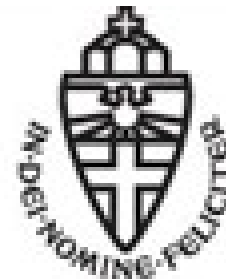


# advanced programming tutorial 2

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assignment 1

# OVERLOADING

# (de)serialize

**class** serialize a **where**

write :: a [String] -> [String]

read :: [String] -> Maybe (a, [String])

result

remaining input

:: Maybe a = Just a | Nothing

- Maybe result of read, it can fail
- do we need

**class** serialize a **where**

write :: a [String] -> [String]

read :: [String] -> (Maybe a, [String]) ?

NO

## instance for [a]

```
instance serialize [a] | serialize a where
  write [] c = [NilString: c]
  write [a:x] c
    = ["(",ConsString: write a (write x [")":c])]
  read [NilString:r]      = Just ([],r)
  read ["(",ConsString:r] =
    case read r of
      Just (a,s) = case read s of Just (x,[")":t]) =
Just ([a:x],t) _ = Nothing
  _ = Nothing
  read _ = Nothing
NilString  ::= "Nil"
ConsString ::= "Cons"
```

## instance for [a] alternative ?

```
instance serialize [a] | toString a where  
  write [] c = [NilString: c]  
  write [a:x] c  
    = ["(", ConsString, toString: a (write x  
[")"] : c)])  
  read ...
```

- is this equally good?
- compiler accepts it
- consider `write [[1],[2,3]]`  
the list `[1]` is transformed by `toString` instead of `write`!
- **this is undesirable**

# continuations

- the second argument of write is a continuation  
it tell what must be written after this element

```
write [] c = [NilString: c]  
write [a:x] c  
= ["(",ConsString: write a (write x [")":c])]
```

- without continuation we would have

```
write :: [a] -> [string] | write a  
write [] = [NilString]  
write [a:x]  
= ["(",ConsString: write a] ++ write x ++ [")"]
```

linear in the size of  
the first argument

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## continuations alternative ?

```
instance serialize [a] | serialize a where  
  write [] c = [NilString: c]  
  write [a:x] c  
    = ["(", ConsString: write a [], ++  
      write x [")"] ++ c  
  read ...
```

linear in the size of  
the first argument

- in general there is nothing wrong with ++  
but recursive calls on long list can be much work

# serialize Bin

```
:: Bin a = Leaf | Bin (Bin a) a (Bin a)
instance serialize (Bin a) | serialize a where
  write Leaf c = ["Leaf": c]
  write (Bin l a r) c
  = ["(", "Bin": write l (write a (write r [")":c)))]
  read [LeafString:r] = Just (Leaf,r)
  read ["(", BinString:r] =
    case read r of
      Just (l,s) = case read s of
        Just (a,t) = case read t of
          Just (r,[")":u]) = Just (Bin l a r,u)
          _ = Nothing
          _ = Nothing
          _ = Nothing
        read _ = Nothing
```

boring and error prone



# kinds

```
:: Bin a = Leaf | Bin (Bin a) a (Bin a)
:: Tree a b = Tip a | Node (Tree a b) b (Tree a b)
:: Rose a = Rose a [Rose a]
:: T1 a b = C11 (a b) | C12 b
:: T2 a b c = C2 (a (T1 b c))
:: T3 a b c = C3 (a b c)
:: T4 a b c = C4 (a (b c))
```

Bool	*
Bin	*→*
Rose	*→*
Bin Int	*
Tree	*→*→*
T1	(*→*)→*→*
T2	(*→*)→(*→*)→*→*
T3	(*→*→*)→*→*→*
T4	(*→*)→(*→*)→*→*

# polykinds

$:: T3\ a\ b\ c = C3\ (a\ b\ c)\ \text{kind: } (* \rightarrow * \rightarrow *) \rightarrow * \rightarrow * \rightarrow *$

$:: Box\ t\ u = Box\ (t\ u)\ \text{kind: } (* \rightarrow *) \rightarrow * \rightarrow *$

- application

$c3a :: T3\ (,) Int Int$

$c3a = C3\ (1,2)$

- polykinds: use  $*$  as a kind variable

$c3b = C3\ (Box\ [1,2,3])$

the actual kind is:  $((* \rightarrow *) \rightarrow * \rightarrow *) \rightarrow (* \rightarrow *) \rightarrow * \rightarrow *$

read the kinds as:  $(a \rightarrow b \rightarrow *) \rightarrow a \rightarrow b \rightarrow *$

- in Clean: accepted; derived type  $c3b :: T3\ Box\ []\ Int$
- specifying the type yields error: conflicting kinds  $* \rightarrow *$  and  $*$ 
  - definition without type is accepted by change

# polykinds 2

$:: T4\ a\ b\ c = C4\ (a\ (b\ c))$  kind:  $(\ast \rightarrow \ast) \rightarrow (\ast \rightarrow \ast) \rightarrow \ast \rightarrow \ast$

$:: IntBox\ t = IntBox\ (t\ Int)$  kind:  $(\ast \rightarrow \ast) \rightarrow \ast$

- application

$c4a :: T4\ []\ []\ Int$

$c4a = C4\ [[1]]$

- polykinds: use  $\ast$  as a kind variable

$c4b = C4\ (IntBox\ (Box\ [5]))$

the actual kind is:  $((\ast \rightarrow \ast) \rightarrow \ast) \rightarrow ((\ast \rightarrow \ast) \rightarrow \ast) \rightarrow (\ast \rightarrow \ast) \rightarrow \ast$

read the kinds as:  $(a \rightarrow \ast) \rightarrow (b \rightarrow a) \rightarrow b \rightarrow \ast$

- in Clean: accepted; derived type  $c4b :: T4\ IntBox\ Box\ []$
- specifying the type yields error: conflicting kinds  $\ast \rightarrow \ast$  and  $\ast$ 
  - different algorithms for checking and deriving kinds

# type constructor class: container

```
class Container t where
```

```
  Cinsert    :: a (t a) -> t a      | <  
  a
```

```
  Ccontains  :: a (t a) -> Bool     | <, Eq  
  a
```

```
  Cshow      :: (t a) -> [String] |  
toString a
```

```
  Cnew       :: t a
```

- $t$  is of kind  $* \rightarrow *$
- hence we need instances like `[]` and `Bin`

# list as container

```
instance Container [] where
```

```
  Cinsert    a c = [a: c]
```

```
  Ccontains a c = isMember a c
```

```
  Cnew      = []
```

```
  Cshow      c = ["{": showElems c ["}"]]
```

using a continuation

a is no type argument of this instance, we cannot impose the restriction here

```
showElems :: [a] [String] -> [String] |  
toString a
```

```
showElems []      c = c
```

```
showElems [x]     c = [toString x: c]
```

```
showElems [x: xs] c  
  = [toString x, ",", ":showElems xs c]
```

# search trees as container

```
:: Bin a = Tip | Bin (Tree a) a (Tree a)
```

```
instance Container Bin where
```

```
  Cinsert a Tip = Bin Tip a Tip
```

```
  Cinsert a (Bin l b r)
```

```
    | a < b = Bin (Cinsert a l) b r
```

```
          = Bin l b (Cinsert a r)
```

```
  Ccontains a Tip = False
```

```
  Ccontains a (Bin l b r)
```

```
    | a < b = Ccontains a l
```

```
    | b > a = Ccontains b r
```

```
          = True
```

```
  Cnew      = Tip
```

```
  Cshow t = [{"":showElems (TreetoList t []) ["}"]}]
```

```
TreetoList Tip          c = c
```

```
TreetoList (Bin l a r) c = TreetoList l [a: TreetoList r c]
```



$O(\log N)$



$O(\log N)$

assignment 2

# GENERIC 1

# serialize using generics

- serialization using classes is boring and error prone
- plan:
  - transform any data type to generics
  - serialize the generic representation
  - read the generic representation
  - transform generic representation to data type

- generic types used

:: UNIT = UNIT

:: EITHER a b = LEFT a | RIGHT b

:: PAIR a b = PAIR a b

:: CONS a = CONS String a



# making generic representations

- we do nothing smart for basic types  
Int, Bool, Char, Real ..
  - we ignore records, arrays, special lists etc. for the moment
- for an algebraic data type
  1. introduce EITHERs to separate the constructors
    - n EITHERs for n+1 constructors
  2. every alternative starts with a CONS
  3. use PAIRs to glue arguments together
    - n PAIRs for n+1 arguments
    - do not transform the arguments!
    - no arguments: use a UNIT

we can replace a  
single argument by  
(PAIR a UNIT)

`:: Li = Ni | Ci Int Li`

`:: LiG ::= EITHER (CONS UNIT) (CONS (PAIR (Int Li)))`

# generic representation

- there is choice in generic representations

```
:: D3 = D3 Int Int Int
```

```
:: D3G1 ::= CONS (PAIR Int (PAIR Int Int))
```

```
:: D3G2 ::= CONS (PAIR (PAIR Int Int) Int)
```

```
:: Val = Int Int | Bool Bool | Char Char
```

```
:: ValG1
```

```
    ::= EITHER (EITHER (CONS Int) (Cons Bool))  
          (Cons Char)
```

```
:: ValG2
```

```
    ::= EITHER (CONS Int)  
          (EITHER (Cons Bool) (Cons Char))
```

both work equally well,  
but be consistent

# serialization for generic types

**instance** serialize UNIT where

```
write _ c          = ["UNIT": c]
read ["UNIT": c] = Just (UNIT, c)
read _            = Nothing
```

- do we need to write:

```
write UNIT c      = ["UNIT": c]
```

- similar for the other generic types

# generic generalization

- idea: transform [7] to

```
RIGHT (CONS "Cons" (PAIR 7 (LEFT (CONS "Nil" UNIT))))
```

- using

```
:: ListG a ::= EITHER (CONS UNIT) (CONS (PAIR a
[a]))
```

- this is printed like

```
[ "RIGHT", "( ", "CONS", "Cons", "( ", "PAIR", "7", "( ", "LEFT",  
    "( ", "CONS", "Nil", "( ", "UNIT", ")", ") ", ") ", ") "]
```

- reading such a list is quite easy for the function `read`

- the strings "LEFT" and "RIGHT" tell exactly what to do if there is any choice (in the case for EITHER)
- the class mechanism will select the appropriate instance!

# prettier printing

- instead of

```
[["RIGHT", "(", "CONS", "Cons", "(", "PAIR", "7", "(", "LEFT",  
    ", \"(\" , \"CONS\", \"Nil\", \"(\" , \"UNIT\", \")\" , \")\" , \")\" , \")\" , \")\" ]]
```

we might prefer the representation

```
[ "( , "Cons" , "( , "7" , "( , "( , "Nil" , ")" , ")" , ")" ]
```

- in the write this is quite easy

- for the read we need to backtrack

- we cannot decide based on "LEFT" and "RIGHT"
- just try LEFT and if the arguments of the constructor are not there try the RIGHT branch

- for CONS we just assume that the string in the input is the constructor name

# reading with backtracking

- something like (as part of a class `serialize`)

```
read :: [String] -> Either a b | read a &  
read b
```

```
read list = case read list of  
    Just (a,m) = Just (LEFT a, m)  
    Nothing    = case read list of  
        Just (b,m) = Just (RIGHT b, m)  
        Nothing    = ..
```



- yes, this looks like magic  
based on the types the compiler selects the right read

## even prettier printing

- after these changes the `[]` will be printed as

```
["(", "Nil", ")"]
```

the result `["Nil"]` is nicer

- but, for `[7]` we want

```
["(", "Cons", "7", "Nil", ")"]
```

- how do we detect this in the generic representation?

- can we do a pattern match? like

```
write (CONS n UNIT) c = [n:c]
```

```
write (CONS n a) c = ["(", n: write a [")"] : c]
```

if not, how can this be done?