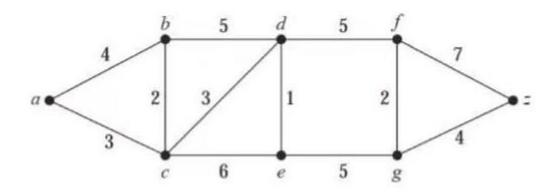
第二次理论作业题解

1 搜索

1. 在下面的带权无向图中找到 a 到 z 的最短路径。请使用带环检测的一致代价搜索算法求解, 并画出搜索树。(*注意:请在搜索树的节点旁边标明扩展顺序)



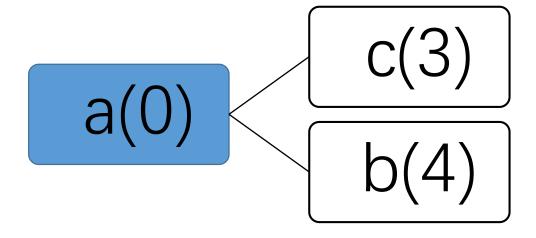
■ 附加题: 试分析Dijkstra算法和带环检测的一致代价搜索算法之间的关系。

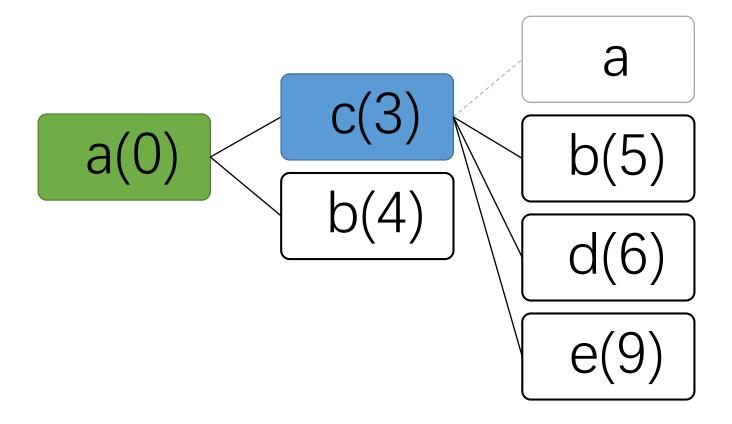
以下为搜索过程。作业画出最后的树即可

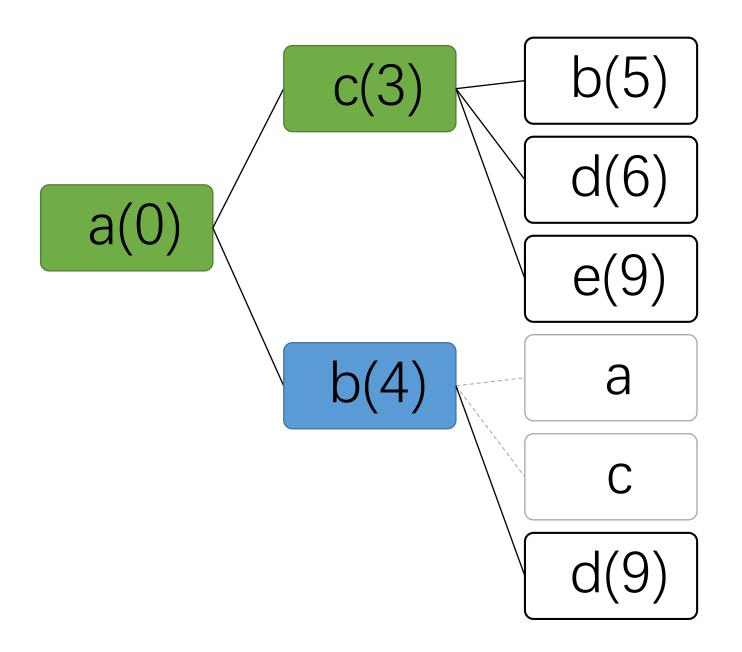
绿:已扩展蓝:正扩展

深灰: 入队但可不扩展

虚线节点: 环节点



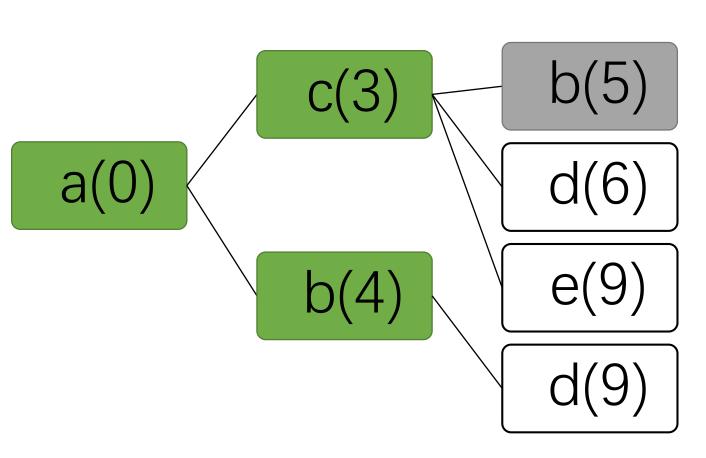


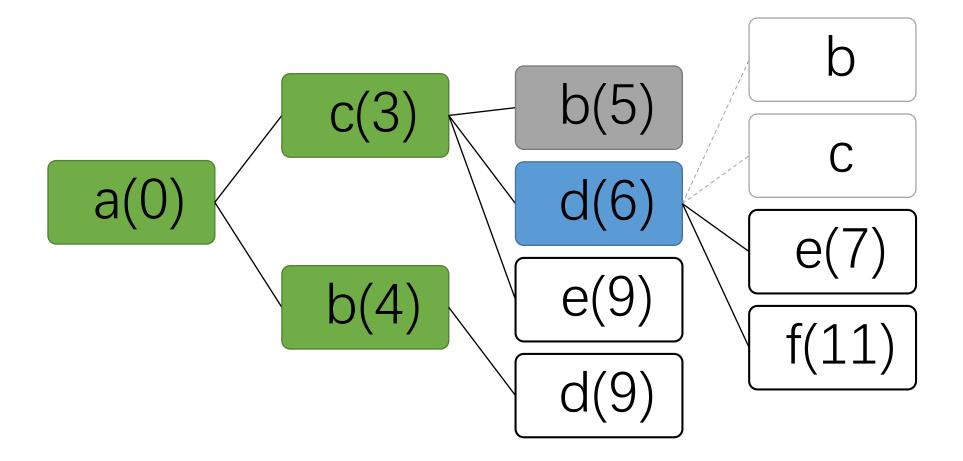


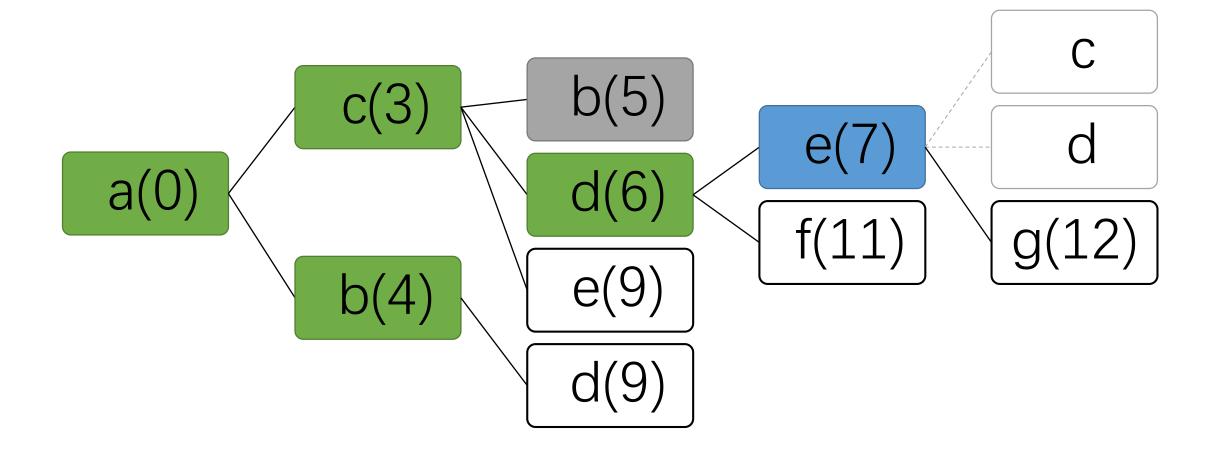
此时无需扩展b(5)。

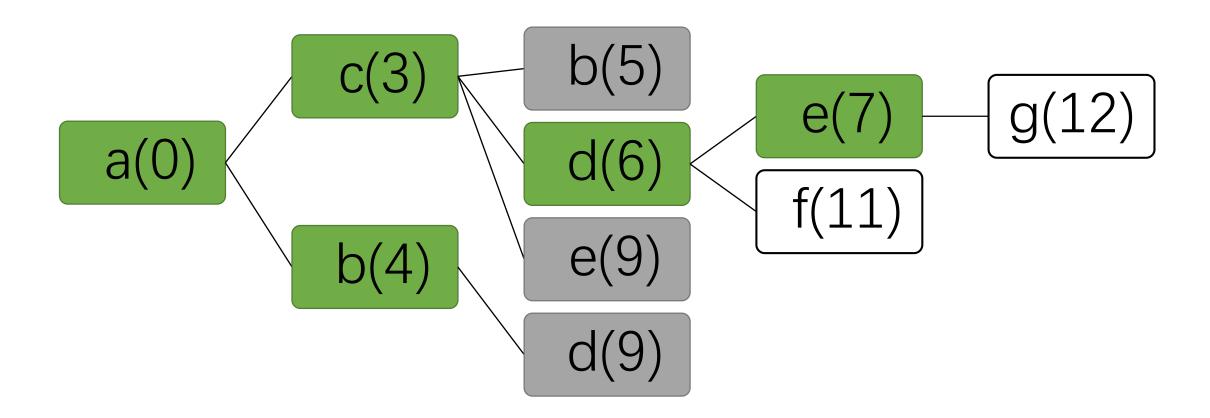
另一种理解是: 把环检测只理解成"后继已被扩展过,则不将该后继入队",也就是说在扩展节点时不再判断它是否被扩展过(换种说法就是节点只要入队就可以被扩展)。这种理解下也可以扩展b(5)。

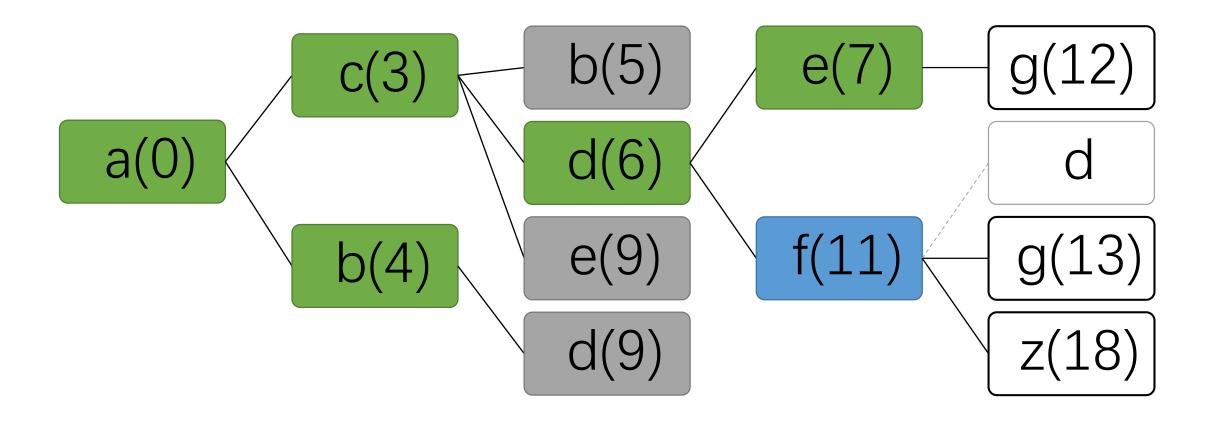
事实上扩展过的节点就不用 再扩展了,因为第一次扩展 的时候已经找到了到该节点 的最小路径。

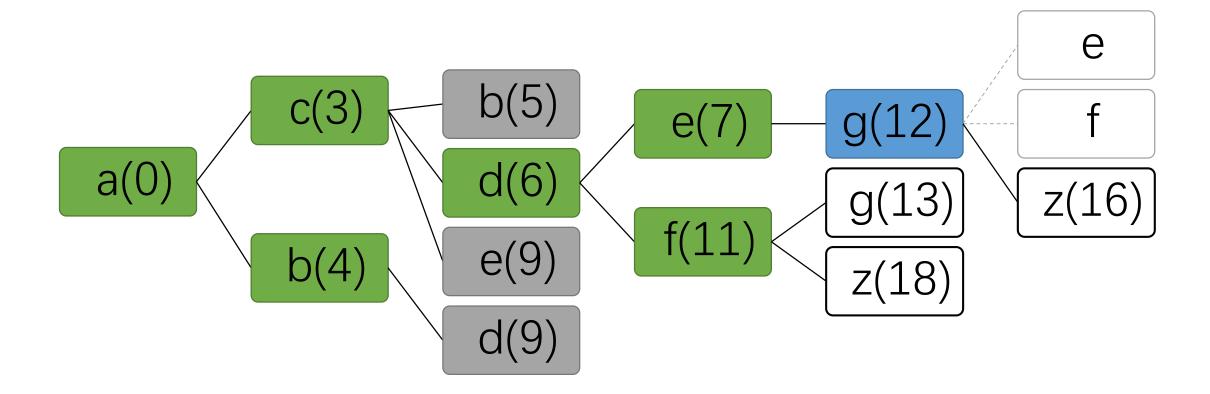


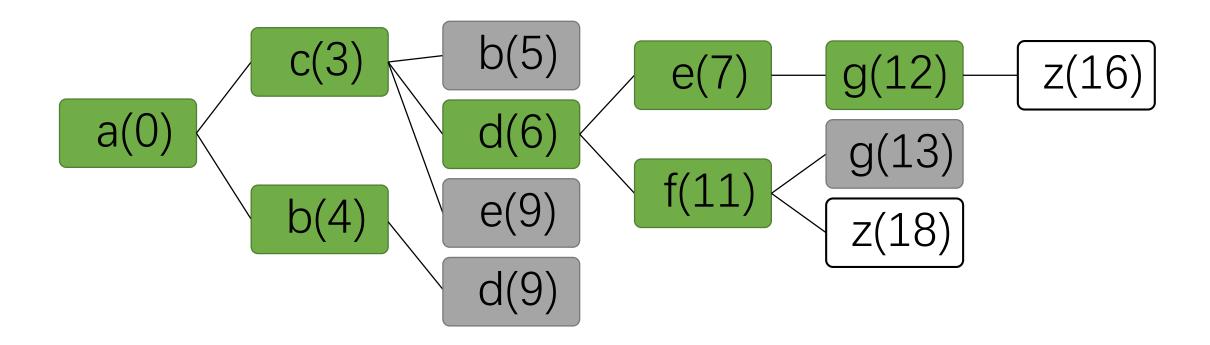




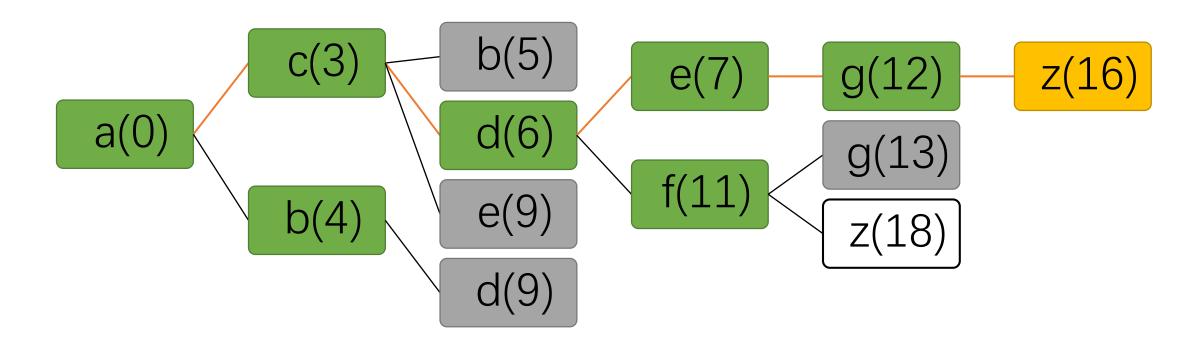


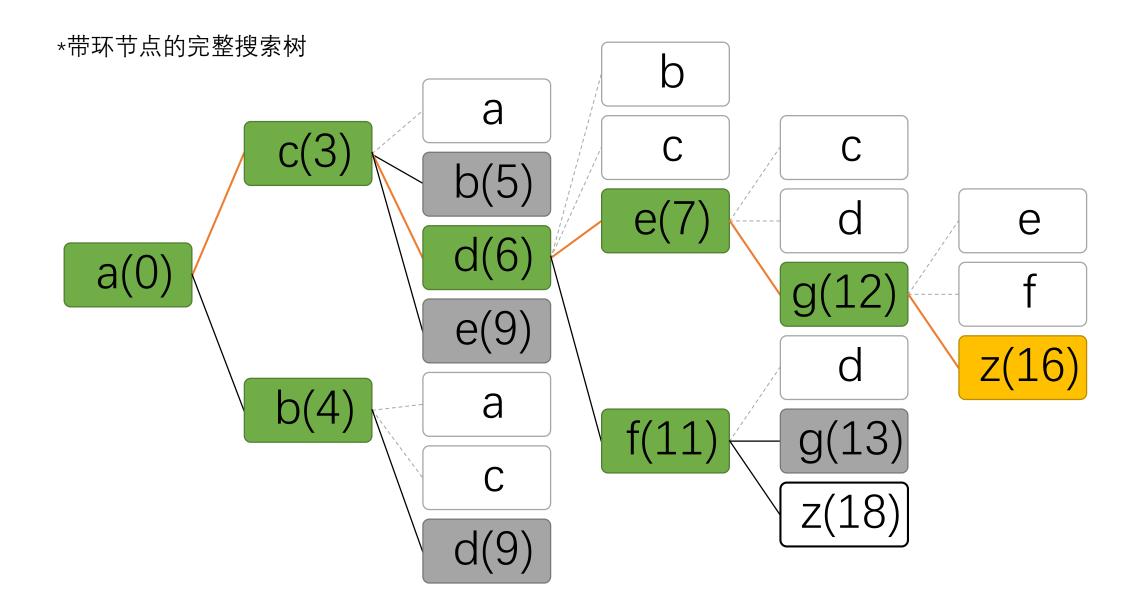






*画出本图即可



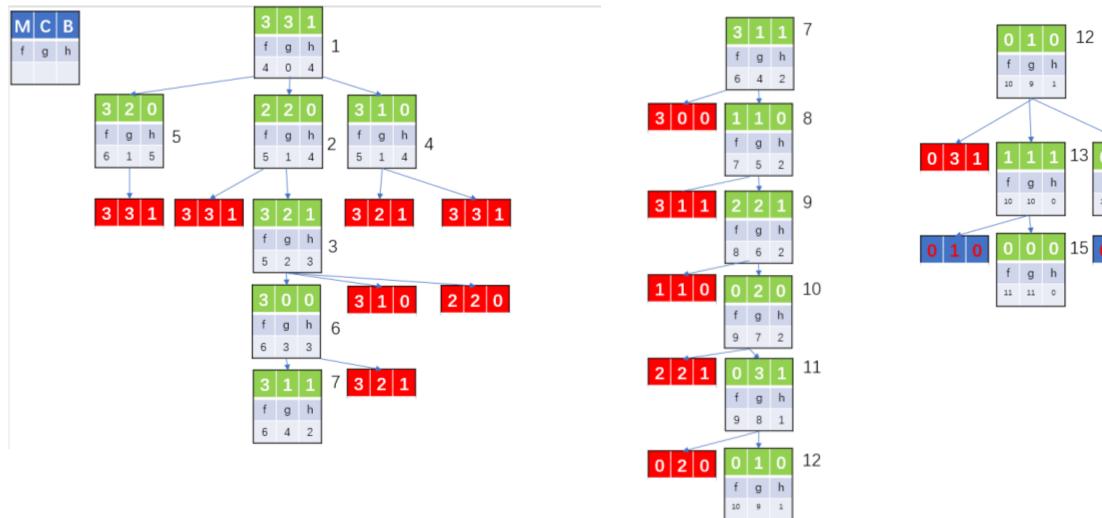


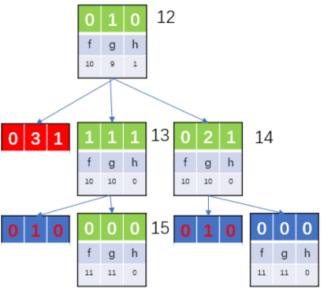
2 过河问题

2. 传教士和食人族过河问题(详见课件): 开始时左岸有M=3个传教士和C=3个食人族。船的容量为K=2。请将h(n)=M+C-2B作为启发式函数,用带环检测的A*搜索来求解。请画出搜索树,并标出每个节点的h值和g值。(*注意:请在搜索树的节点旁边标明扩展顺序)提醒: 1、传教士和食人族都可以划船,但空船不能来回河两岸; 2、请时刻关注状态是否合法。

- 明确在M≥C或C=0的约束下的可行状态: 两岸都要满足"M≥C或C=0"。
- 可行状态: <3,3,B>, <3,2,B>, <3,1,B >, <3,0,B >, <2,2,B >, <2,0,B >, <1,3,B >, <1,1,B >, <0,3,B >, <0,2,B >, <0,1,B >, <0,0,B >
- 禁止状态: <2,3,B >, <2,1,B >, <1,2,B >, <1,0,B >
- •船:K=2,任意组合均可。

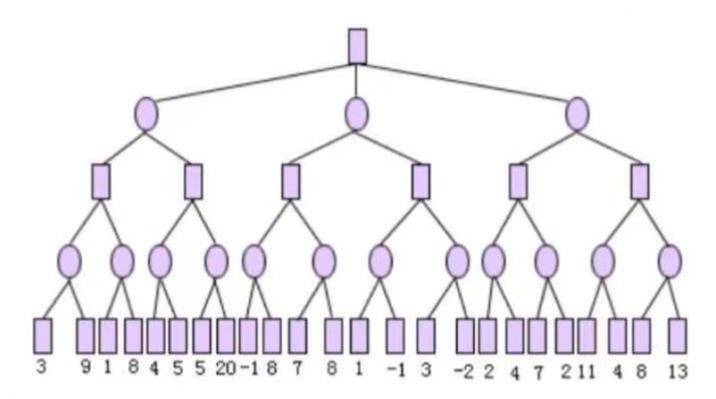
• 感谢谢泽辉同学的答案





3 alpha-beta剪枝

3. 对下面的博弈树使用 $\alpha - \beta$ 剪枝算法,并标出节点值的变化过程。



Implementing Alpha-Beta Pruning

```
AlphaBeta (n, Player, alpha, beta) //return Utility of state
If n is TERMINAL
    return V(n) //Return terminal states utility
ChildList = n.Successors(Player)
If Player == MAX
 for c in ChildList
    alpha = max(alpha, AlphaBeta(c, MIN, alpha, beta))
    If beta <= alpha</pre>
        break
 return alpha
Else //Player == MIN
  for c in ChildList
    beta = min(beta, AlphaBeta(c, MAX, alpha, beta))
    If beta <= alpha
        break
 return beta
```

When AlphaBeta(n,Player,alpha,beta) is called, alpha is the maximum alpha value of n's ancestor Max nodes, and beta is the minimum beta value of n's ancestor Min nodes

Initial call: AlphaBeta(START-NODE, Player, -infinity, +infinity)

