

# 作業系統 - 作業三 (System Call & CPU Scheduling)

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## Motivation

### Problem analysis

當使用NachOS現有的實體記憶體大小，某些測試case會需要大量記憶體，如此將導致core dumped。如下圖所見，三個測試case都發生core dumped的狀況，因為主記憶體中的空間可能不足，因此需要虛擬記憶體來儲存Pages。

```
u14@ubuntu:~/r11921091_nachos2/nachos-4.0/code/userprog$ ./nachos -e ../test/matmult
Total threads number is 1
Thread ../test/matmult is executing.
Assertion failed: line 134 file ../userprog/addrspace.cc
Aborted (core dumped)
u14@ubuntu:~/r11921091_nachos2/nachos-4.0/code/userprog$ ./nachos -e ../test/sort
Total threads number is 1
Thread ../test/sort is executing.
Assertion failed: line 134 file ../userprog/addrspace.cc
Aborted (core dumped)
u14@ubuntu:~/r11921091_nachos2/nachos-4.0/code/userprog$ ./nachos -e ../test/matmult -e ../test/sort
Total threads number is 2
Thread ../test/matmult is executing.
Thread ../test/sort is executing.
Assertion failed: line 134 file ../userprog/addrspace.cc
Aborted (core dumped)
```

### Plan

我們的目標是實現一種虛擬記憶體，當主記憶體沒有空間，Page和data仍可以透過Page swap演算法儲存在其他儲存設備中。也就是當主記憶體中的空間足夠時，Page將存儲在主記憶體中，這樣效能會比較好。但一旦空間不足，Page和data將swap出去。當主記憶體空間已被釋放，且需要資料時，再Page和data swap到主記憶體中。Page替換的演算法的排程方法則由LRU ( Least Recently Used ) 達成。其實這樣的作法就是common linux os的作法，以前有玩過ubuntu server kernel的這個部分，所以有跡可循。

## Implementatoin

先將multi-programming的部分做個修改，增加輔助的記憶體 SynchDisk \*Swap\_Space，用來存放主記憶體放不下的Pages。

Code

code/userprog/userkernel.h

```

1  class UserProgKernel : public ThreadedKernel {
2      public:
3          ...
4          SynchDisk *Swap_Space; // create swap space for virtual memory
5          ...
6      private
7          ...
8  };

```

初始化並分配Disk給Swap\_Space，接著就能透過這個disk做swap space。

Code

code/userprog/userkernel.cc

```

1  void
2  UserProgKernel::Initialize()
3  {
4      ...
5      Swap_Space = new SynchDisk("New_Swap_Space"); //Create swap space for
virtual memory
6      ...
7  }

```

接著，在class Machine增加一些會使用到的變數宣告：

Code

code/machine/machine.h

```

1  class Machine {
2      public:
3          ...
4          int Identity;
5          int SectorNum; //record sector number
6          int FrameName[NumPhysPages];
7          bool usedPhyPage[NumPhysPages]; //record which frame in the main
memory is occupied.
8          bool usedVirPage[NumPhysPages];
9
10         // start for page replacement //
11         int count[NumPhysPages]; //for LRU counter
12         bool reference_bit[NumPhysPages]; //for second chance algorithm.
13         ...
14     private:
15         ...
16 };

```

再於 addrspace.h 的 class AddrSpace 宣告必要變數 ID 和 pageTableLoaded；

ID 用來儲存thread的ID；pageTableLoaded用來存判斷page Table是否有被載入的boolean變數。

Code

code/userprog/addrspace.h

```

1  class AddrSpace {
2      public:
3          ...
4          int ID; // store the ID of the thread
5          ...
6      private:
7          ...
8          bool pageTableLoaded;
9          ...
10 };

```

接下來，這裡將負責載入程式碼到 memory 裡面，所以相當重要。

我在這裡開一個專屬這個 thread 的 page table，在 load 時一直往下找記憶體直到找到沒被用到的或是到底為止，以此來判斷 memory 夠不夠用，若不夠就需要用到輔助記憶體，需要存好一些重要參數以便要用的時候找得到。

修改幾乎是在下面的AddrSpace Load 函數內完成的。首先會將可執行程式碼分配到主記憶體中，並記錄相應的Page Table。如果主記憶體中的frame number足夠，Page 將被儲存到主記憶體中。當frame number不足時，後面的page將被WriteSector寫入虛擬記憶體。相應的page table也會被記錄並設置為無效。

Code

code/userprog/addrspace.cc

```

1  ...
2  AddrSpace::AddrSpace()
3  {
4      ID=(kernel->machine->Identity)++;
5      kernel->machine->Identity=(kernel->machine->Identity)++;
6
7  }
8  ...
9  bool
10 AddrSpace::Load(char *fileName)
11 {
12     openFile *executable = kernel->filesystem->Open(fileName);
13     NoffHeader noffH;
14
15     unsigned int size,tmp;
16
17     if (executable == NULL) {
18         cerr << "Unable to open file " << fileName << "\n";
19         return FALSE;
20     }
21     executable->ReadAt((char *)&noffH, sizeof(noffH), 0);
22     if ((noffH.noffMagic != NOFFMAGIC) &&
23         (wordToHost(noffH.noffMagic) == NOFFMAGIC))
24         SwapHeader(&noffH);
25     ASSERT(noffH.noffMagic == NOFFMAGIC);
26
27     // how big is address space?
28     size = noffH.code.size + noffH.initData.size + noffH.uninitData.size
29           + UserStackSize;    // we need to increase the size
30                               // to leave room for the stack
31     numPages = divRoundUp(size, PageSize);
32     // cout << "number of pages of " << fileName << " is " << numPages << endl;

```

```

33
34
35     pageTable = new TranslationEntry[numPages];
36
37     size = numPages * PageSize;
38
39     //  ASSERT(numPages <= NumPhysPages);    // check we're not trying
40                                           // to run anything too big --
41                                           // at least until we have
42                                           // virtual memory
43
44     //  DEBUG(dbgAddr, "Initializing address space: " << numPages << ", " <<
size);
45
46     // then, copy in the code and data segments into memory
47
48
49
50     if (noffH.code.size > 0) {
51         //      DEBUG(dbgAddr, "Initializing code segment.");
52         //  DEBUG(dbgAddr, noffH.code.virtualAddr << ", " << noffH.code.size);
53
54         for(int j=0,i=0;i < numPages ;i++){
55             j=0;
56             while(kernel->machine->usedPhyPage[j] != FALSE && j <
NumPhysPages)
57                 j += 1;
58
59             //if memory is enough,just put data in without using
virtual memory
60             if(j<NumPhysPages){
61                 pageTable[i].physicalPage = j;    // record in
physical memory position j
62                 pageTable[i].use = FALSE;
63                 pageTable[i].dirty = FALSE;
64                 pageTable[i].ID =ID;
65                 pageTable[i].readOnly = FALSE;
66                 pageTable[i].valid = TRUE;    // TRUE means the
page exists in physical memory
67                 kernel->machine->usedPhyPage[j]=TRUE;
68                 kernel->machine->FrameName[j]=ID;
69                 kernel->machine->main_tab[j]=&pageTable[i];    //
save the page pointer
70                 pageTable[i].count++;
71                 // save data to position j
72                 executable->ReadAt(&(kernel->machine-
>mainMemory[j*PageSize]),PageSize, noffH.code.inFileAddr+(i*PageSize));
73             }
74             //Use virtual memory technique
75             else{
76                 char *buffer;
77                 buffer = new char[PageSize];
78                 tmp=0;
79                 while(kernel->machine->usedVirPage[tmp]!=FALSE)
{tmp++;}
80                 pageTable[i].virtualPage=tmp;
81                 pageTable[i].ID =ID;
82                 pageTable[i].valid = FALSE;

```

```

83         pageTable[i].dirty = FALSE;
84         pageTable[i].readOnly = FALSE;
85         pageTable[i].use = FALSE;
86         kernel->machine->usedVirPage[tmp]=true;
87         executable->ReadAt(buffer, PageSize,
noffH.code.inFileAddr+(i*PageSize));
88         kernel->Swap_Space->WriteSector(tmp, buffer); //call
virtual_disk write in virtual memory
89
90     }
91 }
92 }
93
94
95
96     if (noffH.initData.size > 0) {
97         executable->ReadAt(
98             &(kernel->machine->mainMemory[noffH.initData.virtualAddr]),
99             noffH.initData.size, noffH.initData.inFileAddr);
100     }
101
102     delete executable;          // close file
103     return TRUE;                // success
104 }
105
106 //-----
107 // AddrSpace::Execute
108 // Run a user program. Load the executable into memory, then
109 // (for now) use our own thread to run it.
110 //
111 // "fileName" is the file containing the object code to load into memory
112 //-----
113
114 void
115 AddrSpace::Execute(char *fileName)
116 {
117     Is_ptable_loaded=FALSE;
118     if (!Load(fileName)) {
119         cout << "inside !Load(fileName)" << endl;
120         return;                // executable not found
121     }
122
123     //kernel->currentThread->space = this;
124     this->InitRegisters();      // set the initial register values
125     this->RestoreState();       // load page table register
126     Is_ptable_loaded=TRUE;
127     kernel->machine->Run();      // jump to the user program
128
129     ASSERTNOTREACHED();        // machine->Run never returns;
130                                // the address space exits
131                                // by doing the syscall "exit"
132 }
133
134
135 //-----
136 // AddrSpace::InitRegisters
137 // Set the initial values for the user-level register set.
138 //

```

```

139 // we write these directly into the "machine" registers, so
140 // that we can immediately jump to user code. Note that these
141 // will be saved/restored into the currentThread->userRegisters
142 // when this thread is context switched out.
143 //-----
144
145 void
146 AddrSpace::InitRegisters()
147 {
148     Machine *machine = kernel->machine;
149     int i;
150
151     for (i = 0; i < NumTotalRegs; i++)
152         machine->WriteRegister(i, 0);
153
154     // Initial program LRU_timeser -- must be location of "Start"
155     machine->WriteRegister(PCReg, 0);
156
157     // Need to also tell MIPS where next instruction is, because
158     // of branch delay possibility
159     machine->WriteRegister(NextPCReg, 4);
160
161     // Set the stack register to the end of the address space, where we
162     // allocated the stack; but subtract off a bit, to make sure we don't
163     // accidentally reference off the end!
164     machine->WriteRegister(StackReg, numPages * PageSize - 16);
165     DEBUG(dbgAddr, "Initializing stack pointer: " << numPages * PageSize -
166 16);
167 }
168 //-----
169 // AddrSpace::SaveState
170 // On a context switch, save any machine state, specific
171 // to this address space, that needs saving.
172 //
173 // For now, don't need to save anything!
174 //-----
175
176 void AddrSpace::SaveState()
177 {
178     if(Is_ptable_loaded){
179         pageTable=kernel->machine->pageTable;
180         numPages=kernel->machine->pageTableSize;
181     }
182 }
183 ...

```

在Class TranslationEntry 中增加 ID 和 count。

Code

code/machine/translate.h

```

1 class TranslationEntry {
2     public:
3         ...
4         int ID;
5
6         int count;    //for Least Recently used algorithm
7
8 };

```

下面會是實現 LRU 的主要部分。首先，Page Table的有效位置將被識別以確保Page的位置，就是確認在主記憶體中還是虛擬記憶體中。如果Page在虛擬記憶體中，就會檢查現在主記憶體內是否有閒置空間。

Page Swap發生在主記憶體沒有空間的情況下，將兩個緩存區用於Swap in/Swap out。而為了達成LRU，會用一個for迴圈尋找最少使用的page，從虛擬記憶體中讀取，複製到主記憶體，選擇換出Page到虛擬記憶體，這些步驟會包含bcopy、ReadSector、WriteSector這些動作。

Code

code/machine/translate.cc

```

1  ...
2  ExceptionType
3  Machine::Translate(int virtAddr, int* physAddr, int size, bool writing)
4  {
5      int i, j;
6      ...
7      int Swap_out_page; ///find the page Swap_out_page
8      ...
9      if (tlb == NULL) {        // => page table => vpn is index into table
10         if (vpn >= pageTableSize) {
11             DEBUG(dbgAddr, "Illegal virtual page # " << virtAddr);
12             return AddressErrorException;
13         } else if (!pageTable[vpn].valid) {
14
15             //printf("Page fault Happen!\n");
16             kernel->stats->numPageFaults += 1;    // nachos pagefault counter
17             +1
18             j=0;
19             while(kernel->machine->usedPhyPage[j] != FALSE && j < NumPhysPages)
20                 j += 1;    // find valid frame space
21
22             if( j < NumPhysPages){    // if find valid frame
23                 space, save the page in virtual memory into physical memory
24                 char *buffer; //save page temporary
25                 buffer = new char[PageSize];
26                 pageTable[vpn].physicalPage = j;    // save
27                 physical memory position
28                 pageTable[vpn].valid = TRUE;    // page has
29                 already in physical memory
30                 kernel->machine->usedPhyPage[j]=TRUE;
31                 kernel->machine->FrameName[j]=pageTable[vpn].ID;
32                 kernel->machine->main_tab[j]=&pageTable[vpn];
33
34                 // save the page pointer
35
36                 pageTable[vpn].count++; //for LRU
37
38             }
39         }
40     }
41 }

```

```

32         kernel->Swap_Space-
>ReadSector(pageTable[vpn].virtualPage, buffer);        // read data from swap
area
33         bcopy(buffer,&mainMemory[j*PageSize],PageSize);
// save data into physical memory
34
35     }
36     else{
37         char *buffer1;
38         char *buffer2;
39         buffer1 = new char[PageSize];
40         buffer2 = new char[PageSize];
41
42         //Random
43         //Swap_out_page = (rand()%32);
44
45         //LRU
46
47         int min = pageTable[0].count;
48         swap_out_page=0;
49         for(int cc=0;cc<32;cc++){
50             if(min > pageTable[cc].count){
51                 min = pageTable[cc].count;
52                 swap_out_page = cc;
53             }
54         }
55     }
56     pageTable[Swap_out_page].count++;
57
58     //printf("Page%d swap out!\n",Swap_out_page);
59
60     bcopy(&mainMemory[Swap_out_page*PageSize],buffer1,PageSize);
        kernel->Swap_Space-
>ReadSector(pageTable[vpn].virtualPage, buffer2);
61
62     bcopy(buffer2,&mainMemory[Swap_out_page*PageSize],PageSize);
        kernel->Swap_Space-
>writeSector(pageTable[vpn].virtualPage,buffer1);
63
64         main_tab[Swap_out_page]-
>virtualPage=pageTable[vpn].virtualPage;
65         main_tab[Swap_out_page]->valid=FALSE;
66
67         pageTable[vpn].valid = TRUE;
68         pageTable[vpn].physicalPage = Swap_out_page;
69         kernel->machine-
>FrameName[Swap_out_page]=pageTable[vpn].ID;
70         main_tab[Swap_out_page]=&pageTable[vpn];
71         //printf("Finish the page replacement!\n");
72
73     }
74
75 }
76 ...
77 }
78

```



# Result:

## 1. 執行 `./nachos -e ../test/matmult`

```
u14@ubuntu:~/r11921091_nachos3/nachos-4.0/code/userprog$ ./nachos -e ../test/matmult
Total threads number is 1
Thread ../test/matmult is executing.
return value:7220
No threads ready or runnable, and no pending interrupts.
Assuming the program completed.
Machine halting!

Ticks: total 7691030, idle 1365666, system 6325360, user 4
Disk I/O: reads 80, writes 102
Console I/O: reads 0, writes 0
Paging: faults 80
Network I/O: packets received 0, sent 0
```

## 2. 執行 `./nachos -e ../test/sort`

```
u14@ubuntu:~/r11921091_nachos3/nachos-4.0/code/userprog$ ./nachos -e ../test/sort
Total threads number is 1
Thread ../test/sort is executing.
return value:1
No threads ready or runnable, and no pending interrupts.
Assuming the program completed.
Machine halting!

Ticks: total 440599030, idle 52221566, system 388377460, user 4
Disk I/O: reads 5536, writes 5550
Console I/O: reads 0, writes 0
Paging: faults 5536
Network I/O: packets received 0, sent 0
```

## 3. 執行 `./nachos -e ../test/matmult -e ../test/sort`

```
u14@ubuntu:~/r11921091_nachos3/nachos-4.0/code/userprog$ ./nachos -e ../test/matmult -e ../test/sort
Total threads number is 2
Thread ../test/matmult is executing.
Thread ../test/sort is executing.
return value:7220
return value:1
No threads ready or runnable, and no pending interrupts.
Assuming the program completed.
Machine halting!

Ticks: total 489532030, idle 94826365, system 394705660, user 5
Disk I/O: reads 5647, writes 5715
Console I/O: reads 0, writes 0
Paging: faults 5647
Network I/O: packets received 0, sent 0
```

p.s 如果想看細部 Page swap 的時間點可以去[translate.cc](https://translate.cc) 將第260行的的註解拿掉，如下：

```
1 | printf("Page%d swap out!\n",Swap_out_page);
```

但因為 sort 的時候交換太多次，所以 linux 終端機會把 output 在視窗上洗掉，導致無法看到 matmult 和 sort 同時執行時 matmult 的 return 結果。因此，我選擇將這行先註解掉。

## The difficulties I encountered

還不知道，自己寫得時候還好，畢竟以前做過類似的，唯一的問題就是sort.c的檔案本來是錯的，所以會return 0讓我很confused，但詢問過助教電神宗翰後，電神也更新了sort.c的code，所以就沒問題了！

因為這次太早寫完，所以之後同學問問題後可能會再更新。

## Reference

[Nachos hw3 assignments] <https://cool.ntu.edu.tw/courses/21578/files/3143901?wrap=1>

[Nachos Documentation] <https://homes.cs.washington.edu/~tom/nachos/>