



ETRO  
ELECTRONICS &  
INFORMATICS

# Operating Systems and Security: Security

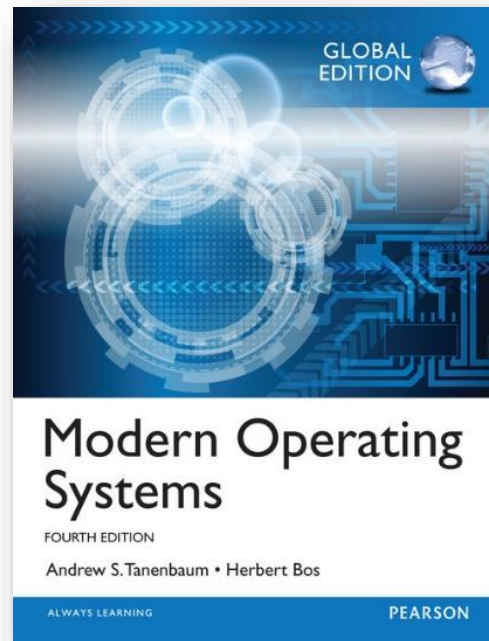
Bruno da Silva

E-mail: [bdasilva@etrovub.be](mailto:bdasilva@etrovub.be)

# CURSUS MATERIAL

## SECURITY

- Protection
- Attacks
- Malware
- Defenses



- Reference Book:
  - **A. S. Tanenbaum, "Modern Operating Systems," Pearson, Chapter 9**  
**(Cryptography or Authentication part are covered in another course)**



ETRO  
ELECTRONICS &  
INFORMATICS

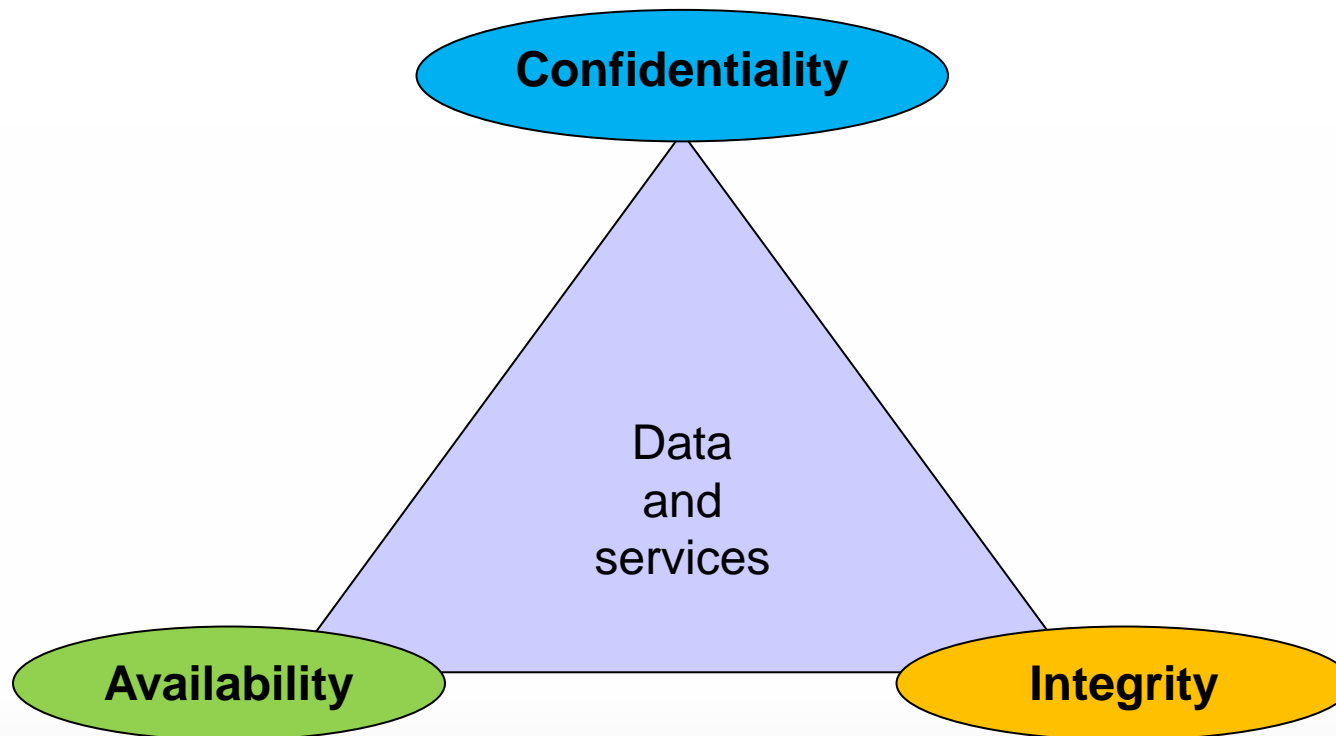
# OPERATING SYSTEMS

SECURITY – PART 1: PROTECTION

# COMPUTER SECURITY TRIAD

## THREADS

Three key objectives are at the heart of computer security



# CAN WE BUILD SECURE SYSTEMS?

Two questions concerning security:

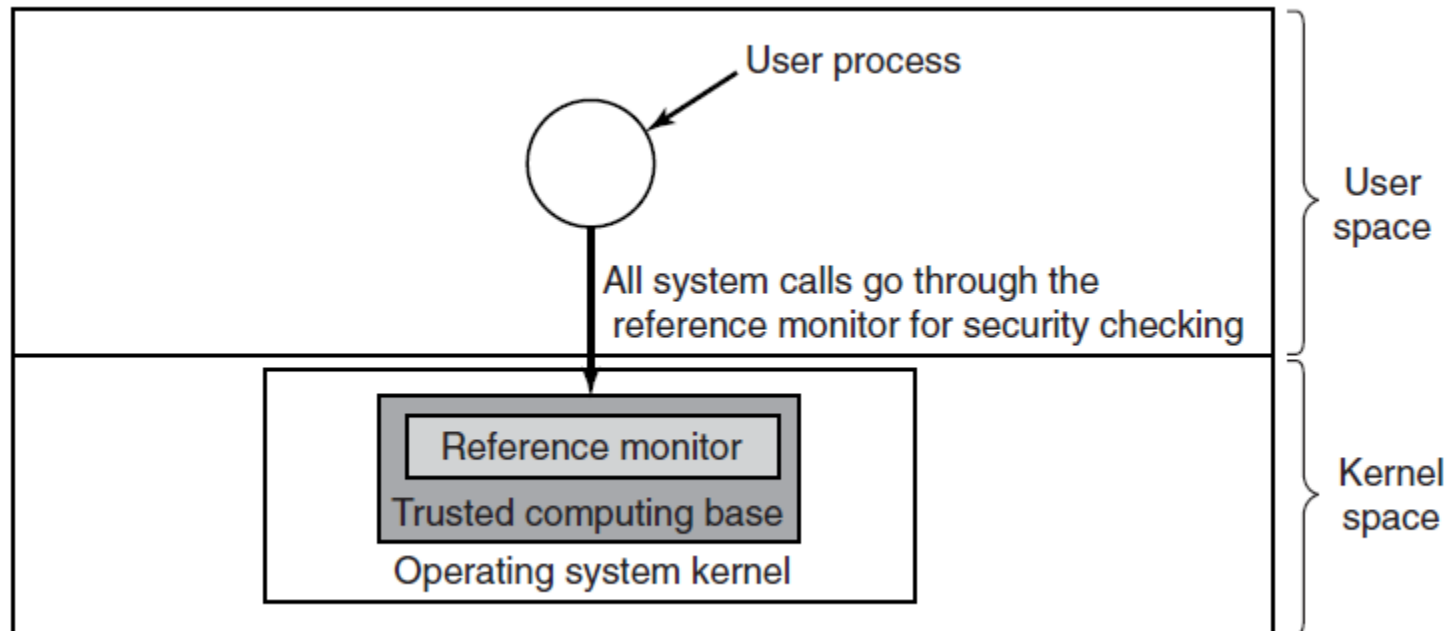
1. Is it possible to build a secure computer system?
2. If so, why is it not done?

# TRUSTED COMPUTING BASE

## DEFINITION

Reference monitor

Accepts all system call security and makes security decisions



# PROTECTION DOMAINS

## MODEL

Domain is a collection of access rights

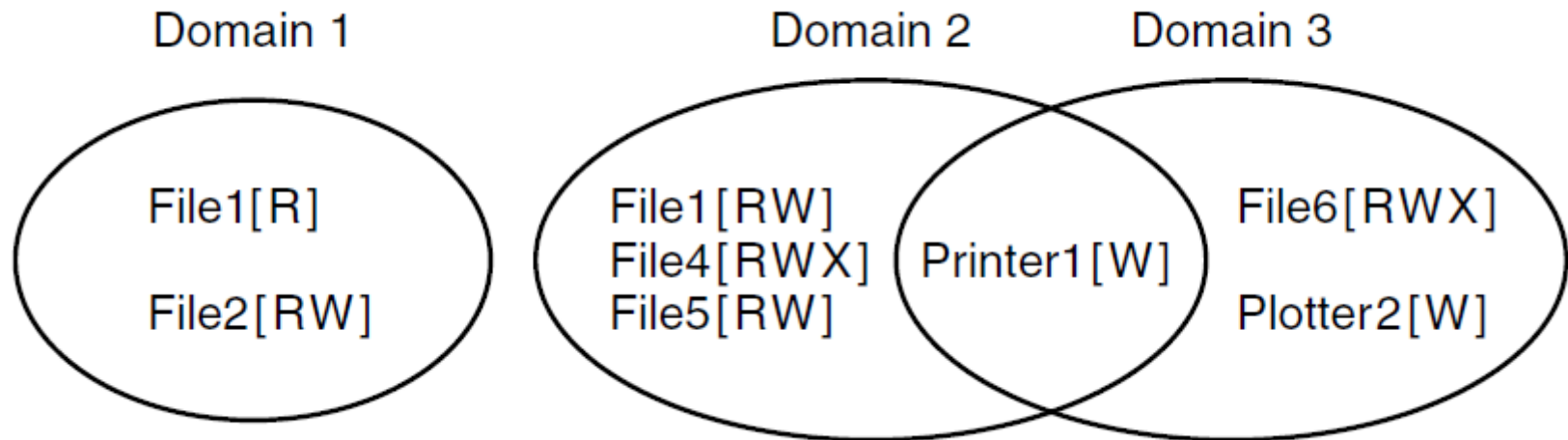
Domain is a set of (object, rights) pairs.

- Each pair specifies an object and some operation that can be performed on it.
- Right, in this context, means the permission to perform one of the operations.

# PROTECTION DOMAINS

## EXAMPLE

Three protection domains.





## CONTROLLING ACCESS TO RESOURCES

### PRINCIPLE OF LEAST AUTHORITY/PRIVILEGE

**Principle of Least Authority:** Security works best when each domain has the minimum objects and privileges to do its work and no more.

At each abstraction layer, every element should be able to access only the resources necessary to perform its task.

# PROTECTION DOMAINS

## A PROTECTION MATRIX

Domain	Object							
	File1	File2	File3	File4	File5	File6	Printer1	Plotter2
1	Read	Read Write						
2			Read	Read Write Execute	Read Write		Write	
3						Read Write Execute	Write	Write

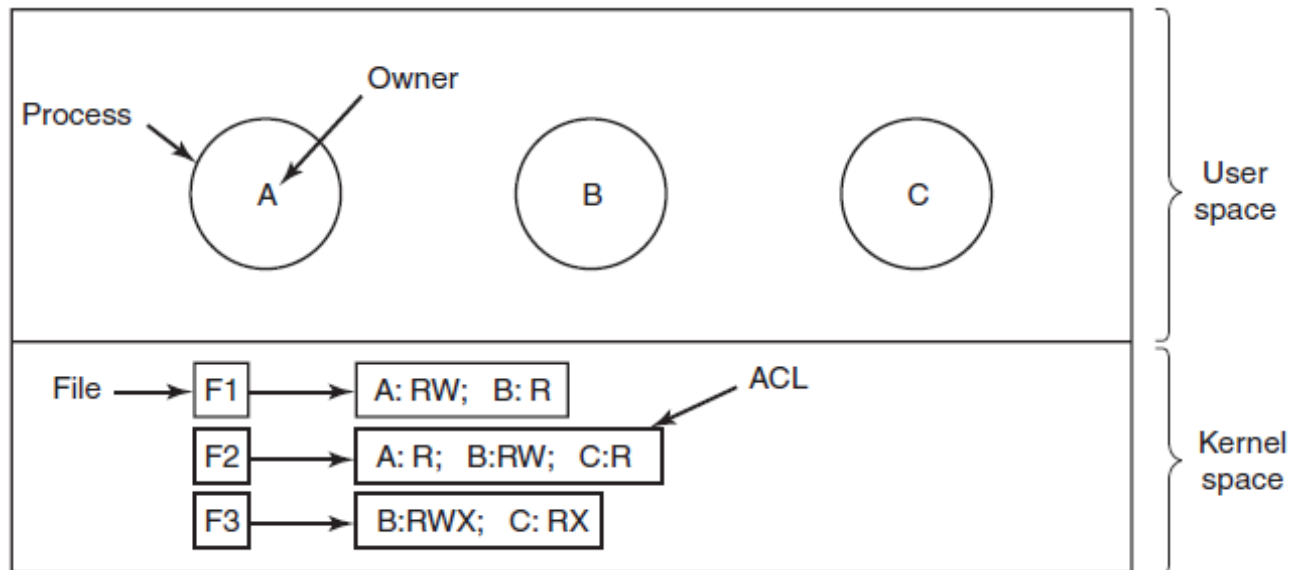
## TWO METHODS OF IMPLEMENTATION

- Access Control lists
- Capabilities list

## CONTROLLING ACCESS TO RESOURCES

### ACCESS CONTROL LISTS

**Access control lists** associate each object with an ordered list containing all the domains that may access the object and how.



Many systems support the concept of a group of users. Groups have names and can be included in ACLs.

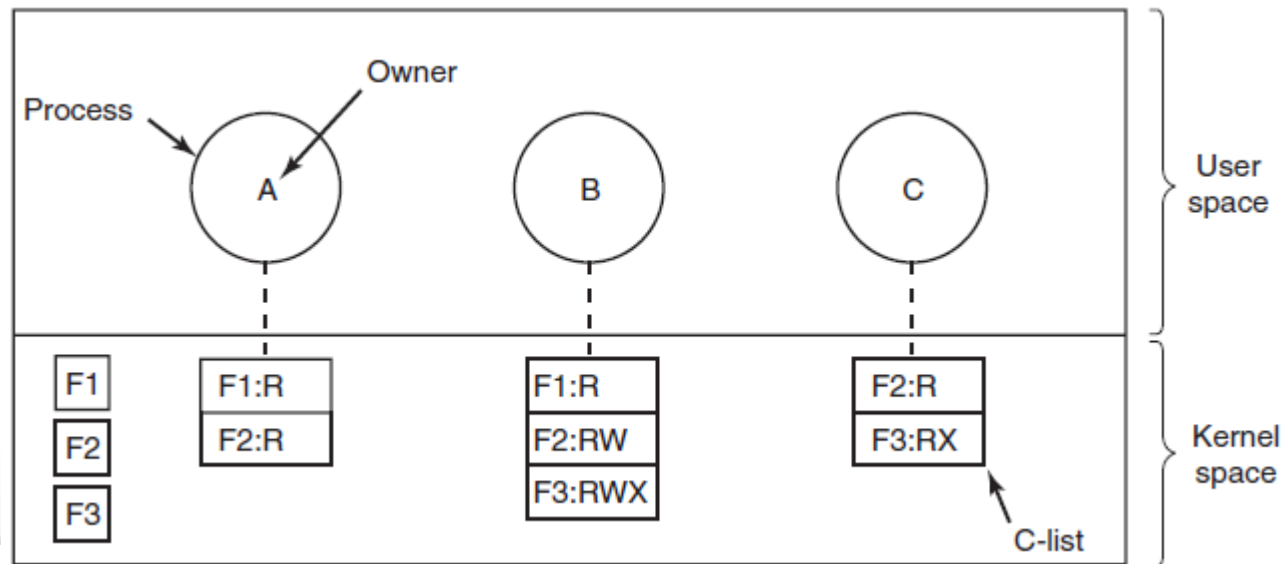
# CONTROLLING ACCESS TO RESOURCES

## CAPABILITIES

When capabilities are used, each process has a capability list.

Each element of this list is called capability, which grants the owner certain rights on a certain object.

Capability lists are themselves objects and may be pointed to from other capability lists, thus facilitating sharing of subdomains.



# FORMAL MODELS OF SECURE SYSTEMS

## PROTECTION MATRICES ARE NOT STATIC

(a) An authorized state.

(b) An unauthorized state.

		Objects		
		Compiler	Mailbox 7	Secret
Eric	Read	Read		
	Execute	Execute		
Henry	Read	Read	Read	
	Execute	Execute	Write	
Robert	Read	Read		Read
	Execute	Execute		Write

(a)

		Objects		
		Compiler	Mailbox 7	Secret
Eric	Read	Read		
	Execute	Execute		
Henry	Read	Read	Read	
	Execute	Execute	Write	
Robert	Read	Read	Read	Read
	Execute	Execute		Write

(b)

Given an initial authorized state and a set of commands, can it be proven that the system can never reach an unauthorized state?

# MULTILEVEL SECURITY

## Discretionary Access Control

- Operating systems allow individual to determine who may read and write their files

## Mandatory Access Control

- Organizational rules state who can see and modify what.

# MULTILEVEL SECURITY

## BELL-LAPADULA MODEL

Bell-LaPadula Model rules for information flow:

1. The simple security property

- Process running at security level  $k$  can read only objects at its level or lower

2. The  $*$  property

- Process running at security level  $k$  can write only objects at its level or higher

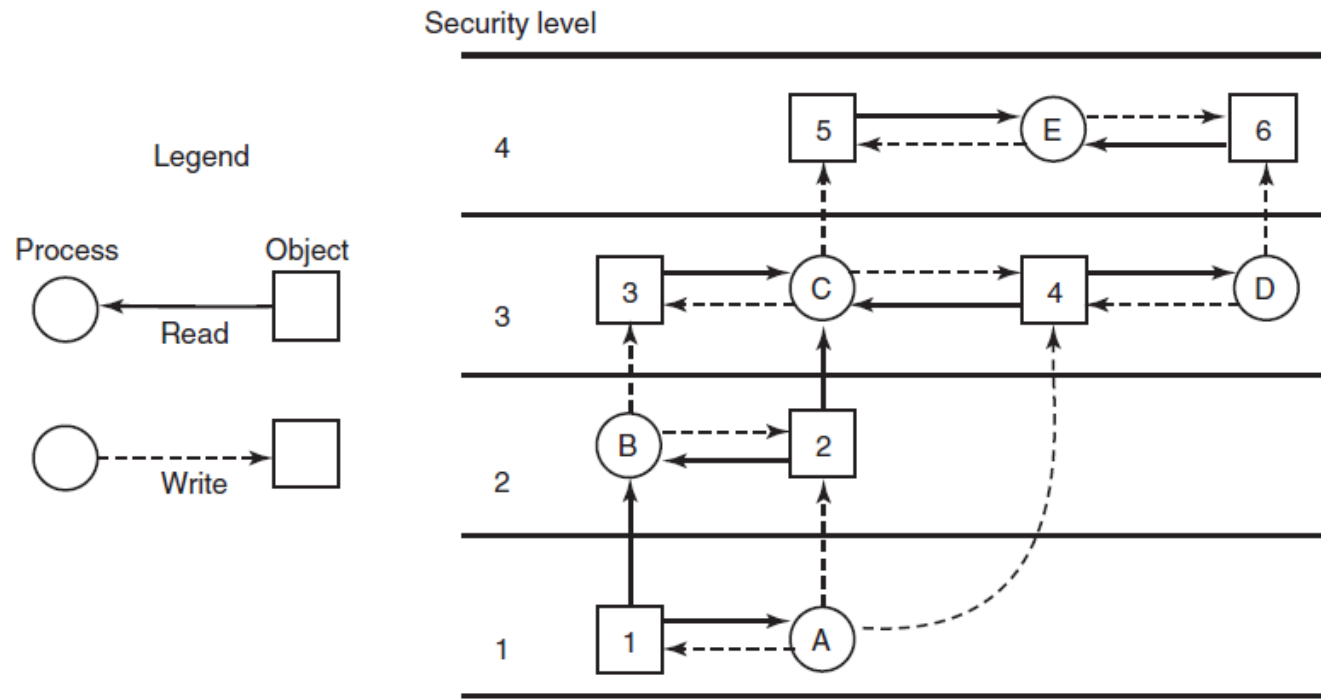
**Fun Fact:** The  $*$  property was so named because in the original report, the authors could not think of a good name for it and used  $*$  as a temporary placeholder until they could devise a better name. They never did and the report was printed with the  $*$ .



# MULTILEVEL SECURITY

## BELL-LAPADULA MODEL

The Bell-LaPadula multilevel security model.



The Bell-LaPadula multilevel security model was devised to keep secrets, not guarantee the integrity of data.

# THE BIBA MODEL

To guarantee the integrity of the data:

## 1. The simple integrity principle

- process running at security level  $k$  can write only objects at its level or lower (no write up).

## 2. The integrity \* property

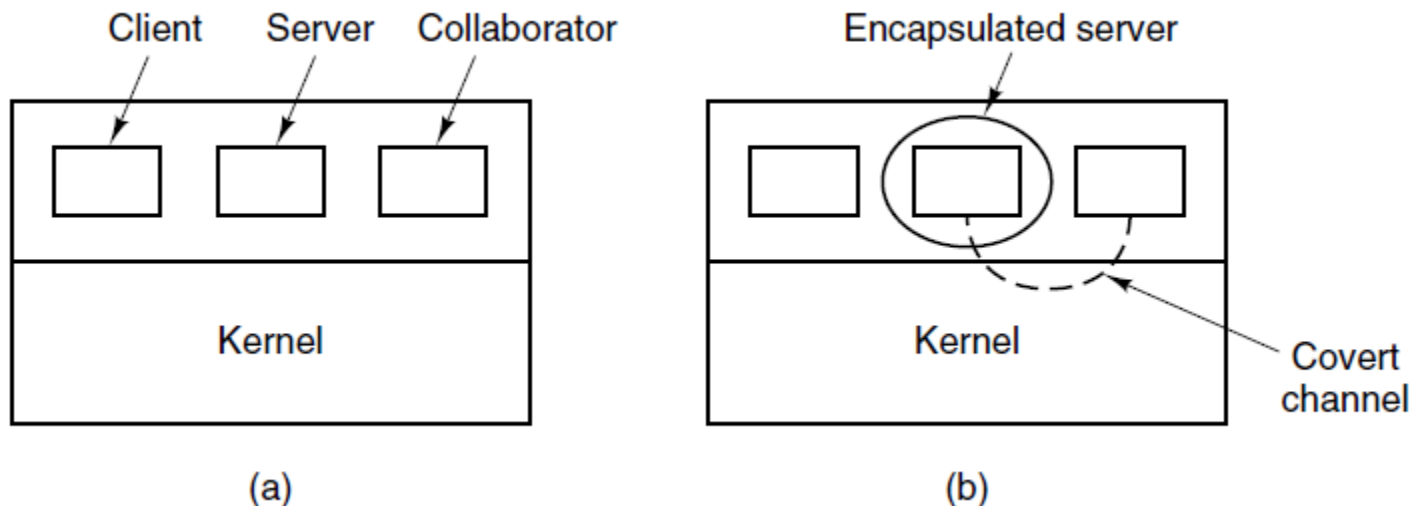
- process running at security level  $k$  can read only objects at its level or higher (no read down).

BUT... security leaks still can occur

# COVERT CHANNELS

**Lampson's model**, showing how information leaks, was originally formulated in terms of a single timesharing system, but the same ideas can be adapted to applications running in the cloud.

(a) The client, server, and collaborator processes. (b) The encapsulated server can still leak to the collaborator via covert channels.



## COVERT CHANNELS (2)

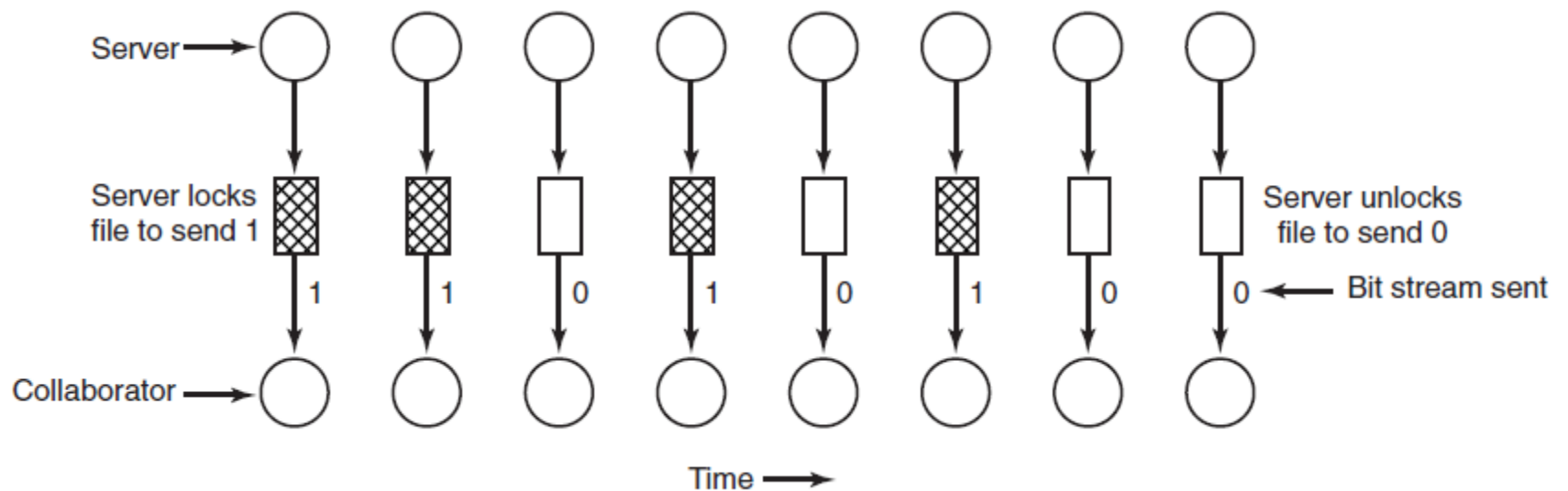


Figure 9-13. A covert channel using file locking.

# STEGANOGRAPHY

## Concealing secrets in plain site

- Hiding secret information in paintings, jpegs, MP3 files... or messages... or the network (e.g. Watermarking)
- Generic description of the pieces of the steganographic process:
  - $\text{cover\_medium} + \text{hidden\_data} + \text{stego\_key} = \text{stego\_medium}$

# STEGANOGRAPHY

## EXAMPLE

The German Embassy in Washington, DC, sent these messages in telegrams to their headquarters in Berlin during World War I (Kahn 1996).

PRESIDENT'S EMBARGO RULING SHOULD HAVE IMMEDIATE NOTICE.  
GRAVE SITUATION AFFECTING INTERNATIONAL LAW. STATEMENT  
FORESHADOWS RUIN OF MANY NEUTRALS. YELLOW JOURNALS  
UNIFYING NATIONAL EXCITEMENT IMMENSELY.

**Do you see hidden the message?**

# STEGANOGRAPHY

## EXAMPLE

The German Embassy in Washington, DC, sent these messages in telegrams to their headquarters in Berlin during World War I (Kahn 1996).

PRESIDENT'S EMBARGO RULING SHOULD HAVE IMMEDIATE NOTICE.  
GRAVE SITUATION AFFECTING INTERNATIONAL LAW. STATEMENT  
FORESHADOWS RUIN OF MANY NEUTRALS. YELLOW JOURNALS  
UNIFYING NATIONAL EXCITEMENT IMMENSELY.

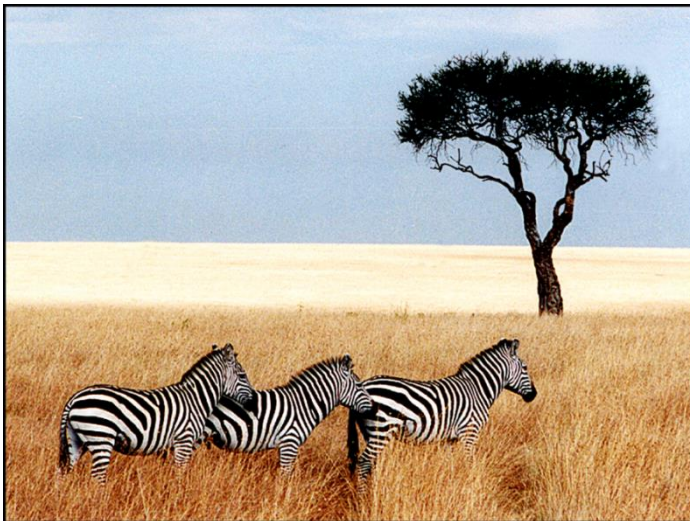
**PERSHING SAILS FROM N.Y. JUNE 1**

More examples: <https://www.garykessler.net/library/steganography.html>

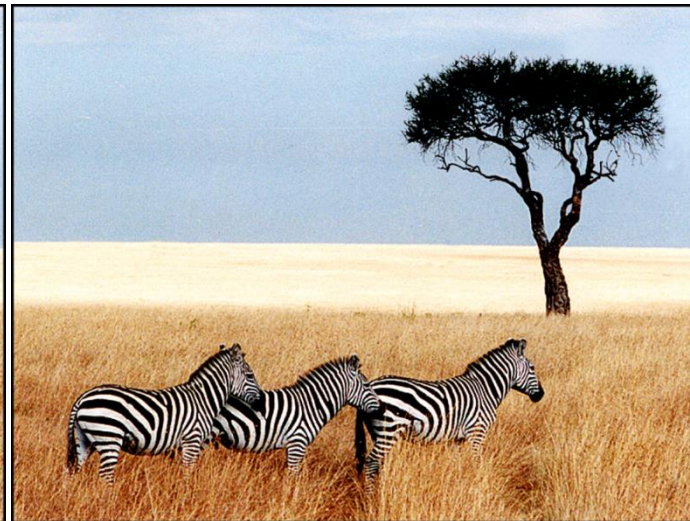
# STEGANOGRAPHY

## EXAMPLE

- Pictures appear the same
- Picture on right has text of 5 Shakespeare plays
  - encrypted, inserted into low order bits of color values



(a)



(b)

a) Three zebras and a tree.

b) Three zebras, a tree, and the complete text of five plays by William Shakespeare.





ETRO  
ELECTRONICS &  
INFORMATICS

# OPERATING SYSTEMS

SECURITY – PART 2: ATTACKS

# WHAT WE ARE COVERING

- Exploiting Software
- Insider Attacks

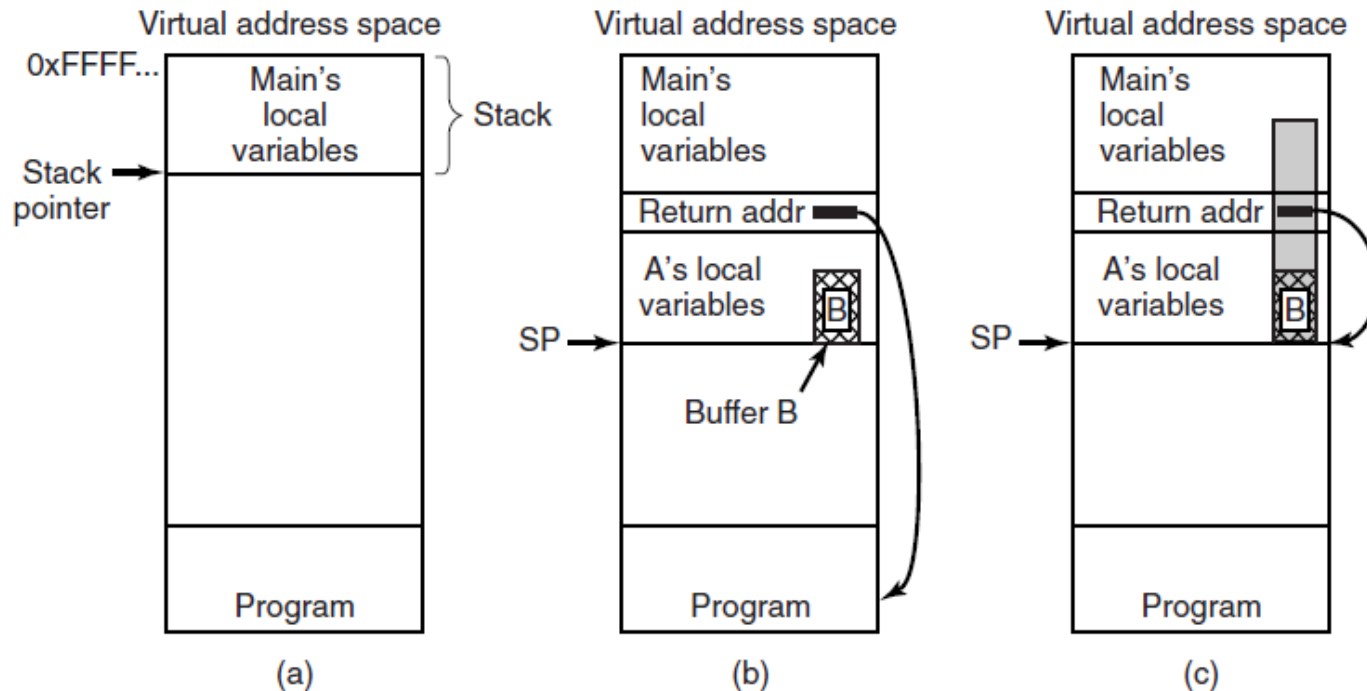
# EXPLOITING SOFTWARE

## TYPES

- Buffer overflow attacks
  - Control Flow Attacks
  - Non-Control Flow Attacks
- Memory Corruption Attacks
  - Format String Attack
  - Dangling pointers
- Null Pointer Dereference Attacks
- Integer Overflow attacks
- Command Injection Attacks
- Time of Check to Time of Use Attacks

# BUFFER OVERFLOW ATTACKS

NO C OR C++ COMPILER DOES ARRAY BOUNDS CHECKING



(a) Situation when the main program is running.

(b) After the procedure A has been called.

(c) Buffer overflow shown in gray.

## DEFENSE: STACK CANARIES

Modern computers use digital 'canaries' as an early warning system, detecting possible buffer overflow attack.

In the code, at places where the program makes a function call, the compiler inserts code to save a random 'canary' value on the stack, just below the return address.

The compiler inserts code at the return to check the canary value.

If the canary has changed... trouble!

# DEFENSE: DATA EXECUTION PREVENTION

## AGAINST CODE INJECTION ATTACKS

Redefine the 'real' problem: The fact that the attacker can inject code and have it executed in the heap or stack!

Defense: prevent the bytes provided by the attacker from being executed as legitimate code.

Modern CPUs have a feature **NX bit** (No-eXecute bit):

- It distinguishes between data segments and the code segments.
- This ensures data segments are writable, but not executable.

A generic name for this security measure is **DEP** (Data Execution Prevention).

# CODE REUSE ATTACKS

Attacker constructs the necessary functions out of existing functions and instructions in the existing binaries and libraries

## Return to libc

Almost all c (C++) programs are linked with a shared library containing the function *system*.

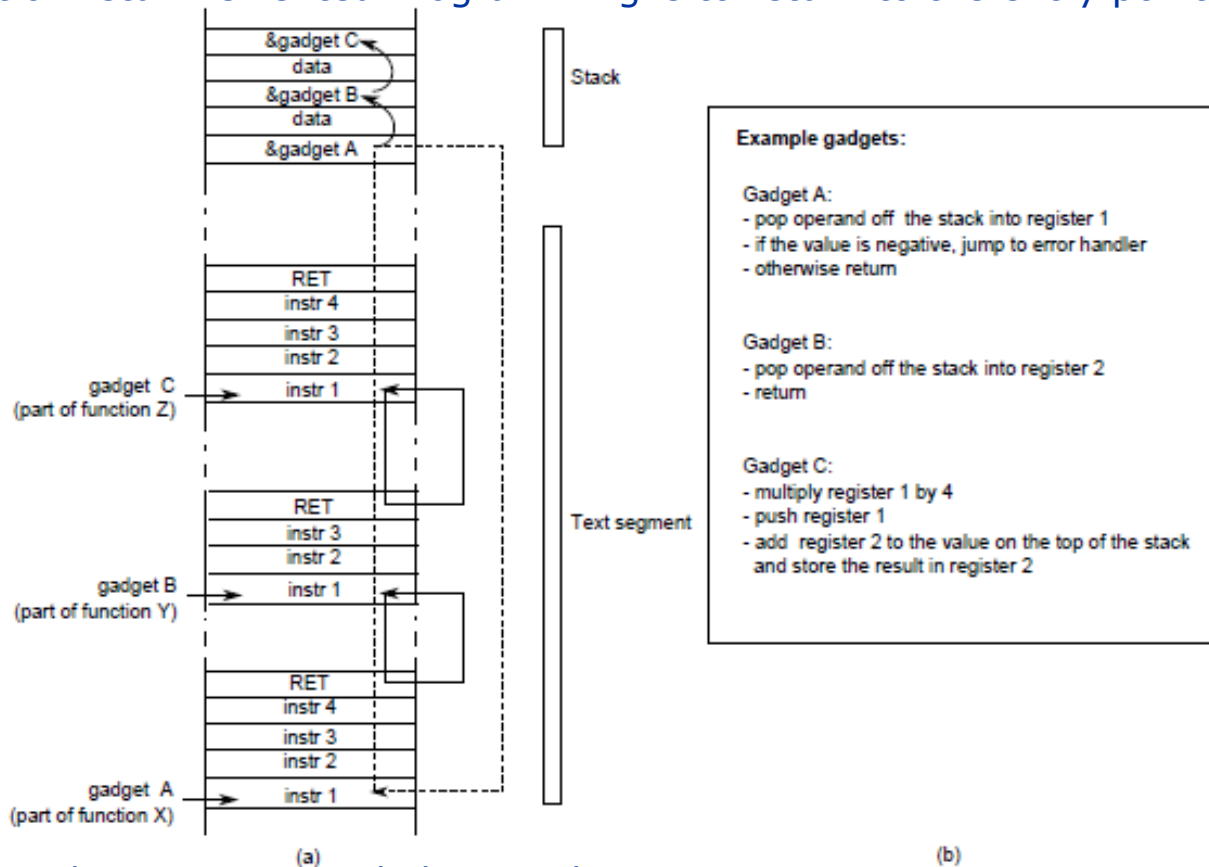
*System* takes a string containing a command and passes it to the shell for execution.

Attack: Place a string containing commands to be executed and divert control the *system* function via the return address.

# CODE REUSE ATTACKS

## RETURN-ORIENTED PROGRAMMING (ROP)

The concept of Return-Oriented Programming is to return to the entry points of library functions.



Return-oriented programming: linking gadgets



# DEFENSE: ADDRESS SPACE LAYOUT RANDOMIZATION (ASLR)

- Randomize the address of functions and data between every run of the program
- Supported with varying granularity.
- Few apply it to the system kernel

# NON-CONTROL FLOW ATTACKS

- Change the data instead of the return addresses!!
- Possible use: change the security credentials granting more or less protection on objects.

```
01. void A() {  
02.   int authorized;  
03.   char name [128];  
04.   authorized = check_credentials (...); /* the attacker is not authorized, so returns 0 */  
05.   printf ("What is your name?\n");  
06.   gets (name);  
07.   if (authorized != 0) {  
08.     printf ("Welcome %s, here is all our secret data\n", name)  
09.     /* ... show secret data ... */  
}
```

# FORMAT STRING ATTACKS

In a C functions that perform formatting, (printf), **when the application doesn't properly validate the input**, the attacker can causes the submitted data to be compromised.

EX: By using exactly the right number of **%08x**, the attacker can use the first four characters of the format string as an address.

# FORMAT STRING ATTACKS

## EXAMPLE

Do you see the problem? Are the codes no equivalent?

```
#include <stdio.h>
```

```
int main(int argc, char *argv[])
```

```
{
```

```
    char* i = argv[1];
```

```
    printf("You wrote: %s\n", i);
```

```
}
```

Vs

```
#include <stdio.h>
```

```
#include <string.h>
```

```
int main(int argc, char *argv[])
```

```
{
```

```
    char test[1024];
```

```
    strcpy(test,argv[1]);
```

```
    printf("You wrote:");
```

```
    printf(test);
```

```
    printf("\n");
```

```
}
```

# FORMAT STRING ATTACKS

## EXAMPLE

So, what happens to the stack when a format string is specified with no corresponding variable on stack?

```
root@kali:~/Desktop/tuts/fmt# gcc fmt_test.c -o fmt_test
```

```
root@kali:~/Desktop/tuts/fmt# ./fmt_test test
```

```
You wrote: test
```

Vs

```
root@kali:~/Desktop/tuts/fmt# ./fmt_wrong testttt
```

```
You wrote:testttt
```

## FORMAT STRING ATTACKS

## EXAMPLE

So, what happens to the stack when a format string is specified with no corresponding variable on stack?

**It will start to pop data off the stack from where the variables should have been located.**

```
root@kali:~/Desktop/tuts/fmt# ./fmt_test $(python -c 'print "%08x"*20')
```

You wrote:

[illegible]

# Vs

```
root@kali:~/Desktop/tuts/fmt# ./fmt_wrong $(python -c 'print "%08x."*20')
```

You

wrote:bfd7469f.000000f0.00000006.78383025.3830252e.30252e78.252e7838.2e783830.78383025.  
3830252e.30252e78.252e7838.2e783830.78383025.3830252e.30252e78.252e7838.2e783830.7838  
3025.3830252e.

# DANGLING POINTERS ATTACK

- User frees memory. But later tries to access it with the pointer.
- Attacker places a specific heap object in the memory location the user frees and re-uses
- Techniques like **heap feng shui** help attackers pull this off.

# NULL POINTER DEREFERENCE ATTACKS

In linux, the kernel space is mapped to every process' address space... and whenever the kernel executes it runs in a process's address space.

If a buggy kernel dereferences a NULL pointer, it usually leads to a crash. Or attacker triggers a NULL pointer dereference from the user process.

Crash occurs because there is no code at page 0.

Using a tool like mmap (Posix function that maps files into memory) attacker can map bad code at that location.

Defense: mmap no longer makes it possible to map to page 0.



# INTEGER OVERFLOW ATTACK

Integer errors can happen when mathematical operations, and external input lead to a result that is too large to fit within the range of values that can be stored in variables of a given data type.

Can be exploited to corrupt applications.

Integer overflow attacks are possible largely due to incorrectly defining numerical data.

# COMMAND INJECTION ATTACKS

- The execution of arbitrary commands on the host operating system via a vulnerable application.
- Command injection attacks are possible largely **due to insufficient input validation** (a complete part of the Defensive Programming approach).

# COMMAND INJECTION ATTACKS

## EXAMPLE

```
int main(int argc, char *argv[])
{
    char src[100], dst[100], cmd[205] = "cp ";
    printf("Please enter name of source file: ");
    gets(src);
    strcat(cmd, src);
    strcat(cmd, " ");
    printf("Please enter name of destination file: ");
    gets(dst);
    strcat(cmd, dst);
    system(cmd);
}
```

/\* declare 3 strings \*/  
/\* ask for source file \*/  
/\* get input from the keyboard \*/  
/\* concatenate src after cp \*/  
/\* add a space to the end of cmd \*/  
/\* ask for output file name \*/  
/\* get input from the keyboard \*/  
/\* complete the commands string \*/  
/\* execute the cp command \*/

Suppose user enters for the destination file: 'file2.txt; rm -rf

Command executed is then: system("cp file1 file 2; rm -rf")

Command injection attacks are possible largely due to insufficient input validation.

# TIME OF CHECK TO TIME OF USE ATTACK

## TOCTOU

Exploits a race condition

Example:

```
int fd;  
if (access ("./my_document", W_OK) != 0)  
    exit (1);  
fd = open ("./my_document", O_WRONLY)  
write (fd, user_input, sizeof (user_input));
```

Between the statements, attacker creates a symbolic link with the same file name to the password file

Then user writes password information to attackers file.

# INSIDER ATTACKS

- Executed by programmers from within the company.
- Attackers have specialized knowledge
- Insider Attacks
  - Logic Bomb
  - Back Doors
  - Login Spoofing

# LOGIC BOMB

A **logic bomb** is a piece of code inserted into an operating system or software application that implements a malicious function after a certain time limit or specific conditions are met.

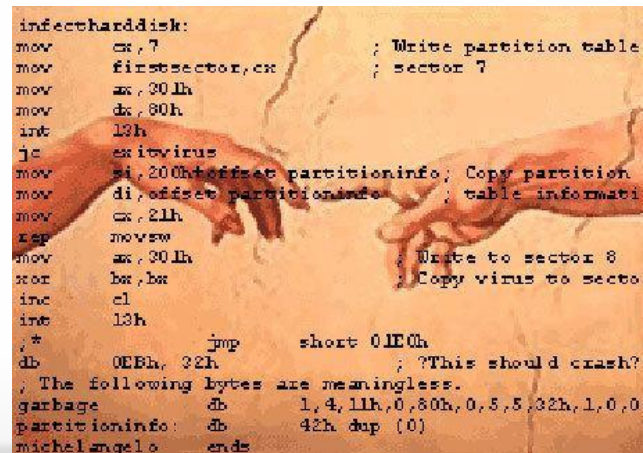
Company programmer writes program:

- potential to do harm
- OK as long as he/she enters password daily
- if programmer fired, no password and bomb explodes

# LOGIC BOMB

## EXAMPLE

- Michelangelo (logic bomb)
  - discovered on 4 February 1991
  - designed to infect DOS systems
  - the virus overwrites the first one hundred sectors of the hard disk with nulls.
  - each year, the virus remained dormant until March 6, the birthday of Renaissance artist Michelangelo.
  - The name was chosen by researchers who noticed the coincidence of the activation date



```
infectharddisk:
mov     cx, 7                ; Write partition table
mov     firstsector, cx      ; sector 7
mov     ax, 301h
mov     dx, 80h
int     13h
jc      exitvirus
mov     si, 200h; offset partitioninfo; Copy partition
mov     di, offset partitioninfo; table informati
mov     cx, 21h
rep     movsw
mov     ax, 301h             ; Write to sector 8
xor     bx, bx               ; Copy virus to secto
inc     cl
int     13h
; *      jmp     short 01E0h
db      0EBh, 32h           ; ?This should crash?
; The following bytes are meaningless.
garbage db      1, 4, 11h, 0, 80h, 0, 5, 5, 32h, 1, 0, 0
partitioninfo: db      42h dup (0)
michelangelo ends
```

# BACK DOORS

```
while (TRUE) {  
    printf("login: ");  
    get_string(name);  
    disable_echoing();  
    printf("password: ");  
    get_string(password);  
    enable_echoing();  
    v = check_validity(name, password);  
    if (v) break;  
}  
execute_shell(name);
```

(a)

(a) Normal code.

```
while (TRUE) {  
    printf("login: ");  
    get_string(name);  
    disable_echoing();  
    printf("password: ");  
    get_string(password);  
    enable_echoing();  
    v = check_validity(name, password);  
    if (v || strcmp(name, "zzzzz") == 0) break;  
}  
execute_shell(name);
```

(b)

(b) Code with a back door inserted.

**Defense:** Code reviews

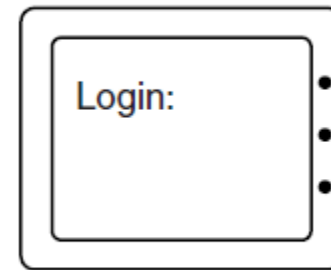


# LOGIN SPOOFING

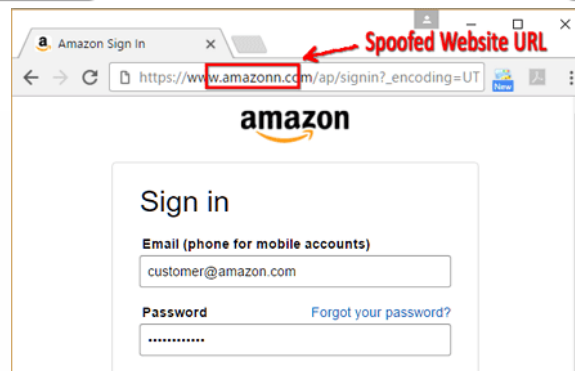
Legitimate user attempting to collect other people's passwords.



(a)



(b)



(a) Correct login screen.

(b) Phony login screen.

**Defense:** have a login sequence start with a key combination the user programs can't catch: CTRL-ALT-DEL in Windows cause the current user to log out and system login program started.



ETRO  
ELECTRONICS &  
INFORMATICS

# OPERATING SYSTEMS

SECURITY – PART 3: MALWARE

# MALWARE

- Malicious software
  - Trojan horses, virus, worms,.. Etc
- Today's malware is all about stealth
- Infected machines report back to attacker, its address, information...??
- Attacker uses backdoor to control the infected machine....
  - Make it a *zombie*.
  - A collection of zombies is called a *botnet*

# MALWARE

- Criminals can rent out botnets
- Keyloggers
- Identity theft
- Malware can lay in wait for something interesting
- Malware can interfere with competition's production process
- Malware could target another person in the company to discredit that person

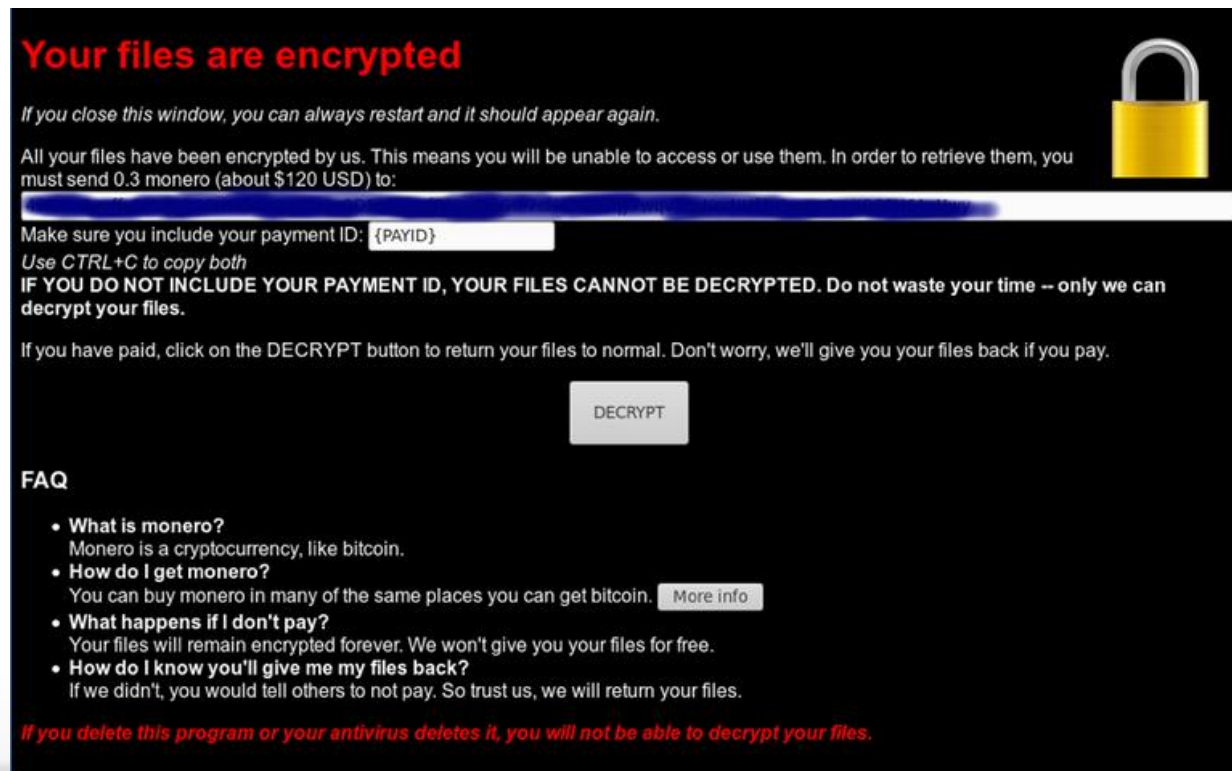
# TYPES OF MALWARE

- Ransomware
- Trojan Horse
- Virus
- Worm
- Spyware
- RootKits

# RANSOMWARE

Malware used for a form of blackmail. Some famous ones: WannaCry, SamSam...

Example: Encrypts files on victim disk (ransomware), then displays message ...



# RANSOMWARE

## EXAMPLES



## 2021 Trends Show Increased Globalized Threat of Ransomware

### SUMMARY

In 2021, cybersecurity authorities in the United States,<sup>[1][2][3]</sup> Australia,<sup>[4]</sup> and the United Kingdom<sup>[5]</sup> observed an increase in sophisticated,

Immediate Actions You Can Take Now to Protect Against Ransomware:

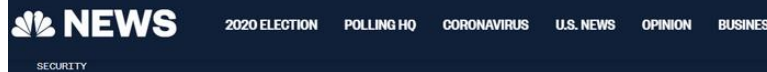
## Cyber-attack on Nvidia linked to Lapsus\$ ransomware gang

Adam Bannister 28 February 2022 at 16:30 UTC

Cyber-attacks US Ransomware



Claims that threat actors said hardware giant had 'hacked back' have surfaced



## Major hospital system hit with cyberattack, potentially largest in U.S. history

Computer systems for Universal Health Services, which has more than 400 locations, primarily in the U.S., began to fail over the weekend.

Sept. 28, 2020, 6:07 PM BST / Updated Sept. 28, 2020, 9:04 PM BST

By Kevin Collier

Sponsored Stories

by Taboola

A major hospital chain has been hit by what appears to be one of the largest medical cyberattacks in United States history.



MedStar Georgetown University Hospital

HACKERS PARALYZED HOSPITAL CHAIN



THIS IS A TOBACCO-FREE CAMPUS

Attacks 2021: <https://illinois.touro.edu/news/the-10-biggest-ransomware-attacks-of-2021.php>  
<https://www.nbcnews.com/tech/security/big-paydays-force-hospitals-prepare-ransomware-attacks-n557176>

# TROJAN HORSE

Transport means...Getting victims to download virus without attacker's intervention.

Free program made available to unsuspecting user

- Actually contains code to do harm

Place altered version of utility program on victim's computer

- trick user into running that program

Example:

- With hacking, music can take control of your car

<https://geeknizer.com/hack-car-with-music/>

<https://web.archive.org/web/20110316080851/http://www.itworld.com-can-take-control-your-car>





# VIRUSES

Virus is a program that can reproduce itself by attaching its code to another program.

Often written in assembler or C.

Attacker infects a program on his own machine, then gets that program distributed.

Once installed on victim's machine, it remains dormant until executed.

Once *activated*...

- Executes its **payload**
- Often waits for a specific date or time
- .... We want to make sure the virus is well distributed before people start noticing it.

# DIFFERENT KINDS OF VIRUSES

## 7 TYPES

1. Companion
2. Executable Program
3. Memory
4. Boot sector
5. Device Driver
6. Macro
7. Source code

# 1. COMPANION VIRUS

Old virus type

Runs with the program is supposed to run

Ex: in old MS-DOS

- We install a program named *prog.com*
- When user enters *prog*, instead of *prog.exe*, our infected program is executed.
- We'll call *prog.exe* after our malicious activity and no one will be the wiser

Can also be done with symbolic links

## 2. EXECUTABLE PROGRAM VIRUS

- Overwrites the executable program with itself.
- Virus written in assembly language
- Inserted into another program
  - use tool called a “dropper”
- Virus dormant until program executed
  - then infects other programs
  - eventually executes its “payload”

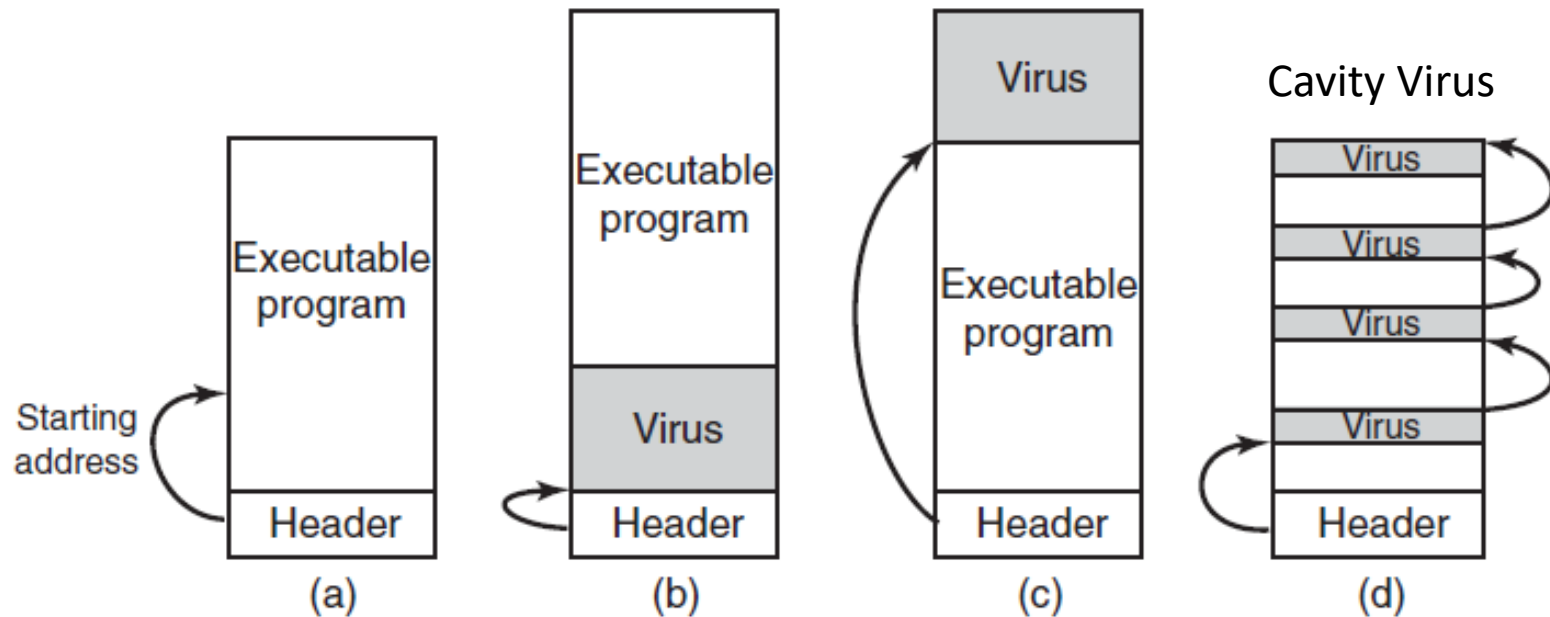
# EXECUTABLE VIRUS

Overwriting virus is easy to detect....

Parasitic virus: this virus attaches itself to the program to do the bad thing, but allows the program to function normally afterward.

# EXECUTABLE PROGRAM VIRUSES (3)

## EXAMPLES



(a) An executable program.

(b) With a virus at the front.

(c) With a virus at the end.

(d) With a virus spread over free space within the program.

### 3. MEMORY-RESIDENT VIRUSES

Stays in RAM, either hiding at the top of memory or down among the interrupt vectors (the last few hundred bytes are generally unused)

Capture one of the interrupt vectors

- Putting it's own address there
- Call the interrupt after it does what it does
- Benefit.. It can run in system mode

## 4. BOOT SECTOR VIRUS

Virus that overwrites the master boot record or boot sector.

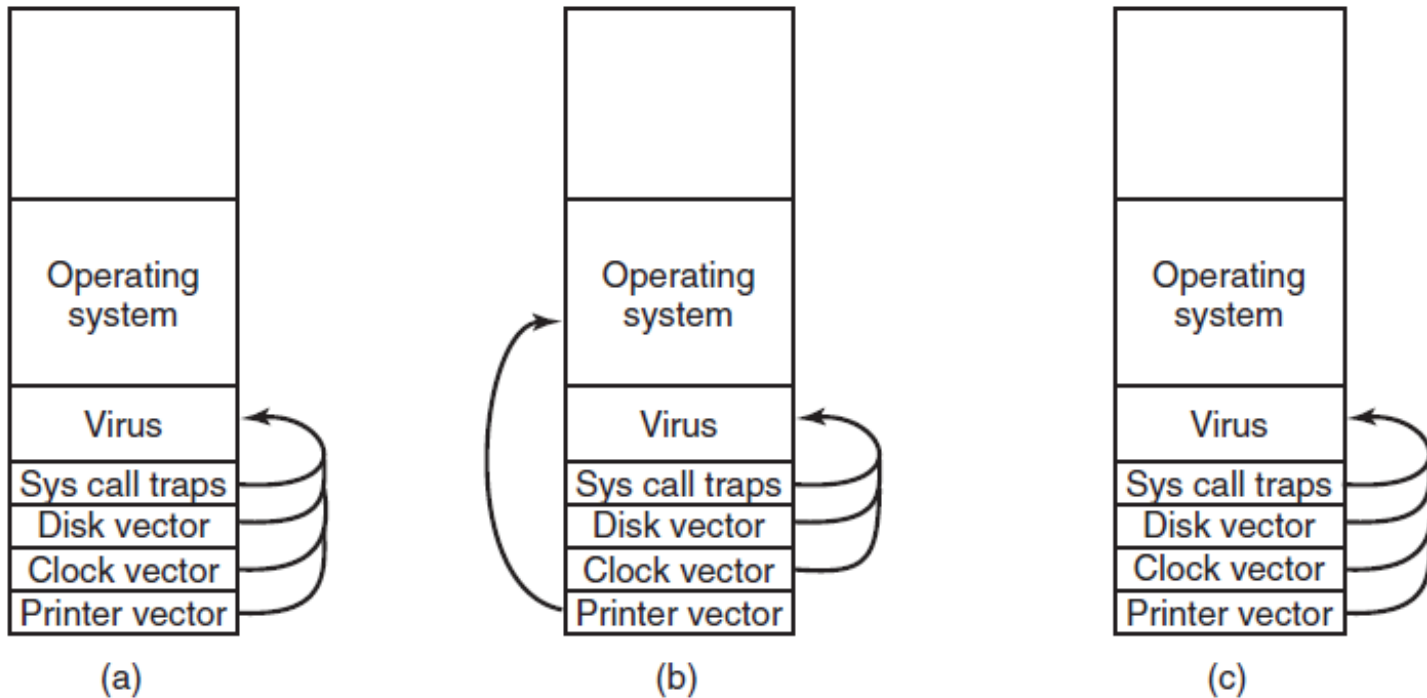
Requires intimate knowledge of the operating system's internal data structure

Copies the first sector of the boot sector to a safe place so it can call it later.

At start-up, it copies the virus to RAM



# BOOT SECTOR VIRUSES



(a) After the virus has captured all the interrupt and trap vectors.

(b) After the operating system has retaken the printer interrupt vector.

(c) After the virus has noticed the loss of the printer interrupt vector and recaptured it.

## 5. DEVICE DRIVER VIRUSES

- Infect the device driver – it's just a executable programs that live on disk
- Device drivers are always loaded at boot time and **may run kernel mode.**

## 6. MACRO VIRUSES

- Virus attached to macros in Microsoft Office (e.g. Excel files)
- Send the infected word document to someone.

## 7. SOURCE CODE VIRUSES

- Very portable
- Looks for C code and changes it to call the virus.

# WORMS

- Self-replicating program
- Moves itself through the network and system without the victims help.
- Consisted of two programs
  - bootstrap to upload worm
  - the worm itself
- Worm first hid its existence
- Next replicated itself on new machines
- Famous Robert Morris Internet worm of 1988
  - Computer Emergency Response Team (CERT)
  - <https://www.cert.be/en>

# SPYWARE

- Description: Runs on the victims machine with victim knowing, doing things behind victim's back
  - Surreptitiously loaded onto a PC without the owner's knowledge
  - Runs in the background doing things behind the owner's back
- Characteristics:
  - Hides, victim cannot easily find
  - Collects data about the user
  - Communicates the collected information back to its distant master
  - Tries to survive determined attempts to remove it
- 3 Broad categories
  - Marketing
  - Surveillance
  - Zombie army

# HOW SPYWARE SPREADS

Possible ways:

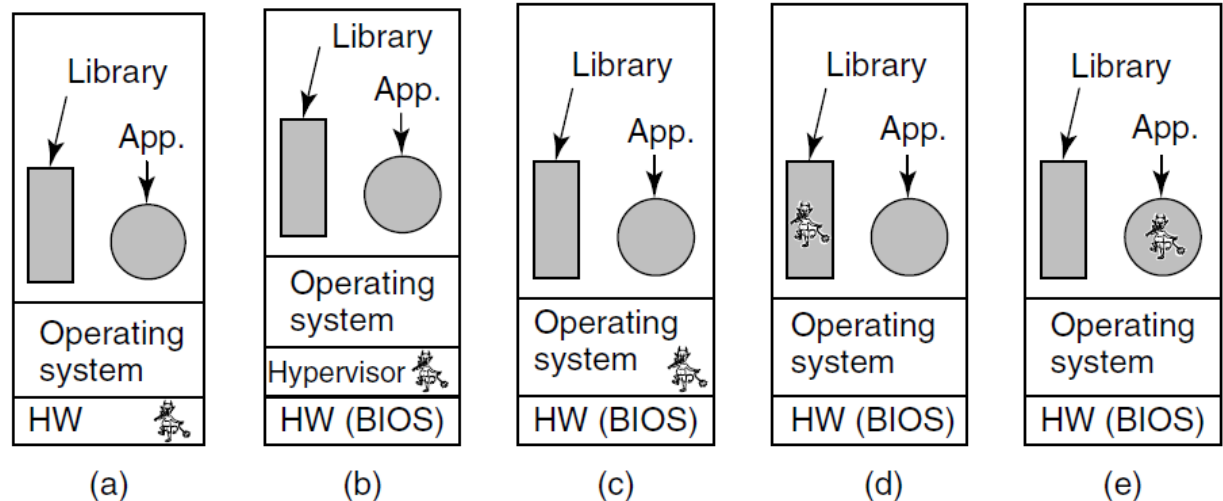
- Same as malware, Trojan horse
- Drive-by download, visit an infected web site
  - Web pages tries to run an .exe file
  - Unsuspecting user installs an infected toolbar, add-ons
  - Malicious activeX controls get installed

# ROOTKITS

- Rootkit: *A rootkit is a collection of computer software, typically malicious, designed to enable access to a computer or an area of its software that is not otherwise allowed (for example, to an unauthorized user) and often masks its existence or the existence of other software\**

- Types of Rootkits:

- (a) Firmware rootkits
- (b) Hypervisor rootkits
- (c) Kernel rootkits
- (d) Library rootkits
- (e) Application rootkits



*\*"Rootkits, Part 1 of 3: The Growing Threat" (PDF). McAfee. 2006-04-17. Archived from [the original](#) (PDF) on 2006-08-23.*



# ROOTKIT DETECTION

## Read the files in the directory

- Unless the dir system call is infected

## Timing related

- Does something take longer than it should

## Sony Rootkit Scandal (2005)

[https://en.wikipedia.org/wiki/Sony\\_BMG\\_copy\\_protection\\_rootkit\\_scandal](https://en.wikipedia.org/wiki/Sony_BMG_copy_protection_rootkit_scandal)



Source: <https://www.wired.com/2006/12/sony-settles-bm/>

# DESIGN PRINCIPLES FOR SECURITY

1. System design should be public
2. Default should be no access
3. Check for current authority
4. Give each process least privilege possible
5. Protection mechanism should be
  - simple
  - uniform
  - in lowest layers of system
6. Scheme should be psychologically acceptable

And ... keep it simple



ETRO  
ELECTRONICS &  
INFORMATICS

# OPERATING SYSTEMS

SECURITY – PART 4: DEFENSES

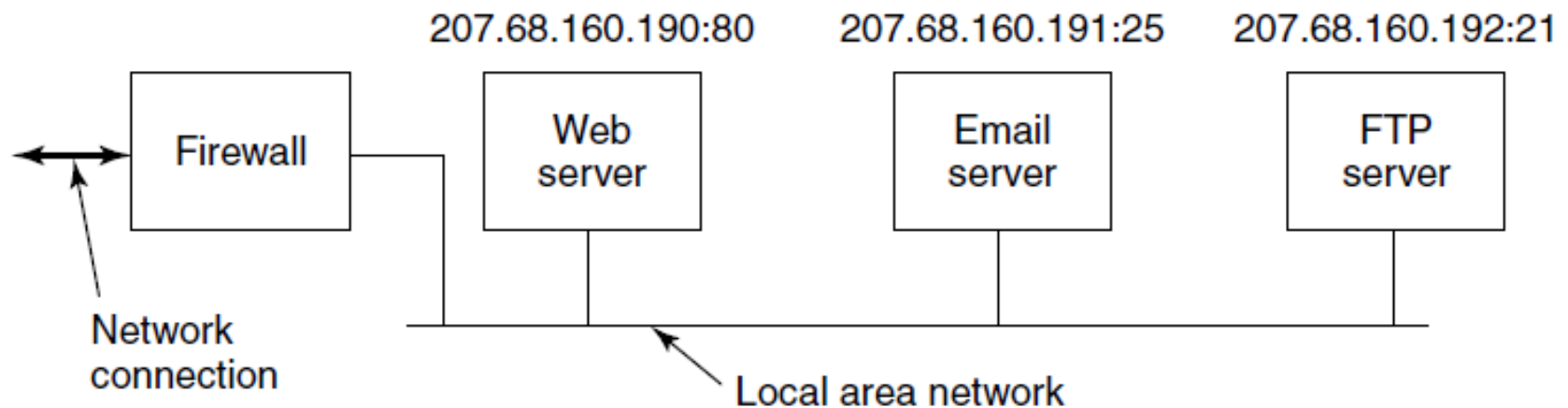
# DEFENSES

1. Firewall
2. Antivirus
3. Code Signing
4. Jailing
5. Model-Based Intrusion Detection
6. Encapsulating Mobile Code
7. Java Security

# DEFENSE 1: FIREWALLS

A simplified view of a hardware firewall protecting a LAN with three computers

- No packets can enter or exit the LAN without approval from Firewall
- Stateless Firewall – Packet header information is used in approval
- Stateful Firewall – Firewall tracks connections ... may inspect packets.



## DEFENSE 2: ANTIVIRUS

### Some techniques

- Virus Scanner
  - Goat file to attract a virus
    - After analysis of virus, add to database known viruses
  - Store file lengths
    - If they change.... Potential problem
  - Hunt for decryption procedure
    - If virus compresses to fit in pgm size..

## DEFENSE 2: VIRUS SCANNERS

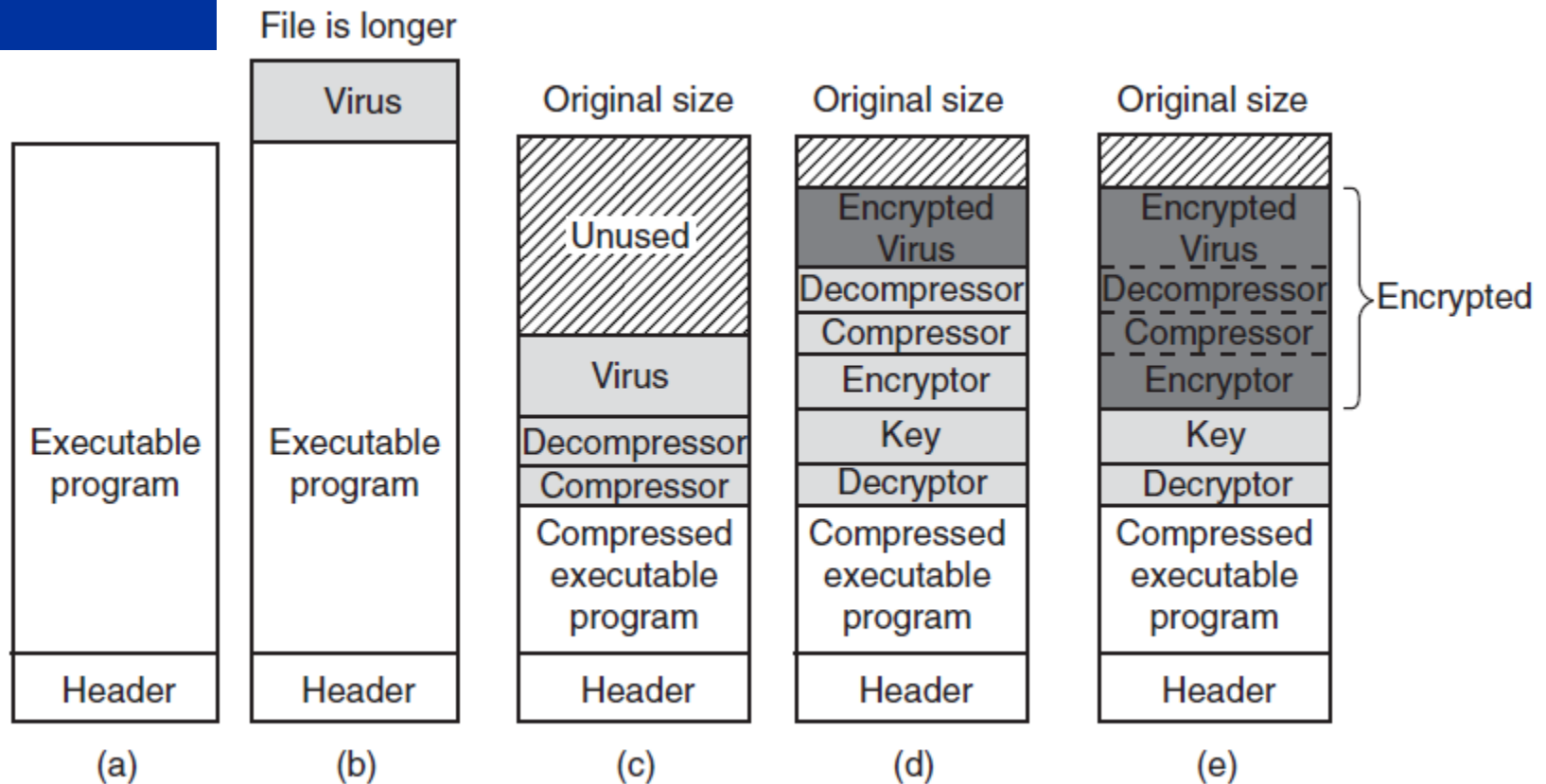


Figure 9-33. (a) A program. (b) An infected program. (c) A compressed infected program. (d) An encrypted virus. (e) A compressed virus with encrypted compression code.

## DEFENSE 2: VIRUS SCANNERS

### POLYMORPHIC VIRUS

All of these examples do the same thing

```
MOV A,R1
ADD B,R1
ADD C,R1
SUB #4,R1
MOV R1,X
```

(a)

```
MOV A,R1
NOP
ADD B,R1
NOP
ADD C,R1
NOP
SUB #4,R1
NOP
MOV R1,X
```

(b)

```
MOV A,R1
ADD #0,R1
ADD B,R1
OR R1,R1
ADD C,R1
SHL #0,R1
SUB #4,R1
JMP .+1
MOV R1,X
```

(c)

```
MOV A,R1
OR R1,R1
ADD B,R1
MOV R1,R5
ADD C,R1
SHL R1,0
SUB #4,R1
ADD R5,R5
MOV R1,X
MOV R5,Y
```

(d)

```
MOV A,R1
TST R1
ADD C,R1
MOV R1,R5
ADD B,R1
CMP R2,R5
SUB #4,R1
JMP .+1
MOV R1,X
MOV R5,Y
```

(e)



# DEFENSE 2: ANTIVIRUS (II)

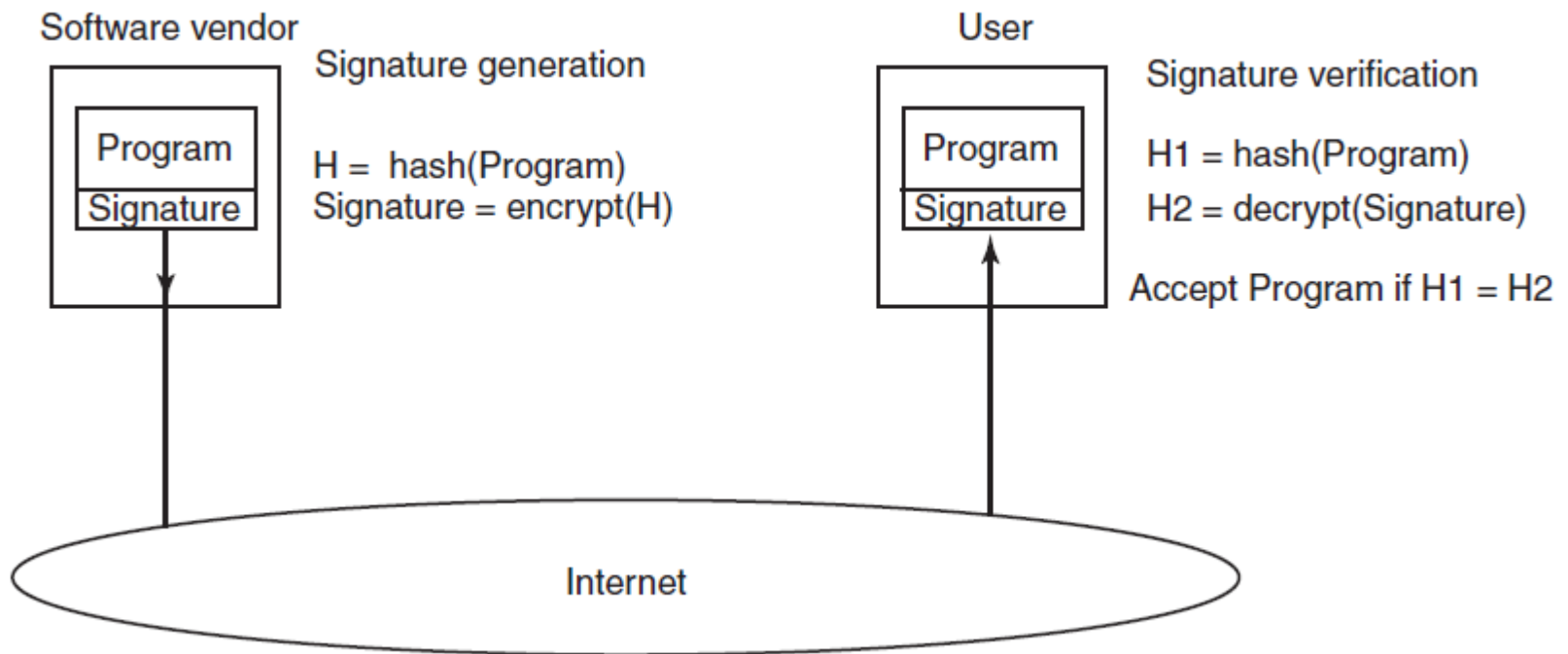
## INTEGRITY CHECKERS

### Some techniques

- Integrity Checkers
  - Compute checksum for clean files
- Behavioral Checkers
  - Monitor all activity
  - Word shouldn't overwrite a file

## DEFENSE 3: CODE SIGNING

Using digital signatures to sign code

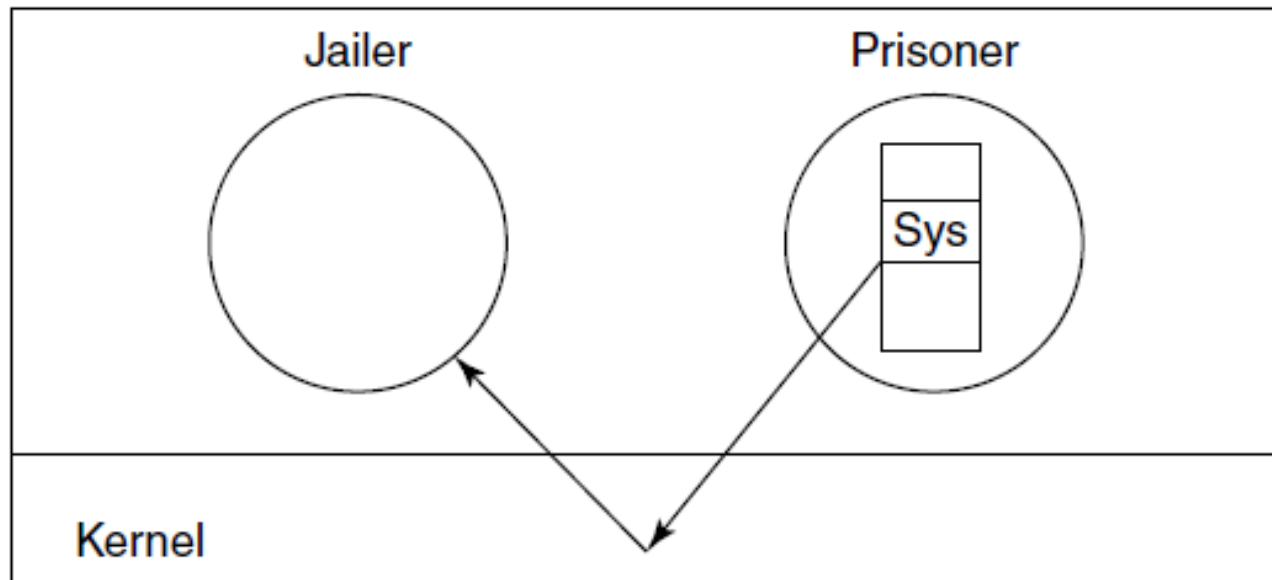


## DEFENSE 4: JAILING

The new program's execution is monitored in a jail.

System call is transferred to jailer who makes the decision if it is allowed.

Like running in a debugger.



# DEFENSE 5: MODEL-BASED INTRUSION DETECTION

## Intrusion Detection System (IDS)

1. Network-Based IDS
  - Focused on incoming packets
2. Host based IDS
  - Static model-based intrusion detection

# MODEL-BASED INTRUSION DETECTION

## Static model-based intrusion detection

- Implemented using jailing technique
- Learn the 'good' behavior of a program from program model.
  - Compiler can generate it and the author certifies it

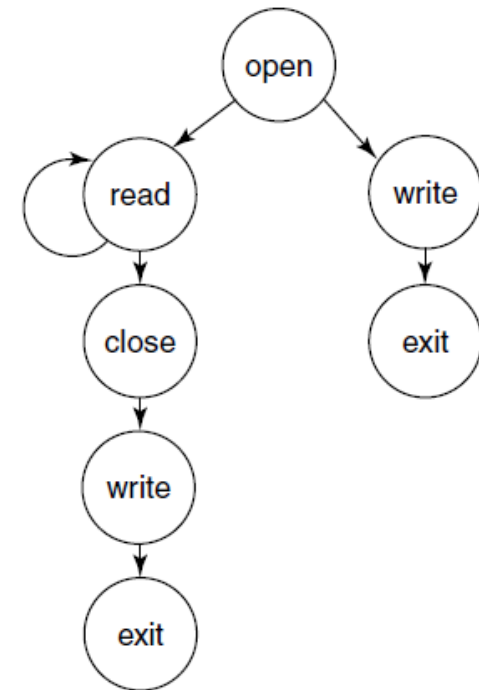
(a) A program.

(b) System call graph for (a).

```
int main(int argc *char argv[])
{
    int fd, n = 0;
    char buf[1];

    fd = open("data", 0);
    if (fd < 0) {
        printf("Bad data file\n");
        exit(1);
    } else {
        while (1) {
            read(fd, buf, 1);
            if (buf[0] == 0) {
                close(fd);
                printf("n = %d\n", n);
                exit(0);
            }
            n = n + 1;
        }
    }
}
```

(a)



(b)

# DEFENSE 6: ENCAPSULATING MOBILE CODE

## Problem:

- Javascript, applets, agents...
- Things that want to execute on our machines
- Things we may want to let execute on our machines

## Defensive methods

- Sandboxing
- Interpretation

# SANDBOXING

Divides virtual address into 2 regions:

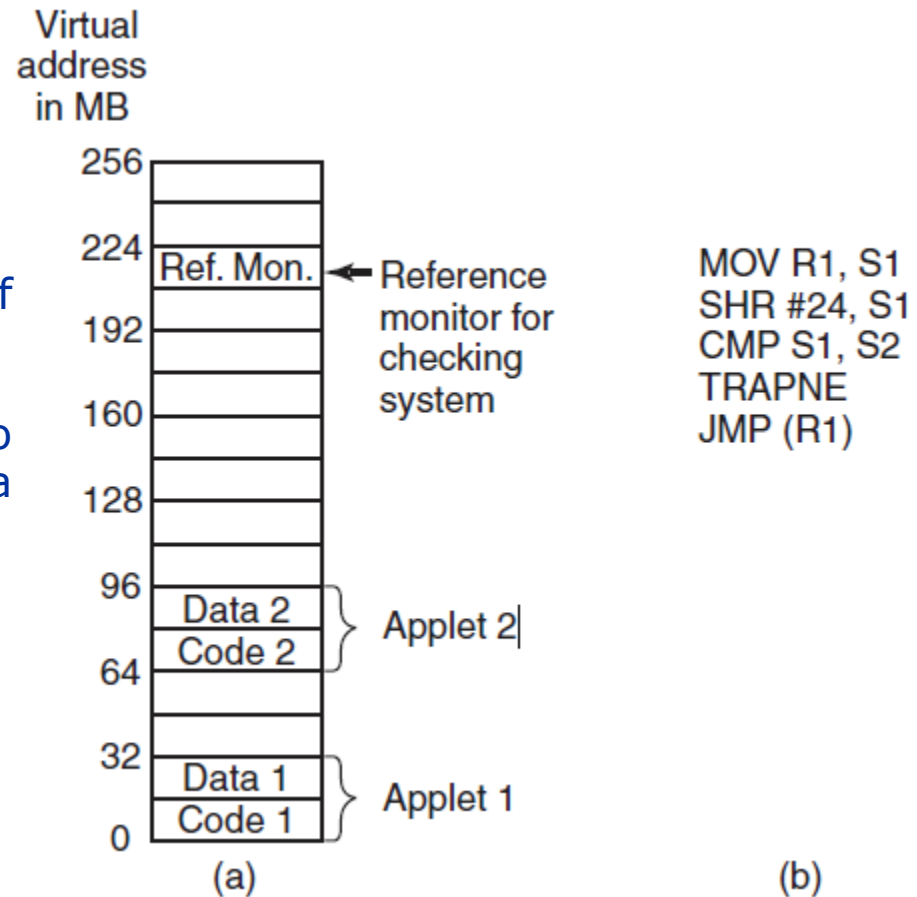
- One for data
- One for code

Confines applet to a limited range of virtual addresses enforced at runtime

Guarantees the applet cannot jump to code outside its code or reference data outside data sandbox.

(a) Memory divided into 16-MB sandboxes.

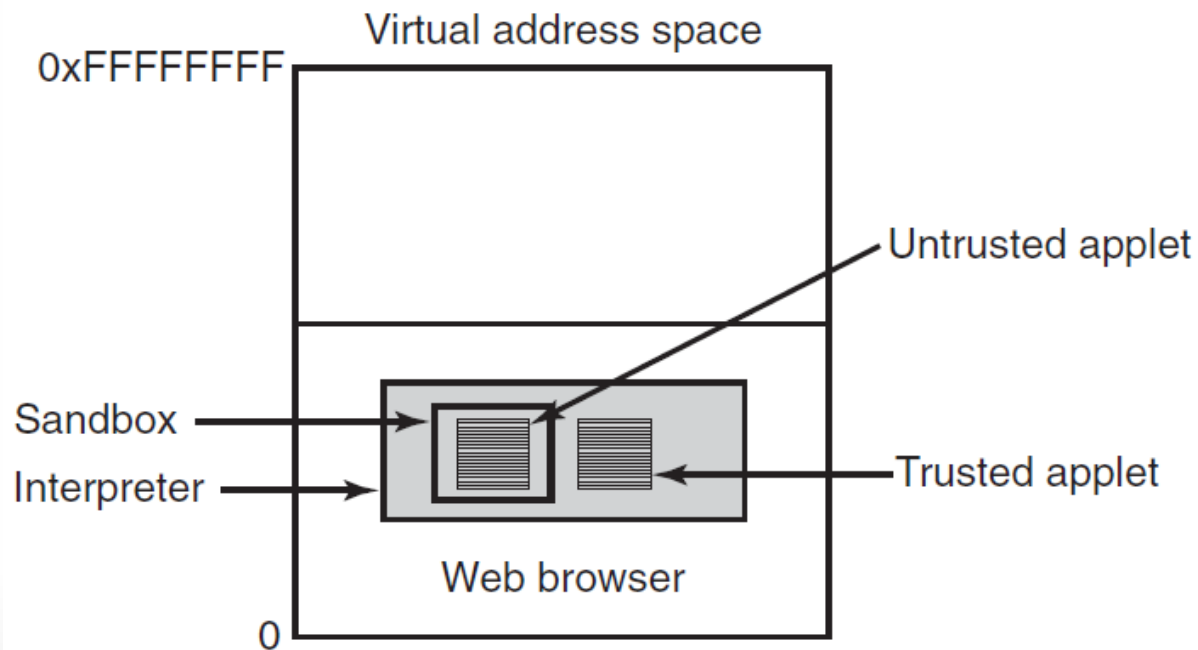
(b) One way of checking an instruction for validity.



# INTERPRETATION

Run applets interpretively.

Every instruction can be examined by interpreter.





## DEFENSE 7: JAVA SECURITY

- A type safe language
  - compiler rejects attempts to misuse variable

Checks on applets include:

1. Does applet attempt to forge pointers?
2. Does it violate access restrictions on private-class members?
3. Does it try to use variable of one type as another?
4. Does it generate stack overflows or underflows?
5. Does it illegally convert variables of one type to another?

## SUMMARY

- Operating systems support mechanisms to satisfy the computer security triad.
- Trusted Computing Base consists of the hardware and software necessary for enforcing all the security rules.
- Protection domains define the operations that can be performed on each object.
- There are different types of attacks by:
  - Exploiting Software
  - Insider Attacks
- Another type of attacks involve malicious software.
  - Trojan horses, viruses, and worms are collectively called malware.
- There are different strategies for defense.
  - From the use of Firewalls, antivirus software to Intrusion Detection Systems.

# Viruses Intruders Spyware Trojans and Adware

"We'll make some piracy  
improvements this year."

-Steve Ballmer  
CEO, Microsoft



Windows Vista™

