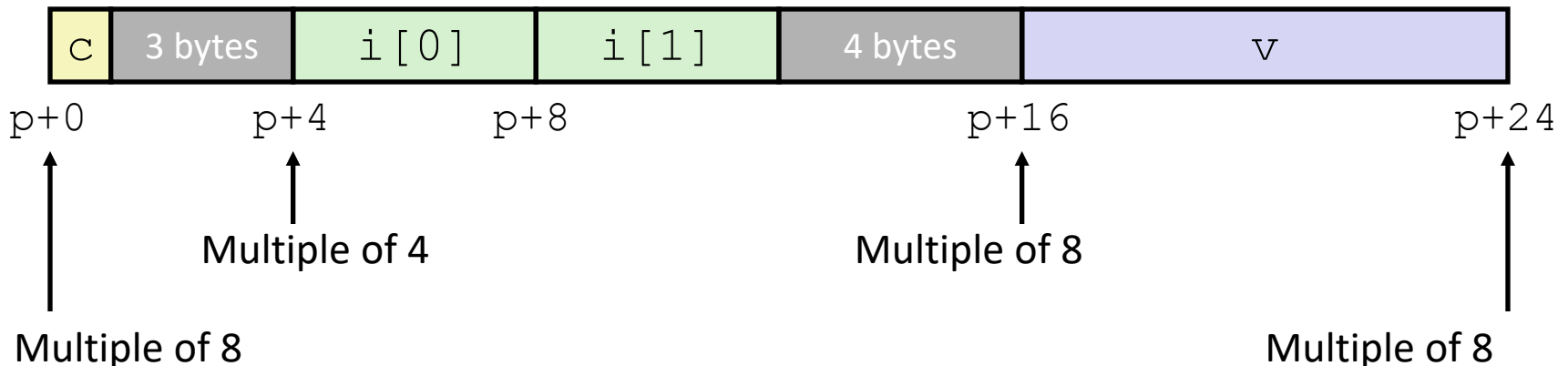
- Unaligned Data



```
struct S1 {  
    char c;  
    int i[2];  
    double v;  
} *p;
```

- Aligned Data

- Primitive data type requires K bytes
- Address must be multiple of K



# Alignment Principles

- Aligned Data
  - Primitive data type requires K bytes
  - Address must be multiple of K
  - Required on some machines; advised on x86-64
- Motivation for Aligning Data
  - Memory accessed by (aligned) chunks of 4 or 8 bytes (system dependent)
    - Inefficient to load or store datum that spans quad word boundaries
    - Virtual memory trickier when datum spans 2 pages
- Compiler
  - Inserts gaps in structure to ensure correct alignment of fields

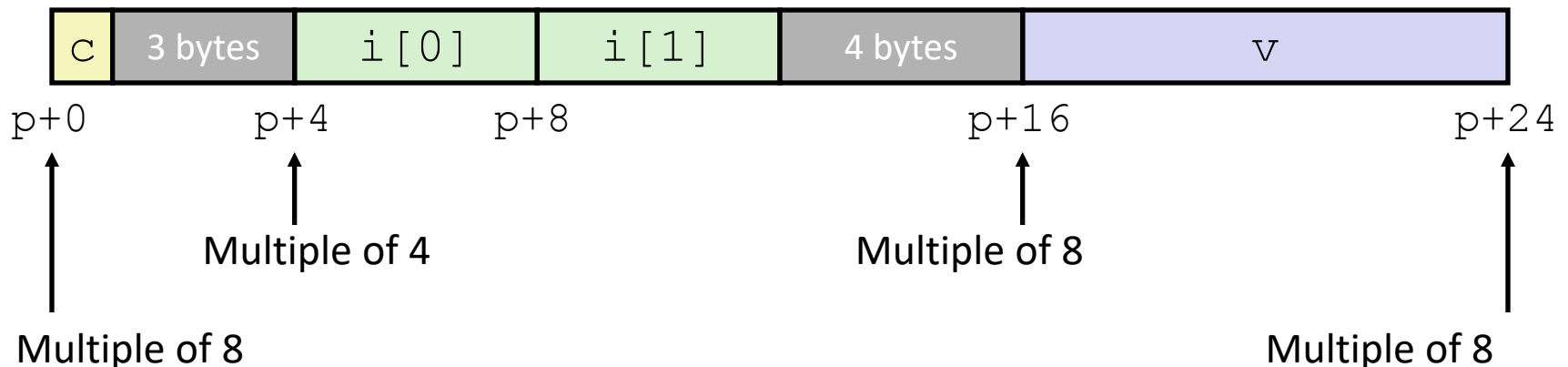
# Specific Cases of Alignment (x86-64)

- 1 byte: **char**, ...
  - no restrictions on address
- 2 bytes: **short**, ...
  - lowest 1 bit of address must be 0<sub>2</sub>
- 4 bytes: **int**, **float**, ...
  - lowest 2 bits of address must be 00<sub>2</sub>
- 8 bytes: **double**, `long`, **char \***, ...
  - lowest 3 bits of address must be 000<sub>2</sub>
- 16 bytes: **long double** (GCC on Linux)
  - lowest 4 bits of address must be 0000<sub>2</sub>

# Satisfying Alignment with Structures

- Within structure:
  - Must satisfy each element's alignment requirement
- Overall structure placement
  - Each structure has alignment requirement K
    - K = Largest alignment of any element
  - Initial address & structure length must be multiples of K
- Example:
  - K = 8, due to **double** element

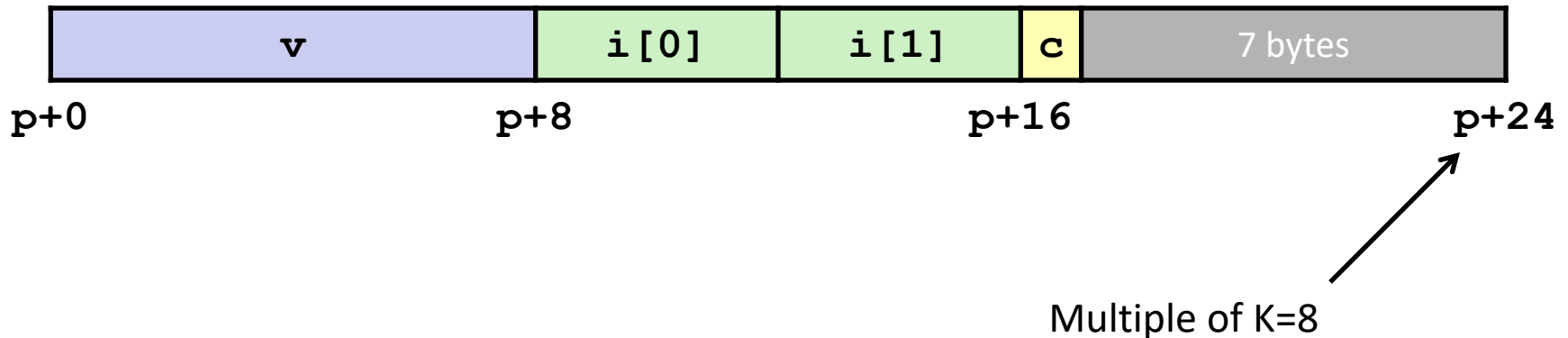
```
struct S1 {  
    char c;  
    int i[2];  
    double v;  
} *p;
```



# Meeting Overall Alignment Requirement

- For largest alignment requirement K
- Overall structure must be multiple of K

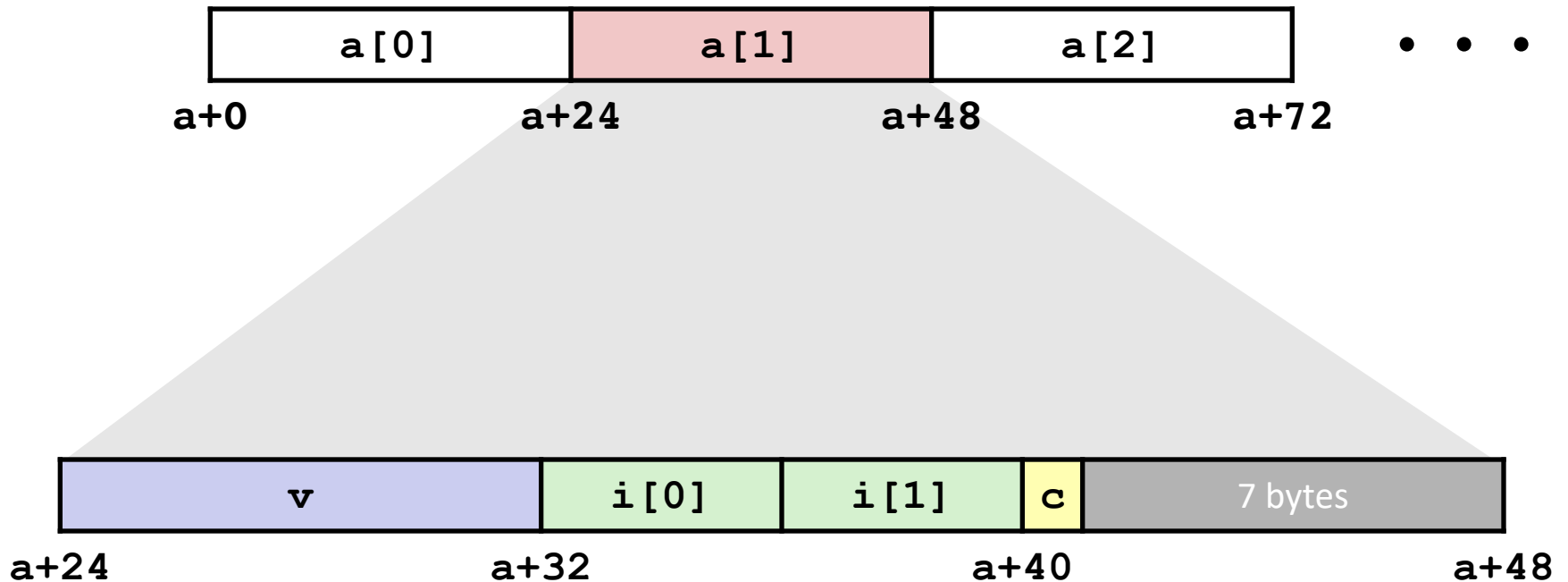
```
struct S2 {  
    double v;  
    int i[2];  
    char c;  
} *p;
```



# Arrays of Structures

- Overall structure length multiple of K
- Satisfy alignment requirement for every element

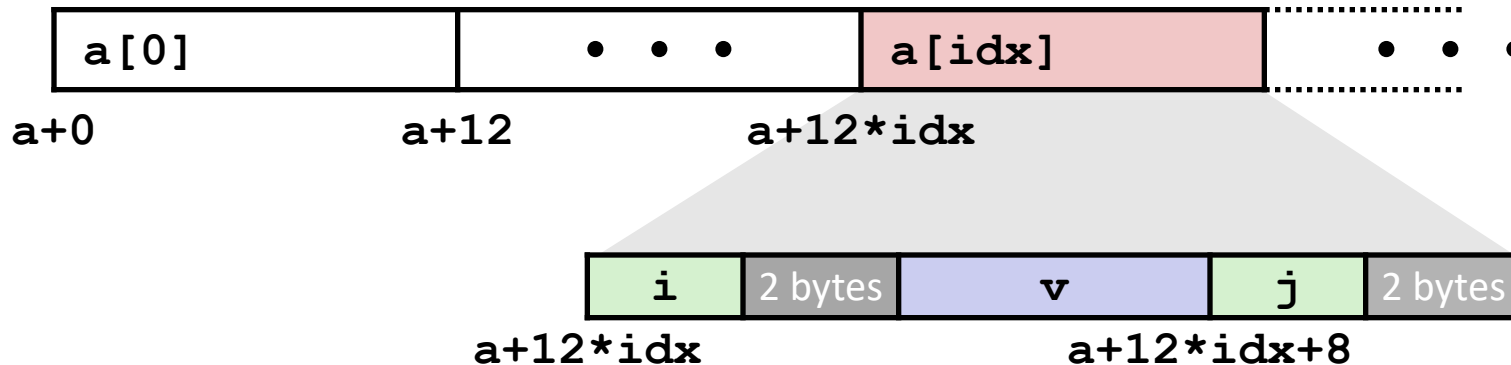
```
struct S2 {  
    double v;  
    int i[2];  
    char c;  
} a[10];
```



# Accessing Array Elements

- Compute array offset  $12 * \text{idx}$ 
  - `sizeof(S3)`, including alignment spacers
- Element `j` is at offset 8 within structure
- Assembler gives offset `a+8`
  - Resolved during linking

```
struct S3 {  
    short i;  
    float v;  
    short j;  
} a[10];
```



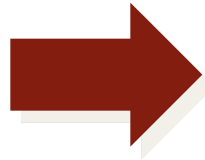
```
short get_j(int idx)  
{  
    return a[idx].j;  
}
```

```
# %rdi = idx  
leaq (%rdi,%rdi,2),%rax # 3*idx  
movzwl a+8(,%rax,4),%eax
```

# Saving Space

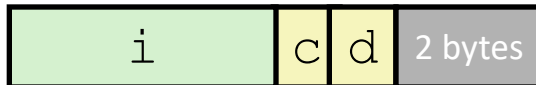
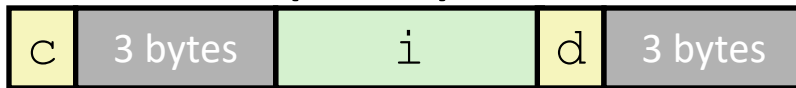
- Put large data types first

```
struct S4 {  
    char c;  
    int i;  
    char d;  
} *p;
```



```
struct S5 {  
    int i;  
    char c;  
    char d;  
} *p;
```

- Effect (K=4)





# Next Class

- Combining Control and Data in Machine Level Programs
- Floating Points